

What's in Main

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Abstract

This document lists the main types, functions and syntax provided by theory *Main*. It is meant as a quick overview of what is available. The sophisticated class structure is only hinted at. For details see <http://isabelle.in.tum.de/dist/library/HOL/>.

1 HOL

The basic logic: $x = y$, *True*, *False*, $\neg P$, $P \wedge Q$, $P \vee Q$, $P \longrightarrow Q$, $\forall x. P$, $\exists x. P$, $\exists!x. P$, *THE* $x. P$.

undefined :: 'a

default :: 'a

Syntax

$x \neq y$ $\equiv \neg (x = y)$ ($\sim =$)

$P \longleftrightarrow Q$ $\equiv P = Q$

if x *then* y *else* z \equiv *If* x y z

let $x = e_1$ *in* e_2 \equiv *Let* e_1 ($\lambda x. e_2$)

2 Orderings

A collection of classes defining basic orderings: preorder, partial order, linear order, dense linear order and wellorder.

op \leq :: 'a \Rightarrow 'a \Rightarrow bool (\leq)

op $<$:: 'a \Rightarrow 'a \Rightarrow bool

Least :: ('a \Rightarrow bool) \Rightarrow 'a

min :: 'a \Rightarrow 'a \Rightarrow 'a

max :: 'a \Rightarrow 'a \Rightarrow 'a

top :: 'a

$bot \quad \quad \quad :: 'a$
 $mono \quad \quad \quad :: ('a \Rightarrow 'b) \Rightarrow bool$
 $strict-mono :: ('a \Rightarrow 'b) \Rightarrow bool$

Syntax

$x \geq y \quad \quad \quad \equiv \quad y \leq x \quad \quad \quad (>=)$
 $x > y \quad \quad \quad \equiv \quad y < x$
 $\forall x \leq y. P \quad \quad \equiv \quad \forall x. x \leq y \longrightarrow P$
 $\exists x \leq y. P \quad \quad \equiv \quad \exists x. x \leq y \wedge P$
 Similarly for $<$, \geq and $>$
 $LEAST x. P \quad \equiv \quad Least (\lambda x. P)$

3 Lattices

Classes semilattice, lattice, distributive lattice and complete lattice (the latter in theory *Set*).

$inf :: 'a \Rightarrow 'a \Rightarrow 'a$
 $sup :: 'a \Rightarrow 'a \Rightarrow 'a$
 $Inf :: 'a set \Rightarrow 'a$
 $Sup :: 'a set \Rightarrow 'a$

Syntax

Available by loading theory *Lattice-Syntax* in directory *Library*.

$x \sqsubseteq y \quad \equiv \quad x \leq y$
 $x \sqsubset y \quad \equiv \quad x < y$
 $x \sqcap y \quad \equiv \quad inf\ x\ y$
 $x \sqcup y \quad \equiv \quad sup\ x\ y$
 $\bigsqcap A \quad \equiv \quad Sup\ A$
 $\bigsqcup A \quad \equiv \quad Inf\ A$
 $\top \quad \equiv \quad top$
 $\perp \quad \equiv \quad bot$

4 Set

Sets are predicates: $'a\ set = 'a \Rightarrow bool$

$\{\} \quad \quad \quad :: 'a\ set$
 $insert :: 'a \Rightarrow 'a\ set \Rightarrow 'a\ set$
 $Collect :: ('a \Rightarrow bool) \Rightarrow 'a\ set$
 $op \in \quad :: 'a \Rightarrow 'a\ set \Rightarrow bool \quad \quad \quad (:)$
 $op \cup \quad :: 'a\ set \Rightarrow 'a\ set \Rightarrow 'a\ set \quad (Un)$

