SIGNETICS BIPOLAR/MOS MICROPROCESSOR

DATA MANUAL



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he present rapid growth of the electronic design revolution is being fueled by the microprocessor/microcomputer. As one of the industry leaders, Signetics gives new impetus to advancing the state-of-the-art with its technological strength, product depth and manufacturing know-how. We continue to develop all aspects of the microprocessor/microcomputer area, simultaneously providing products, ready-to-use training, software and development tools that help you to understand and apply Signetics advanced products.

This Microprocessor Data Manual will enable you to examine Signetics complete line of Bipolar and MOS Microprocessors, as well as a full line of support products. Each data sheet fully explains the operation of each product and contains application hints on how to best utilize the products' design advantages.

You will also find a variety of development systems, educational products, prototyping kits and software for both bipolar and MOS. And with a team of field application engineers throughout the country you will find that Signetics can provide you with the answer to your design problems.

We offer one of the broadest selections available of both bipolar and MOS microprocessor/microcomputers. Combined with our complete line of logic, analog, bipolar and MOS memories, a design engineer can select one source that will provide him with the best possible, cost effective solution to his design problems.

This book contains a compilation of all products currently available. Signetics is continuously developing new products. As you see new product announcements, you should contact your local Signetics sales office, representative or authorized distributor or write Signetics directly at 811 East Arques Avenue, Sunnyvale, California, 94086, for the latest technical information.

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CHAPTER I BIPOLAR Microprocessors Peripherals and Development Products

Microcontrol and Arithmetic Units

The introduction of the Signetics Bit-Slice Microprocessors has brought new levels of high performance to microprocessor applications not previously possible with MOS technology. Combining the Schottky bipolar microprocessors with industry standard memory and support circuits, microinstruction cycle times of 100ns are possible.

In the majority of cases, the choice of a bipolar microprocessor slice, as opposed to an MOS device, is based on speed or flexibility of microprogramming. Starting with these characteristics, the design of the Signetics slice microprocessors has been optimized around the following objectives:

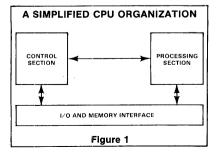
- · Fast cycle time
- All memory and support chips are industry standard
- · Cooler operation
- Lower total system cost

Furthermore, systems built with large-scale integrated circuits are much smaller and require less power than equivalent systems using medium and/or small scale integrated circuits.

Typically, slice microprocessors are employed in the realization of the Central Processing Unit (CPU) of a computer or for implementing dedicated smart controllers. The generalized and simplified structure of a CPU or "Smart" controller can be typically classified into 3 distinct but interactively related functional sections. These sections are generally referred to as the Processing section, the Control section, and the I/O and Memory Interface section. A simplified block diagram of a CPU is illustrated in Figure 1.

The major functions of the Processing section are to:

- provide data transfer paths:
- manipulate data through logic and arithmetic operations;
- provide storage facilities such as a register file; and
- generate necessary status flags based on the kind of operation performed by the ALU.



The major functions of the Control section are to:

- initiate memory or I/O operations:
- decode macroinstructions;
- control the manipulation and transfer of data:
- · test status conditions; and
- · sample and respond to interrupts.

The major functions of the I/O and Memory Interface section are to:

- multiplex data to the proper destination:
- · provide bus driving/receiving capability; and
- · provide latching capability.

With state-of-the-art bipolar Schottky technology, high-performance microprocessors are designed to perform functions of the Processing section. Due to the limitation on the number of pins and chip size, the overall Processing section is partitioned into several functionally equivalent slices. In today's bipolar microprocessor market, 2-bit and 4bit slice architecture predominates. Each architecture type has its uniqueness but, in general, a slice contains a group of general purpose registers, an accumulator, specialpurpose register(s) ALU and related status flags. All of these elements constitute the Processing section of a CPU. The flexibility of slice components allows the designer to construct a processing section of any desired width as required by his application.

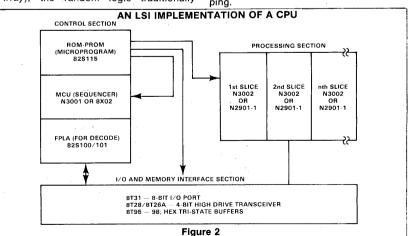
The Control section of the CPU is more complex in design. Typically this section includes the macroinstruction decode logic, test-branch decode, microprogram sequencing logic, and the control store where the microprogram resides. Aside from the microprogram, the remaining portion of the Control section (macroinstruction decode and test-branch decode and sequencing logic), does not lend itself to efficient partitioning into vertical slices. This is due to the random nature of the logic usually found in the Control section. However, horizontal functional grouping is possible. For example, the macroinstruction decode and testbranch decode logic can now be replaced by the FPLA (Field Programmable Logic Array); the random logic traditionally

needed to implement the microprogram sequencing can now be replaced by the Microprogram Control Unit; and, of course, the microprogram can be stored in high density PROMs or ROMs. Since the designer must define his own microstructure, the slice microprocessors permit fundamental optimizations to be made. With slice hardware, the designer may have no macroinstructions at all, placing all of the program in PROM for dedicated control applications. Or he may define, as required, any number of macroinstructions selected specifically for his particular processor purpose. Various minicomputers and several MOS microprocessors have been emulated using slice hardware.

The I/O and Memory Interface section consists mainly of I/O ports, high power bus drivers, receivers, and some temporary register storage facilities. Bidirectional and tri-state devices are the most popular logic elements for implementing this interface structure.

Figure 2 shows an LSI approach to the implementation of the same generalized CPU structure indicated earlier.

Data specifications for Signetics' line of slice microprocessor components are contained within this chapter. Included is the popular 3000 series Microprogram Control Unit and the 2-bit slice Central Processing Element. These Signetics devices feature improved performance specifications over 3000 series components available on the general market. Moreover, the unique Signetics XL plastic package design results in significantly cooler operation of the chip than was previously possible with other plastic package designs. This section also features the 8X02 Control Store Sequencer. This device may be used with any TTL compatible slice processing elements and features extreme ease of use. The 8 simple, yet powerful, instructions permit subroutining and looping (using internal stack), unrestricted jumping, unrestricted conditional branching and conditional instruction skipping



CONTROL STORE SEQUENCER

PRELIMINARY SPECIFICATION

8X02-I.XL

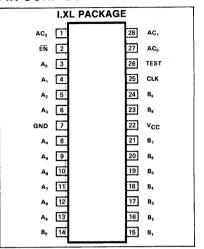
DESCRIPTION

The Signetics 8X02 is a low power Schottky LSI device intended for use in high performance microprogrammed systems to control the fetch sequence of microinstructions. When combined with standard ROM or PROM, the 8X02 forms a powerful microprogrammed control section for computers, controllers, or sequential logic.

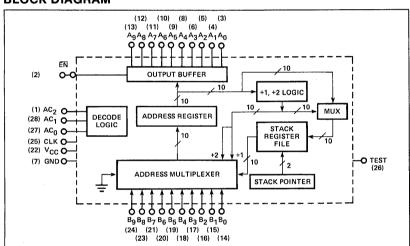
FEATURES

- Low power Schottky process
- 1024 microinstruction addressability
- N-way branch
- · 4-level stack register file (LIFO type)
- · Automatic push/pop stack operation
- "Test and skip" operation on test input line
- 3-bit command code
- · Tri-state buffered outputs
- Auto-reset to address 0 during power-up
- Conditional branching, pop stack, and push stack

PIN CONFIGURATION



BLOCK DIAGRAM



PIN DESIGNATION

PIN	SYMBOL	NAME AND FUNCTION	TYPE
5-6 8-13	A ₀ -A ₉	Microprogram Address outputs	Three-state Active high
1,28,27	AC ₀ -AC ₂	Next Address Control Function inputs All addressing control functions are selected by these command lines.	Active high
14-21 23-24	B ₀ -B ₉	Branch Address inputs Determines the next address of an N-way branch when used with the BRANCH TO SUBROUTINE (BSR) or BRANCH ON TEST (BRT) command.	Active high
2	ĒN	Enable input When in the low state, the Microprogram Address outputs are enabled.	Active low
25	CLK	Clock Input—High to Low transition for stack operations, Low to High transition for address modification.	
26	TEST	Test input Used in conjunction with four NEXT ADDRESS CONTROL FUNCTION commands to effect conditional skips, branches, and stack operations.	Active high
7	GND	Ground	
22	v _{CC}	+5 Volt supply	

PRELIMINARY SPECIFICATION

8X02-I.XL

FUNCTIONAL DESCRIPTION

The Signetics 8X02 Control Store Sequencer is an LSI device using low power Schottky technology and is intended for use in high performance microprogrammed applications. When used alone, the 8X02 is capable of addressing up to 1K words of microprogram. This may be expanded to any microprogram size by conventional paging techniques.

The Address Register consists of 10 D-type, edge-triggered flip-flops with a common clock. A new address is entered into the Address Register on the low-to-high transition of the clock. The next address to be entered into the Address Register is supplied via the Address Multiplexer.

The Address Multiplexer is a 5-input device that is used to select either the branch input, +1 adder, +2 adder, stack register file, or ground (all zeros) as the source of the next microinstruction address. The proper multiplexer channel is automatically selected via the Decode Logic according to the Address Control Function Input and Test Input line.

The +1, +2 logic is used to increment the present contents of the Address Register by 1 or 2, depending on the function input command. Thus, the next address to the Control Store ROM/PROM may be either the current address plus 1 (N+1) or the current address plus 2 (N+2). If the same Microprogram Address is to be used on successive occasions, the clock to the 8X02 must simply be disabled; therefore, no new address is loaded into the Address Register.

The Stack File Register is used to provide a return address linkage whenever a subroutine or loop is executed. The 4X10 stack operates in a last-in, first-out (LIFO) mode, with the stack pointer always pointing to the next address to be read. Operation of the stack pointer is automatically controlled by the Address Control Function Inputs. Since the stack is 4 words deep, up to 4 loops and/or subroutines may be nested.

The branch input is a 10-bit field of direct inputs to the multiplexer which can be selected as the next control store address. Using the appropriate branch command, an N-way branch is possible where N is the

address of any microinstruction within the 1024 word microcode page. Likewise, the RESET command is a special case of an N-way branch in which the multiplexer selects an all zeros input, forcing the next microinstruction address to be zero.

The Test Input line is used in conjunction with the conditional execution of 4 Address Control Function commands. When the Test Input is false (low), the sequencer simply increments to the next address (N+1). When it is true (high), the sequencer executes a branch as defined by the input command, thereby transferring control to another portion of the microprogram.

All Address Output lines of the 8X02 are three-state buffered outputs with a common enable line $(\overline{\text{EN}}).$ When the Enable line is high, all outputs are placed in a high-impedance state, and external access to the control store ROM/PROM is possible. This allows a preprogrammed set of microinstructions to be executed from external or built-in test equipment (BITE), vectored interrupts, and Writable Control Store if implemented.

NEXT ADDRESS CONTROL FUNCTION TABLE

MNEMONIC	DESCRIPTION	FUNC AC ₂			TEST	NEXT ADDRESS	STACK	STACK POINTER
TSK	Test and skip	0	0	0	False True	Current + 1 Current + 2	N.C. N.C.	N.C. N.C.
INC	Increment	0	0	1	x	Current + 1	N.C.	N.C.
BLT	Branch to loop if test input true	0	1	0	False True	Current + 1 Stack reg file	X POP (read)	Decr Decr
POP	POP stack	0	1	1	x	Stack reg file	POP (read)	Decr
BSR	Branch to sub- routine if test input true	1	0	0 .	False True	Current + 1 Branch address	N.C. PUSH (Curr + 1)	N.C. Incr
PLP	Push for looping	1	0	1	x	Current + 1	PUSH (Curr Addr)	Incr
BRT	Branch if test input true	1	1	0	False True	Current + 1 Branch address	N.C. N.C.	N.C. N.C.
RST	Set microprogram address output to zero	1	1	1	×	All O's	N.C.	N.C.

X = Don't care N.C. = No change PRELIMINARY SPECIFICATION

8X02-I,XL

FUNCTIONAL DESCRIPTION

The following is a description of each of the eight Next Address Control Functions (AC₂-AC₀)

MNEMONIC	FUNCTION DESCRIPTION
тѕк	AC ₂₋₀ = 000: TEST AND SKIP Perform test on Test Input Line. If test is Next Address = Current Address + 1 False (Low): Stack Pointer unchanged If test is Next Address = Current Address + 2 True (High) (i.e. Skip next microinstruction) Stack Pointer unchanged
INC	AC ₂₋₀ = 001: INCREMENT Next Address = Current Address + 1 Stack Pointer unchanged
BLT	AC ₂₋₀ = -010: BRANCH TO LOOP IF TEST CONDITION TRUE. Perform test on Test Input Line. If test is Next Address = Current Address + 1 False (Low): Stack Pointer decremented by 1 If test is Next Address = Address from Stack True (High): Register File (POP) Stack Pointer decremented by 1
POP	AC ₂₋₀ = 011: POP STACK Next Address = Address from Stack Register File (POP) Stack Pointer decremented by 1
BSR	AC ₂₋₀ = 100: BRANCH TO SUBROUTINE IF TEST CONDITION TRUE. Perform test on Test Input Line. If test is Next Address = Current Address + 1 False (Low): Stack Pointer unchanged If test is Next Address = Branch Address Input (B ₀₋₉) True (High): Stack Pointer incremented by 1 PUSH (write) Current Address + 1 → Stack Register File
PLP	AC ₂₋₀ = 101: PUSH FOR LOOPING Next Address = Current Address + 1 Stack Pointer incremented by 1 PUSH (write) Current Address → Stack Register File
BRT	AC ₂₋₀ = 110: BRANCH ON TEST CONDITION TRUE Perform test on Test Input Line. If test is Next Address = Current Address + 1 False (Low): Stack Pointer unchanged If test is Next Address = Branch Address Input (B ₀₋₉) True (High): Stack Pointer unchanged
RST	AC ₂₋₀ = 111: RESET TO ZERO Next Address = 0 Stack Pointer unchanged

ABSOLUTE MAXIMUM RATINGS

	PARAMETER	RATING	UNIT
VCC	Power supply voltage	+7	Vdc
VIN	Input voltage	+5.5	Vdc
VO	Off-State output voltage	+5.5	Vdc
TA	Operating temperature range	0° to +70°	°C
TSTG	Storage temperature range	-65° to +150°	°C

PRELIMINARY SPECIFICATION

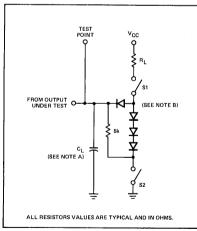
8X02-I,XL

DC ELECTRICAL CHARACTERISTICS TA = 0°C to +70°C, $4.75 \le V \hat{CC} \le 5.25V$

	PARAMETER	TEST CONDITIONS		LIMITS			
	T ANAMETER	TEST CONDITIONS	Min	Typ ¹	Max	UNIT	
VIH	High level input voltage		2			V	
VIL	Low level input voltage				0.8	V	
V,	Input clamp voltage	V _{CC} = 4.75V, I _I = -18mA			-1.5	V	
Voн	High level output voltage	V _{CC} = 4.75V, I _{OH} = -2.6mA	2.4			V	
VOL	Low level output voltage	V _{CC} = 4.75V, I _{OL} = 8mA			0.5	l v	
11	Input current at maximum Input voltage	V _{CC} = 5.25V, V _I = 5.5V			100	μΑ	
'ІН	High level input current AC_2 - AC_0 , \overline{EN} , TEST B_9 - B_0 CLK	V _{CC} = 5.25V, V _I = 2.7V			40 20 60	μΑ μΑ μΑ	
اال	Low level input current AC_2 - AC_0 , \overline{EN} , $TEST$ B_9 - B_0 CLK	V _{CC} = 5.25V, V _I = 0.4V			-0.72 -0.36 -1.08	mA mA mA	
los	Short-circuit output current	V _{CC} = 5.25V	-15	}	-100	mA	
¹ ozh	High-Z state output current	V _{OUT} = 2.7V			20	μΑ	
l OZL	High-Z state output current	V _{OUT} = 0.4V			-20	μΑ	
Icc	Supply current	V _{CC} = 5.25V		165	200	mA	

NOTE

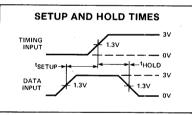
TEST LOAD CIRCUIT

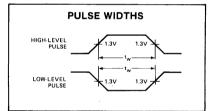


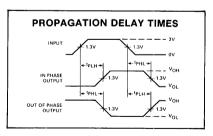
NOTES

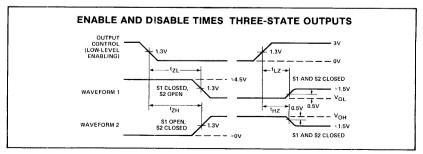
- A. C_L includes probe and jig capacitance.
- B. All diodes are 1N916 or 1N3064.
- C. R_L = 2k, C = 15pF.

VOLTAGE WAVEFORMS









^{1.} All typical values are at V_{CC} = 5V, T_A = 25°C.

CONTROL STORE SEQUENCER

PRELIMINARY SPECIFICATION

8X02-I,XL

AC ELECTRICAL CHARACTERISTICS $T_A = 0^{\circ} -70^{\circ}C$, $V_{CC} = 5.0V \pm 5\%$

	0404447750	Τ0	TO FROM		LIMITS	;	UNIT
	PARAMETER	то	FROM	Min	Typ¹	Max	OI4III
^t HI ⁽¹⁾ (0)	Test			0 0	-10 -24		ns
t _{SF} ¹ (1) (0) t _{Sl} ¹ (1) (0)	Control and data input setup times with respect to CLK (1) for stack related functions (BLT, POP, BSR, PLP) (2) AC ₀ -AC ₂			35 35 28 28	23 22 23 22		ns
tPLZ tPHZ tPZL tPZH tPZH tPHL tPHL	Propagation delay Low to high Z High to high Z High to low High Z to high Propagation delay High to low Low to high	A ₀ -A ₉	EN		12 16 14 15 33 33	35 35 25 35 40 40	ns
t _{PW}	Clock pulse width High Low			50 60	36 42		ns
tsf (1) (0) tsk (1) (0) tsl (1) (0)	Control and data input setup times with respect to CLK (1) for non-stack related functions (TSK, INC, BRT, RST) AC ₀ -AC ₂ B ₀ -B ₉ ² Test			90 90 27 29 60 60	70 70 22 24 45 45		ns
t _{HF} (1) (0) t _{HK} (1) (0)	Control and data input hold times with respect to CLK (1) AC_0 - AC_2 B_0 - B_9 ²			0 0 0	-7 -12 -12 -10		ns

^{1.} Typical values are to T $_{A}$ = 25°C and V $_{CC}$ = 5.0 volts 2. $B_{0}\text{-}B_{9}$ inputs are required to Clock (1) only. See TSK (1) and TSK (0).

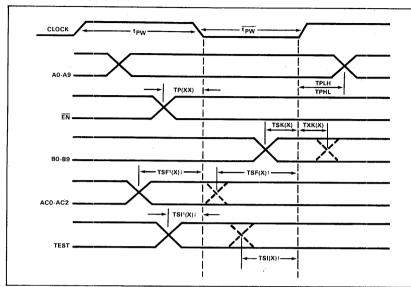
CONTROL STORE SEQUENCER

8X02

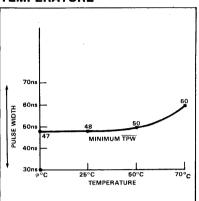
PRELIMINARY SPECIFICATION

8X02-I,XL

TIMING WAVEFORM



PULSE WIDTH (TPW) vs TEMPERATURE



DESCRIPTION

The 4-bit bipolar microprocessor slice is designed as a high-speed cascadable element intended for use in CPUs, peripheral controllers, programmable microprocessors and numerous other applications. The microinstruction flexibility of the 2901-1 will allow efficient emulation of almost any digital computing machine.

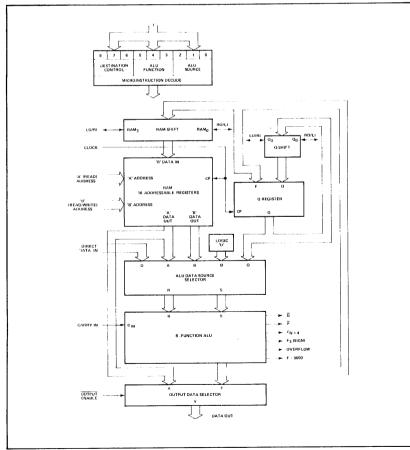
The device, as shown in the block diagram below, consists of a 16-word by 4-bit 2-port RAM, a high-speed ALU, and the associated shifting, decoding and multiplexing circuitry. The 9-bit microinstruction word is organized into 3 groups of 3 bits each and selects the ALU source operands, the ALU function, and the ALU destination register. The microprocessor is cascadable with full look-ahead or with ripple carry, has three-state outputs, and provides various status flag outputs from the ALU. Advanced low-power Schottky processing is used to fabricate this 40-lead LSI chip.

FEATURES

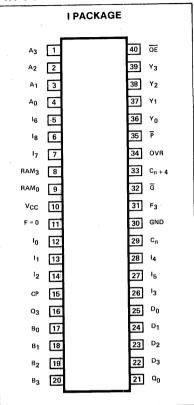
- 80ns cycle time
- 2-address architecture Independent simultaneous access to 2 working registers saves machine cycles
- 8-function ALU
 - Performs addition, 2 subtraction operations, and 5 logic functions on 2 source operands
- Flexible data source selection ALU data is selected from 5 source ports for a total of 203 source operand pairs for every ALU function
- Left/right shift independent of ALU Add and shift operations take only 1 cycle
- 4 status flags
- Carry, overflow, zero, and negative

 Expandable
 - Connect any number of 2901-1's together for longer word lengths
- Microprogrammable
 3 groups of 3 bits each for source operand, ALU function, and destination control

BLOCK DIAGRAM



PIN CONFIGURATION



PIN DESIGNATION

PIN	SYMBOL	NAME AND FUNCTION	TYPE
1-4	A ₀ -A ₃	A Address The 4 address inputs to the register stack used to select 1 register whose contents are displayed through the A port. A_0 is the LSB.	Active high
17-20	B ₀ -B ₃	B Address The 4 address inputs to the register stack used to select 1 register whose contents are displayed through the B port and into which new data can be written when the clock goes LOW. B ₀ is the LSB.	Active high
12-14, 26-28, 5-7	I ₀ -I ₈	Instruction Control The 9 instruction control lines to the 2901-1 used to determined what data sources will be applied to the ALU (I ₀₁₂), what function the ALU will perform (I ₃₄₅), and what data is to be deposited in the Q register or the register stack (I ₆₇₈).	Active high
8 16	RAM ₃ Q ₃	Shift Line A shift line at the MSB of the Q register (Q_3) and the register stack (RAM_3) . Electrically these lines are three-state outputs connected to TTL inputs internal to the 2901-1. When the destination code on I_{078} indicates a left (up) shift $(octal\ 6\ or\ 7)$ the three-state outputs are enabled and the MSB of the Q register is available on the Q_3 and the MSB of the ALU output is available on the RAM $_3$ pin. Otherwise, the three-state outputs are off (high-impedance) and the pins are electrically LS-TTL inputs. When the destination code calls for a right $(down)$ shift, the pins are used as the data inputs to the MSB of the Q register $(octal\ 4)$ and RAM $(octal\ 4\ or\ 5)$.	Three-state Active high
9 21	RAM ₀ Q ₀	Shift Line Shift lines similar to Q ₃ and RAM ₃ , at the LSB of the Q register and RAM. These pins are tied to the Q ₃ and RAM ₃ pins of the adjacent device and are used to transfer data between devices for left and right shifts of the Q register and ALU data.	Active high
22-25	D ₀ -D ₃	Direct Data Inputs A 4-bit data field which may be selected as one of the ALU data sources for entering data into the 2901-1. D ₀ is the LSB.	Active high
36-39	Y ₀ -Y ₃	Data Out The 4 data outputs of the 2901-1. These are three-state output lines. When enabled, they display either the 4 outputs of the ALU or the data on the A port of the register stack, as determined by the destination code I ₆₇₈ . Y ₀ is the LSB.	Three-state Active high
40	ŌĒ	Output Enable When OE is High, the Y outputs are disabled; when OE is Low, the Y outputs are active (high or low).	Active low
32, 35	Ğ, ₱	Carry Generate, Propagate The carry generate and propagate outputs of the 2901-1. These signals are used with the N74S182 for carry-lookahead. See Table 7 for the logic equations.	Active low
34	OVR	Overflow This pin is logically the Exclusive-OR of the carry-in and carry-out of the MSB of the ALU. At the most significant end of the word, this pin indicates that the result of an arithmetic two's complement operation has overflowed into the sign-bit. See Table 7 for logic equation.	Active high
11	F = 0	$F = 0$ This is an open collector output which goes High (off) if the data on the 4 ALU outputs F_{0-3} are all low. In positive logic, it indicates the result of an ALU operation is zero.	Active high
29	Cn	Carry In	Active high
33	C _{n+4}	Carry Out (See Table 7 for logic equations.)	Active high
15	СР	Clock The Q register and register stack outputs change on the clock Low-to-High transition. The clock Low time is internally the write enable to the 16X4 RAM which comprises the "master" latches of the register stack. While the clock is Low, the "slave" latches on the RAM outputs are closed, storing the data previously on the RAM outputs. This allows synchronous master-slave operation of the register stack.	Active high

SYSTEM DESCRIPTION

A detailed block diagram of the bipolar microprogrammable microprocessor structure is shown in Figure 1. The circuit is a 4-bit slice cascadable to any number of bits. Therefore, all data paths within the circuit are 4 bits wide. The two key elements in the Figure 1 block diagram are the 16-word by 4-bit 2-port RAM and the high-speed ALU.

Data in any of the 16 words of the Random Access Memory (RAM) can be read from the A port of the RAM as controlled by the 4-bit A address field input. Likewise, data in any of the 16 words of the RAM as defined by the B address field input can be simultaneously read from the B port of the RAM. The same code can be applied to the A select field and B select field in which case the identical file data will appear at both the RAM A port and B port outputs simultaneously.

When enabled by the RAM write enable (RAM EN), new data is always written into the file (word) defined by the B address field of the RAM. The RAM data input field is

driven by a 3-input multiplexer. This configuration is used to shift the ALU output data (F) if desired. This 3-input multiplexer scheme allows the data to be shifted up (left) 1 bit position, shifted down (right) 1 bit position, or not shifted in either direction.

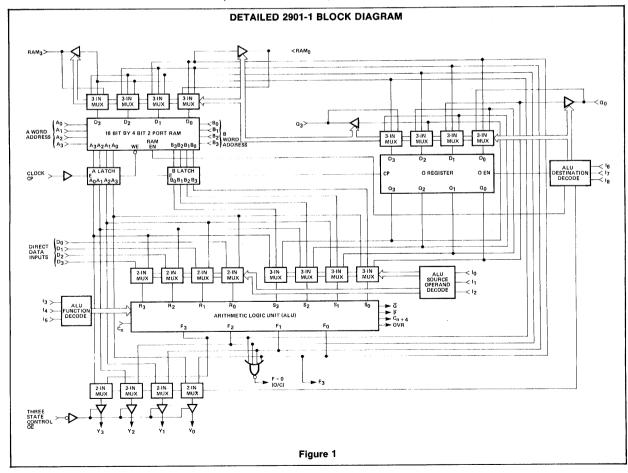
The RAM A port data outputs and RAM B port data outputs drive separate 4-bit latches. These latches hold the RAM data while the clock input is low. This eliminates any possible race conditions that could occur while new data is being written into the RAM.

The high-speed Arithmetic Logic Unit (ALU) can perform 3 binary arithmetic and 5 logic operations on the two 4-bit input words R and S. The R input field is driven from a 2-input multiplexer, while the S input field is driven from a 3-input multiplexer. Both multiplexers also have an inhibit capability; that is, no data is passed. This is equivalent to a "zero" source operand.

Referring to Figure 1, the ALU R-input multiplexer has the RAM A port and the direct

data inputs (D) connected as inputs. Likewise, the ALU S input multiplexer has the RAM A port, the RAM B port and the Q register connected as inputs.

This multiplexer scheme gives the capability of selecting various pairs of the A, B, D, Q and "0" inputs as source operands to the ALU. These 5 inputs, when taken 2 at a time, result in 10 possible combinations of source operand pairs. These combinations include AB, AD, AQ, A0, BD, BQ, B0, DQ, D0 and Q0. It is apparent that AD, AQ and A0 are somewhat redundant with BD, BQ and B0 in that if the A address and B address are the same, the identical function results. Thus, there are only 7 completely non-redundant source operand pairs for the ALU. The 2901-1 microprocessor implements 8 of these pairs. The microinstruction inputs used to select the ALU source operands are the Io, I_1 , and I_2 inputs. The definition of I_0 , I_1 , and I_2 for the 8 source operand combinations are as shown in Table 1. Also shown is the octal code for each selection.



,	MICF	юс	1	OURCE RANDS	
l ₂	l ₂ l ₁ l ₀		Octal Code	R	s
L	L	L	0	Α	Q
L	L	Н	1	Α	В
L	Н	L	2	0	Q
L	Н	Н	3	0	В
Н	L	L	4	0	Α
Н	L	Н	5	D	Α
Н	Н	L	6	D	Q ·
Н	Н	Н	. 7	D	0

Table 1 ALU SOURCE OPERAND CONTROL

The 2 source operands not fully described as yet are the D input and Q input. The D input is the 4-bit wide direct data field input. This port is used to insert all data into the working registers inside the device. Likewise, this input can be used in the ALU to modify any of the internal data files. The Q register is a separate 4-bit file intended primarily for multiplication and division routines but it can also be used as an accumulator or holding register for some applications.

The ALU itself is a high-speed arithmetic/logic operator capable of performing 3 binary arithmetic and 5 logic functions. The I₃, I₄, and I₅ microinstruction inputs are used to select the ALU function. The definition of these inputs is shown in Table 2. The octal code is also shown for reference. The normal technique for cascading the ALU of several devices is in a look-ahead carry mode. Carry generate, G. and carry propagate, P, are outputs of the device for use with a carry-look-aheadgenerator such as the N74S182. A carry-out, C_{n+4} , is also generated and is available as an output for use as the carry flag in a status register. Both carry-in (Cn) and carry-out (Cn+4) are active high.

MICRO C	ODE	ALU			
I ₅ I ₄ I ₃	Octal Code	Function	Symbol		
LLL	0	R Plus S	R+S		
LLH	1	S Minus R	S - R		
LHL	2	R Minus S	R - S		
LHH	3	RORS	R∨S		
HLL	4	RANDS	R∧S		
H L H	5	RANDS	R∧s		
HHL	6	R EX-OR S	R∀S		
ннн	7	R EX-NOR S	R∀S		

Table 2 ALU FUNCTION CONTROL

The ALU has three other status-oriented outputs. These are F_3 , F=0, and overflow (OVR). The F_3 output is the most significant (sign) bit of the ALU and can be used to

M	IICRO CODE			AM CTION	1 4 112415 1211		SUNCTION SURTED SURFEE		-			
I ₈	I ₇	I ₆	Octal Code	Shift	Load	Shift	Load	PUT	RAM₀ LO/RI	RAM ₃ LI/RO	Q₀ LO/RI	Q₃ LI/RO
L	L	۲	0	х	None	None	F→Q	· F	х	х	Х	х
L	L	Н	1	x	None	Х	None-	F	×	х	Х	x
L	Н	L	2	None	F→B	X	None	A	х	х	х	х
L	Н	Н	3	None	F→B	х	None	F	х	х	x	х
Н	L	L	4	Right (Down)	F/2 → B	Right (Down)	Q/2 → Q	F	Fo	IN ₃	Q ₀	IN ₃
Н	L	н	5	Right (Down)	F/2 → B	×	None	F	F ₀	IN ₃	Q_0	х
Н	Н	L	6	Left (Up)	2F → B	Left (Up)	2Q → Q	F	IN _o	F ₃	IN ₀	Q_3
Н	Н	н	7	Left (Up)	2F → B	х	None	F	IN ₀	F ₃	Х	Q_3

X = Don't care. Electrically, the shift pin is a TTL input internally connected to a three-state output which is in the high-impedance state.

Table 3 ALU DESTINATION CONTROL

determine positive or negative results without enabling the three-state data outputs. F_3 is non-inverted with respect to the sign bit output Y_3 . The F=0 output is used for zero detect. It is an open-collector output and can be wire OR'ed between microprocessor slices. F=0 is high when all F outputs are low. The overflow output (OVR) is used to flag arithmetic operations that exceed the available two's complement number range. The overflow output (OVR) is high when overflow exists. That is, when C_{n+3} and C_{n+4} are not the same polarity.

The ALU data output is routed to several destinations. It can be a data output of the device and it can also be stored in the RAM or the Q register. Eight possible combinations of ALU destination functions are available as defined by the I_6 , I_7 , and I_8 microinstructon inputs. These combinations are shown in Table 3.

The 4-bit data output field (Y) features three-state outputs and can be directly bus organized. An output control (OE) is used to enable the three-state outputs. When OE is high, the Y outputs are in the high-impedance state.

A 2-input multiplexer is also used at the data output such that either the A port of the RAM or the ALU outputs (F) are selected at the device Y outputs. This selection is controlled by the $l_{\rm B}$, $l_{\rm T}$, and $l_{\rm B}$ microinstruction inputs. Refer to Table 3 for the selected output for each microinstruction code combination

As was discussed previously, the RAM inputs are driven from a 3-input multiplexer. This allows the ALU outputs to be entered

non-shifted, shifted up (left) one position (X2) or shifted down (right) one position (± 2). The shifter has 2 ports; one is labeled RAM0 and the other is labeled RAM3. Both of these ports consist of a buffer-driver with a three-state output and an input to the multiplexer. Thus, in the shift up mode, the RAM3 buffer is enabled and the RAM0 multiplexer input is enabled. Likewise, in the shift down mode, the RAM0 buffer and RAM3 input are enabled. In the no-shift mode, both buffers are in the high-impedance state and the multiplexer inputs are not selected. This shifter is controlled from the $I_6,\ I_7,\$ and I_8 microinstruction inputs as defined in Table 3.

Similarly, the Q register is driven from a 3-input multiplexer. In the no-shift mode, the multiplexer enters the ALU data into the Q register. In either the shift-up or shift-down mode, the multiplexer selects the Q register data appropriately shifted up or down. The Q shifter also has 2 ports; one is labeled Q_0 and the other is $Q_3.$ The operation of these 2 ports is similar to the RAM shifter and is also controlled from $I_6,\ I_7,\$ and I_8 as shown in Table 3.

The clock input to the 2901-1 controls the RAM, the Q register, and the A and B data latches. When enabled, data is clocked into the Q register on the low-to-high transition of the clock. When the clock input is high, the A and B latches are open and will pass whatever data is present at the RAM outputs. When the clock input is low, the latches are closed and will retain the last data entered. If the RAM-EN is enabled, new data will be written into the RAM file (word) defined by the B address field when the clock input is low.

SOURCE OPERANDS AND ALU FUNCTIONS

There are eight source operand pairs available to the ALU as selected by the $I_0,\,I_1,$ and I_2 instruction inputs. The ALU can perform eight functions; five logic and three arithmetic. The $I_3,\,I_4,\,$ and I_5 instruction inputs control this function selection. The carry input, $C_{\Pi},\,$ also affects the ALU results when in the arithmetic mode. The C_{Π} input has no effect in the logic mode. When I_0 through I_5 and C_{Π} are viewed together, the matrix of Table 4 results. This matrix fully defines the ALU/source operand function for each state.

The ALU functions can also be examined on a "task" basis, i.e., add, subtract, AND, OR, etc. In the arithmetic mode, the carry will

OCTAL I ₅₄₃ , I ₂₁₀	GROUP	FUNCTION
4 0 4 1 4 5 4 6	AND	A \ Q A \ B D \ A D \ Q
3 0 3 1 3 5 3 6	OR	A
6 0 6 1 6 5 6 6	EX-OR	A ♥ Q A ♥ B D ♥ A D ♥ Q
7 0 7 1 7 5 7 6	EX-NOR	A ∀ Q A ∀ B D ∀ A D ∀ Q
7 2 7 3 7 4 7 7	INVERT	о В А
6 2 6 3 6 4 6 7	PASS	Q B A D
3 2 3 3 3 4 3 7	PASS	Q B A D
4 2 4 3 4 4 4 7	"ZERO"	0 0 0 0
5 0 5 1 5 5 5 6	MASK	A A A A A A A A A A

Table 5 ALU LOGIC MODE FUNCTIONS (Cn Irrelevant)

affect the function performed while in the logic mode, the carry will have no bearing on the ALU output. Table 5 defines the various logic operations that the 2901-1 can

perform and Table 6 shows the arithmetic functions of the device. Both carry-in low $(C_n = 0)$ and carry-in high $(C_n = 1)$ are defined in these operations.

0	I ₂₁₀ OCTAL	0	1	2	3	4	5	6	7
C I ₅ T 4 L 3	ALU Source ALU Function	A, Q	A, B	0, Q	0, B	0, A	D, A	D, Q	D, 0
0	C _n = L R Plus S C _n = H	A + Q	A + B A + B + 1	Q O + 1	B B + 1	Α Α + 1	D + A + 1	D + Q D + Q + 1	D D + 1
1	C _n = L S Minus R	Q-A-1	B-A-1	Q-1	B-1	A-1	A-D-1	Q-D-1	-D-1
2	C _n = H C _n = L R Minus S	Q-A A-Q-1	B-A A-B-1	-Q-1	-B-1	A -A-1	A-D D-A-1	Q-D D-Q-1	_D D−1
	C _n = H	A-Q	A-B	-Q	-B	-A	D-A	D-Q	D
3	RORS	AVQ	A∨B	Q	В	Α	D∨A	DVQ	D
4	R AND S	$A \wedge Q$	А∧В	0	0	0	D / A	DAQ	0
5	RANDS	$\overline{A} \wedge Q$	⊼∧B	Q	В	Α	Ď ∧ A	DAQ	0
6	R EX-OR S	A∀Q	A∀B	α	В	Α	D∀A	D∀Q	D
7	R EX-NOR S	A V Q	A∀B	ℚ	В	Ā	D∀A	D∀Q	D

+ = Plus; - = Minus, ∨= OR; ∧= AND, ∀= EX-OR

Table 4 SOURCE OPERAND AND ALU FUNCTION MATRIX

OCTAL	C _n = 0	(LOW)	C _n =	1 (HIGH)
l ₅₄₃ , l ₂₁₀	Group	Function	Group	Function
0 0 0 1 0 5 0 6	ADD	A + Q A + B D + A D + Q	ADD plus one	A + Q + 1 A + B + 1 D + A + 1 D + Q + 1
0 2 0 3 0 4 0 7	PASS B Incren		Increment	Q + 1 B + 1 A + 1 D + 1
1 2 1 3 1 4 2 7	Decrement	Q - 1 B - 1 A - 1 D - 1	PASS	Q B A D
2 2 2 3 2 4 1 7	1's Comp.	-Q - 1 -B - 1 -A - 1 -D - 1	2's Comp. (Negate)	-Q -B -A -D
1 0 1 1 1 5 1 6 2 0 2 1 2 5 2 6	Subtract (1's Comp.)	Q - A - 1 B - A - 1 A - D - 1 Q - D - 1 A - Q - 1 A - B - 1 D - A - 1 D - Q - 1	Subtract (2's Comp.)	Q - A B - A A - D Q - D A - Q A - B D - A D - Q

Table 6 ALU ARITHMETIC MODE FUNCTIONS

LOGIC FUNCTIONS FOR $\bar{G}, \; \bar{P}, \; C_{n+4}, \; \text{AND OVR}$

The four signals \overline{G} , \overline{P} , C_{n+4} , and OVR are designed to indicate carry and overflow conditions when the 2901-1 is in the add or subtract mode. Table 7 indicates the logic equations for these four signals for each of the eight ALU functions. The R and S inputs are the two inputs selected according to Table 1.

I ₅₄₃	FUNCTION	P	G	C n + 4	OVR			
0	R+S	$\overline{P_3P_2P_1P_0}$	$\overline{G_3 + P_3G_2 + P_3P_2G_1 + P_3P_2P_1G_0}$	C ₄	$C_3 \leftrightarrow C_4$			
1	S - R		Same as R+S equations, but substitute R _i for R _i in definitions					
2			Same as R+S equations, but substitute S₁ for S₁ in definitions →					
3-7	All logic operations	High	Low	High	High			

Table 7 LOGIC EQUATIONS

Definitions (+ = OR)

$$\begin{array}{llll} P_0 = R_0 + S_0 & G_0 = R_0 S_0 \\ P_1 = R_1 + S_1 & G_1 = R_1 S_1 \\ P_2 = R_2 + S_2 & G_2 = R_2 S_2 \\ P_3 = R_3 + S_3 & G_3 = R_3 S_3 \end{array}$$

$$C_4 = G_3 + P_3G_2 + P_3P_2G_1 + P_3P_2P_1G_0 + P_3P_2P_1P_0C_1$$

$$C_3 = G_2 + P_2G_1 + P_2P_1G_0 + P_2P_1P_0C_1$$

DC ELECTRICAL CHARACTERISTICS 0°C \leqslant +70°C, 4.75V, V $_{CC} \leqslant$ 5.25V

PARAMETER		TEST CONDITIONS	LIMITS			4
		1 = 0 · OON BINONS	Min	Тур	Max	UNIT
I _{CEX}	High level output voltage Output leakage current for F = 0 output	V _{CC} = 4.75V, I _{OH} = 1.6mA V _{CC} = 4.75V, V _{OH} = 5.25V, V _{IN} = V _{IH} or V _{II}	2.4		250	V μA
VOL	Low level output voltage Y	V _{CC} = 4.75V		1	0.5	
	<u>G</u>	IOL = 16mA			0.5 0.5	V
	C_{n+4} , $F=0$, OVR, \overline{P} , F_3	1 _{OL} = 10mA			0.5	ľv
v _{IH}	RAM _{3,0} , Q _{3,0} High level input voltage	I _{OL} = 6mA	İ	1	0.5	V
VIL	Low level input voltage		2.0			V
VIC	Input clamp voltage	V 00 = 4.75V 1 = 10mA	ĺ		0.8	V
11	High level input current at maximum input voltage	$V_{CC} = 4.75V, I_{ } = -18mA$ $V_{CC} = 5.25V, V_{ } = 5.5V$			-1.5	V mA
liH.	High level input current Clock, OE, A, B, D, I, Cn	V _{CC} = 5.25V, V _I = 2.7V			1.0	μA
1	RAM _{3,0} , Q _{3,0} (Note 1)				20	μΑ
l IF	Low level input current Clock, DE, A, B, D, I, Cn	$V_{CC} = 5.25V, V_1 = 0.5V$				
	RAM _{3,0} , Q _{3,0} (Note 1)				50	μΑ
los	Short circuit output current (Note 2)	V _{CC} = 5.25V	-10		-100 -40	μA mA
loz	High-Z state output current $Y_0 - Y_3$	V ₀ = 2.4V	-10		50	μΑ
	DAM O	$V_0 = 0.5V$			-50	μΑ
	RAM ₃ , ₀ , Q ₃ , ₀	$V_0 = 2.4V$			100	μΑ
¹ CC	Supply current	V ₀ = 0.5V			~100	μΑ
	117 - 2000	V _{CC} = 5.25V		165	265	mΑ

- LO/RI and RO/LI are three-state outputs internally connected to TTL inputs. Input characteristics are measured with I₆₇₈ in a state such that the three-state output is OFF.
- Not more than 1 output should be shorted at a time. Duration of the short-circuit test should not exceed 1 second.

ABSOLUTE MAXIMUM RATINGS

	PARAMETER	RATING	UNIT
VCC	Power supply voltage Input voltage Off-state output voltage Operating temperature range Storage temperature range	+7	Vdc
VIN		+5.5	Vdc
VO		+5.5	Vdc
TA		0° to +70°	°C
TSTG		-65° to +150° C	°C

AC ELECTRICAL CHARACTERISTICS T_A = 0°C to +70°C, V_{CC} = 5.0V ± 5%

		L	UNIT		
	PARAMETER	Min	Тур	Max	
tCY tRMW fQ tCL tCH	Clock Times ¹ Clock cycle time Ready-modify-write cycle Clock frequency to shift Q register Clock low period Clock high period	80 70 25 30 30	65 55 30 22 13		ns ns MHz ns ns
t _H	Hold Times ¹ Any input	0	-3		ns
^t SAB	Setup Times ¹ A, B, setup time ^{2, 3, 4}	90 t _{CL} + 30	65 t _{CL} +15		ns ns
t SBD t SD t SCn t SI012 t SI345	B destination setup time ^{2, 3, 4} D input setup time Cn input setup time ALU source control setup time ALU function setup time	t _{CL} + 15 60 50 80 75	t CL + 5 45 35 60 58		ns ns ns ns ns
^t S1678 ^t SS	ALU destination control setup time ⁵ RAM _{3·0} , input setup time Q _{3·0}	t _{CL} + 25 30 15	t _{CL} +15 20 10		ns

Setup and hold times are defined relative to the clock low-to-high edge. Inputs must be steady at all
times from the setup time prior to the clock until the hold after the clock. The setup times allow
sufficient time to perform the correct operation on the correct operation on the correct data so that
the correct ALU data can be written into one of the registers.

If the B address is used as a source operand, allow for the A, B source setup time; if it is used only for the destination address, use the B dest. setup time.

^{3.} Where 2 numbers are shown, both must be met.

^{4.} t_{CL} is the clock low time.

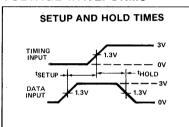
PROPAGATION DELAYS (ns) C_L = 15pF $T_A = 0^{\circ}$ to +70°C, $V_{CC} = 5.0V \pm 5\%$

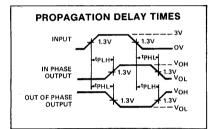
To Output	Υ	F ₃	C _{n+4}	G, F	F = 0 RL =	OVR	SHI OUTP	
Input		_			470		RAM	Q
Clock	60	55	45	50	65	50	65	45
A, B	75	65	60	70	80	60	80	_
D (arithmetic mode)	40	30	30	40	55	40	55	_
D(I = X37, logic mode)	40	30			55	_	55	_
Cn	40	25	25		45	30	50	_
l ₀₁₂	60	50	45	50	60	45	65	
l ₃₄₅	55	40	40	50_	50	45	60	
678	30	_	_				45	45
OE Enable/Disable	25/30	_	_		_	_		_
A bypassing ALU (I = 2xx)	45	_	_		_	-	_	_

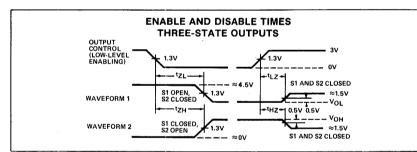
TEST LOAD CIRCUIT

FROM OUTPUT UNDER TEST CL = 15pF (SEE NOTE 1) S1 (SEE NOTE 2) S2

VOLTAGE WAVEFORMS

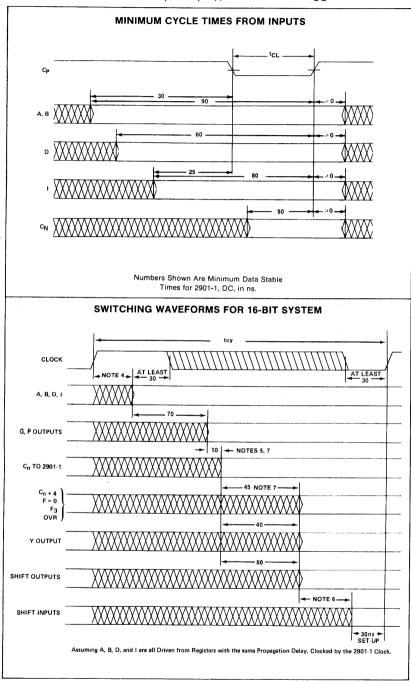






- 1. CL includes probe and jig capacitance.
- 2. All diodes 1N916 or 1N3064.
- 3. R_L = 2K

VOLTAGE WAVEFORMS (Cont'd) $T_A = 0$ °C to +70°C, $V_{CC} = 5.0V \pm 5\%$



- 4. This delay is the max, tpd of the register containing A, B, D, and I.
- 5. 10ns for lookahead carry. For ripple carry over 16 bits use 2 x (Cn \rightarrow Cn + 4).
- This is the delay associated with the multiplexer between the shift outputs and shift inputs on the 2901-1s. Normally applicable only for double length or circular shifts.
- 7. Not applicable for logic operations.

INTRODUCTION

The 2 components of the Series 3000 chip set, when combined with industry standard memory and peripheral circuits, allows the design engineer to construct highperformance processors and/or controllers with a minimum amount of auxillary logic. Features such as the multiple independent address and data buses, tri-state logic, and separate output enable lines eliminate the need for time-multiplexing of buses and associated hardware.

Each Central Processing Element represents a complete 2-bit slice through the data processing section of a computer. Several CPEs may be connected in parallel to form a processor of any desired word length. The Microprogram Control Unit controls the sequence in which microinstructions are fetched from the microprogram memory (ROM/PROM), with these microinstructions controlling the step-by-step operation of the processor.

Each CPE contains a 2-bit slice of 5 independent buses. Although they can be used in a variety of ways, typical connections are:

Input M-bus: Carries data from external

memory

Input I-bus: Carries data from input/output device

Input K-bus:

Used for microprogram mask or literal (constant) value input Output A-bus:

Connected to CPE Memory

Address Register

Output D-bus: Connected to CPE accumulator.

As the CPEs are paralleled together, all buses, data paths, and registers are correspondingly expanded.

The microfunction input bus (F-bus) controls the internal operation of the CPE. selecting both the operands and the operation to be executed upon them. The arithmetic logic unit (ALU), controlled by the microfunction decoder, is capable of over 40 Boolean and binary operations as outlined in the Function Description section of the N3002 data sheet. Standard carry look-ahead outputs (X and Y) are generated by the CPE for use with industry standard devices such as the 74S182.

A typical processor configuration is shown in Figure 1. It should be remembered that in working with slice-oriented microprocessors, the final configuration may be varied to enhance speed, reduce component count, or increase data-processing capability. One method of maximizing a processor's performance is called pipelining. To accomplish this, a group of D-type flip-flops or latches (such as the 74174 Hex D-type Flip-Flop) are connected to the microprogram memory outputs (excluding the address control field AC₀-AC₆) to buffer the current microinstruction and allow the MCU to overlap the fetch of the next instruction with the execution of the current one. The time saved in pipelining operations is the shorter of either the address set-up time to the microprogram memory (ROM/PROM) or the access time of the ROM/PROM, A convenient way of implementing pipelining is to use ROMs with on-board latches, such as the Signetics 82S115.

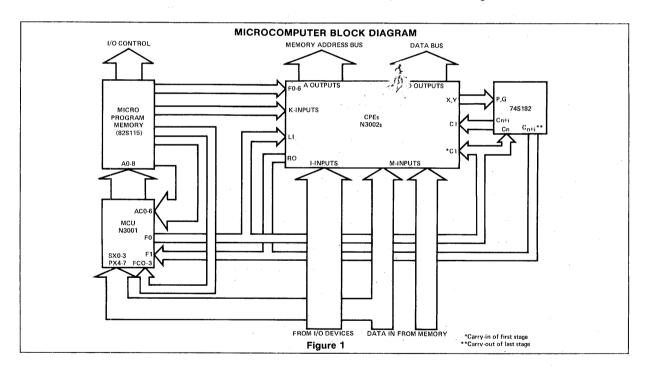


Figure 2 shows a typical microinstruction format using the 82S114 PROMs contained in the Signetics 3000 Microprocessor Designer's Evaluation Kit. Although this particular example is for a 48-bit word (6 PROMs), the allocation of bits for the mask (K-bus) and optional processor functions depends on the specific application of the system and the trade offs which the designer wishes to make.

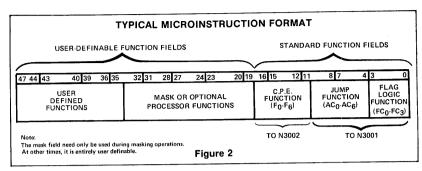
In using the K-bus, it should be kept in mind that the K inputs are always ANDed with the B-multiplexer outputs into the ALU. Bit masking, frequently done in computer control systems, can be performed with the mask supplied to the K-bus directly from the microinstruction.

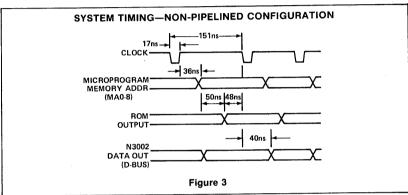
By placing the K-bus in either the all-one or all-zero condition (done with a single control bit in the microinstruction), the accumulator will either be selected or deselected, respectively, in a given operation. This feature nearly doubles the amount of microfunctions in the CPE. A description of these various microfunctions can be found in the N3002 data sheet under the heading Function Description by referring to the K-bus conditions of all-ones (11) and all-zeros (00).

The MCU controls the sequence in which microinstructions are fetched from the microprogram memory (ROM/PROM). In its classical form, the MCU would use a nextaddress field in each microinstruction. However, the N3001 uses a modified classical approach in which the microinstruction field specifies conditional tests on the MCU bus inputs and registers. The next-address logic of the MCU also makes extensive use of a row/column addressing scheme, whereby the next address is defined by a 5bit row address and 4-bit column address. Thus, from a particular address location, it is possible to jump unconditionally to any location within that row or column, or conditionally to other specified locations in one operation. Using this method, the processor functions can be executed in parallel with program branches.

As an example of this flexibility, let us assume a disk controller is being designed. As part of the sequence logic, 3 bits of the disk drive status word must be tested and all 3 must be true in order to proceed with the particular sequencing operation. In any sequencing operation using a status word for conditional branch information, there are innumerable combinations of bits which must be tested throughout the sequencing operation. Using discrete logic techniques, this would involve several levels of gating.

However, the entire operation can be done





in two microinstructions. First, the mask (Kbus) field in the microinstruction format is encoded with a one for each corresponding status bit to be tested and a zero for each bit to be discarded. The status word is input via the I-bus and ANDed with the K-bus mask using the CPE microfunction operation from F-Group 2, R-Group III. Assuming we are using low-true logic (true = 0 volts). we now test the result, which is located in the accumulator AC, for all zeros using the CPE microfunction operation from F-Group 5, R-Group III. Depending on the zero/nonzero status of AC, a one or zero will be loaded into the carryout CO bit. This bit can now be used as a condition for the next address jump calculation within the N3001 MCU. If the AC was zero (status word was true), we will jump to the next address within our controller sequence. If the AC was nonzero (status word not true), then a jump would be made back to the beginning of this 2-microinstruction loop and the test sequence repeated until the status word (all 3 bits) is true.

Figure 3 shows a typical timing diagram for a system operating in the non-pipelined mode. Keep in mind that the maximum clock rate is dependent upon the total of propagation delay times plus required setup times. It is at the designer's discretion to resolve the speed versus complexity tradeoffs

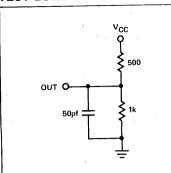
ABSOLUTE MAXIMUM RATINGS

PARAMETER	N3001/N3002	\$3001/\$3002	UNIT
Temperature under bias Storage temperature All output and supply voltages All input voltages Output currents	0 to +70 -60 to +160 -0.5 to +7 -1.0 to +5.5 100	-55 to +125 -65 to +150 -0.5 to +7 -1.0 to +5.5	°C °C V V mA

*NOTE

Stresses above those listed under "Absolute Maximum Ratings" may cause permanent damage to the device. This is a stress rating only and functional operation of the device at these or any other conditions above those indicated in the operational sections of the specification is not implied. Exposure to absolute maximum ratings for extended periods may affect device reliability.

TEST LOAD CIRCUIT



NOTE: All resistor values are typical and in ohms. TEST CONDITIONS: Input pulse amplitude of 2.5 volts Input rise and fall times of 5ns between 1 volt and 2 volts Output load of 10mA and 50pF Speed measurements are taken at the 1.5 volt level

DC ELECTRICAL CHARACTERISTICS

N3001/N3002 T_A = 0°C to +70°C S3001/S3002 T_A = -55°C to +125°C

			N3	001/N30	02	S30	01/S300	02	UNIT
	PARAMETER TEST CONDITIONS		Min	Тур	Max	Min	Тур	Max	
V _{IL} V _{IH} V _{IC}	Input voltage Low level High level Clamp	V _{CC} = 5.0V V _{CC} = 5.0V V _{CC} = 4.75V, I _C = -5mA	2.0	-0.55	0.8	2.0	-0.8	0.8	V V V
V _{OL.} V _{OH}	Output voltage Low level High level	V _{CC} = 4.75V, I _{OL} = 10mA V _{CC} = 4.75V, I _{OH} = -1mA mA ₀ - mA ₈ , ISE, FO		0.35 2.4	0.45 3.0		0.35 2.4	0.45 3.0	V V
lF	Input current Load N3001	V _{CC} = 5.25V, V _F = 0.45V CLK input EN input All other inputs		-0.21 -0.12 -0.05	-0.75 -0.50 -0.25		-0.21 -0.12 -0.05	-0.75 -0.50 -0.25	mA mA mA
I _F	Load N3002	V _{CC} = 5.25V, V _F = 0.45V F ₀ -F ₆ , CLK, K ₀ , K ₁ , EA, ED I ₀ , I ₁ , M ₀ , M ₁ , LI CI		-0.05 -0.85 -2.3	-0.25 -1.5 -4.0		-0.05 -0.85 -2.3		mA mA mA
IR	Input leakage N3001	V _{CC} = 5.25, V _R = 5.25V CLK input EN input All other inputs			120 80 40	\	/ _R = 5.5	V 120 80 40	μΑ μΑ μΑ
I _R	Input leakage N3002	V _{CC} = 5.25V, V _R = 5.25V F ₀ -F ₆ , CLK, K ₀ , K ₁ , EA, ED I ₀ , I ₁ , M ₀ , M ₁ , LI			40 60			40 60	μA μA
los	Output current Short circuit	$V_{CC} = 5.0V$ $mA_0 - mA_8$, ISE, FO	-15	-28	180 -60	-15	-28	180 -60	μA mA
(off)	Off-state	$V_{CC} = 5.25V$ $PR_0 - PR_2, mA_0 - mA_2, FO$ $V_{OUT} = 0.45V$ $mA_0 - mA_8, FO$ $V_{OUT} = 5.25V$			-100 +100	V	OUT = 5	+100	μ Α μ Α
¹cc	Power supply current N3001 Power supply current N3002	V _{CC} = 5.25V ²		170 145	240 190		170 145	250 210	mA mA

SN3001 typical values are for T_A = 25°C, V_{CC} = 5.0V
 SN3002 EN input grounded, all other inputs and outputs open. SN3002 CLK input grounded, other inputs open.

A GUIDE TO THE SELECTION OF SUPPORT COMPONENTS FOR SIGNETICS BIT SLICE MICROPROCESSORS

INTRODUCTION

Signetics family of Bipolar Bit Slice Microprocessor products consists of the 8X02 Control Store Sequencer, the 3001 microprogram control unit (MCU), the 2901-1 four bit central processing element and the 3002 two bit central processing element. These devices, all manufactured with low power Schottky technology, can be readily and efficiently interfaced with all other industry standard components, including the 7400, 74LS, 74S and Signetics' 82S and 8T families.

Figure 1 is a generalized functional block diagram of a typical Bipolar Microprocessor based system. This applications memo categorizes various components that can be used to implement each of the major functional blocks shown in Figure 1.

THE PROCESSING SECTION

The Processing Section of a computer, or "smart" controller, as shown in Figure 2, provides facilities for the transfer of data, logical and arithmetic manipulation of data, and temporary storage of data, address and status information. Cache memories are frequently used to enhance system performance by providing immediately available information and data to the CPU at high speed.

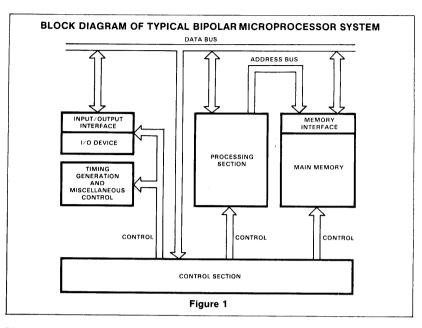
Signetics devices that can be used to implement the Processing Section of the CPU or controller are listed below in Tables 1 and 2.

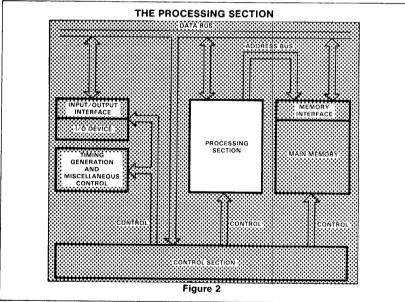
DEVICE	DESCRIPTION
2901-1	4-Bit Register and ALU Central Processing Element
N3002 74S182	Central Processing Element Carry-Look-Ahead Generator

Table 1 ALU, GENERAL PURPOSE REGISTERS, AND CARRY-LOOK-AHEAD CIRCUITS

DEVICE	DESCRIPTION
3101A	RAM (16 words by 4 bits)
82S25	Write-While-Read RAM
	(32 words by 2 bits)
82S25	RAM (16 words by 4 bits)
82S09	RAM (64 words by 9 bits)
82S116/117	RAM (256 words by 1 bit)
74S200/201	RAM (256 words by 1 bit)
74S301	RAM (256 words by 1 bit)
93415A	RAM (1024 words by 1 bit)
93425A	RAM (1024 words by 1 bit)
82S10/11	RAM (1024 words by 1 bit)

Table 2 HIGH SPEED BIPOLAR CACHE MEMORIES





A GUIDE TO THE SELECTION OF SUPPORT COMPONENTS FOR SIGNETICS BIT SLICE MICROPROCESSORS

THE CONTROL SECTION

The Control Section logic as shown in Figure 3, handles most, if not all, of the control functions. These functions are typified by operations such as decoding macroinstructions, testing of hardware or program status, initializing memory and I/O operations, manipulating data, and sampling and responding to external and internal interrupts. These control functions are implemented by executing a single microinstruction or a series of microinstructions. Through microprogramming, a structured form of control can be realized.

Reference Tables 3 through 6 for hardware selection.

DEVICE	DESCRIPTION
82S100/101	Field Programmable Logic Array (FPLA)
82S102/103	Field Programmable Gate Array

NOTE

In addition to the FPLA, all ROMs and PROMs listed below under Conditional Test and Branch Decoding and Microprogram Memory can also be used.

Table 3 MACRO INSTRUCTION DECODING

DEVICE	DESCRIPTION
82S123 82S23 82S229 82S226 82S129 82S126 82S126 82S27	PROM (32 words by 8 bits) PROM (32 words by 8 bits) PROM (256 words by 4 bits) ROM (256 words by 4 bits) PROM (256 words by 4 bits)

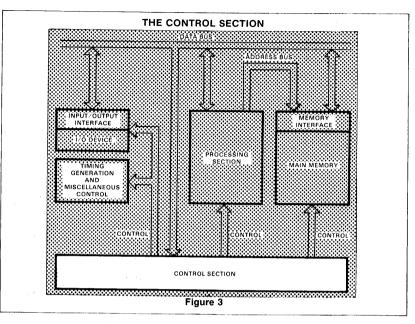
Table 4 CONDITIONAL TEST AND BRAND DECODING

DEVICE	DESCRIPTION
N3001	Microprogram Control Unit (512-word addressability)
8X02	Control Store Sequencer (1024-word addressability)

NOTE

In addition, all ROMs and PROMs listed in Conditional Test and Branch Decoding can also be used.

Table 5 MICROPROGRAM SEQUENCING



DEVICE	DESCRIPTION		
82S130	PROM (512 words by 4 bits)		
82S131	PROM (512 words by 4 bits)		
82S114	PROM with output latches		
	(256 words by 8 bits)		
82S214	ROM with output latches		
	(256 words by 8 bits)		
82\$215	ROM with output latches		
	(512 words by 8 bits)		
82S230/231	ROM (512 words by 4 bits)		
8228	ROM (1024 words by 4 bits)		
82S115	PROM with output latches		
	(512 words by 8 bits)		
82S136/137	PROM (1024 words by 4 bits)		
82S236/237	ROM (1024 words by 4 bits)		
82S280/281	ROM (1024 words by 8 bits)		
82S184/185	PROM (2048 words by 4 bits)		
82S284/285	ROM (2048 words by 4 bits)		

Table 6 MICROPROGRAM MEMORY

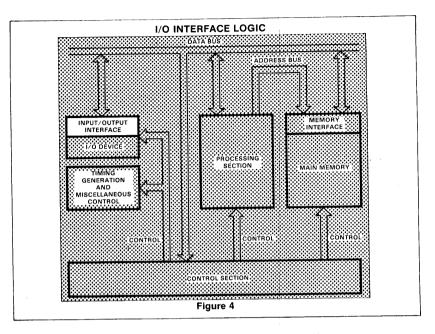
A GUIDE TO THE SELECTION OF SUPPORT COMPONENTS FOR SIGNETICS BIT SLICE MICROPROCESSORS

I/O INTERFACE LOGIC

There are many different types of bus structures (see Figure 4). To save hardware and the number of signal lines, the use of a bidirectional bus is an excellent solution for handling the data bus problem, while a tristate bus structure is ideal for the address bus, (see Table 7). In many instances, when systems of different logic families need to be interfaced with each other, logic level translators are required, (see Table 8).

DEVICE DESCRIPTION		
DEVICE	DESCRIPTION	
8T26A	Tri-state, inverting quad bus transceiver.	
8T28	Tri-state, non-inverting quad bus transceiver.	
8T31	Tri-state 8-bit bidirectional I/O port.	
8T09	Tri-state quad bus driver (40mA), with individual enable/disable	
8T10	Tri-state quad D-type bus (FF) with high drive capability	
8T13	Dual-line driver	
8T14	Triple line receiver with hysteresis	
8T15	Dual Communications EIA/MIL line driver	
8T16	Dual Communication EIA/MIL line receiver with hysteresis.	
8T23	Dual-line driver	
8T24	Triple-line receiver with hysteresis	
8T38	Quad bus transceiver	
8T95	Tri-state, non-inverting hex buffer	
8T96	Tri-state, non-inverting hex buffer	
8T97	Tri-state, inverting hex buffer	
8T98	Tri-state, inverting hex buffer	
8T100	Quad differential line drivers	
8T110	Quad differential line receivers	

Table 7 BUS BUFFERS AND DRIVERS



DEVICE	DESCRIPTION
10124	Quad TTL-to-ECL driver
10125	Quad ECL-to-TTL receiver
8T80	Quad gate TTL to High- voltage
8T90	Hex buffer TTL to High- voltage
8T25	MOS-to-TTL translator
8T18	High-voltage to TTL

Table 8 LOGIC LEVEL TRANSLATORS

MEMORY INTERFACE AND MAIN MEMORY

When dynamic MOS memory devices are used as main memory, see Figure 5, a memory refresh scheme will be needed. Usually sense amplifiers, address and data buffers, and parity generators and checkers are considered as part of the memory interface logic, (see Table 9-12).

DEVICE	DESCRIPTION
7520	Dual core-memory sense amplifiers
7521	Dual core-memory sense amplifiers
7522	Dual core-memory sense amplifiers
7523	Dual core-memory sense amplifiers
7524	Dual core-memory sense amplifiers
7525	Dual core-memory sense amplifiers
3207A-1	Quad TTL-MOS clock drivers
3207A	Quad TTL-MOS clock drivers
8T09	Tri-state quad bus driver
8T380	Tri-state bus receiver, with:
8T25	Tri-state dual sense amplifier/latch

Table 9 SENSE AMPLIFIERS AND DRIVERS

DEVICE	DESCRIPTION	
8262	9-bit parity	
74180	8-bit parity	

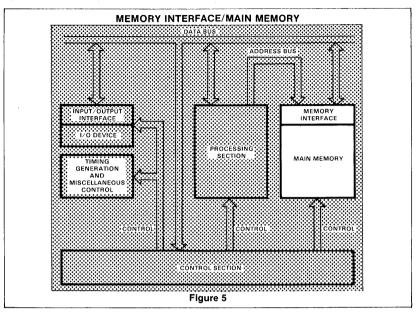
Table 10 PARITY GENERATOR AND CHECKER

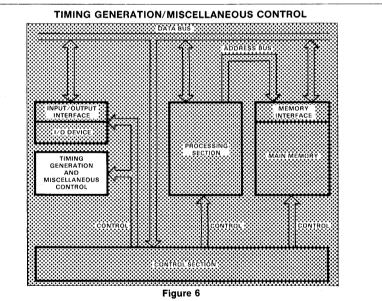
DEVICE	DEVICE DESCRIPTION	
8281	Binary counter	
8291	Binary counter	
74123	Dual one-shot	
82533	Quad 2-to-1 multiplexer	
82S34	Quad 2-to-1 multiplexer	

Table 11 MEMORY REFRESH LOGIC

DESCRIPTION
1K (1024X1) RAM
1K (1024X1) RAM
1K (256X4) RAM,
static 1K (256X4) RAM, static
1K (1024X1) RAM
4K (4096X1) RAM
4K (4096X1) RAM

Table 12 MOS MEMORY





TIMING GENERATION AND MISCELLANEOUS CONTROL

Most timing generation requirements as shown in Figure 6 can be met by using a one-shot (retriggerable monostable multivibrator). If a multiple-phase clocking system is required, additional shift registers may be used.

DEVICE	DESCRIPTION
74123	Dual-monostable multivibrator
9602	Dual-monostable multivibrator
74S194	4-bit bidirectional universal shift register
74S195	4-bit parallel-access shift registers
74S178	4-bit shift register
74S179	4-bit shift register

DESCRIPTION

The S/N3001 MCU is 1 element of a bipolar microcomputer set. When used with the S/N3002, 54/74S182, ROM or PROM memory, a powerful microprogrammed computer can be implemented.

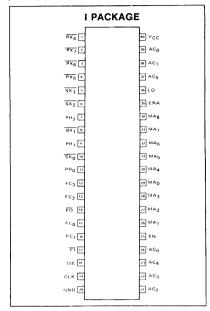
The 3001 MCU controls the fetch sequence of microinstructions from the microprogram memory. Functions performed by the 3001 include:

- · Maintenance of microprogram address register
- Selection of next microinstruction address
- Decoding and testing of data supplied via several input buses
- Saving and testing of carry output data from the central processing (CP) array
- · Control of carry/shift input data to the CP array
- · Control of microprogram interrupts

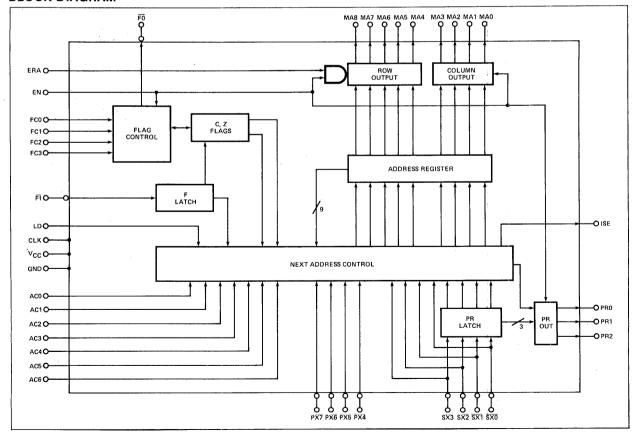
FEATURES

- · Schottky TTL process
- 45ns cycle time (typ.)
- Direct addressing of standard bipolar PROM or ROM
- 512 microinstruction addressability
- Advanced organization:
 - 9-bit microprogram address register and bus organized to address memory by row and column 4-bit program latch
 - 2-flag registers
- 11 address control functions:
 - 3 jump and test latch function 16 way jump and test instruction
- 8 flag control functions:
 - 4 flag input functions
 - 4 flag output functions

PIN CONFIGURATION



BLOCK DIAGRAM



PIN DESIGNATION

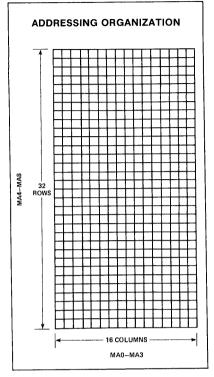
PIN	SYMBOL	NAME AND FUNCTION	TYPE
1-4	$\overline{PX}_4 - \overline{PX}_7$	Primary Instruction Bus Inputs	† · · · · · · · · · · · · · · · · · · ·
		Data on the primary instruction bus is tested by the JPX function to	Active low
		determine the next microprogram address.	
5,6,8,10	\overline{SX}_0 - \overline{SX}_3	Secondary Instruction Bus Inputs	Active low
		Data on the secondary instruction bus is synchronously loaded into the	Active low
		PR-latch while the data on the PX-bus is being tested (JPX). During a	
		subsequent cycle, the contents of the PR-latch may be tested by the JPR	İ
7044		JLL, or JRL functions to determine the next microprogram address.	İ
7,9,11	PR ₀ -PR ₂	PR-Latch Outputs	Open Collecto
		The PR-latch outputs (SX ₀ -SX ₂) are synchronously enabled by the JCE function.	open conector
		They can be used to modify microinstructions at the outputs of the	
10.40		microprogram memory or to provide additional control lines	
12,13	FC ₀ -FC ₃	Flag Logic Control Inputs	Active high
15,16	1	The flat logic control inputs are used to cross-switch the flags (C and Z)	Active riigh
		with the flag logic input (FI) and the flag logic output (FO).	
14	FO	Flag Logic Output	Active low
		The outputs of the flags (C and Z) are multiplexed internally to form the	Three-state
	!	common flag logic output. The output may also be forced to a logical	i inco state
17		0 or logical 1.	
17	FI	Flag Logic Input	Active low
	1	The flag logic input is demultiplexed internally and applied to the inputs	7101170 1017
		of the flags (C and Z). Note: The flag input data is saved in the F-latch	
10		when the clock input (CLK) is low.	
18	ISE	Interrupt Strobe Enable Output	Active high
		The interrupt strobe enable output goes to logical 1 when one of the	riotive mgm
		JZH functions are selected (see Functional Description). It can be used	
10		to provide the strobe signal required by interrupt circuits.	
19 20	CLK	Clock Input	
	GND	Ground	
21-24	AC ₀ -AC ₆	Next Address Control Function Inputs	Active high
37-39		All jump functions are selected by these control lines.	i ioni i i i i i i i i i i i i i i i i i
25	EN	Enable Input	
		When in the high state, the enable input enables the microprogram	
26-29		address, PR-latch and flag outputs.	
30-34	MA ₀ -MA ₃	Microprogram Column Address Outputs	Three-state
35	MA ₄ -MA ₈	Microprogram Row Address Outputs	Three-state
33	ERA	Enable Row Address Input	Active high
		When in the low state, the enable row address input independently	
1		disables the microprogram row address outputs. It can be used to facilitate	
36	LD	the implementation of priority interrupt systems.	
00		Microprogram Address Load Input	Active high
		When the active high state, the microprogram address load input inhibits	ĭ
l		all jump functions and synchronously loads the data on the instruction	
		buses into the microprogram address register. However, it does not inhibit	
40	Voc	the operation of the PR-latch or the generation of the interrupt strobe enable.	
	v _{CC}	+5 Volt supply	

THEORY OF OPERATION

The MCU controls the sequence of microinstructions in the microprogram memory. The MCU simultaneously controls 2 flipflops (C, Z) which are interactive with the carry-in and carry-out logic of an array of CPEs.

The functional control of the MCU provides both unconditional jumps to new memory locations and jumps which are dependent on the state of MCU flags or the state of the "PR" latch. Each instruction has a "jump set" associated with it. This "jump set" is the total group of memory locations which can be addressed by that instruction.

The MCU utilizes a two-dimensional addressing scheme in the microprogram memory. Microprogram memory is organized as 32 rows and 16 columns for a total of 512 words. Word length is variable according to application. Address is accomplished by a 9-bit address organized as a 5-bit row and 4-bit column address.



FUNCTIONAL DESCRIPTION

The following is a description of each of the eleven address control functions. The symbols shown below are used to specify row and column addresses.

MNEMONIC	FUNCTION						
rown	5-bit next row address where n is the decimal row address.						
coln	4-bit next column address where n is the decimal column address.						

Unconditional Address Control (Jump) Functions

The jump functions use the current microprogram address (i.e., the contents of the microprogram address register prior to the rising edge of the clock) and several bits from the address control inputs (ACO-AC6) to generate the next microprogram address.

Flag Conditional Address Control (Jump Test) Functions

The jump/test flag functions use the current microprogram address, the contents of the selected flag or latch, and several bits from the address control function to generate the next microprogram address.

JUMP FUNCTION TABLE

MNEMONIC	NAME AND FUNCTION							
JCC	Jump in current column. AC_0 - AC_4 are used to select 1 of 32 row addresses in the current column, specified by MA_0 - MA_3 , as the next address.							
JZR	Jump to zero row. AC_0 - AC_3 are used to select 1 of 16 column addresses in row ₀ , as the next address.							
JCR	Jump in current row. AC_0 - AC_3 are used to select 1 of 16 addresses in the current row, specified by MA_4 - MA_8 , as the next address.							
JCE	Jump in current column/row group and enable PR-latch outputs. AC ₀ AC ₂ are used to select 1 of 8 row addresses in the current row group specified by MA ₇ -MA ₈ , as the next row address. The current column is specified by MA ₀ -MA ₃ . The PR-latch outputs are asynchronously enabled.							

JUMP/TEST FUNCTION TABLE

MNEMONIC	NAME AND FUNCTION
JFL	Jump/test F-latch. AC ₀ -AC ₃ are used to select 1 of 16 row addresses in the current row group, specified by MA ₈ , as the next row address. If the current column group, specified by MA ₃ , is col ₀ -col ₇ , the F-latch is used to select col ₂ or col ₃ as the next column address. If MA ₃ specifies column group col ₈ -col ₁₅ , the F-latch is used to select col ₁₀ or col ₁₁ as the next column address.
JCF	Jump/test C-flag. AC_0 - AC_2 are used to select 1 of 8 row addresses in the current row group, specified by MA_7 and MA_8 , as the next row address. If the current column group specified by MA_3 is col_0 - col_7 , the C-flag is used to select col_2 or col_3 as the next column address. If MA_3 specifies column group col_8 - col_{15} , the C-flag is used to select col_{10} or col_{11} as the next column address.
JZF	Jump/test Z-flag. Identical to the JCF function described above, except that the Z-flag, rather than the C-flag, is used to select the next column address.
JPR	Jump/test PR-latch. AC_0 - AC_2 are used to select 1 of 8 row addresses in the current row group, specified by MA_7 and MA_8 , as the next row address. The 4 PR-latch bits are used to select 1 of 16 possible column addresses as the next column address.
JLL	Jump/test leftmost PR-latch bits. AC_0 - AC_2 are used to select 1 of 8 row addresses in the current row group, specified by MA_7 and MA_8 , as the next row address. PR_2 and PR_3 are used to select 1 of 4 column addresses in col_4 through col_7 as the next column address.
JRL	Jump/test rightmost PR-latch bits. AC ₀ and AC ₁ are used to select 1 of 4 high-order row addresses in the current row group, specified by MA ₇ and MA ₈ , as the next row address. PR ₀ and PR ₁ are used to select 1 of 4 possible column addresses in col ₁₂ through col ₁₆ as the next column address.
JPX	Jump/test PX-bus and load PR-latch. AC_0 and AC_1 are used to select 1 of 4 row addresses in the current row group, specified by MA_6 - MA_8 , as the next row address. PX_4 - PX_7 are used to select 1 of 16 possible column addresses as the next column address. SX_0 - SX_3 data is locked in the PR-latch at the rising edge of the clock.

PX-Bus and PR-Latch **Conditional Address Control** (Jump/Test) Functions

The PX-bus jump/test function uses the data on the primary instruction bus (PX4-PX₇), the current microprogram address. and several selection bits from the address control function to generate the next microprogram address. The PR-latch jump/ test functions use the data held in the PRlatch, the current microprogram address. and several selection bits from the address control function to generate the next microprogram address

Flag Control Functions

The flag control functions of the MCU are selected by the 4 input lines designated FC₀-FC₃. Function code formats are given in "Flag Control Function summary."

The following is a detailed description of each of the 8 flag control functions.

Flag Input Control **Functions**

The flag input control functions select which flag or flags will be set to the current value of the flag input (FI) line.

Data on FI is stored in the F-latch when the clock is low. The content of the F-latch is loaded into the C and/or Z flag on the rising edge of the clock.

Flag Output Control Functions

The flag output control functions select the value to which the flag output (FO) line will be forced.

FLAG CONTROL FUNCTION TABLE

MNEMONIC	FUNCTION DESCRIPTION
SCZ	Set C-flag and Z-flag to FI. The C-flag and the Z-flag are both set to the value of FI.
STZ	Set Z-flag to FI. The Z-flag is set to the value of FI. The C-flag is unaffected.
STC	Set C-flag to FI. The C-flag is set to the value of FI. The Z-flag is unaffected.
HCZ	Hold C-flag and Z-flag. The values in the C-flag and Z-flag are unaffected.

FLAG OUTPUT CONTROL FUNCTION TABLE

MNEMONIC	FUNCTION DESCRIPTION
FF0	Force FO to 0. FO is forced to the value of logical 0.
FF <u>C</u>	Force FO to C. FO is forced to the value of the C-flag.
FFZ	Force FO to Z. FO is forced to the value of the Z-flag.
FF1	Force FO to 1. FO is forced to the value of logical 1.

FLAG CONTROL FUNCTION SUMMARY

TYPE MNEMONI	MNEMONIC	DESCRIPTION	FC ₁	
	SCZ	Set C-flag and Z-flag to f	0	0
Flag STZ Input STC	Set Z-flag to f	l o	1 1	
	Set C-flag to f	1	0	
	HCZ	Hold C-flag and Z-flag	1	1

TYPE	MNEMONIC	DESCRIPTION	FC ₃	2
Flag Output	FF0 FFC FFZ FF1	Force FO to 0 Force FO to C-flag Force FO to Z-flag Force FO to 1	0 1 0	0 0 1

LOAD FUNCTION		NEXT	ROW		NEXT COL				
LD 0	MA ₈	7 5	6 See Ad	5 dress	4 Contro	MA ₃ I Functi	on Sum	1	0
. 1	0	X ₃	X ₂	X ₁	X ₀	X 1 dileti	X ₆	ттату Х ₅	X_4

f = Contents of the F-latch

xn = Data on PX- or SX-bus line n (active low)

ADDRESS CONTROL FUNCTION SUMMARY

MNEMONIC	DESCRIPTION	FUNCTION AC ₆ 5 4 3 3								NEXT ROW					NEXT COL			
			5	4	. 3	2	1	0	MA ₈	7	6	5	4	MA ₃	2	1	0	
JCC	Jump in current column	0	0	ď₄	d_3	d_2	d₁	· d _o	d₄	d_3	d ₂	d₁	d ₀					
JZR	Jump to zero row	0	1 .	0	d_3	d ₂	d.	d _o	0	. 0	0	0	0	m_3	m ₂	m ₁	m _o	
JCR	Jump in current row	0	1	1	d_3	d_2	d₁		-	-	•	-	٠ ١	d_3	d_2	d ₁	d₀	
JCE	Jump in column/enable	1	1	1	O 3			d₀	m ₈	m ₇	m ₆	m ₅	m ₄	d_3	d_2	d₁	d_0	
JFL	Jump/test F-latch	1	'n	'n	٦	d ₂	d₁	d ₀	m ₈	m ₇	d ₂	. a ₁	d ₀	m_3	m_2	m_1	m_0	
JCF	Jump/test C-flag	i	0	0	d ₃	α ₂	d ₁	d ₀	m ₈	d_3	d_2	d₁	d_0	m_3	0	1	f	
JZF	Jump/test Z-flag	_ '	0		'	d_2 .	d ₁	· d _o	m ₈	m_7	d_2	d₁	d₀	m_3	0	1	С	
		7	Ü	1	1	d_2	d₁	d₀	m ₈ .	m_7	d ₂	d₁	d₀ l	m_3	Ω	4	-	
	Jump/test PR-latch	1	1	0	0	d_2	d₁	d _o	m ₈	m _z	d ₂	d,	1		_	'-	2	
	Jump/test left PR bits	1	1	0	1	d ₂	d₁	ď,	m ₈	m_7			d ₀	p_3	p_2	p_1	p_0	
JRL	Jump/test right PR bits	1	1	1	1	1	d.	.° I		,	d ₂	d₁	d ₀	0	7	p_3	p_2	
JPX	Jump/test PX-bus	1	1	1	1	'n	41	d₀	m ₈	m_7	1	d ₁	d₀	7	1	P1	p_0	
					,	U	u_1	d ₀	m_8	m_7	m_6	d₁	d_0	X_7	Xe	Χs	X.	

NOTE

dn = Data on address control line n

mn = Data in microprogram address register bit n

Pn = Data in PR-latch bit n

xn = Data on PX-bus line n (active low)

f,c,z = Contents of F-latch, C-flag, or Z-flag, respectively

STROBE FUNCTIONS

The load function of the MCU is controlled by the input line designated LD. If the LD line is active high at the rising edge of the clock, the data on the primary and secondary instruction buses, PX_4-PX_7 and SX_0-SX_3 , is loaded into the microprogram address register. PX_4-PX_7 are loaded into MA_0-MA_7 and SX_0-SX_3 are loaded into MA_4-MA_7 . The high-order bit of the microprogram address register MA_8 is set to a logical 0. The bits from the primary instruction bus select 1 of 16 possible column addresses. Likewise, the bits from the secondary instruction bus select 1 of the first 16 row addresses.

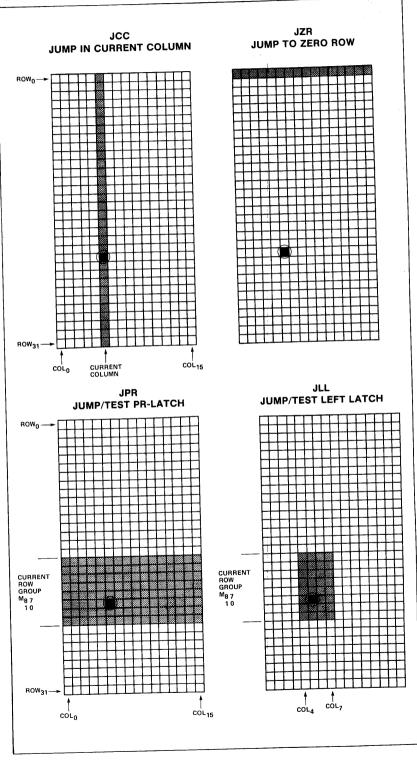
The MCU generates an interrupt strobe enable on the output line designated ISE. The line is placed in the active high state whenever a JZR to col15 is selected as the address control function. Generally, the start of a macroinstruction fetch sequence is situated at row₀ and col₁₅ so the interrupt control may be enabled at the beginning of the fetch/execute cycle. The interrupt control responds to the interrupt by pulling the enable row address (ERA) input line low to override the selected next row address from the MCU. Then by gating an alternative next row address on to the row address lines of the microprogram memory, the microprogram may be forced to enter an interrupt handling routine. The alternative row address placed on the microprogram memory address lines does not alter the contents of the microprogram address register. Therefore, subsequent jump functions will utilize the row address in the register, and not the alternative row address, to determine the next microprogram address.

Note, the load function always overrides the address control function on AC_0 - AC_6 . It does not, however, override the latch enable or load sub-functions of the JCE or JPX instruction, respectively. In addition, it does not inhibit the interrupt strobe enable or any of the flag control functions.

JUMP SET DIAGRAMS

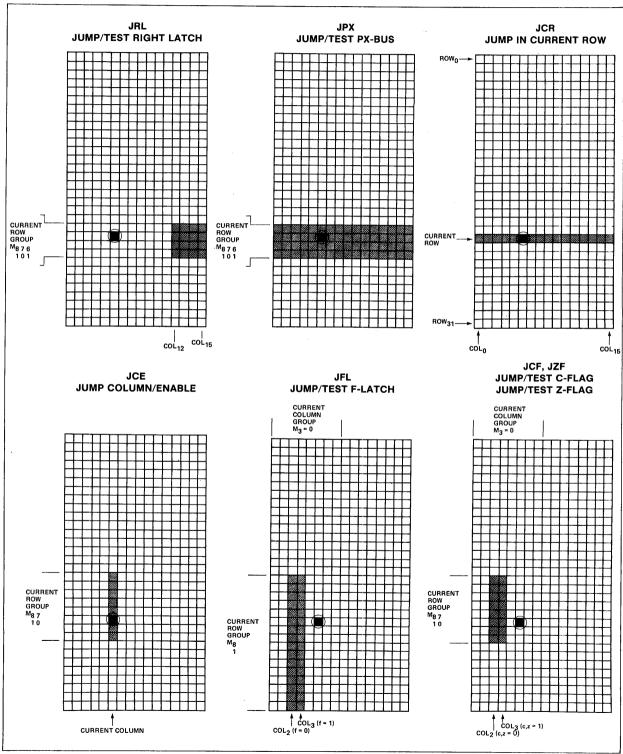
The following 10 diagrams illustrate the jump set for each of the 11 jump and jump/ test functions of the MCU. Location 341 indicated by the circled square, represents 1 current row (row $_{21}$) and current column (col $_{5}$) address. The dark boxes indicate the microprogram locations that may be selected by the particular function as the next address.