$op \cap \quad :: 'a \text{ set} \Rightarrow 'a \text{ set} \Rightarrow 'a \text{ set} \quad (\text{Int})$
 $UNION \quad :: 'a \text{ set} \Rightarrow ('a \Rightarrow 'b \text{ set}) \Rightarrow 'b \text{ set}$
 $INTER \quad :: 'a \text{ set} \Rightarrow ('a \Rightarrow 'b \text{ set}) \Rightarrow 'b \text{ set}$
 $Union \quad :: 'a \text{ set set} \Rightarrow 'a \text{ set}$
 $Inter \quad :: 'a \text{ set set} \Rightarrow 'a \text{ set}$
 $Pow \quad :: 'a \text{ set} \Rightarrow 'a \text{ set set}$
 $UNIV \quad :: 'a \text{ set}$
 $op \text{ ' } \quad :: ('a \Rightarrow 'b) \Rightarrow 'a \text{ set} \Rightarrow 'b \text{ set}$
 $Ball \quad :: 'a \text{ set} \Rightarrow ('a \Rightarrow \text{bool}) \Rightarrow \text{bool}$
 $Bex \quad :: 'a \text{ set} \Rightarrow ('a \Rightarrow \text{bool}) \Rightarrow \text{bool}$

Syntax

$\{x_1, \dots, x_n\} \equiv insert\ x_1\ (\dots\ (insert\ x_n\ \{\})\dots)$
 $x \notin A \equiv \neg(x \in A)$
 $A \subseteq B \equiv A \leq B$
 $A \subset B \equiv A < B$
 $A \supseteq B \equiv B \leq A$
 $A \supset B \equiv B < A$
 $\{x. P\} \equiv Collect\ (\lambda x. P)$
 $\bigcup_{x \in I.} A \equiv UNION\ I\ (\lambda x. A) \quad (\text{UN})$
 $\bigcup x. A \equiv UNION\ UNIV\ (\lambda x. A)$
 $\bigcap_{x \in I.} A \equiv INTER\ I\ (\lambda x. A) \quad (\text{INT})$
 $\bigcap x. A \equiv INTER\ UNIV\ (\lambda x. A)$
 $\forall x \in A. P \equiv Ball\ A\ (\lambda x. P)$
 $\exists x \in A. P \equiv Bex\ A\ (\lambda x. P)$
 $range\ f \equiv f \text{ ' } UNIV$

5 Fun

$id \quad :: 'a \Rightarrow 'a$
 $op \circ \quad :: ('a \Rightarrow 'b) \Rightarrow ('c \Rightarrow 'a) \Rightarrow 'c \Rightarrow 'b$
 $inj\text{-}on \quad :: ('a \Rightarrow 'b) \Rightarrow 'a \text{ set} \Rightarrow \text{bool}$
 $inj \quad :: ('a \Rightarrow 'b) \Rightarrow \text{bool}$
 $surj \quad :: ('a \Rightarrow 'b) \Rightarrow \text{bool}$
 $bij \quad :: ('a \Rightarrow 'b) \Rightarrow \text{bool}$
 $bij\text{-}betw \quad :: ('a \Rightarrow 'b) \Rightarrow 'a \text{ set} \Rightarrow 'b \text{ set} \Rightarrow \text{bool}$
 $fun\text{-}upd \quad :: ('a \Rightarrow 'b) \Rightarrow 'a \Rightarrow 'b \Rightarrow 'a \Rightarrow 'b$

Syntax

$f(x := y) \equiv fun\text{-}upd\ f\ x\ y$
 $f(x_1 := y_1, \dots, x_n := y_n) \equiv f(x_1 := y_1) \dots (x_n := y_n)$

6 Fixed Points

Theory: *Inductive*.

Least and greatest fixed points in a complete lattice $'a$:

$lfp :: ('a \Rightarrow 'a) \Rightarrow 'a$

$gfp :: ('a \Rightarrow 'a) \Rightarrow 'a$

Note that in particular sets $('a \Rightarrow bool)$ are complete lattices.

7 Sum_Type

Type constructor $+$.

$Inl :: 'a \Rightarrow 'a + 'b$

$Inr :: 'a \Rightarrow 'b + 'a$

$op <+> :: 'a\ set \Rightarrow 'b\ set \Rightarrow ('a + 'b)\ set$

8 Product_Type

Types *unit* and \times .

$() :: unit$

$Pair :: 'a \Rightarrow 'b \Rightarrow 'a \times 'b$

$fst :: 'a \times 'b \Rightarrow 'a$

$snd :: 'a \times 'b \Rightarrow 'b$

$split :: ('a \Rightarrow 'b \Rightarrow 'c) \Rightarrow 'a \times 'b \Rightarrow 'c$

$curry :: ('a \times 'b \Rightarrow 'c) \Rightarrow 'a \Rightarrow 'b \Rightarrow 'c$

$Sigma :: 'a\ set \Rightarrow ('a \Rightarrow 'b\ set) \Rightarrow ('a \times 'b)\ set$

Syntax

$(a, b) \equiv Pair\ a\ b$

$\lambda(x, y). t \equiv split\ (\lambda x\ y. t)$

$A \times B \equiv Sigma\ A\ (\lambda.. B)\ (<*>)$

Pairs may be nested. Nesting to the right is printed as a tuple, e.g. (a, b, c) is really $(a, (b, c))$. Pattern matching with pairs and tuples extends to all binders, e.g. $\forall (x, y) \in A. P, \{(x, y). P\}$, etc.

9 Relation

$converse :: ('a \times 'b)\ set \Rightarrow ('b \times 'a)\ set$

$op O :: ('a \times 'b)\ set \Rightarrow ('c \times 'a)\ set \Rightarrow ('c \times 'b)\ set$

$op \text{ `` } :: ('a \times 'b)\ set \Rightarrow 'a\ set \Rightarrow 'b\ set$

$inv\text{-}image :: ('a \times 'a)\ set \Rightarrow ('b \Rightarrow 'a) \Rightarrow ('b \times 'b)\ set$

Id-on :: $'a \text{ set} \Rightarrow ('a \times 'a) \text{ set}$
Id :: $('a \times 'a) \text{ set}$
Domain :: $('a \times 'b) \text{ set} \Rightarrow 'a \text{ set}$
Range :: $('a \times 'b) \text{ set} \Rightarrow 'b \text{ set}$
Field :: $('a \times 'a) \text{ set} \Rightarrow 'a \text{ set}$
refl-on :: $'a \text{ set} \Rightarrow ('a \times 'a) \text{ set} \Rightarrow \text{bool}$
refl :: $('a \times 'a) \text{ set} \Rightarrow \text{bool}$
sym :: $('a \times 'a) \text{ set} \Rightarrow \text{bool}$
antisym :: $('a \times 'a) \text{ set} \Rightarrow \text{bool}$
trans :: $('a \times 'a) \text{ set} \Rightarrow \text{bool}$
irrefl :: $('a \times 'a) \text{ set} \Rightarrow \text{bool}$
total-on :: $'a \text{ set} \Rightarrow ('a \times 'a) \text{ set} \Rightarrow \text{bool}$
total :: $('a \times 'a) \text{ set} \Rightarrow \text{bool}$

Syntax

$r^{-1} \equiv \text{converse } r \quad (\sim^{-1})$

10 Equiv_Relations

equiv :: $'a \text{ set} \Rightarrow ('a \times 'a) \text{ set} \Rightarrow \text{bool}$
op // :: $'a \text{ set} \Rightarrow ('a \times 'a) \text{ set} \Rightarrow 'a \text{ set set}$
congruent :: $('a \times 'a) \text{ set} \Rightarrow ('a \Rightarrow 'b) \Rightarrow \text{bool}$
congruent2 :: $('a \times 'a) \text{ set} \Rightarrow ('b \times 'b) \text{ set} \Rightarrow ('a \Rightarrow 'b \Rightarrow 'c) \Rightarrow \text{bool}$

Syntax

f respects r $\equiv \text{congruent } r \ f$
f respects2 r $\equiv \text{congruent2 } r \ r \ f$

11 Transitive_Closure

rtranc1 :: $('a \times 'a) \text{ set} \Rightarrow ('a \times 'a) \text{ set}$
tranc1 :: $('a \times 'a) \text{ set} \Rightarrow ('a \times 'a) \text{ set}$
reflcl :: $('a \times 'a) \text{ set} \Rightarrow ('a \times 'a) \text{ set}$

Syntax

$r^* \equiv \text{rtranc1 } r \quad (\sim^*)$
 $r^+ \equiv \text{tranc1 } r \quad (\sim^+)$
 $r^= \equiv \text{reflcl } r \quad (\sim=)$