JUMP SET DIAGRAMS



S/N3001-I

JUMP SET DIAGRAMS (Cont'd)



MICROPROGRAM CONTROL UNIT

N3001 T_A = 0°C to +70°C, V_{CC} = 5.0V, ± 5%

S/N3001-I

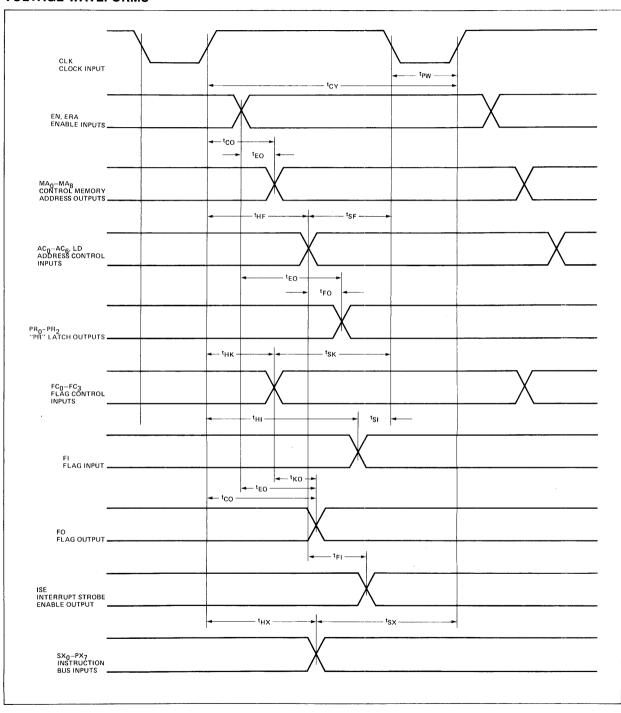
AC ELECTRICAL CHARACTERISTICS S3001 TA = -55°C to +125°C, V_{CC} = 5.0V ± 10%

			N3001					
	PARAMETER	Min	Typ¹	Max	Min	Typ¹	Max	UNIT
tCY	Cycle Time ²	60	45		95	45		ns
t _{PW}	Clock Pulse Width	17	10		40	10		ns
PVV	Control and Data Input Set-Up Times:						ľ	
tSF	LD, AC ₀ -AC ₆ (Set to "1"/"0")	20	3/14		20	3/14		ns
t _{SK}	FC ₀ , FC ₁	7	5		10	5		ns
tSX	PX ₄ -PX ₇ (Set to "1"/"0")	28	4/13		35	4/13		ns
tSI	FI (Set to "1"/"0")	12	-6/0		15	-6/10		ns
t _{SX}		15	5		35	5		ns
-51	Control and Data Input Hold Times:							
t _{HF}	LD, AC ₀ -AC ₆ (Hold to "1"/"0")	4	-3/-14		5	-3/-14		ns
tHK		4	-5		10	-5		ns
tHX	DV DV (11.11) (411.000)	0	-4/-13	1	25	-4/-13		ns
tHI	FI (Hold to "1"/"0")	16	6.5/0		22	6.5/0		ns
tuv	·SX ₀ -SX ₃	0	-5		25	-5		ns
tco			17/24	36	10	17/24	45	ns
1 .00	(mA ₀ -mA ₈ , FO) (tPHL/tPLH)			Ì	1	1		
t _{KO}	Propagation Delay from Control Inputs FC ₂ and FC ₃ to Flag Out (FO)		13	24		13	50	ns
t _{FO}	Propagation Delay from Control Inputs AC ₀ -AC ₆ to Latch Outputs (PR ₀ -PR ₂)		21	32		21	50	ns
t _{EO}	Propagation Delay from Enable Inputs EN and ERA to Outputs (MA ₀ -mA ₈ , FO, PR ₀ -PR ₂)		17	26		17	35	ns
t _{FI}	Propagation Delay from Control Inputs AC ₀ -AC ₆ to Interrupt Strobe Enable Output (ISE)		20	32		20	40	ns

^{1.} Typical values are for TA = 25°C and 5.0 supply voltage. 2. S3001: tCY = tWP + tSF + tCO

S/N3001-I

VOLTAGE WAVEFORMS



DESCRIPTION

The N3002 Central Processing Element (CPE) is one part of a bipolar microcomputer set. The N3002 is organized as a 2-bit slice and performs the logical and arithmetic functions required by microinstructions. A system with any number of bits in a data word can be implemented by using multiple N3002s, the N3001 microcomputer control unit, the N74S182 carry look-ahead unit and ROM or PROM memory.

FEATURES

- 45ns cycle time (typ)
- . Easy expansion to multiple of 2 bits
- 11 general purpose registers
- Full function accumulator
- Useful functions include:

2's complement arithmetic Logical AND, OR, NOT, exclusive-NOR Increment, decrement

Shift left/shift right
Bit testing and zero detection
Carry look-ahead generation
Masking via K-bus
Conditioned clocking allowing non-

Conditioned clocking allowing nondestructive testing of data in accumulator and scratchpad

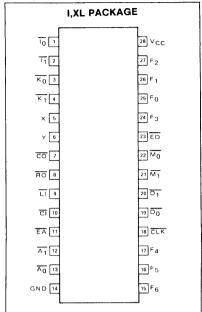
- 3 input buses
- 2 output buses
- Control bus

FUNCTION TRUTH TABLE

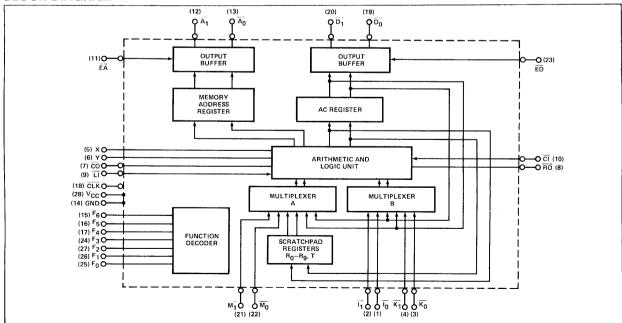
FUNCTION GROUP	F ₆	F ₅	F ₄
0	0	0	0
1	0	0	1
2	0	1	0
3	0	1	1
4	1	0	0
5	1	0	1
6	1	1	0
7	1	1	1

REGISTER GROUP	REGISTER	F ₃	F ₂	F ₁	F ₀
	R₀	0	0	0	0
	R₁	0	0	0	1
	R ₂	0	0	1	0
!	R ₃	0	0	1	1
	R₄	0	1	0	0
	R ₅	0	1	0	1
'	R ₆	0	1	1	0
	R ₇	0	1	1	1
	R ₈	1	0	0	0
	R ₉	1	0	0	1
	T	1	1	0	0
	AC	1	1	0	1
11	Т	1	0	1	0
11	AC	1	0	1	1
111	Т	1	1	1	0
111	AC	1	1	1	1

PIN CONFIGURATION



BLOCK DIAGRAM



PIN DESIGNATION

PIN	SYMBOL	NAME AND FUNCTION	TYPE
1, 2	·	External Bus Inputs	Active low
,		The external bus inputs provide a separate input port for external input devices.	
3, 4	K₀-K₁	Mask Bus Inputs	Active low
		The mask bus inputs provide a separate input port from the microprogram memory,	ĺ
	1	to allow mask or constant entry.	
5, 6	X, Y	Standard Carry Look-Ahead Cascade Outputs	Active high
		The cascade outputs allow high speed arithmetic operations to be performed when they	
		are used in conjunction with the 74S182 Look-Ahead Carry Generator	1
7	CO	Ripple Carry Out	Active low
	·	The ripple carry output is only disabled during shift right operations.	Three-state
8	RO	Shift Right Output	Active low
		The shift right output is only enabled during shift right operations.	Three-state
9		Shift Right Input	Active low
10	CI	Carry Input	Active low
11	ĒĀ	Memory Address Enable Input	Active low
		When in the low state, the memory address enable input enables the memory	
		address outputs (A_0-A_1) .	
12-13	$\overline{A_0}$ - $\overline{A_1}$	Memory Address Bus Outputs	Active low
		The memory address bus outputs are the buffered outputs of the memory	Three-state
		address register (MAR).	
14	GND	Ground	A -45 bi-b
14-17,	F ₀ -F ₆	Micro-Function Bus Inputs	Active high
24-27	5117	The micro-function bus inputs control ALU function and register selection.	
18	CLK	Clock Input	Active low
19-20	$\overline{D_0}$ - $\overline{D_1}$	Memory Data Bus Outputs The memory data bus outputs are the buffered outputs of the full function	Three-state
		·	Tillee-state
21-22	$\overline{M_0}$ - $\overline{M_1}$	accumulator register (AC). Memory Data Bus Inputs	Active low
21-22	1010-1011	The memory data bus inputs provide a separate input port for memory data.	Active low
23	ED	Memory Data Enable Input	Active low
20	"	When in the low state, the memory data enable input enables the memory	,
		data outputs (D_0 - D_1).	
28	V _{CC}	+5 Volt Supply	
	1 100 1	10 York Supply	

SYSTEM DESCRIPTION

Microfunction Decoder and K-Bus

Basic microfunctions are controlled by a 7-bit bus (F_0-F_6) which is organized into 2 groups. The higher 3 bits (F_4-F_6) are designated as F-Group and the lower 4 bits (F_0-F_3) are designated as the R-Group. The F-Group specifies the type of operation to be performed and the R-Group specifies the registers involved.

The F-Bus instructs the microfunction decoder to:

- Select ALU functions to be performed
- Generate scratchpad register address
- Control A and B multiplexer

The resulting microfunction action can be:

- Data transfer
- · Shift operations
- · Increment and decrement
- Initialize stack
- Test for zero conditions
- 2's complement addition and subtraction
- Bit masking
- Maintain program counter

A and B Multiplexers

A and B multiplexers select the proper 2 operands to the ALU.

A multiplexer selects inputs from one of the following:

- M-bus (data from main memory)
- Scratchpad registers
- Accumulator

B multiplexer selects inputs from one of the following:

- I-bus (data from external I/O devices)
- Accumulator
- K-bus (literal or masking information from micro-program memory)

Scratchpad Registers

- Contains 11 registers (R₀-R₉, T)
- Scratchpad register outputs are multiplexed to the ALU via the A multiplexer
- Used to store intermediate results from arithmetic/logic operations
- Can be used as program counter

Arithmetic/Logic Unit (ALU)

The ALU performs the arithmetic and logic operations of the CPE.

Arithmetic operations are:

- 2's complement addition
- Incrementing
- Decrementing
- Shift left
- Shift right

Logical operations are:

- Transfer
- AND
- Inclusive-ORExclusive-NOR
- · Logic complement

ALU operation results are then stored in the accumulator and/or scratchpad registers. For easy expansion to larger arrays, carry look-ahead outputs (X and Y) and cascading shift inputs (LI, RO) are provided.

Accumulator

- Stores results from ALU operations
- The output of accumulator is multiplexed into ALU via the A and B multiplexer as one of the operands

Input Buses

M-bus: Data bus from main memory

- Accepts 2 bits of data from main memory into CPE
- Is multiplexed into the ALU via the A multiplexer

I-bus: Data bus from input/output devices

- Accepts 2 bits of data from external input/output devices into CPE
- Is multiplexed into the ALU via the B multiplexer

K-bus: A special feature of the N3002 CPE

- During arithmetic operations, the K-bus can be used to mask portions of the field being operated on
- Select or remove accumulator from operation by placing K-bus in all "1" or all "0" state respectively
- During non-arithmetic operation, the carry circuit can be used in conjunction with the K-bus for word-wise-OR operation for bit testing
- Supply literal or constant data to CPE

Output Buses

A-bus and Memory Address Register

- Main memory address is stored in the memory address register (MAR)
- Main memory is addressed via the A-bus
- MAR and A-bus may also be used to generate device address when executing I/OO instrucinstructions
- · A-bus has Tri-State outputs

D-bus: Data bus from CPE to main memory or to I/O devices

- Sends buffered accumulator outputs to main memory or the external I/O devices
- · D-bus has Tri-State outputs

FUNCTION DESCRIPTION

FUNCT	FUNCTION DESCRIPTION									
F GROUP	R GROUP	K BUS	NAME	EQUATION	DESCRIPTION					
0	ı	XX	_	$R_n + (AC \wedge K) + CI \rightarrow R_n, AC$	Logically AND AC with the K-bus. Add the result to $R_{\rm n}$ and carry input (CI). Deposit the sum in AC and $R_{\rm n}$.					
		00	ILR	R _n + Cl → R , AC	Conditionally increment R $_{\rm II}$ and load the result in AC. Used to load AC from R $_{\rm II}$ or to increment R $_{\rm II}$ and load a copy of the result in AC.					
		11	ALR	AC + R _n + Cl → R _n , AC	Add AC and CI to $R_{\rm II}$ and load the result in AC. Used to add AC to a register. If $R_{\rm II}$ is AC, then AC is shifted left one bit position.					
0	11	xx		M + (AC ^ K) + CI → AT	Logically AND AC with the K-bus. Add the result to CI and the M-bus. Deposit the sum in AC or T.					
2		00	ACM	M + CI → AT	Add CI to M-bus. Load the result in AC or T, as specified. Used to load memory data in the specified register, or to load incremented memory data in the specified register.					
		11	AMA	M + AC + CI → AT	Add the M-bus to AC and CI, and load the result in AC or T, as specified. Used to add memory data or incremented memory data to AC and store the sum in the specified register.					
0	Ш	xx	-	ATL ^ (ÎL ^ KL) → RO LI ^ ((IH ^ KH) ^ ATH] → ATH [ATL ^ (IL ^ KL)] [ATH ^ (IH ^ KH)] → ATL	None					
		00	SRA	AT _L → RO AT _H → AT _L L _I → AT _H	Shift AC or T, as specified, right one bit position. Place the previous low order bit value on RO and fill the high order bit from the data on LI. Used to shift or rotate AC or T right one bit.					
1	1	XX		$K \lor R_n \to MAR$ $R_n + K + CI \to R_n$	Logically OR R_n with the K-bus. Deposit the result in MAR. Add the K-bus to Rn and Cl. Deposit the result in R _n .					
		00	LMI	Rn→MAR, Rn+Cl→Rn	Load MAR from $\mathbf{R}_{\mathbf{n}}.$ Conditionally increment $\mathbf{R}_{\mathbf{n}}.$ Used to maintain a macro-instruction program counter.					
		11	DSM	11 → MAR, Rn – 1 + Cl → Rn	Set MAR to all ones. Conditionally decrement R_Π by one. Used to force MAR to its highest address and to decrement Rn.					
1	П	XX		KVM → MAR M + K + CI → AT	Logically OR the M-bus with the K-Bus. Deposit the result in MAR. Add the K-bus to the M-bus and CI. Deposit the sum in AC or T.					
		00	LMM	M → MAR, M + CI → AT	Load MAR from the M-bus. Add CI to the M-bus. Deposit the result in AC or T. Used to load the address register with memory data for macro-instructions using indirect addressing.					
		11	LDM	11 → MAR M - 1 + CI → AT	Set MAR to all ones. Subtract one from the M-bus. Add CI to the difference and deposit the result in AC or T, as specified. Used to load decremented memory data in AC or T.					

FUNCTION DESCRIPTION (Cont'd)

F	R	K.	NASAF	FOUNTION	DECORIDATION
GROUP	GROUP	BUS	NAME	EQUATION	DESCRIPTION
1	III.	XX		(AT ∨ K) + (AT ^ K) + CI → AT	Logically OR the K-bus with the complement of AC or T, as specified. Add the result to the logical AND of specified register with the K-bus. Add the sum to CI. Deposit the result in the specified register.
		00	CIA	AT + CI - AT	Add CI to the complement of AC or T, as specified. Deposit the result in the specified register. Used to form the 1's or 2's complement of AC or T.
		11	DCA	AT - 1 + CI → AT	Subtract one from AC or T, as specified. Add CI to the difference and deposit the sum in the specified register. Used to decrement AC or T.
2	I	XX	_	(AC ^K) - 1 + Cl → R _n	Logically AND the K-bus with AC. Subtract one from the result and add the difference to CI. Deposit the sum in R _n .
		00	CSR	CI - 1 → R _n (See Note 1)	Subtract one from CI and deposit the difference in Rn. Used to conditionally clear or set Rn to all 0's or 1's, respectively.
		11	SDR	AC - 1 + Cl → R _n (See Note 1)	Subtract one from AC and add the difference to CI. Deposit the sum in R_n . Used to store AC in R_n or to store the decremented value of AC in R_n .
2	H	XX	-	(AC ^ K) - 1 + CI → AT (See Note 1)	Logically AND the K-bus with AC. Subtract one from the result and add the difference to CI. Deposit the sum in AC or T, as specified.
		00	CSA	CI - 1 → AT (See Note 1)	Subtract one from CI and deposit the difference in AC or T. Used to conditionally clear or set AC or T.
		11	SDA	AC - 1 + CI → AT (See Note 1)	Subtract one from AC and add the difference to CI. Deposit the sum in AC or T. Used to store AC in T, or decrement AC, or store the decremented value of AC in T.
2	. III	XX	- <u>-</u> .	(I ^ K) - 1 + CI → AT (See Note 1)	Logically AND the data of the K-bus with the data on the I-bus. Subtract one from the result and add the difference to CI. Deposit the sum in AC or T, as specified.
		00	CSA	CI − 1 → AT	Subtract one from CI and deposit the difference in AC or T. Used to conditionally clear or set AC or T.
		11	LDI	I – 1 + CI → AT	Subtract one from the data on the I-bus and add the difference to CI. Deposit the sum in AC or T, as specified. Used to load input bus data or decremented input bus data in the specified register.
3	I	xx	_	$R_n + (AC \wedge K) + CI \rightarrow R_n$	Logically AND AC with the K-bus. Add R_{Π} and CI to the result. Deposit the sum in R_{Π} .
		.00	INR	R _n + Cl → R _n	Add CI to R_n and deposit the sum in R_n . Used to increment
		11	ADR	AC + R _n + CI - R _n	R _n . Add AC to R _n . Add the result to CI and deposit the sum in R _n . Used to add the accumulator to a register or to add the incremented value of the accumulator to a register.
3	· 11	xx	-	M + (AC ^ K) + CI → AT	Logically AND AC with the K-bus. Add the result to CI and the M-bus. Deposit the sum in AC or T.
	·	00	ACM	M + CI → AT	Add CI to M-bus. Load the result in AC or T, as specified. Used to load memory data in the specified register, or to load incremented memory data in the specified register.
!		11	AMA	M + AC + CI → AT	Add the M-bus to AC and CI, and load the result in AC or T, as specified. Used to add memory data or incremented memory data to AC and store the sum in the specified register.

NOTE

^{1. 2&#}x27;s complement arithmetic adds 111 . . . 11 to perform subtraction of 000 01. \cdot .

FUNCTION DESCRIPTION (Cont'd)

F GROUP	R GROUP	K BUS	NAME	EQUATION	DESCRIPTION			
3	111	XX		AT + (I ^ K) + CI → AT	Logically AND the K-bus with the I-bus. Add CI and the contents of AC or T, as specified, to the result. Deposit the sum in the specified register.			
		00	INA	AT + CI → AT	Conditionally increment AC or T. Used to increment AC or T.			
		11	AIA	I + AT + CI → AT	Add the I-bus to AC or T. Add CI to the result and deposit sum in the specified register. Used to add input data incremented input data to the specified register.			
4	B	XX		$CI \times (R_n \land AC \land K) \rightarrow CO$ $R_n \land (AC \land K) \rightarrow R_n$	Logically AND the K-bus with AC. Logically AND the result with the contents of R_n . Deposit the final result in R_n . Logically OR the value of CI with the word-wise OR of the bits of the final result. Place the value of the carry OR on the carry output (CO) line.			
		00	CLR	CI → CO, O → R _n	Clear R $_{\mbox{\scriptsize n}}$ to all O's. Force CO to CI. Used to clear a register and force CO to CI.			
		11	ANR	$CI \lor (R_n \land AC) \rightarrow CO$ $R_n \land AC \rightarrow R_n$	Logically AND AC with $R_{\rm fl}$. Deposit the result in $R_{\rm fl}$. Force CO to one if the result is non-zero. Used to AND the accumulator with a register and test for a zero result.			
4	=	XX	_	CI ∨ (M ^ AC ^ K) → CO M ^ (AC ^ K) → AT	Logically AND the K-bus with AC. Logically AND the result with the M-bus. Deposit the final result in AC or T. Logically OR the value of CI with the word-wise OR of the bits of the final result. Place the value of the carry OR on CO.			
		00	CLA	CI → CO, O → AT	Clear AC or T, as specified, to all O's. Force CO to CI. Used to clear the specified register and force CO to CI.			
		11	ANM	CI ∨ (M ∧ AC) → CO M ∧ AC → AT	Logically AND the M-bus with AC. Deposit the result in AC or T. Force CO to one if the result is non-zero. Used to AND M-bus data to the accumulator and test for a zero result.			
4	III	xx	-	CI ∨ (AT ∧1 ∧ K) → CO AT ∧ (I ∧ K) → AT	Logically AND the I-bus with the K-bus. Logically AND the result with AC or T. Deposit the final result in the specified register. Logically OR CI with the word-wise OR of the final result. Place the value of the carry OR on CO.			
		00	CLA	CI → CO, O → AT	Clear AC or T, as specified, to all O's. Force CO to CI. Used to clear the specified register and force CO to CI.			
		11	ANI	CI ∨ (AT ^ I) → CO AT ^ 1 → AT	Logically AND the I-bus with AC or T, as specified. Deposit the result in the specified register. Force CO to one if the result is non-zero. Used to AND the I-bus to the accumulator and test for a zero result.			
5	ı	XX	_	$CI \vee (R_{\Pi} \wedge K) \rightarrow CO$ $K \wedge R_{\Pi} \rightarrow R_{\Pi}$	Logically AND the K-bus with R _n . Deposit the result in Rn. Logically OR CI with the word-wise OR of the result. Place the value of the carry OR on CO.			
		00	CLR	CI → CO, O → R _n	Clear Rn to all O's. Force CO to CI. Used to clear a register and force CO to CI.			
		11	TZR	CI ∨ R _n → CO R _n → R _n	Force CO to one if Rn is non-zero. Used to test a register for zero. Also used to AND K-bus data with a register for masking and, optionally, testing for a zero result.			
5	11	xx	_	CI ∨ (M ∧ K) → CO K ∧ M → AT	Logically AND the K-bus with the M-bus. Deposit the result in AC or T, as specified. Logically OR CI with the word-wise OR of the result. Place the value of the carry OR on CO.			
		00	CLA	CI → CO, O → AT	Clear AC or T, as specified, to all O's. Force CO to CI. Used to clear the specified register and force CO to CI.			
		11	LTM	CI ∨ M → CO M → AT	Load AC or T, as specified, from the M-bus. Force CO to one if the result is non-zero. Used to load the specified register from memory and test for a zero result. Also used to AND the K-bus with the M-bus for masking and, optionally, testing for a zero result.			

FUNCTION DESCRIPTION (Cont'd)

UNCT								
F GROUP	R GROUP	K BUS	NAME	EQUATION	DESCRIPTION			
5	Ш	XX	-	CI ∨ (AT ^K) → CO K ^AT → AT	Logically AND the K-bus with AC or T, as specified. Deposit the result in the specified register. Logically OR CI with the word-wise OR of the result. Place the value of the carry OR on CO.			
		. 00	CLA	CI → CO, O → AT	Clear AC or T, as specified, to all O's. Force CO to CI. Used to clear the specified register and force CO to CI.			
		11	TZA	CI v AT - CO AT - AT	Force CO to one if AC or T, as specified, is non-zero. Used to test the specified register for zero. Also used to AND the K-bus to the specified register for masking and, optionally, testing for a zero result.			
6	I	xx	_	$CI \lor (AC \land K) \to CO$ $R_n \lor (AC \land K) \to R_n$	Logically OR CI with the word-wise OR of the logical AND of AC and the K-bus. Place the result of the carry OR on CO. Logically OR R _n with the logical AND of AC and the K-bus. Deposit the result in R _n .			
		00	NOP	CI → CO, R _n → R _n	Force CO to CI. Used as a null operation or to force CO to CI.			
		11	ORR	$CI \lor AC \to CO$ $R_n \lor AC \to R_n$	Force CO to one if AC is non-zero. Logically OR AC with R _n . Deposit the result in R _n . Used to OR the accumulator to a register and, optionally, test the previous accumulator value for zero.			
6	11	xx	_	$CI \lor (AC \land K) \rightarrow CO$ $M \lor (AC \land K) \rightarrow AT$	Logically OR CI with the word-wise OR of the logical AND of AC and the K-bus. Place the carry OR on CO. Logically OR the M-bus, with the logical AND of AC and the K-bus. Deposit the final result in AC or T.			
		00	LMF	CI → CO, M → AT	Load AC or T, as specified, from the M-bus. Force CO to CI. Used to load the specified register with memory data and force CO to CI.			
		11	ORM	CI ∨ AC → CO M ∨ AC → AT	Force CO to one if AC is non-zero. Logically OR the M-bus with AC. Deposit the result in AC or T, as specified. Used to OR M-bus with the AC and, optionally, test the previous value of AC for zero.			
6	III	xx	_	$CI \lor (I \land K) \rightarrow CO$ $AT \lor (I \land I) \rightarrow AT$	Logical OR CI with the word-wise OR of the logical AND of the I-bus and the K-bus. Place the carry OR on CO. Logically AND the K-bus with the I-bus. Logically OR the result with AC or T, as specified. Deposit the final result in the specified register.			
		00	NOP	CI → CO, AT → AT	Force CO to CI. Used as a null operation or to force CO to CI.			
		11	ORI	CI ∨ I → CO I ∨ AT →	Force CO to one if the data on the I-bus is non-zero. Logically OR the I-bus to AC or T, as specified. Deposit the result in the specified register. Used to OR I-bus data with the specified register and, optionally, test the I-bus data for zero.			
7	1	xx	_	$CI \lor (R_n \land AC \land K) \rightarrow CO$ $R_n \oplus (AC \land K) \rightarrow R_n$	Logically OR CI with the word-wise OR of the logical AND of R _n and AC and the K-bus. Place the carry OR on CO. Logically AND the K-bus with AC. Exclusive-NOR the result with R _n . Deposit the final result in R _n .			
		00	CMR	CI → CO, Rn → Rn	Complement the contents of R _n . Force CO to CI.			
		11	XNR	CI ∨ (R _n ^ AC) → CO R _n ⊕ AC → R _n	Force CO to one if the logical AND of AC and Rn is non-zero. Exclusive-NOR AC with Rn. Deposit the result in Rn. Used to exclusive-NOR the accumulator with a register.			

FUNCTION DESCRIPTION (Cont'd)

F GROUP	R GROUP	K BUS	NAME	EQUATION	DESCRIPTION
7	II ·	XX	_	CI ∨ (M ^ AC ^ K) → CO M ⊕ (AC ^ K) → AT	Logically OR CI with the word-wise OR of the logical AND of AC and the K-bus and M-bus. Place the carry OR on CO. Logically AND the K-bus with AC. Exclusive NOR the result with the M-bus. Deposit the final result in AC or T.
		00	LCM	CI → CO, M → AT	Load the complement of the M-bus into AC or T, as specified. Force CO to CI.
		11	XNM	CI ∨ (M ^ AC) → CO M ⊕ AC → AT	Force CO to one if the logical AND of AC and the M-bus is non-zero. Exclusive-NOR AC with the M-bus. Deposit the result in AC or T, as specified. Used to exclusive-NOR memory data with the accumulator.
7	111	xx	_	CI ∨ (AT ∧ I ∧ K) → CO AT ⊕ (I K) → AT	Logically OR CI with the word-wise OR of the logical AND of the specified register and the I-bus and K-bus. Place the carry OR on CO. Logically AND the K-bus with the I-bus. Exclusive-NOR the result with AC or T, as specified. Deposit the final result in the specified register.
		00	CMA	CI → CO ĀT → AT	Complement AC or T, as specified. Force CO to CI.
		11	XNI	CI ∨ (AT ^ I) → CO I ⊕ AT → AT	Force CO to one if the logical AND of the specified register and the I-bus is non-zero. Exclusive-NOR AC with the I-bus. Deposit the result in AC or T, as specified. Used to exclusive-NOR input data with the accumulator.

FUNCTION DESCRIPTION KEY

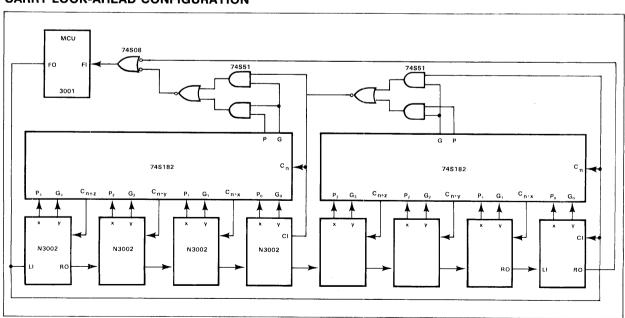
SY	MBOL	MEANING
I,K	(,M	Data on the I, K, and M buses, respectively
CI	,LI	Data on the carry input and left
C	O,RO	input, respectively Data on the carry output and
Rr	ו	right output, respectively Contents of register n includ-
AC	2	ing T and AC (R-Group I) Contents of the accumulator
AT	Г	Contents of AC or T, as specified
M	٩R	Contents of the memory address register
L,I	Ⅎ	As subscripts, designate low
+	į	and high order bit, respectively 2's complement addition
-		2's complement subtraction
^		Logical AND
\ \		Logical OR
Φ		Exclusive-NOR
		Deposit into

AC ELECTRICAL CHARACTERISTICS N3001 = T $_{A}$ = 0°C to +70°C, V $_{CC}$ = 5V ± 5% S3001 = T $_{A}$ = -55°C to +125°C, V $_{CC}$ = 5V ± 10%

		N3002			S3002		
PARAMETER	Min	Тур*	Max	Min	Тур*	Max	UNIT
tCY Clock Cycle Time	70	45		120	45		ns
tWP Clock Pulse Width	17	10		42	10		ns
tFS Function Input Set-Up Time (F ₀ through F ₆)	48	-23 → 35		70	-23 → 35		ns
Data Set-Up Time: t DS I_0 , I_1 , M_0 , M_1 , K_0 , K_1 t SS I_1 , I_2	40 21	12 → 29 0 → 7		60 30	12 → 29 0 → 7		ns ns
Data and Function Hold Time: $tFH = F_0 \ through \ F_6 \\ tDH = 1_0, \ I_1, \ M_0, \ M_1, \ K_0, \ K_1 \\ tSH = LI, \ CI$	4 4 12	0 -28 → -11 -7 → 0		5 5 15	0 -28 → -11 -7 → 0		ns ns ns
Propagation Delay to X, Y, RO from: tXF Any Function Input tXD Any Data Input tXT Trailing Edge of CLK tXL Leading Edge of CLK	13	28 16 → 20 33 18 → 40	52 33 48 70	13	28 16 → 20 33 18 → 40	65 65 75 90	ns ns ns ns
Propagation Delay to CO from: tCL Leading Edge of CLK tCT Trailing Edge of CLK tCF Any Function Input tCD Any Data Input tCC CI (Ripple Carry)	16	24 → 44 30 → 40 25 → 35 17 → 23 9 → 13	70 56 52 55 20		24 - 44 30 - 40 25 - 35 17 - 23 9 - 13	90 100 75 65 30	ns ns ns ns
Propagation Delay to A_0 , A_1 , D_0 , D_1 from: tDL Leading Edge of CLK tDE Enable Input ED, EA		17 → 25 10 → 12	40 20		17 → 25 10 → 12	75 35	ns ns

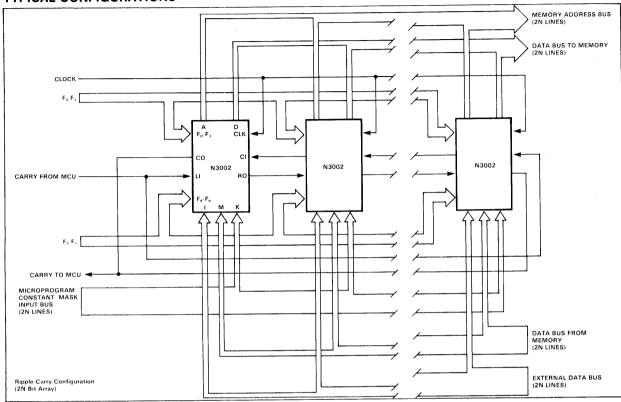
^{*} NOTE

CARRY LOOK-AHEAD CONFIGURATION

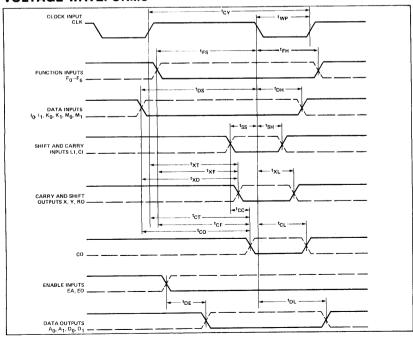


Typical values are for TA = 25°C and typical supply voltage.

TYPICAL CONFIGURATIONS



VOLTAGE WAVEFORMS



INTRODUCTION

A Microcomputer Designed for Control

The 8X300 is a microcomputer designed for control. It features:

Execution Speed

- 250ns instruction execution time
- Direct address capability—up to 8192 16-bit words of program memory
- Eight 8-bit general purpose registers
- Simultaneous data transfer and data edit in a single instruction cycle time
- n-way branch or n-entry table lookup in 2 instruction cycle times
- 8X300 instructions operate with equal speed on 1-bit, 2-bit, 3-bit, 4-bit, 5-bit, 6-bit, 7-bit, or 8-bit data formats

The 8X300 instruction set features controloriented instructions which directly access variable length input/output and internal data fields. These instructions provide very high performance for moving and interpreting data. This makes the 8X300 ideal in switching, controlling, and editing applications

Direct Processing of External Data

The 8X300 I/O system is treated as a set of

internal registers. Therefore data from external devices may be processed (tested, shifted, added to, etc.) without first moving them to internal storage. In fact, the entire concept is to treat data at the I/O interface no differently than internal data. This concept extends to the software which allows variables at the input/output system to be named and treated in the same way as data in storage.

Separate Program Storage and Data Storage

The storage concept of the 8X300 is to separate program storage from data storage. Program storage is implemented in read-only memory in recognition of the fact that programs for control applications are fixed and dedicated. The benefits of using read-only memory are that great speeds may be obtained at lower cost than if read/write memory were used, and that program instructions reside in a non-volatile medium and cannot be altered by system power failures.

8X300 Architecture

Figure 1 of the 8X300 data sheet illustrates the 8X300 architecture. The 8X300 contains an Arithmetic Logic Unit (ALU), Program Counter, and an Address Register. Eight 8bit general purpose registers are also provided, including 7 working registers and an auxiliary register which performs as a working register and also provides an implied operand for many instructions. The 8X300 registers are shown in Figure 1 of the 8X300 data sheet and are summarized below:

Control Registers include:

- Instruction—A 16-bit register containing the current instruction
- Program Storage Address Register (AR)—A 13-bit register containing the address of the current instruction being accessed from Program Storage
- Program Counter (PC)—A 13-bit register containing the address of the next instruction to be read from Program Storage

Data Registers include:

- Working Registers (WR)—Seven 8-bit registers for data storage
- Overflow (OVF)—A 1-bit register that retains the most significant bit position carry from ALU addition operation. Arithmetically treated as
- Auxiliary (AUX)—An 8-bit register. Source of implied operand for arithmetic and logical instructions. May be used as a working register.

A crystal external to the CPU may be used to generate the CPU system clock. The CPU executes 8 instruction types.

DESCRIPTION

The Signetics 8X300 Microcontroller is a monolithic, high-speed microprocessor implemented with bipolar Schottky technology. As the central processing unit, CPU, it allows 16-bit instructions to be fetched, decoded and executed in 250ns. A 250ns instruction cycle requires maximum memory access of 65ns, and maximum I/O device access of 35ns.

Microcontroller instructions operate on 8-bit, parallel data. Logic is distributed along the data path within the Microcontroller. Input data can be rotated and masked before being subject to an arithmetic or logical operation; and output data can be shifted and merged with the input data, before being output to external logic. This allows 1-to 8-bit I/O and data memory fields to be accessed and processed in a single instruction cycle.

PROGRAM STORAGE INTERFACE

Program Storage is typically connected to the A0-A12 (A12 is least significant bit) and I0-I15 signal lines. An address output on A0-A12 identifies one 16-bit instruction word in program storage. The instruction word is subsequently input on I0-I15 and defines the Microcontroller operations which are to follow.

The Signetics 82S115 PROM, or any TTL compatible memory, may be used for program storage.

I/O DEVICES INTERFACE

An 8-bit I/O bus, called the Interface Vector (IV) data bus, is used by the Microcontroller to communicate with 2 fields of I/O devices. The complementary \overline{LB} and \overline{RB} signals identify which field of the I/O devices is selected.

Both I/O data and I/O address information can be output on the IV bus. The SC and WC signals are typically used to distinguish between I/O data and I/O address information as follows:

sc wc

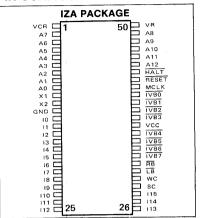
- 0 I/O address is being output on IV bus
- 0 1 I/O data is being output on IV bus
- 0 0 I/O data is expected on the IV,

The Signetics 82SXXX series RAM, and the 8T32/33 may be attached to the IV bus.

FEATURES

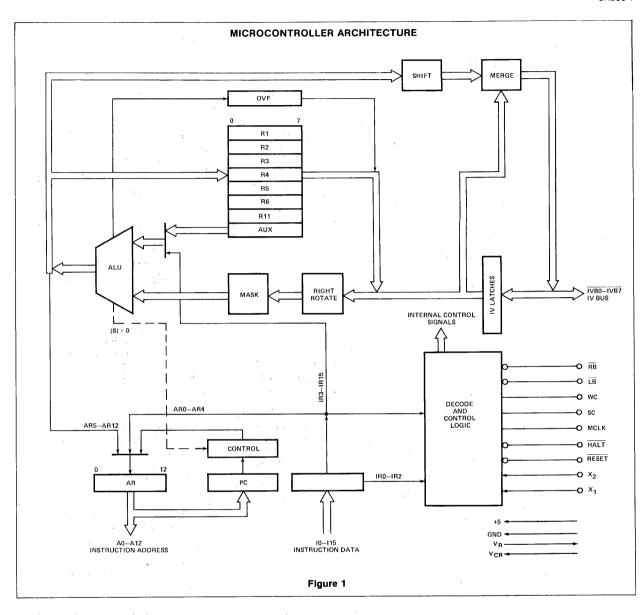
- 185ns instruction decode and execute delay (with Signetics 8T32/33 I/O port)
- . Eight 8-bit working registers
- Single instruction access to 1-bit, 2-bit, 3bit or 8-bit field on I/O bus
- Separate instruction address, instruction, and I/O data buses
- · On-chip oscillator
- Bipolar Schottky technology
- . TTL inputs and outputs
- Tri-state output on I/O data bus
- +5 volt operation from 0° to 70°C

PIN CONFIGURATION



PIN DESIGNATION

PIN	SYMBOL	NAME AND FUNCTION	TYPE
2-9, 45-49	A0-A12:	Instruction address lines. A high level equals "1." These outputs directly address up to 8192 words of program storage. A12 is least significant bit.	Active high
13-28	10-115:	Instruction lines. A high level equals "1." Receives instructions from Program Storage. I ₁₅ is least significant bit.	Active high
33-36, 38-41	ĪVBO-ĪVB7	Interface Vector (IV) Bus. A low level equals "1." Bidirectional tri-state lines to communicate with I/O devices. IVB7 is least significant bit.	Three-state Active low
42	MCLK:	Master Clock. Output to clock I/O devices, and/or provide synchronization for external logic	
30	WC:	Write Command. High level output indicates data is being output on the IV Bus.	Active high
29	SC:	Select Command. High level output indicates that an address is being output on the IV Bus.	Active high
31	LB:	Left Bank. Low level output to enable one of two sets of I/O devices (LB is the complement of RB).	Active low
32	RB:	Right Bank. Low level output to enable one of two sets of I/O devices (\overline{RB} is the complement of \overline{LB}).	Active low
44	HALT:	Low level is input to stop the Microcontroller.	Active Iow
43	RESET:	Low level is input to initialize the Microcontroller.	Active low
10-11	X1, X2:	Inputs for an external frequency determining crystal. May also be interfaced to logic or test equipment.	
50	VR	Reference voltage to pass transistor.	
1	VCR	Regulated output voltage from pass transistor.	
37	V _{CC} :	5V power connection.	
12	GND:	Ground.	



- General working register
- R2 -- General working register
- R3 General working register
- R4 General working register
- General working register
- R6 - General working register
- R11 General working register
- AUX General working register. Contains second term for arithmetic or logical operations.
- OVF The least-significant bit of this register is used to reflect overflow status resulting from the most recent ADD operation (see Instruction Set
 - Summary).
- Instruction Register (IR)
 - Holds the 16-bit instruction word currently being executed.

Program Counter (PC)

- Normally contains the address of the current instruction and is incremented to obtain the next instruction address.

Address Register (AR)

 A 13-bit register containing the address of the current instruction.

Table 1 INTERNAL REGISTERS

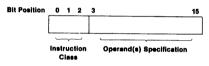
INSTRUCTION CYCLE

Each Microcontroller operation is executed in 1 instruction cycle, which may be as short as 250ns. The Microcontroller generates MCLK to synchronize external logic to the instruction cycle. Instruction cycles are subdivided into quarter cycles. MCLK is an output during the last quarter cycle.

During the third quarter cycle of an instruction, an address is output on A0-A12, identifying the location in program storage o are next instruction word. This instruction word defines the next instruction, which must be input on I0-I15 during the first quarter cycle of the next instruction cycle (see Table 2).

Instruction Set Summary

The 16-bit instruction word input on I0-I15 is decoded by the instruction decode logic to implement events that are to occur during the remainder of the instruction cycle. Generally the 16-bit instruction word is decoded as follows:



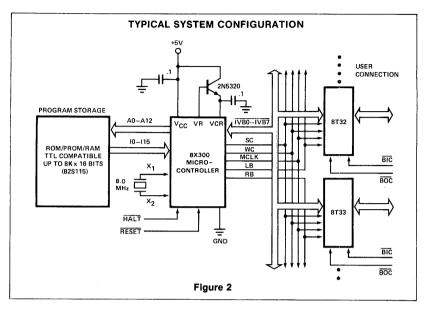
A detailed usage of the 13 "operand(s) specification" bits is given in following sections.

Three operation code bits allow for 8 instruction classes. The 8 instruction classes are summarized in Table 3. Each entry is referred to as an "instruction class" because the unique architecture of the Interpreter allows a number of powerful variations to be specified by the 13 operand (s) specification bits. A complete description of instruction formats and some instruction examples are provided in the Microprocessor Applications manual.

Data Processing

The Microcontroller architecture includes eight 8-bit working registers, an arithmetic logic unit (ALU), an overflow register, and the 8-bit IV Bus. Internal 8-bit data paths connect the registers and IV Bus to the ALU inputs, and the ALU output to the registers and IV Bus. Data processing logic is distributed along these internal 8-bit data paths. Rotate and mask logic precedes the ALU on the data entry path. Shift and merge logic follows the ALU on the data output path. All 4 sets of logic can operate on 8 data bits in a single instruction cycle (See Figure 1).

When less than 8 bits of data are specified for output to the IV bus by the ALU, the data field (shifted if necessary) is inserted into the prior contents of the IV bus latches. The



INSTRUCTION AND IV BUS DATA INPUT	DATA	ADDRESS AND IV BUS CHANGING	ADDRESS AND IV BUS DATA VALID MCLK=HIGH
← ¼ cycle →	← ¼ cycle →	← ¼ cycle →	← ¼ cycle →

Table 2 INSTRUCTION CYCLE

IV bus latches contain data input at the start of an instruction. This data in the IV bus latches will be specified in the instruction as a) IV bus source data or b) data from an automatic read when the IV bus is specified as a destination. Therefore, IV bus bit positions outside an inserted bit field are unmodified.

Data Addressing

Sources and destinations of data are specified using a 5-bit octal number. The source and/or destination of data to be operated upon is specified in a single instruction word.

Referring to Figure 1, the Auxiliary register (address 00) is the implied source of the second argument for ADD, AND or XOR operations.

IVL and IVR are write-only registers used only as a destination. They have addresses and are treated as registers, but in reality they do not exist. When IVL is specified as a destination or the D field = 20-27₈, then LB = 'low', RB = 'high' are generated; when IVR is specified as a destination or the D field = 30-37₈, then RB = low, LB = 'high' are generated.

When IVL or IVR is specified as the destination in an instruction, SC is also activated and data is placed on the IV bus. If IVL or IVR is specified as a source of data, the source data is all zeroes.

INSTRUCTION SEQUENCE CONTROL

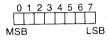
The Address Register and Program Counter are used to generate addresses for accessing an instruction. The Address Register is used to form the instruction address, and in all but 3 instructions (XEC, NZT, and JMP) the address is copied into the Program Counter. The instruction address is formed in 1 of 3 ways:

- For all instructions but the JMP, XEC, and a satisfied NZT, the Program Counter is incremented by 1 and placed in the Address Register.
- For the JMP instruction, the full 13-bit address field from the JMP instruction is placed into the Address Register and copied into the Program Counter.
- For the XEC and NZT instructions, the high order 5- or 8-bits of the Program Counter are combined with 8- or 5lower-order bits of ALU output (XEC or NZT) and placed in the Address Register. For the NZT instruction, it is also copied into the Program Counter.

INSTRUCTION SET

The 8X300 Microcontroller has a repertoire of 8 instruction classes which allow the user to test input status lines, set or reset output control lines, and perform high speed input/output data transfers. All instructions are 16 bits in length and each is fetched, decoded and executed in 250ns.

Data is respresented as an 8-bit byte; bit positions are numbered from left to right, with the least significant bit in position 7.



Within the 8X300, all operations are performed on 8-bit bytes. Arithmetic operations use 8-bit, unsigned 2's complement arithmetic.

INSTRUCTION FORMATS

The general 8X300 instruction format is:

0 1 2	3	4	5	6	7	8	9	10	11	12	13	14	15
Op Code						0	ре	ran	d(s	5)			

All instructions are specified by a 3-bit Operation (Op Code) field. The operand may consist of the following fields: Source

(S) field, Destination (D) field, Rotate/ Length (R/L) field, Immediate (I) field, or (Program Storage) Address (A) field.

The instructions are divided into 5 format types, based on the Op Code and the Operand(s), as shown in Figure 3.

INSTRUCTION FORMATS						
<u>OPERATIONS</u> (REGISTER TO REGISTER)						
MOVE AND ADD XOR						
0 1 2 3 4 5 6 7 8 9 10 11 12 13 14 15						
OP S R D						
Type I						
OPERATIONS (REGISTER TO I/O, I/O TO REGISTER, I/O TO I/O)						
MOVE ADD AND XOR						
0 1 2 3 4 5 6 7 8 9 10 11 12 13 14 15						
OP S (LENGTH) D						
Type II						
OPERATIONS						
XEC XMIT* NZT						
0 1 2 3 4 5 6 7 8 9 10 11 12 13 14 15						
OP S I						
Type III						
<u>OPERATIONS</u>						
XEC XMIT*						
NZT						
0 1 2 3 4 5 6 7 8 9 10 11 12 13 14 15 OP S L						
CODE S (LENGTH) Type IV						
1,001						
OPERATIONS						
JMP						
0 1 2 3 4 5 6 7 8 9 10 11 12 13 14 15						
OP A						
Type V						
'NOTE						
If XMIT, S actually represents the destination.						
Figure 3						

OPERATION	FORMAT	RESULT	NOTES
MOVE	1,11	Content of data field specified by {S, R/L} replaces data in field specified by {D, R/L.}	
ADD	1,11	Sum of AUX and data specified by $\left\{S,R/L\right\}$ replaces data in field specified by $\left\{D,R/L\right\}$.	If S and D both are registers, then R/L specifies a right rotate
AND	1,11	Logical AND of AUX and data specified by $\{S, R/L\}$ replaces data in field specified by $\{D, R/L.\}$	of the register specified by S.
XOR	1,11	Logical exclusive OR of AUX and data specified by {S, R/L} replaces data in field specified by {D, R/L.}	
ХМІТ	III,IV	The literal value I replaces the data in the field specified by {S,L.}	If S is an I/O address then I is limited to range 00-37. Other- wise I is limited to range 000-377.
NZT	III,IV	If the data in the field specified by $\{S, L\}$ equals zero, perform the next instruction in sequence. If the data specified by $\{S, L\}$ is not equal to zero, execute the instruction at address determined by using the literal I as an offset to the Program Counter.	If S is an I/O address then I is limited to range 00-37. Otherwise I is limited to range 000-377.
XEC	III,IV	Perform the instruction at address determined by applying the sum of the literal I and the data specified by {S, L} as an offset to the Program Counter. If that instruction does not transfer control, the program sequence will continue from the XEC instruction location.	The offset operation is performed by reducing the value of PC to the nearest multiple of 32 (if I = 00-37) or 256 (if I = 000-377) and adding the offset.
JMP	V	The address value A replaces contents of the Program Counter.	A limited to the range 0-177778.

Table 3 8X300 INSTRUCTION SUMMARY

INSTRUCTION FIELDS

Op Code Field (3-Bit Field)

The Op Code field is used to specify 1 or 8 8X300 instructions as shown in Table 4.

OP CODE OCTAL VALUE	INS	STRUCTION
0	MOVE	S,R/L,D
1	ADD	S,R/L,D
2	AND	S,R/L,D
3	XOR	S,R/L,D
4	XEC	I,R/L,S or I,S
5	NZT	I,R/L,S or I,S
6 7	XMIT JMP	I,R/L,D or I,D A

Table 4 OP CODE FIELD OCTAL **ASSIGNMENTS**

S.D Fields (5-Bit Fields)

The S and D fields specify the source and destination of data for the operation defined by the Op Code field. The Auxiliary Register is an implied second source for the instructions ADD, AND and XOR, each of which require two source fields. That is, instructions of the form,

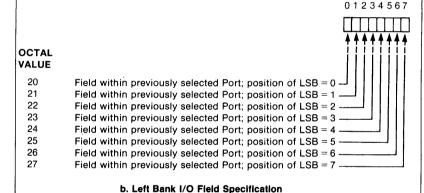
imply a third operand, say Z, located in the Auxiliary Register so that the operation which takes place is actually X + Z, with the result stored in Y.

The S and/or D fields may specify a register, or a 1 to 8-bit I/O field. S and D field value assignments in octal are shown in Table 5. 08-178 is used to specify 1 of 7 working registers (R1-R6, R11), the Auxiliary Register, the Overflow Register, or IVL and IVR write-only registers.

OCTAL VALUE		OCTAL VALUE	
00	AUX-Auxiliary Register	10	OVF-Overflow register-Used only
01	R1		as a source
02	R2	11	R11
03	R3	12	Unassigned
04	R4	13	Unassigned
05	R5	14	Unassigned
06	R6	15	Unassigned
07	IVL Register-Left Bank I/O address	16	Unassigned
	register. Used only as a destination.	17	IVR Register-Right Bank I/O address register. Used only as a destination.

a. Register Specification

208-278 is used to specify the least significant bit of a variable length field within the I/O Port previously selected by the IVL register. The length of the field is determined by R/L.



308-378 is used to specify the least significant bit of a variable length field within the I/O Port previously selected by the IVR Register. The length of the field is determined by R/L.

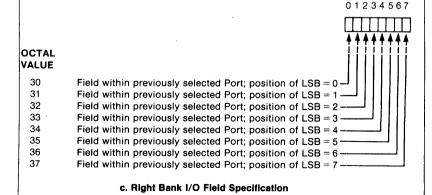


Table 5 S AND D FIELD OCTAL ASSIGNMENTS

R/L Field (3-Bit Field)

The R/L field performs one of two functions, specifying either a field length (L) or a right rotation (R). The function it specifies for a given instruction depends upon the contents of the S and D fields:

- A. When both S and D specify registers, the R/L field is used to specify a right rotation of the data specified by the S field. (Rotation occurs on the bus and not in the source register.) The register source data is right rotated within one instruction cycle time independent of the number of bit positions specified in the R/L field.
- B. When either or both the S and D fields specify a variable length I/O data field, the R/L field is used to specify the length of that data field.
- C. R/L field assignments are shown in Table 6.

R/L FIELD OCTAL VALUE	SPECIFICATION			
0	Field Length = 8 Bits			
1	Field Length = 1 Bit			
2	Field Length = 2 Bits			
3	Field Length = 3 Bits			
4	Field Length = 4 Bits			
5	Field Length = 5 Bits			
6	Field Length = 6 Bits			
7	Field Length = 7 Bits			

Table 6 R/L FIELD OCTAL ASSIGNMENTS

I Field (5/8-Bit Field)

The I field is used to load a literal value (contained in the instruction) into a register, or a variable I/O data field, or to modify the low order bits of the Program Counter.

The length of the I field is based on the S field in XEC, NZT, and XMIT instructions, as follows:

- A. When S specifies a register, the literal I is an 8-bit field (Type III format).
- B. When S specifies variable I/O data field, the literal I is a 5-bit field (Type IV format).

A Field (13-Bit Field)

The A field is a 13-bit Program Storage address field. This allows the 8X300 to directly address 8192 instructions.

REGISTER OPERATIONS

When a register is specified as the source and a variable I/O data field is specified as the destination, the low order bits of the results of the instructions MOVE, ADD, XOR are merged with the original destination data.

When an I/O data field of 1 to 8 bits is specified as the source, and a register as the destination, the 8-bit result of the operations MOVE, ADD, AND, XOR is stored in the register. The operations ADD, AND, XOR actually use the I/O data field (1 to 8 bits) with leading zeros to obtain 8-bit source data for use with the 8-bit AUX data during the operation.

IVL and IVR are write-only pseudo registers, and therefore can be specified as destination fields only. Operations involving IVL and IVR as sources are not possible. For example, it is not possible to increment IVR or IVL in a single instruction, and the contents of IVL or IVR cannot be transferred to a working register, or I/O Port.

The OVF (Overflow) Register can only be used as a source field; it is set or reset *only* by the ADD instruction.

ADDRESSING DATA ON THE INTERFACE VECTOR

I/O data fields are implemented via general purpose 8-bit I/O registers called Interface Vector (IV) Bytes. The IV registers serve to select IV bytes. In order for an instruction to access (read or write) an I/O data field, the address must be output to the IVL or IVR registers.

Thus, two instructions are required to operate on an Interface Vector byte.

XMIT ADDRESS, IVL MOVE LB, RB

Each of the two IV registers (IVL and IVR) may be set to select an IV byte, therefore two I/O ports may be active at one time—one on the Right Bank (IVR) and one on the Left Bank (IVL). Data may be input and output in one instruction following the selection of IV bytes:

XMIT ADDRESS1, IVL XMIT ADDRESS2, IVR ADD LB, RB

Once the IV byte is selected (addressed) it will remain selected until another address is output to the same IV register. Since an IV register (IVL, IVR) can be used only as a destination field of an instruction, any instruction sending data to IVL or IVR can be used to select an IV byte.

From the user's standpoint, however, all IV byte outputs can be read by an external device regardless of whether they are selected or not.

The address range of IVL and IVR is 0-255₁₀.

INSTRUCTION DESCRIPTIONS

The following instruction descriptions employ MCCAP (the 8X300 Cross Assembly Program) programming notation. This notation varies somewhat from the instruction descriptions provided in Tables 3 through 5. Thus, for example, explicit L field definition, as shown in Table 3 and Table 4 is not required by MCCAP instructions; MCCAP can create appropriate variable field addresses from information contained in Data Declaration statements which may be provided by the programmed at the beginning of his program.

The 8X300 instruction set is described below with examples shown in Figures 4 through 11.

MOVE S,D or MOVE S(R),D

Format: Type I, Type II Operation: (S)→(D) Description

Move data. The contents of S are transferred to D; the contents of S are unaffected. If both S and D are registers, R/L specifies a right rotate of the source data before the move. Otherwise, R/L specifies the length of the source and/or destination I/O data field. If the MOVE is between Left Bank and Right Bank I/O field, an 8-bit field must always be moved.

Example

Store the least significant 3 bits of register 5 (R5) in bits 4, 5 and 6 of the I/O Port previously addressed by the IVL register. See Figure 4.

ADD S,D or ADD S(R),D

Format: Type I, Type II

Operation: (S) + (AUX) → D

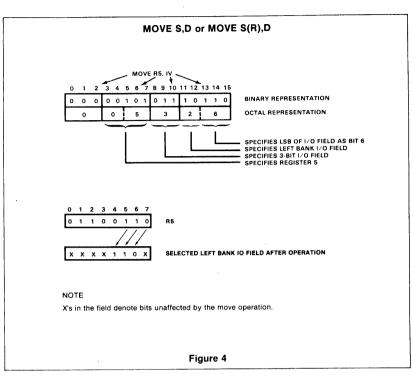
Carry → OVF

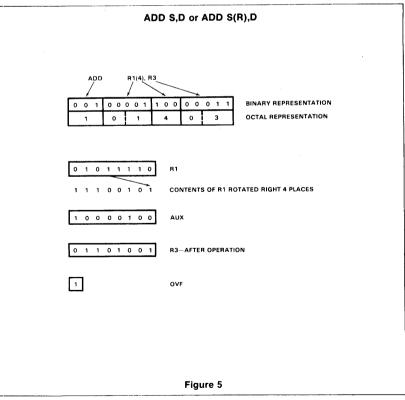
Description

Unsigned 2's complement 8-bit addition. The contents of S are added to the contents of the Auxiliary Register. The result is stored in D; OVF is set to the value of the carry. If both S and D are registers, R/L specifies a right rotate of the source (S) data before the operation. Otherwise, R/L specifies the length of the source and/or destination I/O data fields. S and AUX are unaffected unless specified as the destination.

Example

Add the contents of R1 rotated 4 places to AUX and store the result in R3. See Figure 5.





AND S,D or AND S(R),D

Format: Type I, Type II

Operation: (S) ∧ (AUX)→D

Description

Logical AND. The AND of the source field and the Auxiliary Register is stored into the destination. If both S and D are registers, R/L specifies a right rotate of the source (S) data before the AND operation. Otherwise R/L specifies the length of the source and/or destination I/O data fields. S and AUX are unaffected unless specified as a destination.

Example

Store the AND of the selected right bank I/O field and AUX in R4. The right bank data field is called WSBCD and is 4 bits long and located in bits 2, 3, 4 and 5. See Figure 6.

AND S,D or AND S(R),D ND WSBCD, R4 BINARY REPRESENTATION 0 1 0 1 0 0 0 0 0 0 OCTAL REPRESENTATION SPECIFIES REGISTER 4 SPECIFIES 4-BIT 1/O FIELD SPECIFIES LSB OF RB FIELD AS BIT 5 SPECIFIES RB FIELD 0 1 2 3 4 5 6 7 1 0 0 1 0 1 0 1 SELECTED RIGHT BANK I/O FIELD 0 0 0 0 0 SELECTED FIELD RIGHT JUSTIFIED WITH LEADING ZEROS ADDED. AND OPERATION 0 0 0 0 0 0 1 1 AUX 0 0 0 0 0 0 0 1 Figure 6

XOR S,D or XOR S(R),D

Format: Type I, Type II

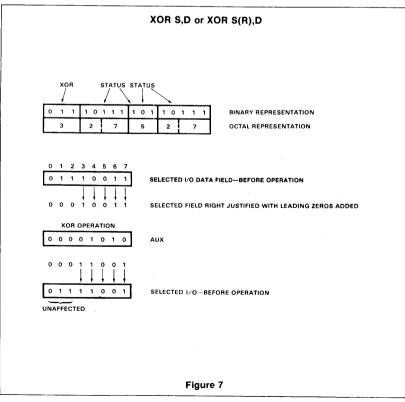
Operation: (S) ⊕ (AUX) → D

Description

Exclusive-OR. The Exclusive-OR of the source field and the Auxiliary Register is stored in the destination. If both S and D are registers, R/L specifies a right rotate of the source (S) data before the XOR operation. Otherwise R/L specifies the length of the source and/or destination I/O data fields. S and AUX are unaffected unless specified as a destination.

Example

Replace the selected I/O data field with the XOR of the field and AUX. The I/O data field is called STATUS and is 5 bits in length and located in bits 3, 4, 5, 6 and 7 of left bank. See Figure 7.



XEC I(S)

Format: Type III, Type IV

Operation

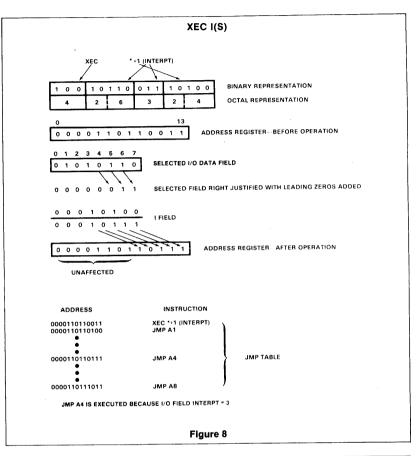
Execute instruction at the address specified by the Address Register with lower 5 or 8 bits replaced by (S) + I.

Description

Execute the instruction at the address determined by replacing the low order bits of the Address Register (AR) with the low order bits of the sum of the literal I and the contents of the source field. If S is a register, the low order 8 bits of AR are replaced; if S is an I/O data field, the low order 5 bits of AR are replaced, resulting in an execute range of 256 and 32 respectively. The Program Counter is not affected unless the instruction executed is a JMP or NZT (whose branch is taken).

Example

Execute one of n JMPs in a table of JMP instructions determined by the value of the selected I/O data field on the left bank. The table follows immediately after the XEC instruction and the I/O field is called INTERPT and is a 3-bit field located in bits 4, 5 and 6. See Figure 8.



XMIT I,D

Format: Type III, Type IV

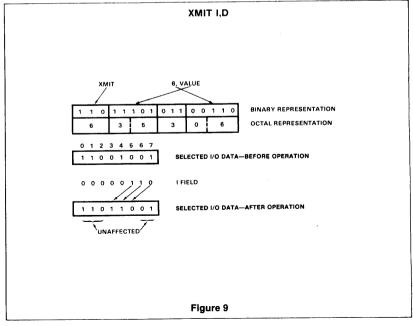
Operation: I→(D)

Description

Transmit literal. The literal field I is stored in D. If D is a register, an 8-bit field is transferred; if D is an I/O data field, up to a 5-bit field is transferred.

Example

Store the bit pattern 110 in the selected I/O data field on the right bank. The field name is VALUE and is located in bits 3, 4 and 5. See Figure 9.



NZT S,I

Format: Type III, Type IV

Operation:

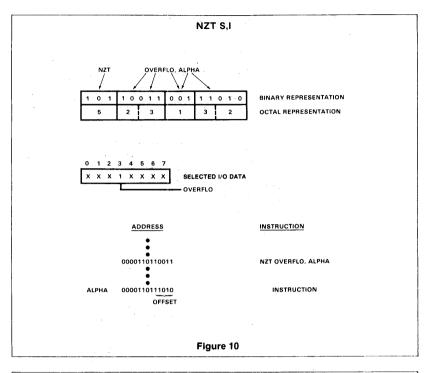
Non-Zero Transfer. If (S) \neq 0, PC offset by I \rightarrow PC; otherwise PC + 1 \rightarrow PC.

Description

If the data specified by the S field is non-zero, replace the low order bits of the Program Counter with I. Otherwise, processing continues with the next instruction in sequence. If S is a register, the low order 8 bits of the PC are replaced; if S is an I/O data field, the low order 5 bits of the PC are replaced, resulting in an NZT range of 256 and 32 respectively.

Example

Jump to Program Address ALPHA if the selected right bank I/O field is non-zero. The field name is OVERFLO and it is a 1-bit field located in bit 3. See Figure 10.



JMP A

Format: Type V

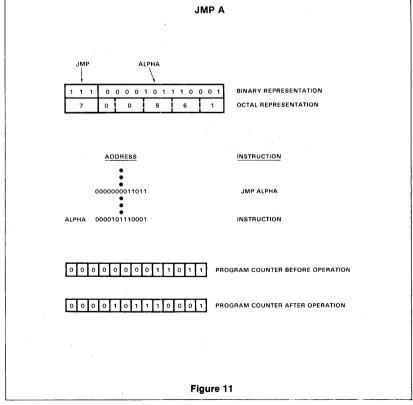
Operation: A → PC

Description

The literal value A is placed in the Program Counter and Address Register, and processing continues at location A. A has a range of 0-177778 (0-8191).

Example

Jump to location ALPHA (0000101110001). See Figure 11.



SYSTEM DESIGN USING THE 8X300 MICROCONTROLLER

Designing hardware around the 8X300 Interpreter reduces to selecting a program storage devicer (ROM, PROM, etc.), selecting I/O devices (IV byte, multiplexers, RAM, etc.), selecting clock mode (system driven or crystal controlled) and interfacing the Microcontroller to these components.

A specific example of a control system using the 8X300 Microcontroller is shown in Figure 12. Only 8 components—four 8T32 I/O Ports, one 82S208 RAM, two 82S215 ROMs, and an 8X300 are required to build this system which contains 512 words of program storage, 32 TTL I/O connection points, 256 bytes of working storage, and operates at a 250ns instruction cycle time.

Halt, Reset Signals

HALT:

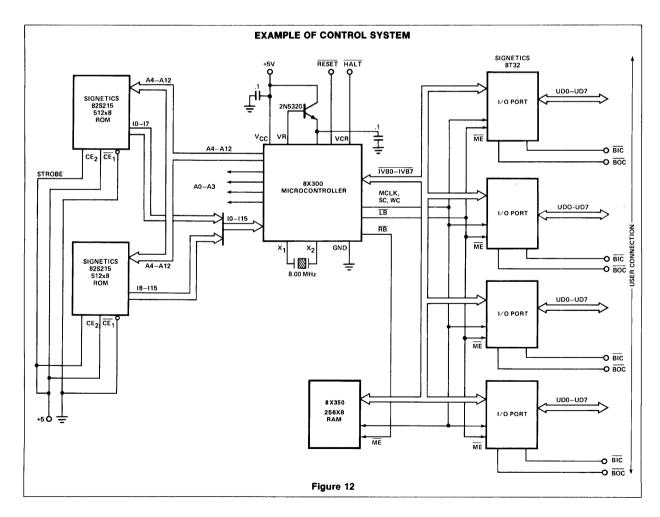
A low level at the HALT input stops internal operation of the Microcontroller at the start of the next instruction after HALT is applied (quarter cycle after MCLK). Since HALT is sampled at the start of each instruction cycle it is possible to prevent a cycle by applying HALT early in that cycle. HALT does not inhibit MCLK or affect any internal registers. Normal operation begins with the next complete cycle after the HALT input goes high.

RESET:

A low level at the RESET input sets the program counter and address register to zero. While RESET is low MCLK is inhibited. If RESET is applied during the last 2 quarter cycles, the MCLK during that cycle may be shortened. RESET should be applied for 1 full instruction cycle time to assure proper operation. When RESET input goes high an MCLK occurs prior to the resumption of normal processing. RESET does not affect the other internal registers.

SYSTEM TIMING

In systems with fast instruction cycle times, most Microcontroller delays are strictly determined by internal gate propagation delays.



Since some events are constrained to occur in certain quarter cycles, as system cycle times become slower, the delays will appear to increase due to gating with internal clocks. In the table of AC Electrical Characteristics, 2 columns are used: 1 to denote times which occur due to internal clock intervention and 1 to denote minimum delays for fast cycle times.

When using Signetics 8T32 I/O Ports, selection of instruction cycle time involves calculating the maximum program storage access time. Assuming the instruction is available when MCLK falls, the I/O control lines are stable 35ns later. Signetics 8T32's require another 35ns to disable a previously selected port and enable the desired port (assumes a change in bank signals). A 10ns margin has been added to the 8T32 enable for this evaluation to reflect the fact that most systems will have more capacitive loading than the 50pF test condition in the 8T32 specification and to allow for line delays.

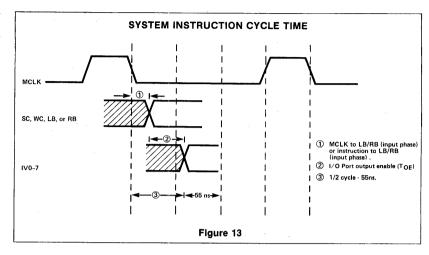
The system instruction cycle time for normal systems such as shown in Figure 12 is determined by Microcontroller propagation delays, program storage access time, and port output enable times. Instruction cycle time is normally constrained by one or more of the following conditions:

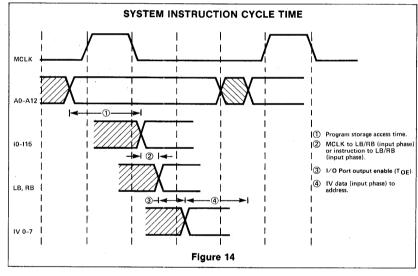
- Instruction to LB/RB (input phase) and I/O Port output enable:
 - TOE $\leq \frac{1}{2}$ cycle 55ns (Figure 13).
- Program storage access time and instruction to LB/RB (input phase) and I/O Port output enable and IV data (input phase) to address ≤ instruction cycle time (Figure 14).
- Program storage access time and instruction to address ≤ instruction cycle time (Figure 15).

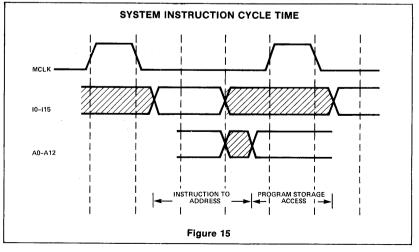
The first constraint can be used to determine the minimum cycle time. Using the inequality 35ns + 35ns $\leq \frac{1}{2}$ cycle - 55ns implies $\frac{1}{2}$ cycle \geq 125ns or an instruction time of 250ns.

Program storage access time for a 250ns instruction cycle can be calculated from the second constraint. Noting that the specification for IV data (input phase) to address is 115ns: Program storage access time + 35ns + 35ns + 115ns \leq 250ns implies program storage access time \leq 65ns.

The third constraint can be used to verify the maximum program storage access time. Noting that the specification for instruction to address is 185ns: Program storage access time + 185ns \leq 250ns confirms that program storage access time 65ns is satisfactory.







System Clock

The Microontroller has an integrated oscillator which generates all necessary clock signals. The oscillator is designed to connect directly to a series resonant quartz crystal via pins X1 and X2. The crystal resonant frequency, f, is related to the desired cycle time, T, by the relationship f = 2/T. For a 250ns system, f = 8.00MHz.

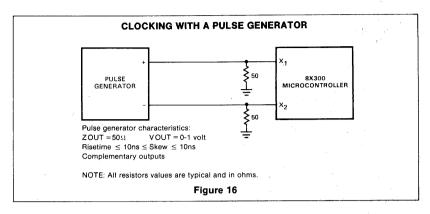
Туре:	Fundamental mode, series resonant
Impedance at Fundamental:	35 ohms maximum
Impedance at harmonics and spurs:	50 ohms minimum

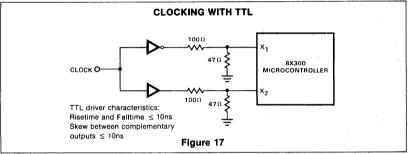
Table 7 CRYSTAL CHARACTERISTICS

In lower speed applications where the cycle time need not be precisely controlled, a capacitor may be connected between X1 and X2 to drive the oscillator. Approximate capacitor values are given in Table 8. If cycle time is to be varied, X1 and X2 should be driven from complementary outputs of a pulse generator. Figure 16 shows a typical configuration. For systems where the Interpreter is to be driven from a master clock the X1 and X2 lines may be interfaced to TTL logic as shown in Figure 17.

Cx,pF	CYCLE TIME
100	300ns
200	500ns
500	1.1 <i>µ</i> s
1000	2.0µs

Table 8 CLOCK CAPACITOR VALVES





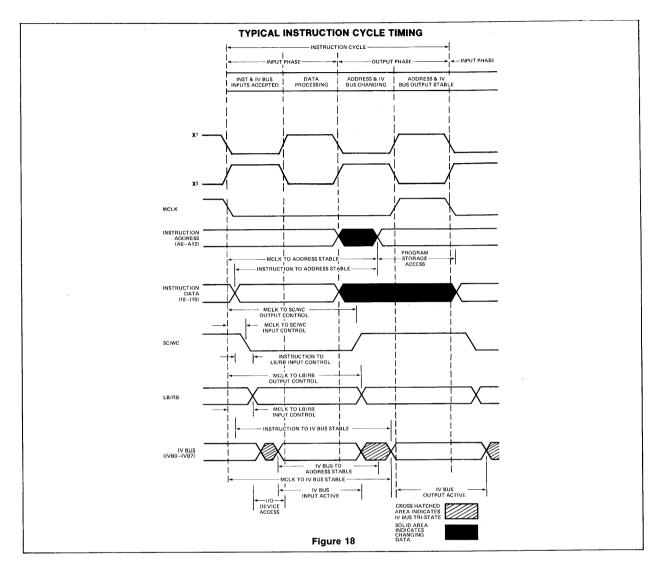
AC ELECTRICAL CHARACTERISTICS V $_{CC}$ = 5V ± 5% and 0°C \leqslant T $_{A}$ < 70°C

DELAY DESCRIPTION	PROPAGATION DELAY TIME	CYCLE TIME LIMIT
X1 falling edge to MCLK (driven from external		
pulse generator)	75ns	
MCLK to SC/WC falling edge (input phase)	25ns	
MCLK to SC/WC rising edge (output phase)		½ cycle + 25ns
MCLK to LB/RB (input phase)	35ns	
Instruction to LB/RB output (input phase)	35ns	
MCLK to LB/RB (output phase)		1/4 cycle + 35ns
MCLK to IV data (output phase)	185ns	½ cycle + 60ns
IV data (input phase) to IV data (output phase)	115ns	
Instruction to Address	185ns	½ cycle + 40ns
MCLK to Address	185ns	½ cycle + 40ns
IV data (input phase) to Address	115ns	
MCLK to IV data (input phase)		½ cycle - 55ns
MCLK to Halt falling edge to prevent		
current cycle		1/ ₄ cycle – 40ns
Reset rising edge to first MCLK		0 to 1 cycle

NOTE

- 1. Reference to MCLK is to the falling edge when loaded with 300pF.
- 2. Loading on Address lines is 150pF.





DC ELECTRICAL CHARACTERISTICS

	DADAMETED	TEAT AGUATIANA		LINUT		
	PARAMETER	TEST CONDITIONS	Min	Тур	Max	UNIT
v _{IH}	High level input voltage X1,X2 All others		.6 2			V V
V _{IL}	Low level input voltage X1,X2 All others				.4 .8	V V
V _{IC}	Input clamp voltage (Note 1)	V _{CC} = 4.75V I _I = -10mA			-1.5	V
ΊΗ	High level input current X1,X2	V _{CC} = 5.25V V _{IU} = .6V		2700		μΑ
	All others	V IH = .6V V CC = 5.25V V IH = 4.5V	•	<1	50	μΑ
IIL.	Low level input current X1,X2 IVBO-7	V _{CC} = 5.25V V _{IL} = .4V		-2500 -140	-200	μA μA
	IO-I15	V _{CC} = 5.25V V _{IL} = .4V V _{CC} = 5.25V V _{vv} = 4V		-880	-1600	μΑ
	HALT, RESET	V CC = 5.25V V IL = .4V V CC = 5.25V V IL = .4V		-230	-400	μΑ
VOL	Low level output voltage A0-A12	V _{CC} = 4.75V		.35	.55	v
	All others	I OL = 4.25mA V CC = 4.75V I OL = 16mA		.35	.55	v
v _{он}	High level output voltage	V _{CC} = 4.75V I _{OH} = 3mA	2.4			V
los	Short circuit output current (Note 2)	V _{CC} = 5.25V	-30		-140	mA
v _{CC}	Supply voltage		4.75	5	5.25	V
1 _{CC}	Supply current	V _{CC} = 5.25V			1.60	mA
REG	Regulator control	V _{CC} = 5.0V	-14		-21	mA
I _{CR}	Regulator current (Note 3)				290	mA
V _{CR}	Regulator voltage (Note 3)		2.2		3.2	V

ABSOLUTE MAXIMUM RATINGS

	PARAMETER	RATING	UNIT
Vcc	Supply voltage	7	ν
	Logic input voltage	5.5	V
	Crystal input	2	V
	voltage		

NOTES

^{1.} Crystal inputs X1 and X2 do not have clamp diodes.
2. Only one output may be grounded at a time.
3. (Limits apply for V_{CC} = 5V ± 5% and 0°C < T_A < 70°C unless specified otherwise.)

INTRODUCTION

Recent trends in system design are focusing on significantly lower costs and greater reliability while simultaneously maintaining or improving system preformance. Consequently, second generation as well as new system designs demand alternatives to conventional logic realizations which substantially reduce the system manufacturing costs. Signetics' LSI Logic Family provides this alternative by employing state-of-theart technology to produce Large-Scale-Integrated, system oriented building blocks. By replacing a relatively large number of conventional logic circuits with a few LSI Logic components, system costs are reduced in proportion to system printed circuit board area, total number of circuits and power requirements. Moreover, system reliability is increased in direct proportion to the decrease in the number of integrated circuits within the system.

SYSTEM LOGIC FAMILY

Signetics' LSI Logic offers the designer a dual technology logic family to achieve cost objectives. First, where speed is a primary objective, low power Schottky TTL technology is employed to yield devices which feature high speed, low power, and medium density. Therefore, significant reduction in the number of high speed logic components is possible without any sacrifice in performance. The distinct advantages of low power Schottky TTL give it one of the best technologies for high performance LSI design. The particular features of LS are:

- Comparable propagation delay to standard TTL 7-10ns/gate average
- Low power dissipation—2mW per gate typical at 50% duty cycle
- Low speed power product—19pj typical
- 45MHz typical maximum J-K flip-flop clock frequency
- High fan-out capability—22 unit load—LS input current requirement is 1/4 that of standard TTL (0.36mA/input). Outputs are capable of sinking 8mA.
- High logic density--70 gates/mm² of silicon
- Higher system reliability due to reduced power density.

Secondly, while lower power Schottky offers increased density at high speeds, substantial component reduction is achieved with devices employing Integrated Injection Logic (I²L). The very high density and low power of I²L makes it extremely attractive for use in realizing relatively large system building blocks. System level functions may each be accomplished at 10MHz speeds by a single I²L integrated circuit. I²L features are:

- Speed—comparable to TTL (10-20ns propagation delay)
- Power—1-2 orders lower than any technology power down mode
- Density—40% smaller than PMOS comparable to NMOS
- Interface—capability of mixing TTL, ECL circuits on the same chip.

The LSI Logic series of products are specially designed to aid system logic designers in the design of high performance, cost effective systems with a minimal number of parts. This is achieved through a line of standard products which are designed with the following objectives:

- Large Scale Building Blocks—devices are partitioned by function with high-level complexity and sophistication. These one chip building blocks are equivalent to 400 to 1500 elementary logic gates.
- Highest Performance Possible—devices in each category are optimized in performance according to their requirements. Devices requiring highest speed are designed using low power Schottky TTL and devices requiring medium speed are designed using I²L to minimize device power dissipation and maximize logic density.
- Off-the shelf items—system logic devices are designed to allow general purpose usage. Devices in this series can be used as stand-alone items to enhance performance on a certain system design or to utilize several components within the family to design a minimal cost system with minimum number of components, such as in microprocessor based systems.

SIGNETICS CAPABILITIES

Signetics is a leader in the development of bipolar LSI logic circuits using both LS and I²L technologies. Our capabilities are demonstrated by an 8-bit Fixed Instruction Microprocessor that Signetics is presently manufacturing with low power Schottky technology.

This circuit, the 8X300 Microcontroller, contains the equivalent of 700 logic gates on a 250X250 mil chip. This capability is being used to develop the LSI Logic Family and produce the substantial cost reduction offered by LSI without paying a penalty in performance.

Signetics is a leader in I²L technology. Signetics has been researching I²L as a standard product technology for over three years. Presently, Signetics possesses one of the fastest I²L processes in the industry. Recognizing the vast potential of I²L technology, Signetics is heavily committed to develop this high-speed line of LSI products using this technology. Signetics' I²L process is basically the same process as its low power Schottky process that utilizes a thin expitaxial layer for speed improvement. Since this process is a highly reliable standard process in Signetics, the confidence

level of producing such LSI circuits is extremely high. In addition to its own development activities, Signetics has access to N.V. Philips' vast development resources for continued development and enhancement of its present capabilities. This massive investment will result in continued improvement in speed/power performance of I2L products.

Table 1 shows expected trends in I²L speed/power curves during the next five years. Clearly I²L will play a major role in the innovation of faster and possibly more complex LSI Logic functions in the support of a continued exploitation of LSI circuitry in systems design.

ADVANTAGES

The advantages of using LSI in system design are multifold. First, note that the cost of an LSI chip is no more than the sum cost of the circuits it replaces, consequently there is a direct saving in manufacturing costs as a result of a net reduction in the number of parts. Moreover an LSI system offers economic advantages over an SSI or MSI approach besides lowering direct manufacturing costs. Some important considerations are:

- Power supply costs are reduced because logic cells internal to the device require less drive capability and consequently consume less power.
- Field repair costs are reduced because reliability is higher with an LSI implementation. This
 higher reliability is achieved because IC failure
 rates are largely proportional to number of
 devices rather than device complexity.
- Another factor to be considered is that the development cost of printed circuit boards will decrease because of smaller number of ICs and that to some extent, this reduction offsets the cost of layout on the LSI devices.
- Applications requiring medium speed performance are often forced to use lower density, high speed semiconductor devices due to lack of availability of medium speed LSI devices. I2L has changed this picture significantly. With the capability of controlling the amount of device injection current into an LSI circuit, the device can be controlled to operate at the desired speed and power for that particular application; thus, power consumption is minimized.

USING STANDARD LS PRO	CESS				
Average Propagation Delay	10-20ns at 200μA/Gate	Injection Current			
Density (Dual Layer Metal)	Shift Register etc.	320 gates/mm ²			
	Random Logic	150 gates/mm ²			
	Single Inverter	1.5 mil ²			
Chip Complexity	2000 gates at maximum speed (Random Logic)				
	4000 gates at maximum speed (Regular Arrays)				
Speed-Power	3pj @ maximum speed				
	0.35pj @ < 1MHz				

Table 1 SIGNETICS I2L CAPABILITY

OBJECTIVE SPECIFICATION

74S182/183-A.F.W

DESCRIPTION

The "182" is a high speed Look-Ahead Carry Generator ordinarily used with the "181" 4-Bit ALU or other arithmetic processing elements. This combination provides high speed Look-Ahead over word lengths of more than 4 bits.

The "182" accepts up to 4 pairs of active Low Carry Propagate and Carry Generate signals, an active High Carry input, and provides anticipated active High carries across 4 groups of binary adders.

Logic equations provided at the outputs are:

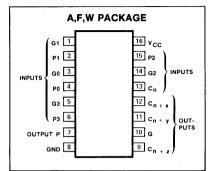
$$\begin{split} &C_{n+x} = G_0 + P_0 C_n \\ &C_{n+y} = G_1 + P_1 G_0 + P_1 P_0 C_n \\ &C_{n+z} = G_2 + P_2 G_1 + P_2 P_1 G_0 + P_2 P_1 P_0 C_n \\ &\overline{G} = \overline{G_3 + P_3 G_2 + P_3 P_2 G_1 + P_3 P_2 P_1 G_0} \end{split}$$

FEATURES

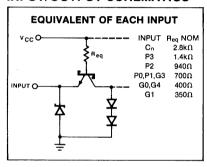
 $\overline{P} = \overline{P_3P_2P_1P_0}$

- Provides look-ahead carries across a group of 4 ALU's
- Multi-level look-ahead for high speed arithmetic operation over long word lengths

PIN CONFIGURATION



INPUT/OUTPUT SCHEMATICS



TYPICAL OF ALL OUTPUTS O VCC SO (INOM O OUTPUT

PIN DESIGNATION

PIN NOS.	DESIGNATION	FUNCTION
G0,G1,G2,G3	3,1,14,5	Active-Low Carry Generate Inputs
P0,P1,P2,P3	4,2,15,6	Active-Low Carry Propagate Inputs
Cn	13	Carry Input
C _{n+x} , C _{n+y} , C _{n+z}	12,11,9	Carry Outputs
G	10	Active-Low Carry Generate Output
Р	7	Active-Low Carry Propagate Output
Vcc	16	Supply Voltage
GND	8	Ground

OBJECTIVE SPECIFICATION

74S182/183-A,F,W

DC ELECTRICAL CHARACTERISTICS

DAGAMETER		PARAMETER TEST CONDITIONS 74S182			74S183			UNIT	
	PARAMETER	1231 CONDITIONS	Min Typ Max		Min Typ		Max	ONII	
VIL	Input Low voltage	V _{CC} = Min			0.8		0.8		٧
VIH	Input High voltage	$V_{CC} = Min$	2.0	-		2.0			V
Vcp	Input clamp diode voltage	$V_{CC} = Min, I_{IN} = -12mA$			-1.5		,	-1.2	V
Vol	Output Low voltage	V _{CC} = Min, I _{OL} = 16mA			0.4			0.5	٧
Vон	Output High voltage	$V_{CC} = Min, I_{OH} = -800 \mu A$	2.4			2.7			V
lıн	Input High current	$V_{CC} = Max, V_{IN} = 2.4V$							
	C _n input				80			50	μA
	P ₂ input				160			150	μA
	P ₃ input				120			100	μA
	P ₀ , P ₁ , G ₃ input		i		200			200	μ A
	G ₀ , G ₂ input	·			360	l		350	μΑ
	G₁ input				400	l		400	μΑ
l _{IL}	Input Low current	$V_{CC} = Max, V_{IN} = 0.4V$			1				
	C _n input				-3.2		1	-2.0	mA
	P ₂ input	·			-6.4	1		-6.0	mA
	P ₃ input	-			-4.8			-4.0	mA
	P ₉ , P ₁ , G ₃ input				-8.0	ŀ	ļ	-8.0	mA
	G ₀ , G ₂ input				-14.4			-14	mA
	G ₁ input				-16			-16	mA
los	Output short circuit current	V _{CC} = Max	-40		-100	-40		-100	mA
lcc	Power supply current	V _{CC} = Max, V _{IN} = 5V		ļ	72			109	mA

AC ELECTRICAL CHARACTERISTICS $T_A = 25^{\circ}C$, $V_{CC} = +5V$ unless otherwise specified.

PARAMETER				TEGT CONDITIONS		54/74		54/74 54/74S		S	— UNIT
		то	FROM	TEST CONDITIONS		Тур	Max	Min	Тур	Max	UNII
tplh tphl tplh tplh tphl	Propagation delay time Low to high High to low Low to high High to low Low to high	G0, G1, G2 G3, P0, P1 P2, P3 G0, G1, G2,	C _{n+x} , C _{n+y} C _{n+z} G	$C_L = 15pF$ $R_L = 400\Omega^1$		11 15	17 22		4.5 4.5 5	7 7 7.5	ns
tphl.	High to low Low to high	G3, P1, P2 P3 P0, P1, P2 P3	Р	$R_L = 280\Omega^2$					4.5	10.5 6.5	
tphl tplh tphl	High to low Low to high High to low	C _n	C _{n+x} , C _{n+y} C _{n+z}						6.5 6.5	10 10 10.5	

NOTES

Load circuit and typical waveforms are shown at the front of section.

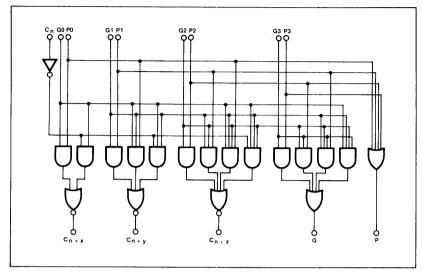
^{1.} $R_L = 400\Omega$ for 54/74.

^{2.} $R_L = 280\Omega$ for 54/74S.

OBJECTIVE SPECIFICATION

74S182/183-A,F,W

LOGIC DIAGRAM



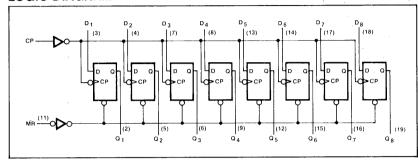
54LS273-N • 74LS273-N

DESCRIPTION

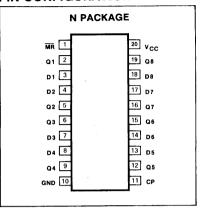
The 54/74LS273 is an Octal D Flip-Flop used primarily as an 8-bit positive edge-triggered storage register. D input informa-

tion is transferred to storage during the Low to High transition of the clock pulse. The Master Reset (MR) input simultaneously clears all flip-flops when low.

LOGIC DIAGRAM



PIN CONFIGURATION



OCTAL D FLIP-FLOP

54/74LS377

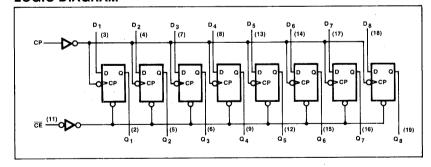
54LS273-N • 74LS273-N

DESCRIPTION

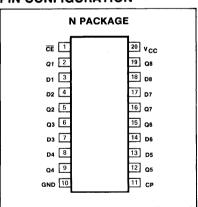
The 54/74LS377 is an Octal D Flip-Flop used primarily as an 8-bit positive edge-triggered storage register. D input information is transferred to storage during the Low

to High transition of the clock pulse. The edge-triggered Clock Enable (\overline{CE}) input enables the clock when Low and holds data when High.

LOGIC DIAGRAM



PIN CONFIGURATION



82S100-I,N • 82S101-I,N

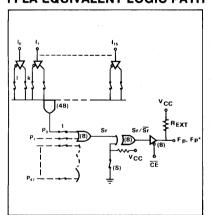
DESCRIPTION

The 82S100 (tri-state outputs) and the 82S101 (open collector outputs) are Bipolar Programmable Logic Arrays, containing 48 product terms (AND terms), and 8 sum terms (OR terms). Each OR term controls an output function which can be programmed either true active-high (Fp), or true activelow (F). The true state of each output function is activated by any logical combination of 16-input variables, or their complements, up to 48 terms. Both devices are field programmable, which means that custom patterns are immediately available by following the fusing procedure outlined in this data sheet.

The 82S100 and 82S101 are fully TTL compatible, and include chip-enable control for expansion of input variables, and output inhibit. They feature either open collector or tri-state outputs for ease of expansion of product terms and application in busorganized systems.

Both devices are available in commercial and military temperature ranges. For the commercial temperature range (0°C to +75°C) specify N82S100/101,I or N, and for the military temperature range (-55°C to +125°C) specify S82S100/101,I.