12 Algebra

Theories *OrderedGroup*, *Ring-and-Field* and *Divides* define a large collection of classes describing common algebraic structures from semigroups up to fields. Everything is done in terms of overloaded operators:

```
0      :: 'a
1      :: 'a
op +   :: 'a ⇒ 'a ⇒ 'a
op -   :: 'a ⇒ 'a ⇒ 'a
uminus :: 'a ⇒ 'a      (-)
op *   :: 'a ⇒ 'a ⇒ 'a
inverse :: 'a ⇒ 'a
op /   :: 'a ⇒ 'a ⇒ 'a
abs    :: 'a ⇒ 'a
sgn    :: 'a ⇒ 'a
op dvd :: 'a ⇒ 'a ⇒ bool
op div :: 'a ⇒ 'a ⇒ 'a
op mod :: 'a ⇒ 'a ⇒ 'a
```

Syntax

```
|x|  ≡  abs x
```

13 Nat

```
datatype nat = 0 | Suc nat
```

```
op +   op -   op *   op ^   op div   op mod   op dvd
op ≤   op <   min    max    Min      Max
of-nat :: nat ⇒ 'a
```

14 Int

Type *int*

```
op +   op -   uminus   op *   op ^   op div   op mod   op dvd
op ≤   op <   min      max    Min      Max
abs    sgn
nat    :: int ⇒ nat
of-int :: int ⇒ 'a
ℤ      :: 'a set      (Ints)
```

Syntax

$int \equiv of\text{-}nat$

15 Finite_Set

$finite :: 'a\ set \Rightarrow bool$
 $card :: 'a\ set \Rightarrow nat$
 $fold :: ('a \Rightarrow 'b \Rightarrow 'b) \Rightarrow 'b \Rightarrow 'a\ set \Rightarrow 'b$
 $fold\text{-}image :: ('b \Rightarrow 'b \Rightarrow 'b) \Rightarrow ('a \Rightarrow 'b) \Rightarrow 'b \Rightarrow 'a\ set \Rightarrow 'b$
 $setsum :: ('a \Rightarrow 'b) \Rightarrow 'a\ set \Rightarrow 'b$
 $setprod :: ('a \Rightarrow 'b) \Rightarrow 'a\ set \Rightarrow 'b$

Syntax

$\sum A \equiv setsum (\lambda x. x) A \quad (\text{SUM})$
 $\sum_{x \in A}. t \equiv setsum (\lambda x. t) A$
 $\sum_{x|P}. t \equiv \sum_{x \in \{x. P\}}. t$
Similarly for \prod instead of \sum (PROD)

16 Wellfounded

$wf :: ('a \times 'a)\ set \Rightarrow bool$
 $acyclic :: ('a \times 'a)\ set \Rightarrow bool$
 $acc :: ('a \times 'a)\ set \Rightarrow 'a\ set$
 $measure :: ('a \Rightarrow nat) \Rightarrow ('a \times 'a)\ set$
 $op\ <*\text{lex}*> :: ('a \times 'a)\ set \Rightarrow ('b \times 'b)\ set \Rightarrow (('a \times 'b) \times 'a \times 'b)\ set$
 $op\ <*\text{mlex}*> :: ('a \Rightarrow nat) \Rightarrow ('a \times 'a)\ set \Rightarrow ('a \times 'a)\ set$
 $less\text{-}than :: (nat \times nat)\ set$
 $pred\text{-}nat :: (nat \times nat)\ set$

17 SetInterval

$lessThan :: 'a \Rightarrow 'a\ set$
 $atMost :: 'a \Rightarrow 'a\ set$
 $greaterThan :: 'a \Rightarrow 'a\ set$
 $atLeast :: 'a \Rightarrow 'a\ set$
 $greaterThanLessThan :: 'a \Rightarrow 'a \Rightarrow 'a\ set$
 $atLeastLessThan :: 'a \Rightarrow 'a \Rightarrow 'a\ set$
 $greaterThanAtMost :: 'a \Rightarrow 'a \Rightarrow 'a\ set$
 $atLeastAtMost :: 'a \Rightarrow 'a \Rightarrow 'a\ set$

Syntax

$\{..<y\}$	\equiv	$lessThan\ y$
$\{..y\}$	\equiv	$atMost\ y$
$\{x<..\}$	\equiv	$greaterThan\ x$
$\{x..\}$	\equiv	$atLeast\ x$
$\{x<..$	\equiv	$greaterThanLessThan\ x\ y$
$\{x..$	\equiv	$atLeastLessThan\ x\ y$
$\{x<..y\}$	\equiv	$greaterThanAtMost\ x\ y$
$\{x..y\}$	\equiv	$atLeastAtMost\ x\ y$
$\bigcup i \leq n. A$	\equiv	$\bigcup i \in \{..n\}. A$
$\bigcup i < n. A$	\equiv	$\bigcup i \in \{..$

Similarly for \bigcap instead of \bigcup

$\sum x = a..b. t$	\equiv	$setsum\ (\lambda x. t)\ \{a..b\}$
$\sum x = a..b. t$	\equiv	$setsum\ (\lambda x. t)\ \{a..b\}$
$\sum x \leq b. t$	\equiv	$setsum\ (\lambda x. t)\ \{..b\}$
$\sum x < b. t$	\equiv	$setsum\ (\lambda x. t)\ \{..b\}$

Similarly for \prod instead of \sum

18 Power

$op\ ^\wedge :: 'a \Rightarrow nat \Rightarrow 'a$

19 Iterated Functions and Relations

Theory: *Relation-Power*

Iterated functions $(f :: 'a \Rightarrow 'a)^\wedge n$ and relations $(r :: ('a \times 'a) \text{set})^\wedge n$.

20 Option

datatype $'a\ option = None \mid Some\ 'a$

$the :: 'a\ option \Rightarrow 'a$
 $Option.map :: ('a \Rightarrow 'b) \Rightarrow 'a\ option \Rightarrow 'b\ option$
 $Option.set :: 'a\ option \Rightarrow 'a\ set$

21 List

datatype $'a\ list = [] \mid op \# 'a\ ('a\ list)$

$op\ @ :: 'a\ list \Rightarrow 'a\ list \Rightarrow 'a\ list$
 $butlast :: 'a\ list \Rightarrow 'a\ list$
 $concat :: 'a\ list\ list \Rightarrow 'a\ list$

$distinct :: 'a\ list \Rightarrow bool$
 $drop :: nat \Rightarrow 'a\ list \Rightarrow 'a\ list$
 $dropWhile :: ('a \Rightarrow bool) \Rightarrow 'a\ list \Rightarrow 'a\ list$
 $filter :: ('a \Rightarrow bool) \Rightarrow 'a\ list \Rightarrow 'a\ list$
 $foldl :: ('a \Rightarrow 'b \Rightarrow 'a) \Rightarrow 'a \Rightarrow 'b\ list \Rightarrow 'a$
 $foldr :: ('a \Rightarrow 'b \Rightarrow 'b) \Rightarrow 'a\ list \Rightarrow 'b \Rightarrow 'b$
 $hd :: 'a\ list \Rightarrow 'a$
 $last :: 'a\ list \Rightarrow 'a$
 $length :: 'a\ list \Rightarrow nat$
 $lenlex :: ('a \times 'a)\ set \Rightarrow ('a\ list \times 'a\ list)\ set$
 $lex :: ('a \times 'a)\ set \Rightarrow ('a\ list \times 'a\ list)\ set$
 $lexn :: ('a \times 'a)\ set \Rightarrow nat \Rightarrow ('a\ list \times 'a\ list)\ set$
 $lexord :: ('a \times 'a)\ set \Rightarrow ('a\ list \times 'a\ list)\ set$
 $listrel :: ('a \times 'a)\ set \Rightarrow ('a\ list \times 'a\ list)\ set$
 $lists :: 'a\ set \Rightarrow 'a\ list\ set$
 $listset :: 'a\ set\ list \Rightarrow 'a\ list\ set$
 $listsum :: 'a\ list \Rightarrow 'a$
 $list-all2 :: ('a \Rightarrow 'b \Rightarrow bool) \Rightarrow 'a\ list \Rightarrow 'b\ list \Rightarrow bool$
 $list-update :: 'a\ list \Rightarrow nat \Rightarrow 'a \Rightarrow 'a\ list$
 $map :: ('a \Rightarrow 'b) \Rightarrow 'a\ list \Rightarrow 'b\ list$
 $measures :: ('a \Rightarrow nat)\ list \Rightarrow ('a \times 'a)\ set$
 $remdups :: 'a\ list \Rightarrow 'a\ list$
 $removeAll :: 'a \Rightarrow 'a\ list \Rightarrow 'a\ list$
 $remove1 :: 'a \Rightarrow 'a\ list \Rightarrow 'a\ list$
 $replicate :: nat \Rightarrow 'a \Rightarrow 'a\ list$
 $rev :: 'a\ list \Rightarrow 'a\ list$
 $rotate :: nat \Rightarrow 'a\ list \Rightarrow 'a\ list$
 $rotate1 :: 'a\ list \Rightarrow 'a\ list$
 $set :: 'a\ list \Rightarrow 'a\ set$
 $sort :: 'a\ list \Rightarrow 'a\ list$
 $sorted :: 'a\ list \Rightarrow bool$
 $splice :: 'a\ list \Rightarrow 'a\ list \Rightarrow 'a\ list$
 $sublist :: 'a\ list \Rightarrow (nat \Rightarrow bool) \Rightarrow 'a\ list$
 $take :: nat \Rightarrow 'a\ list \Rightarrow 'a\ list$
 $takeWhile :: ('a \Rightarrow bool) \Rightarrow 'a\ list \Rightarrow 'a\ list$
 $tl :: 'a\ list \Rightarrow 'a\ list$
 $upt :: nat \Rightarrow nat \Rightarrow nat\ list$
 $upto :: 'a \Rightarrow 'a \Rightarrow 'a\ list$
 $zip :: 'a\ list \Rightarrow 'b\ list \Rightarrow ('a \times 'b)\ list$