FPLA EQUIVALENT LOGIC PATH



LOGIC FUNCTION

Typical Product Term: $P_0 = I_0 \bullet I_1 \bullet \overline{I_2} \bullet I_5 \bullet \overline{I_{13}}$

Typical Output Functions:

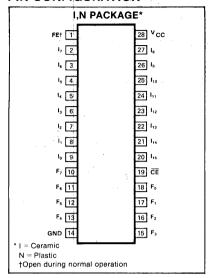
 $F_0 = (\overline{CE}) + (\underline{P_0} + \underline{P_1} + \underline{P_2}) @ S = Closed$ $F_0^* = (\overline{CE}) + (\overline{P_0} \cdot \overline{P_1} \cdot \overline{P_2}) @ S = Open$

For each of the 8 outputs, either the function Fp (active-high) or Fp (active low) is available, but not both. The required function polarity is programmed via link (S)

APPLICATIONS

- CRT display systems
- Random logic
- Code conversion
- Peripheral controllers
- **Function generators**
- Look-up and decision tables
- Microprogramming Address mapping
- Character generators
- Sequential controllers Data security encoders
- **Fault detectors**
- Frequency synthesizers

PIN CONFIGURATION



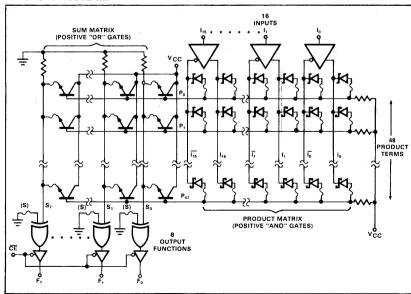
TRUTH RATINGS

MODE	Pn	CE	Sr ? f(Pn)	Fp	Fp
Disabled (82S101)	х		X	1	1
Disabled (82S100)	^	'		Hi-Z	Hi-Z
Read	1 0	0	Yes	1 0	0
Heau	Х	0	No	0	1

THERMAL TABLE

TEMPERATURE	MILI- TARY	COM- MER- CIAL
Maximum		
junction	175°C	150°C
Maximum		· ·
ambient	125°C	75°C
Allowable thermal		
rise ambient	* 4	
to junction	50° C	75°C

LOGIC DIAGRAM



82S100-I,N • 82S101-I,N

ABSOLUTE MAXIMUM RATINGS1

DADA	WETER	RA ⁻	TING	UNIT
PARA	WEIER	Min	Max	UNII
VCC VIN VOUT IIN IOUT	Supply voltage Input voltage Output voltage Input currents Output currents Temperature range Operating	-30	+7 +5.5 +5.5 +30 +100	Vdc Vdc Vdc mA mA °C
T _{STG}	N82S100/101 S82S100/101 Storage	0 55 65	+75 +125 +150	

DC ELECTRICAL CHARACTERISTICS N82S100/101: $0^{\circ} \le T_A \le +75^{\circ}$ C, $4.75V \le V_{CC} \le 5.25V$ S82S100/101: -55° C \leq T_A \leq +125 $^{\circ}$ C, 4.5V \leq V_{CC} \leq 5.5V

		TEGE COMPLETIONS	N82	N82S100/101			S82S100/101		
	PARAMETER	TEST CONDITIONS	Min	Тур	Max	Min	Тур	Max	UNIT
VIH VIL VIC	Input voltage ³ High Low Clamp ^{3,4}	$V_{CC} = Max$ $V_{CC} = Min$ $V_{CC} = Min, I_{IN} = -18mA$	2	-0.8	0.85 -1.2	2.	0.8	0.8 -12	V
V _{OH} V _{OL}	Output voltage High (82S100) ^{3,6} Low ^{3,6}	V _{CC} = Min I _{OL} = 9.6mA I _{OH} = -2mA	2.4	0.35	0.45	2.4	0.35	0.50	V
liH liL	Input current High Low	$V_{IN} = 5.5V$ $V_{IN} = 0.45V$		<1 -10	25 -100		<1 -10	50 -150	μΑ
lolk lo(OFF)	Output current Leakage? Hi-Z state (82S100)? Short circuit (82S100)4.8	V _{CC} = Max V _{OUT} = 5.5V V _{OUT} = 5.5V V _{OUT} = 0.45V V _{OUT} = 0V	-20	1 1 -1	40 40 -40 -70	-15	1 1 -1	60 60 -60 -85	μΑ μΑ mA
lcc	V _{CC} supply current ⁹	V _{CC} = Max		120	170		120	180	mA
Cin Cout	Capacitance ⁷ Input Output	$V_{CC} = 5.0V$ $V_{IN} = 2.0V$ $V_{OUT} = 2.0V$		8 17			8 17		pF

AC ELECTRICAL CHARACTERISTICS $R_1=470\Omega,\ R_2=1k\Omega,\ C_L=30pF$

N82S100/101: $0^{\circ}C \le T_{A} \le +75^{\circ}C$, $4.75V \le V_{CC} \le 5.25V$ S82S100/101: -55° C \leq T_A \leq $+125^{\circ}$ C, 4.5V \leq V_{CC} \leq 5.5V

PARAMETER				N	B2S100/	101	S8	UNIT		
		то	FROM	Min	n Typ ² Max Min Typ ²		Typ ²	Max	UNII	
TIA TCE	Access time Input Output Input Chip enable Output Chip enabl		Input Chip enable		35 15	50 30		35 15	80 40	ns
T _{CD}	Disable time Chip disable	Output	Chip enable		15	30		15	40	ns

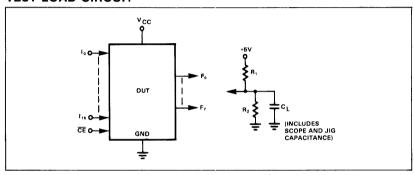
NOTES on following page.

82S100-I,N • 82S101-I,N

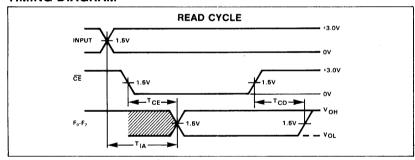
NOTES

- Stresses above those listed under "Absolute Maximum Ratings" may cause permanent damage to the
 device. This is a stress rating only, and functional operation of the device of these or any other
 condition above those indicated in the operation of the device specifications is not implied.
- 2. All typical values are at V_{CC} = 5V, T_A = 25°C.
- 3. All voltage values are with respect to network ground terminal.
- 4. Test one at the time.
- Measured with V_{IL} applied to CE and a logic high stored.
- Measured with a programmed logic condition for which the output test is at a low logic level. Output sink current is applied thru a resistor to Vcc.
- 7. Measured with: VIH applied to CE.
- 8. Duration of short circuit should not exceed 1 second.
- 9. Icc is measured with the chip enable input grounded, all other inputs at 4.5V and the outputs open.

TEST LOAD CIRCUIT



TIMING DIAGRAM



TIMING DEFINITIONS

TCE Delay between beginning of Chip Enable low (with Address valid) and when Data Output becomes valid.

Tco Delay between when Chip Enable becomes high and Data Output is in off state (Hi-Z or high).

TIA Delay between beginning of valid Input (with Chip Enable low) and when Data Output becomes valid.

VIRGIN DEVICE

The 82S100/101 are shipped in an unprogrammed state, characterized by:

- 1. All internal Ni-Cr links are intact.
- Each product term (P-term) contains both true and complement values of every input variable I_m (P-terms always logically "false").

- 3. The "OR" Matrix contains all 48-P-terms.
- 4. The polarity of each output is set to active high (Fp function).
- 5. All outputs are at a low logic level.

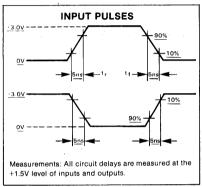
RECOMMENDED PROGRAMMING PROCEDURE

To program each of 8 Boolean logic functions of 16 true or complement variables, including up to 48 P-terms, follow the Program/Verify procedures for the "AND" matrix, "OR" matrix, and output polarity outlined below. To maximize recovery from programming errors, leave all links in unused device areas intact.

SET-UF

Terminate all device outputs with a 10K resistor to +5V. Set GND (pin 14) to 0V.

VOLTAGE WAVEFORM



Output Polarity

PROGRAM ACTIVE LOW (Fp FUNCTION)

Program output polarity before programing "AND" matrix and "OR" matrix. Program 1 output at the time. (S) links of unused outputs are not required to be fused.

- 1. Set FE (pin 1) to VFEL.
- 2. Set Vcc (pin 28) to VccL.
- 3. Set CE (pin 19), and Io through I15 to VIH.
- 4. Apply V_{OPH} to the appropriate output, and remove after a period t_p.
- 5. Repeat step 4 to program other outputs.

VERIFY OUTPUT POLARITY

- Set FE (pin 1) to V_{FEL}; set V_{CC} (pin 28) to V_{CCS}.
- 2. Enable the chip by setting \overline{CE} (pin 19) to V_{II} .
- 3. Address a non-existent P-term by applying V_{IH} to all inputs I₀ through I₁₅.
- 4. Verify output polarity by sensing the logic state of outputs F₀ through F₇. All outputs at a high logic level are programmed active low (F_p function), while all outputs at a low logic level are programmed active high (F_p function).
- 5. Return Vcc to Vccp or Vccl.

"AND" Matrix

PROGRAM INPUT VARIABLE

Program one input at the time and one Pterm at the time. All input variable links of unused P-terms are not required to be fused. However, unused input variables must be programmed as Don't Care for all programmed P-terms.

- 1. Set FE (pin 1) to V_{FEL}, and V_{CC} (pin 28) to V_{CCP}.
- 2. Disable all device outputs by setting CE (pin 19) to V_{IH}.
- 3. Disable all input variables by applying Vix to inputs I₀ through I₁₅.
- Address the P-term to be programmed (No. 0 through 47) by forcing the corresponding binary code on outputs F₀ through F₅ with F₀ as LSB. Use standard TTL logic levels V_{OHF} and V_{OLF}.
- 5a. If the P-term contains neither I_0 nor $\overline{I_0}$ (input is a Don't Care), fuse both I_0 and $\overline{I_0}$ links by executing both steps 5b and 5c, before continuing with step 7.
- 5 b. If the P-term contains I₀, set to fuse the I₀ link by lowering the input voltage at I₀ from V_{IX} to V_{IH}. Execute step 6.
- 5 c. If the P-term contains $\overline{l_0}$, set to fuse the l_0 link by lowering the input voltage at l_0 from V_{IX} to V_{IL} . Execute step 6.
- 6a. After to delay, raise FE (pin 1) from V_{FEL} to V_{FEH}.
- 6b. After t_D delay, pulse the $\overline{\text{CE}}$ input from V_{IH} to V_{IX} for a period t_D.
- 6 c. After to delay, return FE input to VFEL.
- Disable programmed input by returning I₀ to V_{IX}.
- Repeat steps 5 through 7 for all other input variables.
- Repeat steps 4 through 8 for all other Pterms.
- 10. Remove Vix from all input variables.

VERIFY INPUT VARIABLE

- 1. Set FE (pin 1) to V_{FEL}; set V_{CC} (pin 28) to V_{CCP}.
- 2. Enable F_7 output by setting \overline{CE} to V_{IX} .
- 3. Disable all input variables by applying V_{IX} to inputs I₀ through I₁₅.
- Address the P-term to be verified (No. 0 through 47) by forcing the corresponding binary code on outputs F₀ through F₅.

- 5. Interrogate input variable lo as follows:
 - A. Lower the input voltage at I_0 from V_{IX} to V_{IH} , and sense the logic state of output F_7 .
 - B. Lower the input voltage at I₀ from V_{IH} to V_{IL}, and sense the logic state output F₇.

The state of I₀ contained in the P-term is determined in accordance with the following truth table:

I ₀	F ₇	INPUT VARIABLE STATE CONTAINED IN P-TERM
0	1 0	Ī ₀
0	0	I ₀
0	1	Don't Care
0	0 0	(I ₀), (I 0)

Note that 2 tests are required to uniquely determine the state of the input variable contained in the P-term.

- Disable verified input by returning I₀ to Vix.
- 7. Repeat steps 5 and 6 for all other input variables.
- 8. Repeat steps 4 through 7 for all other P-terms.
- 9. Remove VIX from all input variables.

"OR" MATRIX PROGRAM PRODUCT TERM

Program one output at the time for one P-term at the time. All P_n links in the "OR" matrix corresponding to unused outputs and unused P-terms are not required to be fused.

- 1. Set FE (pin 1) to VFEL.
- 2. Disable the chip by setting CE (pin 19) to V_{IH}.
- After to delay, set V_{CC} (pin 28) to V_{CCS}, and inputs I₆ through I₁₅ to V_{IH}, V_{IL}, or V_{IX}.
- Address the P-term to be programmed (No. 0 through 47) by applying the corresponding binary code to input

- variables I₀ through I₅, with I₀ as LSB.
- 5a. If the P-term is contained in output function F₀ (F₀ = 1 or F₀ = 0), got to step 6, (fusing cycle not required).
- 5b. If the P-term is **not** contained in output function F_0 ($F_0 = 0$ or $F_0^* = 1$), set to fuse the P_n link by forcing output F_0 to V_{OPF} .
- 6a. After t_D delay, raise FE (pin 1) from V_{FEL} to V_{FEH}.
- 6b. After t_D delay, pulse the CE input from V_{IH} to V_{IX} for a period t_D.
- 6c. After to delay, return FE input to V_{FEL}.
- 6d. After t_D delay, remove V_{OPF} from output F₀.
- 7. Repeat steps 5 and 6 for all other output functions.
- 8. Repeat steps 4 through 7 for all other P-terms.
- 9. Remove Vccs from Vcc.

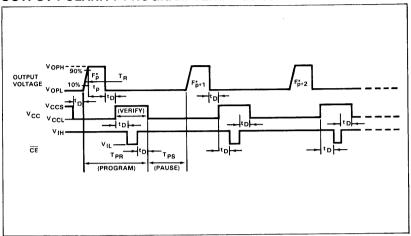
VERIFY PRODUCT TERM

- 1. Set FE (pin 1) to VFEL.
- Disable the chip by setting CE (pin 19) to VIH.
- 3. After t_D delay, set V_{CC} (pin 28) to V_{CCS}, and inputs I₀ through I₁₅ to V_{IH}, V_{IL}, or V_{IX}.
- Address the P-term to be verified (No. 0 through 47) by applying the corresponding binary code to input variables I₀ through I₅.
- 5. After to delay, enable the chip by setting CE (pin 19) to V_{IL}.
- 6. To determine the status of the P_n link in the "OR" matrix for each output function F_p or F_p*, sense the state of outputs F₀ through F₇. The status of the link is given by the following truth table:

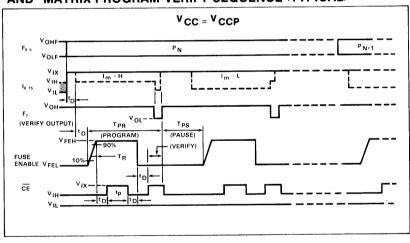
OUTP	UT	
Active High (Fp)	Active Low (Fp)	P-TERM LINK
0 1	1 0	Fused Present

- 7. Repeat steps 4 through 6 for all other P-terms
- 8. Remove Vccs from Vcc.

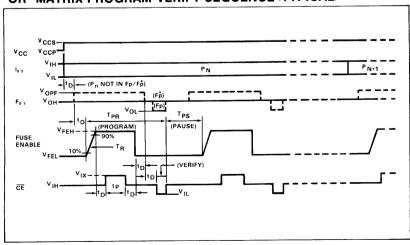
OUTPUT POLARITY PROGRAM-VERIFY SEQUENCE (TYPICAL)



"AND" MATRIX PROGRAM-VERIFY SEQUENCE (TYPICAL)



"OR" MATRIX PROGRAM-VERIFY SEQUENCE (TYPICAL)



PROGRAMMING SYSTEM SPECIFICATIONS 1 (TA = +25°C)

	PARAMETER	TEST CONDITIONS		LIMITS		
	T ANAMETER	TEST CONDITIONS	Min	Тур	Max	ואט
Vccs	V _{CC} supply (program/verify "OR", verify output polarity)2	Iccs = 550mA, min, Transient or steady state	8.5	8.75	9.0	V
V _{CCL} I _{CCS}	Vcc supply (program output polarity) Icc limit (program "OR")	$V_{CCS} = +8.75 \pm .25V$	0 550	0.4	0.8 1,000	V mA
Voph Vopl	Output voltage Program output polarity ³ Idle	I _{OPH} = 300 ± 25mA	16.0 0	17.0 0.4	18.0 0.8	V
Іорн	Output current limit (Program output polarity)	V _{OPH} = +17 ± 1V	275	300	325	m/
VIH VIL	Input voltage High Low		2.4	0.4	5.5 0.8	V
liH liL	Input current High Low	V _{IH} = +5.5V V _{IL} = 0V			50 -500	μА
V _{OHF}	Forced output voltage High Low		2.4	0.4	5.5 0.8	V
IOHF IOLF	Output current High Low	V _{OHF} =: +5.5V V _{OLF} = 0V			100 -1	μA mA
Vix	CE program enable level		9.5	10	10.5	V
l _{IX1}	Input variables current	$V_{IX} = +10V$			2.5	m/
I _{IX2}	CE input current	$V_{IX} = +10V$			5.0	m,
V _{FEH}	FE supply (program)3	I _{FEH} = 300 ± 25mA, Transient or steady state	16.0	17.0	18.0	V
V _{FEL}	FE supply (idle)	I _{FEL} = -1mA, max	1.25	1.5	1.75	V
FEH	FE supply current limit	$V_{FEH} = +17 \pm 1V$	275	300	325	m/
VCCP	Vcc supply (program/verify "AND")	I _{CCP} = 550mA, min, Transient or steady state	4.75	5.0	5.25	V
ICCP	Icc limit (program "AND")	$V_{CCP} = +5.0 \pm .25V$	550		1,000	m/
VOPF	Forced output (program)		9.5	10	10.5	v
IOPF	Output current (program)				10	m/
TR	Output pulse rise time		10		50	με
tp	CE programming pulse width		0.3	0.4	0.5	ms
t _D	Pulse sequence delay		10.			μs
TPR	Programming time			0.6		ms
T _{PR} + T _F	-				50	%
FL	Fusing attempts per link			j	2	сус
Vs	Verify threshold4		1.4	1.5	1.6	V

NOTES

These are specifications which a Programming System must satisy in order to be qualified by Signetics.

^{2.} Bypass Vcc to GND with a $0.01\mu f$ capacitor to reduce voltage spikes.

Care should be taken to ensure that the voltage is maintained during the entire fusing cycle. The recommended supply is a constant current source clamped at the specified voltage limit.

Vs is the sensing threshold of the FPLA output voltage for a programmed link, it normally constitutes
the reference voltage applied to a comparator circuit to verify a successful fusing attempt.

These are new limits resulting from device improvements, and which supersede, but do not obsolete the performance requirements of previously manufactured programming equipment.

16X48X8 FPLA PROGRAM TABLE

	1 1 1	Ī			102	170	7.0			F						NTF	RIES									
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ğ		Н	7	L	1	_	(da	ısh)				Α			•	(per	iod)			Н				L		
N.		NOT								NOTE		indo	onde	ent of	outo	ut pol	arity			TES		rograi	nme	d once	e onl	у.
Sic		Ente P-ter		for un	used	inpu	ts of	used		2. Er	iter (A	A) for	unus	ed ou	tputs	of us	ed P-	terms.				for all				
9										CT 1												IVE		EL*		
E			- ,	- T	-,¬	<u> </u>			JT \	/ARI	ABL	.E.*								_	ا ــــــــــــــــــــــــــــــــــــ		_ 1		i N*	_
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3 2 3	SIGN TOT PRO REV	46 47			+-	7	+	+			Že.									+				1		

[•] Input and Output fields of unused P-terms can be left blank. Unused inputs and outputs are FPLA terminals left floating.

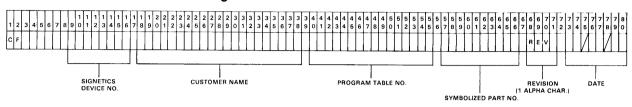
PUNCHED CARD CODING FORMAT

The FPLA Program Table can be supplied directly to Signetics in punched card form.

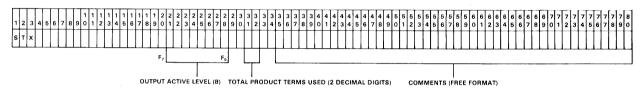
using standard 80-column IBM cards. For each FPLA Program Table, the customer should prepare in input card deck in accordance with the following format. Product Term cards 3 through 50 can be in any

order. Not all 48 Product Terms need to be present. Unused Product Terms require no entry cards, and will be skipped during the actual programming sequence:

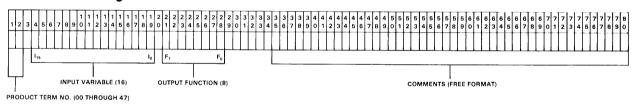
CARD NO.1—Free format within designated fields.



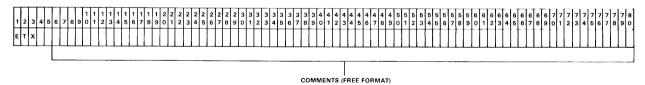
CARD NO. 2-



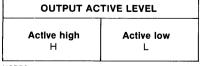
CARD NO. 3 through NO. 50



CARD NO. 51



Output Active Level entries are determined in accordance with the following table:



NOTES

- 1. Polarity programmed once only.
- 2. Enter (H) for all unused outputs.

Input Variable entries are determined in accordance with the following table:

INPUT VARIABLE												
lm	Im	Don't care										
H	L	— (dash)										

NOTE

Enter (--) for lunused inputs of used P-terms.

Output Function entries are determined in accordance with the following table:

OUTPUT F	OUTPUT FUNCTION											
Product term present in F _P A	Product term not present in Fp • (period)											

- Entries independent of output polarity.
- 2. Enter (A) for unused outputs of used P-terms.

BIPOLAR FIELD PROGRAMMABLE LOGIC ARRAY (16X48X8)

82S100-I,N • 82S101-I,N

TWX TAPE CODING FORMAT

The FPLA Program Table can be sent to Signetics in ASCII code format via airmail using any type of 8-level tape (paper, mylar, fanfold, etc.), or via TWX: just dial (910) 339-

9283, tell the operator to turn the paper puncher on, and acknowledge. At the end of transmission instruct the operator to send tape to Signetics Order Entry.

A number of Program Tables can be se-

quentially assembled on a continuous tape as follows, however limit tape length to a roll of 1.75 inch inside diameter, and 4.25 inch outside diameter:

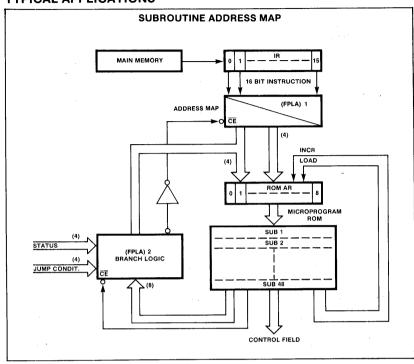
24" LEADE (C/R)	CTRL)R	MAIN HEADING	25 (C/R) MIN.	SUB HEADING (1)	25 RUBOUTS MIN.	 PROGRAM TABLE DATA (1) 	25 (C/R) MIN.	SUB HEADING (N)	25 RUBOUTS MIN.	PROGRAM TAI	3LE	TRAILERK (C/R)
	AIN HEA		he begir	ning of tap	e includes ti	ne following info	ormation,	with each e	ntry preced	ed by a (\$) ch	aracte	er,
1. Custo	omer Na	me				4.	Purchase (Order No				
2. Custo	omer TW	/X No				5.	Number of	Program Ta	ables			
3. Date						6.	Total Num	ber of Parts				
				ontain spec ether used		ation pertinent	to each P	rogram Ta	ble as follo	ws, with eac	n entr	У
1. Signe	etics Dev	ice No				4.	Date					
2. Prog	ram Tab	le No				5.	Customer	Symbolized	Part No			
3. Revis	ion					6.	Number o	Parts				
accord STAR (CO) S STX * A	TOF DAT NTROL A TART OF FIELD ACTIVE IDEN' ACTIVE I (8 DIG	A TEXT or B) DATA ELEVEL HIFIER EVEL DATA AITS, H/L) EF4F3F2F1F	STAR	Ormat: RT OF DATA FI RODUCT TERM 1 SPACE (MA PRODU (2 DE STA INI II 0 • I ₁₅ I ₁₄	IELD MIDENTIFIER MIDATORY) CT TERM NUM EGIMAL DIGIT: INT OF DATA F PUT VARIABLE (16 DIGITS, F I 13 1 12 1 1 1 10 19 1 n accordan	S) FIELD FIDENTIFIER LE DATA 4/(L/-) Fig. 176 514 312 110 * ce with the foll	START OF OUTPUT (8) F F ₇ F ₆ F ₅ owing Ta	DATA FIELD T FUNCTION I FUNCTION D DIGITS. A/•)	DENTIFIER ATA INPI AL P	EN JT AND OUTPUT L PRODUCT TER	ND OF D (CONT) DATA RMS US	FOR FO ETX
	INP	UT VARIA	BLE			OUTPUT FUN	CTION		ου	TPUT ACTIV	/E LE	VEL
I _m H	I _m L	- 1	on't car (dash	- 1	Produc present	t in Fp	oduct terr present in • (period	Fp	Active ⊢	-	Act	ive low L
NOTE					NOTES				NOTES			
	r unused	inputs of use	d P-terms.			lependent of output or unused outputs o				grammed once or all unused outp		

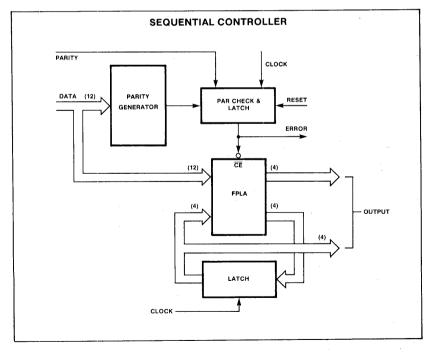
Although the Product Term data are shown entered in sequence, this is not necessary. It is possible to input only one Product Term, if desired. Unused Product Terms require no entry. ETX signalling end of Program Table may occur with less than the maximum number of Product Terms entered.

NOTES

- Corrections to any entry can be made by backspace and rubout. However, limit consecutive rubouts to less than 25.
- P-Terms can be re-entered any number of times. The last entry for a particular P-Term will be interpreted as valid data.
- Any P-Term can be deleted entirely by inserting the character (E) immediately following the P-Term number to be deleted, i.e., *P 25E deletes P-Term 25.
- To facilitate an orderly Teletype print out, carriage returns, line feeds, spaces, rubouts etc. may be interspersed between data groups, but only preceding an asterisk (*).
- Comments are allowed between data fields, provided that an asterisk (*) is not used in any Heading or Comment entry.

TYPICAL APPLICATIONS





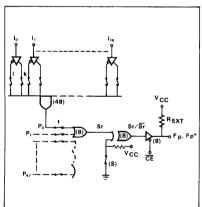
DESCRIPTION

The 82S200 (tri-state outputs) and the 82S201 (open collector outputs) are Bipolar Programmable Logic Arrays, containing 48 product terms (AND terms), and 8 sum terms (OR tems). Each OR term controls an output function which can be programmed either true active-high (Fp), or true active-low (Fp*). The true state of each output function is activated by any logical combination of 16input variables, or their complements, up to 48 terms. Both devices are mask programmable by supplying to Signetics Program Table data in one of the formats specified in this data sheet.

The 82S200 and 82S201 are fully TTL compatible, and include chip enable control for expansion of input variables, and output inhibit. They feature either open collector or tri-state outputs for ease of expansion of product terms and application in busorganized systems.

Both devices are available in commercial and military temperature ranges. For the commercial temperature range (0°C to +75°C) specify N82S200/201, I or N, and for the military temperature range (-55°C to +125°C) specify S82S200/201, I.

PLA EQUIVALENT LOGIC PATH



LOGIC FUNCTION

Typical Product Term: $P_0 = I_0 \cdot I_1 \cdot \overline{I_2} \cdot I_5 \cdot \overline{I_{13}}$

Typical Output Functions:

 $F_0 = (\overline{CE}) + (P_0 + P_1 + P_2) @ S = Closed$ $F_0^* = (\overline{CE}) + (\overline{P_0} \cdot \overline{P_1} \cdot \overline{P_2})$ @ S = Open

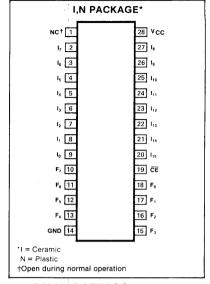
For each of the 8 outputs, either the function Fp (active-high) or Fp (active low) is available, but not both. The required function polarity is programmed via link (S).

APPLICATIONS*

- **CRT** display systems
- Random logic
- Code conversion
- Peripheral controllers
- **Function generators**
- Look-up and decision tables
- Microprogramming
- Address mapping
- Character generators
- Sequential controllers
- Data security encoders
- Fault detectors
- Frequency synthesizers

*For diagrams of Typical Applications reference 82S100 (T.S.)/82S101 (O.C.) Data Sheet.

PIN CONFIGURATION



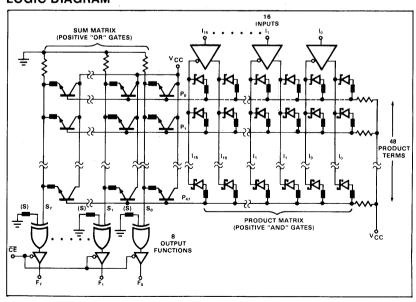
TRUTH TABLE

MODE	Pn	CE	Sr 🖁 f(Pn)	Fp	FĎ
Disabled (82S201)	х	1	×	1	1
Disabled (82S200)	^			Hi-Z	Hi-Z
	1	0	Yes	1	0
Read	0	0		0	1
licau	Х	0	No	0	1

THERMAL RATINGS

TEMPERATURE	MILI- TARY	COM- MER- CIAL
Maximum junction	175° C	150° C
Maximum ambient Allowable thermal	125° C	75°C
rise ambient to junction	50° C	75° C

LOGIC DIAGRAM



ABSOLUTE MAXIMUM RATINGS 1

	PARAMETER	RA*	ring	
	FANAMETEN			UNIT
Vcc	Supply voltage		+7	Vdc
ViN	Input voltage		+5.5	Vdc
Vout	Output voltage		+5.5	Vdc
lin	Input currents	-30	+30	mA
lout	Output currents		+100	mA
	Temperature range	*		°C
TA	Operating			·
	N82S200/201	0	+75	
	S82S200/201	-55	+125	
T _{STG}	Storage	-65	+150	

DC ELECTRICAL CHARACTERISTICS N82S200/201: $0^{\circ} \le T_{A} \le +75^{\circ}C$, $4.75V \le V_{CC} \le 5.25V$

S82S200/201: -55° C \leq T_A \leq +125 $^{\circ}$ C, 4.5 V \leq V_{CC} \leq 5.5 V

	DADAMETEO		N8	2S200/	201	S8	201		
	PARAMETER	TEST CONDITIONS	Min	Typ ²	Max	Min	Typ ²	Max	UNIT
VIH VIL VIC	Input voltage ³ High Low Clamp ^{3,4}	V _{CC} = Max V _{CC} = Min V _{CC} = Min, I _{IN} 7 -18mA	2	-0.8	0.85 -1.2	2	-0.8	0.8 -1.2	V
Vон Vol	Output voltage High (82S200)3,5 Low3,6	V _{CC} = Min I _{OH} = -2mA I _{OL} = 9.6mA	2.4	0.35	0.45	2.4	0.35	0.50	V
lı n lıL	Input current High Low	V _{IN} = 5.5V V _{IN} = 0.45V		<1 -10	25 -100		<1 -10	50 -150	μА
Iolk Io(off) Ios	Output current Leakage ⁷ Hi-Z state (82S200) ⁷ Short circuit (82S200)4.8	V _{CC} = Max V _{OUT} = 5.5V V _{OUT} = 5.5V V _{OUT} = 0.45V V _{OUT} = 0V	-20	1 1 -1	40 40 -40 -70	-15	1 1 -1	60 60 -60 -85	μA μA mA
loc	V _{CC} supply current ⁹	V _{CC} Max		120	170		120	180	mA
C _{IN} Cout	Capacitance ⁷ Input Output	$V_{CC} = 5.0V$ $V_{IN} = 2.0V$ $V_{OUT} = 2.0V$		8 17			8 17		pF

AC ELECTRICAL CHARACTERISTICS $R_1 = 470\Omega$, $R_2 = 1k\Omega$, $C_L = 30pF$

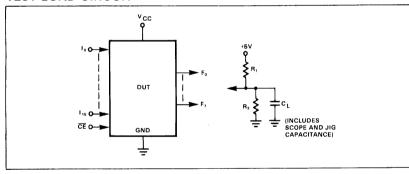
N82S200/201: 0° C \leq T_A \leq +75°C, 4.75V \leq V_{CC} \leq 5.25V S82S200/201: -55° C \leq T_A \leq +125°C, 4.5V \leq V_{CC} \leq 5.5V

				N8	325200/2	01	S8	UNIT		
P	ARAMETER	то	FROM	Min	Typ ²	Max	Min	Typ ²	Max	UNII
T _{IA} T _{CE}	Access time Input Chip enable	Output Output	Input Chip enable		35 15	50 30		35 15	80 50	ns
T _{CD}	Disable time Chip disable	Output	Chip enable		15	30		15	50	ns

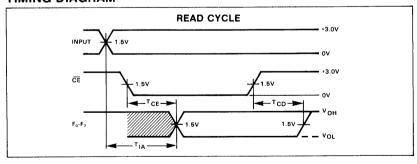
NOTES

- Stresses above those listed under "Absolute Maximum Ratings" may cause permanent damage to the
 device. This is a stress rating only, and functional operation of the device of these or any other
 condition above those indicated in the operation of the device specifications is not implied.
- All typical values are at V_{CC} = 5V, T_A = 25°C.
- 3. All voltage values are with respect to network ground terminal
- 4. Test one at the time.
- Measured with V_{IL} applied to CE and a logic high stored.
- Measured with a programmed logic condition for which the output test is at a low logic level. Output sink current is applied thru a resistor to Vcc.
- 7. Measured with: ViH applied to CE.
- 8. Duration of short circuit should not exceed 1 second.
- I_{CC} is measured with the chip enable input grounded, all other inputs at 4.5V and the outputs open.

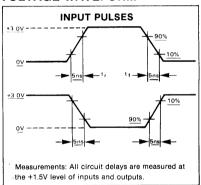
TEST LOAD CIRCUIT



TIMING DIAGRAM



VOLTAGE WAVEFORM



TIMING DEFINITIONS

Tce Delay between beginning of Chip Enable low (with Address valid) and when Data Output becomes valid.

Tcb Delay between when Chip Enable becomes high and Data Output is in off state (Hi-Z or high).

TIA Delay between beginning of valid Input (with Chip Enable low) and when Data Output becomes valid.

16X48X8 PLA PROGRAM TABLE

	1	<u> </u>	T				OV.	101	0 F	LA	rn	_				3LE 3LE		RIE	S								
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^{*}Input and Output fields of unused P-terms can be left blank. Unused inputs and outputs are PLA terminals left floating.

BIPOLAR MASK PROGRAMMABLE LOGIC ARRAY (16X48X8)

82S200-I,N • 82S201-I,N

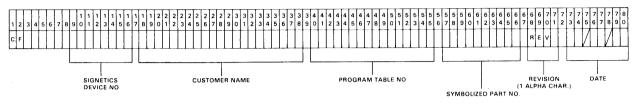
PUNCHED CARD CODING FORMAT

The PLA Program Table can be supplied directly to Signetics in punched card form,

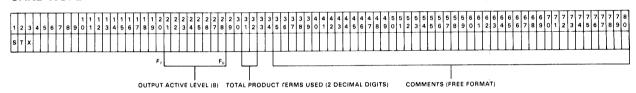
using standard 80-column IBM cards. For each PLA Program Table, the customer should prepare an input card deck in accordance with the following format. Product Term cards 3 through 50 can be in any

order. Not all 48 Product Terms need to be present. Unused Product Terms require no entry cards, and will be skipped during the actual programming sequence:

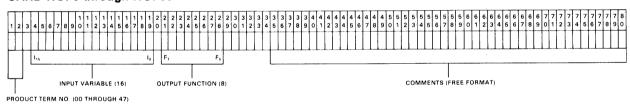
CARD NO.1—Free format within designated fields.



CARD NO. 2-



CARD NO. 3 through NO. 50



CARD NO. 51



Output Active Level entries are determined in accordance with the following table:

NOTES

Polarity programmed once only
 Enter (H) for all unused outputs

Input Variable entries are determined in accordance with the following table:

INPUT VARIABLE									
lm	Īm	Don't care							
H	L	— (dash)							

NOTE

Enter (-) for unused inputs of used P-terms.

Output Function entries are determined in accordance with the following table:

OUTPUT FUNCTION								
Product term present in Fp A	Product term not present in Fp • (period)							

NOTES

- 1. Entries independent of output polarity.
- 2. Enter (A) for unused outputs of used P-terms

TWX TAPE CODING FORMAT

1 ~ 1

The PLA Program Table can be sent to Signetics in ASCII code format via airmail using any type of 8-level tape (paper, mylar, fanfold, etc.), or via TWX: just dial (910) 339-

9283, tell the operator to turn the paper puncher on, and acknowledge. At the end of transmission instruct the operator to send tape to Signetics Order Entry.

A number of Program Tables can be se-

quentially assembled on a continuous tape as follows, however limit tape length to a roll of 1.75 inch inside diameter, and 4.25 inch outside diameter:

LEADER E MAIN 25 SUB 25 PROGRAM C/R) HEADING RUBOUTS PROGRAM DATA	
A. The MAIN HEADING at the beginning of tape includes the followin whether used or not:	
1. Customer Name	4. Purchase Order No.
2. Customer TWX No.	5. Number of Program Tables
3. Date	6. Total Number of Parts
B. Each SUB HEADING should contain specific information pertir preceded by a (\$) character, whether used or not: Signetics Device No	4. Date
2. Program Table No	5. Customer Symbolized Part No
3. Revision	6. Number of Parts
C. Program Table data blocks are initiated with an STX character, an consists of Output Active Level, Product Term, and Output Fundaccordance with the following format:	
START OF DATA TEXT (CONTROL A or B) START OF DATA FIELD ACTIVE LEVEL START OF DATA FIELD PRODUCT TERM IDENTIFIER 1 SPACE (MANDATORY) PRODUCT TERM NUMBER (2 DECIMAL DIGITS)	START OF DATA FIELD END OF DATA TEXT (CONTROL C)

Entries for the 3 Data Fields are determined in accordance with the following Table:

INPUT VARIABLE									
I _m	Ī _m	Don't care							
H	L	— (dash)							

OUTPUT FUNCTION							
Product term present in Fp A	Product term not present in Fp • (period)						

OUTPUT ACTIVE LEVEL

Active high Active low
H

INPUT AND OUTPUT DATA FOR

ALL PRODUCT TERMS USED

P F₀ ETX

NOTE

STX * A

Enter (-) for unused inputs of used P-terms.

IDENTIFIER

ACTIVE LEVEL DATA

(8 DIGITS, H/L)

NOTES

START OF DATA FIELD
INPUT VARIABLE IDENTIFIER

INPUT VARIABLE DATA

F₇F₆F₅F₄F₃F₂F₁F₀* P 00 * | I₁₅I₁₄I₁₃I₁₂I₁₁I₁₀I₉I₈I₇I₆I₅I₄I₃I₂I₁I₀* F F₇F₆F₅F₄F₃F₂F₁F₀* P

- 1. Entries independent of output polarity.
- 2. Enter (A) for unused outputs of used P-terms.

NOTES

OUTPUT FUNCTION IDENTIFIER

OUTPUT FUNCTION DATA

(8 DIGITS, A/•)

1. Polarity programmed once only.

01*

2. Enter (H) for all unused outputs.

Although the Product Term data are shown entered in sequence, this is not necessary. It is possible to input only one Product Term, if desired. Unused Product Terms require no entry. ETX signalling end of Program Table may occur with less than the maximum number of Product Terms entered.

NOTES

- Corrections to any entry can be made by backspace and rubout. However, limit consecutive rubouts to less than 25.
- P-Terms can be re-entered any number of times. The last entry for a particular P-Term will be interpreted as valid data.
- Any P-Term can be deleted entirely by inserting the character (E) immediately following the P-Term number to be deleted, i.e., *P 25E deletes P-Term 25.
- To facilitate an orderly Teletype print out, carriage returns, line feeds, spaces, rubouts etc. may be interspersed between data groups, but only preceding an asterisk (*).
- Comments are allowed between data fields, provided that an asterisk (*) is notused in any Heading or Comment entry.

DESCRIPTION

The 82S102 and 82S103 are Bipolar programmable AND/NAND gate arrays, containing 9 gates sharing 16 common inputs. On-chip input buffers enable the user to individually program for each gate either the True (I_m), Complement ($\overline{I_m}$), or Don't Care (X) logic state of each input. In addition, the polarity of each gate output is individually programmable to implement either AND or NAND logic functions.

Alternately, if desired, OR/NOR logic functions can also be realized by programming for each gate the complement of its input variables, and output (DeMorgan theorem).

Both devices are field-programmable, which means that custom patterns are immediately available by following the fusing procedure outlined in this data sheet.

The 82S102 and 82S103 include chipenable control for output strobing and inhibit. They feature either open collector or tri-state outputs for ease of expansion of input variables and application in busorganized systems.

Both devices are available in the commercial and military temperature ranges. For the commercial range (0°C to +75°C) specify N82S102/103, I or N, and for the military range (-55°C to +125°C) specify S82S102/103, I.

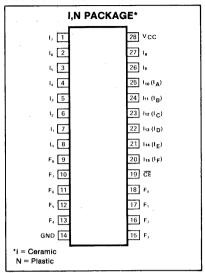
FEATURES

- Field programmable (Ni-Cr link)
- 16 input variables
- 9 output functions
- Chip enable input
- I/O propagation delay: N82S102/103: 30ns max S82S102/103: 50ns max
- Power dissipation: 600mW typ
- Input loading:
 - N82S102/103: -100µA max S82S102/103: -150µA max
- Output options:
 82S102: Open collector
 82S103: Tri-state
- Output disable function: 82S102: Hi 82S103: Hi-Z
- Fully TTL compatible

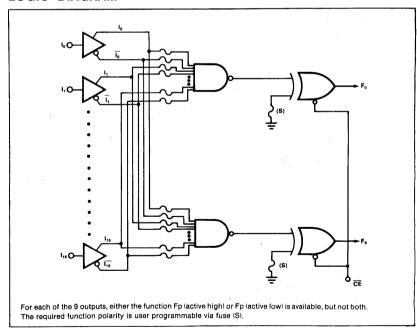
APPLICATIONS

- Random logic
- Address decoders
- Code detectors
- Peripheral selectors
- Fault monitors
- Machine state decoders

PIN CONFIGURATION

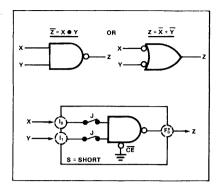


LOGIC DIAGRAM

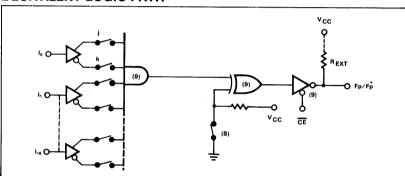


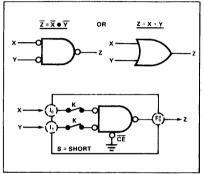
ABSOLUTE MAXIMUM RATINGS

	PARAMETER	RATING	UNIT
Vcc	Supply voltage	+7	Vdc
VIN	Input voltage	+5.5	Vdc
	Output voltage		Vdc
Voн	High (82S102)	+5.5	
Vo	Off-state (82S103)	+5.5	
lin	Input current	±30	l mA
lout	Output current	+100	mA
TA	Temperature range Operating		°C
	N82S102/103	0 to +75	
	S82S102/103	-55 to +125	
T _{STG}	Storage	-65 to +150	



EQUIVALENT LOGIC PATH

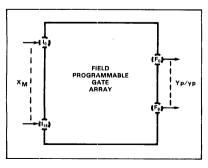




The Field Programmable Gate Array consists of 9 gates with individually programmable inputs and outputs.

The inputs to each gate can be programmed either True (I_m), Complement ($\overline{I_m}$), or Don't Care via corresponding links (j) and (k). The outputs of each gate can be programmed active-high (F_p) or active-low (F_p*) via corresponding links (S). Thus, each gate provides either of 2 output logic functions in terms of external input logic variables X_m as defined below (positive logic):

$$\begin{array}{l} \text{At S} = \text{Open:} \\ \text{Fp} = \overline{CE} + (X_0 \bullet X_1 \bullet X_2 \bullet \dots X_m) = \text{Yp} \\ \text{At S} = \overline{Closed:} \\ \text{Fp} = \overline{CE} + (\overline{X}_0 + \overline{X}_1 + \overline{X}_2 + \dots X_m) = \text{yp} \\ m = 0, 1, 2, \dots ... 15 \end{array}$$

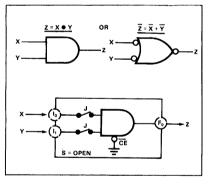


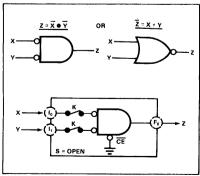
 $p=0,\,1,\,2,\,\ldots\,8$ and where $X_m=I_m,\,\overline{I_m}$, Don't Care, as assigned by programming polarity of inputs $I_{n-1}I_n$

When $\overline{CE}=$ low, all gates are enabled, and $F_p^*=\overline{F}_p$ giving $y_p=\overline{Y}_p$.

PROGRAMMABLE LOGIC FUNCTIONS

All internal links of virgin FPGAs are intact. Therefore, as shown in the Equivalent Logic Path, all symbolic switches are initially closed. Selective programming (opening) of links (J), (K), and (S) enables the user to assign input and output polarities to each gate for implementing NAND, NOR, AND, OR logic functions without changing the routing of input and output wires. This is shown in the following diagrams for a typical gate in terms of 2 input variables, which can be readily extended up to 16.





82S102-I.N • 82S103-I.N

DC ELECTRICAL CHARACTERISTICS N82S102/103: 0° C \leq T_A \leq +75 $^{\circ}$ C, 4.75V \leq V_{CC} \leq 5.25V S82S102/103: -55° C $\leq T_{A} \leq +125^{\circ}$ C, 4.5V $\leq V_{CC} \leq 5.5$ V

			N8	2S102/1	03	S82S102/103			UNIT
	PARAMETER1	TEST CONDITIONS	Min	Typ ²	Max	Min	Typ ²	Max	ONIT
V _I L V _I H V _I C	Input voltage Low¹ High¹ Clamp¹.3	$V_{CC} = Min$ $V_{CC} = Max$ $V_{CC} = Min, I_{IN} = -18mA$	2.0	-0.8	0.85	2.0	-0.8	0.8 -1.2	V
Vol Voh	Output voltage Low1,4 High (82S103)1,5	$V_{CC} = Min$ $I_{OL} = 9.6mA$ $I_{OH} = -2mA$	2.4	0.35	0.45	2.4	0.35	0.50	V
lıL lıн	Input current Low High	$V_{IN} = 0.45V$ $V_{IN} = 5.5V$		-10 <1	-100 25		-10 <1	-150 50	μΑ
lolk lo(OFF)	Output current Leakage (82S102)6 Hi-Z state (82S103)6 Short circuit (82S103)3,7	$V_{CC}=Max$ $V_{OUT}=5.5V$ $V_{OUT}=5.5V$ $V_{OUT}=0.45V$ $V_{OUT}=0V$	-20	1 1 -1	40 40 -40 -70	-15	1 1 -1	60 60 -60 -85	μΑ μΑ mA
Icc	V _{CC} supply current ⁸	V _{CC} = Max		120	170		120	180	mA
Cin Cout	Capacitance Input Output ⁶	$V_{CC} = 5.0V$ $V_{IN} = 2.0V$ $V_{OUT} = 2.0V$		8 15			8 15		pF

AC ELECTRICAL CHARACTERISTICS $R_1 = 470\Omega$, $R_2 = 1k\Omega$, $C_L = 30pF$

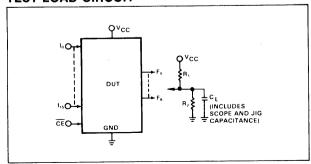
N82S102/103: 0° C \leq T_A \leq +75 $^{\circ}$ C, 4.75 V \leq V_{CC} \leq 5.25 V S82S102/103: -55° C \leq T_A \leq +125° C, 4.5V \leq V_{CC} \leq 5.5V

				N8	2\$102/1	03	S8	UNIT		
P.	ARAMETER	то	FROM	Min	Typ ²	Max	Min	Typ ²	Max	UNII
T _{IA} T _{CE}	Access time Input Chip enable	Output Output	Input Chip enable		20 15	30 30		20 15	50 40	ns
T _{CD}	Disable time Chip disable	Output	Chip enable		15	30	:	15	40	ns

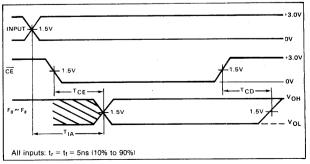
NOTES

- 1. All voltage values are with respect to network ground terminal.
- All typical values are at Vcc = 5V, TA = 25°C.
- Test each output one at a time.
- Measured with a programmed logic condition for which the output under test is at a low logic level.
- Output sink current is supplied through a resistor to Vcc.
- Measured with VIL applied to CE and a logic high at the output.
- Measured with VIH applied to CE.
- Duration of short circuit should not exceed 1 second.
- 8. Icc is measured with the chip enable input grounded, all other inputs at 4.5V and the outputs open.

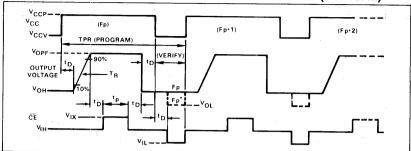
TEST LOAD CIRCUIT



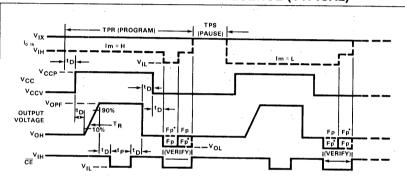
VOLTAGE WAVEFORM



OUTPUT POLARITY PROGRAM-VERIFY SEQUENCE (TYPICAL)



INPUT MATRIX PROGRAM-VERIFY SEQUENCE (TYPICAL)



VIRGIN DEVICE

The 82S102/103 are shipped in an unprogrammed state, characterized by:

- 1. All internal Ni-Cr links are intact.
- Each gate contains both true and complement values of every input variable I_m (logic Null state).
- 3. The polarity of each output is set to active low (F_D* function).
- 4. All outputs are at a high logic level.

RECOMMENDED PROGRAMMING PROCEDURE

To program each of 9 Boolean logic functions of 16 True, Complement, or Don't Care input variables follow the program/verify procedures for the Input Matrix and Output Polarity outlined below. To maximize recovery from programming errors, leave all links of unused gates intact.

SET-UP

Terminate all device outputs with a $10K\Omega$ resistor to +5V.

Output Polarity

PROGRAM ACTIVE HIGH (Fp FUNCTION)

Program output polarity before programming inputs (for convenience). Program one output at a time. (S) links of unused outputs are not required to be fused.

- 1 . Set GND (pin 14) to 0V, and V_{CC} (pin 28) to V_{CCV}.
- 2. Disable all device outputs by setting $\overline{\text{CE}}$ (pin 19) to V_{IH}.
- 3 . Disable all input variables by applying V_{IX} to inputs I_0 through $I_{15}.$
- A.Raise Vcc (pin 28) from Vccv to Vccp.
- B After t_D delay, force output to be programmed to V_{OPF}.
- C.After t_D delay, pulse the CE input from V_{IH} to V_{IX} for a period t_D.
- D.After to delay, remove VopF voltage source from output being programmed.
- E . After t_D delay, return V_{CC} (pin 28) to V_{CCV} , and verify.
- F. Repeat steps A through E for any other output.

VERIFY OUTPUT POLARITY

- 1. Set GND (pin 14) to 0V, and V_{CC} (pin 28) to V_{CCV}.
- 2. Disable all input variables by applying V_{IX} to inputs I₀ through I₁₅.
- A.After to delay, set the CE input to VII.
- B. Verify output polarity by sensing the logic state of outputs F₀ through F₈. All outputs at a low logic level are programmed active low (F_p function), while all outputs at a high logic level are programmed active high (F_p function).

Input Matrix PROGRAM INPUT VARIABLE

Program one input at a time for one gate at a time. Input variable links of unused gates are not required to be fused. However, unused input variables must be programmed at Don't Care for all used gates.

- 1. Set GND (pin 14) to 0V, and V_{CC} (pin 28) to V_{CCV}.
- 2. Disable all device outputs by setting \overline{CE} (pin 19) to V_{IH} .
- 3. Disable all input variables by applying V_{IX} to inputs I₀ through I₁₅.
- A-1.If a gate contains nether I_0 nor $\overline{I_0}$ (input is a Don't Care), fuse both j and k links by executing both steps A-2 and A-3, before continuing with step C.
- A-2.If a gate contains I_0 , set to fuse the k link by lowering the input voltage at I_0 from V_{IX} to V_{IH} . Execute step B.
- A 3.lf a gate contains $\overline{l_0}$, set to fuse the j link by lowering the input voltage at l_0 from V_{IX} to V_{IL} . Execute step B.
- B-1.After t_D delay, raise V_{CC} from V_{CCV} to V_{CCP} .
- B-2.After t_D delay, force output of gate to be programmed to V_{OPF}.
- B-3. After to delay, pulse the CE input from VIH to VIL for a period to.
- B-4.After to delay, remove VOPF voltage source from output of gate being programmed.
- B-5.After t_D delay, return V_{CC} (pin 28) to V_{CCV}, and verify.
- Disable programmed input by returning I₀ to V_{IX}.
- Repeat steps A through C for all other input variables.
- E. Repeat steps A through D for all other gates to be programmed.
- Remove V_{IX} from all input variables.

VERIFY INPUT VARIABLE

Unambiguous verification of the logic state programmed for the inputs of each gate requires prior knowledge of its programmed output polarity. Therefore, the output polarity verify procedure must precede input variable verify.

- 1 . Set GND (pin 14) to 0V, and V_{CC} (pin 28) to V_{CC}V.
- 2. Enable all outputs by setting \overline{CE} (pin 19) to V_{IL} .
- 3. Disable all input variables by applying V_{IX} to inputs I_0 through I_{15} .
- A .Interrogate input variable I_0 as follows: Lower the input voltage to I_0 from V_{IX} to V_{IL} , and sense the logic state of outputs F_{0-8} .

Raise the input voltage to I_0 from V_{IL} to V_{IH} and sense the logic state of outputs F_{0-8} .

BIPOLAR FIELD PROGRAMMABLE GATE ARRAY (16X9) 82S102 (0.C.)/82S103 (T.S.)

82S102-I,N • 82S103-I,N

The state of I₀ contained in each gate is determined in accordance with the given truth table. Note that 2 tests are required to uniquely determine the state of the input variable contained in each gate.

- B.Disable verified input by returning l_0 to V_{IX} .
- C.Repeat steps A and B for all other input variables.
- D.Remove VIX from all input variables.

TRUTH TABLE FOR INPUT VERIFICATION

I ₀	Fp	FĎ	INPUT VARIABLE STATE	LINK FUSED
0	1 0	0	Īo	j
0	0	1 0	l ₀	k
0	1	0	Don't care	Both
0	0	1	(1 ₀), (1 0)	Neither

PROGRAMMING SYSTEMS SPECIFICATIONS1 TA = 25°C

			•		UNIT		
	PARAMETER	TEST CONDITIONS	Min	Тур	Max	01411	
VCCP	Vcc supply Program ²	I _{CCP} = 350 ± 50mA, Transient or steady state	8.5	8.75	9.0	V	
Vccv	Verify	-	4.75	5.0	5.25		
Ісср	ICC limit (program)	$V_{CCP} = +8.75 \pm .25V$, Transient or steady state	400	450	500	mA	
VOPF	Forced output voltage3 (program)	$I_{OP} = 150 \pm 25$ mA, Transient or steady state	16.0	17.0	18.0	V	
IOPF	Output current limit (program)	$V_{OP} = +17 \pm 1V$, Transient or steady state	125	150	175	m.A	
Vih Vil	Input voltage High Low		2.4 0	0.4	5.5 0.8	V	
lin lic	Input current High Low	V _{IH} = +5.5V V _{IL} = 0V			50 -500	μΑ	
VIX IIX1 IIX2	CE program enable level Input variables current CE input current	$\begin{array}{c} V_{IX} = +10V \\ V_{IX} = +10V \end{array}$	9.5	10	10.5 5.0 10.0	M/ m/	
T _R t _P t _D T _{PR}	Output pulse rise time CE programming pulse width Pulse sequence delay Programming time Programming duty cycle		10 0.3 10	0.4	50 0.5	μs ms μs ms	
T _{PR} +T _P F _L V _S	Fusing attempts per link Verify threshold4		1.4	1.5	2 1.6	cyc V	

NOTES

These are specifications which a Programming System must satisfy in order to be qualified by Signetics.

^{2.} Bypass Vcc to GND with a $0.01 \mu F$ capacitor to reduce voltage spikes.

Care should be taken to ensure that the voltage is maintained during the entire fusing cycle. The recommended supply is a constant current source clamped at the specified voltage limit.

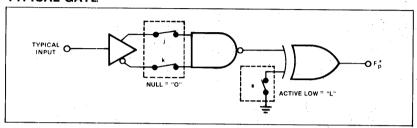
Vs is the sensing threshold of a gate output voltage for a programmed link. It normally constitutes the
reference voltage applied to a comparator circuit to verify a successful fusing attempt.

PROGRAMMING

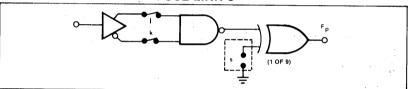
In a virgin device all Ni-Cr links are intact. The initial programmed state of each gate is shown in the Typical Gate illustration.

To program inputs and outputs of each gate for implementing the desired logic function, fuse Ni-Cr links as indicated in the fuse link diagrams.

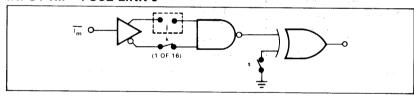
TYPICAL GATE



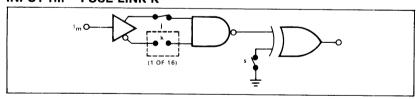
OUTPUT ACTIVE HIGH = FUSE LINK S



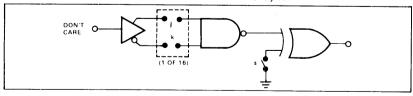
INPUT Im = FUSE LINK J



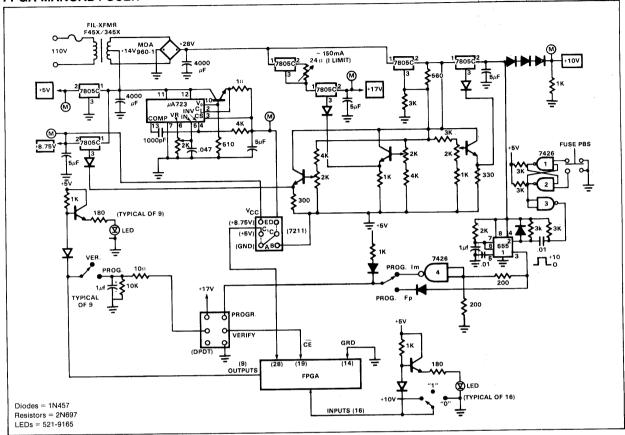
INPUT Im = FUSE LINK K



INPUT DON'T CARE = FUSE BOTH LINKS J, K



FPGA MANUAL FUSER



BIPOLAR FIELD PROGRAMMABLE GATE ARRAY (16X9) 82\$102 (0.C.)/82\$103 (T.S.)

82S102-I,N • 82S103-I,N

16X9 FPGA PROGRAM TABLE

CUSTOMER NAME PURCHASE ORDER # SIGNETICS DEVICE # TOTAL NUMBER OF PARTS PROGRAM TABLE #	CF (XXXX) CUSTOMER SYMBOLIZED PART # DATE RECEIVED
0 =	
1 =	
2 =	
3 =	
4 =	
5 =	
6 =	
7 =	
. -	

OUTF	PUT							IN	PUT VA	RIABL	E						
POLAF	RITY	I _o	l ₁	l ₂	l ₃	I ₄	15	16	I ₇	l ₈	l ₉	IA	lв	Ic	ΙD	ΙE	lF
F ₀	0	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	
F ₁	16	16	17	18	19	20	21	22	23	24	25	26	27	28	29	30	
F ₂	32	32	33	34	35	36	37	38	39	40	41	42	43	44	45	46	
F ₃	48	48	49	50	51	52	53	54	55	56	57	58	59	60	61	62	
F ₄	64	64	65	66	67	- 68	69	70	71	72	73	74	75	76	77	78	7
F ₅	80	80	81	82	83	84	85	86	87	88	89	90	91	92	93	94	
F ₆	96	96	97	98	99	100	101	102	103	104	105	106	107	108	109	110	1
F ₇	112	112	113	114	115	116	117	118	119	120	121	122	123	124	125	126	13
F ₈	128	128	129	130	131	132	133	134	135	136	137	138	139	140	141	142	14
Active-hig Active-low		$I_{\underline{m}} = H$ $I_{\overline{m}} = L$ Don't Care	=	•				L	J								

The number in each cell in the table denotes its address for programmers with a decimal address display.

8T31-N,F

DESCRIPTION

The I/O Port is an 8-bit bidirectional data register designed to function as an I/O interface element in microprocessor systems. It contains 8 clocked data latches accessible from either a microprocessor port or a user port. Separate I/O control is provided for each port. The 2 ports operate independently, except when both are attempting to input data into the I/O Port. In this case, the user port has priority.

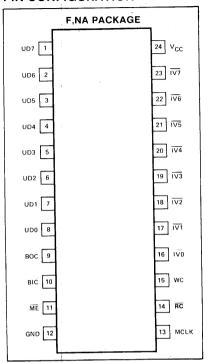
A master enable (ME) is provided that enables or disables the μP bus regardless of the state of the other inputs, but has no effect on the user bus.

A unique feature of this family is its ability to start up in a predetermined state. If the clock is maintained at a voltage less than .8V until the power supply reaches 3.5V, the user port will always be all logic 1 levels, while the microprocessor port will be all logic 0 levels.

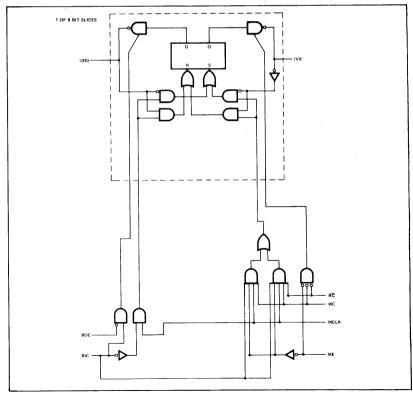
FEATURES

- Each device has 2 ports, one to the user, the other to a microprocessor. I/O Ports are completely bidirectional
- Ports are independent, with the user port having priority for data entry
- User data input synchronous
- The user data bus is available with tristate (8T32, 8T36) or open collector (8T33, 8T35) outputs
- At power up, the user port outputs are high
- Tri-state TTL outputs for high drive capability
- Directly compatible with the 8X300 Microcontroller
- Operates from a single 5V power supply over a temperature range of 0°C to +70°C

PIN CONFIGURATION



BLOCK DIAGRAM



8T31-N.F

PIN DESIGNATION

PIN	SYMBOL	NAME AND FUNCTION	TYPE
1-8	UD0-UD7:	User Data I/O Lines. Bidirectional data lines to communicate with user's equipment.	Acitve high three-state
16-23	IVO-IV7:	Microprocessor Bus. Bidirectional data lines to communicate with controlling digital system.	Active low three-state
10	BIC:	Input Control. User input to control writing into the I/O Port from the user data lines.	Active low
9	BOC:	Output Control. User input to control reading from the I/O Port onto the user data lines.	Active low
11	ME:	Master Enable. System input to enable or disable all other system inputs and outputs. It has no effect on user inputs and outputs.	Active low
15	WC:	Write Command. When WC is high, stores contents of IV0-IV7 as data.	Active high
14	RC	Read Command. When RC is low, data is presented on IV0-IV7.	Active low
13	MCLK:	Master Clock. Input to strobe data into the latches. See function tables for details.	Active high
24	V _{CC} :	5V power connection.	
12	GND:	Ground.	

BIC	вос	MCLK	USER DATA BUS FUNCTION
H	L	X	Output Data
L	X	H	Input Data
H	H	X	Inactive

H = High Level L = Low Level X = Don't care

Table 1 USER PORT CONTROL FUNCTION

ME	RC	wc	MCLK	BIC	MICROPROCESSOR BUS FUNCTION
L X X	L X H X	L H L H	X H X X	X H X L	Output Data Input Data Inactive Inactive Inactive

Table 2 MICROPROCESSOR PORT CONTROL FUNCTION

USER DATA BUS CONTROL

The activity of the user data bus is controlled by the BIC and BOC inputs as shown in Table 1.

The user data input is a synchronous function with MCLK. A low level on the BIC input allows data on the user data bus to be written into the data latches only if MCLK is at a high level. A low level on the BIC input allows data on the user data bus to be latched regardless of the level of the MCLK input.

To avoid conflicts at the data latches, input from the microprocessor port is inhibited when BIC is at a low level. Under all other conditions the 2 ports operate independently.

MICROPROCESSOR BUS CONTROL

As is shown in Table 2, the activity of the microprocessor port is controlled by the ME, RC, WC and BIC inputs, as well as the state of an internal status latch. BIC is included to show user port priority over the microprocessor port for data input.

BUS OPERATION

Data written into the I/O Port from one port will appear inverted when read from the other port. Data written into the I/O Port from one port will not be inverted when read from the same port.

8T31-N,F

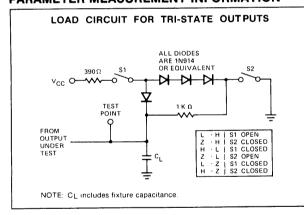
DC ELECTRICAL CHARACTERISTICS V_{CC} = 5V ± 5%, 0°C \leq TA \leq 70°C unless otherwise specified.

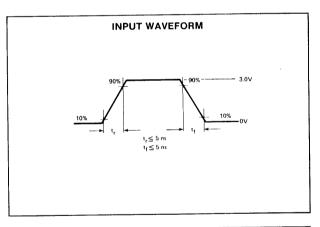
				LIMITS		UNIT	
	PARAMETER	TEST CONDITIONS	Min	Тур	Min	ONIT	
	Input voltage					V	
V _{IH}	High		2.0				
V IH	Low		1		.8		
V _{IL}	Clamp	L ₁ = -5mA			-1		
V IC	Output voltage	I _I = -5mA V _{CC} = 4.75V				V	
V	High	1	2.4				
Vон	Low				.55		
VOL	Input current ¹	V = = 5.25V				μΑ	
•	•	V _{CC} = 5.25V V _{IH} = 5.25V		<10	100	'	
ļн	High	V = 5V	1	-350	-550		
[†] 1∟	Low	V _{IL} = .5V				mA	
	Output current ²	V - 4.75V					
los	Short circuit	V _{CC} = 4.75V	10				
	UD bus						
	IV bus		20	100	150	mA	
Icc	VCC supply current	V _{CC} = 5.25V		100	150	IIIA	

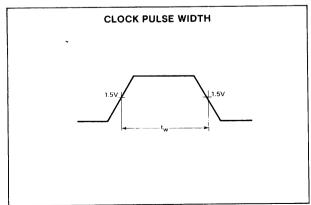
NOTES

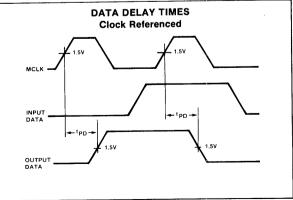
2. Only one output may be shorted at a time.

PARAMETER MEASUREMENT INFORMATION









The input current includes the tri-state/open collector leakage current of the output driver on the data lines.

8T31-N,F

AC ELECTRICAL CHARACTERISTICS $0^{\circ}C \le T_A \le 70^{\circ}C$, $V_{CC} = 5V \pm 5\%$

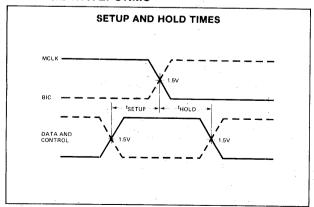
PARAMETER	1212	TEST		LIMITS		
PANAMETER	INPUT	CONDITION	Min	Тур	Max	UNIT
t _{PD} User data delay ¹	UDX MCLK	C _L = 50pF		25 45	38 61	ns ns
t _{OE} User output enable	вос	C _L = 50pF	18	26	47	ns
t _{OD} User output disable	BIC BOC	C _L = 50pF	18 16	28 23	35 33	ns ns
t _{PD} μP data delay ¹	IVBX MCLK	C _L = 50pF		38 48	53 61	ns ns
t _{OE} μP output enable	ME RC WC	C _L = 50pF	14	19	25	ns .
t _{OD} μP output disable	ME RC WC	C _L = 50pF	13	17	32	·ns
t _W Minimum pulse width	MCLK		40			ns
^t SETUP Minimum setup time ²	UDX ³ BIC IVX ME RC WC		15 25 55 30 30 30			ns
HOLD Minimum hold time ²	UDX ³ BIC IVX ME RC WC		25 10 10 5 5			ns

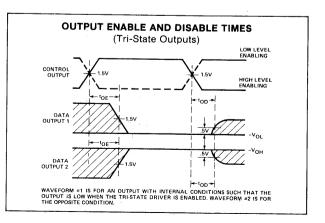
NOTES

- Data delays referenced to the clock are valid only if the input data is stable at the arrival
 of the clock and the hold time requirement is met.
- Set up and hold times given are for "normal" operation. BIC setup and hold times are for a user write operation. RC setup and hold times are for an I/O Port select operation. ME and WC setup and hold times are for a microprocessor bus write operation.

3. Times are referenced to MCLK

VOLTAGE WAVEFORMS





TYPES

- 8T32 Tri-State. Synchronous User Port
- 8T33 Open Collector, Synchronous User
- 8T35 Open Collector, Asynchronous User Port
- 8T36 Tri-State, Asynchronous User Port

DESCRIPTION

The Addressable I/O Port is an 8-bit bidirectional data register designed to function as an I/O interface element in microprocessor systems. It contains 8 data latches accessible from either a microprocessor port or a user port. Separate I/O control is provided for each port. The 2 ports operate independently, except when both are attempting to input data into the I/O Port. In this case, the user port has priority.

A unique feature of the I/O Port is the way in which it is addressed. Each device has an 8bit, field programmable address, which is used to enable the microprocessor port. When the SC control signal is high, data at the microprocessor port is treated as an address. If the address matches the I/O Port's internally programmed address, the microprocessor port is enabled, allowing data transfer through it.

The port remains enabled until an address which does not match is presented, at which time the port is disabled (data transfer is inhibited). A Master Enable input (ME) can serve as a ninth address bit, allowing 512 I/O Ports to be individually selected on a bus, without decoding. The user port is accessible at all times, independent of whether or not the microprocessor port is selected.

A unique feature of this family is their ability to start up in a predetermined state. If the clock is maintained at a voltage less than .8V until the power supply reaches 3.5V, the user port will always be all logic 1 levels, while the port will be all logic 0 levels.

ORDERING

The 8T32/33/35/36 may be ordered in preaddressed form. To order a preaddressed device use the following part number format:

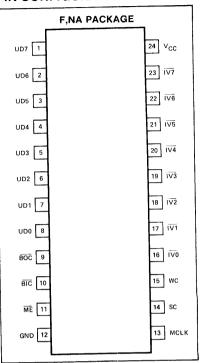
N8TYY-XXX P -P = F Ceramic package NA Plastic package XXX = Any address from 000 through 255 (decimal) -256 available addresses YY = I/O Port version (32, 33, 35, 36)

A stock of 8T32s and 8T36s with addresses 1 through 10 will be maintained. A small quantity of addresses 11 through 50 will also be available with a longer lead time.

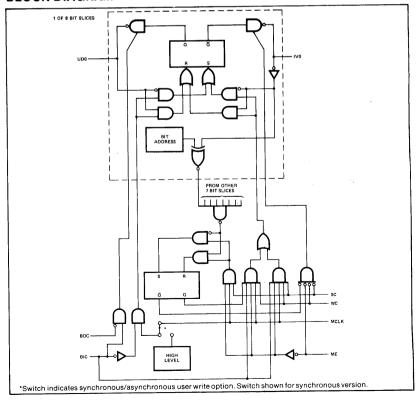
FEATURES

- A field-programmable address allows 1 of 512 I/O Ports on a bus to be selected, without decoders.
- Each device has 2 ports, one to the user, the other to a microprocessor.
- Completely bidirectional.
- Ports are independent, with the user port having priority for data entry.
- A selected I/O Port de-selects itself when another I/O Port address is sensed.
- User data input available as synchronous (8T32, 8T33) or as asynchronous (8T35, 8T36) function.
- The user data bus is available with tristate (8T32, 8T36) or open collector (8T33, 8T35) outputs.
- At power up, the I/O Port is not selected and the user port outputs are high.
- Tri-state TTL outputs for high drive capa-
- Directly compatible with the 8X300 Microcontroller.
- Operates from a single 5V power supply over a temperature range of 0°C to 70°C.

PIN CONFIGURATION



BLOCK DIAGRAM



PIN DESCRIPTION

PIN	SYMBOL	NAME AND FUNCTION	TYPE
1-8	UD0-UD7:	User Data I/O Lines. Bidirectional data lines to communicate with user's equipment. Either tristate or open collector outputs are available.	Active high
16-23	IVO-IV7:	Microprocessor Bus. Bidirectional data lines to communicate with controlling digital system (microprocessor).	Active low three-state
10	BIC:	Input Control. User input to control writing into the I/O Port from the user data lines.	Active low
9	BOC:	Output Control. User input to control reading from the I/O Port onto the user data lines.	Active low
11	ME:	Master Enable. System input to enable or dis- able all other system inputs and outputs. It has no effect on user inputs and outputs.	Active low
15	WC:	Write Command. When WC is high and SC is low, I/O Port, if selected, stores contents of IVO-IV7 as data.	Active high
14	SC:	Select Command. When SC is high and WC is low, data on IV0-IV7 is interpreted as an address. I/O Port selects itself if its address is identical to μP bus data; it de-selects itself otherwise.	Active high
13	MCLK:	Master Clock. Input to strobe data into the latches. See function tables for details.	Active high
24	VCC:	5V power connection.	
12	GND:	Ground.	

BIC	BIC BOC MCLK USER DATA BUS FUNCTION	US FUNCTION		
5.0		WCLK	8T32, 8T33	8T35, 8T36
H L L	Х	X H L X	Input Data Inactive	Input Data Input Data

H = High Level L = Low Level

Table 1 USER PORT CONTROL FUNCTION

ME	sc	wc	MCLK	BIC	STATUS LATCH	I/O PORT FUNCTION
L	L	L	×	Х	SET	Output Data
L	L	Н	Н	Н	SET	Input Data
L	н	L	i H	X	l x	Input Address
L	Н	Н	Н	L	X	Input Address
L	Н	Н	Н	Н	l x	Input Data and Address
L	X	Н	L	Х	l x	Inactive
L	Н	Х	L	X	X	Inactive
L	L	Н	н	L	l x	Inactive
L	L	Х	X	Х	Not Set	Inactive
Н	Х	Х	X	Х	X	Inactive

Table 2 MICROPROCESSOR PORT CONTROL FUNCTION

USER DATA BUS CONTROL

The activity of the user data bus is controlled by the BIC and BOC inputs as shown in Table 1.

For the 8T32 and 8T33, user data input is a synchronous function with MCLK. A low level on the BIC input allows data on the user data bus to be written into the data latches only if MCLK is at a high level. For the 8T35 and 8T36, user data input is an asynchronous function. A low level on the BIC input allows data on the user data bus to be latched regardless of the level of the MCLK input. Note that when 8T35 or 8T36 are used with the 8X300 Microcontroller care must be taken to insure that the Microprocessor bus is stable when it is being read by the 8X300 Microcontroller.

To avoid conflicts at the Data Latches, input from the Microprocessor Port is inhibited when BIC is at a low level. Under all other conditions the 2 ports operate independently.

MICROPROCESSOR BUS CONTROL

As is shown in Table 2, the activity of the microprocessor port is controlled by the ME, SC, WC and BIC inputs, as well as the state of an internal status latch. BIC is included to show user port priority over the microprocessor port for data input.

Each I/O Port's status latch stores the result of the most recent I/O Port select; it is set when the I/O Port's internal address matches the Microprocessor Bus. It is cleared when an address that differs from the internal address is presented on the Microprocessor Bus.

In normal operation, the state of the status latch acts like a master enable; the microprocessor port can transfer data only when the status latch is set.

When SC and WC are both high, data on the Microprocessor Bus is accepted as data, whether or not the I/O Port was selected. The data is also interpreted as an address. The I/O Port sets its select status if its address matches the data read when SC and WC were both high; it resets its select status otherwise.

BUS OPERATION

Data written into the I/O Port from one port will appear inverted when read from the other port. Data written into the I/O Port from one port will not be inverted when read from the same port.

⁼ Low Level X = Don't care

AC ELECTRICAL CHARACTERISTICS 0° C \leq T $_{A} \leq$ 70 $^{\circ}$ C, V $_{CC}$ = 5V \pm 5%

	·	TEST		LIMITS		
PARAMETER	INPUT	CONDITION	Min	Тур	Max	UNIT
t _{PD} User data delay (Note 1)	UDX MCLK* BIC†	C _L = 50pF		25 45 40	38 61 55	ns
^t OE User output enable	вос	C _L = 50pF	18	26	47	ns
t _{OD} User output disable	BIC BOC	C _L = 50pF	18 16	28 23	35 33	ns
t _{PD} μP data delay (Note 1)	IVBX MCLK	C _L = 50pF		38 48	53 61 .	ns
t _{OE} μP output enable	ME SC WC	C _L = 50pF	14	19	25	ns
t _{OD} μP output disable	ME SC WC	C _L = 50pF	13	17	32	ns
t _W Minimum pulse width	MCLK BIC†		40 35			ns
[†] SETUP Minimum setup time	BIC* IVX ME SC WC	(Note 2)	15 25 55 30 30 30			ns
^t HOLD Minimum hold time	UDX ^{II} BIC* IVX ME SC WC	(Note 2)	25 10 10 5 5 5			ns

Applies for 8T32 and 8T33 only.

NOTES:

[†] Applies for 8T35 and 8T36 only

Times are referenced to MCLK for 8T32 and 8T33, and are referenced to BIC for 8T35 and 8T36.

^{1.} Data delays referenced to the clock are valid only if the input data is stable at the arrival

of the clock and the hold time requirement is met.

Set up and hold times given are for "normal" operation. BIC setup and hold times are for a user write operaton. SC setup and hold times are for an I/O Port select operation. WC setup and hold times are for an Microprocessor Bus write operation. ME setup and hold times are for both IV write and select operations.

DC ELECTRICAL CHARACTERISTICS $0^{\circ}C \le T_A \le 70^{\circ}C$, $V_{CC} = 5V \pm 5\%$

	PARAMETER	TEST CONDITIONS		LIMITS		UNITS
		TEOT CONDITIONS	Min	Тур	Max	
VIH	High-level input voltage		2.0			V
VIL	Low-level input voltage				.8	V
V _{CL} .	Input clamp voltage	I _I = -5mA			-1	V
Iн .	High-level input current ¹	$V_{CC} = 5.25V$ $V_{IH} = 5.25V$		<10	100	μА
l _{IL}	Low level input current1	$V_{CC} = 5.25V$ $V_{IL} = .5V$		-350	-550	μΑ
VoL	Low-level output voltage	$V_{CC} = 4.75V$ $I_{OL} = 16mA$.55	٧
Vон	High-level output voltage	$V_{CC} = 4.75V$ $I_{OH} = -3.2$ mA	2.4			٧
los	Short-circuit output current ² UD bus IV bus	$V_{CC} = 4.75V$ $V_{CC} = 4.75V$	10 20			mA mA
Icc	Supply current	V _{CC} = 5.25V		100	150	mA

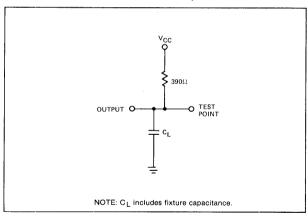
NOTES

- The input current includes the Tri-state/Open Collector leakage current of the output driver on the data lines.
- 2. Only one output may be shorted at a time.
- 3. These limits do not apply during address programming.

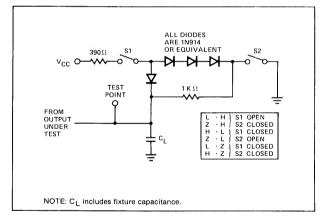
Absolute Maximum Ratings:

Input voltage³ 5.5V

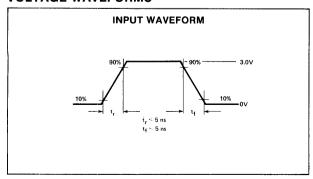
TEST LOAD CIRCUIT (OPEN COLLECTOR OUTPUTS)

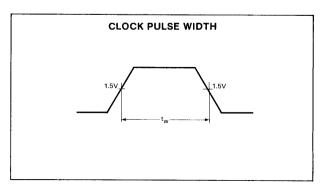


TEST LOAD CIRCUIT (TRI-STATE OUTPUTS)

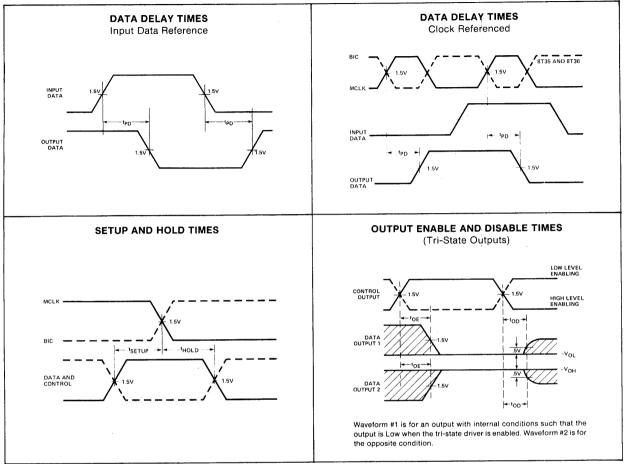


VOLTAGE WAVEFORMS





VOLTAGE WAVEFORMS (Cont'd)

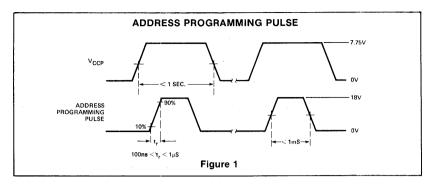


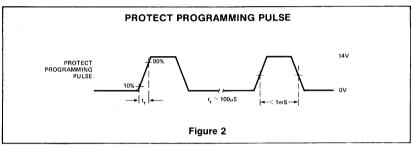
ADDRESS PROGRAMMING

The I/O Port is manufactured such that an address of all high levels (>2V) on the Microprocessor Bus inputs matches the Port's internal address. To program a bit so a low-level input (<0.8V) matches, the following procedure should be used:

- Set all control inputs to their inactive state (BIC = BOC = ME = V_{CC}, SC = WC = MCLK = GND). Leave all Microprocessor Bus I/O pins open.
- 2. Raise VCC to 7.75V ± .25V.
- 3. After V_{CC} has stabilized, apply a single programming pulse to the user data bus bit where a low-level match is desired. The voltage should be limited to 18V; the current should be limited to 75mA. Apply the pulse as shown in Figure 1.
- 4. Return V_{CC} to 0V. (Note 1).
- 5. Repeat this procedure for each bit where a low-level match is desired.
- 6. Verify that the proper address is programmed by setting the Port's status latch (IV0-IV7 = desired address, ME = WC = L, SC = MCLK = H). If the proper address has been programmed, data presented at the μP bus will appear inverted on the user bus outputs. (Use normal V_{CC} and input voltage for verification.)

After the desired address has been programmed, a second procedure must be followed to isolate the address circuitry. The procedure is:





- Set V_{CC} and all control inputs to 0V. (V_{CC} = BIC = BOC = ME = SC = WC = MCLK = 0V). Leave all Microprocessor Bus I/O pins open.
- Apply a protect programming pulse to every user data bus pin, one at a time. The voltage should be limited to 14V; the current should be limited to 150mA. Apply the pulse as shown in Figure 2.
- 3. Verify that the address circuitry is isolated by applying 7V to each user data bus pin and measuring less than 1mA of input current. The conditions should be the same as in step 1 above. The rise time on the verification voltage must be slower than 100µs.

PROGRAMMING SPECIFICATIONS¹

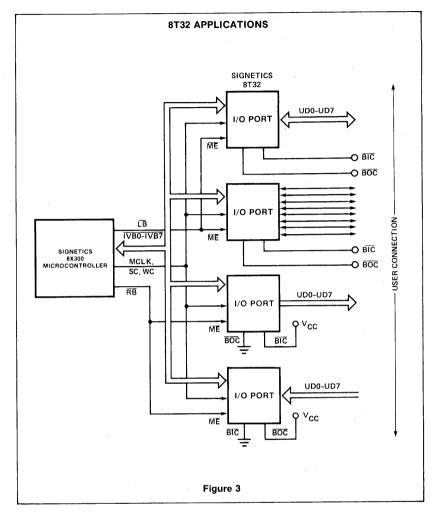
		TEST		LIMITS			
	PARAMETER	CONDITIONS	Min	Тур	Max	UNITS	
V _{CCP}	Programming supply voltage Address Protect		7.5	0	8.0	V V	
ICCP	Programming supply current	V _{CCP} = 8.0V			250	mA	
	Max time V _{CCP} > 5.25V				1.0	s	
	Programming voltage Address Protect		17.5 13.5		18.0 14.0	V V	
	Programming current Address Protect				75 150	mA mA	
	Programming pulse rise time Address Protect		.1 100		1	μs μs	
	Programming pulse width		.5		1	ms	

NOTE

^{1.} If all programming can be done in less than 1 second, VCC may remain at 7.75V for the entire programming cycle

APPLICATIONS

Figure 3 shows some of the various ways to use the I/O Port in a system. By controlling the BIC and BOC lines, the device may be used for the input and output of data, control, and status signals. I/O Port 1 functions bidirectionally for data transfer and I/O Port 2 provides a similar function for discrete status and control lines. I/O Ports 3 and 4 serve as dedicated output and input ports, respectively.



8T39-I.XL

DESCRIPTION

The Bus Expander is specifically designed to increase the I/O capability of 8X300 systems previously limited by fanout considerations. The bus expander serves as a buffer between the 8X300 and blocks of I/O devices. Each bus expander can buffer a block of 16 I/O ports while only adding a single load to the 8X300.

FEATURES

- 15ns max propagation delay
- Bidirectional
- Three-state outputs on both ports
- Pre-programmed address range

APPLICATIONS

The 8T39 Bus Expander is designed to be used with the 8X300 microprocessor to allow increased I/O capability in those systems previously limited by fanout considerations. Figure 1 shows a typical arrangement of the bus expander in an 8X300 system. Each expander services I/O ports whose address is within the range of the expander. Other I/O ports or working storage may be directly connected to the bus as shown.

The bus expander is not limited to use with the 8X300, but may be applied in any system which uses a combined address/data bus.

8T39 ADDRESSING

During normal operation of the 8X300 when an I/O port address is being sent on the IV Bus (SC is high), the I/O port will examine all eight bits of the microprocessor bus for an address compare. Since the 8T39 is used to buffer blocks of I/O ports, only the four most significant bits are examined by the 8T39 for an address compare.

Note that redundant addresses are not programmed into separate devices. Rather, a discrete device (such as the 8T39-03) may be wired for any address requiring two 1 bits and two 0 bits in the address. The various address ranges for this same device are obtained by permuting the high order (DI0 and DO0 are MSB) data lines accordingly. Both input and output lines must be redefined in order to maintain data and address integrity on the extended bus. Table 1 summarizes the 8T39 addressing.

Address functions are specified with the convention that bit 0 is the MSB and bit 7 is the LSB. The DI microprocessor bus address decoding is active low.

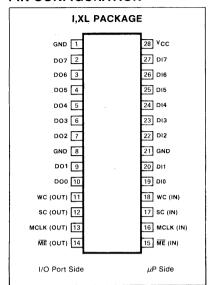
FUNCTIONAL DESCRIPTION

The Bus Expander contains eight sets of non-inverting bidirectional tri-state drivers for the bus data bits, four non-inverting

unidirectional drivers for I/O port control, and necessary control logic. The control logic is required to maintain the proper directional transfer of bus data as dictated by the states of the I/O port control signals and the currently enabled I/O port. Each bus expander is programmed during manufacturing to respond to a specific block of I/O port addresses. Only I/O ports with addresses in the range of a given bus expander may be connected to that expander. A bus expander may be used on either left bank or right bank. Multiple expanders on the same bank must have different address ranges; however, expanders with the same address range can be connected if they are on different banks. Systems may be configured with I/O ports connected directly to the 8X300, as well as I/O ports connected through a bus expander; however, no unbuffered I/O port may have an address within the span of a bus expander on the same bank.

Addition of bus expanders may impact system cycle time due to the added delay in the data path. For the purposes of calculating allowable cycle time as described in the 8X300 data sheet, the bus expander delays

PIN CONFIGURATION



may be considered additive to the I/O port delays so that a buffered I/O port simply appears as a slower I/O port.

PIN DESIGNATION

PIN NO.	SYMBOL	NAME AND FUNCTION	TYPE
2-7,9,10	DO0-DO7	I/O port data bus	Active low, three-state
11	WC(OUT)	Write command output	Active high
12	SC(OUT)	Select command output	Active high
13	MCLK(OUT)	Master clock output	Acitve high
14	ME(OUT)	Master enable output	Active low
15	ME(IN)	Master enable input	Active low
16	MCLK(IN)	Master clock input	Active high
17	SC(IN)	Select command output	Active high
18	WC(IN)	Write command output	Active high
19,20,22-27	DI0-DI7	Microprocessor data bus	Active low, three-state
1,8,21	GND	Ground	
28	Vcc	+5 volt supply	

ADDRESS PATTERN ADDRESS BLOCKS TYPE MSB(0) LSB(7) Octal			
0000XXXX	0-17		
0001XXXX	20-37, 40-57, 100-117, 200-217		
0011XXXX	60-77, 120-137, 220-237, 140-157, 240-257, 300-317		
0111XXXX	160-177, 260-277, 320-337, 340-357		
1111XXXX	360-377		
	MSB(0) LSB(7) 0000XXXX 0001XXXX 0011XXXX 0111XXXX		

Table 1 8T39 ADDRESSING SUMMARY

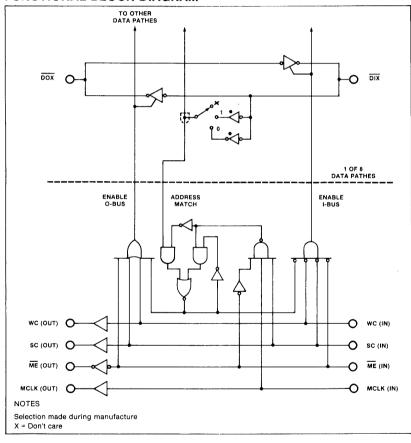
8T39-I,XL

TRUTH TABLE

ME	sc	wc	MCLK	SELECT LATCH	DATA TRANSFER DIRECTION	ADDRESS* COMPARISON
L	L	L	x	Set	DI Bus ← DO Bus	. No
L	L	L	x	Not set	DI Bus → DO Bus	No
L	L	Н	x	X	DI Bus → DO Bus	No .
L	н	х	L	X	DI Bus → DO Bus	No
L	Н	Х	н	X	DI Bus → DO Bus	Yes
Н	Х	×	X	Χ	DI Bus - DO Bus	No

NOTES

FUNCTIONAL BLOCK DIAGRAM



^{*}When an address comparison is made, the select latch is set if the data on the DI Bus is within the manufactured address range of the IV Bus Expander. Otherwise, the select latch is cleared.