Syntax

$[x_1, \dots, x_n] \equiv x_1 \# \dots \# x_n \# []$
 $[m..<n] \equiv upt\ m\ n$

$$\begin{aligned}
[i..j] &\equiv upto\ i\ j \\
[e.\ x \leftarrow xs] &\equiv map\ (\lambda x.\ e)\ xs \\
[x \leftarrow xs.\ b] &\equiv filter\ (\lambda x.\ b)\ xs \\
xs[n := x] &\equiv list-update\ xs\ n\ x \\
\sum x \leftarrow xs.\ e &\equiv listsum\ (map\ (\lambda x.\ e)\ xs)
\end{aligned}$$

List comprehension: $[e.\ q_1, \dots, q_n]$ where each qualifier q_i is either a generator $pat \leftarrow e$ or a guard, i.e. boolean expression.

22 Map

Maps model partial functions and are often used as finite tables. However, the domain of a map may be infinite.

$$'a \multimap 'b = 'a \Rightarrow 'b\ option$$

$$\begin{aligned}
Map.empty &:: 'a \multimap 'b \\
op ++ &:: ('a \multimap 'b) \Rightarrow ('a \multimap 'b) \Rightarrow 'a \multimap 'b \\
op \circ_m &:: ('a \multimap 'b) \Rightarrow ('c \multimap 'a) \Rightarrow 'c \multimap 'b \\
op | ' &:: ('a \multimap 'b) \Rightarrow 'a\ set \Rightarrow 'a \multimap 'b \\
dom &:: ('a \multimap 'b) \Rightarrow 'a\ set \\
ran &:: ('a \multimap 'b) \Rightarrow 'b\ set \\
op \subseteq_m &:: ('a \multimap 'b) \Rightarrow ('a \multimap 'b) \Rightarrow bool \\
map-of &:: ('a \times 'b)\ list \Rightarrow 'a \multimap 'b \\
map-upds &:: ('a \multimap 'b) \Rightarrow 'a\ list \Rightarrow 'b\ list \Rightarrow 'a \multimap 'b
\end{aligned}$$

Syntax

$$\begin{aligned}
Map.empty &\equiv \lambda x.\ None \\
m(x \mapsto y) &\equiv m(x := Some\ y) \\
m(x_1 \mapsto y_1, \dots, x_n \mapsto y_n) &\equiv m(x_1 \mapsto y_1) \dots (x_n \mapsto y_n) \\
[x_1 \mapsto y_1, \dots, x_n \mapsto y_n] &\equiv Map.empty(x_1 \mapsto y_1, \dots, x_n \mapsto y_n) \\
m(xs\ [\mapsto]\ ys) &\equiv map-upds\ m\ xs\ ys
\end{aligned}$$