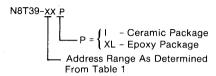
8T39-I.XL

ABSOLUTE MAXIMUM RATINGS¹

	PARAMETER	RATING	UNIT
Vcс	Power supply voltage	+7	Vdc
v_{IN}	Input voltage	+5.5	Vdc
٧o	Off-state output voltage	+5.5	Vdc
T_A	Operating temperature range	0 to +70	°C
^T STG	Storage temperature range	-65 to +150	°C

ORDERING INFORMATION

The Bus Expander is ordered by specifying the following part number:



DC ELECTRICAL CHARACTERISTICS V_{CC} = 5V ± 5%, 0°C \leqslant T_A \leqslant 70°C

	PARAMETER	TEST COMPLETIONS	LIMITS			116117	
	PARAMETER	TEST CONDITIONS	Min	Тур	Max	UNIT	
VIL VIH VIC	Input voltage Low High Clamp	-5mA at V _{CC} min	2.0		.8 -1	V	
V _{OL} VOH	Output voltage Low High	VCC = 4.75V IOL = 16mA IOH = -3.2mA	2.4	- -	.55	V	
I _{IL} ИН	Input current Low* High*	V _{CC} = 5.25V V _{IL} = .5V V _{IH} = 5.25V		< 10	-250 100	uA	
los lcc	Short circuit output current Supply current	V CC = 4.75V V CC = 5.25V	-40		200	mA mA	

AC ELECTRICAL CHARACTERISTICS V_{CC} = 5V ± 5%, 0°C \leqslant T_A \leqslant 70°C, CL¹ = 300pF²

PARAMETER		TO FROM		TEST CONDITIONS	LIMITS			LIMIT
	FARAIWETER	, FROM		1EST CONDITIONS	Min	Тур	Max	UNIT
tpd	Propagation Delay Data	DOX DIX	DIX DOX		-		15	ns
t _{pd}	Control Propagation Delay	ME (out) MCLK (out) SC (out) WC (out)	ME (in) MCLK (in) SC (in) WC (in)				15	
toe	Data Output Enable	DIX DOX	ME (in) SC (in) WC (in)		28		56	ns
tod	Data Output Disable	DIX DOX	ME (in) SC (in) WC (in)		15		-	ns
^t setup	Adverse Setup Time ³	DIX DOX	DIX ME (in) MCLK (in) SC (in) WC (in)		54			ns
^t hold	Address Hold Time ³	DIX DOX	DIX ME (in) MCLK (in) SC (in) WC (in)		3			ns

NOTES

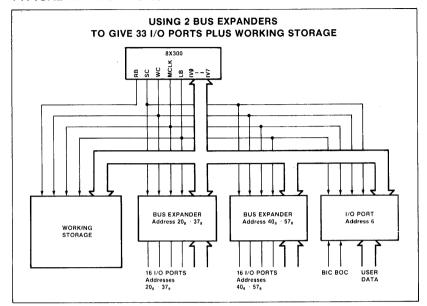
^{1.} Includes tri-state leakage.

^{2.} Minimum clock width \approx 50ns.

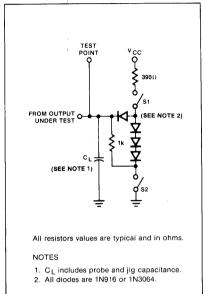
All set up and hold times are referenced to the trailing edge of the clocking input MCLK.

8T39-I,XL

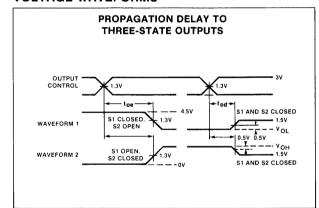
TYPICAL APPLICATIONS

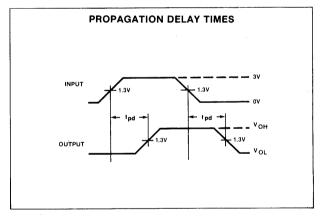


TEST LOAD CIRCUIT



VOLTAGE WAVEFORMS





PRELIMINARY SPECIFICATION

Manufacturer reserves the right to make design changes and improvements.

8T58-I.XL

DESCRIPTION

The Bus Expander is specifically designed to increase the I/O capability of 8X300 systems previously limited by fanout considerations. The bus expander serves as a buffer between the 8X300 and blocks of I/O devices. Each bus expander can buffer a block of 16 I/O ports while only adding a single load to the 8X300.

FEATURES

- 15ns max propagation delay
- Bidirectional
- Three-state outputs on both ports

FUNCTIONAL DESCRIPTION

The Bus Expander contains eight sets of non-inverting bidirectional tri-state drivers for the bus data bits, four non-inverting undirectional drivers for I/O port control, and necessary control logic. The control logic is required to maintain the proper directional transfer of bus data as dictated by the states of the I/O port control signals. A bus expander may be used on either left bank or right bank. Systems may be configured with I/O ports connected directly to the 8X300, as well as I/O ports connected through a bus expander.

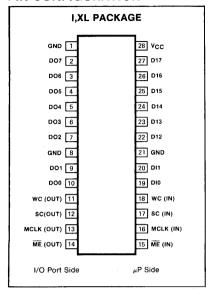
Addition of bus expanders may impact system cycle time due to the added delay in the data path. For the purposes of calculating allowable cycle time as described in the 8X300 data sheet, the bus expander delays may be considered additive to the I/O port delays so that a buffered I/O port simply appears as a slower I/O port.

APPLICATIONS

The 8T39 Bus Expander is designed to be used with the 8X300 microprocessor to allow increased I/O capability in those systems previously limited by fanout considerations. Figure 1 shows a typical arrangement of the bus expander in an 8X300 system. Other I/O ports or working storage may be directly connected to the bus as shown.

The bus expander is not limited to use with the 8X300, but may be applied in any system which uses a combined address/data bus.

PIN CONFIGURATION



TRUTH TABLE

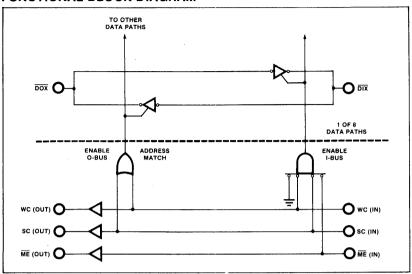
sc	wc	DATA TRANSFER DIRECTION	ADDRESS COMPARISON
L	L	DI Bus ← DO Bus	No
L	Н	DI Bus → DO Bus	No
. н	- X	DI Bus → DO Bus	No
Х	Х	DI Bus → DO Bus	No
	L L	L L L H	SC WC DIRECTION L L DI Bus ← DO Bus L H DI Bus → DO Bus H X DI Bus → DO Bus

PIN DESIGNATION

PIN NO.	SYMBOL	NAME & FUNCTION	TYPE
2-7,9,10	D00-D07	I/O port data bus	Active low, three-state
11	WC(OUT)	Write command output	Active high
12	SC(OUT)	Select command output	Active high
13	MCLK(OUT)	Master clock input	Active high
14	ME(OUT)	Master enable output	Active low
15	ME(IN)	Master enable input	Active low
16	MCLK(IN)	Master clock input	Active high
17	SC(IN)	Select command output	Active high
18	WC(IN)	Write command output	Active high
19,20,22-27	DI0-DI7	Microprocessor data bus	Active low, three-state
1,8,21	GND	Ground	
28	Vcc	+5 volt supply	

8T58-I,XL

FUNCTIONAL BLOCK DIAGRAM



ABSOLUTE MAXIMUM RATINGS

	PARAMETER	RATING	UNIT
Vcc	Power supply voltage	+7.	Vdc
VIN	Input voltage	+5.5	Vdc
Vo	Off-state output voltage	+5.5	Vdc
T_A	Operating temperature range	0 to +70	l °C
TSTG	Storage temperature range	-65 to +150	°C

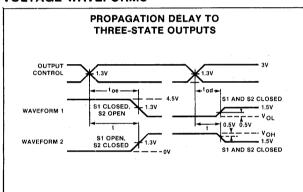
NOTE Includes tri-state leakage.

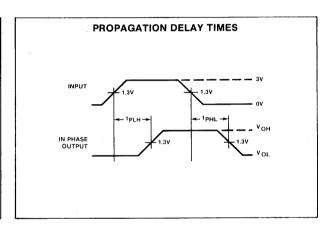
AC ELECTRICAL CHARACTERISTICS $V_{CC} = 5V \pm 5\%$, 0°C $\leq T_A \leq 70$ °C, $C_L = 300 pF$

						LIMITS Min Typ		
PA	RAMETER	то	FROM	TEST CONDITIONS	Min			UNIT
ťpd	Path delay Data	DOX DIX	DIX DOX				15	ns
t _{pd}	Control	ME(OUT) MCLK(OUT) SC(OUT) WC(OUT)	ME(IN) MCLK(IN) SC(IN) WC(IN)				15	ns
t _{oe}	Data Output Enable	DIX DOX	ME(IN) SC(IN) WC(IN)		28		56	ns
tod	Data Output Disable	DIX DOX	ME(IN) SC(IN) WC(IN)		15			

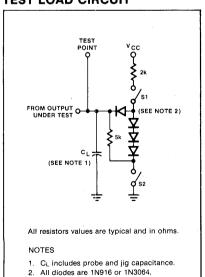
PARAMETER		PARAMETER TEST CONDITIONS			LIMITS				
		TEST CONDITIONS	Min	Тур	Тур Мах				
	Input voltage					l v			
VIL	Low		1		.8				
VIH	High		2.0						
V_{IC}	Clamp	-5mA at V _{CC} min			-1				
	Output voltage	V _{CC} = 4.75V				V			
V_{OL}	Low	$I_{OL} = 50 \text{mA}$		}	.55				
Vон	High	$I_{OH} = -3.2 \text{mA}$	2.4		1				
	Input current	V _{CC} = 5.25V				μΑ			
l _{IL}	Low1	V _{IL} = .5V			-250				
Iн	High ¹	$V_{IH} = 5.25V$		<10	100				
los	Short circuit output current	V _{CC} = 4.75V	-40			mA			
lcc	Supply current	$V_{CC} = 5.25V$			200	mA			

VOLTAGE WAVEFORMS

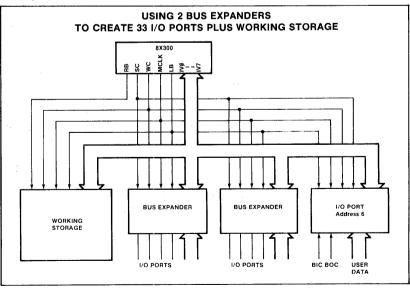




TEST LOAD CIRCUIT



TYPICAL APPLICATION



8X01-A.F

OBJECTIVE SPECIFICATION

DESCRIPTION

The CRC Generator/Checker circuit is used to provide an error detection capability for serial digital data handling system. The serial data stream is divided by a selected polynomial and the division remainder is transmitted at the end of the data stream, as a Cyclic Redundancy check character (CRCC). When the data is received, the same calculation is performed. If the received message is error-free, the calculated remainder should satisfy a predetermined pattern. In most cases, the remainder is zero except in the case where Synchronous Data Link Control type protocols are used whereby the correct remainder is checked for 1111000010111000 (x0 - x15)

8 polynomials are provided and can be selected via a 3-bit control bus. Popular polynomials such as CRC-16 and CCITT are implemented. Polynomials can be programmed to start with either all zeros or all ones.

Automatic right justification for polynomials of degree less than 16 is provided.

FUNCTIONAL DESCRIPTION

The CRC Generator/Checker circuit provides a means of detecting errors in a serial data communications environment. A binary message can be interpreted as a binary polynomial H(x). This polynomial can be divided by a generator polynomial P(x) such that H(x) = P(x) Q(x) + R(x) whereby Q(x) is the quotient and P(x) is the remainder. During transmission, the remainder is appended to the end of the message as check bits. For a given message, a unique remainder is generated. Hardware implementation of division is simply a feedback shift register with Exclusive-OR gating. Subtraction and addition in modulo 2 is implemented by the Exclusive-OR function. The number of shift register stages is equal to the degree of the divisor polynomial.

Table 1 shows the polynomials implemented in the CRC circuit. Each polynomial can be selected via the 3-bit polynomial control inputs S_0 , S_1 and S_2 . To generate the check bits, the data stream is entered via the Data (D) input, using the high to low transition of the Clock (CP) input. This data is gated with the most significant output (Q) of the register, and controls the exclusive OR gates. The Check Word Enable (CWE) must be held high while the data is being entered. After the last data bit is entered, the CWE is brought low and the check bits are shifted out of the register and appended to the data bits using external gating.

FEATURES

- I²L technology
- TTL inputs/outputs
- 10MHz (max) data rate
- Total power dissipation = 175mw (max)
- V_{CC} = 5.0V
- V j j = 1.0V
- . Separate preset and reset controls
- SDLC specified pattern match
- Automatic right justification

TYPICAL APPLICATIONS

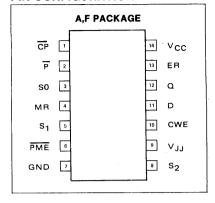
- · Floppy and other disc systems
- · Digital cassette and cartridge systems
- · Data communication systems

To check an incoming message for errors, both the data and check bits are entered through the D input with the CWE input held high. The 8X01 is not in the data path. but only monitors the message. The Error output becomes valid after the last check bit has been entered into the 8X01 by a high to low transition of CP. If no detectable errors have occurred during the data transmission, the resultant internal register bits are all low and the Error output (ER) is low. If a detectable error has occurred, ER is high. ER remains valid until the next high to low transition of CP or until the device has been preset or reset. PME must be high if ER output is used to reflect all zero result.

For data communications using the Synchronous Data Link Control protocol (SDLC), the 8X01 is first preset to all ones before any accumulation is done. This applies to both transmitter and receiver.

A special pattern of 1111000010111000 ($x^0 - x^{15}$) is used in place of all zeros during receiving for valid message check. PME is incorporated to select this option. If PME is low during the last bit time of the message, ER output is low if result matches this special pattern. When ER is high, error has occurred.

PIN CONFIGURATION



PIN DESIGNATION

PIN NO.	FUNCTION
S ₀ , S ₁ , S ₂	Polynominal Select inputs
D	Data input
CP	Clock (operates on high to low transition) input
CWE	Check Word Enable
P	Preset (active low) input
MR	Master Reset (active high) input
Q	Data output
ER	Error (active high) output
PME	Pattern match enable (active low)

A high level on the Master Reset (MR) input asynchronously clears the register. A low level on the Preset (P) input asynchronously sets the entire register if the control code inputs specify a 16-bit polynomial; in the case of the 12 or 8-bit check polynomials only the most significant 12 or 8 register bits are set and the remaining bits are cleared.

For SDLC, the user must invert the check sum shifted out of the 8X01.

RECOMMENDED OPERATING CONDITIONS

	ADAMETED		UNIT		
P.	ARAMETER	Min	Тур	Max	
V _{CC}	Supply voltage Supply current	4.75 40	5.0	5.25 100	V mA
CP	Clock input	0	,	5	MHz

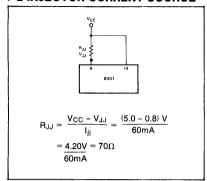
OBJECTIVE SPECIFICATION

8X01-A,F

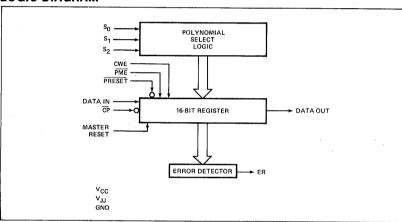
TRUTH TABLE

SE	LECT CO	DDE		
S ₂	S ₁	s ₀	POLYNOMIAL	REMARKS
L	L	L	$\chi^{16} + \chi^{15} + \chi^{2} + 1$	CRC-16
L	L	н	$\chi^{16} + \chi^{14} + \chi^{+1}$	CRC-16 REVERSE
L	н	L	$\chi^{16} + \chi^{15} + \chi^{13} + \chi^{7} + \chi^{4} + \chi^{2} + \chi^{1} + 1$	
L	н	Н	$\chi^{12} + \chi^{11} + \chi^{3} + \chi^{2} + \chi + 1$	CRC-12
н	L	L	X8+X7+X5+X4+X+1	
Н	L	н	X8+1	LRC-8
Н	Н	L	χ16 ₊ χ12 ₊ χ5 ₊₁	CRC-CCITT
Н	Н	н	$\chi^{16} + \chi^{11} + \chi^{4} + 1$	CRC-CCITT REVERSE

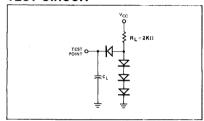
I²L INJECTOR CURRENT SOURCE



LOGIC DIAGRAM



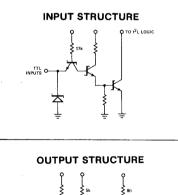
TEST CIRCUIT

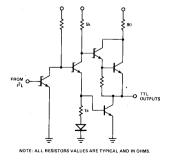


DC CHARACTERISTICS OVER OPERATING TEMPERATURE RANGE (Unless Otherwise Noted)

PARAMETER		TEST CONDITIONS				
	PANAWIE I EN	TEST CONDITIONS	Min	Тур	Max	UNIT
V _{IH}	Input high voltage		2.0			٧
VIL	Input low voltage				0.8	V
VIС	Input clamp diode voltage	V _{CC} = MIN, I _{IN} = -18mA			-1.5	V
Vон	Output high voltage	V _{CC} = MIN, I _{OH} = -400μA, I _{JJ} = MIN	2.7			٧
VOL	Output low voltage	V _{CC} = MIN, I _{OL} = 8mA,			0.5	٧
IJL	Input low current	V _{CC} = MAX, V _{IN} = 0.4V			-0.36	mA
11H	Input high current	V _{CC} = MAX, V _{IN} = 2.7V			20	μΑ
ЧН	Max. input current	V _{CC} = MAX			0.1	mA
los	Output short circuit current	CC = MAX, VOUT = OV, I _{JJ} = MIN	-10		-42	mA
Toc	Supply current	V _{CC} = MAX, inputs open	10		18	mA
IJ	Injection current	V _{CC} = MAX, inputs open V _{CC} = MAX, inputs open	60		100	mA

INPUT/OUTPUT CIRCUITS



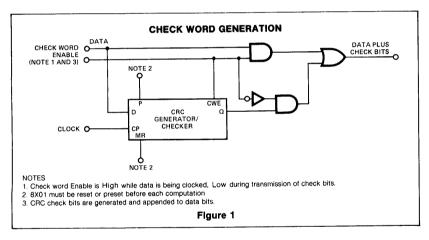


OBJECTIVE SPECIFICATION

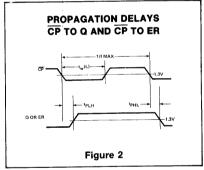
8X01-A,F

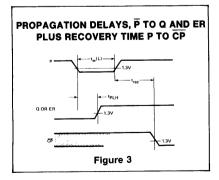
AC ELECTRICAL CHARACTERISTICS

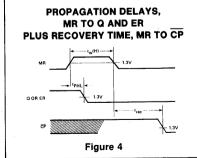
			TOT CONDITIONS		LIMITS	3	UNIT
PARAMETER	то	FROM	TEST CONDITIONS	Min	Тур	Max	Olti.
f _{max} Maximum clock frequency Pulse width				5	7		MHz ns
t _w CP(L) Clock low t _w P(L) Preset low t _w MR(H) Master reset high			See Figure 2 See Figure 3 See Figure 4	100 120 150			
Setup and hold time t _S D Setup time t _S CWE Setup time t _D Hold time	Clock Clock Clock	Data CWE Data, CWE	See Figure 5	0	120 75 -30	150 100 0	ns
Propagation delay tPLH Low to high tPHL High to low	Data output	Clock, preset	See Figures 1,2,3		100	160	ns
tpLH Low to high tpHL High to low	Data output	MR	See Figure 4		75	100	
tpLH Low to high tpHL High to low	Error output	Clock, MR, preset	See Figures 2,3,4		150	ļ	
tREC Recovery time	Clock	MR, preset	See Figures 3,4		60	90	ns

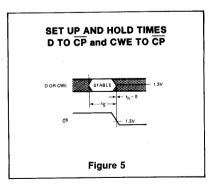


VOLTAGE WAVEFORMS









8X08-F.XA

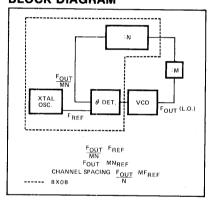
DESCRIPTION

This LSI integrated circuit performs the digital control functions required for generating AM/FM radio frequency local oscillator signals using digital phase locked loop techniques. By the use of low power Schottky and ECL technologies on the same substrate it is possible to operate at 80MHz input frequencies with an average system power of 1.6mW per gate typical.

FEATURES

- 80MHz input frequency
- ECL prescaler
- LS process
- Single 5V supply
- Power dissipation—600mW (max)
- External components—
 - 1 crystal
 - 2 capacitors

PHASE LOCKED LOOP BLOCK DIAGRAM



PHASE LOCKED LOOP PRINCIPLES

Digital phase locked loops are comprised of 4 basic building blocks: A fixed reference frequency generator (crystal oscillator and divider), a phase comparator, a voltage controlled oscillator (VCO) and a programable counter (÷N).

In cases where very high frequencies must be generated, a fixed prescaler (÷M) is employed to divide the local oscillator frequency down to a frequency compatible with the programmable counter. Fout from the VCO is divided down by the prescaler and programmable counters and compared to the reference frequency by the phase detector. If Fout is not equal to Fref in phase and

frequency, the phase detector generates a signal which causes the VCO frequency to

increase or decrease until F ref = $\frac{\text{Fout.}}{\text{MN}}$ When

this occurs, the local oscillator is essentially as stable as the crystal reference oscillator.

The local oscillator frequency (Fout) is changed by programming a different number into the programmable counter. The distance between discrete frequencies or the channel spacing is determined by the reference frequency.

For the AM/FM circuit, up to 200 channels are possible with selectable channel spacing of 10kHz for AM operation and 2000 channels at 100kHz for FM operation.

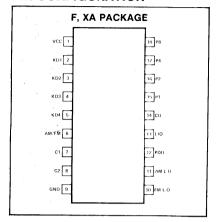
AM/FM Frequency Synthesizer Circuit Description

The frequency synthesizer circuit logic diagram is shown below. Following is a description of each of the major blocks.

Programmable Counter

The programmable counter consists of 3 stages of decade counter plus a divide by 1 or 0 counter to divide by numbers up to 1999. BCD programming data is presented to the dividers in parallel form, one digit at a time. Parallel data is strobed into internal latches via strobe signals; one strobe for each digit. A ÷ 5 80MHz ECL prescaler precedes the programmable counter for FM operation. This prescaler plus an external 160MHz ÷ 2 flip-flop provide a ÷ 10 160MHz prescaler (÷M) function to scale the programmable counter input frequency down to 16MHz maximum. A logic control circuit bypasses the +M prescaler and the first decade counter for AM operation. By this technique, the channel spacing is programmable to 10kHz for AM operation and 100kHz for FM operation.

PIN CONFIGURATION



VCO

An externally provided integrator and voltage controlled oscillator must be provided to perform the complete frequency synthesizer function. The integrator converts the pulses that come from the phase detector into a dc signal that controls the output frequency of the voltage controlled oscillator. It is in the integrator part of the circuit that the critical loop constants are determined. The voltage controlled oscillator is normally a LC tuned oscillator with varactor diode tuning that is controlled by the do signals from the integrator. In this case, two are required, one for the AM band and one for the FM band. The FM oscillator output must be +5V ECL compatible while the AM oscillator must be TTL compatible.

RECOMMENDED OPERATING CONDITIONS

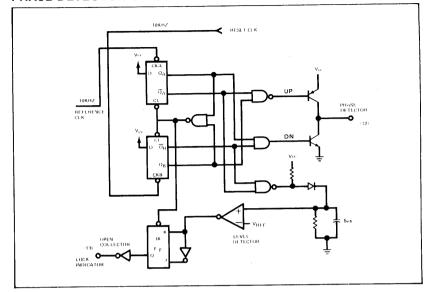
PARAMETER					
		Min	Тур	Max	UNIT
[™] A V _{CC}	Operating free air temperature Supply voltage Max input voltage	-40 4.75	5.0	+85 5.25	°C V
	data, strobe Max AM local oscillator input			16	V
	operating frequency (Pin 11) Max FM local oscillator input		20		MHz
	operating frequency (Pin 10) Maximum reference frequency		100	80	MHz
	oscillator operating frequency		6		MHz

Phase Detector Circuit

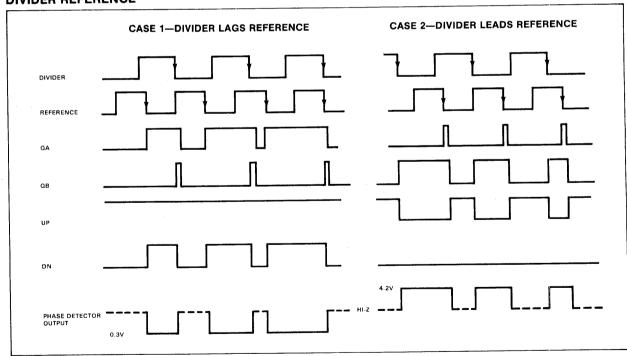
The phase detector is a digital edge-detecting device that provides an output three-state signal that is in a high impedance state when the 2 input signals are equal in phase and/or frequency. The output of the phase detector is a series of pulses that swing from the high impedance state to .3V typical or from the high impedance state to 4.2V typical. If the positive edge of the divider input leads the reference, the pulses will go to 4.2V. If it lags they will go to .3V.

The width of the output pulses is a function of the time between the positive edges (phase) of the 2 signals. An example of the operation of the device is shown where the reference signal is twice the frequency of the divider signal and has a phase lead of 270°. The output pulses are converted to a dc signal by integrating amplifiers causing the output frequency of the voltage controlled oscillator to increase or decrease (increase in this case) until the divider output and the reference output are equal in phase and frequency.

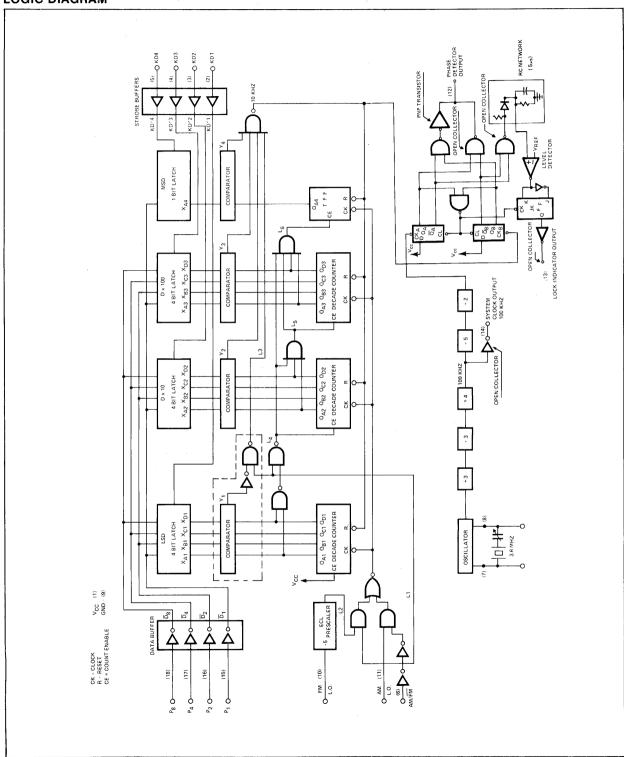
PHASE DETECTOR CIRCUIT



DIVIDER REFERENCE



LOGIC DIAGRAM



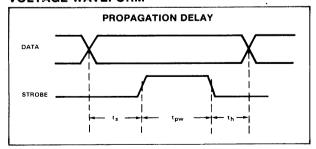
DC ELECTRICAL CHARACTERISTICS OVER RECOMMENDED OPERATING CONDITIONS

PARAMETER			LII	UNITS		
		TEST CONDITIONS	Min	Тур	Max	ONITS
V _{IH}	High level input voltage					
	P, KD, AM/FM inputs		5.25			V
	AM L.O. input		2			V
	FM L.O. input		4.1		5.25	V
V II	Low level input voltage					
	P, KD, AM/FM inputs				3.75	٧
	AM L.O. input				0.8	V
	FM L.O. input				3.3	V
IH	High level input current					
11 (P, Kp, AM/FM inputs	VCC = max, V ₁ = 16V			200	μ A
	, 3,	VCC = max, V ₁ = 5.25V			40	μΑ
	AM L.O. input					
	(with 5k Ω pullup to V _{CC})	VCC = max, V _I = 5.25V			200	μΑ
	FM L.O. input	VCC = max, V _I = 5.25V			400	μΑ
IL.	Low level inputs current					
16	P. Kp. AM/FM inputs	V _{CC} = max, V _I = 3.75V			-40	μ A
	AM L.O. input					
	(with 5k Ω pullup to VCC)	V _{CC} = max, V _I = 0.4V	7		-1.6	mA
	FM L.O. input	V _{CC} = max, V _I = 0.4V			-40	μ A
/OL	Low level output voltage					
0_	System clock output	VCC = min, IOL = 16mA			0.8	V
	Lock indicator output	V _{CC} = min, l _{OL} = 16mA			0.8	V
	Phase detector output	$V_{CC} = min, I_{OL} = 40\mu A$	1		0.5	V
/он	High level output voltage					
	Phase detector output	V _{CC} = min, I _{OH} = -40μA	VCC-0.5V			
loL	High level output current					
~ L	System clock output	VCC = min, VOH = 16V			250	μ A
	Lock indicator output	VCC = min, VOH = 16V			250	μΑ
CC	Supply current	V _{CC} = max			150	mA

AC ELECTRICAL CHARACTERISTICS

	PARAMETER		FROM		LIMITS		
	PARAMETER	то			Тур	Max	UNIT
	Strobe pulse width Setup and hold time			200	100		ns ns
ts	Logic high Logic low	Strobe	Data	150 50			
th	Logic high Logic low	Strobe	Data	40 0			

VOLTAGE WAVEFORM



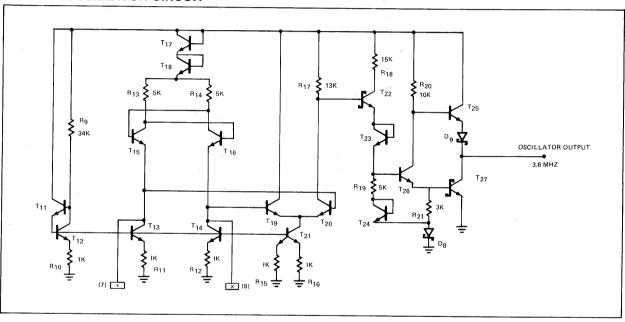
Crystal Oscillator Circuit

In this circuit, the cross-coupled transistor pair form a bistable circuit. The crystal provides positive feedback between the emitters of T₁₅ and T₁₆ which causes the circuit to oscillate at the crystal frequency.

Recommended Crystal Characteristics

Туре	Fundamental Mode Series Resonant		
Series resistance	< 100 Ω		

CRYSTAL OSCILLATOR CIRCUIT



DESCRIPTION

The Signetics Series 3000 Bipolar Microprocessor Chip Set provides new levels of high performance to microprocessor applications not previously possible with MOS technology. Combining the Schottky bipolar N3001 Microprogram Control Unit (MCU) and N3002 Central Processing Element (CPE) with industry standard memory and support circuits microinstruction cycle times of 100ns are possible.

In the majority of cases, the choice of a bipolar microprocessor slice, as opposed to an MOS device, is based on speed or flexibility of microprogramming. Starting with these characteristics, the design of the Signetics Series 3000 Microprocessor has been optimized around the following objectives:

- Fast cycle time
- All memory and support chips are industry
- Cooler operation
- Lower total system cost

Systems built with large-scale integrated circuits are much smaller and require less power than equivalent systems using medium and/or small scale integrated circuits.

The 2 components of the Series 3000 chip set, when combined with industry standard memory and peripheral circuits, allows the design engineer to construct highperformance processors and/or controllers with a minimum amount of auxiliary logic. Features such as the multiple independent address and data buses, tri-state logic, and separate output enable lines eliminate the need for time-multiplexing of buses and associated hardware.

Each Central Processing Element represents a complete 2-bit slice through the data processing section of a computer. Several CPEs may be connected in parallel to form a processor of any desired word length. The Microprogram Control Unit controls the sequence in which microinstructions are fetched from the microprogram memory (ROM/PROM), with these microinstructions controlling the step-by-step operation of the processor.

Each CPE contains a 2-bit slice of 5 independent buses. Although they can be used in a variety of ways, typical connections are:

Input M-bus:

Carries data from external

memory

Input I-bus:

Carries data from input/output device

Input K-bus:

Used for microprogram mask or literal (constant) value in-

put

Output A-bus:

Connected to CPE Memory

Address Register Connected to CPE accumu-Output D-bus:



As the CPEs are paralleled together, all buses, data paths, and registers are correspondingly expanded.

The microfunction input bus (F-bus) controls the internal operation of the CPE, selecting both the operands and the operation to be executed upon them. The arithmetic logic unit (ALU), controlled by the microfunction decoder, is capable of over 40 Boolean and binary operations as outlined in the FUNCTION DESCRIPTION section of the N3002 data sheet. Standard carry look-ahead outputs (X and Y) are generated by the CPE for use with industry standard devices such as the 74S182.

FEATURES

- Bipolar Schottky Technology
- Multiple Input/Output Bus Structure
- Fastest Microprocessor Available
- 512 Microinstruction Addressibility
- Full Function Accumulator

AVAILABILITY

Immediate delivery for Signetics Rep. or Distributors.

CONTENTS

1 ea--N3001 Microprogram Control Unit 4 ea-N3002 Central Processing Element

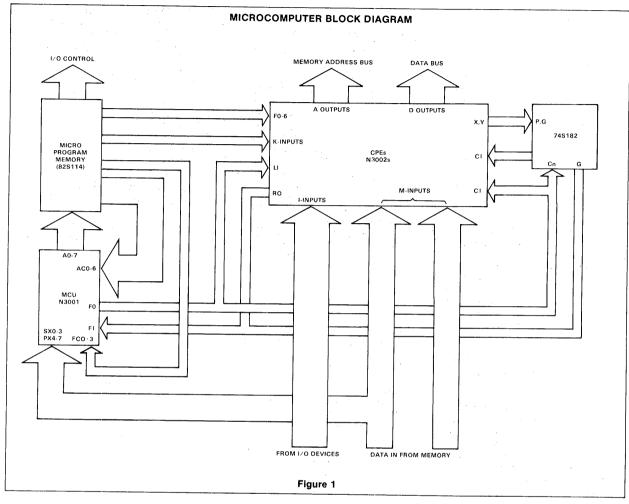
1 ea-74S182 Look-Ahead Carry

3 ea-82S114 256 × 8 PROM

1 ea-8T31 Bidirectional I/O Port

2 ea-8T26A Quad Bus Transceiver

1 ea-Introductory Manual



COMPATIBLE PRODUCTS 82S100, 82S101 FPLA

- Field programmable (Ni-Cr Link)
- Input variables—16
- Output functions—8
- Product terms—48
- Address access time—50ns
- Tri-state (82S100) or open collector (82S101) outputs
- 28-pin ceramic dip

82S115/123/129 PROMs

- Schottky TTL technology
- Single +5V power supply
- 32 × 8 organization (82S123)
- 256 × 8 organization (82S129)
- 512 × 8 organization (82S115)
- Field programmable (Nichrome)
- On-chip storage latches (82S115 only)
- · Low current pnp inputs

- Tri-state outputs
- 35ns typical access time
- Standard 24-pin DIP (82S115)
- Standard 16-pin DIP (82S123, 82S129)

82S25/82S116/82S11 RAMs

- Schottky TTL technology
- 16 × 4 organization (82S25)
- 256 × 1 organization (82S116)
- 1024 × 1 organization (82S11)
- On-chip address decoding
- 16-pin ceramic dip

8T26A/8T28 Quad Transceiver

- Schottky TTL technology
- · 4 pairs of bus drivers/receivers
- Separate drive and receive enable lines
- Tri-state outputs
- Low current pnp inputs
- High fan out-driver sinks 40mA
- 20ns maximum propagation delay
- Standard 16-pin DIP

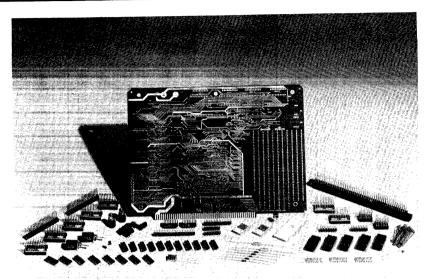
8T31 8-Bit Bidirectional Port

- · Schottky TTL technology
- 2 independent bidirectional buses
- · 8-bit latch register
- Independent read, write controls for each bus
- Bus A overrides if a write conflict occurs
- Register can be addressed as a memory location
- via Bus B Master Enable
- · 30ns maximum propagation delay
- Low input current: 500μA
- · High fan out-sinks 20mA
- Standard 24-pin DIP

DESCRIPTION

The 8080 Emulation Kit is a microprogrammable microprocessor utilizing Schottky LSI components to implement an emulation of an Intel 8080A microcomputer system. The emulation is functionally equivalent to a microprocessor system incorporating the following Intel devices: 8080A, 8228, 8224 and 8212. The kit provides the standard address, data, status and control buses as defined in the Intel 8080 Microcomputer System Manual. Since the kit uses bipolar LSI elements, the emulator lacks the twophase non-overlapping clock. Furthermore, those signals emanating from the 8080 during SYNC time are not provided, but rather the useful status signals provided by the 8228 system controller are implemented. The emulation also provides an extension of the 8228 operation during multi-byte interrupts. This is realized by allowing any 8080 program branch instruction to be inserted during interrupts rather than restricting multi-byte instructions to CALL during interrupts. Finally, a nonstandard status signal, RTRAP, is provided which indicates that the present instruction is a reserved or undefined instruction. After this indication, the processor will enter the normal HALT routine and await an interrupt. (Intel 8080A operation during undefined instructions is undefined.) Thus all 12 of the unused instructions in the 8080 instruction set are reserved for future instruction set expansion. These unused codes may be used at any time to extend the usual instruction set without requiring any reprogramming of the bipolar PROMs used for microprogram memory. Finally, the emulator is fully static so that the clock may be adjusted from a typical cycle time of 150ns to dc.

The kit contains all the parts necessary to construct the emulator and includes preprogrammed PROMs. The kit is designed to be assembled by a skilled technician in about 8 hours.



FEATURES

- Full emulation of 8080A system
- Speed increase by factor of 2 to 9.2 over 8080A system
- Static operation; microcycle time dc to 150ns
- Operation from single +5V supply
- **Executes all 8080 instructions**
- Hardware multiply and divide
- Microprogram expandable
- Includes single phase clock
- Full vectored interrupt to any location within 64K memory

Part Number: Availability:

3000KT8080SK Immediate delivery from Signetics, Rep or Distributors

KIT CON	ITENTS
1 each	N74123
1 each	N3001
8 each	N3002
7 each	N82S115
1 each	N82S23
2 each	N82S123
2 each	N82S126
3 each	N8263
3 each	N74S182
1 each	N74S280
2 each	N7475
1 each	DM8613
11 each	N74S174
2 each	N8T28
3 each	N8T97
1 each	N74S153
2 each	N74S157
1 each	N7400
1 each	N74S02
3 each	N74S04
1 each	N74S08
1 each	N74S10
1 each	N74S133
2 each	Resistor n
1 each	P.C. board
1 020h	Manual

etworks 1K, 16-pin

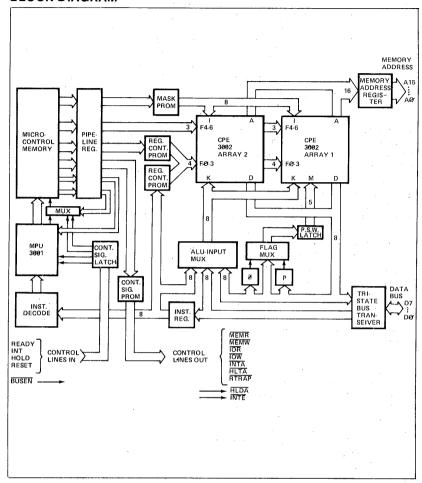
1 each Manual 1 each Schematic

Set of microprogram listings 1 each Over 25 miscellaneous Plus:

resistors, capacitors and other

parts

BLOCK DIAGRAM



INSTRUCTION EXECUTION

TIMES (150ns microinstruction cycle time, 150ns RAM access time)

l				Γ	
INSTRUCTION	ı	CLI	ES		CUTION ME (µs)
LXI	7			.05	
PUSH .	9			.35	
PUSH PSW	9			.35	
POP	9		1	.35	
POP PSW		. 8		l	.20
STA		9		l	.35 .20
LDA XCHG		7			.05
XTHL	13			.95	
SPHL		3			.45
PCHL		6			.90
DAD		4			.60
STAX		6			.90
LDAX		5			.75
INX		2			.30
DCX:		3			.45
CMA		2			.30
STC		3			.45
CMC DAA		10			.60 1.50
SHLD		11			i.65
LHLD		12			1.80
EI		2			.30
DI		2			.30
NOP		2			.30
MUL		26		3	3.90
MOV r1, r2		3			.45
MOV M, r	6			.90	
MOV r, M	6			.90	
HLT	4			.60	
MVIr	4 7		٠,	.60 .05	
MVI M INR	3	- 1	'	.45	
DCR		4			.60
INR M, DCR M		8	.	1	.20
Arithmetic reg		4	ı		.60
Arithmetic men	n	7		1	.05
Arithmetic imm	ed.	4			.60
RLC		4			.60
RRC		3	-		.45
RAL		4			.60
RAR		3			.45
JMP JMP	6			.90 .60	
CALL ₁	13	ł	1	.95	
CALL ¹		4		'	.60
RET		8		1	.20
RET ¹		2			.30
RST		13		1	.95
IN	8			.20	
OUT		8		1	.20
DIV	M	in	T	yp²	Max
DIV	4	80	F	6.60	8.40
					0.40

NOTES

- 1. Conditional branch with condition not met.
- 2. Depends upon value of divisor.

MICROASSEMBLER 8X02AS1000SS

DESCRIPTION

The Signetics Microassembler is a complete software package designed to support general slice system architectures which employ Signetics bipolar microprocessor components. It provides a flexible microassembler language which is adaptable to the numerous configurations and formats of systems using bipolar microprocessor chips. The Microassembly lanquage allows specification of the microinstruction formats and bipolar devices utilized in a user system. This enables the Microassembly language to be tailored to the specific configuration for which the microprogram is written. In particular, the Mircoassembly language provides intrinsic support for the 3002 and 2901 Central Processing Elements and the 8X02 Control Store Sequencer. Through the inclusion of explicit definitions, similar support can be obtained for the 3001 Microprogram Control Unit, as well as other bipolar chips. Although specifically intended for use with Signetics' bipolar microprocessor products, the flexibility of the Microassembler enables it to handle virtually all microprogrammed applications.

The Microassembler consists of two independent programs. The first reads the microprogram and the appropriate configuration and format descriptions written in the Microassembly language. It produces a listing of the source input and the resulting binary form of the microinstructions in the microprogram. The listing also includes diagnostics for any errors found in the source and a cross reference for symbols used in the microprogram.

The second program punches paper tapes that can be used to program microcontrol

store PROMS. It reads an object form of the microprogram produced by the first program. The microprogram object is partitioned into PROM modules and separate output is produced for each PROM.

MICROASSEMBLY LANGUAGE

The Microassembly language input to the Microassembler is divided into two sections. The first section defines the microinstruction formats and device configurations of the user system. The second section is the microprogram in the form of symbolic microinstructions.

The definition section specifies the name and length of each field in the microinstruction. Microinstructions may be any length, and multiple microinstruction formats may be defined. In addition, the definition section allows definition of opcodes used in writing symbolic microinstructions.

The microprogram section is the sequence of symbolic microinstructions which comprise the microprogram. A symbolic microinstruction is a statement specifying values for each field in the microinstruction. User defined opcodes allow mnemonic specification of field values. Field values not specified may be defaulted.

The Microassembly language also supports the standard assembler directives—ORG, EQU, SET, SPACE, TITLE, EJECT. Values in the Microassembly language may be decimal, hexadecimal, octal or binary constants, ASCII character constants, or symbolic values. These may also be combined into expressions using arithmetic, binary shift, and logical operators. The supported operators include +, -, SHR, SHL, NOT, AND, OR, XOR.

PROM FORMATTING

The microprogram object is formatted into paper tapes for programming microcontrol store PROMS using a microformat command file. The microformat commands specify the partitioning of the microprogram into PROMS and specify the format of the paper tape output. Each bit of a microinstruction may be individually allocated to a PROM and optionally inverted. The addressing space of the microprogram may also be partitioned with provisions for inverted addresses. A typical application would be the allocation of a 1024X40 microprogram to ten 512X8 PROMS.

The following paper tape formats are supported—2650 Absolute Object Code SMS format and various BNPF formats. The 2650 Absolute Object Code, can read by the Signetics TWIN (Test Ware Instrument—A microprocessor system development instrument) when it is used to program PROMS.

USING THE MICROASSEMBLER

The Microassembler is written in ANSI FORTRAN IV and it may be run on any machine of sufficient memory size which has the requisite standard FORTRAN compiler. It is available in source form on magnetic tape. Or, if desired, the Microassembler may be accessed via TYMSHARE, GE or NCSS Timesharing Services. The Microassembler also provides special commands called toggles which adapt the Microassembler to specific operating environments.

A full description of the Microassember is contained in the Signetics Microassembler Reference Manual.

DESCRIPTION

The Signetics 8X300 Fixed Instruction Bipolar Microprocessor provides new levels of high performance to microprocessor applications not previously possible with MOS technology.

In the majority of cases, the choice of a bipolar microprocessor slice, as opposed to an MOS device, is based on speed. The 8X300 processor, combined with high-speed memory and I/O devices, is capable of executing all instruction in 250ns.

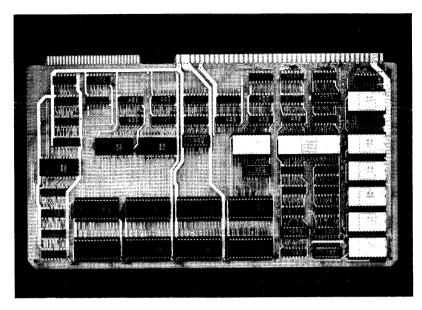
The 8X300 is optimized for control and data movement applications. It has a 13-bit address bus for selecting instructions from program storage and a separate input bus for entering a 16-bit instruction words. Data handling and I/O device addressing are accomplished via the 8-bit Interface Vector (IV) bus. The IV bus is supported by four additional control lines and a clock.

The unique features of the 8X300 IV bus and instruction set permit 8-bit parallel data to be rotated or masked before undergoing arithmetic or logical operations. Then, the data may be shifted and merged into any set of from 1 to 8 contiguous bits at the destination. The entire process of input, shifting, processing and output is done in 1 instruction cycle time. The 250ns cycle time makes the 8X300 ideally suited for high-speed applications.

The evaluation board contains all the elements which a designer needs to judge the suitability of the 8X300 for his systems applications. Included with the 8X300 are 4 I/O ports for external device interface, 256 bytes of temporary (working) data storage, and 512 words of program storage, all properly connected to the 8X300 to allow immediate exercising of the board. For this purpose, the PROMs are preprogramed with the I/O control, RAM control, and RAM integrity diagnostic programs. With the remaining PROM space, the designer may enter his own benchmark, test, or development routines.

The board design allows complete flexibility in access to the address, instruction, and IV busses as well as all controls and signals of the 8X300. The IV bus, I/O port user connection, clock signals control lines, address bus and instruction bus are wired to output pins, the board edge connector and flat cable connectors.

The board layout permits variations and/or expansions of the basic design. In addition to the access to all signals for transfer off the board, a wire wrap area is provided so that the designer may add to the board circuitry as he desires. The addition may include memory, additional interfaces, or



special circuits which meet specific user requirements.

Controls are also provided for diagnostic and instructional purposes by allowing various operating modes. In the WAIT mode, the program may be single stepped for ease of checkout. The one-shot instruction jamming allows control of the program start location, changes of program flow, changing or examining the internal registers, or testing of simple sequences. The repeated instruction jamming provides a means of repetitive execution of an instruction so that the I/O bus and the control lines may be examined without software changes. In both of these jam cases, the jammed instruction is selected by boardmounted switches.

FEATURES

- . 250ns CPU with Crystal
- 4 I/O Ports (32 Lines)
- 256 Bytes Data Storage
- 512 Words Program Storage
- Run/Wait Control
- Single Step
- Instruction Jamming One Shot Instruction Jam Repeated Jam
- All Buses to Output Pins
- Firmware Diagnostics
- Wire-Wrap Area
- Edge Connector
- Flat Cable Connectors
- Wire Wrap Posts for Bus Lines

CONTENTS

1 ea- 8X300

8 ea- 82S116 (256 x 1 RAM)

2 ea- 82S115 (512 x 8 PROM)

4 ea— 8T32 Addressable Bidirectional I/O Port

1 ea-8T31 (Bidirectional I/O Port)

2 ea— 8T26A (Quad Bus Transceiver)

4 ea — 74157 (Quad 2-Input Data Selector)

2 ea- 7474 (Dual D Flip Flop)

2 ea- 7400 (Quad Nand Gate)

1 ea- 7427 (3-Input NOR Gate)

1 ea— P.C. Board

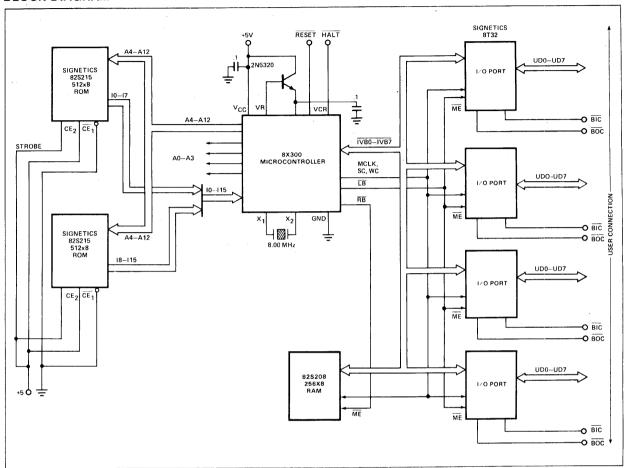
Misc. Parts

 1 ea— Introductory Manual, assembly instructions, code listings and schematics

AVAILABILITY

Immediate delivery from Signetics Rep. or Distributors.

BLOCK DIAGRAM



Auxiliary Circuits

The 8X300 can be used with any bipolar (or TTL-compatible) ICs. It can directly address 8192 program instruction locations and up to 512 I/O ports. The memory paging feature may be employed for larger working storage. Typical auxiliary circuits include:

Program I/O	Storage 8T32/33	82S115 (512x8 PROM) (8-Bit Synchronous Bidirectional I/O Port)
	8T35/36	(8-Bit Asynchronous Bidirectional I/O Port)
	8T31	(8-Bit Bidirectional I/O Port)
Working	8T39 Storage	(Quad Bus Extender) 82S116 RAM

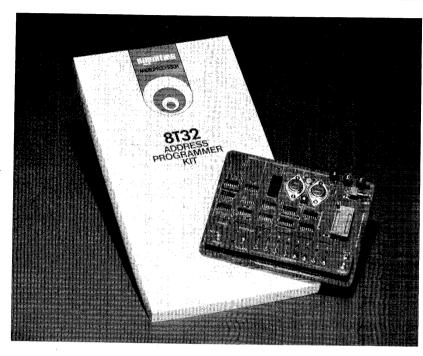
DESCRIPTION

A new kit has been developed which allows for convenient programming of the select addresses for the 8T32 family of devices, including the 8T33, 8T35 and 8T36. The kit consists of all parts necessary for the construction of the kit including voltage regulators, address program switches and status lights. All that is required is to construct the kit, apply +5 volts and approximately +28 volts and the 8T32 may be programmed. The programming sequence consists of three steps. First, the 8T32 select address is programmed. Second, the kit verifies the programming by testing for positive response to the address as well as the disable for the address complement. After verification, the programmer is switched to the protect mode in which the isolation is performed. This step prevents further address alteration. The kit includes documentation on assembly and use.

The kit is made available for those designers using 8T32 family devices so that they may easily program the address in their lab or at a conveniently located Signetics or Distributer facility. For prototyping, the designer may order unprogrammed devices and then put in addresses as required rather than ordering several addresses of the preprogrammed type. This will save the designer time and money.

AVAILABILITY

Immediate delivery from Signetics Rep. or Distributors



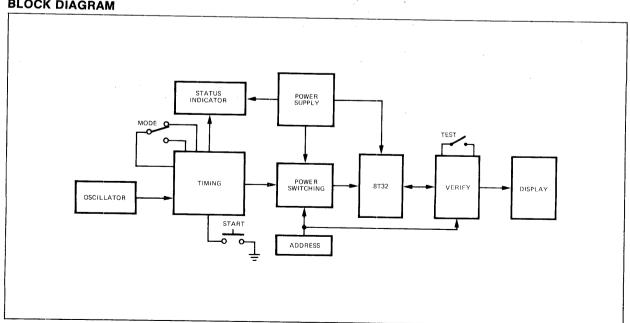
CONTENTS

PC board All IC's, discrete parts, and hardware Assembly instruction manual

FEATURES

- Programs, isolates, and verifies select addresses for 8T32 family
- Kit contains all components necessary to build programmer
- Two hour assembly time

BLOCK DIAGRAM



DESCRIPTION

The 8X300 Assembler is a program which accepts source code written in the 8X300 assembly language and which produces both a listing of the symbolic program and an object module in one of three formats: MCSIM, ROM Simulator or BPNF.

The assembler is written in ANSI standard Fortran and is approximately 2800 card images in length. It requires some rewindable I/O medium such as disk or tape for temporary storage, an input device to read source code and two output devices, one to output the assembler listing and one to output the object module. If desired, the assembler may be compiled and linked to execute in overlays.

The assembler consists of two passes which build a symbol table, issue helpful error messages, produce an easily readable program listing and output a computer readable object module.

FEATURES

- Macros nested to three levels
- Conditional assembly
- Address arithmetic
- Automatic procedure/subroutine handling
- Reserved symbols for registers
- Multiple entry to procedures
- Symbolic machine operation codes
- Symbolic address assignment references
- Symbolic data creation statements (IV **Byte References)**
- Free format source code
- Syntax error checking
- Self-defining constants
- Assembly listing control statements
- ASCII character code
- Comment statements and comment areas
- Forward referencing

MCCAP LANGUAGE

The assembler recognizes three types of statements: comment statements, machine instructions, and pseudo-ops (directives).

Comment statements are used to document a listing. They must contain an asterisk in column 1 and may contain any legal character in the other 79 columns.

Machine instruction statements are those statements that generate instruction(s) to be executed by the 8X300 processor. Machine instruction statements include:

Move byte from one location to MOVE another.

Add the byte in one location to the ADD byte in the auxiliary register and store the result in another location.

AND the byte in one location with AND the byte in the auxiliary register and store the result in another location.

XOR Exclusive OR the byte in one location with the byte in the auxiliary register and store the result in another location.

Execute the instruction at the XEC specified address.

Transmit the literal (immediate **XMIT** data) contained in this instruction to the specified location.

Jump to the specified address if NZT the value in the specified location is not zero.

Jump to the specified address. JMP Transfer control to a procedure or CALL an entry point.

Returns control from a procedure RTN or an entry point.

Places the address of an IV Byte SEL variable into the IVL or the IVR register.

Pseudo-Ops (or Directives) are commands to the assembler, but have no meaning to the 8X300 processor.

Assigns a value to the symbol EQU in the label field.

Identical to EQU, except that SET the symbol may be redefined by another SET statement.

Assigns a symbolic name to LIV the address, position and precision of the "left bank" Interface Vector Byte variable.

Identical to LIV, except that the RIV assignment is made to the "right bank" I/O Field.

Introduces and names the PROG main program.

and names Introduces **PROC** procedure.

Specifies an additional entry

ENTRY point to a procedure. Terminates a procedure, a pro-END

gram or a MACRO. Changes the value of the loca-ORG

tion counter. Specifies the format of the ob-OBJ iect module.

Introduces source statements ΙF subject to conditional assem-

blv. Terminates the source state-**ENDIF** ments subject to conditional

assembly. Specifies the list and punch LIST

options. Negates the list and punch op-NLIST tions.

Advances listing to next page. **EJCT** Advances listing to next line. SPAC Defines the source statements **MACRO** generated a macro call.

AVAILABILITY

The assembler is available on NCSS, GE and TYMSHARE timesharing services. It is also available from Signetics on 9-track magnetic tape written in EBCDIC in 80character unblocked records at a density of 800 bpi.

FEATURES

- Totally self-contained with keyboard alpha-numeric display, tape reader, TTY output
- Dual MicroControllers: one to run instrument, one dedicated to execute user's program
- Real time instruction execution
- Control of program execution-Halt Single Step and Run Modes
- Direct program/register/IV examination and modification through keyboard
- Three breakpoint types—Normal, Halt and Insert Instruction
- Up to 4k words of high speed read/write program storage
- Format compatible with papertape output of MicroController Cross Assembly Program



The SMS MicroController Simulator, MCSIM, is a hardware development instrument designed to perform in the user's system exactly as an SMS MicroController. It directly supports the SMS 300 (8X300) as well as MicroController prototyping systems. MCSIM gives the user an Interpreter system with a modifiable Program Storage and a Control Panel which together provide a means of entering, running, monitoring, debugging and changing a program. It features ease of use due to its high level interactive display/keyboard and simple operating system. MCSIM allows the user to run his program on-line in real time. Through the Control Panel, data stored in internal registers, Working Storage and IV Bytes can be examined and modified. An extensive breakpoint capability permits the user to quickly locate program faults.

MCSIM contains two MicroControllers: one for running the instrument and a second totally dedicated to the user's program and prototyping system. This insures that when program development is complete and the design specifications have been met, production systems will perform identically with the prototype.

OPERATION

MCSIM operation is separated into six modes, each mode consisting of a related set of functions and status displays. Access to a mode is made using one of the six mode selection keys. (See Table 1). The currently selected mode is displayed at the extreme left part of the display; the currently selected function is displayed at the extreme right of the display. Selection among the available functions is made using the 1 and 1 keys, incrementing and decrementing to the next function. The function selected at last exit is automatically re-selected when the mode is re-entered except for the Break-



point Mode which selects the Current display.

An operation is generally performed using the following four steps:

- Select a mode by depressing one of the six Mode Select Keys.
- 2. Select a function in that mode by posi-

tioning the function roll with one of the two Function Select Keys.

- Set up any necessary conditions, such as entering data or readying papertape in the reader.
- Activate the function by depressing the Function Acknowledge Key (FCN/ACK).

MODE	FUNCTIONS
Manual Examination and alteration of Program Storage	Change Address Store Instruction
Tape Load or dump all or part of Program Storage	Load Verify Begin Punch Address End Punch Address Punch Program Identification
Register Examination and alteration of the Interpreter's internal registers	Change Address Store Octal Data Store Binary Data
Interface Vector Examination and alteration of the locations on the IV Bus	Complement Overflow Display Current Enabled Bytes Change IV Address Store Binary Data in IV Location
Set one of three types of breakpoint for program debugging	Display Currently Active Breakpoint Set Normal (Sync) Breakpoint Set Stop Breakpoint Replace Instruction at Breakpoint
Execute Control and monitor execution of a program	Halt Run Program Insert Instruction Single Step

Table 1 MCSIM OPERATOR CONTROLLED FUNCTIONS

SPECIFICATIONS Control Panel

The Mode Select Keys and the Function Select Keys used in combination give the user 26 possible functions to accomplish complete loading and testing of the program. Each Mode Select Key accesses a set of functions, one of which is chosen using the Function Select Keys.

Status Messages

Message displays are used to indicate special conditions which may result from the operation of MCSIM:

POWER ON, MCSIM SYSTEM START ???MORE THAN ONE KEY PRESSED UNABLE TO READ TAPE. RESTART UNDEFINED MEMORY ADDRESS (octal address)

INVALID INSTRUCTION (instruction code)
LOADING INTERRUPTED, RESTART,
V'FYING INTERRUPTED, RESTART
P'CHING INTERRUPTED, RESTART

NO VERIFY (octal address, tape data, memory data)

CLEAR RESET LINE TO START RUNNING CLEAR HALT LINE TO START RUNNING INVALID 8-BIT OCTAL VALUE (octal value) UNDEFINED IV BYTE ADDRESS (octal address)

Sync Output

Output pulse whenever address breakpoint is reached.

System Interface

MCSIM is connected to the user's prototyping system via a single ribbon cable. This cable terminates in either a MicroController Simulation Module or a dual in-line plug which is pin for pin compatible with the SMS 300 Interpreter. There are two different simulation modules to exactly match presently available MicroController systems. Selection of the proper input/output connector makes a MCSIM appear to be physically and electrically equivalent to a production MicroController or Interpreter.

MCCAP

The MicroController Cross Assembly Program (MCCAP) is designed to translate symbolic instructions into object code that can be executed by the SMS 300. This program will run on most computers that have a FORTRAN compiler with a computer word length of at least 16 bits and a random access capability. MCCAP features:

- Symbolic address assignment and references
- Automatic Procedure Handling
- Predefined System Procedures
- Forward References
- Expression Evaluation
- Flexibility to handle MicroController component configurations as well as standard systems
- Generation of object code for MCSIM, and SMS ROM Simulator, or most PROM programmers

Data Input/Output

COMPONENTS	DESCRIPTION		
Panel	36 character self scan display for messages and data readout 12 key numeric keyboard for data entry/modification		
Tape Reader	120 character/second tape reader for high speed data entry (ASCII format)		
TTY	20mA current loop output for listing of program code on Teletype® terminal		

PHYSICAL CHARACTERISTICS

Power	115V or 230V \pm 10% 50 or 60Hz \pm 10% 350 watts/min. to 750 watts/max. (Power dissipation dependent on the number of Simulation Modules configured in the system)			
Dimensions Installation	7 inches high x 17 inches wide x 19 inches deep May be used on table or may be installed in standard 19 inch rack with mount adapters			
Weight	65 lbs. net, 75 lbs. shipping			
Ventilation	Air Flow 120 CFM			
Environmental 0° to 55°C, Relative Humidity to 90%.				

FEATURES

- Real time monitoring instrument for SMS 300
- · Totally self contained
- Displays IV address and data
- Displays current program address
- Displays current instruction
- Control of RESET and HALT
- Single step capability
- Real time instruction insertion
- Two real time breakpoints
- · Breakpoint output signal

DESCRIPTION

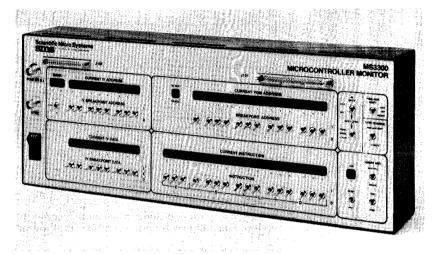
The MicroController Monitor is a self contained debug and maintenance tool for use with all systems containing the SMS 300 (8X300). It provides a control panel allowing an operator to observe, modify, and control program execution in real time permitting system faults to be rapidly traced. The Monitor displays the status of the Address, Instruction, and IV Bus of an Interpreter. Switch registers associated with each display can be used to set breakpoint addresses and enter instructions or data from the Monitor. Program execution can be halted, single stepped, or altered from Monitor controls.

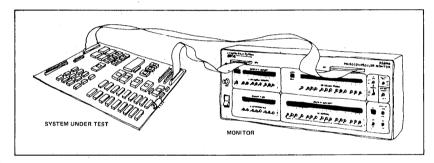
The Monitor can be used to insert instructions thereby examining or modifying the contents of internal registers, Working Store, and IV Bytes. In addition, it can be used to start program execution from any address and enable the operator to set program loops allowing a program segment to be checked independently of normal program flow. Breakpoints on IV data and program address allow the operator to halt program execution or insert instructions in real time.

As shown in the diagram below, the Monitor connects to the system under test through flat ribbon cables (provided) or an adaptor (optional) which plugs into an Interpreter socket. These connections are buffered and present a negligible load to the system. Other features include a Sync output which provides pulses at a selectable program address or Interface Vector event.

OPERATION

The MicroController Monitor provides the user with two basic modes of system operation, RUN and STEP. When in the RUN mode, indicator LED's provide a continuous real time display of the three major system busses: Address, Instruction, and Data. Program execution can be halted by one of two actions: depressing the Halt switch or selecting the Halt on Breakpoint function. When halted, it is possible to execute one instruction at a time by pressing the RUN/





STEP control to the STEP position. The readout shows the current (static) bus information and the next instruction to be executed.

For checkout purposes it is useful to be able to set a program loop or modify the normal progam flow in some other manner. This is accomplished by inserting in real time the instruction set up on the Instruction Switch Register whenever the Program Address Breakpoint is reached. To check for system noise or similar malfunction a separate Insert Instruction switch can be set to the Multiple position, forcing one instruction repetitively. When halted, a single instruction can be inserted by toggling the Insert Instruction switch to the Single position. When an instruction is being inserted, system ROM is temporarily disabled.

To aid in the diagnosis of a malfunction an advanced breakpoint capability is provided. In addition to the Halt and Insert functions when a Program Address Breakpoint is reached, a breakpoint can be generated on the Interface Vector (IV) Bus data or address. This permits program execution to be halted or a Sync pulse generated in re-

sponse to external data or the contents of RAM. Three modes of IV Breakpoint generation are selectable: Breakpoint on IV Address; Breakpoint on IV Address and IV Write Data; and Breakpoint on IV Address and IV Read Data. The IV Breakpoint Data switches have a "don't care" position to permit generation of a breakpoint on a data subfield. A Sync pulse output is provided on either the IV Breakpoint or Program Address Breakpoint (switch selectable).

SPECIFICATIONS

Controls	Function	
RUN/STEP	Sets operation mode, continuous execution (RUN) or single step	
INSERT INST	(STEP). Unconditionally causes execution of the instruction in the Instruc-	
RESET/HALT	tion Switch Register. Three position switch used to unconditionally HALT program execution or RESET Interpreter's Program Counter to zero.	
PGM ADR BKPT	Selects appropriate function to be performed when Program Address Breakpoint is reached: HALT—halt on breakpoint, INSERT—replace normal instruction with instruction set into Instruction Switch Register, OFF—program execution unchanged by breakpoint.	
IV BKPT	Selects appropriate function to be performed when IV Breakpoint is reached: HALT—halt on breakpoint, OFF—program execution unchanged by breakpoint.	
BKPT UPON	Selects IV Breakpoint on one of three conditions: specified IV address, specified write data at IV address, specified read data at the IV address.	
SYNC SOURCE	Selects source for output Sync pulse, IV Breakpoint or Program Address Breakpoint. Sync pulse always available at breakpoint regardless of function selected by breakpoint controls.	
Switch Registers		
INSTRUCTION	Sixteen two position switches for entering instruction to be inserted (during breakpoint or insert instruction control).	
BREAKPOINT ADDRESS	Thirteen two position switches to select program breakpoint address.	
IV BREAKPOINT ADDRESS	Eight two position switches to select IV byte breakpoint address or Working Store breakpoint address.	
BANK	Selects display of LB (Left Bank) or RB (Right Bank) IV address information. Also selects either LB or RB for IV breakpoint.	
IV BREAKPOINT DATA	Eight three position switches, 0, 1, X (don't care), to select IV Data Breakpoint.	
Displays		
RUN	LED display indicating when Interpreter not halted.	
CURRENT INSTRUCTION	LED display of the binary value of next instruction to be executed.	
CURRENT PROGRAM ADDRESS	LED display of the binary address of the next instruction to be performed.	
IN XEC RANGE	LED display indicating that the value of the Interpreter's program counter and address register may not match because the previous instruction command was "Execute."	
BANK	Displays band addressed by the current instruction (left bank or right bank).	
CURRENT IV ADDRESS	Shows the binary address of the currently enabled IV Byte or Working Store location.	
CURRENT IV DATA	Binary display of the data on the IV bus.	
Outputs		
SYNC PGM EN(L)	Pulse output whenever the switch selected breakpoint occurs. A low true signal used to enable or disable system program ROM (required only when Interpreter adaptor is used).	
Cycle Time		
	onitor can be adjusted to work with any MicroControllersystem with a cycle time of 200ns to 2µs.	

PHYSICAL CHARACTERISTICS

Power	115Vac \pm 10% at 0.75 amp max. 230Vac \pm 10% at 0.38 amp max. (optional) 47-63Hz
Dimensions	7" H x 17.5" W x 3" D
Installation	May be used on table or may be installed in standard 19 inch rack (using optional adaptors).
Weight	10 lbs. net, 13 lbs. shipping
Ventilation	Forced air, 120 CFM.
Environmental	0° to+50°C, relative hu- midity to 90%
Cables	2 provided, 3 feet long

CHAPTER 2 MOS Microprocessors, Peripherals and Development Products

signetics

2650A-I • 2650A-1-I

DESCRIPTION

The 2650A and -1 are additional members of the Signetics family of 8-bit, NMOS microprocessors.

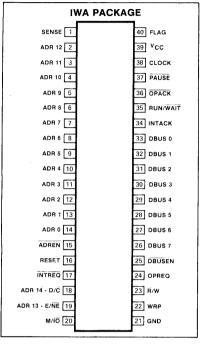
The 2650A is a functional equivalent of the 2650 with a new mask design which provides improved device operating margins.

The 2650A-1 is a high speed version of the 2650A.

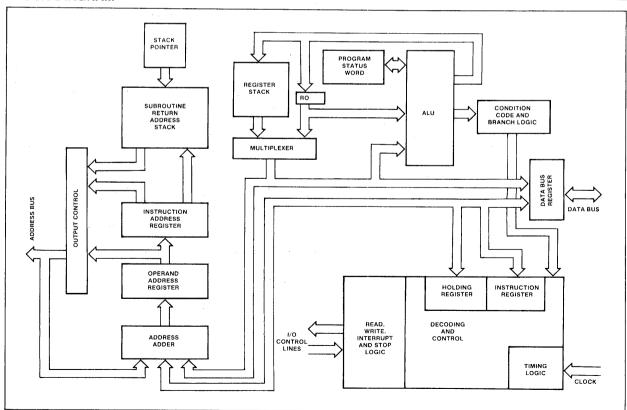
FEATURES

- Static 8-bit parallel NMOS microprocessor
- Single power supply of +5 volts
- TTL level single phase clock
- Standard 40 pin dual in-line package
- TTL compatible inputs and outputs
- 75 variable length instructions of 1, 2 or 3 bytes
- 32k byte address range
- Coding efficiency with multiple addressing modes
- Synchronous or asynchronous memory and I/O interface
- Interfaces directly with industry standard memories
- Single bit serial I/O path
- Seven 8-bit addressable general purpose registers
- Vectored interrupt
- Subroutine return address stack
- 2.4 µs machine cycle time (2650A)
- 1.5μs machine cycle time (2650A-1)

PIN CONFIGURATION



BLOCK DIAGRAM



2650A-I • 2650A-1-I

PIN DESIGNATION

MNEMONIC	NUMBER	NAME	TYPE	FUNCTION
ADR0-ADR12	14-2	Address lines	0	Low order memory address lines for instruction or operand fetch. ADR0 is the least significant bit and ADR12 is the most significant bit. ADR0 through ADR7 are also used as the I/O device address for extended I/O instructions.
ADR13-E/NE	19	Address 13- Extended/Non Extended	0	Low order memory page address line during memory reference instructions. For I/O instructions this line discriminates between extended and non-extended I/O instructions.
ADR14-D/C	18	Address 14- Data/Control	0	High order memory page address line during memory reference instructions. It also serves as the I/O device address for non-extended I/O instructions.
ADREN	15	Address enable	ı	Active low input allowing tri-state control of the address bus ADR0-ADR12.
DBUS0-DBUS7	33-26	Data bus	I/O	These lines provide communication between the CPU, Memory, and I/O devices for instruction and data transfers.
DBUSEN	25	Data bus enable	1	This active low input allows tri-state control of the data bus.
OPREQ	24	Operation request	0	Indicates to external devices that all address, data and control information is valid.
OPACK	36	Operation acknowledge	ı	Active low input indicating completion of an external operation. This allows asynchronous functioning of external devices.
M/ IO	20	Memory/input-output	0	Indicates whether the current operation references memory or I/O.
R∕W	23	Read/Write	0	Indicates a read or a write operation.
WRP	22	Write pulse	0	This is a timing signal from the 2650 that provides a positive-going pulse during each requested write operation (memory or I/O) and a high level during read operations.
SENSE	1	Sense	ŀ	The sense bit in the PSU reflects the logic state of the sense input to the processor at pin #1.
FLAG	40	Flag	0	The flag bit in the PSU is tied to a latch that drives the flag output at pin #40.
INTREQ	17	Interrupt request	I	This active low input line indicates to the processor that an external device is requesting service. The processor will recognize this signal at the end of the current instruction if the interrupt inhibit status bit is zero.
INTACK	34	Interrupt acknowledge	0	This line indicates that the 2650 is ready to receive the interrupt vector (relative address byte) from the interrupting device.
PAUSE	37	Pause	1	This active low input is used to suspend processor operation at the end of the current instruction.
RUN/WAIT	35	Run/Wait	0	This output is a processor status indicator. During normal operation this line is high. If the processor is halted either by executing a halt instruction or by a low input on the pause line, the run/wait line will go low.
RESET	16	Reset	1	Resets the instruction address register to zero and clears the interrupt inhibit bit.
CLOCK	38	Clock	ı	A positive going pulse train that determines the instruction execution time.
Vcc	39	+5V	1	+5V power
GND	21	GND	1	Ground

FUNCTIONAL DESCRIPTION

The 2650 series processors are general purpose, single chip, fixed instruction set, parallel 8-bit binary processors. A general purpose processor can perform any data manipulations through execution of a stored sequence of machine instructions. The processor has been designed to closely resemble conventional binary computers, but executes variable length instructions of one to three bytes in length.

The 2650 series contains a total of seven general purpose registers, each eight bits long. They may be used as source or destination for arithmetic operations, as index registers, and for I/O transfers.

The processor can address up to 32,768 bytes of memory in four pages of 8,192 bytes each. The processor instructions are one, two, or three bytes long, depending on the instruction. Variable length instructions tend to conserve memory space since a one-or-two byte instruction may often be used rather than a three byte instruction. The first byte of each instruction always specifies the operation to be performed and the addressing mode to be used. Most instructions use six of the first eight bits for this purpose, with the remaining two bits forming the register field. Some instructions use the full eight bits as an operation code.

The Data Bus and Address signals are tristate to provide convenience in system design. Memory and I/O interface signals are asynchronous so that Direct Memory Access (DMA) and multiprocessor operations are easy to implement.

The block diagram for the 2650 series shows the major internal components and the data paths that interconnect them. In order for the processor to execute an instruction, it performs the following general steps:

- The Instruction Address Register provides an address for memory.
- The first byte of an instruction is fetched from memory and stored in the Instruction Register.
- The Instruction Register is decoded to determine the type of instruction and the addressing mode.
- If an operand from memory is required, the operand address is resolved and loaded into the Operand Address Register.
- 5. The operand is fetched from memory and the operation is executed.
- 6. The first byte of the next instruction is fetched.

The Instruction Register (IR) holds the first byte of each instruction and directs the subsequent operations required to execute each instruction. The IR contents are decoded and used in conjunction with the timing information to control the activation and sequencing of all the other elements on the chip. The Holding Register is used in some multiple-byte instructions to contain further instruction information and partial absolute addresses.

The Arithmetic Logic Unit (ALU) is used to perform all of the data manipulation operations, including Load, Store, Add, Subtract, And, Inclusive Or, Exclusive Or, Compare, Rotate, Increment and Decrement. It contains and controls the Carry bit, the Overflow bit, the Interdigit Carry and the Condition Code Register.

The Register Stack contains six registers that are organized into two banks of three registers each. The Register Select bit picks one of the two banks to be accessed by instructions. In order to accommodate the register-to register instructions, register zero (R0) is outside the array. Thus, register zero is always available along with one set of three registers.

The Address Adder is used to increment the instruction address and to calculate relative and indexed addresses.

The Instruction Address Register holds the address of the next instruction byte to be

accessed. The Operand Address Register stores operand addresses and sometimes contains intermediate results during effective address calculations.

The Return Address Stack (RAS) is a Last In, First Out (LIFO) storage which receives the return address whenever a Branch-to-Subroutine instruction is executed. When a Return instruction is executed, the RAS provides the last return address for the processor's IAR. The stack contains eight levels of storage so that subroutines may be nested up to eight levels deep. The Stack Pointer is a three bit wraparound counter that indicates the next available level in the stack. It always points to the current address.

PROGRAM STATUS WORD

The Program Status Word (PSW) is a major feature of the 2650 which greatly increases its flexibility and processing power. The PSW is a special purpose register within the processor that contains status and control hits

It is divided into two bytes called the Program Status Upper (PSU) and Program Status Lower (PSL). The PSW bits may be tested, loaded, stored, preset, or cleared using the instructions which affect the PSW. The bits are utilized as shown in Table 1.

BYTE	MNEMONIC	FUNCTION
PSU0,1,2	SP	Pointer for the Return Address Stack.
PSU3,4		Not used. These bits are always zero.
PSU5	11	Used to inhibit recognition of additional Interrupts.
PSU6	F	Flag is a latch directly driving the flag output.
PSU7	S	Sense equals the state of the sense input.
PSL0	С	Carry stores any carry from the high-order bit of ALU.
PSL1	СОМ	Compare determines if a logical or arithmetic comparison is to be made.
PSL2	OVF	Overflow is set if a two's complement overflow occurs.
PSL3	WC	With Carry determines if the carry is used in arithmetic and rotate instructions.
PSL4	RS	Register Select identifies which bank of 3 GP registers is being used.
PSL5	IDC	Inter Digit Carry stores the bit-3 to-bit-4 carry in arithmetic operations.
PSL6,7	CC 1	Condition Code is affected by compare, test and arithmetic instructions.

PSU

7	6	5	4	3	2	1	0
S	F	П			SP2	SP1	SP0

S	Sense
F	Flag
- 11	Interrupt Inhibit
SP2	Stack Pointer Two
SP1	Stack Pointer One
SP0	Stack Pointer Zero

PSL

7	6	5	4	3	2	1	0
CC1	CC0	IDC	RS	WC	OVF	СОМ	С

CC1	CCU	IDC	RS	WC	OVF	COM
CC1	Co	nditio	n Co	de On	e	
CC0	Co	nditio	n Co	de Ze	ro	
IDC	Inte	erdigi	t Car	Y		

RS Register Bank Select WC With/Without Carry

OVF Overflow
COM Logical/Arithmetic Compare

C Carry/Borrow

Table 1 PROGRAM STATUS WORD

9

2650A-I • 2650A-1-I

INPUT/OUTPUT INTERFACE

The 2650 series microprocessor has a set of versatile I/O instructions and can perform I/O operations in a variety of ways. One- and two-byte I/O instructions are provided, as well as a special single-bit I/O facility. The I/O modes provided by the 2650 are designated as Data, Control, and Extended I/O.

Data or Control I/O instructions, also called Non-Extended I/O instructions, are one byte long. Any general purpose register can be used as the source or destination. A special control line indicates if either a Data or Control instruction is being executed.

Extended I/O is a two-byte read or write instruction. Execution of an extended I/O instruction will cause a 8-bit address, taken from the second byte of the instruction, to be placed on the low order eight address lines. The data, which can originate or terminate with any general purpose register, is placed on the data bus. This type of I/O can be used to simultaneously select a device and send data to it.

Memory reference instructions that address data outside of physical memory may also

be used for I/O operations. When an instruction is executed, the address may be decoded by the I/O device rather than memory.

MEMORY INTERFACE

The memory interface consists of the address bus, the 8-bit data bus and several signals that operate in an interlocked or handshaking mode.

The Write Pulse signal is designed to be used as a memory strobe signal for any memory type. It has been particularly optimized to be used as the Chip Enable or Read/Write signal.

INTERRUPT HANDLING CAPABILITY

The 2650 series has a single level hardware vectored interrupt capability. When an interrupt occurs, the processor finishes the current instruction and sets the Interrupt Inhibit bit in the PSW. The processor then executes a Branch to Subroutine Relative to location Zero (ZBSR) instruction and sends out Interrupt Acknowledge and Operation Request signals. On receipt of the INTACK

signal the interrupting device inputs an 8-bit address, the interrupt vector, on the data bus. The relative and relative indirect addressing modes combined with this 8-bit address allow interrupt service routines to begin at any addressable memory location.

INSTRUCTION SET

It may be seen from examination of the 2650 instruction set that there are many powerful instructions which are all easily understood and are typical of larger computers. There are one-, two-, and three-byte instructions as a result of the multiplicity of addressing models. See Table 2 for a complete listing and Block Diagram for instruction formats.

Automatic incrementing or decrementing of an index register is available in the arithmetic indexed instructions. All of the branch instructions except indexed branching can be conditional.

Register-to-register instructions are one byte; register-to-storage instructions are two or three bytes long. The two-byte register-to-memory instructions are either immediate or relative addressing types.

_	1			T .	OP	co	DE	1 :-		P	SW E	BITS						r
	MON		DESCRIPTION OF OPERATION		R	or C	C			AF	FEC	TED		,	BYTES	CYCLES	FOR-	NOTE
LOAD/STORE	LOD	ZIRA	Load register zero Load immediate Load relative Load absolute	03 07 0B 0F	06 0A	05	I —	•	IDC	C	OVI	SP	11	F	1 2 2 3	2 2 2 3 4	Z I R A	1 1 1,6 6
LOA	STR	Z R A	Store register zero Store relative Store absolute	CB		C	1 — 9 C8 0 CC	•							1 2 3	2 3 4	Z R A	1 6 6
ETIC	ADD	Z I R A	Add to register zero w/wo carry Add immediate w/wo carry Add relative w/wo carry Add absolute w/wo carry	87 8B 8F	86 8A 8E	85 89 80	88 8C	•	•	•	•				1 2 2 3	2 2 3 4	Z I R A	1 1 1,6 1,6
ARITHMETIC	SUB	Z I R A	Subtract from register zero w/wo borrow Subtract immediate w/wo borrow Subtract relative w/wo borrow Subtract absolute w/wo borrow Decimal adjust register	A7 AB AF	A6 AA AE	AS AS	A0 A4 A8 AC AC	•	•	•	•				1 2 2 3 1	2 2 3 4 3	Z I R A	1 1,6 1,6 1,10
	AND	{Z I R A	AND to register zero AND immediate AND relative AND absolute	43 47 4B	42	41 45 49	 44 48	•							1 2 2 3	2 2 3 4	Z I R	1,10 1 1,6 1,6
LOGICAL	IOR	ZIRA	Inclusive-OR to register zero Inclusive-OR immediate Inclusive-OR relative Inclusive-OR absolute		6A	65 69	64	•							1 2 2 3	2 2 3 4	Z I R A	1 1 1,6 1,6
	EOR	Z I R A	Exclusive-OR to register zero Exclusive-OR immediate Exclusive-OR relative Exclusive-OR absolute		22 26 2A 2E	25 29	24	• • • •				-			1 2 2 3	2 2 3 4	Z I R A	1 1 1,6 1,6
PARE	сом	(z	Compare to register zero arithmetic/logical Compare immediate arithmetic/ logical				E0 E4	•							1 2	2	Z I	3
ROTATE/COMPARE	JU	R	Compare relative arithmetic/ logical Compare absolute arithmetic/		EA EE		E8 EC	•							2	3 4	R A	3,6 3,6
ROT	RRR RRL		logical Rotate register w/wo carry Rotate register left w/wo carry	1	52 D2		50 D0		•	•	•				,1 1	2	z z	1
	DOT	ſR	Branch on condition true	1B	1A	19	18								2	3	R	7,8
	BCT	۱۸	relative Branch on condition true absolute	1F	1E	1D	1C		:						3	3	В	7,8
CH	BCF	${R \choose A}$	Branch on condition false relative Branch on condition false absolute		9A 9E										2	3 3	R B	7
BRANCH	BRN	{ ^R	Branch on register non-zero relative	5B	5A	59	58								2	3	R	7,8
		l _A	Branch on register non-zero absolute	5F	5E	5D	5C								3	3	В	7,8
	BIR	${R \choose A}$	Branch on incrementing register relative Branch on incrementing	1			D8								2	3	R	7,8
		. ^	register absolute		DE	טט	DC		·						3	3	В	7,8

Table 2 INSTRUCTION SET SUMMARY

2650A-I • 2650A-1-I

_	MNE-		1	PC						W BI						FOR-	NOTE
	MONIC	DESCRIPTION OF OPERATION	3	R 01 2	r C(0	CC			OVF		П	F	BYTES	CYCLES	MAT.*	.1012
	(R	Branch on decrementing		FA					_					2	3	R	7,8
(Cont'd)	BDR {	register relative Branch on decrementing register absolute	FF	FE	FD	FC								3	3	В	7,8
BRANCH	ZBRR	Zero branch relative, uncon- ditional	9В		_	_								2	3	ER	6
BRA	вха	Branch indexed absolute, unconditional	9F	_	_	_			-					3	3	EB	5,6
	BST R	Branch to subroutine on con- dition true, relative	3В	3A	39	38					•			2	3	R	7,8
_	ا ۸	Branch to subroutine on con- dition true, absolute	3F	3E	3D	3C					•			3	3	В	7,8
BRANCH/RETURN	BSF (R	Branch to subroutine on con- dition false, relative				B8					•			2	3	R	7
CH/RE	A	Branch to subroutine on con- dition false, absolute				BC					•			3	3	В	7
BAN	BSN (R	Branch to subroutine on non- zero register, relative				78					•	l		3	3	R B	7,8 7,8
INE B	\ A	Branch to subroutine on non- zero register, absolute			/ [7C					•					ER	6
SUBROUTINE	ZBSR	Zero branch to subroutine relative, unconditional	BB		_	_								2	3		
SUB	BSXA	Branch to subroutine, indexed, absolute unconditional	BF											3	3	EB _	5,6
	RET C	Return from subroutine, con- ditional				34					•			1	3	z z	8
	lE.	Return from subroutine and enable interrupt, conditional	\perp				-	ļ			Ľ	<u> </u>					
INPUT/OUTPUT	WRTD REDD	Write data Read data	73	F2 72	71	70								1	2	Z Z	1
Õ	WRTC REDC	Write control Read control	B3	B2	B ⁻ 31									1 1	2 2	Z Z	1
5	WRTE	Write extended				5 D4	1							2	3	1	ŀ
ž	REDE	Read extended	57	56	55	54	•		ļ	<u> </u>	<u> </u>	_	↓	2	3	l .	1
MISC.	HALT	Halt, enter wait state	-	_	-	40 C0								1 1	1	E E	
Ĭ	NOP TMI	No operation Test under mask immediate	F7	F6	F	F4								2	3	ī	4
	LPS {U	Load program status, upper Load program status, lower		92 93							•	•	•	1 1	2 2	E E	9
	SPS {U	Store program status, upper		12			•	!						1	2 2	E	1 1
S	(L	Store program status, lower Clear program status, upper,		13 74			•						•	2	3	EI	'
PROGRAM STATU		masked Clear program status, lower, masked		75	i		•	•	•	•				2	3	ΕI	9
RAM	, (U	Preset program status, upper,		76	6						•	•	•	2	3	Εl	
PROG	PPS {	masked Preset program status, lower, masked		77	,		•	•	•	•				2	3	El	9
	TPS {U	Test program status, upper, masked		B4	1		•							2	3	EI	4
	ا ا	Test program status, lower, masked		B	5		•							2	3	EI	4

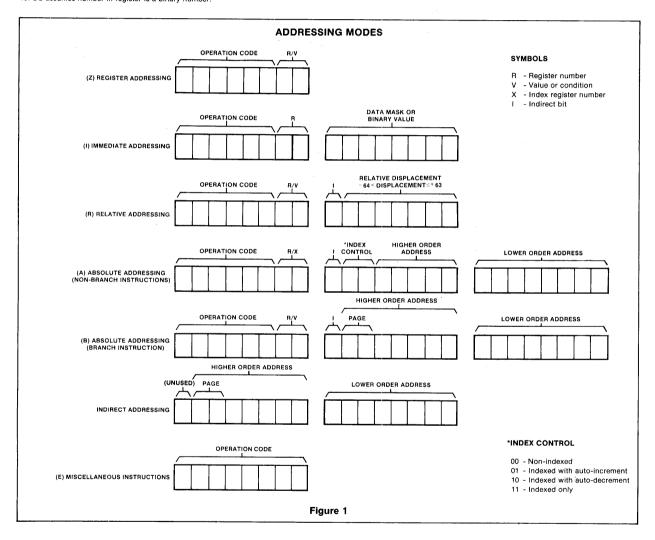
*See Figure 1

Table 2 INSTRUCTION SET SUMMARY (Cont'd)

2650A-I • 2650A-1-I

PRELIMINARY SPECIFICATION

- 1. Condition code (CC1, CC0): 01 if positive, 00 if zero, 10 if negative. 2. Condition code (CC1, CC0): 01 if R0 > r, 00 if R0 = r, 10 if R0 < r.
- 3. Condition code (CC1, CC0): 01 if r > V, 00 if r = V, 10 if r < V.
- 4. Condition code (CC1, CC0): 00 if all selected bits are 1s, 10 if not all the selected bits are 1s
- 5. Index register must be register 3 or 3.
- 6. Requires two additional cycles if indirection is specified.
- 7. Requires two additional cycles if indirection is specified and branch is taken.
- 8. Specify CC = 11 for unconditional branch.
- 9. RS, WC and COM bits in PSW are also affected.
- 10. CC assumes number in register is a binary number.



2650A-I • 2650A-1-I

ABSOLUTE MAXIMUM RATINGS¹

	PARAMETER	RATING	UNIT		
T _A T _{STG} P _D	Operating temperature Storage temperature Package power dissipation ² All input, output, and supply voltages with respect to GND ³	0 to 70 -65 to +150 1.6 5 to +6	۸ ٥٥°		

DC ELECTRICAL CHARACTERISTICS $T_A = 0$ °C to 70°C, $V_{CC} = 5V \pm 5\%$.

				LIMITS		
PARAMETER		TEST CONDITIONS	Min	Тур	Max	TINU
	Current					μΑ
l _{IL}	Input load	$V_{IN} = 0 \text{ to } 5.25V$	1		10	1
Ісон	Output high leakage	ADREN, DBUSEN = $2.2V V_{OUT} = 4V$			10	
LOL	Output low leakage	ADREN, DBUSEN = $2.2V V_{OUT} = 0.45V$			10	
	Voltage levels		İ			V
VIH	Input high		2.2		Vcc	
VIL	Input low		-0.5		0.8	
Voh	Output high	$I_{OH} = -100\mu A$	2.4			
Vol	Output low	I _{OL} = 1.6ma	0.0		0.45	
lcc	Power supply current	$V_{CC} = 5.25V T_A = 0^{\circ}C$	1.		150	mA
	Capacitance					pf
CIN	Input	$V_{IN} = 0V$	ļ		10	
Соит	Output	$V_{OUT} = 0V$	1	1	10	

NOTES

- Stresses above those listed under "Absolute Maximum Ratings" may cause
 permanent damage to the device. This is a stress rating only and functional operation
 of the device at these or at any other condition above those indicated in the operation
 sections of this specification is not implied.
- For operating at elevated temperatures the device must be derated based on +150°C maximum junction temperature and thermal resistance of 50°C/W junction to ambient (40 pin IW package).
- This product includes circuitry specifically designed for the protection of its internal devices from the damaging effects of excessive static charge. However, it is suggested that conventional precautions be taken to avoid applying any voltages larger than the rated maxima.
- 4. Parameters valid over operating temperature range unless otherwise specified.
- 5. All voltage measurements are referenced to ground.

PRELIMINARY SPECIFICATION: Manufacturer reserves the right to make design and process changes and improvements

2650A-I • 2650A-1-I

AC ELECTRICAL CHARACTERISTICS $T_A = 0$ °C to +70°C, $V_{CC} = +5V \pm 5\%$.

	PARAMETER		2650A			2650A-1	1	UNIT	
		Min	Тур	Max	Min	Тур	Max	CIVIT	
Тсн	Clock high phase	400			250			ns	
T _{CL}	Clock low phase	400			250			ns	
TCP	Clock period	800			500		1	ns	
TPC	Processor cycle time ^{5,7}	2400			1500			ns	
TOR	OPREQ pulse width7	2Tc++		2TcH+	2T _{CH} +		2T _{CH} +	ns	
TCOR	Clock to OPREQ time	T _{CL} -100		T _{CL} +50	T _{CL} - 100		T _{CL} + 50 200	ns	
TAS	Address stable	50			50			ns	
TAD	Address delay	50	50	1				ns	
Tcs	Control signal stable	50			50			ns	
TDIS	Data in setup	0			0			ns	
T _{DIH}	Data in hold	10		1	10			ns	
T_{DD}	Data out delay	50		1	50			ns	
T_{DS}	Data out stable	50			50			ns	
TOAS	OPACK setup time	100			100		†	ns	
TOAH	OPACK hold time	150			150		1	ns	
TWPD	Write pulse delay	100		450	100		300	ns	
TWPW	Write pulse width7	T _{CL} -100		TCL	T _{CL} - 100		TCL	ns	
TIRS	INTREQ setup time			150			150		
TIRH	INTREQ hold time	0			0				
T_{ABD}	Address bus tri-state delay			180			180	ns	
τ_{DBD}	Data bus tri-state delay			150	1		150	ns	

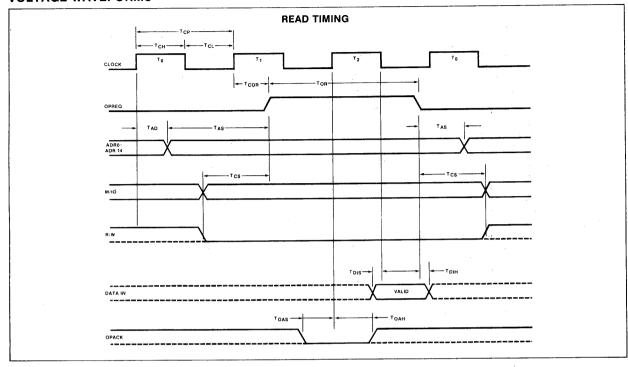
NOTES

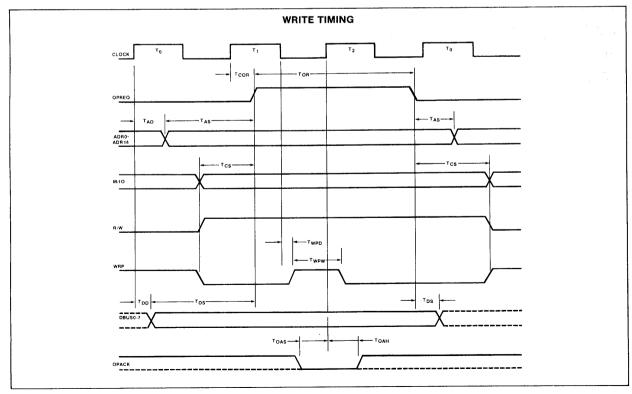
- 1. Input levels swing between 0.80 and 2.2 volts.
- 2. Input signal transition times are 20ns.
- 3. Timing reference level is 1.5 volts.
- Output load is -100μA at 100pF and 1 TTL load.
- 5. Processor cycles time consists of three clock periods.
- Output buffer rise time is 150ns maximum.
- These values assume that OPACK is returned in time to not cause the processor to idle. Otherwise, the specified maximum will increase by an integral number of clock cycles.

PRELIMINARY SPECIFICATION: Manufacturer reserves the right to make design and process changes and improvements

2650A-I • 2650A-1-I

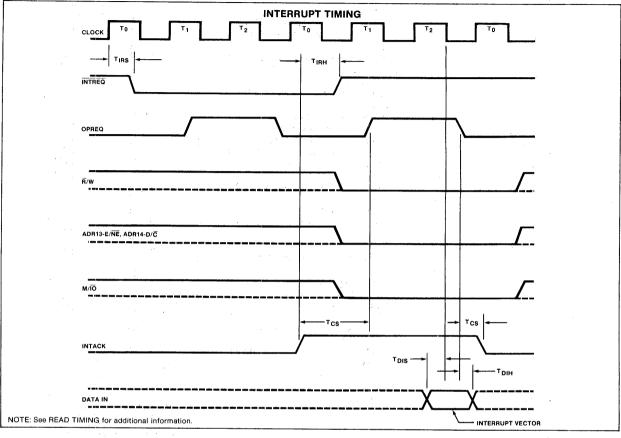
VOLTAGE WAVEFORMS

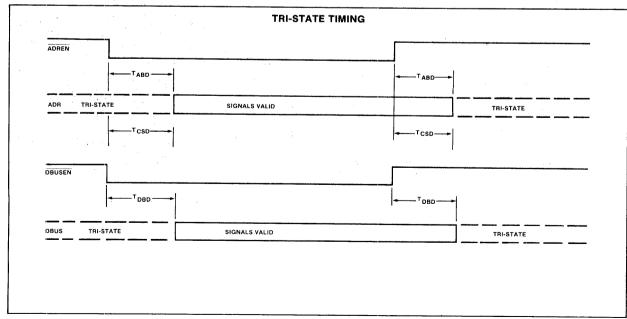




2650A-I • 2650A-1-I

VOLTAGE WAVEFORMS (Cont'd)





DESCRIPTION

SC/MP (Simple Cost-effective MicroProcessor) is a single-chip 8-bit microprocessor packaged in a standard, 40-pin, dual-in-line package.

N-channel, silicon gate, depletion mode standard-process technology ensures high performance, high reliability, and high producibility.

SC/MP is intended for use in generalpurpose applications where cost per function is a most significant criterion. But cost efficiency is only a part of SC/MP's story. It goes on to include a variety of useful functions that are not even provided by some of the expensive microprocessors, like selfcontained timing circuitry, 16-bit (65K) addressing capability, serial or parallel datatransfer capability and common memory/peripheral instructions. The builtin features in conjunction with the low initial cost describe what SC/MP really is-a microprocessor specifically designed to provide the simplest and most efficient solution to many application requirements.

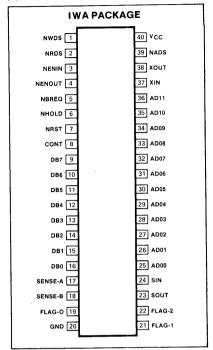
APPLICATIONS

- Test systems and instrumentation
- Machine tool control
- Small business machines
- Word processing systems
- · Educational systems
- Multiprocessor systemsProcess controllers
- Terminals
- Traffic controls
- · Laboratory controllers
- Sophisticated games
- Automotive

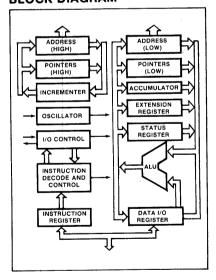
FEATURES

- Simpler interfacing
 - Bidirectional Tri-state 8-bit data bus
 - TTL-compatible input/output interface
- Si-gate N-channel ion-implant process
- Direct Memory Access (DMA) and multiprocessor capabilities
 - Handshake bus-access control on chip
- · Simplified programming
 - Multiple addressing modes program-counter-relative, immediate data, indexed, auto-indexed, and implied
- Direct control output
- Three user-accessible control-flag outputs
- Simpler I/O hardware
 - Separate serial-data input and output ports
 - Two sense inputs
 - Direct interfacing to standard memory parts
- Simplified timing hardware
- On-chip clock generator
- Interface flexibility
 - Capability to interface with memories or peripherals of any speed
- Large system capability
- Address capability to 65K bytes of memory
- · Simplified power requirements
 - Single 5-volt supply
 - Low power

PIN CONFIGURATION



BLOCK DIAGRAM



ABSOLUTE MAXIMUM RATINGS*

	PARAMETER	RATING	UNIT
T _A T _{STG}	Voltage at any pin Operating temperature range Storage temperature range Lead temperature (solder,10 sec)	-0.5 to +7.0 0 to +70 -65 to +150 +300	°°°°°°°°°°°°°°°°°°°°°°°°°°°°°°°°°°°°°°

*NOTE

Maximum ratings indicate limits beyond which damage may occur. Continuous operation at these limits is not intended and should be limited to those conditions specified under electrical characteristics.

DC ELECTRICAL CHARACTERISTICS $T_A = 0$ °C to +70°C, $V_{CC} = +5V \pm 5\%$

PARAMETER	TEST COMPLETIONS		LIMITS		
FARAWETER	TEST CONDITIONS	Min	Тур	Max	UNIT
INPUT SPECIFICATIONS All input pins except V _{CC} and GND Logic "1" input	· E	2.0		Vcc	V
Logic "0" input		-0.5		0.8	V
Capacitance (All pins except V _{CC} and GND)				10	pF
Supply current lcc	T _A = 25°C outputs unloaded T _A = 0°C outputs unloaded			45 50	mA mA
OUTPUT SPECIFICATIONS TRI-STATE pins (NWDS, NRDS, DB0-DB7, AD00-AD11) Logic "1" output	$I_{OUT} = -100 \mu A$	2.4			v
Logic "0" output	I _{OUT} = 2.0mA			0.4	V
NADS, FLAG 0-2, SOUT, NENOUT Logic "1" output Logic "1" output	$I_{OUT} = -100 \mu A$ $I_{OUT} = -1 mA$	Vcc - 1 1.5			v v
Logic "0" output	lout = 2.0mA			0.4	V
NBREQ1 Logic "0" output	I _{OUT} = 2.0mA			0.4	V
Logic "1" output	0 ≤ V _{OUT} ≤ V _{CC}			±10	μA
XOUT Logic "1" output Logic "0" output	$I_{OUT} = -100\mu A$ $I_{OUT} = 1.6mA$	2.4		0.4	V

FUNCTIONAL DESCRIPTION

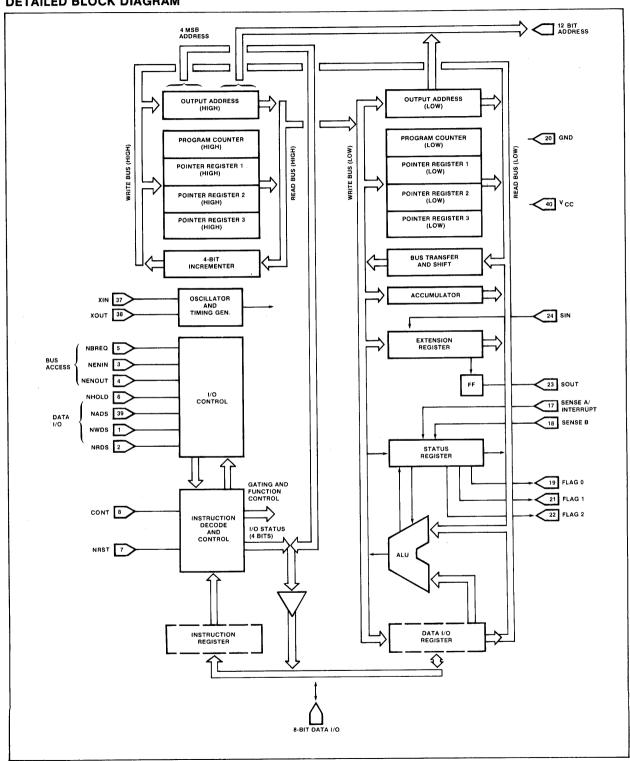
SC/MP is a self-contained general-purpose microprocessor designed for ease of implementation in stand-alone, DMA (Direct Memory Access), and multiprocessor applications. Communications between SC/MP and external memory/peripheral devices are effected via a 12-bit dedicated address bus and an 8-bit bidirectional data bus. During the address interval of each input/output cycle, SC/MP employs both buses to provide a 16-bit address output: the 12 least

significant address bits are sent out over the 12-bit address bus and the 4 most significant address bits are sent out over the 8-bit data bus along with 4 status bits. Separate strobe outputs from SC/MP (NADS, NWDS, NRDS) indicate when valid address information is present on the two buses, and when valid input/output memory or peripheral data are present on the 8-bit bus. To further extend flexibility of application, serial data input/output ports are also provided

so that serial data transfers can be effected under program control. The remaining input/output signals shown in the Detailed Block Diagram are dedicated to general-purpose control and status functions, including initialization, bus management, microprocessor halt, interrupt request, input/output cycle extension, and user-specified hardware/software interface functions. A detailed description of each input/output signal is provided in Table 1.

NBREQ is an input/output signal that requires an external resistor to V_{CC}.

DETAILED BLOCK DIAGRAM



AC ELECTRICAL CHARACTERISTICS $T_A = 0$ °C to +70°C, $V_{CC} = +5V \pm 5\%$, 1 TTL Load.^{1, 3}

PARAMETER	TEST CONDITIONS		LIMITS		
ANAMETER	TEST CONDITIONS	Min	Тур	Max	UNI
fx	$R = 240\Omega \pm 5\%$ (Figure 2B) $C = 300pF \pm 10\%$	0.1 2.0		4.0 4.0	MH MH
T _C ² Microcycle		500 1			ns μs
External clock input Two Tw1 XOUT/ADS timing relationship T _H (ADS)	See Figure 2A See Figure 3	120 120 100		225	ns
Address and input/output status TD1(ADS) TW(ADS) TS(ADDR) TH(ADDR) TS(STAT) TH(STAT) TH(NBREQ)	See Figures 5 and 6	(T _C /2) - 50 (T _C /2) - 165 50 (T _C /2) - 150 50		3T _C /2	ns
Data input cycle TD(RDS) TW(RDS) TS(RD) TH(RD) TACC(RD)	See Figure 5	0 T _C + 50 175 0 2T _C - 200			ns
Data output cycle TD(WDS) TW(WDS) TS(WD) TH(WD)	See Figure 6	T _C - 50 T _C - 75 (T _C /2) - 200			ns
Input/output cycle extend T _S (HOLD) T _{D1} (HOLD) T _{D2} (HOLD) T _W (HOLD) T _H (HOLD)	See Figure 7	200 130		275 350 ∞	ns
Bus access T _D (NENOUT) T _{D2} (ADS) T _H (NENIN)	See Figure 4	T _C /2		150 3T _C /2	ns
Output load capacitance XOUT All other output pins				30 75	pF

NOTES

^{1.} All times measured from valid Logic "0" level = 0.8V or valid Logic "1" level = 2.0V.

^{2.} T_C is the time period for two-clock cycles of the on-chip or external oscillator ($T_C = 2/f_x$). Refer to paragraph titled Timing Control for detailed definition.

^{3.} All times measured with a 50% duty cycle on the external clock.

PIN DESIGNATIONS*

SIGNAL MNEMONIC	FUNCTIONAL NAME	DESCRIPTION
NRST .	Reset input	Set high for normal operation. When set low, aborts in-process opera- tions. When returned high, internal control circuit zeroes all programmer-accessible registers; then, first instruction is fetched from memory location 0001 ₁₆ .
CONT	Continue input	When set high, enables normal execution of program stored in external memory. When set low, SC/MP operation is suspended (after completion of current instruction) without loss of internal status.
NBREQ	Bus request input/output	Associated with SC/MP internal allocation logic for system bus. Can be used as bus request output or bus busy input. Requires external load resistor to Vcc.
NENIN	Enable input	Associated with SC/MP internal allocation logic for system bus. When set low, SC/MP is granted access to system buses. When set high, places system buses in high-impedance (Tri-state) mode.
NENOUT	Enable output	Associated with SC/MP internal allocation logic for system bus. Set low when NENIN is low and SC/MP is not using system buses (NBREQhigh). Set high at all other times.
NADS	Address strobe output	Active-Low strobe. While low, indicates that valid address and status output are present on system buses.
NRDS	Read strobe output	Active-Low strobe. On trailing edge, data are input to SC/MP from 8-bit bidirectional data bus. High-impedance (Tri-state) output when input/output cycle not in progress.
NWDS	Write strobe output	Active-Low strobe. On trailing edge, data are input to SC/MP from 8-bit bidirectional data bus. High-impedance (Tri-state) output when input/output cycle not in progress.
NHOLD	Input/ouput cycle extend input	When set low prior to trailing edge of NRDS or NWDS strobe, stretches strobe to extend input/output cycle; that is, strobe is held low until NHOLD signal is returned high.
SENSE A	Sense/interrupt request input	Serves as interrupt request input when SC/MP internal IE (Interrupt Enable) flag is set. When IE flag is reset, serves as user-designated sense condition input. Sense condition testing is effected by copying status register to accumulator.
SENSE B	Sense input	User-designated sense-condition input. Sense-condition testing is effected by copying status register to accumulator.
SIN	Serial input to E register	Under software control, data on this line are right-shifted into E registe by execution of SIO instruction.
SOUT	Serial output from E register	Under software control, data are right-shifted onto this line from E register by execution of SIO instruction. Each data bit remains latched until execution of next SIO instruction.
FLAGS 0, 1, 2	Flag outputs	User-designated general-purpose flag outputs of status register. Unde program control, flags can be set and reset by copying accumulator to status register.
AD00- AD11	Address bit 00 through address bit 11	Twelve Tri-state address output lines. SC/MP outputs 12 least significant address bits on this bus when NADS strobe is low. Address bits are then held valid until trailing edge of read (NRDS) or write. (NWDS strobes. After trailing edge of NRDS or NWDS strobe, bus is set to high impedance (Tri-state) mode until next NADS strobe.

2

PIN DESIGNATIONS (Cont'd)

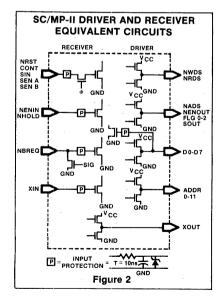
MNEMONIC	FUNCTIONAL NAME Output at NADS Time 2.3.4	DESCRIPTION
DB0	Address Bit 12	Fourth most significant bit of 16-bit address.
DB1	Address Bit 13	Third most significant bit of 16-bit address.
DB2	Address Bit 14	Second most significant bit of 16-bit address
DB3	Address Bit 15	Most significant bit of 16-bit address.
DB4	R-Flag	When high, data input cycle is starting; when low, data output cycle is starting.
DB5	I-Flag	When high, first byte of instruction is being fetched.
DB6	D-Flag	When high, indicates delay cycle is starting; that is, second byte of DLY instruction is being fetched.
DB7	H-Flag	When high, indicates that Halt instruction has been executed. (In some system configurations, the H-Flag output is latched and, in conjunction with the Continue input, provides a programmed Halt.)

NOTES

- The 8-bit bidirectional data bus is set to the high-impedance (tri-state) mode except
 when it is actually in use by SC/MP (NADS, NRDS, or NWDS low). During the
 addressing interval of each input/output cycle (NADS low). SC/MP provides address
 and status outputs over the bus. During the ensuing data-transfer interval (NRDS or
 NWDS low), 8-bit input or output data bytes are routed over the bus.
- The DB0 through DB7 (AD12-HFLG) lines are a high-impedance (open circuit) load when SC/MP does not have access to the input/output bus.
- Input at NRDS time: Input data are expected on the eight (DB0-DB7) lines.
- 4. Output at NWDS time: Output data are valid on the eight (DB0-DB7) lines.

DRIVERS AND RECEIVERS

Equivalent circuits for SC/MP drivers and receivers are shown in Figure 2. All inputs have static charge protection circuits consisting of an RC filter and voltage clamp. These devices still should be handled with care, as the protection circuits can be destroyed by excessive static charge.



TIMING CONTROL

All necessary timing signals are provided by a three-stage inverter ring oscillator contained on the SC/MP chip. Two control pins, XIN and XOUT, permit the frequency of the oscillator to be controlled by any of the following methods:

- By leaving the XOUT pin unterminated and driving the XIN pin with an externally generated TTL clock that conforms to the parameters shown in Figure 3A. For this method, the frequency of the oscillator is equalt to the frequency of the external clock input.
- By connecting a resistor-capacitor feedback network between the XIN and XOUT pins and GND as shown in Figure 3B.
- 3. By connecting a crystal with low-pass filter network between the XIN and XOUT pins and GND as shown in Figure 3C (for above 1MHz or Figure 3D (for 1MHz or below). For this method, the frequency of the oscillator is equal to the resonant frequency of the crystal and the low-pass filter prevents unwanted harmonic oscillations.

In addition to illustrating appropriate frequency-control networks for the on-chip oscillator, Figures 3A through 3D also show how an optional driver may be used to derive a system clock from the oscillator signal present at the XOUT pin. For reference purposes, the timing relationship between the XOUT signal and the NADS strobe is shown in Figure 5.

In the discussions that follow, instruction execution and input/output timing are described in terms of microcycles.

The time interval of a microcycle is four times the period of the oscillator, that is:

period of one microcycle = 2Tc

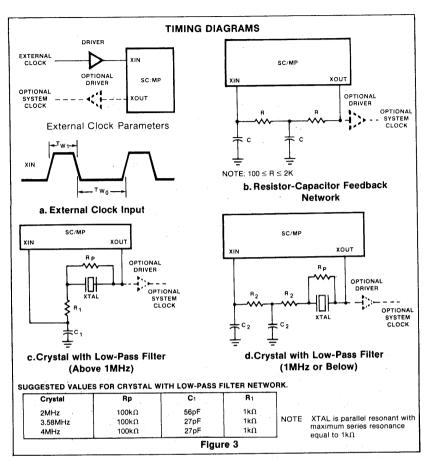
$$T_C = 2\left(\frac{1}{f_{OSC}}\right) = 2\left(\frac{1}{f_{res}}\right) = 2\left(\frac{1}{f_{XIN}}\right)$$

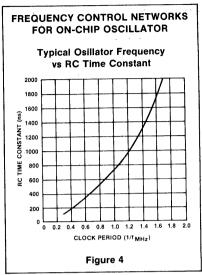
Where:

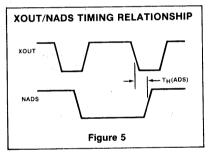
T_C = time period for two cycles of on-chip or external oscillator

f_{osc} = frequency of on-chip oscillator f_{res} = resonant frequency of crystal connected between XIN and XOUT pins

fxin = frequency of external clock applied to XIN pin







INSTRUCTION FORMAT

The SC/MP instruction repertoire includes both single-byte and double-byte instructions. A single-byte instruction consists of an 8-bit operation code that specifies an operation that SC/MP can execute without further reference to memory. A double-byte instruction consists of an 8-bit operation code and an 8-bit data or displacement field. When the second byte represents a data field, the data are processed by SC/MP during execution of the instruction, thereby eliminating the need for further memory references. When the second byte represents a displacement value, it is used to calculate a memory address that will be accessed (written into or read from) during execution of the instruction (refer to Addressina).

DATA STORAGE

As shown in the Detailed Block Diagram, SC/MP provides ten internal registers, seven of which are accessible to the programmer. The purpose and function of these registers are described below.

Program Counter—The program counter is a 16-bit register that contains the address of the instruction being executed. The contents of this register are automatically incremented by one just before each instruction is fetched from memory to enable sequential execution of the stored instructions. Under program control, the contents of this register also may be modified or exchanged with the contents of a pointer register to effect subroutine calls and program branches.

NOTE

The 16-bit address output of the program counter consists of a 4-bit high-order address and a 12-bit low-order address. When the program counter is incremented at the start of each instruction fetch input/output cycle, only the 12 low-order bits are affected; no carry is provided to the 4 high-order bits. For systems employing memories of 4K or less, the high-order bits can be ignored as they are set to 0000.6 following initialization. For systems employing larger memories, the contents of a pointer register can be modified to select the desired 4K block of memory.

Pointer Registers—The pointer registers are 16-bit general-purpose registers that normally are loaded under program control with reference addresses that serve as page

pointers, stack pointers, and subroutine pointers. In applications having minimal memory addressing requirements, these registers may be used alternately as data storage registers.

NOTE

When interrupt requests are enabled, pointer register 3 is automatically referenced by the internal microprogram for formation of the starting address of the user-generated interrupt service routine. (See Figure 10.) In this case, the contents of pointer register 3 must be set to one less than the memory location of the first instruction in the interrupt service routine.

Accumulator—The 8-bit accumulator (AC) is the primary working register of SC/MP. It is used for performing and storing the results of arithmetic and logic operations as well as for data transfers, shifts, rotates, and data exchanges with the program counter, the pointer registers, and the status register.

Extension Register—The extension register is used both for serial input/output data transfers and with the accumulator to effect arithmetic, logic, and data-transfer operations. If the second byte of an indexed or auto-indexed memory reference instruction

(refer to Addressing) equals -128₁₀, the contents of the extension register are used as the displacement value for address formation.

Status Register—The status register provides storage for arithmetic, control, and software status flags. For more detailed information on the function of this register, refer to Status Register under the description of the Arithmetic and Logic Unit.

Instruction Register—The 8-bit instruction register is not accessible to the programmer. During the fetch phase of each instruction cycle, this register is loaded with the 8-bit instruction operation code retrieved from memory (for a single-byte instruction or the first byte of a double-byte instruction).

Data Input/Output Register—The data input/output register is not accessible to the programmer. It is used for temporary storage of all input/output data received via or transmitted over the 8-bit bidirectional data bus during the data-transfer interval of each input/output cycle (NRDS or NWDS low).

Address Register—The 16-bit address register is not accessible to the programmer. It is used for temporary storage of the 16-bit address transmitted during an input/output cycle.

ARITHMETIC AND LOGIC UNIT

The Arithmetic and Logic Unit (ALU) provides the data-manipulation capability that is an essential feature of any microprocessor. The operations provided by the ALU include OR, XOR, increment, decrement, binary addition, and decimal addition. For decimal addition, the data inputs to the ALU are treated as two 4-bit BCD digits, thereby eliminating the program-storage and execution time required to perform BCD to binary conversion.

BUS TRANSFER LOGIC

The bus transfer logic processes the gating and function control outputs of the instruction-decode logic to provide the shift-right (with link, without link, or with serial input data), rotate (with or without link), and bus-exchange functions necessary for data movement between the SC/MP internal read and write buses. A general summary of the data-manipulation capabilities available to the programmer follows.

1. Either the low-order or the high-order byte of any pointer register can be exchanged with the contents of the 8-bit accumulator. Thus, data exchanges between the pointer registers can be effected one byte at a time via the accumulator.

- 2. The contents of the program counter can be directly exchanged with the contents of any pointer register.
- 3. The contents of the extension register can be loaded into the accumulator or can be exchanged with the contents of the accumulator. When the accumulator is loaded from the extension register, the original contents of the accumulator are lost.
- 4. The contents of the status register can be copied into the accumulator to enable status modification or conditional branch testing. When the status register is copied into the accumulator, the contents of the status register are not altered but the original contents of the accumulator are lost.
- 5. The contents of the accumulator can be copied into the status register to change the outputs of the status register, except for status bits 4 and 5 (Sense A and B inputs to SC/MP). Since these are read-only bits, they are not affected by data movements internal to SC/MP. Copying the accumulator into the status register does not alter the contents of the accumulator.

NOTE

The flag 0, 1, and 2 outputs of the status register serve as latched flags, in other words, they are set to the specified state when the contents of the accumulator are copied into the status register, and they remain in the specified state until the contents of the status register are modified again under program control.

STATUS REGISTER

The function of each bit in the status register is described briefly below.

7	6	5	4	3	2	1	0
CY/L	ΟV	SB	SA	IE	F ₂	F ₁	F ₀

User Flag 0—User-assigned general-purpose status bit for implementation as software status bit or in system control applications. This status bit is available as an external output from SC/MP.

User Flag 1-Same as User Flag 0.

User Flag 2-Same as User Flag 0.

Interrupt Enable Flag—Internal status bit that is set and reset under program control. When set, SC/MP recognizes external interrupt requests received via Sense A input. When reset, inhibits SC/MP from recognizing interrupt requests.

Sense A—General purpose status input for sensing external conditions. When IE flag is reset, this bit can be tested by copying status register to accumulator. When IE flag is set, this bit serves as interrupt request input causing SC/MP to automatically branch to user-generated interrupt-service routine in response to high input.

Sense B—Same as Sense A except that it is not tested for interrupt status.

NOTE

Sense A and B inputs are read-only bits. Thus, they are not affected when the contents of the accumulator are copied into the status register.

Overflow (OV)—This bit is set if an arithmetic overflow occurs during an add (ADD, ADI, or ADE) or a complement-and-add instruction (CAD, CAI, or CAE). It is not affected by the decimal-add instructions (DAD, DAI, or DAE).

Carry/Link (CY/L)—This bit is set if a carry from the most significant bit occurs during an add, complement-and-add, or decimal-add instruction. Thus, it serves as a carry input to the next add instruction. In addition, it is included in the Shift Right with Link (SRL) and Rotate Right with Link (RRL) instructions.

CONTROL

The operation of the SC/MP microprocessor consists of repeatedly accessing or fetching instructions from the program stored in external memory and executing the operations specified by the instructions. These two steps are carried out under the control of an internal microprogram. (SC/MP is not user-microprogrammable.) The microprogram is similar to a state table specifying the series of states of system control signals necessary to carry out each instruction. Microprogram storage is provided in the instruction decode and control logic, and microprogram routines are implemented to fetch and execute instructions. The fetch routine first increments the program counter, and then causes the instruction address to be transferred from the program counter to the system buses via the output address register. The microprogram next initiates an input data transfer. When the instruction operation code is subsequently placed on the 8-bit data bus (singlebyte instruction or first byte of double-byte instruction), the operation code is loaded into the instruction register. The operation code is then partially decoded to determine whether the instruction contains a second byte. If it does, a second input data transfer is effected to load the next byte in the data input/output register.

After the complete instruction is stored in the instruction and/or data input/output register(s), the instruction decoder transforms the instruction operation code into the address of the appropriate instruction-execution routine contained in the internal microprogram. The microprogram then branches to the specified internal address to initiate execution of the instruction. The

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resulting execution routine comprises one or more microinstructions that implement the required functions. For example, the first microcycle of an Extension Register Add Instruction (ADE) causes the contents of the extension register to be gated onto the read bus, transferred to the write bus via the bus control logic, and then written into the data input/output register. The next microcycle causes the contents of the accumulator to be gated onto the read bus, the contents of the read bus to be added to the contents of the data input/output register via the ALU, and the resultant output of the ALU to be written into the accumulator via the write bus. The final step of the execution routine is a jump back to the fetch routine to access the next instruction.

INITIALIZATION

Since SC/MP may power up in a random condition, the following power-up and initialization procedure is recommended.

1. Apply power (GND and V_{CC}) and set NRST low.

NOTE

Allow ample time (typically, 250ms) for the oscillator and the internal clocks to stabilize. In systems where NRST is

set low after turning on power, NRST must remain low for a minimum of 4Tc. While NRST is low, any in-process operations are aborted automatically. When NRST is low, strobes and address and data buses are in the Non-I/O state (high-Z state).

2. Set NRST high. If the rise time of this input is too slow, the processor, first, will initialize and execute a few instructions and, then, will reinitialize. If the application is such that multiple initialization is undesireable, NRST should be brought high at a minimum rate of 2 volts per microcycle.

NOTE

This causes the SC/MP internal control circuit to set the contents of all programmer-accessible registers to zero. Thus, when SC/MP is granted access to the system buses following initialization, the first instruction is fetched always from memory location 0001 fs. The NBREQ output goes low, indicating the start of this input/output cycle; this occurs at a time within 13Tc after NRST is set high. Normal execution of the program continues as long as NRST remains high.

PARALLEL DATA TRANSFERS

Parallel data transfers occur during each instruction fetch and during the ensuing read/write cycle associated with execution of the memory-reference instructions. This class of instruction could perhaps more properly be called the "Input/Output Refer-

ence Class" in the case of the SC/MP microprocessor, since all data tranfers, whether with memory, peripheral devices, or a cenrral processor data bus, occur through the execution of these instructions. This unified bus structure is in contrast with many other microprocessors and minicomputers that have one instruction type (input/output class) for communication with peripheral devices and another instruction type (memory reference class) for communication with memories. The advantage of the approach taken by SC/MP is that a wider variety of instructions (the entire memory-reference class) is available for communications with peripherals. Thus, the LD and ST (Load and Store) instructions can be used for basic transfers, the ILD and DLD (increment/decrement and load) instructions can be used for indexing peripheral registers, and the remaining memory reference instructions can be used, as required, for "one-step" retrieval and processing of peripheral input data

BUS UTILIZATION

The bus utilization of SC/MP is shown in Table 2.

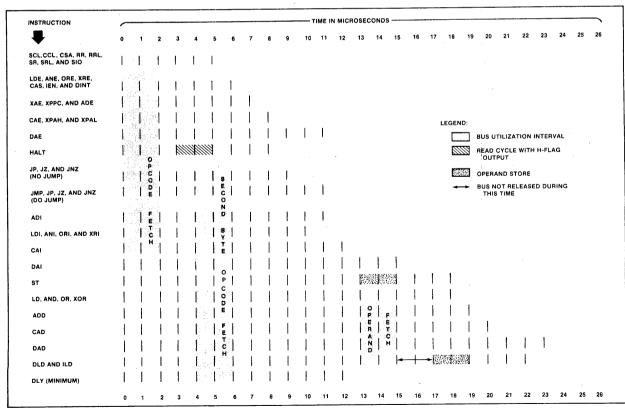


Table 2 BUS UTILIZATION OF EACH INSTRUCTION

NBREQ, NENIN, and NENOUT are active and bus access is controlled as shown in Figure 5. If NENIN is returned high during an input/output cycle, the input/output cycle is repeated when NENIN is again returned low.

During an ILD or DLD instruction, SC/MP does not relinquish the bus between the loading of the data and the storing of the modified data. If NENIN is brought high after the data have been loaded, the load portion of the cycle is not repeated when NENIN is returned low.

BUS ACCESS

Before SC/MP can initiate parallel data transfers with memory or peripheral devices, it must have access to the system address and data buses. Three of the SC/MP input/output signals are associated with bus control: NBREQ, NENIN, and NENOUT. For simple stand-alone applications, the NENOUT signal can be ignored and the NENIN signal can be tied to GND to allow the SC/MP microprocessor to have continual access to the system buses. The NBREQ input/output line then goes low during each input/output cycle as shown in Figures 6 and 7 to indicate when SC/MP is actually using the system buses.

NOTE

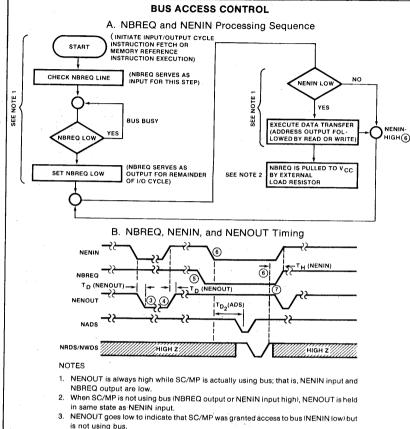
The NBREQ input/output line must be tied to VCC via an external load resistor to allow normal operation of the SC/MP microprocessor

For DMA and multiprocessor applications, the NBREQ, NENIN, and NENOUT signals can be interconnected in various configurations to allow bus access to be granted to requesting devices according to userspecified priorities. Figure 5 illustrates the general sequence in which these signals are processed by SC/MP to gain access to the system buses and to indicate when the buses are actually being used.

INPUT/OUTPUT CYCLE

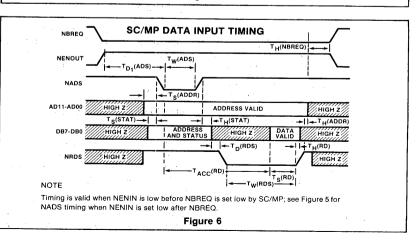
Once SC/MP has control of the system buses, the actual input/output cycle begins. As shown in Figures 6 and 7, the functions of memory addressing, data reading, and data writing are implemented, respectively, by the address strobe (NADS), the read strobe (NRDS), and the write strobe (NWDS). Note that the NBREQ signal is reset high at the end of the input/output cycle to indicate that the system buses are now free for use by the highest-priority requesting device.

The first operation that SC/MP performs for each input/output cycle is to load the 12 least significant address bits onto the 12-bit address bus, and the 4 most significant



- is not using bus
- NENOUT goes high in response to high NENIN input.
- SC/MP generates bus request; bus access not granted because NENIN high.
- NENIN goes low. Bus access now granted and input/output cycle actually initiated. If NENIN is set high while SC/MP has access to the bus, the address and data ports will go to the high-impedance (Tri-State) state, but NBREQ will remain low. When NENIN is subsequently set low, the input/output cycle will begin again.
- Input/output cycle completed. NENOUT goes low to indicate that SC/MP granted access to bus but not using bus. If NENIN had been set high before completion of input/output cycle. NENOUT would have remained high

Figure 5



address bits along with 4 status bits onto the 8-bit data bus. At the same time, SC/MP sets the NADS output low to indicate that the address and the status information are valid. The low-order address on the 12-bit bus is then held valid for the duration of the input/output cycle; the high-order address and the status information on the 8-bit bus remain valid only while NADS is low. While valid, the status bits have the following significance:

RFLG—When high, indicates that input/output cycle is read cycle; when low, indicates that input/output cycle is write cycle.

IFLG—Set high to indicate that instruction operation code (single-byte instruction or first byte of double-byte instruction) will be output from memory following NADS.

DFLG—Set high only when second byte of Delay Instruction is to be read from memory following NADS. Execution of the Delay Instruction then starts at trailing edge of NRDS. Upon completion, SC/MP provides NADS output to initiate next input/output cycle if bus access is granted. Time in microcycles from leading edge of delay flag to leading edge of subsequent NADS output is computed from the following formula:

Delay = [9 + 2(AC) + 2 disp + 29 disp]microcycles

where:

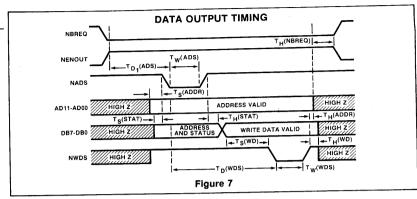
(AC) = unsigned contents of accumulator disp = unsigned displacement value contained in second byte of Delay Instruction

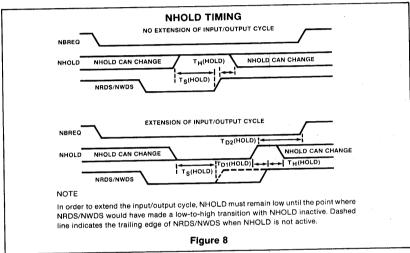
The time derived from the above formula does not include the four microcycles required to fetch the first byte of the Delay Instruction. Thus, when the Delay Instruction is used for software timing, total instruction execution time equals [13 + 2(AC) + 2 disp + 29 disp] microcycles.

NOTE

When Halt Instruction is executed, instruction decode and control logic inhibits incrementing of program counter for one input/output cycle. Thus, Halt Instruction is read from memory a second time to enable generation of HFLG output, but no further processing of Halt Instruction occurs. In effect, this procedure ensures HFLG is output in advance of the next instruction to be fetched from memory.

HFLG—Set high only during addressing interval of read cycle that follows Halt Instruction. HFLG may be used to cause user-provided external logic to set the CONT input low, and thereby to effect a programmed halt. Since HFLG read cycle precedes the next instruction fetch, termination of programmed halt enables fetch of first instruction that follows Halt Instruction.





After resetting the NADS output, SC/MP generates an NRDS or NWDS strobe, respectively, to initiate a data-input (read) or data-output (write) operation. For a read operation, input data are strobed into SC/MP from the 8-bit bus on the trailing edge of the NRDS strobe. For a write operation, SC/MP places valid output data on the 8-bit bus on the leading edge of the NWDS strobe. After resetting the NRDS or NWDS strobe to complete the data transfer, SC/MP then resets the NBREQ signal to indicate that the system buses are free for use by another controller.

INPUT/OUTPUT CYCLE EXTENSION

As shown in Figure 8, the NHOLD signal may be set low prior to the trailing edge of the NRDS or NWDS strobe to cause SC/MP to lengthen the input/output cycle by holding the strobe active until after the NHOLD signal is returned high. Since there is no restriction on the maximum duration of the NHOLD signal, it can be used in a variety of

applications ranging from accommodation of memories/peripherals with long access times to single-cycle control of the operating program for software debug purposes.

Figure 9 illustrates a typical circuit that may be used to generate an NHOLD signal of repeatable duration. The circuit shown employs an N74165 8-Bit Parallel In/Serial Out Shift Register to allow selection of an input/ output cycle extend time that ranges from T_C/2 to 2T_C in increments of T_C/2. Functional operation of the circuit is controlled by the NADS strobe and XOUT signals. Each time that the NADS strobe goes low, the data present at the A through H terminals are loaded into the shift register in parallel. When the NADS strobe subsequently returns high, the data are then shifted out serially on the positive-to-negative transitions of XOUT. Thus, the NHOLD output of the circuit is set low on the leading edge of each NADS strobe and, as shown in the chart that accompanies the circuit diagram, it remains low for a time period ranging from three clock cycles minimum (B, C,

D, and E inputs = Logic "1") to seven clock cycles maximum (B, C, D, and E inputs = Logic "0").

It is important to note that instruction execution time is increased whenever an input/ output cycle is extended via the NHOLD signal. For purposes of computing the increase in instruction execution time, it is necessary to distinguish between the terms Input/Output Cycle Delay Period and Input/ Output Cycle Extend Time. The term Input/Output Cycle Delay Period refers to the time that the NRDS/NWDS strobe is actually "stretched" to provide the required memory or peripheral access time. The term Input/Output Cycle Extend Time refers to the additional number of microcycles required by the internal SC/MP microprogram to complete the extended input/output cycle: that is:

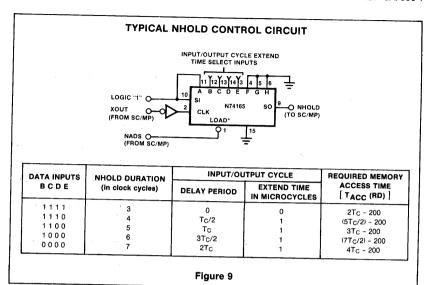
INPUT/OUTPUT CYCLE					
Delay Period	Extend Time				
T _C /2 through 2T _C					
(> 0 ≤ 1μcycle)	1 μcycle				
5T _C /2 through 4T _C	1				
(> 1 ≤ 2μcycles)	2μcycles				
9T _C /2 through 6T _C					
(> 2 ≤ 3µcycles)	3μcycles				
etc.	etc.				

The total increase in instruction execution time, therefore, is equal to the Input/Output Cycle Extend Time multiplied by the total number of input/output cycles associated with the instruction. For example, a DLD Instruction is normally executed in 22 microcycles. Since this instruction employs three read input/output cycles and one write input/output cycle, an Input/Output Cycle Extend Time of one microcycle would increase total DLD Instruction execution time to 26 microcycles.

SERIAL DATA TRANSFERS

Serial input/output data transfers can be used efficiently with very slow input/output peripherals such as X-Y plotters, teletypewriters, slow-speed printers, and so forth. Such transfers can be effected in any of the following manners:

1. By assigning serial input/output functions to the extension register via the SIO (Serial Input/Output) Instruction. When this instruction is executed, the contents of the extension register are shifted right one bit. At the same time, data present on the SIN line are shifted into bit position 7 of the extension register and the original contents of bit position 0 are shifted into a flip-flop to provide a latched output of the SOUT line. The SOUT data are then held latched until the next SIO instruction is executed.



- 2. By using one of the status flags as an output data bit and one of the sense lines as an input data bit.
- 3. By implementing external logic such that only one line of the 8-bit data input/output bus is used.

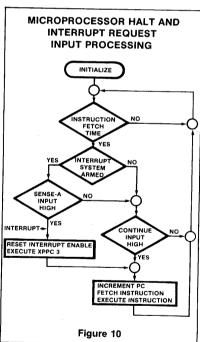
For synchronous systems, serial data input/output timing may be provided by program loops that employ the delay instruction, or by using one or more of the transfer instructions (see Table 2) to test the output of an external timing circuit. For asynchronous systems, one of the sense inputs can be used for testing bit-received/ready status and a pulsed flag output can be provided, under program control, for peripheral indexing each time that a data bit is actually shifted in or out.

Systems that have several input/output devices must be multiplexed, device selection can then be accomplished using the status flag outputs of SC/MP, or by using parallel input/output commands to load an external latch. Systems that do not require serial input/output capability can employ the SIN and SOUT lines as a sense input and flag output, respectively.

INTERRUPTS

When the internal interrupt enable (IE) flag is set under program control, the Sense A line is enabled to serve as an interrupt request input; when the IE flag is reset, SC/MP is inhibited from detecting interrupts. Thus, while the IE flag is set, the Sense A input is tested prior to the fetch phase of each instruction as shown in Figure 10. Upon detection of an interrupt re-

quest (Sense A high), the following events occur automatically.



1. The status register IE flag is reset to prevent SC/MP from responding to any further interrupt requests. Interrupt request capability can then be reenabled during or at the end of the ensuing user-generated interrupt service routine via the IEN (Enable Interrupt) Instruction or by copying the accumulator into the status register.

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- 2. The contents of the program counter are exchanged with the contents of the pointer register 3.
- 3. The contents of the program counter are incremented by one to address the first instruction of the user-generated interrupt service routine.

The interrupt system must be armed before interrupts are enabled. This is accomplished as follows:

- 1. First, the Interrupt Enable Bit in the Status Register is set true by executing either an Enable Interrupt Instruction (IEN) or a Copy Accumulator to Status Register Instruction (CAS).
- 2. Second, one additional instruction is fetched and executed.

A return from interrupt is accomplished by executing two instructions: Enable Interrupt (IEN) immediately followed by Ex-

change Pointer 3 with Program Counter (XPPC 3).

MICROPROCESSOR HALT

The CONT input to SC/MP is provided to enable suspension of operation without loss of internal status. Processing of the CONT input is shown in Figure 9. Since this is an asynchronous input, it can be controlled by external timing logic, or as stated previously, the HALT flag output that appears on the 8-bit data bus (during the read cycle that follows execution of a Halt Instruction) can be used with an external circuit to effect a programmed halt condition. Note that when an interrupt request is detected while the CONT input is low, the first instruction of the user-generated interrupt service routine is automatically executed. Thus, the first instruction of the interrupt service routine can be used to reset the external CONT input logic and, thereby, to terminate the microprocessor halt condition if so desired.

After execution of an instruction, the CONT input must be high for a minimum time of 2T_C (1 microcycle) in order to fetch and execute the next instruction.

INSTRUCTION SET

The SC/MP instruction set provides the general-purpose user of microprocessors a powerful programming capability along with above-average flexibility and speed. The instruction set consists of 46 instructions, which comprise eight general categories. A listing of the complete instruction set is provided in Table 2; typical instruction execution times are given in Table 3, and notations and symbols used as shorthand expressions of instruction capability are defined in Table 4.

MNEMONIC	DESCRIPTION	OBJECT FORMAT	OPERATION	MICRO- CYCLES
LD ST AND OR XOR DAD ADD CAD	Memory reference instructions Load Store AND OR Exclusive-OR Decimal add Add Complement and add	7 6 5 4 3 2 1 0 7 6 5 4 3 2 1 0 1 1 0 0 0 m ptr disp 1 1 0 0 1 1 1 1 1 1 0 0 1 1 1 1 1 1 1	(AC)←(EA) ((EA)←(AC) (AC)←(AC) Λ (EA) (AC)←(AC) V (EA) (AC)←(AC) Ψ (EA) (AC)←(AC)+(EA)+(CY/L); (CY/L) (AC)←(AC)+(EA)+(CY/L); (CY/L), (OV) (AC)←(AC)+(EA)+((CY/L); (CY/L), (OV)	18 18 18 18 18 23 19
ILD DLD LDI ANI ORI XRI	Memory increment/ decrement instructions Increment and load Decrement and load Immediate instructions Load immediate AND immediate OR immediate Exclusive-OR immediate	7 6 5 4 3 2 10 7 6 5 4 3 2 1 0 10 1 0 10 ptr disp	(AC), (EA)(EA) + 1 (AC), (EA)(EA) - 1 (AC)data (AC)(AC) Λ data (AC)(AC) V data (AC)(AC) V data	22 22 10 10 10
DAI ADI CAI	Decimal add immediate Add immediate Complement and add immediate	11101100 11110100 11111100	(AC)←(AC) ₁₀ +data ₁₀ +(CY/L); (CY/L) (AC)←(AC)+data+(CY/L); (CY/L), (OV) (AC)←(AC) + ~data + (CY/L); (CY/L), (OV)	15 11 12
JMP JP JZ JNZ	Transfer instructions Jump Jump if positive Jump if zero Jump if not zero	7 6 5 4 3 2 1 0 7 6 5 4 3 2 1 0 1 0 0 1 0 0 ptr disp 1 0 0 1 1 0 0 1 1 0 0 1 1 1 1 0 0 1 1 1 1	(PC)-EA If (AC) ≥ 0, (PC)-EA If (AC) = 0, (PC)-EA If (AC) ≠ 0, (PC)-EA	11 9, 11 9, 11 9, 11
DLY	Double-byte miscellaneous instructions Delay	76543210 76543210 10001111 disp	count AC to -1. delay = 13 + 2(AC) + 2 disp + 29 disp microcycles	13 to 131, 59

Table 3 SC/MP INSTRUCTION SUMMARY

MNEMONIC	DESCRIPTION	OBJECT FORMAT	OPERATION	MICRO-
	Extension register instructions	76543210		
LDE	Load AC from extension	0100000	(AC)←(E)	6
XAE	Exchange AC and extension	0000001	(AC)←(E)	7
ANE	AND extension	01010000	(AC)←(AC) Λ (E)	6
ORE	OR extension	01011000	(AC)←(AC) V (E)	6
XRE	Exclusive-OR extension	01100000	(AC)←(AC) V (E)	6
DAE	Decimal add extension	01101000	(AC)-(AC) ₁₀ +(E) ₁₀ +(CY/L);(CY/L)	11
ADE	Add extension	01110000	(AC)-(AC)+(E)+(CY/L);(CY/L),(OV)	7
CAE	Complement and add extension	01111000	(AC) \leftarrow (AC) $+$ $+$ (AC) $+$ $+$ $+$ $+$ $+$ $+$ $+$ $+$ $+$ $+$	8
6	Pointer register move instructions	76543210		
XPAL	Exchange pointer low	0 0 1 1 0 0 ptr	(AC)←(PTR _{7:0})	
XPAH	Exchange pointer high	001101	(AC)-(PTR _{15:8})	8
XPPC	Exchange pointer with PC	001111	(PC)-(PTR)	8 7
		[001111]	(PC)=(PTH)	/
	Shift, rotate, serial I/O			
010	instructions	76543210		
SIO	Serial input/output	00011001	$(E_i) \rightarrow (E_{i-1}), SIN \rightarrow (E_7), (E_0) \rightarrow SOUT$	5
SR	Shift right	00011100	$(AC_i) \rightarrow (AC_{i-1}), 0 \rightarrow (AC_7)$	5
SRL	Shift right with link	00011101	$(AC_i) \rightarrow (AC_{i-1}), (CY/L) \rightarrow (AC_7)$	5
RR	Rotate right	00011110	$(AC_i) \rightarrow (AC_{i-1}), (AC_0) \rightarrow (AC_7)$	5
RRL	Rotate right with link	00011111	$(AC_i)\rightarrow (AC_{i-1}), (AC_0)\rightarrow (CY/L)\rightarrow (AC_7)$	5
	Single-byte miscellaneous			
	instructions	76543210		
HALT	Halt	0000000	Pulse H-flag	8
CCL	Clear carry/link	0000010	(CY/L)-0	5
SCL	Set carry/link	00000011	(CY/L)-1	5
DINT	Disable interrupt	00000100	(IE)←0	6
IEN	Enable interrupt	00000101	(IE)←1	6
CSA	Copy status to AC	00000110	(AC)←(SR)	5
CAS	Copy AC to status	00000111	(SR)(AC)	6
NOP	No operation	00001000	None	5

Table 3 SC/MP INSTRUCTION SUMMARY (Cont'd)

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INSTRUCTION	READ CYCLES	WRITE CYCLES	TOTAL MICROCYCLES
ADD	3	0	19
ADE	1	0	7
ADI	2	0	11
AND	3	0	18
ANE	1	0	6
ANI	2	0	10
CAD	3	0	20
CAE	1	0	8
CAI	2 .	0	12
CAS	1	0	6
CCL	1	0	5
CSA	1	0	5
DAD	3	0	23
DAE	1	0	11
DAI	2	0	15
DINT	1	0	6
DLD	3	1	22
DLY	2	0	13 - 131593
HALT	2	0	8
IEN	1	0	6
ILD	3	1	22
JMP	2	0	11
JNZ	2	0	9,11 for Jump
JP	2	0	9,11 for Jump 9,11 for Jump
JZ	2	0	
LD	3	0	18 6
LDE	1	0	10
LDI	2	0	5
NOP	1	0	18
OR	3	0	6
ORE	1	0	10
ORI	2	0	5
RR	1	0	5
RRL	1	0	5
SCL	1	0	5
SIO	1 1		5
SR	1	0	5
SRL	2	1	18
ST	1	Ö	7
XAE	3	0	18
XOR	1	0	8
XPAH	1	0	8
XPAL	1 1	0	7
XPPC		0	6
XRE XRI	2	0	10

NOTE

If slow memory is being used, the appropriate delay should be added for each read or write cycle.

Table 4 INSTRUCTION EXECUTION TIME

ADDRESSING

During execution, instructions and data defined in a program are stored into and loaded from specific memory locations, the accumulator, or selected registers. Because SC/MP, memory (read/write and read-only),

and peripherals are on a common data bus, any instruction used to address memory may be used to address the peripherals. The formats of the instruction groups that reference memory are shown below.

2	1,0		7 , 0
m	ptr		disp
,	ptr		disp
	m	4	mptr

Memory Reference Instructions Memory Increment/ Decrement Instructions and Transfer Instructions

Memory-reference instructions use the PC-relative, indexed, or auto-indexed methods of addressing memory. The memory-increment/decrement instructions and the transfer instructions use the PC-relative or indexed methods of addressing.

The various methods of addressing memory and peripherals are shown below.

Immediate addressing is an addressing format specific to the immediate instruction group.

TYPE OF	0	PERAND	FORMATS
ADDRESSING	m	ptr	disp
PC-relative Indexed Immediate	0 0 1	0 1,2, or 3 0	-128 to +127 -128 to +127 -128 to +127
Auto- indexed	1	1,2, or 3	-128 to +127

For PC-relative, indexed, and auto-indexed memory-reference instructions, another feature of the addressing architecture is that the contents of the extension register are substituted for the displacement if the instruction displacement equals –128 (–X'80).

NOTE

All arithmetic operations associated with address formation affect only the 12 low-order address bits; no carry is provided to the 4 high-order bits. For systems employing memories of 4K or less, the high-order bits can be ignored as they are set to 0000 following initialization. For systems employing larger memories, the high-order bits must be set to the starting address of the desired 4K block of memory. For example:

0001₂ enables memoory locations 1000₁₆ - 1FFF₁₆ to be addressed.

 0010_2 enables memory locations 2000_{16} - $2FFF_{16}$ to be addressed and so forth.

PC-Relative Addressing—A PC-relative address is formed by adding the displacement value specified in the operand field of the instruction to the current contents of the program counter. The displacement is an 8-bit twos-complement number, so the range of the PC-relative addressing format is -128₁₀ to +127₁₀ locations from the current contents of the program counter.

Immediate Addressing—Immediate addressing uses the value in the second byte of a double-byte instruction as the operand for the operation to be performed (see below)

For example, compare a Load (LD) instruction to a Load Immediate (LDI) instruction. The Load instruction uses the contents of

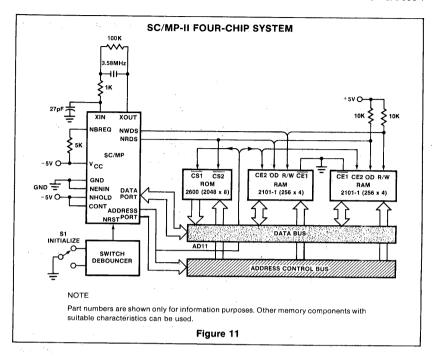
the second byte of the instruction in computing the effective address of the data to be loaded. The Load Immediate instruction uses the contents of the second byte as the data to be loaded.

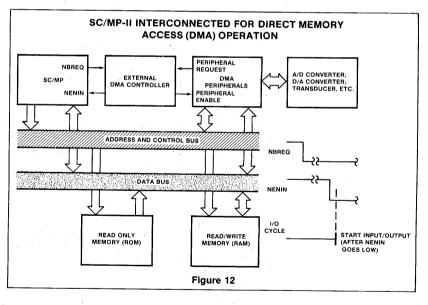
Indexed Addressing—Indexed addressing enables the programmer to address any location in memory through the use of the pointer register and the displacement. When indexed addressing is specified in an instruction, the contents of the designated pointer register are added to the displacement to form the effective address. The contents of the pointer register are not modified by indexed addressing.

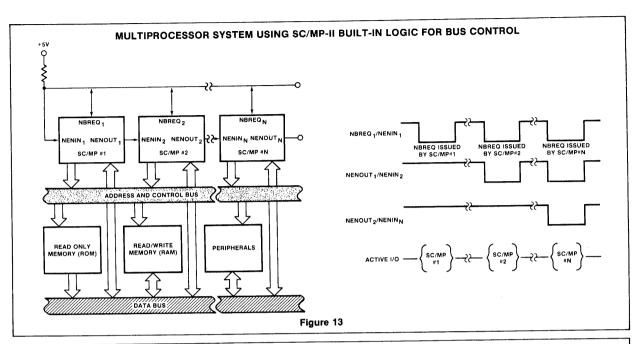
Auto-Indexed Addressing—Auto-indexed addressing provides the same capabilities as indexed addressing along with the ability to increment or decrement the designated pointer register by the value of the displacement. If the displacement is less than zero, the pointer register is decremented by the displacement before the contents of the effective address are fetched or stored. If the displacement is equal to or greater than zero, the pointer register is used as the effective address, and the pointer register is incremented by the displacement after the contents of the effective address are fetched or stored.

SYSTEM IMPLEMENTATION

Figures 11 and through 13 illustrate typical SC/MP system configurations. In Figure 10, SC/MP is shown interconnected to three memory devices to form a stand-alone 4-device system that provides 256 words of read/write memory and 2,048 words for program storage. Figure 12 shows SC/MP interconnected to an external controller for Direct Memory Access (DMA) operation, and Figure 13 illustrates a multiprocessor application using SC/MP's built-in logic to control bus access.







SYMBOL AND NOTATION	MEANING
AC	8-bit Accumulator.
CY/L	Carry/Link Flag in the Status Register.
data	Signed, 8-bit immediate data field.
disp	Displacement; represents an operand in a nonmemory reference instruction or an address modifier field in a memory reference instruction. It is a signed twos-complement number.
EA	Effective Address as specified by the instruction.
E	Extension Register; provides for temporary storage, variable displacements and separate serial input/output port.
i	Unspecified bit of a register.
IE	Interrupt Enable Flag.
m	Mode bit, used in memory reference instructions. Blank parameter sets $m = 0$, at sets $m = 1$.
l ov	Overflow Flag in the Status Register.
PC	Program Counter (Pointer Register 0); during address formation, PC points to the last byte of the instruction being executed.
ptr	Pointer Register (ptr = 0 through 3). The register specified in byte 1 of the instruction.
ptr _{n:m}	Pointer register bits; n:m = 7 through 0 or 15 through 8.
SIN	Serial Input pin.
SOUT	Serial Output pin.
SR	8-bit Status Register.
()	Means "contents of." For example, (EA) is contents of Effective Address.
[]	Means optional field in the assembler instruction format.
~	Ones complement of value to right of ~.
-	Means "replaces."
-	Means "is replaced by."
+	Means "exchange."
@	When used in the operand field of the instruction, sets the mode bit (m) to 1 for auto-incrementing/auto-decrementing indexing.
10+	Modulo 10 addition.
Λ	AND operation.
V	Inclusive-OR operation.
₩	Exclusive-OR operation.
≥	Greater than or equal to.
=	Equals.
≠	Does not equal.

Table 5 GLOSSARY

MP8080A-I

DESCRIPTION

The MP8080A is an 8-bit microprocessor housed in a standard, 40-pin dual-in-line package. The chip, which is fabricated using N-channel silicon gate MOS technology, functions as the central processing unit (CPU) in Signetics' 8080 microcomputer family.

The MP8080A has a 16-bit address bus that is capable of addressing up to 65k bytes of memory and up to 256 input and 256 output devices. Data is routed to and from the MP8080A on a separate bidirectional 8-bit bus. This data bus is also Tri-State, making direct memory addressing (DMA) and multi-processing applications possible. The MP8080A directly provides signals to control the interface to memory and I/O ports. All buses, including control, are TTL compatible.

An asynchronous interrupt capability is included in the MP8080A to allow external signals to change the instruction sequence. The interrupting device may vector the program to a particular service routine location (or some other direct function) by specifying an interrupt instruction to be executed.

FEATURES

- 2μ s instruction cycle
- Variable length instructions
- General purpose registers—six plus an accumulator
- Direct addressing up to 65k bytes
- Variable length stack accessed by 16-bit stack pointer
- Addresses 256 input and 256 output ports
- Provisions for vectored interrupts
- Tri-state bus for DMA and multiprocessing capability
- Tri-state TTL drive capabilities for address and data buses
- Decimal arithmetic capability
- Multiple addressing modes
 Direct

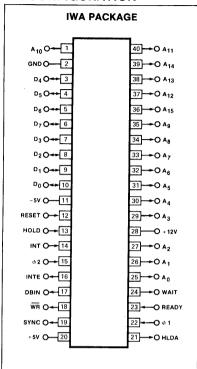
Register

Register indirect

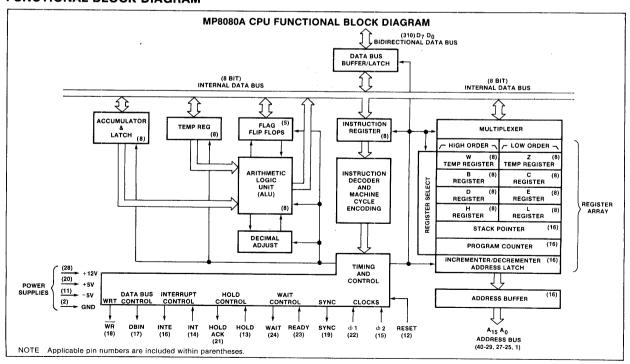
Immediate

 Direct plug-in replacement for Intel MP8080A

PIN CONFIGURATION



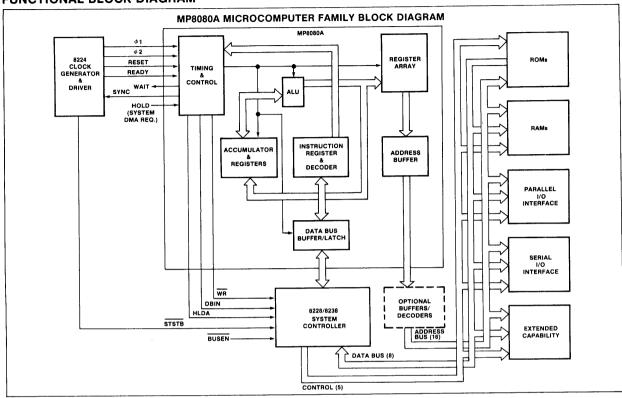
FUNCTIONAL BLOCK DIAGRAM



G

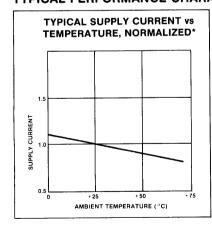
MP8080A-I

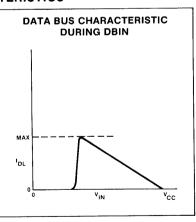
FUNCTIONAL BLOCK DIAGRAM



*NOTE \[\Delta 1 \] supply/\(\Delta T_A = 0.45\%/\) \(^{\text{C}} \)

TYPICAL PERFORMANCE CHARACTERISTICS





PIN DESIGNATION

MNEMONIC	PIN NO.	TYPE	NAME AND FUNCTION
	<u> </u>		
READY	23	,	Ready: When high (logic 1), indicates that valid memory or input data are available to the CPU on the MP8080A data bus. The Ready signal is used to synchronize the CPU with slower memory or input/output devices. If the MP8080A does not receive a high Ready input after sending out an address to memory or an input/output device, the MP8080A enters a Wait mode for as long as the Ready input remains low (logic 0). The CPU may also be single stepped by the use of the Ready signal.
HOLD	13	l	Hold: When high, requests that the CPU enter the Hold mode. When the CPU is in th Hold mode, the CPU address and the data buses both will be in the high-impedanc state. The Hold mode allows an external device to gain control of the MP8080A addres and data buses immediately following the completion of the current machine cycle b the CPU. The CPU acknowledges the Hold mode via the HLDA output line. The Hold request is recognized under the following conditions: 1. The CPU is in the Halt mode. 2. The Ready signal is active and the CPU is in the t2 or tw state.
INT	14	1	Interrupt Request: When high, the CPU recognizes an interrupt request on this line afte completing the current instruction or while in the Halt mode. An interrupt request is no honored if the CPU is in the Hold mode (HLDA = logic 1) or the INTE Flip-flop is rese (logic 0).
RESET	12	1	Reset: When activated (high) for a minimum of three clock periods, the content of the Program Counter is cleared and the Interrupt Enable and HLDA Flip-flops are reset Following a Reset, program execution starts at memory location 0. It should be noted that the status flags, accumulator, stack pointer, and registers are not cleared during the Reset sequence.
$\phi_{1,}\phi_{2}$	22,15	I	φ ₁ and φ ₂ Clocks: Two non-TTL compatible clock phases which provide nonoverlap ping timing references for internal storage elements and logic circuits of the CPU.
SYNC	19	0	Sychronizing Signal: When activated (high), the beginning of a new machine cycle i indicated and the status word is outputted on the Data Bus.
A ₁₅ -A ₀	25-27, 29-40	0	Address Bus: This bus comprises sixteen tri-state output lines. The bus provides the address to memory (up to 65k bytes) or denotes the input/output device number for up to 256 input and 256 output peripherals.
WAIT	24	0	Wait: When high, acknowledges that the CPU is in the Wait mode.
WR	18	0	Write: When low, the data on the data bus are stable for Write memory or outpu operation.
HLDA	21	Ö	 Hold Acknowledge: Goes high in response to a logic 1 on the Hold line and indicates that the data and address bus will go to the high-impedance state. The HLDA begins at one of the following times: 1. The ts state of the Read memory input operation. 2. The clock period following the ts state of a Write memory output operation. In both cases, the HLDA signal starts after the rising edge of the φ1 clock, and high impedance occurs after the rising edge of the φ2 clock.
INTE	16	0	Interrupt Enable: Indicates the content of the internal INTE Flip-flop. The Enable and Disable Interrupt (EI and DI) Instructions cause the INTE Flip-flop to be set and reset respectively. When the Flip-flop is reset (INTE = logic 0), it inhibits interrupts from being accepted by the CPU. In addition, the INTE Flip-flop is automatically reset (thereby disabling further interrupts) at the t ₁ state of the instruction fetch cycle, when are interrupt is accepted; it is also reset by the Reset Signal.
D ₇ -D ₀	3-10	I/O	Data Bus: This bus comprises eight tri-state input/output lines. The bus provided bidirectional communication between the CPU, memory, and input/output devices for instructions and data transfers. A status word (which describes the current machine cycle) is also outputted on the data bus during the first state of each machine cycle (SYNC = logic 1).
+12V	28	ı	+2 Volts: V _{DD} Supply.
+5V	20	1	+5 Volts: V _{CC} Supply.
-5V	11	1	-5 Volts: VBB Supply.
GND	2	1	Ground: VSS (0 volt) reference.

MP8080A-I

ABSOLUTE MAXIMUM RATINGS

PARAMETER	RATING	UNIT
Operating temperature range Storage temperature range	0 to 70 -65 to +150	°C
All input or output voltages with respect to V _{BB} V _{DD} and V _{SS} with respect to V _{BB}	-0.3 to +20 -0.3 to 20	V V W
	Operating temperature range Storage temperature range All input or output roltages with respect to VBB VDD and VSS with respect to VBB	Operating temperature range 0 to 70 Storage temperature range -65 to +150 All input or output voltages with respect to VBB -0.3 to +20

NOTE Maximum ratings indicate limits beyond which permanent damage may occur. Continuous operation at these limits is not intended and should be limited to those conditions specified under dc electrical characteristics.

DC ELECTRICAL CHARACTERISTICS $T_A = 0$ °C to +70°C, $V_{DD} = +12$ V \pm 5%, $V_{CC} = \pm5$ V \pm 5%, $V_{BB} = -5V \pm 5\% \ V_{SS} = 0V$, unless otherwise specified.

				LIMITS					
	PARAMETER	TEST CONDITIONS	Min	Тур	Max	UNIT			
VIHC VILC VIH VIL	Input voltage Clock high Clock low High Low		9.0 V _{SS} -1 3.3 V _{SS} -1		V _{DD} +1 V _{SS} +0.8 V _{CC} +1 V _{SS} +0.8	V			
V _{OH} V _{OL}	Output voltage High Low	$I_{OH} = 150 \mu A$ $I_{OL} = 1.9 \text{mA}$ on all inputs	3.7		0.45	V			
IDD(AV) ICC(AV) IBB(AV)	Supply current Avg. (V _{DD}) Avg. (V _{CC}) Avg. (V _{BB})	Operation $t_{CY} = 0.48 \mu s$ Operation $t_{CY} = 0.48 \mu s$		40 60 0.01	70 80 1	mA			
lil Icl	Leakage current Input Clock	Vss ≤ Vin ≤ Vcc Vss ≤ Vclock ≤ Vdd			±10 ±10	μΑ			
IDL*	Data bus (in input mode)	$V_{SS} \le V_{IN} \le V_{SS} + 0.8V$ $V_{SS} \le +0.8V \le V_{IN} \le V_{CC}$			-100 -2.0	μΑ			
lfL	Address and data bus (during hold)	V _{ADDR/DATA} = V _{CC} V _{ADDR/DATA} = V _{SS} +0.45V			+10 -100	μΑ			
Cφ Cin Cout	Capacitance Clock Input Output	$(T_A=25^{\circ}C,V_{CC}=V_{DD}=V_{SS}=0V,$		17 6 10	25 10 20	pF			

*NOTE

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When DBIN is high and $V_{\text{IN}} > V_{\text{IH}}$ an internal active pull up will be switched onto the Data

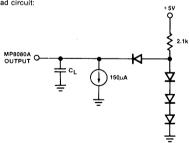
signetics

AC ELECTRICAL CHARACTERISTICS (Cont'd) $T_A = 0$ °C to +70°C, $V_{DD} = +12V \pm 5\%$, $V_{CC} = \pm 5V \pm 5\%$,

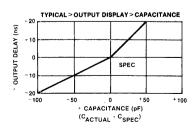
 $V_{BB} = -5V \pm 5\% \ V_{SS} = 0V$, unless otherwise specified

	DADAMETED			LIMITS		
	PARAMETER	TEST CONDITIONS	Min	Тур	Max	UNIT
tcy3	Clock period		0.48		2.0	μS
$t_{r,tf}$	Clock rise and fall time		0		50	ns .
t_{ϕ} 1	φ ₁ Pulse width		60			l ns
tφ2	φ ₂ Pulse width	•	220			ns
t _{D1}	Delay φ1 to φ2					ns
t _{D2}	Delay φ 2 to φ 1		70			ns
t _{D3}	Delay φ₁ to φ₂ leading edges		80			ns
t_{DA}^2	Address output delay from ϕ_2	$C_L = 100pF$	1		200	ns
t_{DD}^2	Data output delay from φ ₂	$C_L = 100pF$	1 1		220	ns
t _{DC} ²	Signal output delay from φ ₁ or φ ₂ (SYNC, WR, WAIT, HLDA)	$C_L = 50pF$			120	ns
t _{DF} 2	DBIN delay from φ ₂	C _L = 50pF	25		140	ns
t _{DI} 1	Delay for input bus to enter input mode				toF	ns
tps1	Data setup time during ϕ_1 and DBIN		30		'0'	ns
t _{DS2}	Data setup time to φ ₂ during DBIN		150		1	ns
t _{DH} 1	Data hold time from \$\dot\delta_2\$ during DBIN		1			ns
t _{IE} 2	INTE output delay from ϕ_2	$C_L = 50pF$	1 1		200	ns
t _{RS}	Ready setup time during φ ₂		120		1 200	ns
ths	Hold setup time to ϕ_2		140			ns
tis	INT setup time during ϕ_2		120			ns
	(during ϕ_1 in halt mode)	the state of the s				
tH	Hold time from ϕ_2 (READY, INT, HOLD)					ns
tFD	Delay to float during hold	C _L = 100pF: Address, data			120	ns
	(address and data bus)		l i			
t _{AW} 2	Address stable prior to WR ⁵	•	(Note 5)		(Note 5)	ns
t _{DW} 2	Output data stable prior to WR 6	$C_L = 50pF$: \overline{WR} , HLDA, DBIN	(Note 6)		(Note 6)	ns
twp2	Output data stable from WR 7		(Note 7)		(Note 7)	ns
tw _A 2	Address stable from WR ⁷		(Note 7)		(Note 7)	ns
t _{HF} 2	HLDA to float delay®		(Note 8)		(Note 8)	ns
twF2	WR to float delay ⁹		(Note 9)		(Note 9)	ns
t _{AH} 2	Address hold time after DBIN during HLDA		-20			ns

- 1. Data input should be enabled with DBIN status. No bus conflict can then occur and data hold time is assured. tpH = 50ns or tpF, whichever is less.
- 2. Typical load circuit:



3. $t_{CY} = t_{D3} + t_{r}\phi_{2} + t_{!}\phi_{2} + t_{D2} + t_{f}\phi_{2} + t_{r}\phi_{1} \geq$ 480ns.

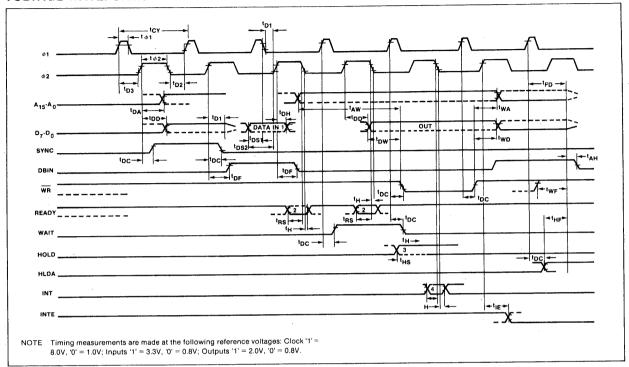


- 4. The following are relevant when interfacing the MP8080A to devices having $V_{IH} = 3.3V$:
- a) Maximum output rise time from 0.8V to 3.3V = 100ns @ C_L = SPEC. b) Output Delay when measured to 3.0V = SPEC + 60ns @ C_L = SPEC.
- c) If $C_L \neq SPEC$, add 0.6ns/pF if $C_L > C_{SPEC}$, subtract 0.3ns/pF (from modified delay) if
- C₁ < C₅PEC. add 0.6ns/P.H.L₂ > C₅PEC, subtract 0.3ns/pF from mod C₁ < C₅PEC. 5. taw = $2t_{CY} t_{D3} t_{r}\phi_2 140ns$. 6. $t_{DW} = t_{CY} t_{D3} t_{r}\phi_2 170ns$. 7. If not HLDA, $t_{WD} = t_{WA} = t_{D3} + t_{r}\phi_2 + 10ns$. If HLDA, $t_{WD} = t_{WA} = t_{WF}$. 8. $t_{HF} = t_{D3} + t_{r}\phi_2 50ns$.

- 9. $t_{WF} = t_{D3} + t_r \phi_2 10 ns$.

MP8080A-I

VOLTAGE WAVEFORMS



SYMBOLS	DATA BUS BIT	DEFINITION
OUT	D4	Indicates that the address bus contains the address of an output device and the data bus will contain the output data when WR is active.
M ₁	D ₅	Provides a signal to indicate that the CPU is in the fetch cycle of the first byte of an instruction.
INP*	D ₆	Indicates that the address bus contains the address of an input device and the input data should be placed on the data bus when DBIN is active.
MEMR*	D ₇	Designates that the data bus will be used for memory read data.
INTA*	D ₀	Acknowledge signal for Interrupt request. Signal should be used to gate a restart instruction onto the data bus when DBIN is active.
WO	D ₁	Indicates that the operation in the current machine cycle will be a Write memory or Output function (WO = 0). Otherwise, a Read memory or Input operation will be executed.
STACK	D ₂	Indicates that the address bus holds the pushdown stack address from the Stack Pointer.
HLTA	D ₃	Acknowledge signal for Halt Instruction.

Table 1 STATUS INFORMATION DEFINITION

*These three status bits can be used to control the flow of data onto the MP8080A data bus.

NOTES

- Data in must be stable for this period during DBIN-T₃. Both t_{DS1} and t_{DS2} must be satisfied.
 - Ready signal must be stable for this period during t₂ or T_W. (Must be externally synchronized.)
 - Hold signal must be stable for this period during T₂ or T_W when entering hold mode, and during T₃, T₄, T₅, and T_{WH} when in hold mode. (External synchronization is not required.)
 - Interrupt signal must be stable during this period of the last clock cycle of any instruction in order to be recognized on the following instruction. (External synchronization is not required.)

MACHINE CYCLE	TYPE				DATA	BUS BIT			
MACHINE CYCLE	1176	D7	D ₆	D ₅	D4	D ₃	D ₂	D ₁	D ₀
Instruction fetch	1	1	0	1	0	0	0	1	0
Memory read	2	1	0	Ó	0	0	0	1	ň
Memory write	3	0	0	ō	0	0	l n	0	l ň
Stack read	4	1 1	lõ	0	0	0	1 1	1	۱ ۾
Stack write	5	0	j ,	0	ň	n	;	ا أ	ا ا
Input read	6	0	1	l ŏ	0	0	1 6	1	1 0
Output write	7	1 0	ا أ	0	1	1 0	١	١ ٪	0
Interrupt acknowledge	8	١٥	0	1	ا أ	0	١	1 1	1
Halt acknowledge	9	1 1	o o	ا أ	0	1	0		'
Interrupt acknowledge while halt	10	Ö	ŏ	1	ő	1	0 -	1	1

Table 2 STATUS WORD CHART

MNEN	MONIC	DESCRIPTION	OPERATION	<u></u>	D.			COD		D-	Do	No. of Bytes	No. of Machine (M) Cycles	No. of States (T)	C		HOITION		
DATA	TRANSF	ER GROUP		107	D 6	U 5	D 4	D3	D 2	וט	טט	L	Cycles	1	S		AC	<u> </u>	CY
LDA		Load Accumulator Direct	(A) ← ((byte 3) (byte 2))	10	0	1.	. 1	1	0	1	0	3	4	13	Τ				
LDAX	В	Load Accumulator Indirect	(A) ← ((B) (C))	١ŏ	0	0	0	1	0	1	0	1	2	7					
LDAX	D	Load Accumulator Indirect	(A) ((D) (E))	1 0	. 0	ō	1	1	0	1	0	1	2	7	İ				
LHLD		Load H and L Direct	(L) ((byte 3) (byte 2))	0	0	1	o	1	0	1	0	3	5	16	ł				
			(H) ← ((byte 3) (byte 2) + 1)	"	•	•	·	i '	٠		v	ľ	١	10	i				
LXI	В	Load Immediate, Registers B and C	(B) ← (byte 3)	0	0	0	0	0	0	0	1	3	3	10					
			(C) (byte 2)					1				,	"		ł				
LXI	D	Load Immediate, Registers D and E	(D) (byte 3)	0	0	0	1	0	0	- 0	1	3	3	10					
			(E) (byte 2)					l											
ĽXI	Н	Load Immediate, Registers H and L	(H) ← (byte)	0	0	1	0	0	0	. 0	1	3	3	10					
			(L) ← (byte 2)	1				İ											
LXI	SP	Load Immediate, Stack Pointer	(SPH) ← (byte 3)	0	0	1	1	0	0	0	1	3	3	10					
MOV	M -	Mana da Mana da	(SPL) (byte 2)	1												(F)	lags N	lot	
MOV	M,r	Move to Memory	((H) (L)) ← (r)	0	1	1	1	0	S	S	S	1	2	7		Αt	ffecte	d)	
MOV	r,M r1, r2	Move from Memory	(r) -((H) (L))	0	. 1	D	D	D	1	1	0	1	2	7					
MVI	M M	Move Registers	(r1) ← (r2)	0	1	D	D	D.	S	S	s	1	1	5	ŀ				
MVI	r	Move to Memory Immediate Move Immediate	((H) (L)) ← (byte 2)	0	0	-1	1	0	1	_ 1	0	2	3	10					
SHLD	,	Store H and L Direct	(r) — (byte 2)	0	0	D	D	D	1	1	0	2	2	7					
OI ILD		Store H and L Direct	((byte 3) (byte 2)) ← (L)	0	0	1	0	0	0	1	0	3	5	16					
STA		Store Accumulator Direct	((byte 3) (byte 2) + 1) ← (H)	١.	_			_	_		_	_							
STAX	В	Store Accumulator Indirect	((byte 3) (byte 2)) ← (A) ((B) (C)) ← (A)	0	0	. 1	1	0	0	1	0	3	4	13					
STAX	D	Store Accumulator Indirect	((D) (E)) ← (A)	1 -	-	0	0	0	0	1	0	1	2	7	1				
XCHG		Exchange H and L with D and E	(H) ↔ (D)	0	0	0	0	0	0	1	0	1	2	7					
,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,		Exonange II and E with D and E	(L) (E)	l '	١.	1	U.	١,	0	1	1	1	1	4					
ARITH	METIC G	BOUR		L				L											
	WILL TO C			1													т—	т	т
ACI		Add Immediate with Carry	(A) ← (A) + (byte 2) + (CY)	1	1	0	0	1	1	1	0	2	2	7	1	1	1	1	1
ADC ADC	M r	Add Memory with Carry	(A) (A) + ((H) (L))+ (CY)	1	0 -	0	0	1	1	1	0	1	2	7	1	1	1	1	1
ADD	м -	Add Register with Carry Add Memory	$(A) \leftarrow (A) + (r) + (CY)$		0	0	0	1	S	S	S	1	1	4 1	1	1 :	1	1	1
ADD	IVI I			1															1 .
ΔΠΠ	- 1		(A) ← (A) + ((H) (L))	1	0	0	0	0	1	1	0	1	2	7	1	1	1	1	1
	r	Add Register	$(A) \leftarrow (A) + ((H) (L))$ $(A) \leftarrow (A) + (r)$	1	0	0,,	0	0	1 S	1 S	0 S	1 1	1	7 4	1	1	ī	1	1
ADI	r	Add Flegister Add Immediate	(A) ← (A) + ((H) (L)) (A) ← (A) + (r) (A) ← (A) + (byte 2)	1 1 1	0 0 1	0 0	0	0	1 S 1	1 S 1	0 S 0	1 1 1	1 2	7 4 7.	1 1	1	1	1	1
ADI	r	Add Register	(A) ← (A) + ((H) (L)) (A) ← (A) + (r) (A) ← (A) + (byte 2) 8-bit number in Accumulator	1	0	0,,	0	0	1 S 1	1 S	0 S	1 1	1	7 4	1	1	ī	1	1
ADI	r	Add Flegister Add Immediate	(A) (A) + ((H) (L)) (A) (A) + (r) (A) (A) + (byte 2) 8-bit number in Accumulator is converted to two 4-bit BCD	1 1 1	0 0 1	0 0	0	0	1 S 1	1 S 1	0 S 0	1 1 1	1 2	7 4 7.	1 1	1	1	1	1
ADI DAA	r B	Add Register Add Immediate Decimal Adjust Accumulator	(A) (A) + ((H) (L)) (A) (A) + (r) (A) (A) + (byte 2) 8-bit number in Accumulator is converted to two 4-bit BCD digits	1 1 0	0 0 1 0	0 0 1	0 0	0 0	1 S 1 1	1 S 1 1	0 S 0 1	1 1 1	1 2 1	7 4 7. 4	1 1	1	1	1	1 1
ADI DAA DAD		Add Flegister Add Immediate	(A) (A) + ((H) (L)) (A) (A) + (r) (A) (A) + (byte 2) 8-bit number in Accumulator is converted to two 4-bit BCD digits (H) (L) (H) (L) + (B) (C)	1 1 1 0	0 0 1 0	0 0 1	0 0 0	0 0 .	1 S 1 1	1 S 1 1	0 S 0 1	1 1 1 1	1 2 1	7 4 7 4	1 1	1	1	1	1 1
ADI DAA DAD DAD	В	Add Register Add Immediate Decimal Adjust Accumulator Add B and C to H and L	(A) (A) + ((H) (L)) (A) (A) + (r) (A) (A) + (byte 2) 8-bit number in Accumulator is converted to two 4-bit BCD digits (H) (L) (H) (L) + (B) (C) (H) (L) (H) (L) + (D) (E)	1 1 1 0 0 0	0 0 1 0 0 0	0 0 1 0 0	0 0 0	0 0 0 1	1 S 1 1	1 S 1 1	0 S 0 1	1 1 1 1	1 2 1 3 3	7 4 7 4 10	1	1	1	1	1 1 1 1
ADI DAA DAD DAD DAD	B D	Add Register Add Immediate Decimal Adjust Accumulator Add B and C to H and L Add D and E to H and L Add H and L to H and L	(A) (A) + ((H) (L)) (A) (A) + (r) (A) (A) + (byte 2) 8-bit number in Accumulator is converted to two 4-bit BCD digits (H) (L) (H) (L) + (B) (C) (H) (L) (H) (L) + (D) (E) (H) (L) (H) (L) + (H) (L)	1 1 1 0 0 0 0 0	0 0 1 0	0 0 1 0 0 0	0 0 0	0 0 0 1 1	1 S 1 1	1 S 1 1	0 S 0 1	1 1 1 1 1	1 2 1 3 3	7 4 7 4 10 10	1 1	1	1	1	1 1 1 1 1 1
ADI DAA DAD DAD DAD DAD	В D Н	Add Register Add Immediate Decimal Adjust Accumulator Add B and C to H and L Add D and E to H and L	(A) (A) + ((H) (L)) (A) (A) + (r) (A) (A) + (byte 2) 8-bit number in Accumulator is converted to two 4-bit BCD digits (H) (L) (H) (L) + (B) (C) (H) (L) (H) (L) + (D) (E)	1 1 1 0 0 0 0 0 0	0 0 1 0 0 0 0 0	0 0 1 0 0	0 0 0 1 0 1	0 0 0 . 1 1 1 1 1	1 S 1 1 0 0 0	1 S 1 1 0 0	0 S 0 1 1 1 1	1 1 1 1 1 1	1 2 1 3 3 3	7 4 7. 4 10 10 10	1 1 1	1	1 1	1 1	1 1 1 1
DAD DAD DAD DAD DAD DAD DCR	В D H SP	Add Register Add Immediate Decimal Adjust Accumulator Add B and C to H and L Add D and E to H and L Add H and L to H and L Add Stack Pointer to H and L	(A) (A) + ((H) (L)) (A) (A) + (r) (A) (A) + (r) (A) (A) + (byte 2) 8-bit number in Accumulator is converted to two 4-bit BCD digits (H) (L) (H) (L) + (B) (C) (H) (L) (H) (L) + (D) (E) (H) (L) (H) (L) + (H) (L) (H) (L) (H) (L) + (H) (L)	1 1 1 0 0 0 0 0 0 0	0 0 1 0 0 0 0 0 0	0 0 1 0 0 0 1 1 1	0 0 0 1 0 1 1	0 0 0 1 1 1 0	1 S 1 1 0 0 0 0	1 S 1 1 0 0 0 0	0 5 0 1 1 1 1	1 1 1 1 1 1 1 1	1 2 1 3 3 3 3	7 4 7 4 10 10 10	1 1	1	1	1 1 1	1 1 1 1 1 1
DAD DAD DAD DAD DCR DCR DCX	В D Н SP М r В	Add Register Add Immediate Decimal Adjust Accumulator Add B and C to H and L Add D and E to H and L Add H and L to H and L Add Stack Pointer to H and L Decrement Memory	(A) (A) + ((H) (L)) (A) (A) + (r) (A) (A) + (r) (A) (A) + (byte 2) 8-bit number in Accumulator is converted to two 4-bit BCD digits (H) (L) (H) (L) + (B) (C) (H) (L) (H) (L) + (D) (E) (H) (L) (H) (L) + (H) (L) (H) (L) (H) (L) + (SP) ((H) (L) (H) (L) - 1	1 1 1 0 0 0 0 0 0	0 0 1 0 0 0 0 0	0 0 1 0 0 0 0 1 1	0 0 0 1 0 1 1 D	0 0 0 . 1 1 1 1 1	1 S 1 1 0 0 0 0 1 1	1 S 1 1 0 0	0 S 0 1 1 1 1	1 1 1 1 1 1 1 1 1 1	1 2 1 3 3 3 3 3	7 4 7 4 10 10 10 10 10	1 1 1	1	1 1	1 1	1 1 1 1 1 1
DAD DAD DAD DAD DCR DCR DCX DCX	в D H P M г в D	Add Register Add Immediate Decimal Adjust Accumulator Add B and C to H and L Add D and E to H and L Add H and L to H and L Add Stack Pointer to H and L Decrement Memory Decrement Register Decrement Registers B and C Decrement Registers D and E	(A) (A) + ((H) (L)) (A) (A) + (r) (A) (A) + (r) (A) (A) + (byte 2) 8-bit number in Accumulator is converted to two 4-bit BCD digits (H) (L) (H) (L) + (B) (C) (H) (L) (H) (L) + (D) (E) (H) (L) (H) (L) + (H) (L) (H) (L) (H) (L) + (SP) (H) (L) (H) (L) 1 (r) (r) 1	1 1 1 0 0 0 0 0 0 0 0 0	0 0 1 0 0 0 0 0 0 0	0 0 1 0 0 1 1 1 D	0 0 0 1 0 1 1	0 0 0 1 1 1 0 D	1 S 1 1 0 0 0 0	1 S 1 1 0 0 0 0	0 S 0 1 1 1 1 1	1 1 1 1 1 1 1 1	1 2 1 3 3 3 3	7 4 7 4 10 10 10	1 1	1	1	1 1 1	1 1 1 1 1 1
DAD DAD DAD DCR DCR DCX DCX DCX	вономивон	Add Register Add Immediate Decimal Adjust Accumulator Add B and C to H and L Add D and E to H and L Add H and L to H and L Add Stack Pointer to H and L Decrement Memory Decrement Registers Decrement Registers B and C Decrement Registers D and E Decrement Registers H and L	(A) (A) + ((H) (L)) (A) (A) + (r) (A) (A) + (byte 2) 8-bit number in Accumulator is converted to two 4-bit BCD digits (H) (L) (H) (L) + (B) (C) (H) (L) (H) (L) + (H) (L) (H) (L) (H) (L) + (H) (L) (H) (L) (H) (L) + (SP) ((H) (L) (H) (L)) - 1 (r) (r) - 1 (B) (C) (B) (C) - 1	1 1 1 0 0 0 0 0 0	0 0 1 0 0 0 0 0 0 0 0 0	0 0 1 0 0 1 1 1 D	0 0 0 1 0 1 1 D	0 0 1 1 1 0 D	1 S 1 1 0 0 0 0 1 1 0	1 S 1 1 0 0 0 0 0 0	0 5 0 1 1 1 1 1 1	1 1 1 1 1 1 1 1 1 1 1	1 2 1 3 3 3 3 1	7 4 7 4 10 10 10 10 5 5	1	1	1	1 1 1	1 1 1 1 1 1
DAD DAD DAD DCR DCX DCX DCX DCX	вонямгвоня	Add Register Add Immediate Decimal Adjust Accumulator Add B and C to H and L Add D and E to H and L Add H and L to H and L Add Stack Pointer to H and L Decrement Memory Decrement Register Decrement Registers B and C Decrement Registers D and E Decrement Registers H and L Decrement Registers H and L	(A) (A) + ((H) (L)) (A) (A) + (r) (A) (A) + (r) (A) (A) + (byte 2) 8-bit number in Accumulator is converted to two 4-bit BCD digits (H) (L) (H) (L) + (B) (C) (H) (L) (H) (L) + (D) (E) (H) (L) (H) (L) + (H) (L) (H) (L) (H) (L) + (SP) ((H) (L) (H) (L) -1 (F) (F) -1 (B) (C) (B) (C)1 (D) (E) (D) (E) -1 (H) (L) (H) (L) -1 (SP) (SP)1	1 1 1 0 0 0 0 0 0 0	0 0 1 0 0 0 0 0 0 0	0 0 1 0 0 1 1 1 D 0	0 0 0 1 0 1 1 D 0	0 0 0 1 1 1 0 D 1	1 S 1 1 0 0 0 0 1 1 0 0 0	1 S 1 1 0 0 0 0 0 0 0 1 1	0 5 0 1 1 1 1 1 1	1 1 1 1 1 1 1 1 1 1 1 1	1 2 1 3 3 3 3 3 1	7 4 7 4 10 10 10 10 5 5	1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1	1	1	1 1 1	1 1 1 1 1 1
DAD DAD DAD DAD DAD DCR DCR DCX DCX DCX DCX DCX	в D H 9 ^M г в D H 9 ^M	Add Register Add Immediate Decimal Adjust Accumulator Add B and C to H and L Add D and E to H and L Add H and L to H and L Add Stack Pointer to H and L Decrement Memory Decrement Register Decrement Registers B and C Decrement Registers D and E Decrement Registers H and L Decrement Register H and L Decrement Stack Power Increment Memory	(A) (A) + ((H) (L)) (A) (A) + (r) (A) (A) + (r) (A) (A) + (r) (A) (A) + (r) (A) (A) + (r) (A) (A) + (byte 2) 8-bit number in Accumulator is converted to two 4-bit BCD digits (H) (L) (H) (L) + (B) (C) (H) (L) (H) (L) + (B) (C) (H) (L) (H) (L) + (B) (C) (H) (L) (H) (L) + (B) (H) (L) (H) (L) -1 (r) (r) 1 (B) (C) (B) (C) 1 (D) (E) (D) (E) 1 (H) (L) (H) (L) 1 (SP) (SP) 1 ((H) (L) (H) (L) +- 1	1 1 1 0 0 0 0 0 0 0 0	0 0 1 0 0 0 0 0 0 0 0	0 0 1 1 0 0 0 1 1 D 0 0 1 1	0 0 0 1 0 1 1 D 0	0 0 0 1 1 1 1 0 D 1 1	1 S 1 1 0 0 0 0 1 1 0 0 0 0	1 S 1 1 0 0 0 0 0 0 0 1 1 1	0 S O 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1	1 1 1 1 1 1 1 1 1 1 1 1 1	1 2 1 3 3 3 3 3 1 1 1	7 4 7 4 10 10 10 10 10 5 5 5	1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1	1	1	1 1 1	1 1 1 1 1 1
DAD DAD DAD DAD DAD DCR DCR DCX DCX DCX DCX INR INR	воном гвоном г	Add Register Add Immediate Decimal Adjust Accumulator Add B and C to H and L Add D and E to H and L Add H and L to H and L Add Stack Pointer to H and L Decrement Memory Decrement Registers B and C Decrement Registers D and E Decrement Registers H and L Decrement Stack Power Increment Memory Increment Memory Increment Register	(A) (A) + ((H) (L)) (A) (A) + (r) (A) (A) + (r) (A) (A) + (r) (A) (A) + (r) (A) (A) + (r) (A) (A) + (byte 2) 8-bit number in Accumulator is converted to two 4-bit BCD digits (H) (L) (H) (L) + (B) (C) (H) (L) (H) (L) + (D) (E) (H) (L) (H) (L) + (H) (L) (H) (L) (H) (L) + (SP) (H) (L) (H) (L) -1 (r) (r) -1 (B) (C) (B) (C) -1 (D) (E) (D) (E) -1 (H) (L) (H) (L) -1 (SP) (SP) -1 ((H) (L) (H) (L) +1 (r) (r) +1	1 1 1 1 0 0 0 0 0 0 0 0 0 0 0	0 0 1 0 0 0 0 0 0 0 0 0	0 0 1 1 1 D 0 0 1 1 1	0 0 0 1 0 1 1 0 1 0 1	0 0 0 1 1 1 1 0 D 1 1 1	1 S 1 1 0 0 0 0 1 1 0 0 0 0 0 0	1 S 1 1 1 0 0 0 0 0 0 1 1 1 1 1	0 S 0 1 1 1 1 1 1 1 1	1 1 1 1 1 1 1 1 1 1 1 1	1 2 1 3 3 3 3 1 1 1 1 1 1	7 4 7 4 10 10 10 10 5 5 5 5	1 1 1	1 1			1 1 1 1 1 1
DAD DAD DAD DAD DAD DCR DCX DCX DCX DCX INR INR	вонямгвонямгв	Add Register Add Immediate Decimal Adjust Accumulator Add B and C to H and L Add D and E to H and L Add H and L to H and L Add Stack Pointer to H and L Decrement Memory Decrement Register Decrement Registers B and C Decrement Registers D and E Decrement Registers H and L Decrement Stack Power Increment Memory Increment Register Increment Register	(A) (A) + ((H) (L)) (A) (A) + (r) (A) (A) + (r) (A) (A) + (r) (A) (A) + (byte 2) 8-bit number in Accumulator is converted to two 4-bit BCD digits (H) (L) (H) (L) + (B) (C) (H) (L) (H) (L) + (D) (E) (H) (L) (H) (L) + (D) (E) (H) (L) (H) (L) + (SP) ((H) (L) (H) (L) -1 (F) (F) -1 (B) (C) (B) (C) -1 (D) (E) (D) (E) -1 (H) (L) (H) (L) -1 (SP) (SP) -1 ((H) (L) (H) (L) +1 (r) (r) +1 (B) (C) (B) (C) +1	1 1 1 1 1 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0	0 0 1 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0	0 0 1 0 0 1 1 1 D 0 0 1 1 1 1	0 0 0 1 0 1 1 D 0 1	0 0 0 1 1 1 1 0 0 1 1 1 1	1 S 1 1 0 0 0 0 1 1 0 0 0 0 1	1 S 1 1 0 0 0 0 0 1 1 1 1 0 0	0 S 0 1 1 1 1 1 1 1 1	1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1	1 2 1 3 3 3 3 3 1 1 1 1 1 3 3	7 4 7 4 10 10 10 10 5 5 5 5 5	I I I I I I I I I I I I I I I I I I I	1	1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1	1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1	1 1 1 1 1 1
DAD DAD DAD DAD DAD DCR DCX DCX DCX DCX INR INR	воном гвоном г	Add Register Add Immediate Decimal Adjust Accumulator Add B and C to H and L Add D and E to H and L Add H and L to H and L Add Stack Pointer to H and L Decrement Memory Decrement Registers B and C Decrement Registers D and E Decrement Registers H and L Decrement Stack Power Increment Memory Increment Memory Increment Register	(A) (A) + ((H) (L)) (A) (A) + (r) (A) (A) + (r) (A) (A) + (r) (A) (A) + (r) (A) (A) + (r) (A) (A) + (byte 2) 8-bit number in Accumulator is converted to two 4-bit BCD digits (H) (L) (H) (L) + (B) (C) (H) (L) (H) (L) + (D) (E) (H) (L) (H) (L) + (H) (L) (H) (L) (H) (L) + (SP) (H) (L) (H) (L) -1 (r) (r) -1 (B) (C) (B) (C) -1 (D) (E) (D) (E) -1 (H) (L) (H) (L) -1 (SP) (SP) -1 ((H) (L) (H) (L) +1 (r) (r) +1	0 0 0 0 0 0 0 0 0 0	0 0 1 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0	0 0 1 1 1 D 0 0 1 1 1 D D	0 0 0 1 0 1 1 D 0 1 1 0	0 0 0 1 1 1 1 0 0 0 1 1 1 1 0 0	1 S 1 1 0 0 0 0 1 1 0 0 0 1 1	1 S 1 1 0 0 0 0 0 1 1 1 0 0 0	0 S 0 1 1 1 1 1 1 1 1 1 0 0	1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1	1 2 1 3 3 3 3 3 1 1 1 1 3 3 1	7 4 7 4 10 10 10 10 5 5 5 5 5 5 5 5		1	1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1	1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1	1 1 1 1 1 1

Table 3 INSTRUCTION SET

8-BIT N-CHANNEL MICROPROCESSOR

MP8080A-I

MNEMONIC	DESCRIPTION	OPERATION			c	P C	ODE	.			No. of Bytes	No. of Machine (M)	No. of States (T)	co	DNDI.	TION	FLAC	GS
MILLMOING			D7	D6	D 5	D 4	D3	D2	D 1	Dο	,	Cycles	(1)	S	Z	AC	Р	С
ARITHMETIC	GROUP (Cont'd)						,								,			_
INX SP	Increment Stack Pointer	(SP) (SP) + 1	0	0	1	1	0	0	1	1	1	1	5	:			1	
SBB M	Subtract Memory with Borrow	$(A) \leftarrow (A) - ((H) (L)) + (CY)$	1	0	0	1	1	1	1	0	1	2	7	1	;	1	;	1
SBB r	Subtract Register with Borrow	(A) (A) - (r) - (CY)	1	0	0	1	1	S	S	S	1 2	2	7	1	1	1	i	
SBI	Subtract Immediate with Borrow	(A) - (A) - (byte 2) - (CY)	1 1	1	0	1	0	1	1	0	1	2	7		l i	i	i	
SUB M	Subtract Memory	(A) ← (A) – ((H) (L))	1 1	0	0	1	0	s	s	S	1	1	4	l i	1		1	
SUB r	Subtract Register	(A) ← (A) – (r)	Hi	1	0	1	0	1	1	o	2	2	7	1	1	1	1	١
SUI	Subtract Immediate	(A) - (A) - (byte 2)	Ŀ				Ľ				<u> </u>				L		<u> </u>	L
LOGICAL GF	ROUP											·	т			1	·	Т
ANA M	AND Memory	(A) (A) ((H) (L))	1	0	1	0	0	1	1	0	1	2	7	!	1 1	ld td	1	
ANA r	AND Register	(A) ← (A) (r)	1	0	1	0	0	S	S	S	1	1	4	!	1 1	1q	;	
ANI	AND Immediate	(A) (A) (byte 2)	1	1	1	0	0	1	1	0	2	2	7		Ι'	1,0	1:	
CMA	Complement Accumulator	$(A) \leftarrow (\overline{A})$	0	0	1	0	1	1	1	1	1	1 1	4			10	١.	
CMC	Complement Carry	$(CY) \leftarrow (\overline{CY})$	0	0	1	1	1	1	1	1	1 1	2	7	1	1 a	1 1	1	1
CMP M	Compare Memory	(A)—((H) (L))	1	0	1	1	1	1 S	1 S	S	1	1	4	1	16	1	l i	
CMP r	Compare Register	(A)—(r) (A)—(byte 2)	1	1	1	1	1	1	1	0	2	2	7	1	10	1	1 1	1
CPI	Compare Immediate	(A) — (B) V ((H) (L))	1	ò	1	1	0	1	1	0	1	2	7	1	1	0	1	l
ORA M	OR Memory	(A) - (A) V (r)	1	ō	1	1	0	s	s	s	1	1	4	1	1	0	1	ı
ORA r	OR Register OR Immediate	(A) - (A) V (byte 2)	1	1	1	1	0	1	1	0	2	2	7	1	1	0	1	1
ORI	Rotate Left through Carry	$(A_{n+1}) \leftarrow (A_n); (CY) \leftarrow (A_7)$	Ó	0	0	1	0	1	1	1	1 1	1	4		·			1
RAL	Hotate Left through Carry	(A ₀) ← (CY)					1						1			1	1	1
RAR	Rotate Right through Carry	$(A_n) \leftarrow (A_{n+1}); (CY) \leftarrow (A_0)$	0	0	0	1	1	1	1	1	1	1	4					
nan	Tiotato riigite tinoagii carry	(A ₇) ← (CY)					1				1	l			1	1		
RLC	Rotate Left	$(A_{n+1}) \leftarrow (A_n); (A_0) \leftarrow (A_7)$	0	0	0	0	0	1	1	1	1	1 1	4	1			'	
		(CY) - (A7)										١.	1.		1		١.	
RRC	Rotate Right	$(A_n) \leftarrow (A_{n-1}); (A_7) \leftarrow (A_0)$	0	0	0	0	1	1	1	1	1	1	4			'	1	l
		(CY) - (A ₀)	_	_								1	4	١.	1.	١.,		١
STC	Set Carry	(CY) - 1	0	0	1	1		1		1	1 1	2	7	1	1	0		
XRA M	Exclusive OR Memory	(A) — (A) ∀ ((H) (L))	1	0		0		1 S		S	1	1	4	1	1	0	1	1
XRA r	Exclusive OR Register	(A) - (A) ∀ (r)		0 1							2	2	7			0	ŀ	
XRI	Exclusive OR Immediate	(A) ← (A) ∀ (byte 2)					L.					L		1		ــــــــــــــــــــــــــــــــــــــ	ــــــــــــــــــــــــــــــــــــــ	_
BRANCH GF	ROUP					-							T	_				
	T	((OD) 1) (DC(I)		- 1	0	0	1	1	0	1	1 3	5	17					
CALL	Call Unconditional	((SP) - 1) ← (PCH)	1				- 1				1		1 "					
CALL	Call Unconditional	((SP) - 2) ← (PCL)	1				ĺ						''					
CALL	Call Unconditional	((SP) - 2) ← (PCL) (SP) ← (SP) - 2	1	•														
		((SP) - 2) ← (PCL) (SP) ← (SP) - 2 (PC) ← (byte 3) (byte 2)			0	1		1	n	n	3							
CC	Call Unconditional Call on Carry	((SP) - 2) ← (PCL) (SP) ← (SP) - 2 (PC) ← (byte 3) (byte 2) If CY = 1,			0	1	1	1	0	0	3	3/5	11/17					
		((SP) - 2) (PCL) (SP) (SP) - 2 (PC) (byte 3) (byte 2) If CY = 1, ((SP) - 1) (PCH)			0	1	1	1	0	0	3							
		((SP) - 2) (PCL) (SP) (SP) - 2 (PC) (byte 3) (byte 2) If CY = 1, ((SP) - 1) (PCH) ((SP) - 2) (PCL)			0	1	1	1	0	0	3							
		((SP) - 2) (PCL) (SP) (SP) - 2 (PC) (byte 3) (byte 2) If CY = 1, ((SP) - 1) (PCH) ((SP) - 2) (PCL) (SP) (SP) - 2			0	1	1	1	0	0	3							
СС	Call on Carry	((SP) - 2) (PCL) (SP) (SP) - 2 (PC) (byte 3) (byte 2) If CY = 1, ((SP) - 1) (PCH) ((SP) - 2) (PCL) (SP) (SP) - 2 (PC) (byte 3) (byte 2)	1	1							3							
		((SP) - 2) (PCL) (SP) (SP) - 2 (PC) (byte 3) (byte 2) If CY = 1, ((SP) - 1) (PCH) ((SP) - 2) (PCL) (SP) (SP) - 2 (PC) (byte 3) (byte 2) If S = 1,	1	1	0				0			3/5	11/17					
СС	Call on Carry	((SP) - 2) (PCL) (SP) (SP) - 2 (PC) (byte 3) (byte 2) If CY = 1, ((SP) - 1) (PCH) ((SP) - 2) (PCL) (SP) (SP) - 2 (PC) (byte 3) (byte 2)	1	1								3/5	11/17					
СС	Call on Carry	((SP) - 2) - (PCL) (SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If CY = 1 ((SP) - 1) - (PCH) ((SP) - 2) - (PCL) (SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If S = 1, ((SP) - 1) - (PCH)	1	1								3/5	11/17					
СС	Call on Carry	((SP) - 2) (PCL) (SP) (SP) - 2 (PC) (byte 3) (byte 2) If CY = 1, ((SP) - 1) (PCH) ((SP) - 2) (PCL) (SP) (SP) - 2 (PC) (byte 3) (byte 2) If S = 1, ((SP) - 1) (PCH) ((SP) - 2) (PCL)	1	1							3	3/5	11/17					
СС	Call on Carry	((SP) - 2) - (PCL) (SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If CY = 1, ((SP) - 1) - (PCH) ((SP) - 2) - (PCL) (SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If S = 1, ((SP) - 1) - (PCH) ((SP) - 2) - (PCL) (SP) - (SP) - 2	1	1		1	1		0	0		3/5	11/17					
СС	Call on Carry Call on Minus	((SP) - 2) (PCL) (SP) (SP) - 2 (PC) (byte 3) (byte 2) If CY = 1, ((SP) - 1) (PCH) ((SP) - 2) (PCL) (SP) (SP) - 2 (PC) (byte 3) (byte 2) If S = 1, ((SP) - 1) (PCH) ((SP) - 2) (PCL) (SP) (SP) - 2 (PC) (byte 3) (byte 2) If CY = 0, ((SP) - 1) (PCH)	1	1	1	1	1	1	0	0	3	3/5	11/17			(Flags	s Not	
СС	Call on Carry Call on Minus	((SP) - 2) (PCL) (SP) (SP) - 2 (PC) (byte 3) (byte 2) If CY = 1, ((SP) - 1) (PCH) ((SP) - 2) (PCL) (SP) (SP) - 2 (PC) (byte 3) (byte 2) If S = 1, ((SP) - 1) (PCH) ((SP) - 2) (PCL) (SP) (SPI) - 2 (PC) (byte 3) (byte 2) If CY = 0,	1	1	1	1	1	1	0	0	3	3/5	11/17			Flags Affec		
СС	Call on Carry Call on Minus	((SP) - 2) (PCL) (SP) (SP) - 2 (PC) (byte 3) (byte 2) If CY = 1, ((SP) - 1) (PCH) ((SP) - 2) (PCL) (SP) (SP) - 2 (PC) (byte 3) (byte 2) If S = 1, ((SP) - 1) (PCH) ((SP) - 2) (PCL) (SP) (SP) - 2 (PC) (byte 3) (byte 2) If CY = 0, ((SP) -1) (PCH) ((SP) -1) (PCH) ((SP) -2) (PCL)	1	1	1	1	1	1	0	0	3	3/5	11/17			∜Flags Affec		
СС	Call on Carry Call on Minus	((SP) - 2) (PCL) (SP) (SP) - 2 (PC) (byte 3) (byte 2) If CY = 1, ((SP) - 1) (PCH) ((SP) - 2) (PCL) (SP) (SP) - 2 (PC) (byte 3) (byte 2) If S = 1, ((SP) - 1) (PCH) ((SP) - 2) (PCL) (SP) (SP) - 2 (PC) (byte 3) (byte 2) If CY = 0, ((SP) - 1) (PCH) ((SP) - 2) (PCH) ((SP) - 2) (PCL) ((SP) - (SP) - 2 (PC) (byte 3) (byte 2)	1	1	1	1	0	1	0	0	3	3/5 3/5 3/5	11/17					
СС	Call on Carry Call on Minus	((SP) - 2) - (PCL) (SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If CY = 1, ((SP) - 1) - (PCH) ((SP) - 2) - (PCL) (SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If S = 1, ((SP) - 1) - (PCH) ((SP) - 2) - (PCL) ((SP) - 2) - (PCL) ((SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If CY = 0, ((SP) - 1) - (PCH) ((SP) - 2) - (PCL) ((SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If CY = 0, ((SP) - 1) - (PCH) ((SP) - 2) - (PCL) ((SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If Z = 0,	1	1	1	1	0	1	0	0	3	3/5	11/17					
CC CM CNC	Call on Carry Call on Minus Call on No Carry	((SP) - 2) - (PCL) (SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If CY = 1, ((SP) - 1) - (PCH) ((SP) - 2) - (PCL) ((SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If S = 1, ((SP) - 1) - (PCH) ((SP) - 2) - (PCL) ((SP) - 1) - (PCL) ((SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If CY = 0, ((SP) - (PCL) ((SP) - (PCL) ((SP) - (SP) - (PCL) ((SP) - (SP) - 2) If CY = 0, ((SP) - 1) - (PCH) ((SP) - (SP) - 2) ((SP) - (SP) - 2 ((SP) - (SP) - 2 ((SP) - 1) - (PCH)	1	1	1	1	0	1	0	0	3	3/5 3/5 3/5	11/17					
CC CM CNC	Call on Carry Call on Minus Call on No Carry	((SP) - 2) (PCL) (SP) (SP) - 2 (PC) (byte 3) (byte 2) If CY = 1, ((SP) - 1) (PCH) ((SP) - 2) (PCL) (SP) (SP) - 2 (PC) (byte 3) (byte 2) If S = 1, ((SP) - 1) (PCH) ((SP) - 2) (PCL) (SP) (byte 3) (byte 2) If CY = 0, ((SP) 1) (PCH) ((SP) 2) (PCL) (SP) (SP) 2 (PC) (byte 3) (byte 2) If Z = 0, ((SP) (SP) 2 ((SP) (SP) 2 ((SP) (SP) 2) ((SP) (SP) 2) ((SP) (SP) 2) ((SP) (SP) 2)	1	1	1	1	0	1	0	0	3	3/5 3/5 3/5	11/17					
CC CM CNC	Call on Carry Call on Minus Call on No Carry	((SP) - 2) - (PCL) (SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If CY = 1, ((SP) - 1) - (PCH) ((SP) - 2) - (PCL) (SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If S = 1, ((SP) - 1) - (PCH) ((SP) - 2) - (PCL) ((SP) - 2) - (PCL) ((SP) - 2) - (PCL) ((SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If CY = 0, ((SP) - 1) - (PCH) ((SP) - 2) - (PCL) ((SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If CY = 0, ((SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If Z = 0, ((SP) - 1) - (PCH) ((SP) - 2) - (PCL) ((SP) - 2) - (PCL) ((SP) - 1) - (PCH)	1	1	1	1	0	1	0	0	3	3/5 3/5 3/5	11/17					
CC CM CNC CNZ	Call on Carry Call on Minus Call on No Carry Call on Not Zero	((SP) - 2) - (PCL) (SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If CY = 1, ((SP) - 1) - (PCH) ((SP) - 2) - (PCL) ((SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If S = 1, ((SP) - 1) - (PCH) ((SP) - 2) - (PCL) ((SP) - 1) - (PCH) ((SP) - 2) - (PCL) ((SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If CY = 0, ((SP) - 1) - (PCH) ((SP) - 2) - (PCL) ((SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If Z = 0, ((SP) - 1) - (PCH) ((SP) - 2) - (PCL) ((SP) - 2) - (PCL) ((SP) - 2) - (PCH) ((SP) - 2) - (PCH) ((SP) - 2) - (PCH) ((SP) - 2) - (PCH) ((SP) - 2) - (PCH) ((SP) - 2) - (PCH) ((SP) - 2) - ((SP) - 2 ((PC) - (byte 3) (byte 2)	1 1 1 1 1	1 1 1	0 0	1 0	0 0	1 1	0	0	3	3/5 3/5 3/5	11/17					
CC CM CNC	Call on Carry Call on Minus Call on No Carry	((SP) - 2) (PCL) (SP) - (SP) - 2 (PC) (byte 3) (byte 2) If CY = 1, ((SP) - 1) (PCH) ((SP) - 2) (PCL) ((SP) - (SP) - 2 (PC) (byte 3) (byte 2) If S = 1, ((SP) - 1) (PCH) ((SP) - 2) (PCL) ((SP) - (SP) - 2 (PC) (byte 3) (byte 2) If CY = 0, ((SP) - 1) (PCH) ((SP) - 2) (PCL) ((SP) - (SP) - 2 (PC) (byte 3) (byte 2) If CY = 0, ((SP) - 1) (PCH) ((SP) - 2) (PCL) ((SP) (SP) - 2 (PC) (byte 3) (byte 2) If Z = 0, ((SP) - 1) (PCH) ((SP) - 2) (PCL) ((SP) (SP) - 2 (PC) (byte 3) (byte 2) If S = 0,	1 1 1 1 1	1 1 1	0 0	1 0	0	1 1	0	0	3	3/5 3/5 3/5	11/17					
CC CM CNC CNZ	Call on Carry Call on Minus Call on No Carry Call on Not Zero	((SP) - 2) - (PCL) (SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If CY = 1, ((SP) - 1) - (PCH) ((SP) - 2) - (PCL) (SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If S = 1, ((SP) - 1) - (PCH) ((SP) - 2) - (PCL) ((SP) - 2) - (PCL) ((SP) - 2) - (PCL) ((SP) - 1) - (PCH) ((SP) - 2) - (PCL) ((SP) - 1) - (PCH) ((SP) - 1) - (PCH) ((SP) - 1) - (PCH) ((SP) - 1) - (PCH) ((SP) - 2) - (PCL) ((SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If Z = 0, ((SP) - 1) - (PCH) ((SP) - 2) - (PCL) ((SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If Z = 0, ((SP) - 1) - (PCH) ((SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If S = 0, ((SP) - 1) - (PCH)	1 1 1 1 1	1 1 1	0 0	1 0	0 0	1 1	0	0	3	3/5 3/5 3/5	11/17					
CC CM CNC CNZ	Call on Carry Call on Minus Call on No Carry Call on Not Zero	((SP) - 2) - (PCL) (SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If CY = 1, ((SP) - 1) - (PCH) ((SP) - 2) - (PCL) ((SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If S = 1, ((SP) - 1) - (PCH) ((SP) - 2) - (PCL) ((SP) - 1) - (PCH) ((SP) - 2) - (PCL) ((SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If CY = 0, ((SP) - 1) - (PCH) ((SP) - 1) - (PCH) ((SP) - 2) - (PCL) ((SP) - 1) - (PCL) ((SP) - 1) - (PCL) ((SP) - 1) - (PCH) ((SP) - 2) - (PCL) ((SP) - 1) - (PCH) ((SP) - 2) - (PCL) ((SP) - 1) - (PCH) ((SP) - 2) - (PCL) ((SP) - 1) - (PCH) ((SP) - 1) - (PCH) ((SP) - 1) - (PCH) ((SP) - 1) - (PCH)	1 1 1 1 1	1 1 1	0 0	1 0	0 0	1 1	0	0	3	3/5 3/5 3/5	11/17					
CM CMC CNC	Call on Carry Call on Minus Call on No Carry Call on Not Zero	((SP) - 2) (PCL) (SP) - (SP) - 2 (PC) (byte 3) (byte 2) If CY = 1, ((SP) - 1) (PCH) ((SP) - 2) (PCL) (SP) - (SP) - 2 (PC) (byte 3) (byte 2) If S = 1, ((SP) - 1) (PCH) ((SP) - 1) (PCH) ((SP) - 1) (PCL) (SP) - (SP) - 2 (PC) (byte 3) (byte 2) If CY = 0, ((SP) - 11) (PCH) ((SP) - 2) (PCL) ((SP) - (SP) - 2 (PC) (byte 3) (byte 2) If CY = 0, ((SP) - 11) (PCH) ((SP) - (SP) - 2 (PC) (byte 3) (byte 2) If Z = 0, ((SP) - 1) (PCH) ((SP) - 2) (PCL) ((SP) (SP) - 2 (PC) (byte 3) (byte 2) If S = 0, ((SP) (byte 3) (byte 2) If S = 0, ((SP) (byte 3) (byte 2) If S = 0, ((SP) 1) (PCH) ((SP) 2) (PCL) ((SP) 2) (PCL)	1 1 1 1 1	1 1 1	0 0	1 0	0 0	1 1	0	0	3	3/5 3/5 3/5	11/17					
CC CM CNC CNZ	Call on Carry Call on Minus Call on No Carry Call on Not Zero Call on Positive	((SP) - 2) - (PCL) (SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If CY = 1, ((SP) - 1) - (PCH) ((SP) - 2) - (PCL) (SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If S = 1, ((SP) - 1) - (PCH) ((SP) - 2) - (PCL) ((SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If CY = 0, ((SP) - 1) - (PCH) ((SP) - 2) - (PCL) ((SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If CY = 0, ((SP) - 1) - (PCH) ((SP) - 2) - (PCL) (SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If Z = 0, ((SP) - 1) - (PCH) ((SP) - 2) - (PCL) ((SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If S = 0, ((SP) - 1) - (PCH) ((SP) - 2) - (PCL) ((SP) - 1) - (PCH) ((SP) - 2) - (PCH) ((SP) - 2) - (PCH) ((SP) - 1) - (PCH) ((SP) - 1) - (PCH) ((SP) - 1) - (PCH) ((SP) - 2) - (PCH) ((SP) - 1) - (PCH) ((SP) - 1) - (PCH) ((SP) - 2) - (PCH) ((SP) - 2) - (PCH) ((SP) - 1) - (PCH) ((SP) - 2) - (PCH) ((SP) - 1) - (PCH) ((SP) - 2) - (PCH) ((SP) - 1) - (PCH)	1 1 1 1 1 1 1 1 1	1 1 1	0 0	1 0	0 0	1 1 1	0 0	0 0	3 3	3/5 3/5 3/5	11/17					
CM CMC CNC	Call on Carry Call on Minus Call on No Carry Call on Not Zero	((SP) - 2) - (PCL) (SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If CY = 1, ((SP) - 1) - (PCH) ((SP) - 2) - (PCL) ((SP) - 2) - (PCL) ((SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If S = 1, ((SP) - 1) - (PCH) ((SP) - 2) - (PCL) (SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If CY = 0, ((SP) - 1) - (PCH) ((SP) - 2) - (PCL) (SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If Z = 0, ((SP) - 1) - (PCH) ((SP) - 2) - (PCL) ((SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If Z = 0, ((SP) - 1) - (PCH) ((SP) - 2) - (PCL) ((SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If S = 0, ((SP) - 1) - (PCH) ((SP) - 2) - (PCL) ((SP) - 2) - (PCL) ((SP) - 2) - (PCL) ((SP) - 2) - (PCL) ((SP) - 2) - (PCL) ((SP) - 2) - (PCL) ((SP) - 2) - (PCL) ((SP) - (SP) - 2 ((SP) - 1) - (PCH) ((SP) - 2) - (PCL) ((SP) - (SP) - 2 ((1 1 1 1 1 1 1 1 1	1 1 1	0 0	1 0	0 0	1 1 1	0 0	0 0	3 3	3/5 3/5 3/5 3/5	11/17					
CM CNC CNZ	Call on Carry Call on Minus Call on No Carry Call on Not Zero Call on Positive	((SP) - 2) - (PCL) (SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If CY = 1, ((SP) - 1) - (PCH) ((SP) - 2) - (PCL) ((SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If S = 1, ((SP) - 1) - (PCH) ((SP) - 1) - (PCH) ((SP) - 1) - (PCL) ((SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If CY = 0, ((SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If CY = 0, ((SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If Z = 0, ((SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If Z = 0, ((SP) - 1) - (PCH) ((SP) - 2) - (PCL) ((SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If S = 0, ((SP) - 1) - (PCH) ((SP) - 2) - (PCL) ((SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If S = 0, ((SP) - 1) - (PCH) ((SP) - 2) - (PCL) ((SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If P = 1, ((SP) - 1) - (PCH)	1 1 1 1 1 1 1 1 1	1 1 1	0 0	1 0	0 0	1 1 1	0 0	0 0	3 3	3/5 3/5 3/5 3/5	11/17					
CM CNC CNZ	Call on Carry Call on Minus Call on No Carry Call on Not Zero Call on Positive	((SP) - 2) - (PCL) (SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If CY = 1, ((SP) - 1) - (PCH) ((SP) - 2) - (PCL) (SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If S = 1, ((SP) - 1) - (PCH) ((SP) - 2) - (PCL) ((SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If CY = 0, ((SP) - 1) - (PCH) ((SP) - 2) - (PCL) ((SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If CY = 0, ((SP) - 1) - (PCH) ((SP) - 2) - (PCL) ((SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If Z = 0, ((SP) - 1) - (PCH) ((SP) - 2) - (PCL) ((SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If S = 0, ((SP) - 1) - (PCL) ((SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If S = 0, ((SP) - 1) - (PCH) ((SP) - 2) - (PCL) ((SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If P = 1, ((SP) - 1) - (PCH) ((SP) - 2) - (PCL)	1 1 1 1 1 1 1 1 1	1 1 1	0 0	1 0	0 0	1 1 1	0 0	0 0	3 3	3/5 3/5 3/5 3/5	11/17					
CM CNC CNZ	Call on Carry Call on Minus Call on No Carry Call on Not Zero Call on Positive	((SP) - 2) - (PCL) (SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If CY = 1, ((SP) - 1) - (PCH) ((SP) - 2) - (PCL) ((SP) - 2) - (PCL) ((SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If S = 1, ((SP) - 1) - (PCH) ((SP) - 2) - (PCL) (SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If CY = 0, ((SP) - 1) - (PCH) ((SP) - 2) - (PCL) (SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If Z = 0, ((SP) - 1) - (PCH) ((SP) - 2) - (PCL) ((SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If Z = 0, ((SP) - 1) - (PCH) ((SP) - 2) - (PCL) ((SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If S = 0, ((SP) - 1) - (PCH) ((SP) - 2) - (PCL) ((SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If S = 0, ((SP) - 1) - (PCH) ((SP) - 2) - (PCL) ((SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If P = 1, ((SP) - 1) - (PCH) ((SP) - 2) - (PCL) ((SP) - (SP) - 2	1 1 1 1 1 1 1 1 1	1 1 1	0 0	1 0	0 0	1 1 1	0 0	0 0	3 3	3/5 3/5 3/5 3/5	11/17					
CC CM CNC CNZ CP	Call on Carry Call on Minus Call on No Carry Call on Not Zero Call on Positive Call on Parity Even	((SP) - 2) - (PCL) (SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If CY = 1, ((SP) - 1) - (PCH) ((SP) - 2) - (PCL) (SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If S = 1, ((SP) - 1) - (PCH) ((SP) - 2) - (PCL) ((SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If CY = 0, ((SP) - 1) - (PCH) ((SP) - 2) - (PCL) ((SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If CY = 0, ((SP) - 1) - (PCH) ((SP) - 2) - (PCL) ((SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If Z = 0, ((SP) - 1) - (PCH) ((SP) - 2) - (PCL) ((SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If S = 0, ((SP) - 1) - (PCL) ((SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If S = 0, ((SP) - 1) - (PCH) ((SP) - 2) - (PCL) ((SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If P = 1, ((SP) - 1) - (PCH) ((SP) - 2) - (PCL)	1 1 1 1 1 1 1 1 1	1 1 1 1	1 0 0	1 0 0	0 0	1 1 1	0 0	0 0	3 3	3/5 3/5 3/5 3/5	11/17					
CM CNC CNZ	Call on Carry Call on Minus Call on No Carry Call on Not Zero Call on Positive	((SP) - 2) - (PCL) (SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If CY = 1, ((SP) - 1) - (PCH) ((SP) - 2) - (PCL) ((SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If S = 1, ((SP) - 1) - (PCH) ((SP) - 2) - (PCL) ((SP) - 1) - (PCH) ((SP) - 2) - (PCL) ((SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If CY = 0, ((SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If CY = 0, ((SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If Z = 0, ((SP) - 1) - (PCL) ((SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If Z = 0, ((SP) - 1) - (PCL) ((SP) - 2) - (PCL) ((SP) - 2) - (PCL) ((SP) - 1) - (PCH) ((SP) - 2) - (PCL) ((SP) - 1) - (PCH) ((SP) - 2) - (PCL) ((SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If P = 1, ((SP) - 1) - (PCH) ((SP) - 2) - (PCL) ((SP) - 1) - (PCH) ((SP) - 1) - (PCH) ((SP) - 1) - (PCH) ((SP) - 1) - (PCH) ((SP) - 1) - (PCH) ((SP) - 1) - (PCH) ((SP) - 1) - (SP) - 2 (PC) - (byte 3) (byte 2) (PC) - (byte 3) (byte 2)	1 1 1 1 1 1 1 1 1	1 1 1 1	1 0 0	1 0 0	0 0	1 1 1	0 0	0 0	3 3	3/5 3/5 3/5 3/5	11/17 11/17 11/17 11/17 11/17					
CC CM CNC CNZ CP	Call on Carry Call on Minus Call on No Carry Call on Not Zero Call on Positive Call on Parity Even	((SP) - 2) - (PCL) (SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If CY = 1, ((SP) - 1) - (PCH) ((SP) - 2) - (PCL) (SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If S = 1, ((SP) - 1) - (PCH) ((SP) - 2) - (PCL) ((SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If CY = 0, ((SP) - 1) - (PCH) ((SP) - 2) - (PCL) ((SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If CY = 0, ((SP) - 1) - (PCH) ((SP) - 2) - (PCL) ((SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If Z = 0, ((SP) - 1) - (PCH) ((SP) - 2) - (PCL) ((SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If S = 0, ((SP) - 1) - (PCL) ((SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If S = 0, ((SP) - 1) - (PCH) ((SP) - 2) - (PCL) ((SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If P = 1, ((SP) - 1) - (PCH) ((SP) - 2) - (PCL) ((SP) - 1) - (PCH) ((SP) - 2) - (PCL) ((SP) - 1) - (PCH) ((SP) - 2) - (PCL) ((SP) - 1) - (PCH) ((SP) - 2) - (PCL) ((SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If P = 0,	1 1 1 1 1 1 1 1 1	1 1 1 1	1 0 0	1 0 0	0 0	1 1 1	0 0	0 0	3 3	3/5 3/5 3/5 3/5	11/17 11/17 11/17 11/17 11/17					
CC CM CNC CNZ CP	Call on Carry Call on Minus Call on No Carry Call on Not Zero Call on Positive Call on Parity Even	((SP) - 2) - (PCL) (SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If CY = 1, ((SP) - 1) - (PCH) ((SP) - 2) - (PCL) ((SP) - 2) - (PCL) ((SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If S = 1, ((SP) - 1) - (PCH) ((SP) - 2) - (PCL) (SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If CY = 0, ((SP) - 1) - (PCH) ((SP) - 2) - (PCL) (SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If Z = 0, ((SP) - 1) - (PCH) ((SP) - 2) - (PCL) ((SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If Z = 0, ((SP) - 1) - (PCH) ((SP) - 2) - (PCL) ((SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If S = 0, ((SP) - 1) - (PCH) ((SP) - 2) - (PCL) ((SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If P = 1, ((SP) - 1) - (PCH) ((SP) - 2) - (PCL) ((SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If P = 1, ((SP) - 1) - (PCH) ((SP) - 2) - (PCL) ((SP) - (SP) - 2 (PC) - (byte 3) (byte 2) If P = 0, ((SP) - 1) - (PCH)	1 1 1 1 1 1 1 1 1	1 1 1 1	1 0 0	1 0 0	0 0	1 1 1	0 0	0 0	3 3	3/5 3/5 3/5 3/5	11/17 11/17 11/17 11/17 11/17					

Table 3 INSTRUCTION SET (Cont'd)

MP8080A-I

ý	

MNEMONIC	DESCRIPTION	OPERATION				OP	co	DE				No. of Bytes	No. of Machine (M)	No. of States (T)	CONDITION FLAG
BRANCH GE	ROUP (Cont'd)		D,	D	6 D	5 D	1 D	3 D	2 [01 0	0		Cycles	(')	S Z AC P
cz	Call on Zero	If Z = 1.	T ₁	- 1		0		1	1	0	٥١	3	2/5	11/17	
	Sun 817 2610	((SP) - 1) ← (PCH) ((SP) - 2) ← (PCL) (SP) ← (SP) - 2	'					,	ij.	U		3	3/5	11/17	
JC	Jump on Carry	(PC) ← (byte 3) (byte 2) If CY = 1,	1	1	0	1		1 (0	1	0	3	3	10	
JM	Jump on Minus	(PC) — (byte 3) (byte 2) If S = 1, (PC) — (byte 3) (byte 2)	1	1	1	1	-	1 (0	1	0	3	3	10	
JMP JNC	Jump Unconditional Jump on No Carry	(PC) ← (byte 3) (byte 2) If CY = 0, (PC) ← (byte 3) (byte 2)	1	1							1 0	3 3	3 3	10 10	
JNZ	Jump on Lot Zero	If $Z = 0$, $(PC) \leftarrow (byte 3) (byte 2)$	1	1	0	0	() ()	1 .	0	3	3	10	
JP	Jump on Positive	If S = 0, (PC) — (byte 3) (byte 2)	1	1	1	1	0) ()	1	0	3	3	10	
JPE	Jump on Parity Even	<pre>if P = 1, (PC) ← (byte 3) (byte 2)</pre>			1		1				0	3	3	10	
JPO	Jump on Parity Odd	If P = 0, (PC) (byte 3) (byte 2)	İ		1						0	3	3	10	
JZ PCHL	Jump on Zero H and L to Program Counter	If Z = 1, (PC) ← (byte 3) (byte 2) (PCH) ← (H)			0		1				0	3	3	10 5	
RC	Return on Carry	(PCL) (L) If CY = 1,			0		1				0	1	1/3	5/11	(Flags Not
RET	Return	(PCL) - ((SP)) (PCH) - ((SP) + 1) (SP) - (SP) + 2 (PCL) - ((SP)) (PCH) - ((SP) + 1) (SP) - (SP) + 2;	1	1	0	0	1	C)	0	1	1	3	10	Affected)
RM	Return on Minus	If $S = 1$, (PCL) ((SP)) (PCH) ((SP) + 1) (SP) (SP) + 2	1	1	1	1	1)	0 (1	1/3	5/11	
RNC	Return on No Carry	If CY = 0, (PCL) ((SP)) (PCH) ((SP) + 1) (SP) (SP) + 2	1	1	0	1	0	C)	0 (1	1/3	5/11	
RNZ	Return on Not Zero	If $Z = 0$, $(PCL) \leftarrow ((SP))$ $(PCH) \leftarrow ((SP) + 1)$ $(SP) \leftarrow (SP) + 2$	1	1	0	0	0	. 0	' '	0 (1	1/3	5/11	•
RP	Return on Positive	If S = 0, (PCL) ← ((SP)) (PCH) ← ((SP) + 1) (SP) ← (SP) + 2	1	1	1	1	0	. 0	(0 (1	1/3	5/11	
RPE	Return on Parity Even	If P = 1, (PCL) ← ((SP)) (PCH) ← ((SP) + 1) (SP) ← (SP) + 2	1	1	1	0	1	0	(0 (1	1/3	5/11	
RPO	Return on Parity Odd	If P = 0, (PCL) ← ((SP)) (PCH) ← ((SP) + 1) (SP) ← (SP) + 2	1	1	1	0	0	0	() (1	1/3	5/11	
RST	Restart	((SP) - 1) ← (PCH) ((SP) - 2) ← (PCL) (SP) ← (SP) - 2 (PC) ← 8* (NNN)			N		N	1	1	1		1	3	11	
RZ	Return on Zero	H Z = 1, (PCL) ((SP)) (PCH) ((SP) + 1) (SP) (SP) + 2	1	1	0	0	1	0	O	0		1	1/3	5/11	

Table 3 INSTRUCTION SET (Cont'd)

8-BIT N-CHANNEL MICROPROCESSOR

MP8080A-I

MNEMO	NIC	DESCRIPTION	OPERATION				OP	COD	E			No. of Bytes	No. of Machine (M)	States	cc	NDI	TION	FLA	GS
IIII C	,,,,,			D7	D ₆	D ₅	D4	Dз	D ₂	D ₁	Dο		Cycles	(T)	s	Z	AC	P	CY
STACK	, I/O, A	ND MACHINE CONTROL GROUP												,					т—
DI		Disable Interrupts	The Interrupt system is dis- abled following the execution of the DI instruction.	1	1	1	1	0	0	1	1	1	1	4		. !			
EI		Enable Interrupts	The interrupt system is en- abled following the execution of next instruction.	1	1	1	1	1	0	1	1	1	1	4					
HLT		Halt	Processor is stopped; registers and flags are unaffected.	0	1	1	1	0	1	1	0	1	1	7		•			
IN NOP		Input No operation	 (A)	0	0	0	1	0	0	0	0	1	3	10					
OUT POP	В	Output Pop Registers B and C off Stack	(data) (A) (C) ((SP)) (B) ((SP) + 1) (SP) (SP) + 2	1 1	1	0	0	0	0	0	1	1	3	10					
POP	D	Pop Registers D and E off Stack	$(SP) \leftarrow (SP) + 2$ $(E) \leftarrow ((SP))$ $(D) \leftarrow ((SP) + 1)$ $(SP) \leftarrow (SP) + 2$	1	1	0	1	0	0	0	1	1	3	10					
POP	Н	Pop Registers H and L off Stack	(L) \leftarrow ((SP)) (H) \leftarrow ((SP) + 1) (SP) \leftarrow (SP) + 2	1	1	1		0	0	0	1	1	3	10					
POP	PSW	Pop Accumulator and Flags off Stack	(CY) - ((SP)) ₀ (P) - ((SP)) ₂ (AC) - ((SP)) ₄ (Z) - ((SP)) ₆ (S) - ((SP)) ₇ (A) - ((SP) + 1) (SP) - (SP) + 2	1		1		0	0	0	1	1	3	10	1		1	1	1
PUSH	В	Push Registers B and C on Stack	((SP) - 1) (B) ((SP) - 2) (C) (SP) (SP) 2	1	1		0	0	1	0	1	1	3	11					'
PUSH	D	Push Registers D and E on Stack	((SP) - 1) ← (D) ((SP) - 2) ← (E) (SP) ← (SP) - 2	1	1	0	1	0	1	0	1	1	3	11				ľ	
PUSH	Н	Push Registers H and L on Stack	((SP) - 1) (H) ((SP) - 2) (L) (SP) (SP) 2	1	1	1	0	0	1	0	1	1	3	11					1
PUSH	PSW	Push Accumulator and Flags on Stack	((SP) - 1) - (A) ((SP) - 2)0 - (CV) ((SP) - 2)1 - 1 ((SP) - 2)2 - (P) ((SP) - 2)3 - 0 ((SP) - 2)4 - (AC) ((SP) - 2)5 - 0 ((SP) - 2)6 - (Z) ((SP) - 2)7 - (S) ((SP) - 2)7 - (S)	1	1	1	1	0	1	0	1	1	3	11					
SPHL XTHL		Move H and L to Stack Pointer Exchange Top of Stack with H and L	(SP) ← (H) (L) (L) → ((SP)) (H) → ((SP) + 1)	1						0	1	1	1 5	5 18					

NOTES

- a. Z = 1 if (A) = (H) (L); CY = 1 if (A) < (H) (L)b. Z = 1 if (A) = (r);
- CY = 1 if (A) < (r)
- c. Z = 1 if (A) = (byte 2); CY = 1 if (A) < (byte 2)
- d. As if an arithmetic operation were performed.

Table 3 INSTRUCTION SET (Cont'd)

MP8080A-I

CONDITION FLAGS AND STANDARD RULES

There are five condition flags associated with the execution of instructions on the MP8080A. They are Zero, Sign, Parity, Carry, and Auxiliary Carry, and each flag is represented by a 1-bit register in the CPU. A flag is "set" by forcing the bit to 1, "reset" by forcing the bit to 0. The bit positions of the flags are indicated in the PUSH and POP PSW instructions.

Unless indicated otherwise, when an instruction affects a flag, it affects it in the following manner:

ZERO (Z)

If the result of an instruction has the value 0, this flag is set; otherwise, it is reset.

SIGN (S)

If the most significant bit of the result of the operation has the value 1, this flag is set; otherwise, it is reset.

PARITY (P)

If the modulo 2 sum of the bits of the result of the operation is 0 (that is, if the result has even parity), this flag is set; otherwise, it is reset (that is, if the result has odd parity).

CARRY (CY)

If the instruction resulted in a carry (from addition) or a borrow (from subtraction or a comparison) out of the high-order bit, this flag is set; otherwise, it is reset.

AUXILIARY CARRY (AC)

If the instruction caused a carry out of bit 3 and into bit 4 of the resulting value, the auxiliary carry is set; otherwise, it is reset. This flag is affected by singleprecision additions, subtractions, increments, decrements. comparisons. and logical operations: however, AC is used principally with additions and increments preceding a DAA (Decimal Adjust Accumulator) Instruction.

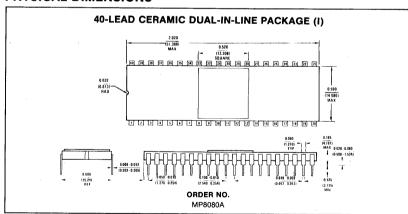
SYMBOLS AND ABBREVIATIONS

The following symbols and abbreviations are used in the subsequent description of the MP8080A instructions:

Symbols	Meaning				
A B C	Register A (Acc Register B	cumulator)			
D	Register C Register D	· ·			
H	Register H Register L	4.5			
DDD, SSS	The bit patter one of the regis	sters A, B, C, D,			
	SSS = source):	= destination,			
	DDD or SSS	Register Name			
	111	A			
	000 001	B C			
	010	D			
	011 100	E H			
	101	Ľ.			
byte 2	The second byte	e of the instruc-			
byte 3	The third byte of the instruc-				
port r,r1,r2	8-bit address o One of the regis E, H, L				

Symbols	Meaning
PC	16-bit program counter regis- ter (PCH and PCL are used to refer to the high-order and
SP	low-order 8 bits respectively.) 16-bit stack pointer register (SPH and SPL are used to refer to the high-order and low-
()	order 8 bits respectively.) The contents of the memory location or registers enclosed
← ∧ ∀ ∨	in the parentheses "Is replaced by" Logical AND Exclusive-OR
+	Inclusive-OR Addition
- *	Twos complement subtraction Multiplication "Exchange"
_	The ones complement (for example, (A))
n NNN	The restart number 0 through 7 The binary representation 000 through 111 for restart number 0 through 7 resepctively
• 0 1	"Not affected" "Reset" "Set"
X ‡	Unknown Flags affected according to Standard Rules, except as noted.

PHYSICAL DIMENSIONS



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2651-I

DESCRIPTION

The Signetics 2651 PCI is a universal synchronous/asychronous data communications controller chip designed for microcomputer systems. It interfaces directly to the Signetics 2650 microprocessor and may be used in a polled or interrupt driven system environment. The 2651 accepts programmed instructions from the microprocessor and supports many serial data communication disciplines, synchronous and asynchronous, in the full or half-duplex mode.

The PCI serializes parallel data characters received from the microprocessor for transmission. Simultaneously, it can receive serial data and convert it into parallel data characters for input to the microcomputer.

The 2651 contains a baud rate generator which can be programmed to either accept an external clock or to generate internal transmit or receive clocks. Sixteen different baud rates can be selected under program control when operating in the internal clock mode.

The PCI is constructed using Signetics nchannel silicon gate depletion load technology and is packaged in a 28-pin DIP.

FEATURES

- Synchronous operation
 to 8-bit characters
 Single or double SYN operation
 Internal character synchronization
 Transparent or non-transparent mode
 Automatic SYN or DLE-SYN insertion
 SYN or DLE stripping
 Odd, even, or no parity
 Local or remote maintenance loop
 back mode
 Baud rate: dc to 0.8M baud (1X clock)
- Asynchronous operation
 5 to 8-bit characters
 1, 1 1/2 or 2 stop bits
 Odd, even, or no parity
 Parity, overrun and framing error detection
 Line break detection and generation
 False start bit detection
 Automatic serial echo mode

back mode
Baud rate: dc to 0.8M baud (1X clock)
dc to 50k baud (16X clock)
dc to 12.5k baud (64X clock)

Local or remote maintenance loop

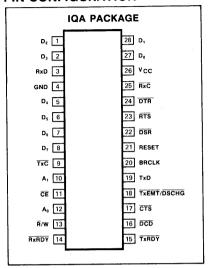
OTHER FEATURES

- Internal or external baud rate clock
- 16 internal rates-50 to 19,200 baud
- Double buffered transmitter and receiver
- · Full or half duplex operation
- Fully compatible with 2650 CPU
- TTL compatible inputs and outputs
- Single 5V power supply
- No system clock required
- 28-pin dual in-line package

APPLICATIONS

- Intelligent terminals
- Network processors
- Front end processors
- Remote data concentrators
- Computer to computer links
- Serial peripherals

PIN CONFIGURATION



PIN DESIGNATION

PIN NO.	SYMBOL	NAME AND FUNCTION	TYPE
27,28,1,2, 5-8	D ₀ -D ₇	8-bit data bus	1/0
21	RESET	Reset	1
12,10	A ₀ -A ₁	Internal register select lines	1
13	R/W	Read or write command	l I
11	CE	Chip enable input	1
22	DSR	Data set ready	ı
24	DTR	Data terminal ready	. 0
23	RTS	Request to send	0
17	CTS	Clear to send	l
16	DCD	Data carrier detected	į į
18	TxEMT/DSCHG	Transmitter empty or data set change	0
9	TxC	Transmitter clock	I/O
25	RxC	Receiver clock	I/O
19	TxD	Transmitter data	0
3	RxD	Receiver data	1
15	TxRDY	Transmitter ready	0
14	RxRDY	Receiver ready	0
20	BRCLK	Baud rate generator clock	1
26	Vcc	+5V supply	1
4	GND	Ground	1
		1	

ABSOLUTE MAXIMUM RATINGS1

PARAMETER	RATING	UNIT
Operating ambient temperature ² Storage temperature All voltages with respect to ground ³	0 to +70 -65 to +150 -0.5 to +6.0	°C °C V

PRELIMINARY SPECIFICATION

Manufacturer reserves the right to make design and process changes and improvements.

BAUD RATE	THEORETICAL FREQUENCY 16X CLOCK	ACTUAL FREQUENCY 16X CLOCK	PERCENT ERROR	DIVISOR
50	0.8 KHz	0.8 KHz		6336
75	1.2	1.2		4224
110	1.76	1.76		2880
134.5	2.152	2.1523	0.016	2355
150	2.4	2.4		2112
300	4.8	4.8		1056
600	9.6	9.6		528
1200	19.2	19.2		264
1800	28.8	28.8		176
2000	32.0	32.081	0.253	158
2400	38.4	38.4		132
3600	57.6	57.6		88
4800	76.8	76.8		66
7200	115.2	115.2		44
9600	153.6	153.6		33
19200	307.2	316.8	3.125	16

NOTE

Table 1 BAUD RATE GENERATOR CHARACTERISTICS
Crystal Frequency = 5.0688MHz

PIN NAME	PIN NO.	INPUT/OUTPUT	FUNCTION
Vcc	26		+5V supply input
GND	4	1	Ground
RESET	21	1	A high on this input performs a master reset on the 2651. This signal asynchronously terminates any device activity and clears the Mode, Command and Status registers. The device assumes the idle state and remains there until initialized with the appropriate control words.
A1-A0	10,12		Address lines used to select internal PCI registers.
<u>R</u> ∕w	13		Read command when low, write command when high.
CE	11	1	Chip enable command. When low, indicates that control and data lines to the PCI are valid and that the operation specified by the \overline{R}/W , A_1 and A_0 inputs should be performed. When high, places the D_0 - D_7 lines in the tri-state condition.
D ₇ -D ₀	8,7,6,5, 2,1,28,27	1/0	8-bit, three-state data bus used to transfer commands, data and status between PCI and the CPU. Do is the least significant bit; D7 the most significant bit.
TxRDY	15	0	This output is the complement of Status Register bit SR0. When low, it indicates that the Transmit Data Holding Register (THR) is ready to accept a data character from the CPU. It goes high when the data character is loaded. This output is valid only when the transmitter is enabled. It is an open drain output which can be used as an interrupt to the CPU.
RxRDY	14	0	This output is the complement of Status Register bit SR1. When low, it indicates that the Receive Data Holding Register (RHR) has a character ready for input to the CPU. It goes high when the RHR is read by the CPU, and also when the receiver is disabled. It is an open drain output which can be used as an interrupt to the CPU.
TxEMT/DSCHG	18	O	This output is the complement of Status Register bit SR2. When low, it indicates that the transmitter has completed serialization of the last character loaded by the CPU, or that a change of state of the DSR or DCD inputs has occurred. This output goes high when the Status Register is read by the CPU, if the TxEMT condition does not exist. Otherwise, the THR must be loaded by the CPU for this line to go high. It is an open drain output which can be used as an interrupt to the CPU.

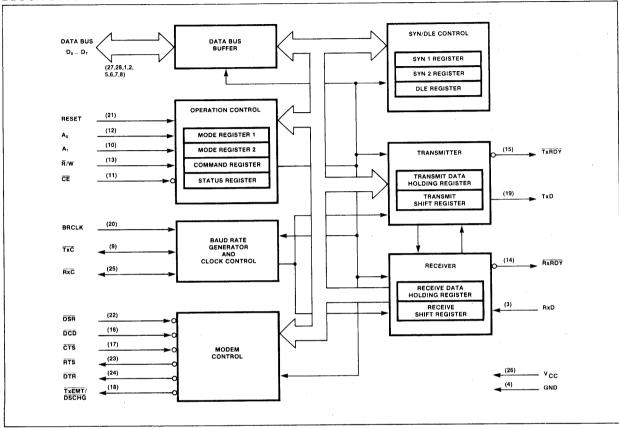
Table 2 CPU-RELATED SIGNALS

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¹⁶X clock is used in asynchronous mode. In synchronous mode, clock multiplier is 1X.

2651-1

BLOCK DIAGRAM



BLOCK DIAGRAM

The PCI consists of six major sections. These are the transmitter, receiver, timing, operation control, modem control and SYN/DLE control. These sections communicate with each other via an internal data bus and an internal control bus. The internal data bus interfaces to the microprocessor data bus via a data bus buffer.

Operation Control

This functional block stores configuration and operation commands from the CPU and generates appropriate signals to various internal sections to control the overall device operation. It contains read and write circuits to permit communications with the microprocessor via the data bus and contains Mode Registers 1 and 2, the Command Register, and the Status Register. Details of register addressing and protocol are presented in the PCI Programming section of this data sheet.

Timing

The PCI contains a Baud Rate Generator (BRG) which is programmable to accept external transmit or receive clocks or to divide an external clock to perform data communications. The unit can generate 16 commonly used baud rates, any one of which can be selected for full duplex operation. See Table 1.

Receiver

The Receiver accepts serial data on the RxD pin, converts this serial input to parallel format, checks for bits or characters that are unique to the communication technique and sends an "assembled" character to the CPU.

Transmitter

The Transmitter accepts parallel data from the CPU, converts it to a serial bit stream, inserts the appropriate characters or bits (based on the communication technique) and outputs a composite serial stream of data on the TxD output pin.

Modem Control

The modem control section provides interfacing for three input signals and three output signals used for "handshaking" and status indication between the CPU and a modem.

SYN/DLE Control

This section contains control circuitry and three 8-bit registers storing the SYN1, SYN2, and DLE characters provided by the CPU. These registers are used in the synchronous mode of operation to provide the characters required for synchronization, idle fill and data transparency.

INTERFACE SIGNALS

The PCI interface signals can be grouped into two types: the CPU-related signals (shown in Table 2), which interface the 2651 to the microprocessor system, and the device-related signals (shown in Table 3), which are used to interface to the communications device or system.

PIN NAME	PIN NO.	INPUT/OUTPUT	FUNCTION
BRCLK	20	1	5.0688MHz clock input to the internal baud rate generator. Not required if external receiver and transmitter clocks are used
RxC	25	1/0	Receiver clock. If external receiver clock is programmed, this input controls the rate at which the character is to be received. Its frequency is 1X, 16X or 64X the baud rate, as programmed by Mode Register 1. Data is sampled on the rising edge of the clock. If internal receiver clock is programmed, this pin becomes an output at 1X the programmed baud rate.
TxC	9	1/0	Transmitter clock. If external transmitter clock is programmed, this input controls the rate at which the character is transmitted. Its frequency is 1X, 16X or 64X the baud rate, as programmed by Mode Register 1. The transmitted data changes on the falling edge of the clock. If internal transmitter clock is programmed, this pin becomes an output at 1X the programmed baud rate.
RxD	3	1	Serial data input to the receiver. "Mark" is high, "Space" is low.
TxD	19	0	Serial data output from the transmitter. "Mark" is high, "Space" is low. Held in Mark condition when the transmitter is disabled.
DSR	22	1	General purpose input which can be used for Data Set Ready or Ring Indicator condition. Its complement appears as Status Register bit SR7. Causes a low output on TxEMT/DSCHG when its state changes.
DCD	16		Data Carrier Detect input. Must be low in order for the receiver to operate. Its complement appears as Status Register bit SR6. Causes a low output on TXEMT/DSCHG when its state changes.
CTS	17	1 1	Clear to Send input. Must be low in order for the transmitter to operate.
DTR	24	O	General purpose output which is the complement of Command Register bit CR1. Normally used to indicate Data Terminal Ready.
RTS	23	0	General purpose output which is the complement of Command Register bit CR5. Normally used to indicate Request to Send.

Table 3 DEVICE-RELATED SIGNALS

OPERATION

The functional operation of the 2651 is programmed by a set of control words supplied by the CPU. These control words specify items such as synchronous or asynchronous mode, baud rate, number of bits per character, etc. The programming procedure is described in the PCI Programming section of this data sheet.

After programming, the PCI is ready to perform the desired communications functions. The receiver performs serial to parallel conversion of data received from a modem or equivalent device. The transmitter converts parallel data received from the CPU to a serial bit stream. These actions are accomplished within the framework specified by the control words.

Receiver

The 2651 is conditioned to receive data when the DCD input is low and the RxEN bit in the command register is true. In the asynchronous mode, the receiver looks for a high to low transition of the start bit on the RxD input line. If a transition is detected, the state of the RxD line is sampled again after a delay of one-half of a bit time. If RxD is now high, the search for a valid start bit is begun again. If RxD is still low, a valid start bit is assumed and the receiver continues to sample the input line at one bit time intervals

until the proper number of data bits, the parity bit, and the stop bit(s) have been assembled. The data is then transferred to the Receive Data Holding Register, the RxRDY bit in the status register is set, and the RxRDY output is asserted. If the character length is less than 8 bits, the high order unused bits in the Holding Register are set to zero. The Parity Error, Framing Error, and Overrun Error status bits are set if required. If a break condition is detected (RxD is low for the entire character as well as the stop bit [s]), only one character consisting of all zeros (with the FE status bit set) will be transferred to the Holding Register. The RxD input must return to a high condition before a search for the next start bit begins.

When the PCI is initialized into the synchronous mode, the receiver first enters the hunt mode. In this mode, as data is shifted into the Receiver Shift Register a bit at a time, the contents of the register are compared to the contents of the SYN1 register. If the two are not equal, the next bit is shifted in and the comparison is repeated. When the two registers match, the hunt mode is terminated and character assembly mode begins. If single SYN operation is programmed, the SYN DETECT status bit is set. If double SYN operation is programmed, the first character assembled after SYN1 must be SYN2 in order for the SYN DETECT bit to be set.

Otherwise, the PCI returns to the hunt mode. (Note that the sequence SYN1-SYN1-SYN2 will not achieve synchronization). When synchronization has been achieved. the PCI continues to assemble characters and transfer them to the Holding Register, setting the RxRDY status bit and asserting the RxRDY output each time a character is transferred. The PE and OE status bits are set as appropriate. Further receipt of the appropriate SYN sequence sets the SYN DE-TECT status bit. If the SYN stripping mode is commanded, SYN characters are not transferred to the Holding Register, Note that the SYN characters used to establish initial synchronization are not transferred to the Holding Register in any case.

Transmitter

The PCI is conditioned to transmit data when the CTS input is low and the TxEN command register bit is set. The 2651 indicates to the CPU that it can accept a character for transmission by setting the TxRDY status bit and asserting the TxRDY output. When the CPU writes a character into the Transmit Data Holding Register, these conditions are negated. Data is transferred from the Holding Register to the Transmit Shift Register when it is idle or has completed transmission of the previous character. The TxRDI conditions are then asserted

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again. Thus, one full character time of buffering is provided.

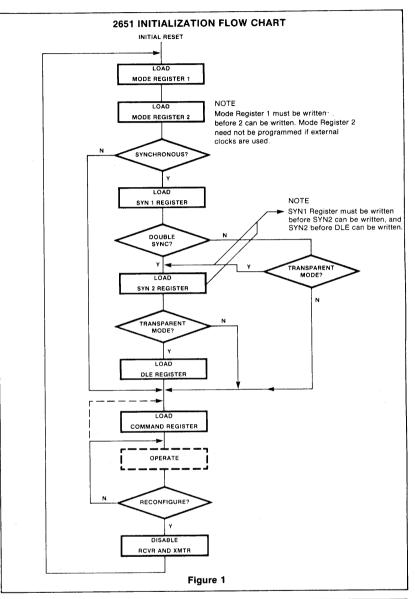
In the asynchronous mode, the transmitter automatically sends a start bit followed by the programmed number of data bits, the least significant bit being sent first. It then appends an optional odd or even parity bit and the programmed number of stop bits. If, following transmission of the stop bits, a new character is not available in the Transmit Holding Register, the TxD output remains in the marking (high) condition and the TxEMT/DSCHG output and its corresponding status bit are asserted. Transmission resumes when the CPU loads a new character into the Holding Register. The transmitter can be forced to output a continuous low (BREAK) condition by setting the Send Break command bit high.

In the synchronous mode, when the 2651 is initially conditioned to transmit, the TxD output remains high and the TxRDY condition is asserted until the first character to be transmitted (usually a SYN character) is loaded by the CPU. Subsequent to this, a continuous stream of characters is transmitted. No extra bits (other than parity, if commanded) are generated by the PCI unless the CPU fails to send a new character to the PCI by the time the transmitter has completed sending the previous character. Since synchronous communications does not allow gaps between characters, the PCI asserts TxEMT and automatically "fills" the gap by transmitting SYN1s, SYN1-SYN2 doublets, or DLE-SYN1 doublets, depending on the command mode. Normal transmission of the message resumes when a new character is available in the Transmit Data Holding Register. If the SEND DLE bit in the command register is true, the DLE character is automatically transmitted prior to transmission of the message character.

PCI PROGRAMMING

Prior to initiating data communications, the 2651 operational mode must be programmed by performing write operations to the mode and command registers. In addition, if synchronous operation is programmed, the appropriate SYN/DLE registers must be loaded. The PCI can be reconfigured at any time during program execution. However, the receiver and transmitter should be disabled if the change has an effect on the reception or transmission of a character. A flowchart of the initialization process appears in Figure 1.

The internal registers of the PCI are accessed by applying specific signals to the $\overline{\text{CE}}$, $\overline{\text{R}}/\text{W}$, A_1 and A_0 inputs. The conditions necessary to address each register are shown in Table 4.



CE	A 1	A 0	R/W	FUNCTION		
1	Х	х	Х	Tri-state data bus		
1 0	0	0	0	Read receive holding register		
1 0	0	0	1 1	Write transmit holding register		
0	0	1	0	Read status register		
0	1 0	1	1 1	Write SYN1/SYN2/DLE registers		
0	1	0	0	Read mode registers 1/2		
l 0	1 1	0	1 1	Write mode registers 1/2		
О	1	1	0	Read command register		
0	1	1	1	Write command register		

NOTE

See AC Characteristics section for timing requirements.

Table 4 2651 REGISTER ADDRESSING

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PROGRAMMABLE COMMUNICATIONS INTERFACE (PCI)

PRELIMINARY SPECIFICATION

Mode Register 1 (MR1)

The SYN1, SYN2, and DLE registers are accessed by performing write operations with the conditions $A_1=0$, $A_0=1$, and $\overline{R}/W=$ 1. The first operation loads the SYN1 register. The next loads the SYN2 register, and the third loads the DLE register. Reading or loading the mode registers is done in a similar manner. The first write (or read) operation addresses Mode Register 1, and a subsequent operation addresses Mode Reqister 2. If more than the required number of accesses are made, the internal sequencer recycles to point at the first register. The pointers are reset to SYN1 Register and Mode Register 1 by a RESET input or by performing a "Read Command Register" operation, but are unaffected by any other read or write operation.

The 2651 register formats are summarized in Tables 5, 6, 7 and 8. Mode Registers 1 and 2 define the general operational characteristics of the PCI, while the Command Register controls the operation within this basic frame-work. The PCI indicates its status in the Status Register. These registers are cleared when a RESET input is applied.

Table 5 illustrates Mode Register 1. Bits MR11 and MR10 select the communication format and baud rate multiplier, 00 specifies synchronous mode and 1X multiplier. 1X, 16X, and 64X multipliers are programmable for asynchronous format. However, the multiplier in asynchronous format applies only if the external clock input option is selected by MR24 or MR25.

MR13 and MR12 select a character length of 5, 6, 7, or 8 bits. The character length does not include the parity bit, if programmed, and does not include the start and stop bits in asynchronous mode.

MR14 controls parity generation. If enabled, a parity bit is added to the transmitted character and the receiver performs a parity check on incoming data. MR15 selects odd or even parity when parity is enabled by MR14.

In asychronous mode, MR17 and MR16 select character framing of 1, 1.5, or 2 stop bits. (If 1X baud rate is programmed, 1.5 stop bits defaults to 1 stop bits on transmit). In synchronous mode, MR17 controls the number of SYN characters used to establish synchronization and for character fill when the transmitter is idle. SYN1 alone is used if MR17 = 1, and SYN1-SYN2 is used when MR17 = 0. If the transparent mode is specified by MR16, DLE-SYN1 is used for character fill, but the normal synchronization sequence is used.

Mode Register 2 (MR2)

Table 6 illustrates Mode Register 2. MR23, MR22, MR21, and MR20 control the frequency of the internal baud rate generator (BRG). Sixteen rates are selectable. When driven by a 5.0688 MHz input at the BRCLK input (pin 20), the BRG output has zero error except at 134.5, 2000, and 19,200 baud. which have errors of +0.016%, +0.235%, and +3.125% respectively.

MR25 and MR24 select either the BRG or the external inputs TxC and RxC as the clock source for the transmitter and receiver, respectively. If the BRG clock is selected. the baud rate factor in asynchronous mode is 16X regardless of the factor selected by MR11 and MR10. In addition, the corresponding clock pin provides an output at 1X the baud rate.

MR17	MR16	MR15	MR14	MR13	MR12	MR11	MR10
		Parity Type Parity Control Character Length		er Length	Mode and Baud Rate Factor		
ASYNCH: STOP 00 = INVALID 01 = 1 STOP BIT 10 = 11/2 STOP BIT 11 = 2 STOP BIT	uts	0 = ODD 1 = EVEN	0 = DISABLED 1 = ENABLED	01 = 6 10 = 7	5 BITS 6 BITS 7 BITS 8 BITS	10 = ASYNCHR	NOUS 1X RATE ONOUS 1X RATE ONOUS 16X RATE ONOUS 64X RATE
SYNCH: NUMBER OF SYN CHAR	SYNCH: TRANS- PARENCY CONTROL						
0 = DOUBLE SYN 1 = SINGLE SYN	0 = NORMAL 1 = TRANSPARENT						

Baud rate factor in asynchronous applies only if external clock is selected. Factor is 16X if internal clock is selected.

Table 5 MODE REGISTER 1 (MR1)

MR27	MR26	MR25	MR24	MR23	MR22	MR21	MR20
		Transmitter Clock	Receiver Clock		Baud Rate	Selection	
NOT USED		0 = EXTERNAL 1 = INTERNAL	0 = EXTERNAL 1 = INTERNAL	0001 = 0010 = 0011 = 0100 = 0101 = 0110 =	= 110 = 134.5 = 150 = 300	1000 = 1800 1001 = 2000 1010 = 2400 1011 = 3600 1100 = 4800 1101 = 7200 1111 = 19,20))))

Table 6 MODE REGISTER 2 (MR2)

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PRELIMINARY SPECIFICATION

Command Register (CR)

Table 7 illustrates Command Register. Bits CR0 (TxEN) and CR2 (RxEN) enable or disable the transmitter and receiver respectively. If the transmitter is disabled, it will complete the transmission of the character in the Transmit Shift Register (if any) prior to terminating operation. The TxD output will then remain in the marking state (high). If the receiver is disabled, it will terminate operation immediately. Any character being assembled will be neglected.

Bits CR1 (DTR) and CR5 (RTS) control the DTR and RTS outputs. Data at the outputs is the logical complement of the register data.

In asynchronous mode, setting CR3 will force and hold the TxD output low (spacing condition) at the end of the current transmitted character. Normal operation resumes when CR3 is cleared. The TxD line will go high for a least one bit time before beginning transmission of the next character in the Transmit Data Holding Register. In synchronous mode, setting CR3 causes the transmission of the DLE register contents prior to sending the character in the Transmit Data Holding Register. CR3 should be reset in response to the next TxRDY.

Setting CR4 causes the error flags in the Status Register (SR3, SR4, and SR5) to be cleared. This bit resets automatically.

The PCI can operate in one of four submodes within each major mode (synchronous or asynchronous). The operational sub-mode is determined by CR7 and CR6. CR7-CR6 = 00 is the normal mode, with the transmitter and receiver operating independently in accordance with the Mode and Status Register instructions.

In asynchronous mode, CR7-CR6 = 01 places the PCI in the Automatic Echo mode. Clocked, regenerated received data is automatically directed to the TxD line while normal receiver operation continues. The receiver must be enabled (CR2 = 1), but the

transmitter need not be enabled. CPU to receiver communications continues normally, but the CPU to transmitter link is disabled. Only the first character of a break condition is echoed. The TxD output will go high until the next valid start is detected. The following conditions are true while in Automatic Echo mode:

- Data assembled by the receiver is automatically placed in the Transmit Holding Register and retransmitted by the transmitter on the TxD output
- 2. Transmit clock = receive clock.
- 3. \overline{TxRDY} output = 1.
- The TxEMT/DSCHG pin will reflect only the data set change condition.
- 5. The TxEN command (CR0) is ignored.

In synchronous mode, CR7-CR6 = 01 places the PCI in the Automatic SYN/DLE Stripping mode. The exact action taken depends on the setting of bits MR17 and MR16:

- In the non-transparent, single SYN mode (MR17-MR16 = 10), characters in the data stream matching SYN1 are not transferred to the Receive Data Holding Register (RHR).
- In the non-transparent, double SYN mode (MR17-MR16 = 00), characters in the data stream matching SYN1, or SYN2 if immediately preceded by SYN1, are not transferred to the RHR. However, only the first SYN1 of an SYN1-SYN1 pair is stripped.
- In transparent mode (MR16 =1), characters in the data stream matching DLE, or SYN1 if immediately preceded by DLE, are not transferred to the RHR. However, only the first DLE of a DLE-DLE pair is stripped.

Note that Automatic Stripping mode does not affect the setting of the DLE Detect and SYN Detect status bits (SR3 and SR5).

Two diagnostic sub-modes can also be configured. In Local Loop Back mode (CR7-CR6 = 10), the following loops are connected internally:

- The transmitter output is connected to the receiver input.
- 2. DTR is connected to DCD and RTS is connected to CTS.

- 3. Receive clock = transmit clock.
- 4. The DTR, RTS and TxD outputs are held high.
- 5. The CTS, DCD, DSR and RxD inputs are ignored.

Additional requirements to operate in the Local Loop Back mode are that CR0 (TxEN), CR1 (DTR), and CR5 (RTS) must be set to 1. CR2 (RxEN) is ignored by the PCI.

The second diagnostic mode is the Remote Loop Back mode (CR7-CR6 = 11). In this mode:

- Data assembled by the receiver is automatically placed in the Transmit Holding Register and retransmitted by the transmitter on the TxD output.
- 2. Transmit clock = receive clock.
- 3. No data is sent to the local CPU, but the error status conditions (PE, OE, FE) are set.
- The RxRDY, TxRDY, and TxEMT/DSCHG outputs are held high.
- 5. CR1 (TxEN) is ignored.
- 6. All other signals operate normally.

Status Register

The data contained in the Status Register (as shown in Table 8) indicate receiver and transmitter conditions and modem/data set status.

SR0 is the Transmitter Ready (TxRDY) status bit. It. and its corresponding output, are valid only when the transmitter is enabled. If equal to 0, it indicates that the Transmit Data Holding Register has been loaded by the CPU and the data has not been transferred to the Transmit Shift Register. If set equal to 1, it indicates that the Holding Register is ready to accept data from the CPU. This bit is initially set when the Transmitter is enabled by CR0, unless a character has previously been loaded into the Holding Register. It is not set when the Automatic Echo or Remote Loop Back modes are programmed. When this bit is set, the TxRDY output pin is low. In the Automatic Echo and Remote Loop Back modes, the output is held high.

CR7	CR6	CR5	CR4	CR3	CR2	CR1	CR0
Operati	ng Mode	Request to Send	Reset Error		Receive Control (RxEN)	Data Terminal Ready	Transmit Control (TxEN)
01 = ASYNCI ECHO M SYNCH DLE ST 10 = LOCAL	L OPERATION H: AUTOMATIC MODE SYN AND/OR RIPPING MODE LOOP BACK E LOOP BACK	0 = FORCE RTS OUTPUT HIGH 1 = FORCE RTS OUTPUT LOW	ERROR FLAG	ASYNCH: FORCE BREAK 0 = NORMAL 1 = FORCE BREAK SYNCH: SEND DLE 0 = NORMAL 1 = SEND DLE	0 = DISABLE 1 = ENABLE	0 = FORCE DTR OUTPUT HIGH 1 = FORCE DTR OUTPUT LOW	0 = DISABLE 1 = ENABLE

Table 7 COMMAND REGISTER (CR)

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SR7	SR6	SR5	SR4	9R3	SR2	SR1	SRO
Data Set Ready	Data Carrier Detect	FE/SYN Detect	Overrrun	PE/DLE Detect	TxEMT/DSCHG	RxRDY	TxRDY
0 = DSR INPUT IS HIGH 1 = DSR INPUT IS LOW	0 = DCD INPUT IS HIGH 1 =DCD INPUT IS LOW	ASYNCH: 0 = NORMAL 1 = FRAMING ERROR SYNCH: 0 = NORMAL 1 = SYN CHAR DETECTED	0 = NORMAL 1 = OVERRUN ERROR	ASYNCH: 0 = NORMAL 1 = PARITY ERROR SYNCH: 0 = NORMAL 1 = PARITY ERROR OR DLE CHAR RECEIVED	0 = NORMAL 1 = CHANGE IN DSR OR DCD, OR TRANSMIT SHIFT REGIS- TER IS EMPTY	0 = RECEIVE HOLDING REG EMPTY 1 = RECEIVE HOLDING REG HAS DATA	0 = TRANSMIT HOLDING REG BUSY 1 = TRANSMIT HOLDING REG EMPTY

Table 8 STATUS REGISTER (SR)

SR1, the Receiver Ready (RxRDY) status bit, indicates the condition of the Receive Data Holding Register. If set, it indicates that a character has been loaded into the Holding Register from the Receive Shift Register and is ready to be read by the CPU. If equal to zero, there is no new character in the Holding Register. This bit is cleared when the CPU reads the Receive Data Holding Register or when the receiver is disabled by CR2. When set, the RxRDY output is low.

The TxEMT/DSCHG bit, SR2, when set, indicates either a change of state of the DSR or DCD inputs or that the Transmit Shift Register has completed transmission of a character and no new character has been loaded into the Transmit Data Holding Register. Note that in synchronous mode this bit will be set even though the appropriate "fill" character is transmitted. It is cleared when the transmitter is enabled by CR0 and does not indicate transmitter condition until at

least one character is transmitted. It is also cleared when the Status Register is read by the CPU. When SR2 is set, the TxEMT/-DSCHG output is low.

SR3, when set, indicates a received parity error when parity is enabled by MR14. In synchronous transparent mode (MR16 = 1), with parity disabled, it indicates that a character matching the DLE Register has been received. However, only the first DLE of two successive DLEs will set SR3. This bit is cleared when the receiver is disabled and by the Reset Error command. CR4.

The Overrun Error status bit, SR4, indicates that the previous character loaded into the Receive Holding Register was not read by the CPU at the time a new received character was transferred into it. This bit is cleared when the receiver is disabled and by the Reset Error command, CR4.

In asynchronous mode, bit SR5 signifies that the received character was not framed by the programmed number of stop bits. (If 1.5 stop bits are programmed, only the first stop bit is checked.) In synchronous nontransparent mode (MR16 = 0), it indicates receipt of the SYN1 character is single SYN mode or the SYN1-SYN2 pair in double SYN mode. In synchronous transparent mode (MR16 = 1), this bit is set upon detection of the initial synchronizing characters (SYN1 or SYN1-SYN2) and, after synchronization has been achieved, when a DLE-SYN1 pair is received. The bit is reset when the receiver is disabled, when the Reset Error command is given in asynchronous mode, and when the Status Register is read by the CPU in the synchronous mode.

 $\underline{SR6}$ and $\underline{SR7}$ reflect the conditions of the \overline{DCD} and \overline{DSR} inputs respectively. A low input sets its corresponding status bit and a high input clears it.

DC ELECTRICAL CHARACTERISTICS $T_A = 0$ °C to 70°C, $V_{CC} = 5.0$ V ± 5 % 4,5,6

PARAMETER		TEST CONDITIONS				
	r Aname I En	TEST CONDITIONS	Min	Тур	Max	UNIT
V _{IL} V _{IH}	Input voltage Low High			0.8 2.0		V
V _{OL} V _{OH}	Output voltage Low High	I _{OL} = 1.6mA I _{OH} = -100uA		0.25 2.8		V
l _{IL}	Input load current	V _{IN} = 0 to 5.5V		10		μА
Tristate I _{LH} I _{LL}	Output leakage current Data bus high Data bus low	$V_{O} = 4.0V$ $V_{O} = 0.45V$		10 10		μΑ
lcc	Power supply current			90		mA

PRELIMINARY SPECIFICATION

Manufacturer reserves the right to make design and process changes and improvements.

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PROGRAMMABLE COMMUNICATIONS INTERFACE (PCI)

PRELIMINARY SPECIFICATION

2651-I

CAPACITANCE TA = 25°C, VCC = 0V

				UNIT		
P.	ARAMETER	TEST CONDITIONS	Min	Тур	Max	ONII
Cin	Capacitance Input	(- 454) I-			20	pF
Соит	Output	fc = 1MHz Unmeasured pins tied to ground	į		20	
C _{I/O}	Input/Output	to ground			20	

PRELIMINARY SPECIFICATION

Manufacturer reserves the right to make design and process changes and improvements.

AC ELECTRICAL CHARACTERISTICS TA = 0°C to +70°C, Vcc = 5.0V ±5%4.5.6

		TEST CONDITIONS		LIMITS			
PARAMETER		TEST CONDITIONS	Min	Тур	Max	UNIT	
tres toe	Pulse width Reset Chip enable			1000 200		ns	
tas tah tcs tch tds tdb tdb trxs	Setup and hold time Address setup Address hold R/W control setup R/W control hold Data setup for write Data hold for write Rx data setup Rx data hold			10 10 10 10 150 80 150 280		ns	
t _{DD}	Data delay time for read Data bus floating time for read	$C_L = 100pF$ $C_L = 100pF$		180 70		ns ns	
fBRG fR/T	Input clock frequency Baud rate generator TxC or RxC		5.0637 dc	5.0688 1.0	5.0738	MH	
tBRH tBRL tR/TH tR/TL	Clock state Baud rate high Baud rate low TxC or RxC high (duty cycle) TxC or RxC low (duty cycle)		44% 44%	90 90 50% 50%	56% 56%	ns	
tTXD tTCS	TxD delay from falling edge of TxC Skew between TxD changing and falling edge of TxC output ⁸	C _L = 100pF C _L = 100pF	·	300 0		ns ns	

NOTES

- 1. Stresses above those listed under Absolute Maximum Ratings may cause permanent damage to the device. This is a stress rating only and functional operation of the device at these or at any other condition above those indicated in the operation section of this specification is not implied.
- 2. For operating at elevated temperatures, the device must be derated based on +150°C maximum junction temperature and thermal resistance of 60° C/W junction to ambient (IQ ceramic package).
- 3. This product includes circuitry specifically designed for the protection of its internal devices from the damaging effects of excessive static charge. Nonetheless, it is suggested that conventional precautions be taken to avoid applying any voltages larger than the rated maxima.
- 4. Parameters are valid over operating temperature range unless otherwise specified.
- All voltage measurements are referenced to ground. All time measurements are at the Voh, Vol., V_{IH} , V_{OL} , V_{IH} , V_{IH} , V_{OL} , V_{IH} , V_{OL} , V_{IH} , V_{OL} , V_{IH} , VIL levels as appropriate.
- Typical values are at +25°C, typical supply voltages and typical processing parameters.

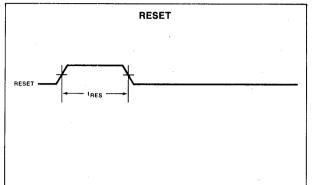
 TxRDY, RxRDY and TxEMT/DSCHG outputs are open drain.
- Parameter applies when internal transmitter clock is used.

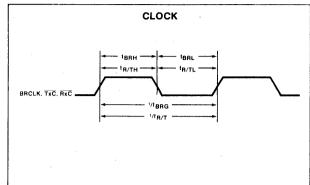
PRELIMINARY SPECIFICATION

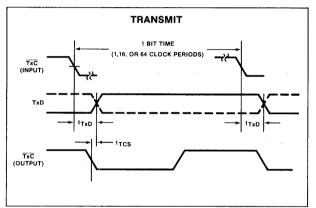
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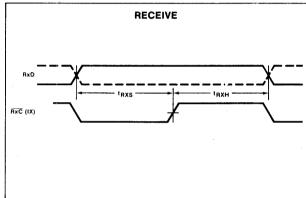
2651-I

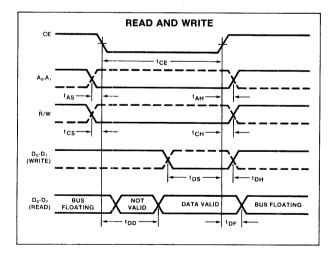
TIMING DIAGRAMS



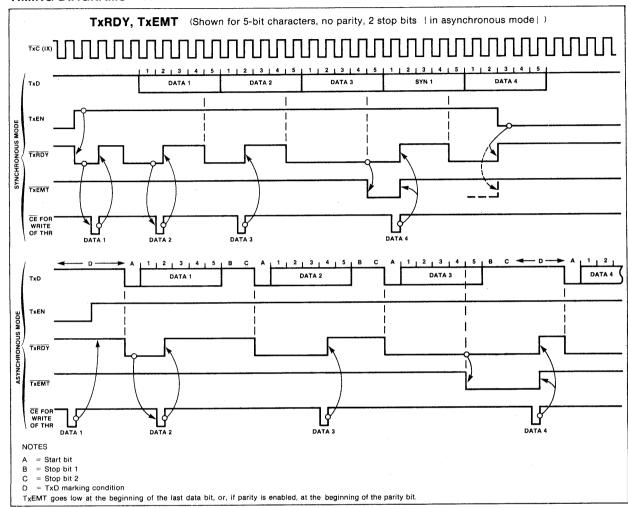




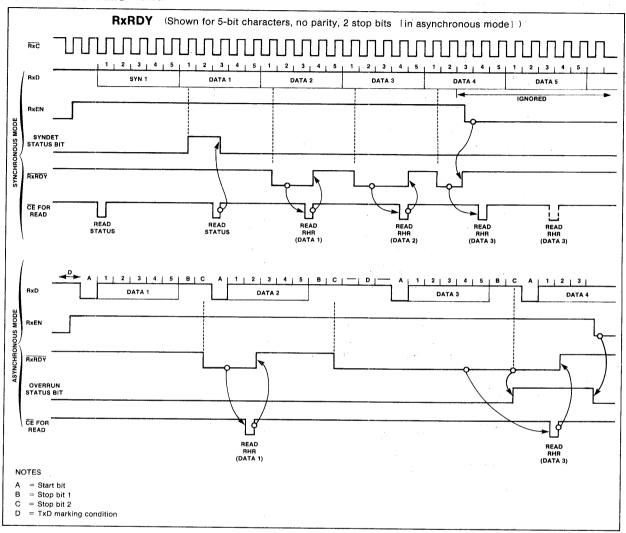




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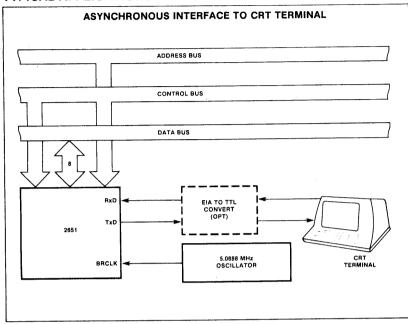


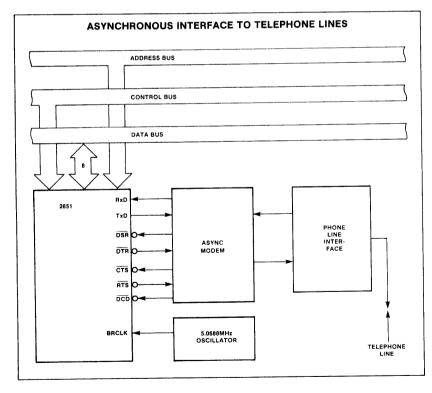
2651-1



2651-I

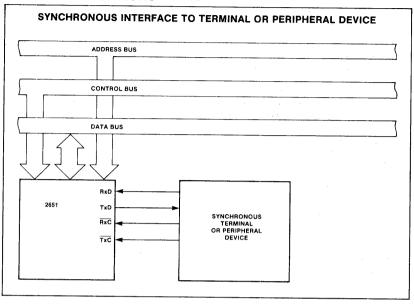
TYPICAL APPLICATIONS

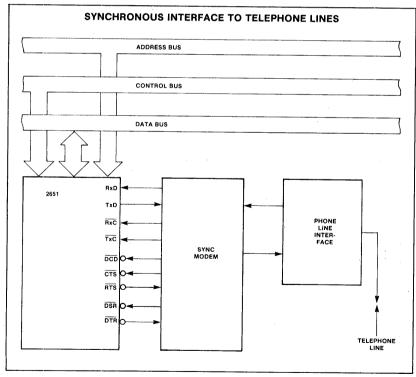




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TYPICAL APPLICATIONS (Cont'd)





2652-1

DESCRIPTION

The 2652 Multi-Protocol Communications Controller (MPCC) is a monolithic nchannel MOS LSI circuit that formats, transmits and receives synchronous serial data while supporting bit-oriented or byte control protocols. The chip is TTL compatible, operates from a single +5V supply, and can interface to a processor with an 8 or 16-bit bidirectional data bus.

FEATURES

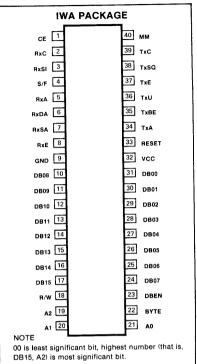
- . DC to 500K bps data rate
- Protocol management Bit-oriented protocols (BOP): SDLC, ADCCP, HDLC Byte-control protocols (BCP): BI-SYNC, **DDCMP**
- Programmable operation 8 or 16-bit tri-state data bus Protocol selection—BOP or BCP Error control-CRC or VRC or no error check Character length—1 to 8 bits for BOP or 5 to 8 bits for BCP SYNC or secondary station address comparison for BCP-BOP Idle transmission of SYNC/FLAG or MARK for BCP-BOP

- Automatic detection and generation of special BOP control sequences, i.e., FLAG. ABORT, GA
- Zero insertion and deletion for BOP
- Short character detection for last BOP data character
- SYNC generation, detection, and stripping for BCP
- Maintenance Mode for self-testing
- Common parameter control registers
- Independent status and data registers for receive and transmit
- Status indicator signals can be used as **CPU** interrupts
- TTL compatible
- 40-pin package
- Single +5V supply

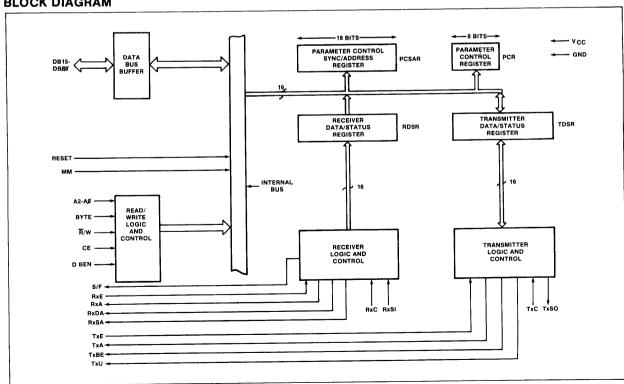
APPLICATIONS

- Intelligent terminals
- Line controllers
- **Network processors**
- Front end communications
- Remote data concentractors
- Communication test equipment
- Computer to computer links

PIN CONFIGURATION



BLOCK DIAGRAM



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PIN DESIGNATION

MNEMONIC	PIN NO.	TYPE	NAME AND FUNCTION
DB15-DB00	17-10 24-31	1/0	Data Bus: DB07-DB00 contain bidirectional data while DB15-DB08 contain control and status information to or from the processor. Corresponding bits of the high and low order bytes can be WIRE OR'ed onto an 8-bit data bus.
A2-A0	19-21	ı	Address Bus: A2-A0 select internal registers. The four 16-bit registers can be addressed on a word or byte basis. See Register Address section.
BYTE	22	1	Byte: Single byte (8 bit) data bus transfers are specified when this input is high. A low level specifies 16 bit data bus transfers.
CE	1	1	Chip Enable: A high input permits a data bus operation when DBEN is activated.
R/W	18	I	Read/Write: R/W controls the direction of data bus transfer. When high, the data is to be loaded into the addressed register. A low input causes the contents of the addressed register to be presented on the data bus.
DBEN	23	ı	Data Bus Enable: After A2-A0, CE, BYTE and R/W are set up, DBEN may be strobed. During a read, the tri-state data bus (DB) is enabled with information for the processor. During a write, the stable data is loaded into the addressed register and TxBE will be reset if TDSR was addressed.
RESET	33	I	Reset: A high level initializes all internal registers and timing.
MM	40	I	Maintenance Mode: MM internally gates TxSO back to RxSI and TxC to RxC for off line diagnostic purposes. The RxC input is disabled when MM is asserted.
RxE	8	I	Receiver Enable: A high level input permits the processing of RxSI data. A low level disables the receiver logic and initializes all receiver registers and timing.
RxA	5	0	Receiver Active: RxA is asserted when the first data character of a message is ready for the processor. In the BOP mode this character is the address. The received address must match the secondary station address if the MPCC is a secondary station. In BCP mode, if strip-SYNC (PCSAR ₁₃) is set, the first non-SYNC character is the first data character; if strip-SYNC is zero, the character following the second SYNC is the first data character. In the BOP mode, the closing FLAG resets RxA. In the BCP mode, RxA is reset by a low level at RxE.
RxDA*	6	0	Receiver Data Available: RxDA is asserted when an assembled character is in RDSR _L and is ready to be presented to the processor. This output is reset when RDSR _L is read.
RxC	2	i	Receiver Clock: RxC(1X) provides timing for the receiver logic. The positive going edge shifts serial data into the RxSR from RxSI.
S/F	4	0	SYNC/FLAG: S/F is asserted for one RxC clock time when a SYNC or FLAG character is detected.
RxSA*	7	0	Receiver Status Available: RxSA is asserted when there is a zero to one transition of any bit in RDSR _H except for RSOM. It is cleared when RDSR _H is read.
RxSI	3	l i	Receiver Serial Input: RxSI is the received serial data. Mark = '1', space = '0'.
TxÉ	37	l	Transmitter Enable: A high level input enables the transmitter data path between TDSR _L and TxS0. At the end of a message, a low level input causes $TxS0 = 1$ (mark) and $TxA = 0$ after the closing FLAG (BOP) or last character (BCP) is output on TxS0.
TxA	34	0	Transmitter Active: TxA is asserted when TxE is high and TSOM (TDSR ₈) is set. This output will reset when TxE is low and the closing FLAG (BOP) or last character (BCP) has been output on TxS0.
TxBE*	35	0	Transmitter Buffer Empty: TxBE is asserted when the TDSR is ready to be loaded with new control information or data. The processor should respond by loading the TDSR which resets TxBE.
TxU*	36	0	Transmitter Underrun: TxU is asserted during a transmit sequence when the service of TxBE has been delayed for more than one character time. This indicates the processor is not keeping up with the transmitter (TxS0 depends on PCSAR ₁₁). TxU is reset by RESET or setting of TSOM (TDSR ₈).
ГхC	39	1	Transmitter Clock: TxC (1X) provides timing for the transmitter logic. The positive going edge shifts data out of the TxSR to TxS0.
ΓxSO	38	0	Transmitter Serial Output. TxS0 is the transmitted serial data. Mark = '1', space = '0'.
Vcc	32	1	+5V: Power supply.
GND	9	1	Ground: 0V reference ground.

^{*}Indicates possible interrupt signal.

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	REGISTERS	NO. OF BITS	DESCRIPTION*
Addressat PCSAR PCR		16 8	PCSARH and PCR contain parameters common to the receiver and transmitter. PCSARL contains a programmable SYNC character (BCP) or secondary station address (BOP).
RDSR	Receive Data/Status Register	16	RDSR _H contains receiver status information. RDSR _L = RxDB contains the received assembled character.
TDSR	Transmit Data/Status Register	16	$TDSR_H$ contains transmitter command and status information. $TDSR_L = TxDB$ contains the received assembled character.
Internal CCSR HSR RXSR TXSR RXCRC TXCRC	Control Character Shift Register Holding Shift Register Receiver Shift Register Transmitter Shift Register Receiver CRC Accumulation Register Transmitter CRC Generation Register	8 16 8 8 16 16	These registers are used for character assembly (CCSR, HSR, RxSR), disassembly (TxSR), and CRC accumulation/generation (RxCRC, TxCRC).

NOTE

*H = High byte - bits 15-8

L = Low byte - bits 7-0

Table 1 GLOSSARY

CHARACTER	DESCRIPTION
FCS	Frame Check Sequence is transmitted/received as 16 bits following the last data character of a BOP message and is usually CRC-CCITT (X16 + X12 + X5 +1) with dividend preset to 1's.
BCC	Block Check Character is transmitted/received as two successive characters following the last data character of a BCP message. Either CRC-16 (X¹6 + X¹5 + X² + 1) with dividend preset to 0's or LRC (X²8 + 1) as computed by the processor, is polynomial. CRC-16 is used with 8-bit EBC DIC. LRC is used with 7-bit ASCII in conjunction with VRC. The CRC-16 is computed on all characters beginning with the first nonsync character at the start of the message.

Table 2 ERROR CONTROL

FUNCTIONAL DESCRIPTION

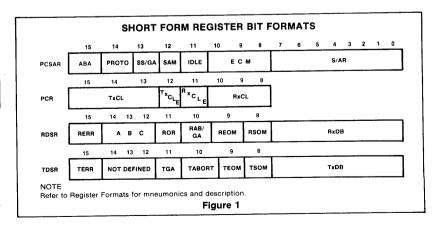
The MPCC can be functionally partitioned into receiver logic, transmitter logic, registers that can be read or loaded by the processor, and data bus control circuitry. The MPCC block diagram is shown in Figure 1 while the receiver and transmitter data paths are depicted in Figures 2 and 3.

OPERATION	BIT PATTERN	FUNCTION
ВОР		_
FLAG	01111110	Frame message
ABORT	11111111 generation	Terminate communication
	01111111 detection	
l GA	01111111	Terminate loop mode repeater
	1	function
Address	(PCSAR _L)1	Secondary station address
BCP		-
SYNC	(PCSAR _L) or (TxDB) ² generation	Frame message

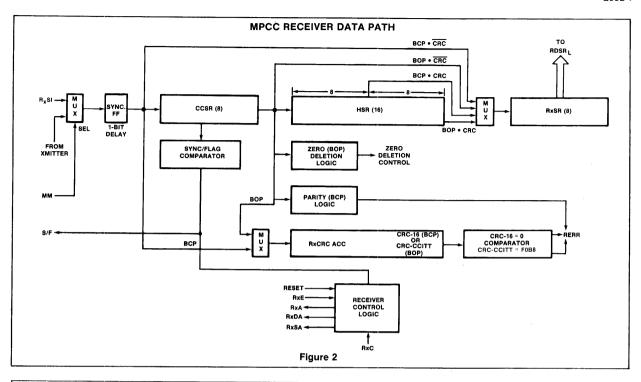
NOTES

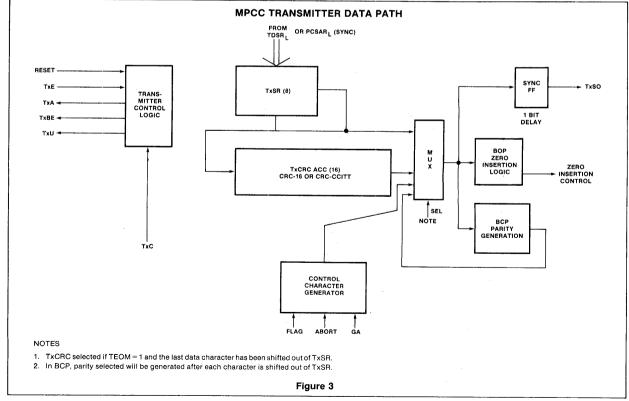
- 2. For IDLE = 0 or 1 respectively

Table 3 SPECIAL CHARACTERS



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RECEIVER OPERATION General

After initializing the parameter control registers (PCSAR and PCR), the RxE input must be set high to enable the receiver data path. The serial data on the RxSI is synchronized and shifted into an 8-bit Control Character Shift Register (CCSR) on the rising edge of RxC. A comparison between CCSR contents and the FLAG (BOP) or SYNC (BCP) character is made until a match is found. At that time, the S/F output is asserted for one RxC time and the 16-bit Holding Shift Register (HSR) is enabled. The receiver then operates as described below.

BOP Operation

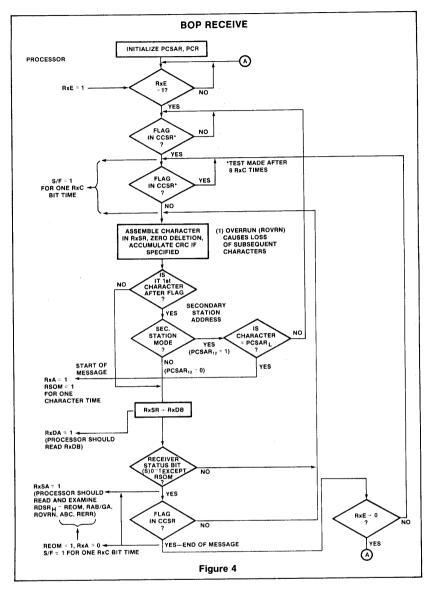
A flow chart of receiver operation in BOP mode appears in Figure 4. Zero deletion (after five ones are received) is implemented on the received serial data so that a data character will not be interpreted as a FLAG, ABORT, or GA, Bits following the FLAG are shifted through the CCSR, HSR, and into the Receiver Shift Register (RxSR). A character will be assembled in the RxSR and transferred to the RDSRL for presentation to the processor. At that time the RxDA output will be asserted and the processor must take the character before the next character is assembled in the RxSR. If not, an overrun (RDSR₁₁ = 1) will occur and succeeding characters will be lost.

The first character following the FLAG is the secondary station address. If the MPCC is a secondary station (PCSAR $_{12} = 1$), the contents of RxSR are compared with the address stored in PCSAR $_{L}$. A match indicates the forthcoming message is intended for the station; the RxA output is asserted, the character is loaded into RDSR $_{L}$, RxDA is asserted and the Receive Start of Message bit (RSOM) is set. No match indicates that another station is being addressed and the receiver searches for the next FLAG.

If the MPCC is receiving data from a secondary station (PCSAR₁₂ = 0), no secondary address check is made; RxA is asserted and RSOM is set once the first non-FLAG character has been loaded into RDSR_L and RxDA has been asserted.

When the 8 bits following the address character have been loaded into RDSR_L and RxDA has been asserted, RSOM will be cleared. The processor should read this 8-bit character and interpret it as the Control field.

Received serial data that follows is read and interpreted as the Information field by the processor. It will be assembled into character lengths as specified by PCR₈₋₁₀. As



before, RxDA is asserted each time a character has been transferred into RDSR_L and is cleared when RDSR_L is read by the processor. RxSA is asserted when there is a zero to one transition for any bit in RDSR_H except for RSOM. RxSA and all bits in RDSR_H except RSOM are cleared when RDSR_H is read. The processor should check RDSR₉₋₁₅ each time RxSA is asserted. If RDSR₉ is set, then RDSR₁₂₋₁₅ should be examined.

Receiver character length may be changed dynamically in response to RxDA: read the character in RxDB and write the new character length into RxCL. The character length will be changed on the next RxC positive going edge.

The CRC-CCITT, if specified by PCSAR₈₋₁₀, is accumulated in RxCRC on each character following the FLAG. When the closing FLAG is detected in the CCSR, the received CRC is in the 16-bit HSR. At that time, the Receive End of Message bit (REOM) will be set; RxSA and RxDA will be asserted. The processor should read the last data character in RDSR₂ and the receiver status in RDSR₉₋₁₅. If RDSR₁₅ = 1, there has been a

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PRELIMINARY SPECIFICATION

transmission error; the accumulated CRC-CCITT is incorrect. If RDSR₁₂₋₁₄ \neq 0, the last data character is not of prescribed length. Neither the received CRC nor closing FLAG are presented to the processor. The processor may drop RxE or leave it active at the end of the received message.

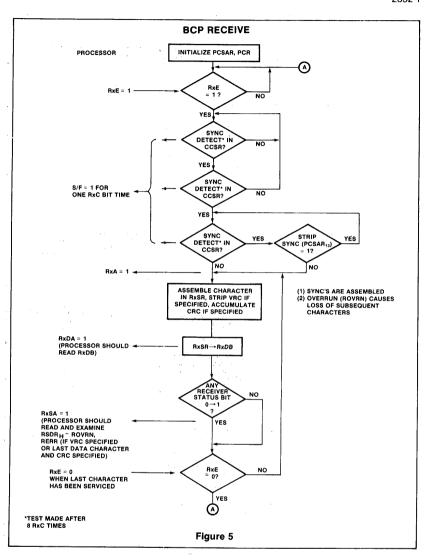
BCP Operation

The operation of the receiver in BCP mode is shown in Figure 5. The receiver initially searches for two successive SYNC characters. of length specified by PCR₈₋₁₀, that match the contents of PCSARL. The next non-SYNC character or next SYNC character if stripping is not specified (PCSAR₁₃ = 0), causes RxA to be asserted and enables the receiver data path from CCSR through HSRL to RxSR. All characters following the first non-SYNC are assembled in RxSR and loaded into RDSRL. RxDA is active when a character is available in RDSRL. RxSA is active on a 0 to 1 transition of any bit in RDSR_H. The signals are cleared when RDSRL or RDSRH are read respectively.

If CRC-16 error control is specified by PCSAR₈₋₁₀, the processor must determine the last character received prior to the CRC field. When that character is loaded into RDSRL and RxDA is asserted, the received CRC will be in CCSR and HSRL. To check for a transmission error, the processor must read the receiver status (RDSRH) and examine RDSR₁₅. This bit will be set for one character time if an error free message has been received. If RDSR₁₅ = 0, the CRC-16 is in error. Note that this bit should be examined only at the end of a message and that the accumulated CRC will include all characters starting with the first non-SYNC character at the start of the message. In particular, SYNC's in the middle of a message, DLE characters, and the first SOH or STX after line turn around are subject to CRC.

If VRC had been selected for error control, parity (odd or even) is regenerated on each character and check with the parity bit received. A discrepancy causes RDSR₁₅ to be set and RxSA to be asserted. This must be sensed by the processor. The received parity bit is stripped before the character is presented to the processor. The processor should compute and check LRC if required.

When the processor has read the last character of the message, it should drop RxE which disables the receiver logic and initializes all receiver registers and timing.



TRANSMITTER OPERATION General

After the parameter control register (PCSAR and PCR) have been initialized, TxE must be set high to enable the transmitter data path. TxSO is held to mark until TSOM (TDSR₈) is set. Then, transmitter operation depends on protocol mode.

BOP Operation

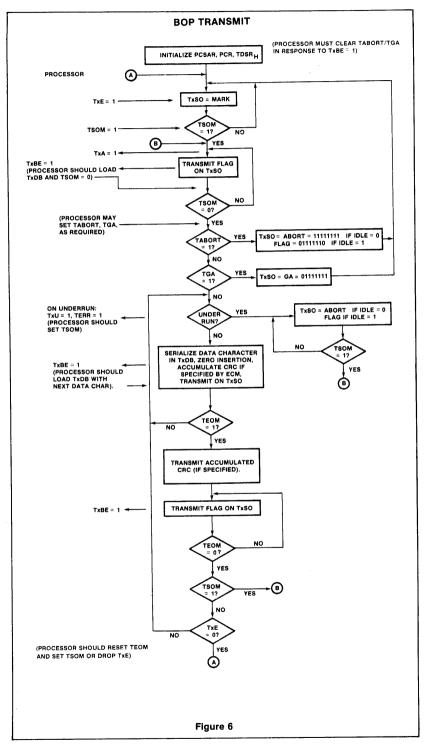
Transmitter operation for BOP is shown in Figure 6. A FLAG is sent when the processor sets the Transmit Start of Message bit (TSOM). The FLAG is used to synchronize the message that follows. TxA will be asserted after TSOM is set. When TxBE is asserted by the MPCC, the processor should load

TDSR_L with the first character of the message. TSOM should be cleared at the same time TDSR_L is loaded (16-bit data bus) or immediately thereafter (8-bit data bus). FLAGs are sent as long as TSOM = 1.

All succeeding characters are loaded into TDSR_L by the processor when TxBE =1. Each character is serialized in TxSR and transmitted on TxSO. Internal zero insertion logic stuffs a "0" into the serial bit stream after five successive "1s" are sent. This insures a data character will not match a FLAG, ABORT, or GA reserved control character. As each character is transmitted, the Frame Check Sequence (FCS) is generated as specified by Error Control Mode (PCSAR₈₋₁₀). The FCS should be the

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CRC-CCITT polynomial)X16 + X12 + X5 + 1) preset to 1s. If an underrun occurs (processor is not keeping up with the transmitter), TxU and TERR (TDSR₁₅) will be asserted with ABORT or FLAG used as the TxSO line fill depending on the state of IDLE (PCSAR₁₁). The processor must set TSOM and retransmit the message to recover.

A residual character of 3 to 7 bits may be transmitted at the end of BOP information field to make sure that field is a multiple of 8 bits. In response to TxBE, write the residual character length into TxCL and load TxDB with the residual character. Dynamic alteration of character length should be done in exactly the same sequence.

After the last data character has been loaded into TDSRL and sent to TxSR (TxBE = 1), the processor should set TEOM (TDSR9). The MPCC will finish transmitting the last character followed by the FCS and the closing FLAG. The processor should clear TEOM and drop TxE when the next TxBE is asserted. This corresponds to the start of closing FLAG transmission. When TxE has been dropped, TxA will be low 11/2 bit times after the last bit of the closing FLAG has been transmitted. TxSO will be marked after the closing FLAG has been transmitted.

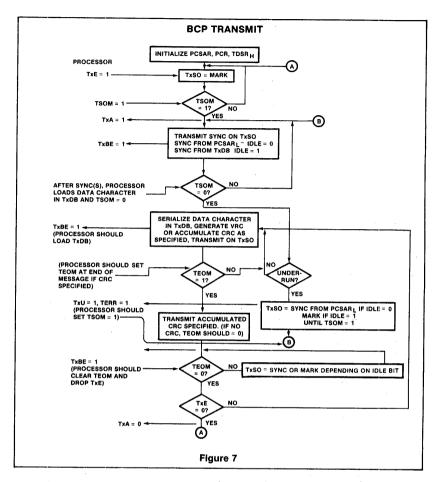
If TxE and TEOM are high, the transmitter continues to send FLAGs. The processor may initiate the next message by resetting TEOM and setting TSOM, or by loading TDSR_L with a data character and then simply resetting TEOM (without setting TSOM).

BCP Operation

Transmitter operation for BCP mode is shown in Figure 7. If TxE is high, TxA will be asserted when TSOM = 1. At that time SYNC characters are sent from PCSARL or TDSRL (IDLE = 0 or 1) as long as TSOM = 1. TxBE is asserted at the start of transmission of the first SYNC character. For more than one SYNC, the processor should reassert TSOM in response to the assertion of TxBE. When TSOM = 0 transmission is from TDSRL, which must be loaded with characters from the processor each time TxBE is asserted. If this loading is delayed for more than one character time, an underrun results: TxU and TERR are asserted and the TxSO line fill depend on IDLE (PCSAR11). The processor must set TSOM and retransmit the message to recover.

CRC-16, if specified by PCSAR₈₋₁₀, is generated on each character transmitted from TDSR_L when TSOM = 0. The processor must set TEOM = 1 after the last data character has been sent to TxSR (TxBE = 1). The MPCC will finish transmitting the last data

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bits. If a residual character is to be transmitted, TxCL should be changed to the residual character length prior to transmission of that character.

DATA BUS CONTROL

The processor must set up the MPCC register address (A2-A0), chip enable (CE), byte select (BYTE), and read/write (R/W) inputs before each data bus transfer operation.

During a read operation $(\bar{R}/W=0)$, the leading edge of DBEN will initiate an MPCC read cycle. The addressed register will place its contents on the data bus. If BYTE=1, the 8-bit byte is placed on DB15-08 or DB07-00 depending on the H/L status of the register addressed. Unused bits in RDSRL are zero. If BYTE=0, all 16 bits (DB15-00) contain MPCC information. The trailing edge of DBEN will reset RxDA and/or RxSA if RDSR_H or RDSR_L is addressed respectively.

DBEN acts as the enable and strobe so that the MPCC will not begin its internal read cycle until DBEN is asserted.

During a write operation $(\bar{R}/W=1)$, data must be stable on DB₁₅₋₀₈ and/or DB₀₇₋₀₀ prior to the leading edge of DBEN. The stable data is strobed into the addressed register by DBEN. TxBE will be cleared if the addressed register was TDSR_H or TDSR_L.

character and the CRC-16 field before sending SYNC characters which are transmitted as long as TEOM = 1. If SYNCs are not desired after CRC-16 transmission, the processor should clear TEOM and lower TxE when the TxBE corresponding to the start of CRC-16 transmission is asserted. When TEOM = 0, the line is marked and a new message may be initiated by setting TxE and TSOM.

If LRC is required, it must be generated by the processor and transmitted after the last data character. TEOM should not be set under this condition. If VRC is specified, it is generated on each data character and the data character length must not exceed 7 bits. For software LRC or CRC TEOM should be set only if SYNCs are required at the end of the message block.

SPECIAL CASE

The capability to transmit 16 spaces is provided for line turnaround in half duplex mode or for a control recovery situation. This is achieved by setting TSOM and TEOM, clearing TEOM when TxBE = 1, and proceeding as required.

PROGRAMMING

Prior to initiating data transmission or reception, PCSAR and PCR must be loaded with control information from the processor. The contents of these registers (see Register Format section) will configure the MPCC for the user's specific data communication environment. These registers should be loaded during power-on initialization and after a reset operation. They can be changed at any time that the respective transmitter or receiver is disabled.

The default value for all registers is zero. This corresponds to BOP, primary station mode, 8-bit character length, FCS = CRC-CCITT preset to 1s.

For BOP mode the character length register (PCR) may be set to the desired values during system initialization. The address and control fields will automatically be 8-

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MULTI-PROTOCOL COMMUNICATIONS CONTROLLER (MPCC)

PRELIMINARY SPECIFICATION

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	A2	A1	Α0	REGISTER
BYTE = 0	16-BIT DATA	BUS = DB ₁₅ -	DB 00	
	0	0	X	RDSR
	0	1	X	TDSR
	1	0	X	PCSAR
	1	1	x	PCR*
BYTE = 1	YTE = 1 8-BIT DATA BUS = DB ₇₋₀ or DB ₁₅₋₈ **			
	0	0 .	Ö	RDSRL
	0	0	1	RDSRH
	0	1	0	TDSRL
	0	1	1	TDSRH
	1	0	0	PCSARL
	1	0	1	PCSARH
ł	1	1	0	PCR _L *
	1	1	1	PCRH

NOTES

- * PCR lower byte does not exist. It will be all "0"s when read.
- ** Corresponding high and low order pins should be tied together.

Table 4 MPCC REGISTER ADDRESSING

BIT	NAME	MODE				FUNCTION			
00-07	Not Defined								
08-10	RxCL	BOP/BCP	Receiver RxCLE =		r Length	is loaded by the processor depending on RxBC when			
			10	9	8	Char. length (bits)			
			0	0	0	8			
	1 1		0	0	1	1			
			0	1	0	2			
1			0	1	1	3			
			1	0	0	4			
			1	0	1	5			
			1	1	0	6			
	1		1	1	1	7			
11	RxCLE	BOP/BCP	Receiver remaining	Receiver Character Length Enable should be zero when the processor loads RxCL. The remaining bits of PCR are not affected during loading.					
12	TxCLE	BOP/BCP	Transmitter Character Length Enable should be zero when the processor loads TxCL. The remaining bits of PCR are not affected during loading.						
13-15	TxCL	BOP/BCP				th is loaded by the processor when TxCLE = 0. Character nat is identical to RxCL.			

Table 5 PARAMETER CONTROL REGISTER (PCR)-(R/W)

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BIT	NAME	MODE			FUNC	TION		
00-07	S/AR	BOP BCP	SYNC/ADDRESS Register. (secondary station. The conte FLAG character to determine SYNC character is loaded intransmit bit synchronization	nts of the e if the o this re	is regi messa egister	ster is c ge is m by the	ompared with eant for this processor. It	h the first received non- s station. is used for receive and
08-10	ECM	BOP/BCP	Error Control Mode	10	9	8	Mode	Char. length
			CRC-CCITT preset to 1's CRC-CCITT preset to 0's Not used CRC-16 preset to 0's VRC old VRC even Not used No error control ECM should be loaded by the are idle.					
11	IDLE		Determines line fill character and TERR set) and transmiss	to be u	sed if	transmi	tter underru	n occurs (TxU asserted
		ВОР	IDLE = 0, transmit ABORT c IDLE = 1, transmit FLAG cha IDLE = 0 transmit initial SYN S/AR. IDLE = 1 transmit initial SYN run.	haracte racters C char	rs dur durin acters	ing und g under and un	errun and w run and whe derrun line f	hen TABORT = 1. en TABORT = 1. fill characters from the
12	SAM	ВОР	Secondary Address Mode = automatic recognition of the the processor must load the SAM = 0 inhibits the received the receiver after the first no	receive second second	d seco ary ad ary ad	ndary s dress ir dress c	tation addre: nto TxDB. omparison w	ss. When transmitting,
13	SS/GA	ВОР	Strip SYNC/Go Ahead. Operation depends on mode. For loop mode only. SS/GA = 1 permits GA character to terminate a received message. When a GA is detected REOM and RAB/GA will be set and the processor should terminate the repeater function. SS/GA = 0 permits only a FLAG or ABORT character to terminate a message. SS/GA = 1, causes the receiver to strip SYNC's immediately following the first two SYNC's detected. SYNC's in the middle of a message will not be stripped. SS/GA = 0, presents any SYNC's after the initial two SYNC's to the processor.					
14	PROTO	BOP BCP	Determines MPCC Protocol r PROTO = 0 BOP PROTO = 1 BCP					
15	APA	ВОР	All Parties Address. If this bit of '111111111' as well as the n	s set, th	econd	iver dat ary sta	a path is ena tion address	bled by an address field

Table 6 PARAMETER CONTROL SYNC/ADDRESS REGISTER (PCSAR)-(R/W)

BIT	NAME	MODE	FUNCTION
00-07	TxDB	BOP/BCP	Transmit Data Buffer. Contains processor loaded characters to be serialized in TxSR and transmitted on TxSO.
08	TSOM	ВОР	Transmitter Start of Message. Set by the processor to initiate message transmission provided TxE = 1. TSOM = 1 generates FLAGs. When TSOM = 0 transmission is from TxDB and FCS generation begins. FCS, as specified by PCSAR ₈₋₁₀ , should be CRC-CCITT preset to 1's. TSOM = 1 generates SYNCs from PCSAR _L or transmits from TxDB for IDLE = 0 or 1 respectively. When TSOM = 0 transmission is from TxDB and CRC generation (if specified) begins.

Table 7 TRANSMIT DATA/STATUS REGISTER (TDSR) (R/W except TDSR 15)

MULTI-PROTOCOL COMMUNICATIONS CONTROLLER (MPCC)

PRELIMINARY SPECIFICATION

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BIT	NAME	MODE	FUNCTION
09	TEOM	ВОР	Transmit End of Message. Used to terminate a transmitted message when CRC error checking is used. TEOM = 1 causes the FCS and the closing FLAG to be transmitted following the transmission of the data character in TxSR. FLAGs are transmitted until TEOM = 0. ABORT or GA are transmitted if TABORT or TGA are set when TEOM = 1.
		ВСР	TEOM = 1 causes CRC-16 to be transmitted (if selected) followed by SYNCs from PCSAR _L or TxDB (IDLE = 0 or 1). Clearing TEOM prior to the end of CRC-16 transmission (when TxBE = 1) causes TxSO to be marked following the CRC-16. TxE must be dropped before a new message can be initiated. If CRC is not selected, TEOM should not be set.
10	TABORT	ВОР	Transmitter Abort = 1 will cause ABORT or FLAG to be sent (IDLE = 0 or 1) after the current character is transmitted. (ABORT = 11111111)
11	TGA	ВОР	Transmit Go Ahead (GA) instead of FLAG when TEOM = 1. This facilitates repeater termination in loop mode. (GA = 01111111)
12-14	Not Defined		
15	TERR	Read only BOP BCP	Transmitter Error = 1 indicates the TxDB has not been loaded in time to maintain continous transmission. TxU will be asserted to inform the processor of this condition. TERR is cleared by setting TSOM. ABORT's or FLAG's are sent as fill characters (IDLE = 0 or 1) SYNC's or MARK's are sent as fill characters (IDLE = 0 or 1). For IDLE = 1 the last character before underrun is not valid.

Table 7 TRANSMIT DATA/STATUS REGISTER (TDSR) (R/W except TDSR 15) (Cont'd)

BIT	NAME	MODE	FUNCTION
00-07	RxDB	BOP/BCP	Receiver Data Buffer. Contains assembled characters from the RxSR. If VRC is specified, the parity bit is stripped.
08	RSOM	ВОР	Receiver Start of Message = 1 when a FLAG followed by a non-FLAG has been received and the latter character matches the secondary station address if SAM = 1. RxA will be asserted when RSOM = 1. RSOM resets itself after one character time and has no effect on RxSA.
09	REOM	BOP	Receiver End of Message = 1 when the closing FLAG is detected and the last data character is loaded into RxDB or when an ABORT/GA character is received. REOM is cleared on reading RDSRH, reset operation, or dropping of RxE.
10	RAB/GA	ВОР	Received ABORT or GA character = 1 when the receiver senses an ABORT character if SS/GA = 0 or a GA character if SS/GA = 1. RAB/GA is cleared on reading RDSR _H , reset operation, or dropping of RxE. A received ABORT inhibits RxDA.
11	ROR	BOP/BCP	Receiver Overrun = 1 indicates the processor has not read the last character in the RxDB within one character time. Subsequent characters will be lost. ROR is cleared on reading RDSR _H , reset operation, or dropping or RxE.
12-14	ABC	ВОР .	Assembled Bit Count. Specifies the number of bits in the last received data character of a message and should be examined by the processor when REOM = 1 (RxDA and RxSA asserted). ABC = 0 indicates the message was terminated (by a FLAG or GA) on a character boundary as specified by PCSCR ₈₋₁₀ . Otherwise, ABC = number of bits in the last data character. ABC is cleared when RDSR _H is read, reset operation, or dropping RxE. The residual character is right justified in RDSR _L .
15	RERR	BOP BCP	Receiver Error indicator should be examined by the processor when REOM = 1 in BOP, or when the processor determines the last data character of the message in BCP with CRC or when RxSA is set in BCP with VRC. CRC-CCITT preset to 1's should be specified by PCSAR ₈₋₁₀ : RERR = 1 indicates FCS error (CRC ≠ F0B8) RERR = 0 indicates FCS received correctly (CRC = F0B8) CRC-16 preset to 0's on 8-bit data characters specified by PCSAR ₈₋₁₀ : RERR = 1 indicates CRC-16 received correctly (CRC-0). RERR = 0 indicates CRC-16 error (CRC ≠ 0) VRC specified by PCSAR ₈₋₁₀ : RERR = 1 indicates VRC error RERR = 0 indicates VRC is correct

Table 8 RECEIVER DATA/STATUS REGISTER (RDSR)-(Read Only)

2652-1

ABSOLUTE MAXIMUM RATINGS1

	PARAMETER	RATING	UNIT
T _A T _{STG}	Operating ambient temperature ² Storage temperature Input or output voltages	0 to +70 -65 to +150	°C °C
Vcc	with respect to GND ³ With respect to GND	-0.3 to +15 -0.3 to +7	V V

DC ELECTRICAL CHARACTERISTICS $T_A = 0$ °C to +70°C, $V_{CC} = +5V \pm 5\%4,5,6$

PARAMETER		PARAMETER TEST CONDITIONS		LIMITS	LIMITS		
Input voltage VIL Low VIH High Output voltage VOL Low VOH High ICC Power supply current Leakage current IIL Input	1201 CONDITIONS	Min Typ Max		Max	UNIT		
V _{IL} Low V _{IH} High			2.0		0.8	V	
V _{OL} V _{OH}	Low	I _{OL} = 1.6mA I _{OH} = -100 <i>µ</i> A	2.4		0.4	V	
lcc	Power supply current	V _{CC} = 5.25V, T _A = 0°C			150	mA	
lıL loL		$V_{IN} = 0 \text{ to } 5.25V$ $V_{OUT} = 0 \text{ to } 5.25V$			10 10	μА	
Cin Cout	Capacitance Input Output	V _{IN} = 0V, f = 1MHz V _{OUT} = 0V, f = 1MHz			20 20	pF	

AC ELECTRICAL CHARACTERISTICS $T_A=25^{\circ}$ C, $V_{CC}=5$ V \pm 5%, AC timing indicated is with outputs unloaded.4.5.6

	PARAMETER		LIMITS		
	- Allameter	Min	Тур	Max	UNIT
	Setup and hold time				ns
tacs	Address/control setup	50			113
tach	Address/control hold	lo			İ
tos	Data bus setup (write)	50			
tDH	Data bus hold (write)	0			
t _{RXS}	Receiver serial data transfer	150			
t _{RxH}	Receive serial data hold	150			
	Pulse width				ns
tres	RESET	250			
toben	DBEN	250			
	Delay time				ns
t_{DD}	Data bus (read)			200	
t _{TxD}	Transmit serial data			300	
tor	Data bus float time (read)			150	ns
f	Clock (RxC, TxC) frequency			500	kHz
t _{CLK1}	Clock high	1000			ns
t _{CLK0}	Clock low	1000			ns

MULTI-PROTOCOL COMMUNICATIONS CONTROLLER (MPCC)

PRELIMINARY SPECIFICATION

2652-1

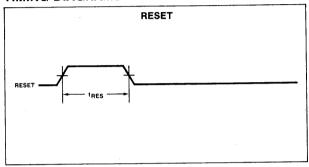
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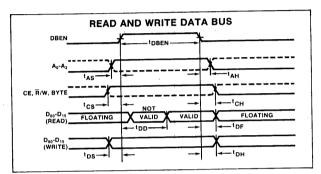
- Stresses above those listed under "Absolute Maximum Ratings" may cause permanent damage to the device. This is a stress rating only and functional operation of the device at these or at any other condition above those indicated in the operation sections of this specification is not implied.
- For operating at elevated temperatures the device must be derated based on +150°C
 maximum junction temperature and thermal resistance of 60°C/W junction to ambient
 (IQ ceramic package).
- 3. This product includes circuitry specifically designed for the protection of its internal devices from the damaging effects of excessive static charge. Nonetheless, it is suggested that conventional precautions be taken to avoid applying any voltages larger than the rated maxima.
- 4. Parameters are valid over operating temperature range unless otherwise specified.
- All voltage measurements are referenced to ground. All time measurements are at the VOH, VOL, VIH, or VIL levels as appropriate.
- Typical values are at +25°C, nominal supply voltages, and nominal processing parameters.

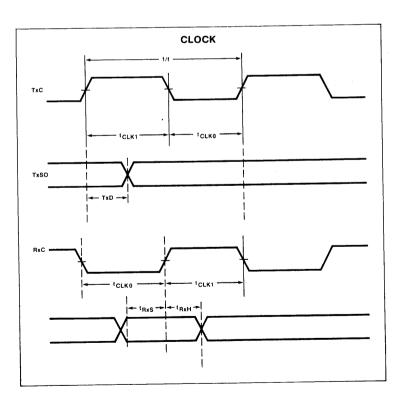
PRELIMINARY SPECIFICATION

Manufacturer reserves the right to make design and process changes and improvements.

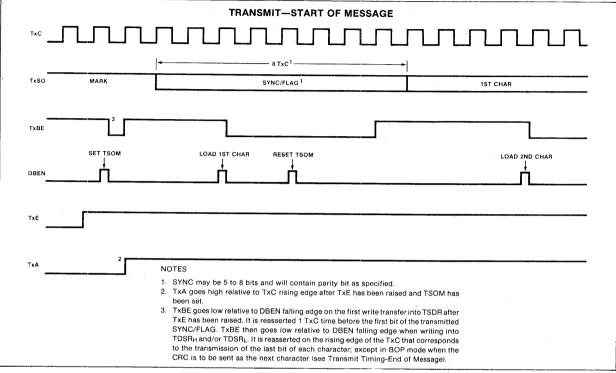
TIMING DIAGRAMS

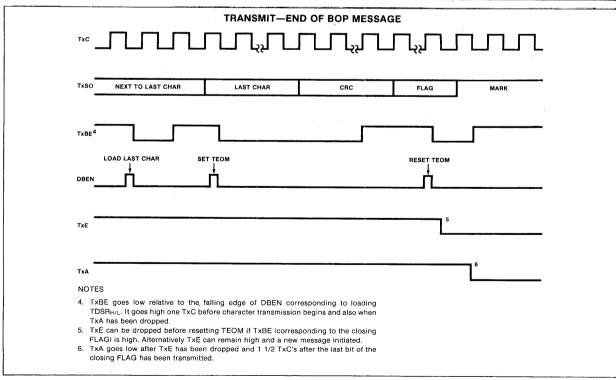




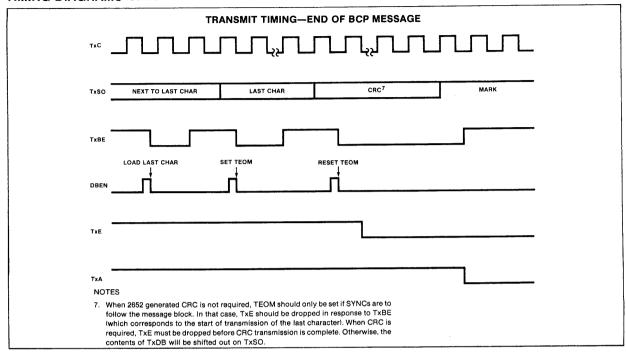


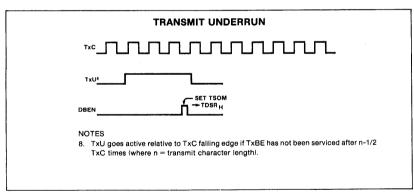
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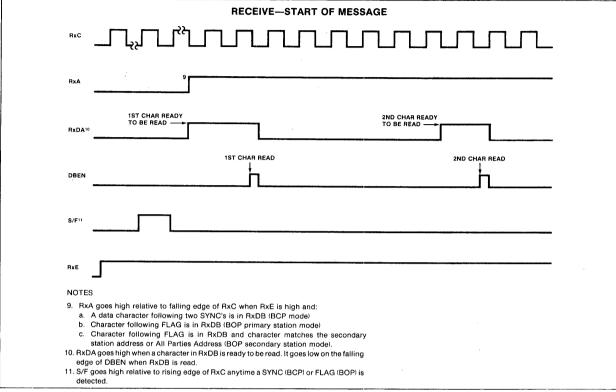


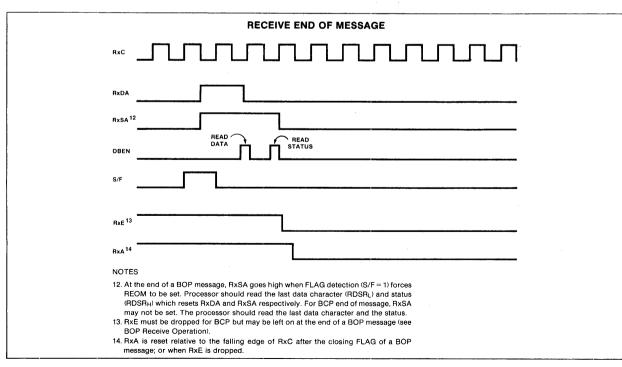
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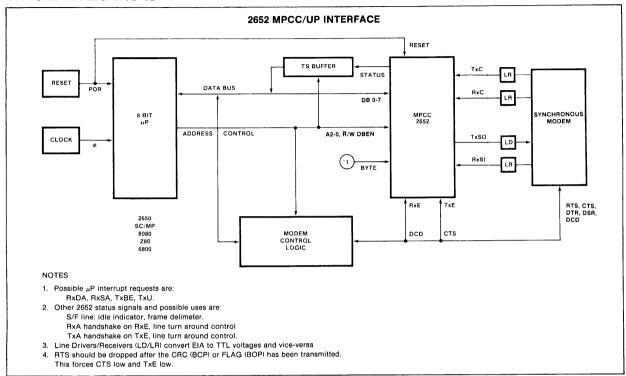
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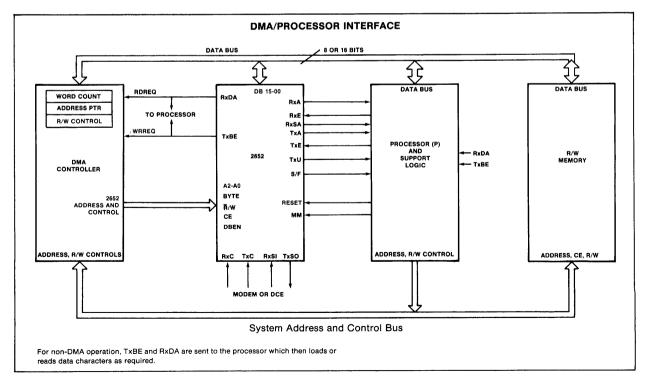




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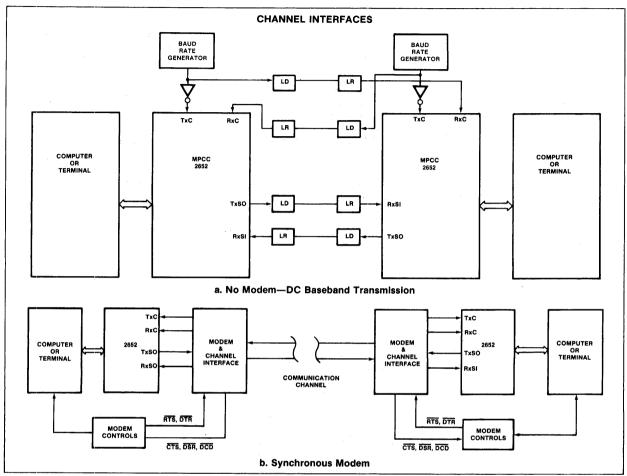
TYPICAL APPLICATIONS





2652-1

TYPICAL APPLICATIONS (Cont'd)



PRODUCT PREVIEW

2655

DESCRIPTION

The Signetics 2655 PPI is designed for 2650 microcomputer systems. It consists of three ports (24 I/O pins), which can be individually programmed to function as input, output or bidirectional ports. Interface with the 2650 is via an eight-bit bidirectional data bus.

The PPI may be programmed for five major modes of operation: static I/O, strobed I/O, bidirectional I/O, serial I/O, or serial/timer I/O. In the serial/timer mode, parallel to serial or serial to parallel conversion of data operates simultaneously with the timer on one of the three ports.

FEATURES

- Five selectable major operating modes: Static I/O Strobed I/O Bidirectional
 - Serial I/O Serial/timer I/O

- Three ports (A, B, and C) with 24 programmable I/O pins
- Completely TTL compatible
- Three MHz programmable timer or event counter
- Fully compatible with the 2650 microprocessor
- Direct bit set/reset capability of each bit for all three ports
- 300ns port read/write access time
- Operates in a polled or interrupt driven system environment
- Forty pin dual in-line package
- Single +5 volt supply

OPERATION

The following is a functional description of the five operating modes of the PPI. Each mode is selected via a mode control word. Interrupt generation and interrupt enable/ disable functions are available with each mode except the static mode which operates entirely under program control.

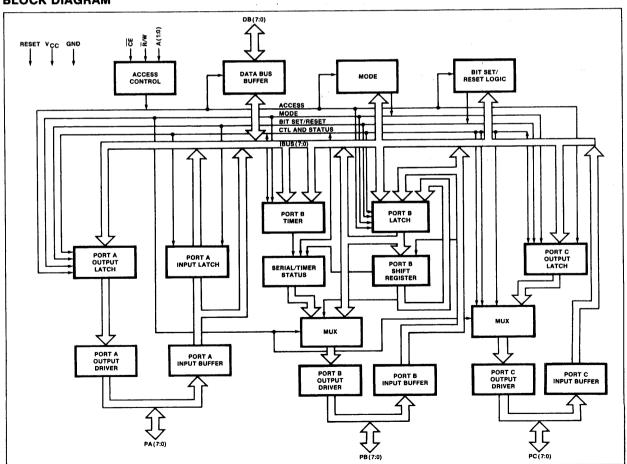
STATIC MODE

All three ports can operate in the static I/O mode. This mode allows each pin of each port to be either an input pin or an output pin. A logic "1" written to a pin of a selected port from the 2650 will condition that pin to be an input or output pin. Writing a logic "0" to the pin conditions that pin to be an output pin only. Outputs are latched while inputs are not. Each pin may be set or reset on an individual basis by a "set/reset" command.

STROBED I/O MODE

In this mode, data may be transferred to or from a specified port in conjunction with strobe or "handshaking" signals. Ports A and B can operate in the strobed I/O mode and port C bits are used as control and status bits. In this mode both inputs and outputs are latched, and each port can be either an input or output.

BLOCK DIAGRAM



2655

PRODUCT PREVIEW

BIDIRECTIONAL I/O MODE

This mode provides a means for communicating with a peripheral device over a single eight-bit bus with both transmitting and receiving capability. Port A operates in this mode with Port C pins providing "handshaking" signals for status and control. Both inputs and outputs are latched and port direction is determined by a control signal from the peripheral.

SERIAL I/O MODE

This mode provides a means for communicating with a peripheral device on a bit serial

basis through Port B. Parallel data from the CPU will be shifted out to the peripheral over the least significant bit of Port B (PB0). Eight clocks will be required for a complete character transfer. The eight-bit character will be repeatedly shifted out until the CPU presents another character to Port B.

For the serial in mode, data is input from the peripheral at the most significant bit of Port B (PB7). Eight clocks will be required to assemble the eight-bit character. An interrupt request will signal the CPU for character transfer.

TIMER MODE

This mode enables the PPI to perform time interval measurements, pulse width measurements, and event counting. This timing function is performed during the serial/timer mode, and is restricted to Port B only. The mode is initiated by selecting the desired operation and loading a 16-bit down counter with an initial value. The counter does not start counting until the upper eight bits have been loaded. An interrupt can be generated to signal the CPU when the timer reaches a zero count.

PIN DESIGNATION

PIN NO.	MNEMONIC	MONIC TYPE	NAME AND FUNCTION
27-34	D7-D0	7-D0 I/O	Data Bus: Eight-bit tri-state bidirectional data bus. All data and command transfers are made using this bus. D0 is the least-significant bit; D7 is the most-significant bit.
35	RESET	SET I	Reset: Resets all internal storage elements, including the data latches and command registers. Resets ports A, B and C to accept input data, and operating mode to Static mode. A functionally equivalent on chip power-on reset is also provided.
8,9	A1, A0	, A0 I	Address: Address lines used to select internal PPI modes or registers. Indicates control or data words to be placed on the data bus. Used in conjunction with the \overline{R}/W line.
5	R/W	/W I	Read/Write: When low, gates the selected register to the data bus. When high, gates the contents of the data bus into the selected register.
6	CE	Σ I	Chip Enable: When low, identifies that control and data lines to the PPI are valid.
36	SCLK	I I	Serial Clock: Provides a serial clock for the parallel-to-serial or serial-to-parallel conversion.
37-40 1-4	I/O PA7-PA0	A7-PA0 I/O	Port A: An eight-bit tri-state quasi-bidirectional port.* PA0 is the least-significant bit; PA7 is the most-significant bit.
25-18	I/O PB7-PB0	37-PB0 1/O	Port B: An eight-bit quasi-bidirectional port.* PB0 is the least-significant bit; PB7 is the most-significant bit.
	-		Port B also has parallel-to-serial, or serial-to-parallel conversion capability with PB0, 7 being either the serial output or input respectively. Data is double buffered.
			Port B can also operate as a 16-bit binary timer, as an event counter, or as a pulse width indicator. An output is generated whenever the counter is decremented to the all-zero state.
10-13 17-14	I/O PC7-PC0	07-PC0 I/O	Port C: An eight-bit quasi-bidirectional port.* PC0 is the least-significant bit; PC7 is the most significant bit. Port C bits are also used as control and status signals in conjunction with ports A and B. When a pin is used as a strobe input, the line receives an external strobe input which clocks information from port A or port B into the port A or port B data latches.
			When a pin is used as a status line, the line indicates port A or port B status condition which may be used as an interrupt input to the 2650.
26		1	+5 Volt: Power supply.
7		1	Ground: 0V reference.
		1	When a pin is used as a status line, the line indicates port A or port B status condit be used as an interrupt input to the 2650. +5 Volt: Power supply.

[&]quot;A quasi-bidirectional port allows each bit to be designated as input or output under program control. If any bit of the port is set to a "1," that bit becomes an input or output depending on the usage of the port pin. If the peripheral is driving the port bit (i.e., overriding the logic "1" condition produced by the internal port pull up resistor), then the bit is an input. If the peripheral is receiving from the port bit, then a "0" or "1" written to the port will be transmitted to the peripheral.

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2656-1

DESCRIPTION

The Signetics 2656 System Memory Interface (SMI) is a mask programmable circuit with on-chip memory, I/O, and timing (clock) functions. It is useable either in 2-chip or multi-chip microcomputer systems. Used with the 2650 microprocessor, it provides a 2-chip microcomputer. This 2-chip microcomputer offers the user 2KX8 bits of ROM, 128X8 bits of RAM, and an 8-bit inputoutput port.

Used as a system interface in a multi-chip microcomputer, with larger memory and/or additional peripheral requirements, the programmable versatility of the SMI provides decoded chip enable outputs. These outputs connect directly to other memory or I/O functional blocks with few and often no requirement for additional interfacing chips. This reduces both chip count and cost in complex microcomputer systems.

The 2656 is processed using Signetics nchannel silicon gate technology. Only a single power supply of +5 volts is needed for operation.

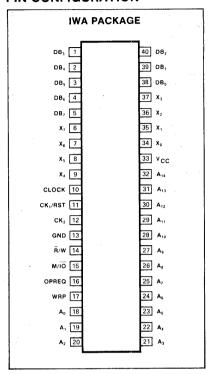
FEATURES

- 2KX8 mask programmable ROM
- 128X8 static RAM
- 8 multi-purpose pins for either chip enables or I/O bits
- . 8-bit latch for either I/O or MPU storage
- Internal clock generator with crystal, RC, or external timing source
- System power-on reset
- 40-pin dual-in-line package
- Single +5 volt supply

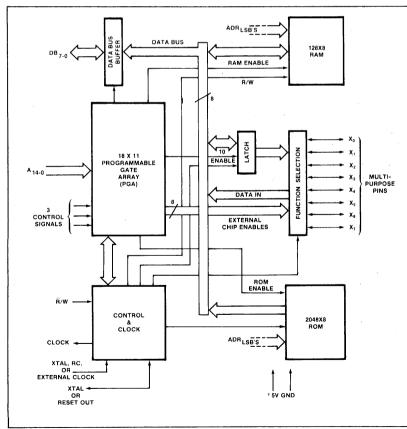
APPLICATIONS

- 2-chip microcomputer
- System control for multi-chip microcomputers—eliminates or reduces TTL suppport circuitry for memory and I/O device selection.
- From small (2K-2 chip) to 32K microprocessor-based systems

PIN CONFIGURATION



BLOCK DIAGRAM



FUNCTIONAL BLOCK DESCRIPTIONS Data Bus Buffer

A tri-state bidirectional 8-bit bus transceiver for data transfer between the SMI and MPU.

Programmable Gate Array (PGA)

Provides select signal outputs for the internal ROM, RAM, Latch, and up to 8 multipurpose I/O pins that are mask programmed as Chip Enables. A PGA output is active when the input variables match any one of 11 corresponding mask programmed product terms. The 18 input variables are normally address and control bus signals from the MPU and may be programmed as "1", "0", and "don't care." Each product term is a specified combination of the input variables.

Control and Clock

Generates the Clock output signal to the MPU and control signals for the ROM, RAM, and Latch. A mask programmable frequency divider provides input frequency division by 1, 2, 3, or 4. The timing source is mask programmable and may be a crystal, RC, or external oscillator. If either of the latter two are designated as a timing source, the second timing pin becomes a Reset output to the MPU.

2656-1

PRELIMINARY SPECIFICATION

PIN DESIGNATION

MNEMONIC	PIN NO.	TYPE	NAME AND FUNCTION
DB ₀ -DB ₇	38-40 1-5	1/0	8-bit Bidirectional Data Bus: All data transfers between the MPU and ROM, RAM, Latch, and X pins are made using this bus.
A ₀ -A ₁₄	18-32	1	MPU Address Bus: Address bus inputs occupy contiguous bit positions with A_0 as the least significant address bit.
OPREQ	16	- 1	Control Signal: A signal that specifies the valid state of address and control bus.
M/ĪŌ, WRP	15, 17	1	Optional Signals: Possibilities include Memory or I/O (M/IO), Write Pulse (WRP), external control signals, or additional high order address bits.
R/W	14	'	Read/Write Control: A control signal from the MPU that indicates whether the requested operation is to be a Read or Write (0 or 1 respectively). This signal must not change while OPREQ is true.
CLOCK	10	0	Clock Output to the MPU: The frequency is determined by the timing element and the mask programmable divisor (divided by 1, 2, 3, 4).
CK ₁ /RST, CK ₂	11,12	1/0, 1	Connections for the Timing Element: Only CK ₂ is necessary for an RC or external timing source. The CK ₁ /RST pin then becomes a power-on Reset output. Two pins are necessary for direct connection of a crystal.
Vcc	33	1	+5V: Power supply.
GND	13	1	Ground: 0V reference ground.
X ₀ -X ₇	34-37 9-6	1/0	Multi-purpose I/O Pins: These pins can be mask programmed as external memory or I/O chip enables, or bidirectional I/O port data bits, or any combination of the two.

ROM

2,048 bytes of mask-programmable Read Only Memory for storage of instructions and constants. The ROM base address is PGA mask programmable over the entire MPU address range. The ROM can be disabled by a mask option.

RAM

128 bytes of Read/Write Memory for MPU data storage and retrieval. The RAM base address is PGA mask programmable over the entire MPU address range. RAM dominates over ROM if address overlap is intentionally mask programmed. The RAM can be disabled by a mask option.

Function Select

A 1X8 Function Select array of mask-programmable contacts determine the function of each of the multi-purpose I/O pins (X_0-X_7) . Two modes exist:

- CE The X pin is an active low Chip Enable (CE) for either external memory or an I/O port. PGA inputs receive the external address and MPU control signals required to generate the CE output.
- P The X pin is a bidirectional I/O port data bit. A portion of the PGA provides the control signal to select the port.

Latch

Holds output data for the multi-purpose I/O pins mask programmed as a mode P. The latch continues to function as a read/write element even if all multi-purpose I/O pins are programmed as chip enables. Thus, any X pin that is programmed as an external chip enable can have corresponding latch bits available for temporary data storage or software flags. To read an input pin, the corresponding latch bit must first be written to a "one" by the MPU program. This is done to disable all active outputs, changing them to passive pullup outputs. This permits inputs to be sensed on the same pin. Subsequent reads of the same pin do not have to be preceded by a write if the state of the latch pin remains a "one."

2656-1

ABSOLUTE MAXIMUM RATINGS1

	PARAMETER	RATING	UNIT
TA TSTG	Operating ambient temperature ² Storage temperature All voltages with respect to ground ³	0 to +70 -65 to +150 -0.5 to +6.0	°C °C V

NOTES

- 1. Stress above those listed under Absolute Maximum Ratings may cause permanent damage to the device. This is a stress rating only and functional operation of the device at these or at any other condition above those indicated in the operations section of this specification is not implied.
- 2. For operating at elevated temperatures, the device must be derated based on +150°C
- maximum junction temperature and thermal resistance of 55° C/W junction to ambient (IWA ceramic package.)
- 3. This product includes circuitry specifically designed for this protection of its internal devices from the damaging effects of excessive static charge. Nonetheless, it is suggested that conventional precautions be taken to avoid applying any voltages larger than the rated maxima.

DC ELECTRICAL CHARACTERISTICS $T_A = 0$ °C to +70°C, $V_{CC} = 5.0V \pm 5\%1,2,3,4,10$

		TEST CONDITIONS		LIMITS		UNIT
	PARAMETER	TEST CONDITIONS	Min	Тур	Max	CIVIT
	Input voltage					V
VIL	Low		5		0.8	
ViH	High		2.2			
	Output voltage			İ		V
Vol	Low	$I_{OL} = 1.6 mA$			0.45	
Voн	High	$I_{OH} = -100 \mu A6$	2.4	ļ		ļ
VOHP	I/O port output high voltage	$I_{OHP} = -50\mu A^7$	2.4			V
lıL	Input load current	$V_{IN} = 0 \text{ to } 5.5V$		1	10	μΑ
ILD	Data bus leakage current	$V_{OUT} = 4.0V$			10	μΑ
Icc	Power supply current				150	mA
	Capacitance	$T_A = 25^{\circ} C$, $V_{CC} = 0V$	i			
CIN	Input	$f_c = 1MHz$		4	10	рF
Соит	Output	Unmeasured pins		4	10	pF
C _{1/O}	1/0	tied to ground		6	10	pF

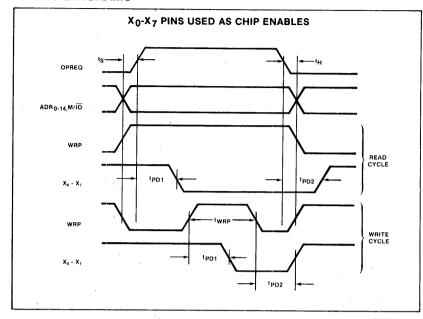
AC ELECTRICAL CHARACTERISTICS $T_A = 0$ °C to +70°C, $V_{CC} = 5.0 V \pm 5\%$, 1,2,3,4,9,10

				LIMITS		·	
	PARAMETER	TEST CONDITIONS	Min Typ		Max	UNIT	
ts	Address and control setup time		0			ns	
tн	Address and control hold time		0			ns	
t _{PD1}	Propagation delay time - high to low5	C _L = 100pF	50		230	ns	
tPD2	Propagation delay time - low to high5	C _L = 100pF	50		230	ns	
twap	Write pulse width		200		1	ns	
top	Data bus delay time for read	C _L = 100pF	200		580	ns	
tor	Data bus floating time for read	C _L = 100pF			580	ns	
tobs	Data bus setup time for write		200	•		ns	
tobe	Data bus hold time for write	Į.	200			ns	
top	Output delay time	C _L = 100pF	100		400	ns	
fc	Crystal or external clock frequency				4.0	MH:	
tc	External clock high or low state		110			ns	
trst	Reset output pulse width8	1	30		300	μS	
taws	Address to write pulse setup time		50			ns	

- 1. Parameters are valid over operating temperature range unless otherwise specified.
- 2. All voltage measurements are referenced to ground. All time measurements are at the VOH, VOL, VIH, VIL levels as appropriate.
- 3. Manufacturer reserves the right to make design and process changes and
- 4. AC characteristics assume the following input signals from a Signetics 2650 microprocessor:
- Pins 18-32 = ADR₀-ADR₁₄
- = M/10
- = OPREQ = WRP Pin 17
- 5. For X₀-X₇, time reference is OPREQ or time input changes, whichever is later. 6. Data Bus, Chip Enable, CK₁/RST or Clock outputs.
- 7. Port outputs only.
- 8. V_{CC} rise time must be less than 100 microseconds.
- 9. The delay times (Minima and Maxima) tend to track each other for any single device.
- 10. This is not a final specification. Parametric limits are subject to change.

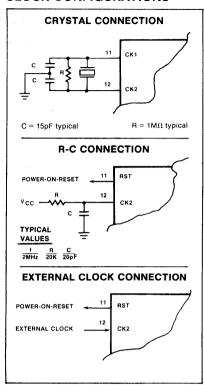
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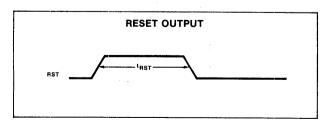
TIMING DIAGRAMS

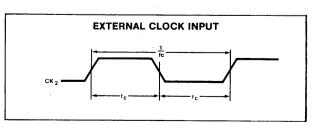


ADR₀₋₁₄. M/IO ADR₀₋₁₄. M/IO FLOATING OR V₀-X₇ PINS USED AS INPUTS) OPREO ADR₀₋₁₄. M/IO FLOATING OR NOT VALID t_{DD} t_{DF}

CLOCK CONFIGURATIONS

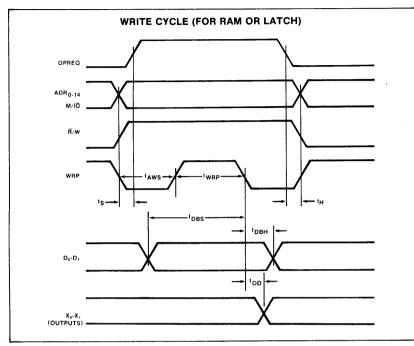






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TIMING DIAGRAMS (Cont'd)



CUSTOM PATTERN PROGRAMMING INSTRUCTIONS

A computer-aided technique utilizing punched computer cards is employed to specify a custom version of the 2656. This technique requires that the customer supply Signetics with a deck of standard 80-column computer cards describing the data to be stored in the ROM array, the Programmable Gate Array (PGA), and the Function Select Array.

On receipt of a card deck, Signetics will translate the card deck to a truth table using the Signetics Computer-Aided Design (CAD) facility. The truth table will then be sent to the customer for final approval. On receipt of final approval, Signetics will produce masks and proceed with manufacturing.

The contents of each card in the deck are:

Customer Identification Cards are always labeled with a "C" in column 1. For customer identification any number of cards is permitted. A 4-card example is shown below. The following data should be included:

COLUMN	DATA
1	С
2	Blank
3-66	Company name (Card #1)
	Street address (Card #2)
1	City, State, Zip (Card #3)
i	Contact person name (Card
	#4)
67	Blank
68-70	SMI
71	Blank
72-75	2656
76-78	Blank
79-80	2-digit decimal number indi-
	cating the truth table number.
	Must be the same on all cards in the deck.

(Explanatory Note: The next card for this example is card No. 5)

CARD #5—SMI FUNCTIONAL PARAMETERS

Ì	COLUMN		DATA	
	1-9	Blank		ı

The next eight columns specify whether multi-purpose pins (X_0-X_7) are to be chip enables or I/O port data bits. Each column must contain either the character "E" for chip enable or the character "P" for I/O port. Card column 10 specifies X_0 , card column 17 specifies X_7 .

COLUMN	DATA		
10	E or P for X ₀		
11	E or P for X ₁		
12	E or P for X ₂		
13	E or P for X ₃		
14	E or P for X ₄		
15	E or P for X ₅		
16	E or P for X ₆		
17	E or P for X ₇		
18-19	Blank		

Select one of the next three columns to specify the type of clock source; crystal, external or R/C network. An "X" designates the selected clock source. The other two columns must be blank.

COLUMN	DATA		
20	X or blank (Crystal)		
21	X or blank (R/C Network)		
22	X or blank (External)		
23-29	Blank		

Select one of the next four columns to specify the clock source divider value to divide by 1, 2, 3 or 4. An "X" designates the selected divisor. The other three columns must be blank.

COLUMN	DATA			
30	X or blank (÷ 1)			
31	X or blank (÷ 2)			
32	X or blank (÷ 3)			
33	X or blank (÷ 4)			
34-39	Blank			

Access to ROM, RAM or I/O Port bits may be disabled. The next three columns are used to disable ROM, RAM or I/O Port. Specify an "X" (disable) or blank (do not disable).

COLUMN	DATA				
40	X or blank (ROM)				
41	X or blank (RAM)				
42	X or blank (PORT)				
43-78	Blank				
79-80	Two-digit decimal number in- dicating the truth table num- ber.				

2656-1

CARD #6—PGA SPECIFICATION

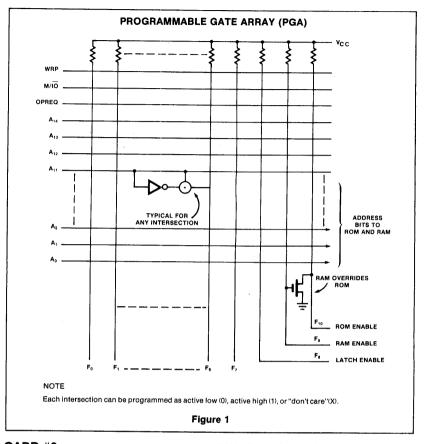
The PGA is an 18 X 11 input-output structure. The first eight outputs are available as chip enables (if selected by the appropriate function—select parameters on card #5). The last three outputs are internally connected to enable the SMI internal I/O Port, RAM and ROM. (See Figure 1.) Ascending order of card columns correspond to ascending PGA outputs. Thus, columns 5 through 15 correspond, respectively, to PGA chip enable outputs 0 through 7, the latch enable, the RAM enable and the ROM enable. Valid characters are 1, 0 or X. "1" indicates active high, "0" indicates active low, and "X" indicates don't care.

COLUMN	DATA
1	Α
2	0
3-4	Blank
5-15	1, 0 or X (A ₀ for F ₀ to F ₁₀)
16	Blank
17-27	1, 0 or X (A ₁ for F ₀ to F ₁₀)
28	Blank
29-39	1, 0 or X (A ₂ for F ₀ to F ₁₀)
40	Blank
41-51	1, 0 or X (A ₃ for F ₀ to F ₁₀)
52	Blank
53-63	1, 0 or X (A ₄ for F ₀ to F ₁₀)
64	Blank
65-75	1, 0 or X (A ₅ for F ₀ to F ₁₀)
76-78	Blank
79-80	2-digit decimal number indi-
	cating the truth table number.

CARD #7—

PGA SPECIFICATION (Cont'd)

COLUMN	DATA
1	Α
2	6
3-4	Blank
5-15	1, 0 or X (A ₆ for F ₀ to F ₁₀)
16	Blank
17-27	1, 0 or X (A ₇ for F ₀ to F ₁₀)
28	Blank
29-39	1, 0 or X (A ₈ for F ₀ to F ₁₀)
40	Blank
41-51	1, 0 or X (A ₉ for F ₀ to F ₁₀)
52	Blank
53-63	1, 0 or X (A ₁₀ for F ₀ to F ₁₀)
64	Blank
65-75	1, 0 or X (A ₁₁ for F ₀ to F ₁₀)
76-78	Blank
79-80	2-digit decimal number indi-
	cating the truth table number.



CARD #8—
PGA SPECIFICATION (Cont'd)

	OII IOATIOIT TOOILE
COLUMN	DATA
1	Α
2-3	12
4	Blank
5-15	1, 0 or X (A ₁₂ for F ₀ to F ₁₀)
16	Blank
17-27	1, 0 or X (A ₁₃ for F ₀ to F ₁₀)
28	Blank
29-39	1, 0 or X (A ₁₄ for F ₀ to F ₁₀)
40	Blank
41-51	1, 0 or X (M/ \overline{IO} for F ₀ to F ₁₀)
52	Blank
	1,0 or X (OPREQ for F ₀ to F ₁₀)
	Blank
	1, 0 or X (WRP for F ₀ to F ₁₀)
	Blank
79-80	2-digit decimal number indi-
	cating the truth table number.
	1 2-3 4 5-15 16 17-27 28 29-39 40 41-51

BINARY ROM CODE FORMAT CARD #9 THROUGH CARD #264

These remaining cards specify the 2048 8-bit locations of the SMI ROM. The first data field of card #9 is ROM location 0, the last data field of card #264 is ROM location 2047. For each field, the leftmost character is the MSB, the rightmost character is the LSB. Valid characters are 1or 0. Alternately, Por N characters may be used. A "1" or "P" represents a logic high at the SMI data bus connection. A "0" or "N" represent a logic low. Each of these cards is configured as follows:

COLUMN	DATA
	Right justified decimal address of first data field of the card (columns 7-14).
6 7-14	Blank 1 or 0; or, P or N
7-14	or o, or, P or N

2656-I

PRELIMINARY SPECIFICATION

COLUMN	DATA
15	Blank
16-23	1 or 0; or, P or N
24	Blank
25-32	1 or 0; or, P or N
33	Blank
34-41	1 or 0; or, P or N
42	Blank
43-50	1 or 0; or, P or N
51	Blank
52-59	1 or 0; or, P or N
60	Blank
61-68	1 or 0; or, P or N
69	Blank
70-77	1 or 0; or, P or N
78	Blank
79-80	2-digit decimal number indi-
	cating the truth table number.

HEXADECIMAL ROM CODE FORMAT

CARD #9 THROUGH CARD #72

The customer may also submit ROM code in hexadecimal format. 32 locations (bytes) are specified on each card. Columns 7 and 8 of card #9 contain the data for ROM location 0 while columns 69 and 70 of card #72 contain the data for ROM location H'7EF'. In

each data pair, the left character is the hexadecimal equivalent of bits D₇ to D₄; the right character is the hexadecimal equivalent of bits D₃ to D₀. Table 1 shows the conversion from binary to hexidecimal.

Each of these cards is configured as follows:

COLUMN	DATA
1-4	Right justified hexadecimal address of first data pair of the card (columns 7-8). For example, card #9 will contain '0000,' card #10 will contain '0020' and so forth until card #72, which contains '07E0.'
5-6	Blank
7-8	First data pair
9-10	Next data pair
11-12	Next data pair
11-12	Next data pair
11-12	Next data pair
11-12	Next data pair
67-68	Next to last data pair
69-70	Last data pair of card
71-78	Blank
79-80	2-digit decimal number indicating the truth table number.

BINARY			HEXA-	
COMBINATION			DECIMAL	
D ₇ -D ₄ OR D ₃ -D ₀			CHARACTER	
0 0 0 0 0 0 0 0 1 1 1 1 1 1	0 0 0 0 1 1 1 1 0 0 0	0 0 1 1 0 0 1 1 0 0 1 1 0 0	0 1 0 1 0 1 0 1 0 1 0 1 0 1	0 1 2 3 4 5 6 7 8 9 A B C D E F

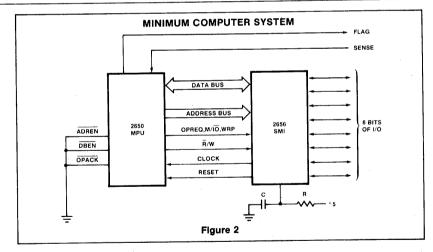
Table 1 BINARY TO HEXADECIMAL CONVERSION

APPLICATION EXAMPLES A TWO CHIP MICROCOMPUTER

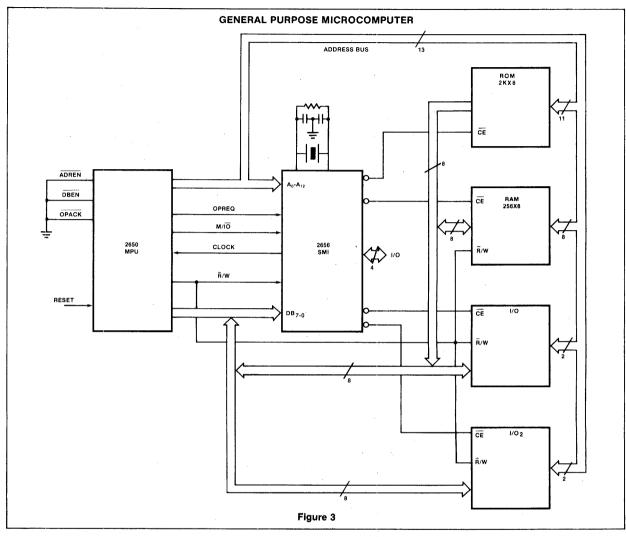
The minimum system depicted in Figure 2 is composed of the 2650 MPU, a 2656 SMI, and a resistor/capacitor network for a timing input. The SMI provides MPU CLOCK and RESET signals and eight I/O Port bits (X₀-X₇).

Note that because all eight I/O bits share the same device address, the MPU program must "housekeep" unwanted bits when reading or writing to any of them. If a one-bit write operation is desired, the MPU must preserve the state of all other output pins. This can be done by reading the latch, changing the state of the selected bit, and outputing the result.

Any latch bit used for input must first be written to a Logic "1" before a read operation. After reading the data, the bits not desired must be masked out by the program.



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A GENERAL PURPOSE MICROCOMPUTER

A general purpose microcomputer is illustrated in Figure 3. The SMI provides ROM RAM and the MPU clock from a crystal source. It also provides chip enables for external ROM, RAM and LSI peripheral controllers 1 and 2 as well as four bidirectional data bits.

The state of M/\overline{O} in the Address Map (Table 6) indicates that I/O_1 and I/O_2 are memory mapped while the SMI port is referenced using any I/O instruction (extended or non-extended) on the 2650. See the SMI Program Table (Table 7).

The 2650 addresses the SMI port to input or output data on X₄-X₇. SMI latch bits 0-3 are available for MPU temporary storage since

the corresponding X pins are used for external chip enables.

Note that the 13 MPU address bits are assigned to the first 13 PGA inputs (A₀-A₁₂) and that only two MPU control signals (OPREQ, M/IO) are input to the PGA (WRP is

not required).

The SMI uses the \overline{R}/W signal from the 2650 to read or write from the internal RAM and port. External devices are enabled when their address (A₀-A₁₂, M/ \overline{IO}) have been decoded by the SMI and OPREQ is valid.

	PIN	# BYTES	ADDRESS	M/IO
Internal ROM		2K	0000-07FF	1
External ROM	X ₀	2K	0800-0FFF	1
*Spare ROM		1K	1000-13FF	1
External RAM	X ₁	256	1400-14FF	1
Internal RAM	_	128	1500-157F	1
*Spare RAM	_	512	1600-17FF	1
I/O ₁	X ₂	4	1800-1803	1
I/O ₂	Х3	4	1804-1807	1
SMI Port	_	- 1	ANY	0

*Space left for possible system expansion.

Table 6

ADDRESS MAP FOR GENERAL PURPOSE MICROCOMPUTER

2656-I

							P	RODU	CT TE	RM										
NO.	WRP	A 14	A 13	M/IO	OPREQ	A ₁₂	A ₁₁	A 10	A ₉	A 8	A 7	A ₆	A ₅	A ₄	A ₃	A ₂	Α1	A ₀	FS*	_
0	X	Х	Х	1	1	0	1	х	Х	х	Х	х	Х	Х	х	Х	х	x	E	X_0
1	X	Х	X	1	1	1	0	1	0	0	Х	Х	Х	Х	Х	Х	X	X	E	X ₁
2	Х	X	Х	. 1	1	1	1	0	0	0	0	0	0	0	0	0	X	X	LE_	X2
3	X	X	X	1	1	. 1	1	0	0	0	0	0	0	0	0	1	X	X	E	X3
4	X	Х	X	X	Х	. X	X	Х	Х	Х	Х	X	X	Х	X	X	X	X	P_	_ X₄
5	Х	Х	Х	Х	X	Х	X	Х	X	Х	Х	X	X	X	X	X	X	X	P	X ₅
6	X	Х	Х	X	Х	X	Х	Х	Х	X	Х	X	X	X	X	X	X	X	P	_ X ₆
7	Х	Х	X	Х	Х	Х	Х	Х	X	X	Х	X	X	X	X	X	X	X	L P	_ X ₇
8	X	Х	X	0	1	Х	Х	X	X	X	Х	X	X	X	X	X	X	X	PORT	Γ
9	Х	X	X	1	1	1	0	1	0	1	0	Х	X	X	X	X	X	X	RAM	l
10	X	Х	Х	1	1	0	0	Х	Х	Х	Х	Х	Х	X	X	X	X	X	ROM	ı
	DI	SABLE	E					PGA				XTL	R	/C	EXT		1	2	3 4	4
								٦	rimin(g sou	RCE		1					\checkmark		
RO	M I	RAM	PO	RT			Tal	ole 7 S	MI PR	OGRA	M TAI	RIF				CK Divide by				

*P = Port E = Enable

Table 7 SMI PROGRAM TABLE FOR GENERAL PURPOSE

MICROCOMPUTER

HIGH PERFORMANCE MICROCOMPU-**TER SYSTEM**

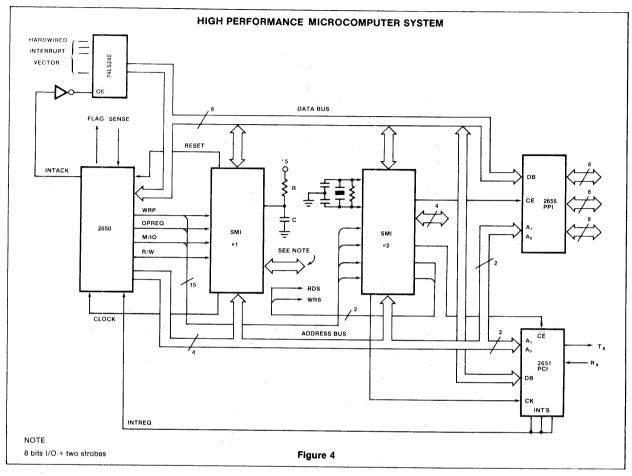
The six-chip high performance system of Figure 4 contains a 2655 Programmable Peripheral Interface and a 2651 Programmable Communication Interface. Two SMIs are used in the system. The first SMI provides CLOCK and RESET signals to the 2650, initial increments of ROM and RAM, and eight bits of bidirectional I/O to the user's external interface. The second SMI provides a crystal controlled clock for the PCI, additional amounts of ROM and RAM. read (RDS) and write (WRS) strobes for the first SMI's eight-bit port, chip enables for PPI and PCI, and four bits of bidirectional I/O data.

The PPI and PCI addresses are decoded by SMI#2 to generate the chip selects on X7 and X₆. The two least significant MPU address bits specify one of three PPI eight bit I/O ports or one of four PCI internal registers. Three PCI interrupt conditions are WIRE-ORed (active low) to the 2650 interrupt request input (INTREQ). The interrupt acknowledge (INTACK) is inverted and used to enable the interrupt vector onto the data bus through the 74LS240 tri-state buffer.

									RM	CT TE	RODU	P							
FS*	A ₀	A ₁	A ₂	A ₃	A4	A 5	A 6	Α7	A 8	Ag	A ₁₀	A ₁₁	A 12	OPREQ	WRP	A 13	A 14	M/ĪŌ	NO.
Р	Х	Х	Х	Х	Х	Х	Х	Х	х	Х	Х	Х	Х	Х	х	Х	Х	Х	0
Р	Х	Х	X	Х	Х	Х	Х	X	Х	Χ	Х	Х	X	Х	Х	Х	Х	Х	1
E	X	X	Х	1	0	0	0	0	0	1	0	1	1	1	1	X	X	Х	2
E	1	1	0	0	0	0	0	0	0	. 1	0	1	1	1	1	Х	Х	Х	3
E	0	1	0	0	0	0	0	0	0	1	0	1	1	1	1	Χ	X	Х	4
E	1	0	0	0	0	0	0_	0	0	1	0	1	1	1	1	Х	X	Х	5
E	X	X	X	X	X	Х	Х	Х	0	0	1	0	1	1	X	Х	X.	Х	6
E	X	X	Х	X	Χ	Х	Χ	Х	X	Х	0	1	0	1	Х	Х	Х	Х	7
PORT	0	0	0	0	0	0	0	0	0	1	0	1	1	1	1	Х	Х	Х	8
RAM	X	X	Х	Х	X	Х	Х	0	1	0	1	0	1	1	1	Х	Х	Х	9
ROM	X	X	X	Х	X	Х	Х	Х	- X	X	Х	0 .	. 0	1	1	Х	Х	Х	10
3 4	2	1		EXT	C E	R/C	XTL				PGA					E	ISABL	D	
	\checkmark		Γ	$\sqrt{}$		T		IRCE	sou	rimine	7								
le hv	K Divi		-						000		hla R (_			RT	DC.	RAM		RO

= Port E = Enable Table 8 SMI PROGRAM TABLE FOR HIGH PERFORMANCE MICROCOMPUTER SYSTEM

2656-1



MICROCOMPUTER SYSTEM WITH MSI INPUT-OUTPUT

The system of Figure 5 incorporates MSI integrated circuits for input and output. It contains a total of 3k bytes of ROM/PROM, 384 bytes of RAM, six bits of SMI latch storage, one latching input port (74LS374), one gated input port (74LS240), one byte-oriented output port (74LS240), one bit-oriented output port (74LS259/93L34), two single bit inputs and two single bit outputs (one each from the SMI and the 2650). The SMI divides the external clock frequency by 1, 2, 3, or 4 and provides CLOCK and RESET outputs to the 2650. Six X pins are used for chip enables and two for data bits.

Included are system memory map (Table 9) and the SMI Program Table. The I/O device and the SMI port are all memory mapped; thus, the 2650 M/IO signal is not connected to the SMI.

WRP from the 2650 is connected to the SMI to control the two output ports:

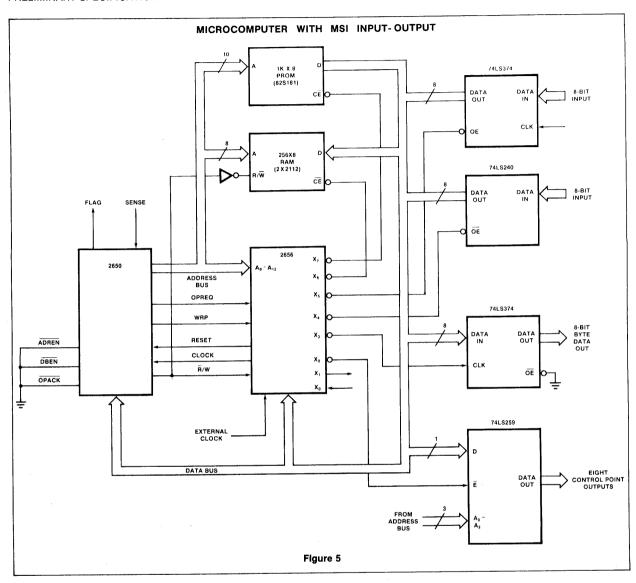
- When output-port-1's address and OPREQ are valid, the WRP input provides a gated chip enable on X₃. The trailing edge of this signal is used to clock 2650 data into the 74LS374.
- 2. One bit of 2650 data is gated into the 74LS259/93L34 addressable latch (output port 2) by means of X₂ which serves as the enable input. This signal is active when output port 2 has been selected. At the same time, three address bus lines select the latch to be written into.

	PIN	BYTES	ADDRESS
Internal ROM	_	2k	0000-07FF
External PROM (82S181)	X7	1k	0800-0BFF
External RAM (2-2112's)	X ₆	256	1400-14FF
Internal RAM	-	128	1500-157F
SMI Port	1 – 1	1	1A00
nput Port 1 (74LS374)	X ₅	1	1A01
nput Port 2 (74LS240)	X4	1	1A02
Output Port 1 (74LS374)	X3	1	1A03
Output Port 2 (74LS259/93L34)	X ₂	8	1A08-1A0F
	1 1	i i	

Table 9 ADDRESS MAP FOR MSI I/O SYSTEM

Table 0 ADDDESS MAD

2656-1



DESCRIPTION

The MP8251 is a programmable Universal Synchronous/Asynchronous Receiver/ Transmitter (USART) chip contained in a standard 28-pin dual-in-line package. The chip, which is fabricated using N-channel silicon gate technology, functions as a serial data input/output interface in the MP8080A microcomputer family. The functional configuration of the MP8251 is programmed by the system software for maximum flexibility, thereby allowing the system to receive and transmit virtually any serial data communication signal presently in use (including IBM Bisync).

The MP8251 can be programmed to receive and transmit either synchronous or asynchronous serial data. The MP8251 performs serial-to-parallel conversion on data characters received from an input/output device or a MODEM, and parallel-to-serial conversion on data characters received from the CPU. The CPU can read the complete status of the MP8251 at any time during the functional operation. Status information reported includes the type and the condition of the

transfer operations being performed by the MP8251, as well as any transmission error conditions (parity, overrun, or framing).

FEATURES

- Synchronous and asynchronous full duplex operations
- Synchronous Mode Capabilities
 Selectable 5- to 8-Bit characters
 Internal or external character synchronization
 - Automatic sync insertion
- Asynchronous mode capabilities
 Selectable 5- to 8-Bit characters
 3 selectable clock rates (1x, 16x or 64x the baud rate)

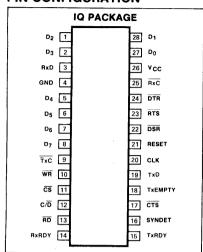
Line break detection and generation 1-, 11/2-, or 2-Stop bit detection and generation

- False start bit detection
- Baud rates

DC to 56k baud (synchronous mode) DC to 9.6k baud (asynchronous mode)

- Transmission error detection capabilities
 Parity
 - Overrun Framing

PIN CONFIGURATION



- Double buffering of data
- TTL compatible
- Single TTL clock
- · Reduces system component count

PIN DESIGNATION

MNEMONIC	TYPE	NAME AND FUNCTION
CS	ı	Chip select: When low (logic 0), the chip is selected. This enables communication between the MP8251 and the MP8080A microprocessor.
RD	1	Read: When low, allows the MP8080A to read data or status information from the MP8251.
WR	1	Write: When low, allows the MP8080A to write data or control words into the MP8251.
C/D	1	Control/data: Used in conjunction with an active RD or WR input (logic 0) to determine overall device operation as indicated below. CS C/D RD WR OPERATION
		0 0 0 1 Data character read from MP8251 0 0 1 0 Data character written into MP8251 0 1 0 1 Status information read from MP8251 0 1 1 0 Control word written into MP8251 1 x x x Device not selected
RESET	. !	Reset: When high (logic 1), places the MP8251 in the idle mode. The device remains in this mode until a new set of control words is written into the MP8251 to program its functional definition. Minimum Reset pulse width is 6 t _{CY} .
CLK	ļ	Clock: TTL clock that is used to generate internal timing signals for the MP8251. The minimum frequency of the CLK input is 30 times the receiver/transmitter clock frequency for the synchronous mode, and 4.5 times the receiver/transmitter clock frequency for the asynchronous mode. The CLK input is normally connected to the ϕ_2 (TTL) output of the 8224 Clock Generator and Driver device.
DSR	1	Data set ready: General purpose input whose condition can be tested by the MP8080A using a status read operation. However, a low level DSR input is normally used to test data set ready conditions.
CTS	í	Clear to send: If low when the TxEn bit (D ₀) of the command Instruction Control Word (see figure) is set high, enables the MP8251 to transmit serial data.

PIN DESIGNATION (Cont'd)

MNEMONIC	TYPE	NAME AND FUNCTION	
TxC	1	Transmitter Clock: This clock input controls the rate at which a data character is to be transmitted. The frequency of the TxC input is equal to the baud rate for the synchronous mode, and is a multiple (1x, 16x or 64x) of the baud rate for the asynchronous mode. A portion of the Mode Instruction Word (see figure) selects the value of the baud rate factor when in the asynchronous mode. Transmitter data are clocked out of the MP8251 on the falling edge of the TxC input.	
RxD	1	Receiver data: Serial data input from a MODEM or an input/output device.	
RxC		Receiver clock: This clock input controls the rate at which a data character is to be received. The frequency and selection of the RxC input is as described above for the TxC input. Receiver data are clocked into the MP8251 on the rising edge of the RxC input.	
Vcc		+5V power supply.	
GND		Ground: 0-V reference.	
DTR	0	Data terminal ready General purpose output which can be set to an active low by programming the DTR bit (D ₁) of the Command Instruction Control Word. However, a low level DTR output is normally used for data terminal ready or rate select control.	
RTS	0	Request to send: General purpose output which can be set to an active low by programming the RTS bit (D_5) of the Command Instruction Control Word. However, the RTS output is normally used for request to send control in the transmit mode.	
TxD	0	Transmitter data: Composite serial data output to a MODEM or input/output device. The TxD output is held in the marking state (logic 1) upon a Reset operation.	
TxRDY	0	Transmitter ready: When high, alerts the MP8080A that the transmitter is ready to accept a data character. The TxRDY output, which is automatically reset whenever a character is written into the MP8251, can be used as an interrupt to the system. For polled operation, the condition of the TxRDY signal can be tested by the MP8080A using a status read operation.	
TxE	0	Transmitter empty: Goes high to indicate the end of a transmit mode. The TxE output is automatically reset whenever a character is written into the MP8251. In the synchronous mode, a high-level TxE output indicates that a character has not been loaded, the transmitter buffer is empty, and the sync character(s) of a data block are soon to be transmitted automatically as fillers.	
RxRDY	0	Receiver ready: When high, alert the MP8080A that the receiver contains a data character that is ready to be input to the CPU. The RxRDY output, which is automatically reset whenever a character is read from the MP8251, can be used as an interrupt to the system. For polled operation, the condition of the RxRDY signal can be tested by the MP8080A using a status read operation.	
D7-D0	I/O	Data bus: This bus comprises eight Tri-state input/output lines. The bus provides bidirectional communications between the MP8251 and the MP8080A. Data are routed to or from the internal data bus buffer upon execution of an MP8080A OUT or IN instruction, respectively. In addition, control words, command words and status information are transferred through the data bus buffer.	
SYNDET	I/O	Sync detect: This pin may be used in the synchronous mode only. System software can program SYNDET as either an input or an output. When used as an output (Internal sync detect mode), a high level SYNDET output is automatically reset upon a status read operation by the MP8080A. When used as an input (external sync detect mode) a high level SYNDET causes the MP8251 to start assembling data characters on the falling edge of the next RxC input.	

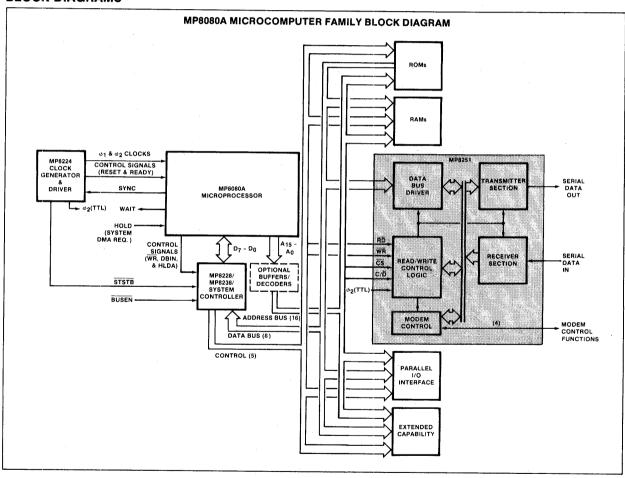
ABSOLUTE MAXIMUM RATINGS*

	PARAMETER	RATING	UNIT
T _A T _{STG} P _D	Ambient temperature under bias Storage temperature Voltage on any pin with respect to ground Power dissipation	0 to +70 -65 to +150 -0.5 to +7	°C V W

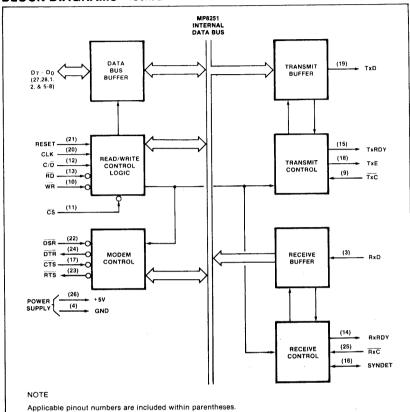
*NOTE

Maximum ratings indicate limits beyond which permanent damage may occur. Continuous operation at these limits is not intended and should be limited to those conditions specified under dc electrical characteristics.

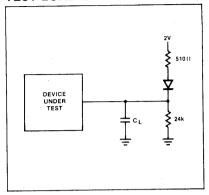
BLOCK DIAGRAMS

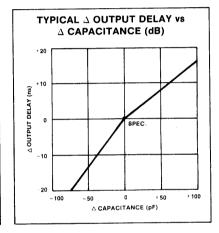


BLOCK DIAGRAMS (Cont'd)



TEST LOAD CIRCUIT





DC ELECTRICAL CHARACTERISTICS $T_A = 0$ °C to +70°C; $V_{CC} = 5.0$ V ± 5 %; GND = 0V.

				UNIT		
	PARAMETER	TEST CONDITIONS	Min	Тур	Max	UNIT
VIH VIL	Input voltage High Low		2.0 -0.5		V _C C 0.8	V
Voh Vol	Output voltage High Low	$I_{OH} = -100 \mu A$ $I_{OL} = 1.6 mA$	2.4		0.45	V
IDL IIL	Leakage Data bus Input	V _{OUT} = 0.45V V _{OUT} = V _{CC} V _{IN} = V _{CC}			-50 10 10	μΑ
Icc	Power supply current			45	80	mA
Cin Ci/o	Capacitance Input I/O	$T_A = 25^{\circ}C; V_{CC} = GND = 0V$ $f_c = 1MHz$ Unmeasured pins returned to GND			10 20	pF

AC ELECTRICAL CHARACTERISTICS $T_A = 0$ °C to +70°C; $V_{CC} = 5.0 V \pm 5\%$; $GND = 0 V_{CC} = 5.0 V \pm 5\%$; $COC = 5.0 V \pm 5\%$

	PARAMETER	TEST CONDITIONS		UNIT			
		1231 CONDITIONS	Min	Тур	Max	UNI	
BUS PA Read C	RAMETERS1 ycle						
tar	Address stable before \overline{READ} (\overline{CS} , C/\overline{D})		50			ns	
t _{RA}	Address hold time for READ (CS, C/D)		5		<u> </u>	ns	
t _{RR}	READ pulse width		430			ns	
t _{RD}	Data delay from READ	C _L = 100pF			350	ns	
tDF	READ to data floating	C _L = 100pF			200	ns	
		$C_L = 15pF$	25			ns	
tav	Recovery time between WRITES2		6		†	tcy	
Write Cy	rcle					101	
taw	Address stable before WRITE		20			ns	
twA	Address hold time for WRITE		20			ns	
tww	WRITE pulse width		400			ns	
t _{DW}	Data set-up time for WRITE		200			ns	
two	Data hold time for WRITE		40			ns	
OTHER	TIMINGS			·		113	
tcy	Clock period ³		0.420		1.35		
tφw	Clock pulse width		220		0.7tcy	μS	
t _R , t _F	Clock rise and fall time		0		50	ns	
totx	TxD delay from falling edge of TxC	C ₁ = 100pF	-		1	ns	
tsRx	Rx data set-up time to sampling pulse	C _L = 100pF	2		<u> </u>	μS	
tHRx	Rx data hold time to sampling pulse	C _L = 100pF	2			μS	
f _{Tx}	Transmitter input clock frequency	OL - 100pi				μS	
	1x baud rate		DC		56	kHz	
	16x and 64x baud rate		DC		520	kHz	
trew	Transmitter input clock pulse width						
	1x baud rate 16x and 64x baud rate		12			tcy	
t _{TPD}	Transmitter input clock pulse delay		1			tcy	
	1x baud rate		15				
	16x and 64x baud rate		3			t _{CY}	
f_{Rx}	Receiver input clock frequency						
	1x baud rate 16x and 64x baud rate		DC		56	kHz	
t _{RPW}	Receiver input clock pulse width		DC		520	kHz	
VIII VV	1x baud rate		40				
	16x and 64x baud rate		12			toy	
tRPD	Receiver input clock pulse delay			· · · · · · · · · · · · · · · · · · ·		tcy	
	1x baud rate		15		ĺ	tcy	
	16x and 64x baud rate		3			tcy	
t _{Tx}	TxRDY delay from center of data bit	C _L = 50pF			16	tcy	
t _{Rx}	RxRDY delay from center of data bit				20	tcy	
tis	Internal SYNDET delay from center of data bit				25	tcy	
tes	Internal SYNDET set-up time before falling edge of RxC				16	tcy	
t _{TxE}	TxEMPTY delay from center of data bit	$C_L = 50pF$			16	tcy	
twc	Control delay from rising edge of WRITE (TxE, DTR, RTS)				16	tcy	
tor TES	Control to READ set up time (DSR, CTS)				16	tcy	

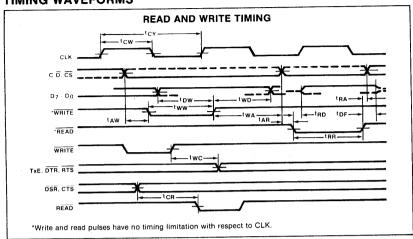
^{1.} AC timings measured at V_{OH} = 2.0V, V_{OL} = 0.8V, and with test load circuit.

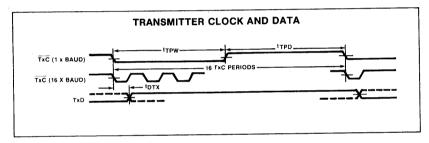
^{2.} This recovery time is for initialization only, when MODE, SYNC1, SYNC2, COMMAND and first DATA BYTES are written into the USART. Subsequent writing of both

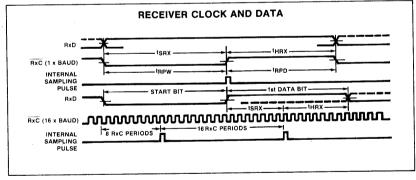
COMMAND and DATA are only allowed when TxRDY = 1.

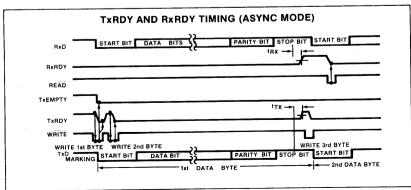
^{3.} The $\overline{\text{TxC}}$ and $\overline{\text{RxC}}$ frequencies have the following limitations with respect to CLK: for 1x baud rate, f_{Tx} or $f_{Rx} \le 1/30 \ t_{CY}$) for 16x and 64x baud rate, f_{Tx} or $f_{Rx} \le 1/4.5 t_{CY}$)

TIMING WAVEFORMS

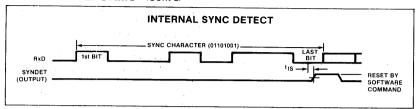


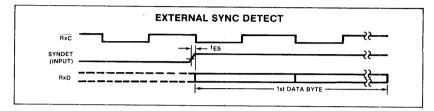




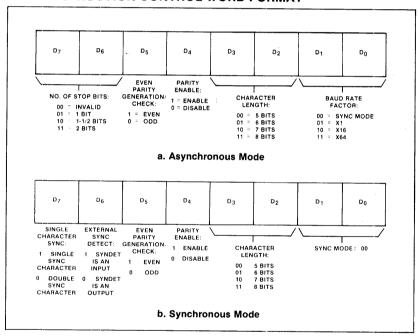


TIMING WAVEFORMS (Cont'd)

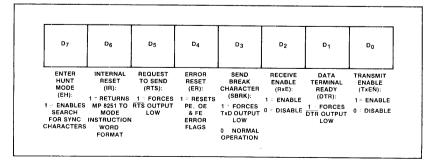




MODE INSTRUCTION CONTROL WORD FORMAT

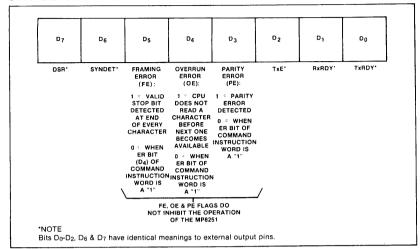


COMMAND INSTRUCTION CONTROL WORD FORMAT



9

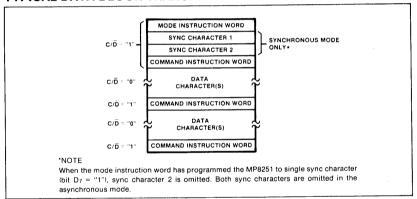
STATUS READ WORD FORMAT



MP8251 STATUS

The MP8251 has provisions for allowing the programmer to read the status of the device at any time during the functional operation. When the C/D input is a high-level, a normal read operation is executed to read this status information. The figure below shows the bits in the Status Read Word format. Since some of the status word bits have identical meaning to external output pins, the MP8251 can be used in a completely polled environment or in an interrupt driven environment.

TYPICAL DATA BLOCK TRANSFER



DESCRIPTION

The MP8255 is a programmable peripheral interface contained in a standard, 40-pin dual-in-line package. The chip, which is fabricated using N-channel silicon gate technology, functions as a general-purpose parallel input/output interface in Signetics MP8080A microcomputer family. The functional configuration of the MP8255 is programmed by the system software so that normally no external logic is required to interface peripheral devices.

The MP8255 has three basic modes of operation that can be selected by the system software. In the first mode (Mode 0), the MP8255 provides simple input and output

operations for three 8-bit ports. Data is simply written to or read from a specified port (Port A, B, or C) without the use of "handshaking" signals. In the second mode (Mode 1), the MP8255 enables the transfer of input/output data to or from a specified 8-bit port (Port A or B) in conjunction with strobes or "handshaking" signals. Ports A and B use the lines of Port C in this mode to generate or accept the "handshaking" signals with the peripheral device. In the third mode (Mode 2), the MP8255 enables communications with a peripheral device or structure via one bidirectional 8-bit bus port (Port A). "Handshaking" signals are provided over the lines of Port C in this mode to maintain proper bus flow discipline.

ABSOLUTE MAXIMUM RATINGS

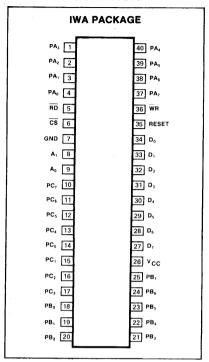
	PARAMETER	RATING	UNIT
TA TSTG	Ambient temperature under-bias Storage temperature Voltage on any pin with respect to ground	0 to +70 -65 to +150 -0.5 to +7	°C °C V

AC ELECTRICAL CHARACTERISTICS $T_A = 0$ °C to +70°C; $V_{CC} = +5V \pm 5\%$;

 $V_{SS} = 0V$

		V 55 - U	•		
	PARAMETER		LIMITS		
	FARAMETER	Min	Тур	Max	UNIT
twp	Pulse width of WR	430			ns
t_{DW}	Time D.B. stable before WR	10	l	ŀ	ns
t_{WD}	Time D.B. stable after WR	65			ns
taw	Time address stable before WR	20	Ī		ns
twa	Time address stable after WR	35			ns
tcw	Time CS stable before WR	20			ns
twc	Time CS stable after WR	35			ns
twB	Delay from WR to output			500	ns
trp	Pulse width of RD	430			ns
tin	RD set-up time	50			ns
tHR	Input hold time	50			ns
t _{RD}	Delay from $\overline{RD} = 0$ to system bus			350	ns
top	Delay from $\overline{RD} = 1$ to system bus	150			ns
tar	Time address stable before RD	50			ns
tcr	Time CS stable before RD	50			ns
tak	Width of ACK pulse	500			ns
tst	Width of STB pulse	350			ns
tps .	Set-up time for peripheral	150			ns
tpH	Hold time for peripheral	150			ns
tra	Hold time for A_1 , A_0 after $RD = 1$	379			ns
trc	Hold time for \overline{CS} after $\overline{RD} = 1$	5			ns
tAD	Time from $\overline{ACK} = 0$ to output (Mode 2)			500	ns
tĸD	Time from ACK = 1 to output floating	l l		300	ns
two	Time from $\overline{WR} = 1$ to $\overline{OBF} = 0$			300	ns
tao	Time from $ACK = 0$ to $\overline{OBF} = 1$			500	ns
tsı	Time from STB = 0 to IBF			600	ns
tRI	Time from $\overline{RD} = 1$ to $IBF = 0$			300	ns

PIN CONFIGURATION



FEATURES

- Outputs source 1mA at 1.5 volts
- 24 programmable input/output pins
- Direct bit set reset capability
- TTL compatible
- Reduces system component count

PIN DESIGNATION

MNEMONIC	PIN NO.	TYPE			NAME AND FUNCTION					
<u>cs</u>	6	1	Chip Select: When low, the chip is selected. This enables communication between the MP8255 and the MP8080A microprocessor.							
RD	5	1	Read: When lo	w, allows the MP	8080A to read data or status inform	ation from the MP8255.				
WR	36	ı	Write: When lo	w, allows the MP	8080A to write data or control word	ds into the MP8255.				
A ₀ ,A ₁	9,8	l	the A ₁₅ - A ₀ Ad	Port Select: These two inputs, which are normally connected to the least significant bits of the A ₁₅ - A ₀ Address Bus, control the selection of one of three 8-bit ports (A, B, and C) or the internal control word register as indicated below.						
			A ₁	Ao	Selected					
			0	0	Port A					
			0	1	Port B					
			1	0	Port C					
			1	1	Control word register					
RESET D7 - D0	35 27 - 34	1/0	input high impu Data Bus: This al communicat internal data b MP8080A. In ac bus buffer.	edance mode. bus comprises eightion between the sus buffer upon eightion, control wo	nternal registers of the chip and set ght tri-state input/output lines. The b MP8255 and the MP8080A. Data i execution of an Out or In instruction ords and status information are trans	ous provides bidirections routed to or from the on, respectively, by the sferred through the data				
PA ₇ - PA ₀	37 - 40, 1 - 4	1/0	data input latci	h.*	ort comprises one 8-bit data output l					
PB7 - PB0	25 - 18	1/0	and 8-bit data	buffer.	port comprises one 8-bit data input					
PC ₇ - PC ₄ PC ₃ - PC ₀	10 - 13 17 - 14	1/0	bit data input b these 4-bit por conjunction wi (see figure) for a status/contro	Port C: This 8-bit input/output port comprises one 8-bit data output latch/buffer and one 8-bit data input buffer. The port can be split into two 4-bit ports under the mode control. Each of these 4-bit ports contains a 4-bit latch that may be used for the control and status signals, in conjunction with Ports A and B. The system software includes a Bit Set/Reset Control Word (see figure) for setting or resetting any of the eight bits of Port C. When Port C is being used as a status/control for Port A and B, the Port C bits can be set or reset by using the Bit Set/Reset Control Word as the second byte of OUT Instruction(s).						
V _{CC} GND	26 7	I	V _{CC} Supply: Ground: 0-vol	+5 volts t reference						

^{*}NOTE

DC ELECTRICAL CHARACTERISTICS $T_A = 0$ °C to +70°C; $V_{CC} = \pm 5V \pm 5\%$; $V_{SS} = 0V$

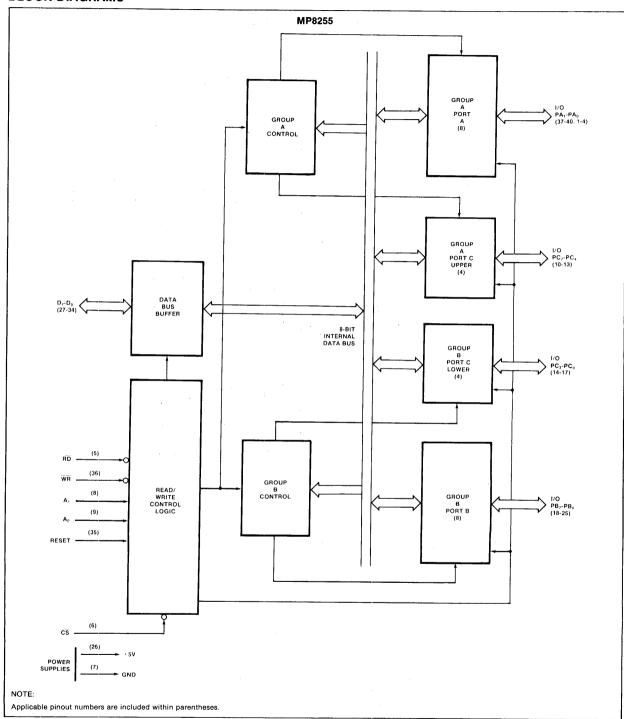
				UNIT		
	PARAMETER	TEST CONDITIONS	Min	Тур	Max	UNI
· · · · · · · · · · · · · · · · · · ·	Input voltage		2.0			٧
V _{IH} V _{IL}	High Low				0.8	
Voh Vol	Output voltage High Low	$I_{OH} = -50\mu A (-100\mu A \text{ for D,B, Port})$ $I_{OL} = 1.6\text{mA}$	2.4		0.4	٧
Ioн*	Darlington drive current Power supply current	$V_{OH} = 1.5V, R_{EXT} = 390\Omega$		2.0 40		mA mA

^{*}NOTE

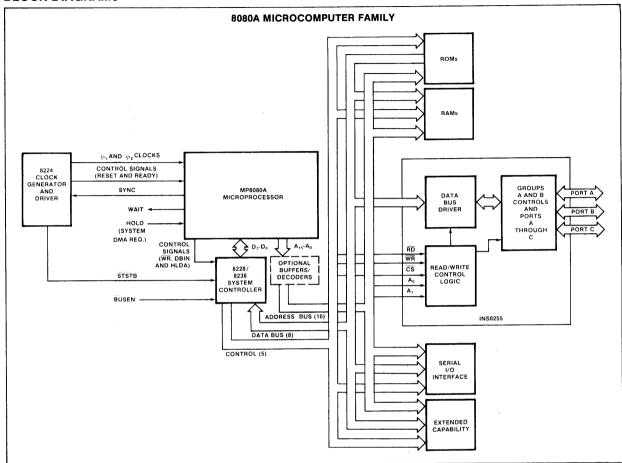
Available on 8 pins only of ports A and C, selected randomly.

The system software uses a Mode Definition Control Word (see figure) as the second byte of OUT Instruction(s) to program the functional configuration of Ports A through C. Whenever the mode is changed, all output registers (and status flip flops) are reset.

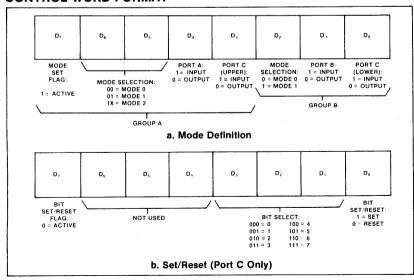
BLOCK DIAGRAMS



BLOCK DIAGRAMS



CONTROL WORD FORMAT

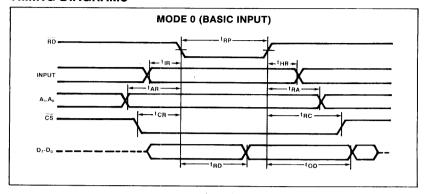


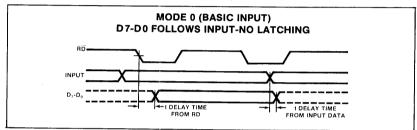
OPERATING MODES

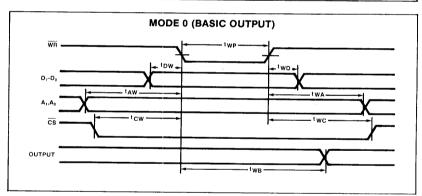
Mode 0 (Basic Input/Output)

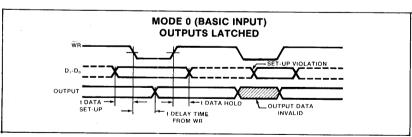
In this mode, simple input and output operations for each of the three ports are provided. No "handshaking" is required; data is simply written to or read from a specified port.

TIMING DIAGRAMS









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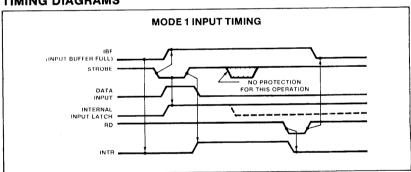
MODE 0 PORT DEFINITION CHART

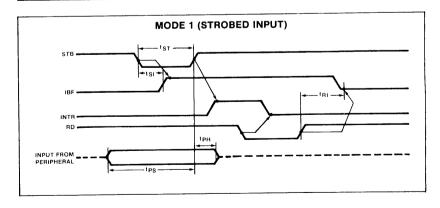
	CONTROL WORD BITS							GRO	OUP A	GROUP B		
NO.	D ₇	D ₆	D ₅	D ₄	D ₃	D ₂	D ₁	D ₀	Port A	Port C (Upper)	Port B	Port C (Lower)
0	1	0	0	0	0	0	0	0	Output	Output	Output	Output
1	1	0	0	0	0	. 0	0	1	Output	Output	Output	Input
2	1	0	0	0	0	0	1	0	Output	Output	Input	Output
3	1	0	0	0	0	0	1	1	Output	Output	Input	Input
4	1	0	0	0	1	0	0	0	Output	Input	Output	Output
5	1	0	0	0	1	0	0	1	Output	Input	Output	Input
6	1	0	0	0	1	0	1	0	Output	Input	Input	Output
7	1	0	0	0	1	0	1	1	Output	Input	Input	Input
8	1	0	0	1	0	0	0	0	Input	Output	Output	Output
9	1	0	0	1	0	0	0	1	Input	Output	Output	Input
10	1	0	0	1	0	0	1	0	Input	Output	Input	Output
11	1	0	0	1	0	0	1	1	Input	Output	Input	Input
12	1	0	0	1	1	0	0	0	Input	Input	Output	Output
13	1	0	0	1	1	0	0	1	Input	Input	Output	Input
14	1	0	0	1	1	0	1	0	Input	Input	Input	Output
15	1	0	0	1	1_	0	1	1	Input	Input	Input	Input

Mode 1 (Strobed Input/Output)

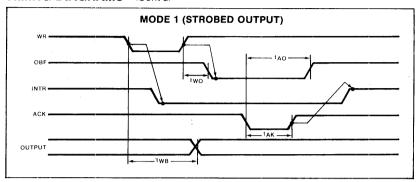
In this mode, a means for transferring input/output data to or from a specified port in conjunction with strobes or "handshaking" signals is provided. Port A and Port B use the lines on Port C to generate or accept these "handshaking" signals in Mode 1. The programmer can read the contents of Port C to test or verify the status of each peripheral device. Since no special instruction is provided in the MP8080A microcomputer system to read the Port C status information, a normal read operation must be executed to perform this function.

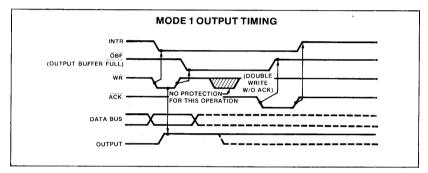
TIMING DIAGRAMS

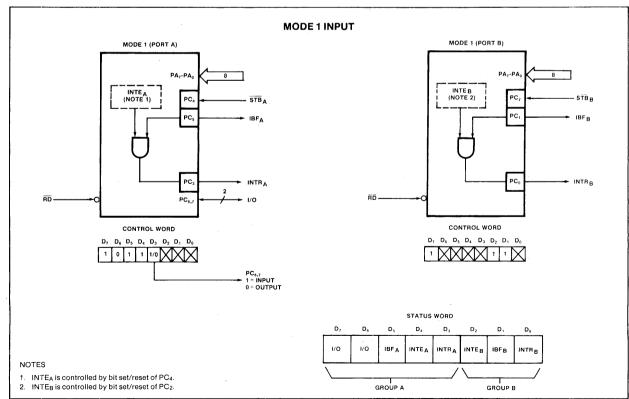




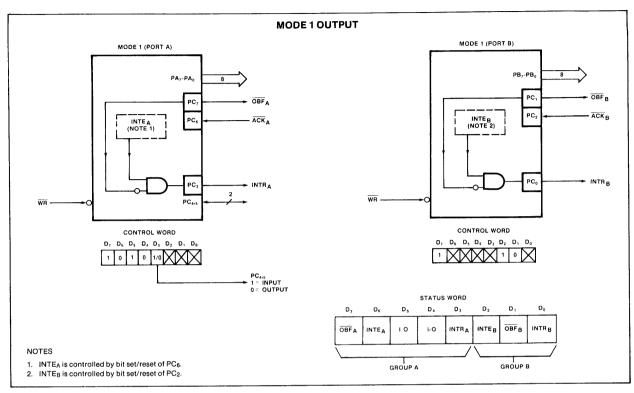
TIMING DIAGRAMS (Cont'd)

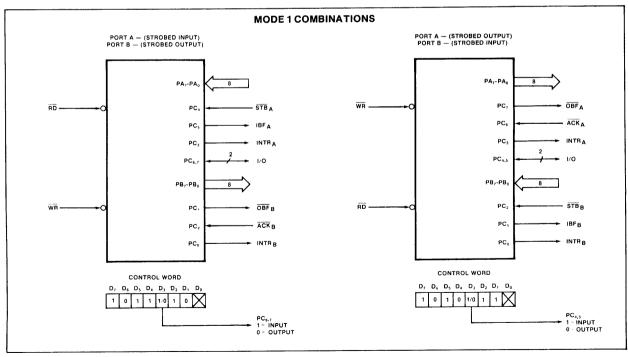






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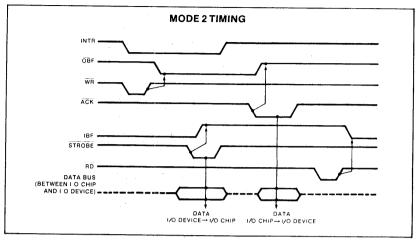


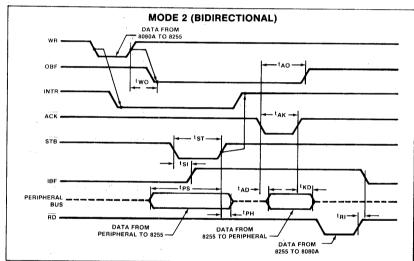


Mode 2 (Strobed Bidirectional Bus Input/Output)

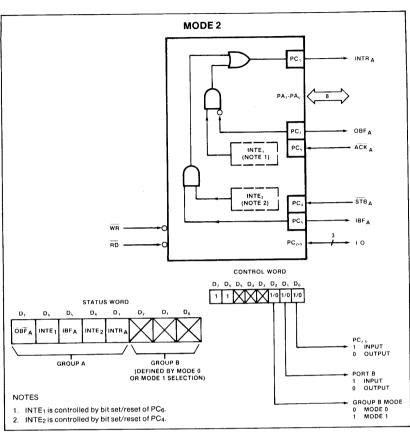
This mode enables communication with a peripheral device or structure on a single 8-bit bus for both transmitting and receiving data (bidirectional bus input/output). "Handshaking" signals are provided to maintain proper bus flow discipline in a manner similar to Mode 1. In addition, interrupt generation and enable/disable functions are available in Mode 2.

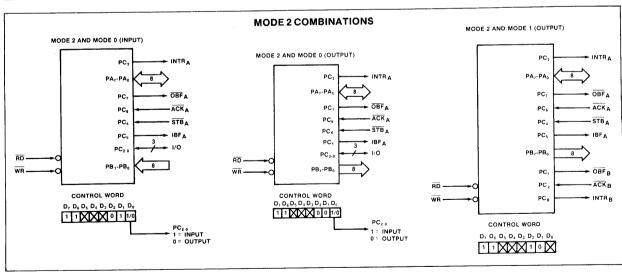
TIMING DIAGRAMS



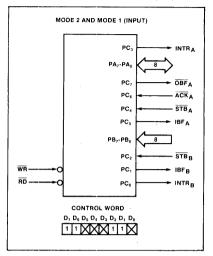


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MODE 2 COMBINATIONS (Cont'd)



MODE DEFINITION SUMMARY TABLE

PORT	MO	DE 0	мог	DE 1	MODE 2
BITS	IN	OUT	IN	OUT	GROUP A ONLY
PA ₀	ln'	Out	In	Out	Bidirectional
PA ₁	In	Out	In	Out	I • • • • • • • • • • • • • • • • • • •
PA ₂	In	Out	ln	Out	· ·
PA ₃	In	Out	In	Out	
PA ₄	¹ In	Out	In	Out	
PA ₅	In	Out	In	Out	
PA ₆	In	Out	In	Out	1
PA ₇	In	Out	In	Out	Bidirectional
PB ₀	In	Out	In	Out	
PB ₁	In	Out	In	Out	•
PB ₂	. In	Out	In	Out	
PB ₃	In	Out	In	Out	(Mode 0 or Mode 1 only)
PB ₄	In	Out	In	Out	(Mode of Mode 1 only)
PB ₅	In	Out	In	Out	4
PB ₆	- In	Out	In	Out	
PB ₇	In	Out	In	Out	
PC ₀	In	Out	INTRB	INTRB	1/0
PC ₁	In	Out	IBFB	OBFB	1/0
PC ₂	In	Out	STBB	ACKB	1/0
PC ₃ .	In	Out	INTRA	INTRA	INTRA
PC ₄	In	Out	STBA	1/0	STBA
PC ₅	In	Out	IBF₄	1/0	<u>IBF</u> _A
PC ₆	In	Out	1/0	ACKA	ACKA
PC ₇	In	Out	1/0	OBFA	OBFA

2650 MICROPROCESSOR SERIES

DESCRIPTION

The KT9000 kit contains a 2650 microprocessor and enough chips to allow for the implementation of a small developmental system. Since the interface requirements of the 2650 are completely TTL compatible, no attempt has been made to limit the user's flexibility by dictating a fixed logic configuration. There is complete freedom in using standard SSI or MSI logic to adapt the microprocessor to the memory, I/O devices, or clock.

Several minimal system examples are presented to enable quick set up and evaluation. Other configurations to adapt to individual requirements should become evident from these examples.

PARTS DESCRIPTIONS

2112: The 2112 is a static 1024-bit Random Access Memory organized as 256 words by 4 Bits/Word. It is fabricated with N-Channel, Silicon Gate, MOS technology and achieves an access time of less than 800 nanoseconds. No clocks are required, and the chip is powered from a single 5 volt source.

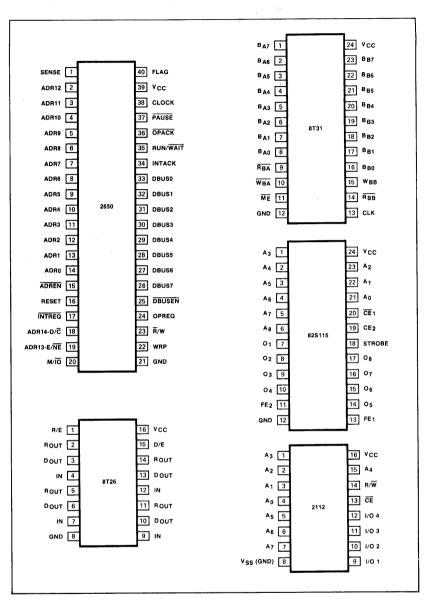
82S115I: The 82S115I is a 4096-bit Schottky-Clamped, Bipolar Read Only Memory, incorporating on-chip data output registers. It is field-programmable and fully TTL compatible with on-chip decoding and two chip enable inputs for ease of memory expansion. Inputs to the device are PNP transistors with a maximum current requirement of $100\,\mu\text{A}$.

8T31: The 8T31 is an 8-bit Bidirectional I/O Port designed to function as a general purpose I/O interface element. It consists of 8 clocked latches with two sets of bidirectional Inputs/Outputs. The capability exists for various hook-up schemes allowing master control from either the microprocessor or from the I/O device.

8T26B: The 8T26B consists of four pairs of inverting Tri-State Logic elements configured as a Quad Bus Drivers/Receivers with separate buffered receiver enable and driver enable lines. Both the driver and receiver gates have Tri-State outputs and low-current PNP inputs.

CIRCUIT EXAMPLES

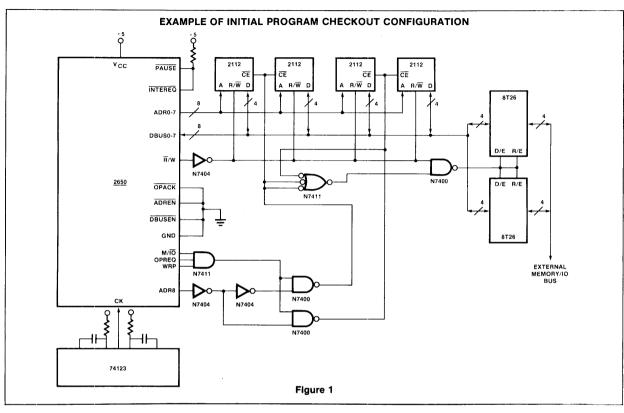
Two circuit configurations are presented to indicate a possible program checkout approach. Figure 1 is hooked up to allow the use of RAM for program debugging. The second figure represents a possible final system configuration with the program fixed in PROM. Both circuits use the 8T26's as bus buffers.

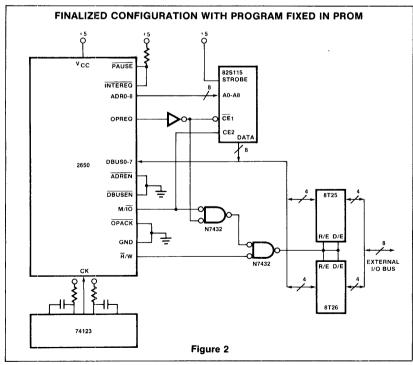


PARTS LIST

PART NO.	QTY	DESCRIPTION	REFERENCE DATA SHEET
2650	1	CPU	_
2112	4	256X4 RAM	MOS Products
82S115I	1	4K PROM (Unprogrammed) 512X8	Bipolar Memories
8T31I	2	8-bit Bidirectional I/O Port	8000 Product
8T26B	4	Quad Bus DR/REC	8000 Product
2650BM1000	1	Basic Manual	_

2650 MICROPROCESSOR SERIES

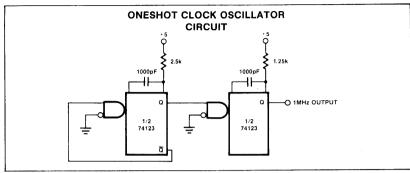


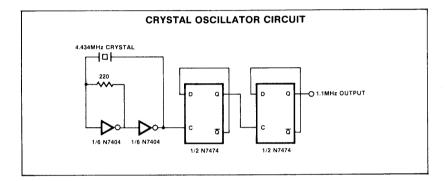


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2650 MICROPROCESSOR SERIES

TYPICAL APPLICATIONS





DESCRIPTION

The Signetics Adaptable Board Computer, ABC 1500, is a modular microcomputer containing a CPU, memory, I/O ports and support circuitry. It is designed to cover a broad range of applications from software development to system hardware prototyping. Cost performance trade offs have been carefully considered to achieve maximum flexibility and allow the card to be tailored to a variety of individual requirements.

The basic configuration consists of the 2650 8-bit microprocessor, 512 bytes of read/write memory (four 2112 static RAM's), 1K bytes of 2608 ROM with PIPBUG*, two 8T31 I/O ports and buffering on data, address and control lines. A single +5 volt supply will be required to power the card and communicate with a serial 20mA current loop terminal.

Modifications to the basic system can be easily made to allow for various memory configurations and operating modes. Unused plated through holes are provided for the PROM memory chips (82S115's). Other options are jumper selectable.

The ABC 1500 is sold either as a completely assembled and tested card (2650PC1500) or in kit form (2650KT9500).

*PIPBUG is a loader, editor, and debug program. See Table 1.

FEATURES

- Expandable printed circuit card: Unused area on card filled with plated through holes on .300 in. centers for wirewrap sockets
- 1K bytes of PIPBUG* ROM (in socket)
- 512 bytes of RAM
- Two latched I/O ports
 - —four non-extended I/O read/write user strobes
- Tri-state buffers on data, address and control lines
- · Serial input/output port
- Single +5 volt supply requirement (1.7A max) for card and 20mA current loop interface (±12 volt supply for RS232 interface)
- Interrupt and single step capability
- Simple clock configured from dual monostable multivibrator
- 24K memory expansion capability
- Directly compatible with 4K RAM card (2650PC2000) and power supply demonstration base (2650DS2000)
- Card dimensions: 8 in. by 6.875 in. with a 100 pin connector along the 8 in. dimension

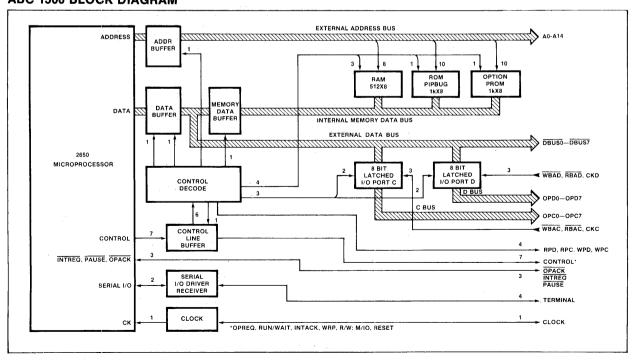
OPTIONS

- 1K bytes of PROM in place of ROM
- 512 bytes of PROM or ROM in place of RAM
- · Asynchronous operation capability
- External clock input
- Interrupt vector from Port C

INTERFACE

- Terminal interface jumper selectable
 A. W4 to W5 and W6 to W7 jumpers select
 the 20mA current loop mode
 - B. W3 to W4 and W7 to W8 jumpers select the RS232 mode
- Normally high input lines (10k pull up resistor on each): INTREQ, PAUSE, RESET, WBAC, WBAD, CKC, CKD
- Plated through holes are available at each connector pin to allow for insertion of wirewrap pins.
- · Edge connector supplied with card
- To allow for external clock input, remove iumper W9-W10
- Asynchronous operation by removing jumper W1 to W2 and driving OPACK
- During vectored interrupts, it is possible to allow port C to place the interrupt address on the data bus by removing jumper W21-W22 and jumpering W22-W23

ABC 1500 BLOCK DIAGRAM



ALPHA CHAR- ACTER INPUT	COMMAND							
Α	Alter memory							
В	Set breakpoint							
С	Clear breakpoint							
D	Dump memory to papertape							
G	Go to address							
L	Load memory from papertape							
S	See and alter registers							

NOTE

The program is entered by resetting the card. The terminal will then respond with an asterisk (*).

Table 1 PIPBUG COMMANDS

MEMORY CONFIGURATION

All of page 0 is reserved for on card memory (0 to 8191₁₀). Address lines A9, A11 and A12 are not decoded (Don't care signals) allowing two IC's to perform not only memory decoding but also I/O port decoding. As an added benefit, usable memory space exists at the top of page 0 (see Table 2) due to the interleaving effect between the ROM and RAM memories. This memory space can be used as interrupt vector address locations in a negative direction from address location

"0" (a negative relative instruction from address location "0" wraps around the first 8K page).

There are a total of two blocks in the RAM structure, each of which contains 256 bytes of RAM. Since PIPBUG uses the first 63 RAM locations for temporary storage, the first actual user location is 1087₁₀ (43F₁₆) (there are seven other address locations corresponding to the first user location—see memory map). Starting at 43F₁₆, the range for on card RAM extends to address 1535₁₀ (5FF₁₆), giving a total usable on-card space of 449 bytes.

The first external memory location for addon memory is 8192₁₀ (2000₁₆). All of page 1, 2, and 3 are available, giving a total memory expansion capability of 24K.

MEMORY OPTIONS

Modifications to the basic configuration can be made to provide a mix of RAM/PROM/ROM memories. PROM memories can be used in place of the PIPBUG ROM by removing the ROM from its socket and adding one or two 82S115 PROMs (512X8). Area and plated through holes are provided on the card for insertion of sockets for the

PROMs or the PROMs themselves. Decoding for the PROMs has been provided by ABC 1500 logic.

Data and address lines for 2112 RAMs and 82S129 PROMs or 82S229 ROMs are identical. It is, therefore, possible to use PROMs and/or ROMs in place of RAM. This option will require removal of the RAMs (two per block), and changing the jumper for each 256 byte block of PROM or ROM added. The jumper needed for each block of memory is as follows:

MEMORY SECTION	RAM JUMPER	PROM/ROM JUMPER
First Block	W12-W13	W11-W13
Second Block	W15-W16	W14-W16

I/O PORT CONFIGURATION

Two ports (C and D) are implemented with 8T31 bidirectional ports and can be used for general purpose I/O. Each consist of 8 clocked latches with two sets of bidirectional inputs/outputs. Data written into one side of the port will appear inverted at the other side.

One side of each port (Bus B of the 8T31) is tied to the external data bus (DBUSX). The 2650 communicates with each port over this

		ADDRES	LINES			DECIMAL ORGANIZATION		HEX ADDRESS
A14	A13	* A12	* A11	A10	** A9	8K	Second block RAM	1FFF
0	0	х	х	1	x		First block RAM Second block RAM	
0	0	×	х	0	х	7 K	First block RAM PIPBUG ROM	1BFF
0	0	x	X	1	X	6K	Second block RAM First block RAM	17FF
_						5K	Second block RAM First block RAM	13FF
0	0	X	X	0	Х	4K	PIPBUG ROM Second block RAM	0FFF
0	0	х	х	1	х	3K	First block RAM Second block RAM First block RAM	
0	0	×	×	0	x		PIPBUG ROM	0BFF 07FF
0	0	×	х	1	х	2K	Second block RAM First block RAM Second block RAM	06FF 05FF
0		x	х		X	1K	First block RAM PIPBUG ROM	04FF 03FF
	0	<u> </u>	^	0	<u> </u>	0	PIPBUG ROW	0000

NOTES

- 1. * = Don't care for ROM and RAM; ** = Don't care for RAM.
- 2. Each block of RAM = 256 bytes.

Table 2 PAGE 0 MEMORY MAP

bus with one byte, non-extended I/O instructions. During 2650 activity with the ports, the ABC 1500 will provide four output strobes indicating the nature of the operation.

STROBE	FUNCTION	STROBE PULSE WIDTH
WPC	Write to Port C	Duration of WRP
WPD	Write to Port D	Duration of WRP
RPC	Read Port C	Duration of OPREQ
RPD	Read Port D	Duration of OPREQ

The other side of each port (Bus A of the 8T31) is controlled by the user. Four control lines are used to read, write or tri-state the buses.

CONTRO	L LINE	FUNCTION	
WBAD (1)	RBAD	(0)	Read Port D
WBAD (0)	RBAD	(0)	Write to Port D
WBAD (1)	RBAD	(1)	Tri-state D Bus
WBAC (1)	RBAC	(0)	Read Port C
WBAC (0)	RBAC	(O)	Write to Port C
WBAC (1)	RBAC	(1)	Tri-state C Bus

If no external logic is connected each port will be in the "read" mode (WBAX lines pulled high). The RBAX lines are tied to ground to allow read/write control of the buses with just the WBAX lines. To allow control for tri-stating the buses, the following jumpers must be removed:

LINE	JUMPER
RBAC	W19-W20
RBAD	W17-W18

The clock for each port (CKC—Port C clock, CKD—Port D clock) is available at a connector pin for external control. These normally "high" lines can be pulled low to disable writing to the ports from either the 2650 or the external device.

ABC 1500 EDGE CONNECTOR SIGNAL LIST

BC 1300 I	EDGE CONNECTOR
PIN #	FUNCTION
1	GND
2	GND
3	_NC*
4	DBUS0
5	DBUS1
6	DBUS2
7	DBUS3
8 9	DBUS4
9 10	DBUS5 DBUS6
11	DBUS7
12	NC*
13	A14—D/C
14	NC*
15	A13—E/NE
16	INTACK
17	R/W
18	WRP
19	RUN/WAIT
20	OPR <u>E</u> Q
21	M/IO
22	OPACK
23	CLOCK
24	TS
25 26	RESET
27	INTREQ PAUSE
28	NC*
29	RBAD
30	NC*
31	RBAC
32	NC*
33	All
34	A13—E/NE
35	A12 _
36	A14—D/ C
37 38	A9
39	A10 A8
40	A0 A7
41	A6
42	A5
43	A3
44	Α0
45	A1
46	A4
47	A2
48	+12V
49 50	-12V
30	+5V

IAL LIST	
PIN#	FUNCTION
A B C D E F H J K L M N P R S T U V W X Y Z a b c d e f g h j k E n p r s t u v W x y z a b c d e f	GND GND NC* OPD 0 OPD 1 OPD 2 OPD 3 OPD 4 OPD 5 OPD 6 OPD 7 NC* TTY SERIAL IN + TTY SERIAL OUT + TTY SERIAL OUT + TTY SERIAL OUT - R\$232 GROUND R\$232 OUTPUT NC* NC* R\$232 INPUT NC* OPC 0 OPC 1 OPC 2 OPC 3 OPC 4 OPC 5 OPC 6 OPC 7 NC* RPD WBAD WPD CKD NC* NC* NC* NC* NC* NC* NC* NC* NC* NC*

*NC = No connection

ABC 1500 PARTS LIST

QUANTITY	DESCRIPTION
1	PC1500 printed circuit board
1	Edge connector - AMP 225-804-50
1	2650 8-bit static microprocessor
1	N7402 quad 2 input NOR gate - I/O strobe logic
1	N7416 hex inverter buffer - current loop interface
1	N74123 monostable multivibrator - clock for 2650
1	N74S138 3 line to 8 line decode - control decode
1	8T15—EIA line driver - RS232 driver
4	8T26 quad Tri-state driver/receiver - data and memory data buffer
2	8T31 8-bit latched bidirectional I/O port
4	8T97 hex Tri-state driver - address and control line buffer
1	2608 static ROM (1024X8) - PIPBUG ROM - CN0035
4	2112 static RAM (256X4) - organized as 512 byte RAM
1	82S123 PROM (32X8) coded PROM CD 1500 - control decode
4	1N914 diode
1	2N2222 transistor
1	50pF capacitor
1 1	300pF capacitor
13	.1μF capacitor
3	4.7μF capacitor
2	220 OHM resistor
6	1K resistor
2	2K resistor
1	3.3K resistor
8	10K resistor
1	20K resistor
1	24 pin 2608 ROM socket - Robinson-Nugent ICN 246-54

NOTE

All resistors 1/4 watt.

DESCRIPTION

The Signetics PC-4000 is an emulation of the Signetics 2656, a 40-pin NMOS-LSI system memory interface chip. The PC-4000, in circuit board form, offers the engineer a system design aid. By designing with the PC-board emulator the specific ROM and PGA (programmable gate array) patterns required for the user's application, can be determined.

In utilizing the PC-4000, the engineer-designer is able to implement the same functions as the 2656. Through a 40-pin plug and cable attached to the PC-4000, which has identical pin-outs to the 2656, the user connects the PC-4000 directly into his prototype system. He can access the 128X8 RAM and the PROM. The user can use the eight multi-function ports either as I/O ports or chip enables. Also he may use the poweron reset, and the clock generator and divider which are included on the PC-4000.

This data sheet contains:

- Circuit Emulation Description
- Functional Block Diagram
- Functional Block Descriptions
- Operation and Use

Load and Drive Characteristics PC-4000 and 2656 Timing Comparisons

PC-4000 Component Positioning Diagram

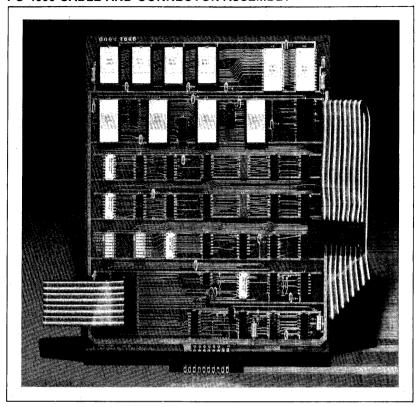
- Programming
 - System Address Map for Coding the FPLAs
 - FPLA Coding Forms
 Coding Switches and Jumpers
- Circuit Diagram and Connectors
- Parts List
- Shipping Configuration

CIRCUIT EMULATION DESCRIPTION

The PC-4000 emulator contains the same circuits that are available to the user on the 2656 Systems Memory Interface chip. Functionally the PC-4000 replaces the 2656 in the user's prototyping system through a pinfor-pin compatible plug and its 40-wire ribbon cable attached to the PC-4000.

In emulating the 2656, the speed of the board circuitry is equivalent to or faster than the on-chip circuitry. ROM is implemented with bipolar PROMs and on-chip RAM is implemented with bipolar RAM packages. The programmable gate array (PGA) for selecting of the ROM and RAM enables and the I/O function selects are implemented with field programmable logic arrays (FPLAs).

PC-4000 CABLE AND CONNECTOR ASSEMBLY



The oscillator provided on the PC-4000 has the same frequency dividers available as on the 2656. The oscillator divide ratio is selected by toggle switch settings.

FUNCTIONAL BLOCK DESCRIPTIONS

Input Buffer: This circuitry buffers the incoming signals from the MPU (specifically for the 2650 MPU, the addresses, A0-A14, and the 4 control signals, OPREQ, WRP, \overline{R}/W and M/\overline{IO}).

Data Bus Transcelver: This circuitry buffers the incoming 8-bit data bus from the MPU, and provides output drivers to the data bus.

FPLAs 1 and 2: FPLAs 1 and 2 represent a portion of the programmable gate array of the 2656 SMI chip. These 2 FPLAs decode the chip enable signals.

FPLAs 3 and 4: These FPLAs represent the remainder of the gate array. FPLA 3 and a portion of FPLA 4 generate the on-board ROM, RAM and port enable signals. The remainder of FPLA 4 is used to generate the data bus control signals.

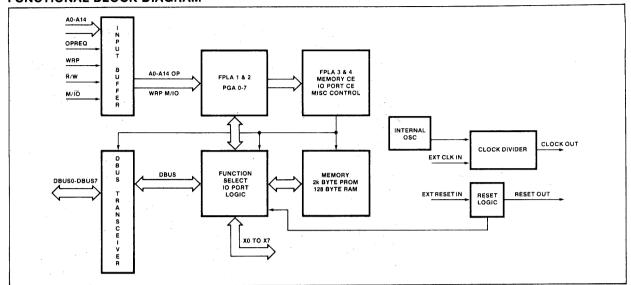
Function Select and I/O Port Logic (Via FPLAs 1-4): The function select and I/O port logic allow each of the multi-purpose pins of the 2656, to be individually selected as either an I/O port or a chip enable via eight switches, FS0-FS7. When assigned as chip enables, they must be programmed in FPLAs 1 and 2. When a multi-purpose pin is assigned as I/O or as an input port, the port address is programmed in FPLAs 1 and 2.

When the FS switch is On, the corresponding pin of the 2656 is selected as an I/O port. When the FS switch is On, the corresponding pin is selected as a chip enable.

Memory: Both ROM and RAM memory functions are implemented in bipolar PROM and RAM respectively. The memory chip enables are programmed in FPLAs 3 and 4. (See the Programming section)

Clock Divider, Reset Logic, and Internal Oscillator: This circuitry provides an internal oscillator with switches for frequency divide by 1, 2, 3, or 4. Reset logic is provided on the PC-4000 to allow the user to reset the system during debug via an external switch

FUNCTIONAL BLOCK DIAGRAM



closure, without the need for powering down. The reset logic is used when the RC or external oscillator modes are selected.

If the RC or crystal internal oscillator mode is desired, the frequency determining components must be installed on the emulator PCB. The pins on the header that correspond to the RC and Crystal pins are not used in this mode, and only the external Reset function is provided.

OPERATION AND USE

After preprogramming the on-board PROMs and FPLA with the desired user data, the operation of the PC-4000 consists of supplying electrical power from the external source and connecting into the user's system.

Power Supply Requirements: The PC-4000 requires 5Vdc at 2.2Adc typical. The maximum current requirement is 3.5Adc.

Load and Drive Characteristics: The load characteristics of PC-4000 inputs presented to any bus input are $400\mu Adc$ for a zero and $40\mu Adc$ for a 1. The PC-4000's output lines will drive a minimum of 5 TTL loads, which far exceeds the drive characteristics of the 2656 SMI. Thus the user must exercise care when changing from the emulator to the 2656 to avoid inadvertent loading problems.

PC-4000 and **2656** Timing Comparisons: The PC-4000 is implemented in TTL logic. All timing delays and Rise/Fall times are less than those of corresponding signals in the 2656.

Component Positioning Diagram: The

components, sockets, switches, and jumper connections that allow the user to adapt the PC-4000 to his specific system requirements are shown in Figure 1. Positioning Diagram. Each section of the PC-board containing components that the user must either insert, adjust, or program are shown as callouts. Only pins, 1,2/A, B (GND) and 9, 10 /K,L (+5Vdc) of the edge connector should be used. All other pins must be left open.

Programming and Operating Summary

- 1. Programming Address Functions:
 - A. Define Memory and I/O address map for system.
 - B. Program FPLA 1 and FPLA 2 with minimum programming configuration and X pin chip select information. See Coding Table for FPLA 1 and FPLA 2.
 - C. Program FPLA 3 and FPLA 4 with Memory ROM and RAM and I/O port addresses and minimum programming configuration. See Coding table for FPLA 3 and FPLA 4.

See Tables for Coding Switches:

- D. ts there ROM on SMI?

 If yes, program switch S2-2 (No ROM)

 OFF.

 If no, program switch S2-2 (No ROM)

 ON.
- E. Is there RAM on SMI?
 If yes, program switch S2-3 (No RAM)
 OFF.
 If no, program switch S2-3 (No RAM)

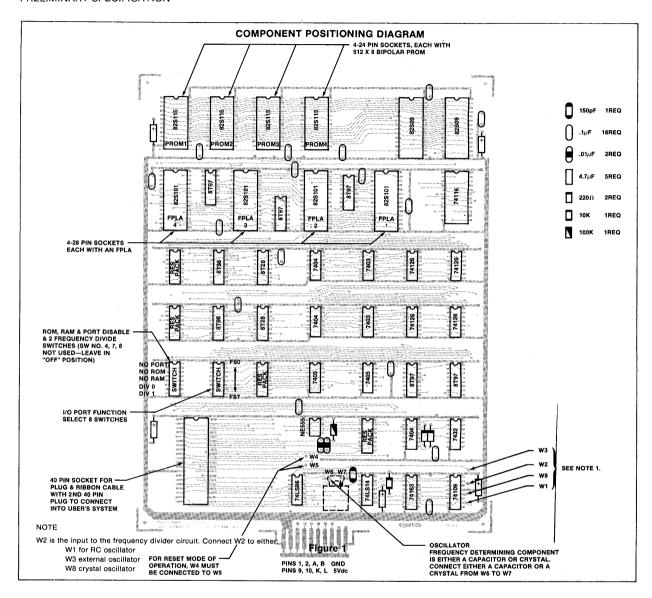
- F. Is the I/O Port on SMI used?
 If yes, program switch S2-1 (No Port)
 OFF.
 If no, program switch S2-1 (No Port)
- G. Program Switch S2-4 (1k2k) Off if using 2656 SMI chip.
- 2. Clock Divider and Reset Programming:
 - A. Is SMI Oscillator XTAL Driven?
 If yes, attach XTAL to board at W6 and
 W7. Jumper W8 to W2, skip to Step 2E.
 - B. Is SMI Oscillator RC Driven?
 If yes, calculate capacitor value and attach capacitor to board at W6 and W7. Jumper W1 to W2, skip to Step 2D.
 - C. Is SMI Oscillator Driven by an external source?
 - If yes, jumper W2 to W3.
 - D. Is SMI to provide system power on reset?
 - If yes, jumper W4 to W5.
 - E. Set Frequency Divider Switches
 Divide Oscillator Frequency by 1 set
 switches S2-5 Off, S2-6 Off
 Divide Oscillator Frequency by 2 set
 switches S2-5 On, S2-6 Off
 Divide Oscillator Frequency by 3 set
 switches S2-5 Off, S2-6 On
 Divide Oscillator Frequency by 4 set
 switches S2-5 On, S2-6 On
- 3. Programming X Pin Functions:
 - A. Is X0 Pin chip Select?

 If yes, program switch S1-1 Off.

 Is X0 Pin I/O Port?

 If yes, program switch S1-1 On.
 - B. Is X1 Pin chip Select?
 If yes, program switch S1-2 Off.

ON.



Is X1 Pin I/O Port?

If yes, program switch S1-2 On.

C. Is X2 Pin chip Select?

If yes, program switch S1-3 Off. Is X2 Pin I/O Port?

If yes, program switch S1-3 On.

D. Is X3 Pin chip Select?

If yes, program switch S1-4 Off. Is X3 Pin I/O Port?

If yes, program switch S1-4 On.

E. Is X4 Pin chip Select?

If yes, program switch S1-5 Off.

Is X4 Pin I/O Port?

If yes, program switch S1-5 On.

F. Is X5 Pin chip Select?

If yes, program switch S1-6 Off. Is X5 Pin I/O Port?

If yes, program switch S1-6 On.

G. Is X6 Pin chip Select?

If yes, program switch S1-7 Off. Is X6 Pin I/O Port?

If yes, program switch S1-7 On.

H. Is X7 Pin chip Select?

If yes, program switch S1-8 Off.

Is X7 Pin I/O Port? If yes, program switch S1-8 On.

4. PROM Programming:

A. Develop Application Software

B. Program PROMS

5. Actual Use

A. Connect power to Edge Connector

B. Connect umbilical cable to socket of 2656 in system

C. Turn power on and debug system

PROGRAMMING

To program the PC-4000 the user must have available both FPLA programming equipment and PROM programming equipment. Equipment for programming the PROMs and FPLAs may be obtained from Data I/O Corporation, Curtis Electro Devices, and Signetics, among others. Normally, in defining the user's system hardware, the information needed to program the FPLAs, the Function Select Switches, the Clock and the Reset emerge from the system design. The following programming procedures are suggested for each of the programmable functions. The PGA in its FPLA implementation is considered first. It consists of:

- · PGA input and output signal definitions. (See Table 1)
- System Address Map for coding the FPI As
- FPLA Coding Forms

For detailed instructions on the programming procedures and equipment required for using FPLAs refer to "Signetics Field Programmable Logic Arrays," February

The Programmable Gate Array (PGA): The PGA decodes the X0 thru X7 chip select signals in FPLA 1 and FPLA 2. The memory chip enable signals are decoded in FPLA 3. The I/O port select signals are decoded in FPLA 3 and FPLA 4, and the Data Bus Control Signal in FPLA 4. The PGA has a minimum programming configuration as shown in the coding tables. (Reference: Coding Forms and Schematic Diagram)

System Address Map for Coding the FPLAs: The system address map is used in conjunction with the FPLA coding forms that follow. These two sets of documents provide the source information for direct transfer into the customer pattern selected for the Programmable Gate Array (PGA) of the 2656. (Note: The specifics of programming the punch card deck for all programmable functions of the 2656 are covered in the Signetics 2656 data sheet. See CUS-TOM PATTERN PROGRAMMING IN-STRUCTIONS)

FPLA Coding Forms: (Tables 2-6) Four coding forms are provided respectively for FPLAs #1, #2, #3, and #4. These forms show the minimum coding required to make the PC-4000 functional. When additional coding from the systems address map is inserted on the FPLA coding forms, the PGA of the PC-4000 is defined.

MNEMONIC	FPLA NO.	TYPE	FUNCTION					
A0 to A14		1	Address Input					
MEM		1	Memory or 1/O					
OP		1	Operation Request					
WRPB	_	1	Write Pulse					
FS0 to FS7	_	1	Function Select for Pins X0 to X7					
NOROMN	_	1	Disable the ROM Chip enables					
NORAMN		1	Disable the RAM Chip enable					
NOPORTN	_	1	Disable the I/O Port					
1k2k	_	- 1	Sets the configuration of the ROM Memory to					
			2k (should be true for 2656 emulation)					
READ	_	. 1	Read Signal					
SELROM	_	1	Enable the Data Bus. A ROM operation has					
			been selected.					
PORT		1	The I/O Port operation has been selected.					
RAMCEN	_	1	The RAM Chip enable.					
PGA0 to PGA7	1,2	0	Chip Enable Decode for Pins X0 to X7					
ROMCE0N	3	0	Chip Enable for ROM 0					
ROMCE1N	3	0	Chip Enable for ROM 1					
ROMCE2N	3	0	Chip Enable for ROM 2					
ROMCE3N	3	0	Chip Enable for ROM 3					
SELROM	3	0	ROM Operation has been selected					
PORT	3,4	0	An I/O Operation has been selected					
RAMCEN	3	0	Chip Enable for the RAM					
RE	4	0	Enable the Data Bus Receiver					
DE	4	0	Enable the Data Bus Driver					
PWCN	4	0	Port Write Clock					
REDPORTN	4	0	Enable the Port Read Driver to the Internal					
			Data Bus.					
REDRAMN	4	0	Turn on the RAM output Driver					

NOTE

Signals ending in . . . N are active low

NORAM, NOROM, NOPORT disable the control signals for these functions. The Switch is on to disable these functions.

Table 1 PGA SIGNAL DESIGNATION

If the RAM, ROM and Port addresses are not programmed in the FPLAs (don't care all terms) the output signals will always be valid. The switch options along with the minimum programming of the FPLA prevent these signals from causing bus conflicts.

The symbols for the entries on the FPLA coding forms 1 through 4 (Tables 2-6) are:

> Input Variable H = Input High

L = Input Low

- = Don't Care

Blank = User Option

Output Function A - Active

* - Not Present in Term

Active Level

H - Active High L - Active Low

The tables that follow all use the standard coding format used for Signetics FPLAs. The format provides a matrix expression for logical expressions of the forms:

Output Signals =
$$(I_1 \bullet I_2 \bullet I_3 - - I_n)_1 + (I_1 \bullet I_2 \bullet I_3 \bullet - - - I_n)_2 +$$

(11 • 12 • 13--1n)k

n = 18 (for the 2656) k = 1-4 depending upon

the particular use

Ii = an input signal or its complement

Example:

Where:

The output function:

 $\overline{\text{Output}} = A \bullet B \bullet D + D \bullet \overline{A} + C \bullet \overline{B} \bullet E$

would be shown in the table as:

	[INPUT VARIABLE											UTP	UT F	UNC	TION	V
Sign Nam or No.	^{1e} / ⊲		ر م	/_) / u												
	PRODUCT TERM											AC	TIVIT	E LE	VEL		-
				nob	001	IEN	IVI					L				T	1
											7	6	5				1
00	-	-	-	-	-							*					1
01	Н	Н	-	Н	-							A			1		1
02	Ļ	-	-	Н	-							Α				1	1
03	-	L	Н	-	Н			1				Α					1
04	Н	Η	Н	-	-							*					1

Asterisks indicate that these terms are not present in this output.

The minimum necessary fixed programming for a 2650 application is shown below:

DEVICE	MEMORY ADDRESS	I/O PORT ADDRESS	WRP	ОР	M/IO	A14	A13	A12	A11	A10	A09	A08	A07	A06	A05	A04	A03	A02	A01	A00
X0					ļ															
X1	- !			-																
X2																				
Х3										, i			-							
X4																				
X5																				
X6																				
X7																				
I/O Port																				
RAM					1								-	Х	Х	Х	Х	Х	Х	Х
PROM					1					X	Х	Х	Х	Х	Х	Х	Х	Х	Χ	X

NOTE

Memory address and I/O Port address columns remain blank, unless X-pins are chip enables. The memory I/O columns for memory are programmed as "1" and for I/O are programmed as "0." For extended I/O operation A13 = 1, A12 thru A8 and A14 are

programmed Don't Care and A7 thru A0 as I/O address. For non-extended I/O operation, there are two modes, data and control. For data A13 = 0, A14 = 1, A12 thru A0 are Don't Care. For control, A13 = 0, A14 = 0, A12 thru A0 are Don't Care.

Table 2 SYSTEM ADDRESS MAP FOR CODING FPLA

	Γ						1	NPU	T VA	RIAE	BLE			-				Γ		ΟU	ITPU	T FU	INCT	ION		
Signa Nam or No.	- 1	\int	\int	FS3	FS2	FS1	FS0	48	47	46	45	44	43	42	41	40	//	<u> </u>	CA6	CAS	C44	CA3	C.R.	CA1	CAO	
																	-			ACT	ΓIVE	LEVI	EL			l
	PRODUCT TERM														1	Н	Н	Н	Н	Н	Н	Н				
	15	14	13	12	11	10	09	08	07	06	05	04	03	02	01	00		7	6	5	4	3	2	1	0	ĺ
00																		•	*	*	*	*	*	*	Α	
01																		• [*	*	*	*	*	Α	*	
02																		*	*	*	*	*	Α	*	*	
03																		*	*	*	*	Α	*	*	*	
04																		*	*	*	Α	*	*	*	*	
05						1												*	*	Α	*	*	*	*	*	
06																		*	Α	*	*	*	*	*	*	
07						Ī	Ì											4	*	*	*	*	*	*	*	

Table 3 CODING TABLE FOR FPLA NO. 1

	[···			ı	NPU	T VA	RIAE	BLE						\Box		ΟL	JTPU	T FU	NCT	ION		
Signa Nam or No.				F\$7	FS6	FSS	FS4	WRP	ලී	MEM	414	413	472	411	A10	409	%	CA6	CAS	CA4	CA3	CA2	54.7	CAO	
	PRODUCT TERM															AC	TIVE	LEVI	EL			ĺ			
	PRODUCT TERM													Н	Н	Н	Н	Н	Н	Н	H	ı			
				12	11	10	09	08	07	06	05	04	03	02	01	00	7	6	5	4	3	2	1	0	
00																	*	*	*	*	*	*	*	Α	ı
01																	*	*	*	*	*	*	Α	*	
02		1															*	*		*	*	Α	*	*	
03																	*	*	*	*	Α	*	*	*	
04							Н										*	*	*	Α	*	*	*	*	
05		Ī				Н											*	*	Α	*	*	*	*	*	
06					Н												*	Α	*	*	*	*	*	*	
07				Н													Α	*	*	*	*	*	*	*	

Table 4 CODING TABLE FOR FPLA NO. 2

	[i	NPU	T VA	RIA	BLE					······································		Γ		OL	JTPU	T FL	JNCT	TION		
Signa Nam or No.	,	READ	NORAL	NORON	WRP	ළි	MEM	A14	473	472	471	470	409	408	407	RAMCEA		MAMCEN		PORT	SELROLL	ROMOF	ROMOFE	ROMCE	ROMCFOLL	Also /
	PRODUCT TERM																		AC	TIVE	LEV	EL			1	
							PAC	יטענ	<i>-</i>	EK IVI							1	П	Н	Н	Н	L	L	L	L	
	15	14	13	12	11	10	09	08	07	06	05	04	03	02	01	00	-	7	6	5	4	3	2	1	0	
00	-	Н	-	Н	Н	H	Н					L	L	-	-	Н	Γ,		*	*	Α	*.	. *	*	Α	
01	-	Н	-	Н	Н	Н	Н					L	Н	-	-	Н		ľ	*	*	Α	*	*	Α	*	
02	Н	Ι	-	Н	Н	Н	Н					Н	L	-	-	Н			*	*	Α	*	Α	*	*	
03	Н	Ι	-	Н	Ι	Н	Н					Н	Н	-	-	Н			*	*	Α	Α	*	*	*	
04	-	ı	Н	-	Η	Н	Н									-]	1	T	*	*	*	*	*	*	*	
05	-	1	-	-				-								-	F	-	*	Α	*	*	*	*	*	
																										i.

Table 5 CODING TABLE FOR FPLA NO. 3

							. 1	NPU	T VA	RIAE	BLE		•					1 [IT FL	JNCT	TION		
Signa Nam or No.					NOPOR	RAMOE	PORT	SELROIL	READ	A6	45	44	43	42	A. /	40				REDPORT	PWCN	REDRAL	DE	AR.	PORT	
	PRODUCT TERM ACTIVE LEVEL																									
							PHU	יטטכ	<i>)</i>	EKM										L	L	L	Н	L	Н	
					11	10	09	08	07	06	05	04	03	02	01	00		7.	6	5	4	3	2	1	0	
00					Н	-	-	-	-								ſ	*	*	*	*	*	*	*	Α	WRITE RAM
01					-	L	L	L	L	-	-	-	-	-	-	-	Γ	*	*	*	*	*	*	Α	. *	READ ROM
02					-	Н	L	Н	Н	-	-	-	-	-	-	-		*	*	*	*	*	Α	*	*	READ PORT
03					-	Н	Н	L	Н	-	-	-	-	-	-	-		*	*	Α	*	*	Α	*	*	WRITE PORT
04					-	Н	Н	L	L	-	-	-	-	1	1		. [* .	*	*	Α	*	*	Α	*	READ RAM
05					-	L	L	L	Н	-	-	-	-	-	-	-		*	*	*	*	Α	Α	*	*	
											L					LJ	-1	4								

Table 6 CODING TABLE FOR FPLA NO. 4

CODING SWITCHES AND JUMPERS: The switches and jumpers are set by the user for the particular application of the 2656 in his system. Purposes of the switches and jumpers are described in the functional block descriptions. Specifics of coding are given in 3 groups.

Tables 7, 8 and 9 show the switch positions and jumper connections for:

- 1. Memory and clock divider functions.
- 2. Select port or chip enable.
- 3. Jumpers

These tables, when properly filled-in, may be used as source data for the codes to be transferred into the 2656 coding forms.

Oscillator Frequency: output clock frequency equals oscillator frequency:

- ÷ 1 set S2-5 Off, S2-6 Off
- ÷ 2 set S2-5 On, S2-6 Off
- ÷ 3 set S2-5 Off, S2-6 On
- ÷ 4 set S2-5 On, S2-6 On

SWITCH	ON	OFF	FUNCTION
S2-2			NO ROM
S2-3		l	NO ROM
S2-1			NO I/O PORT
S2-4		×	1k/2k
S2-5			DIV0
S2-6		}	DIV1

Table 7 MEMORY AND CLOCK DIVIDER

SWITCH NO.	I/O PORT SWITCH ON	CHIP ENABLE SWITCH OFF	X PIN
S1-1			X0
S1-2			X1
S1-3			X2
S1-4			Х3
S1-5			X4
S1-6			X5
S1-7			X6
S1-8			X7

Table 8 PORT OR CHIP ENABLE

	JUMPER	COMPONENT
XTAL OSC	W2 to W8	XTAL FROM W6 to W7
RC OSC	W1 to W2	CAPACITOR FROM W6 to W7
EXT DRIVEN	W2 to W3	VV6 to VV7
OUTPUT	W4 to W5	

Table 9 JUMPERS AND ADDITIONAL COMPONENTS

40-Conductor Cable Pin-Out: The Interface from the PC-4000 to the users system is a 40-pin cable with 40-pin dip plugs on both ends.

NO.	MNEMONIC	STATE	OPERATION
SWITCHES			
S1-2	FS0	ON OFF	Select X0 As I/O Port Bit Select X0 As Chip Select
S1-3	FS1	ON OFF	Select X1 As I/O Port Bit Select X1 As Chip Select
S1-1	FS2	ON OFF	Select X2 As I/O Port Bit Select X2 As Chip Select
S1-4	FS3	ON OFF	Select X3 As I/O Port Bit Select X3 As Chip Select
S1-5	FS4	ON OFF	Select X4 As I/O Port Bit Select X4 As Chip Select
S1-6	FS5	ON OFF	Select X5 As I/O Port Bit Select X5 As Chip Select
S1-7	FS6	ON OFF	Select X6 As I/O Port Bit Select X6 As Chip Select
S1-8	FS7	ON OFF	Select X7 As I/O Port Bit Select X7 As Chip Select
S2-1	NOROM	ON OFF	ROM Chip Selects Disabled ROM Chip Selects Enable
S2-2	NORAM	ON OFF	RAM Chip Selects Disabled RAM Chip Selects Enable
S2-3	NOPORT	ON OFF	Port Select Disabled Port Select Enabled
S2-4	1k2k	ON OFF	ROMCE2N AND ROMCE3N Disabled ROMCE2N AND ROMCE3N Enabled
S2-5 S2-6	CLKDIV0 CLKDIV1		
JUMPERS			•
W2 to W8 W1 to W2 W2 to W3 W4 to W5			Use Internal Xtal Oscillator Use Internal RC Oscillator CK2 Pin is External Clock Input CK1 is Reset Output

Table 10 SWITCH AND JUMPER SUMMARY

MNEMONIC	NO.	MNEMONIC	NO.
DBUS3	1	DBUS2	40
DBUS4	2	DBUS1	39
DBUS5	3	DBUS0	38
DBUS6	4	X3	37
DBUS7	5	X2	36
X7	6	X1	35
X6	7	X0	34
X5	8	NC ²	33
X4	9	ADDR14	32
CLK OUT	10	ADDR13	31
CLK 1	11	ADDR12	30
CLK 2	12	ADDR11	29
VSS GND	13	ADDR10	28
R/W	14	ADDR9	27
M/I/O	15	ADDR8	26
OPREQ	16	ADDR7	25
WRP	17	ADDR6	24
ADDR0	18	ADDR5	23
ADDR1	19	ADDR4	22
ADDR2	20	ADDR3	21
L .		1	

Table 11 SMI CABLE PINOUTS¹

NOTES

- Timing element pins have alternate functions Clock 1 is external reset out. To use, connect W (4) to W (5). Clock 2 is external Clock in. To use, connect W (2) to W (3). Clock 2 not used: to use RC oscillator connect W (1) to W (2). To use internal XTAL oscillator connect W (2) to W (8).
- 2. V_{DD} for chip, no connect for board.

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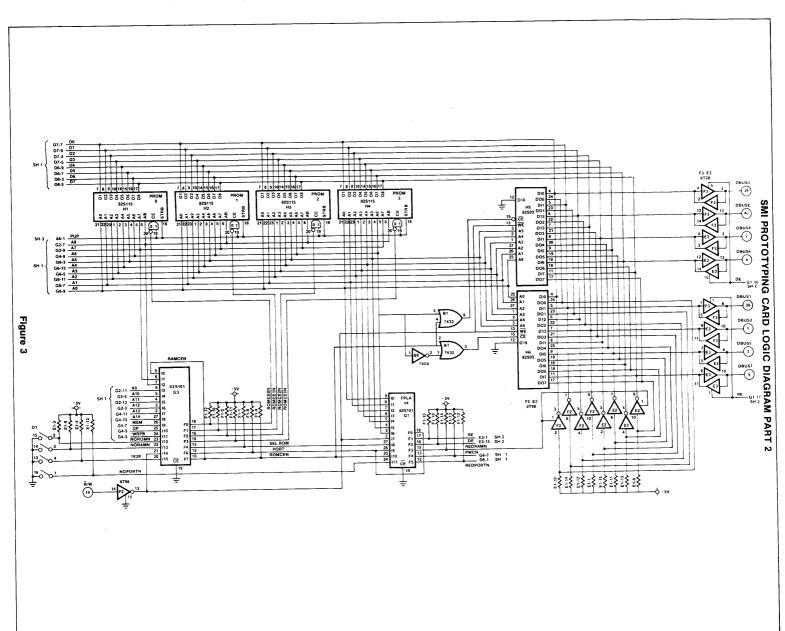
Figure 2

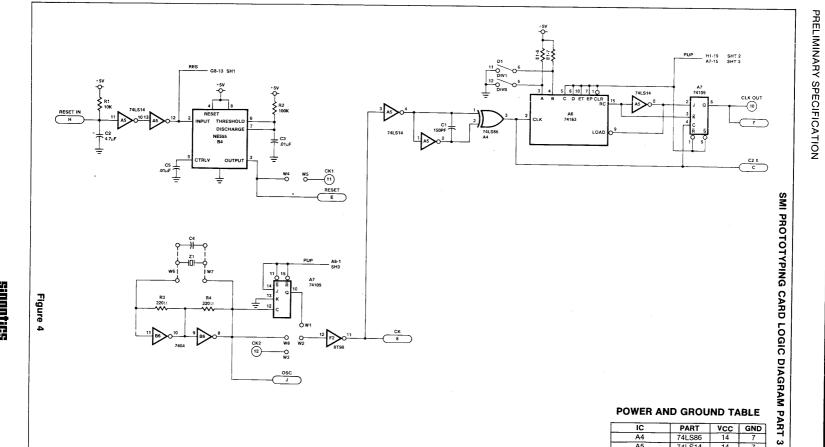
FS7N

G7 825101

G6-11 G6-5 G6-13 G6-3 G4-9 G2-9 G2-7 E4-9 E4-5 D4-1 D4-9

G2-11 A3 9 1
G2-23 A10 8 1
G2-13 A12 6 1
G2-3 A12 6 1
G4-11 A14 4 1
G4-7 OD 2 1
G4-7 OD 2 1
G4-3 WRPB 27 8
G4-5 WRPB 27 8
G4-5 S5 25 1
G4-1 E5-14 F5S 23 11
E2-12 F5S 23 112





POWER AND GROUND TABLE

- ao	,	\DLL
PART	VCC	GND
74LS86	14	7
74LS14	14	7
74163	16	8
74109	16	8
NE555	8	1
7404	14	7
7432	14	7
7405	14	7
8T97	16	8
8T98	16	8
8T28	16	8
7403	14	7
74126	14	7
825101	28	14
74116	24	12
825115	24	12
82509	28	14
RES PACK	16	
	74LS86 74LS14 74163 74109 NE555 7404 7432 8T97 8T98 8T28 87403 74126 825101 74116 82509	74LS86 14 74LS14 14 74LS14 16 74109 16 NE555 8 7404 14 7405 14 8T97 16 8T98 16 8T98 16 8T98 16 8T08 14 8403 14 74126 14 825101 28 74116 24 82509 28

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Figure 4

PRELIMINARY SPECIFICATION

EDGE CONNECTOR

The edge connector for the 2650/PC-4000 is an AMP 225-21021-4-01-117 edge connector. This edge connector has 20 pins on 0.156 inch centers.

EDGE CONNECTOR PIN-OUT				
GND	1	А	GND	
GND	2	В	GND	
	3	C	C2X	
	4	D		
	5	E	RESET OUT*	
	6	F	CLK OUT*	
	7	н	RESET IN'	
СК	8	J	OSC	
Vcc	9	K	Vcc	
Vcc	10	L	Vcc	
*Factory	*Factory Test Only Do not Use			
Figure 5				

SHIPPING CONFIGURATION

Included with the PC-4000:
Blank 82S115 PROMS 4 Each.
Blank 82S101 FPLAS 4 Each.
Edge Connector
2-foot cable with 40 pin headers on each end.
PC-4000 Data Sheet

REQUIRED BUT NOT SUPPLIED

Timing element for oscillator

PART NUMBER	QUANTITY	DESCRIPTION		
7403	2	Open Collector Quad Nand		
7404	3	Hex Inverter		
7405	2	Open Collector Hex Inverter		
74LS14	1	Hex Schmitt Trigger Inverter		
7432	1	Quad 2 Input or Gate		
74LS86	1	Quad 2 Input XOR		
74109	1	Dual JK Flop		
74116	1	Dual Quad D Latch with Clear		
74126	4	Quad Tristate Buffer		
74163	. 1	4 Bit Binary Counter		
8T28B	2	Bidirectional Data Bus Driver Receiver		
8T97B	5	Tristate Hex Buffer		
8T98B	2	Tristate Hex Inverter Buffer		
82S101	4	Open Collector FPLA		
82S115	4	512X8 PROM		
82S09	2	64X9 BIPOLAR RAM		
NE555	1	Timer		
761-1-51.0KOHM	4	Resistor Dip Pak		
	1	10k OHM Resistor		
	1	100k OHM Resistor		
	2	220 OHM Resistor		
İ		Note: All Resistors 1/4 Watt		
	1	150pF Capacitor		
	2	0.01μF Capacitor		
	5	4.7μF Capacitor		
	14	0.1μF Capacitor		
XTAL or Capacitor	1	Timing Element for Oscillator (Not Supplied)		
	1	Edge Connector AMP 225-21021-401-117		
	2	8 Position Dip Switch		
	1	PC Board 2650/PC-4000		
	4	24 Pin Dip Socket		
	4	28 Pin Dip Socket		
	1	40 Pin Dip Socket		
	1	Cable Assembly		
	ļ			

Table 12 PARTS LIST

DESCRIPTION

The 2650PC1001 is a complete microcomputer on a single printed circuit board. The heart of this computer is Signetics' 2650 Microprocessor; a single chip, N-Channel MOS Integrated Circuit which contains the CPU and control sections of the classical general purpose computer architecture.

In addition to the Microprocessor, the 2650PC1001 contains both control and read/write memory, i/O ports, clock, and all the necessary buffering and interface circuits to permit data transfer both on and off the p.c.b.

FEATURES

- 2650 Microprocessor
- 1K bytes of ROM with PIPBUG*
- 1K bytes of RAM (off-board expandable)
- 1MHz crystal oscillator
- Serial I/O (either TTY 20mA current loop or RS232—selectable by jumper wire)
- Two 8-bit output ports
- Two 8-bit input ports
- DMA capability
- LED display indicators
- Data bus and address bus test points
- Buffered data and address outputs
- Single power supply (+5 volts)**

MEMORY

The memory of the 2650PC1001 is divided into two segments:

- a. ROM with PIPBUG
- b. RAM (Read/Write Memory)

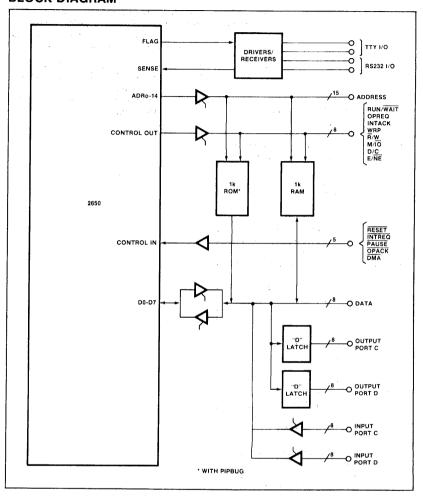
The Read-Only Memory (ROM) supplied with the card is the Signetics' 82S129 Field Programmable type (PROM). Eight of these 256X4 devices are arranged to provide a 1KX8 memory array. The 2650PC1001 is supplied with the PIPBUG loader and debugger already programmed into the ROM. Since the devices are loaded into sockets, however, they can be easily replaced with other ROMs or PROMs programmed by the user.

The 1KX8 array is constructed with Signetics' 2606 NMOS RAM devices. Since the 2606 is a 256X4 device, again 8 devices are used in the array.

SERIAL I/O

The serial I/O capability of the 2650PC1001 utilizes a unique serial I/O feature of the basic 2650 microprocessor. This feature

BLOCK DIAGRAM



allows serial data to be transferred directly into the 2650 under program control by using the sense and flag pins on the microprocessor.

Two types of serial I/O ports are available. The first is a teletype interface which can be directly connected to a teletype 20mA current loop. The second is an RS232 interface which provides a connection for voltage driven peripheral equipment. The selection of the particular interface to be used is made by connecting a jumper wire directly from the microprocessor flag and sense lines to the appropriate output port. If the RS232 interface is used, +12 and -12 volt supplies are required in addition to the +5 volt supply which operates the rest of the board.

PARALLEL I/O

Parallel I/O channels using the 2650's unique Non-Extended I/O mode are also provided. This mode allows a single byte instruction to select one of two distinct I/O devices. On the 2650PC1001, these two devices are represented by four separate data channels; two for reading and two for writing. The output (or write) channels are fully latched and buffered. The input (or read) channels are fully buffered. One read and one write channel represents a single I/O device. In addition to the Non-Extended I/O ports, the data and address buses, plus the appropriate control signals, are also available to provide the full extended I/O capability.

^{*}Signetics Loader and Debugging Program. (See appl. note SS50)

^{**}Assumes RS232 I/O port is not used.

OTHER I/O

A complete listing of the I/O pins, plus a brief description of any I/O signal not detailed above, is as follows:

PIN	DESCRIPTION
1,2	Ground
4-11	Processor Data Bus
12	Strobe to Enable Input Data Port
13	D/C Output
14	DMA Control Input
15	Extended/Non-Extended Output
16	Interrupt Acknowledge Output
17	R/W Output
18	Write Pulse Output
19	Run/Wait Output
20	Operation Request Output
21	Memory/IO Output
22	Operation Acknowledge Input
	Clock Output (or Input if on-board
23	clock output of input if on-board
0.4	Operation Request Input for DMA
24	
25	Reset Input
26	Interrupt Request Input
27	Pause Input
28-32	Unused
33-47	Address Bus
48	+12 Volts for RS232
49	-12 Volts for RS232
50	+5 Volts
A,B	Ground
C	Not used
D-M	Non-Extended Output Port "D"
N	Clock to load data into Output
1_	Port "D"
P	TTY serial data input (+)
R	TTY serial data input (-)
S	TTY serial data Output pull up
1_	resistor (current loop +)
T	TTY serial data Output; TTL Level,
	open collector (current loop re-
1	turn)
U	RS232 ground
V	RS232 Output
W	TTY tape reader Output; TTL Lev-
1.,	el, open collector (+)
X	TTY tape reader Output pull up
	resistor (-)
Y	RS232 Input
Z	Clock to load data into Output
1	Port "C"
a-h	Non-Extended Output Port "C"
j	Strobe to enable input Port Con-
1.	trol
k-u	Non-Extended Input Port "D"
v-c	Non-Extended Input Port "C"
d	+12 for RS232
е	-12 for RS232
f	+5 Volts

NOTE

Italic items indicate buffered 2650 Microprocessor Outputs.

SUMMARY

The above is intended to provide a brief description of Signetics' 2650PC1001 Prototyping Board. More detailed information can be obtained from the following:

M20	PIPBUG Application Note	
M14	2650PC1001 Manual (Detailed	d
	Description)	
M1	Serial I/O using Sense and Flag	g
	Application Note	

DESCRIPTION

The 2650 PC2000 is a 4K Memory Card designed to be compatible with the 2650 microprocessor. It is composed of 32, 21L02 NMOS, 1K by 1 bit static RAMs, and organized in four groups of one kilo-byte each. Decoding is provided to select one of the four groups and also distinguish the card in multi-card configurations. In a system application utilizing up to 8 cards (32K), each card is uniquely identified by hardwired jumpers. No external decoding is required.

The decoding logic is sectioned into two blocks. The first block determines if the address identifies that card as being part of the 8K page address. (The 2650 memory scheme is organized into 4 pages of 8K each.) The second block uniquely locates 1K bytes of memory on the board in the 8K bytes of memory of the selected page. Each 1K bank is individually selected by hardwired jumpers to the decoder.

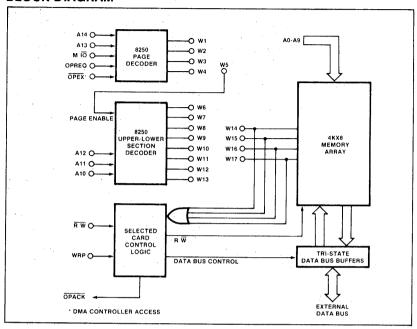
FEATURES

- Requires only single +5V supply
- Industry standard 21L02 memories
- Fully decoded for 32K memory organization
- Data bus buffered with tri-state drivers/receivers
- Accessible from microprocessor or DMA controller
- TTL compatible
- Dimensions are 8" X 6.875" with a 50 pin edge connector along the 8" dimension
- Typical power consumption of 4.5 watts

SIGNAL DEFINITION

Memory control signals and address lines between the 2650 microprocessor and the 2650 PC2000 are indicated in the block diagram. The OPEX control line is reserved for use with DMA controllers. Its function is similar to that of the OPREQ line from the 2650. When either of these lines are true and a memory operation is specified $(M/\overline{IO} =$ High) the memory card is enabled to decode address lines A0 through A14. When a bank is selected, the selected card control logic block allows the read-write line (\overline{R}/W) and write pulse (WRP) to pass to the memory array and also enable the external data bus drivers. When the operation is complete the memory card responds with a true condition on OPACK.

BLOCK DIAGRAM



JUMPER ADDRESS DECODING

Jumpers are applied to designated platedthrough holes identified by a 'Wn' mnemonic. To identify the card to be part of a particular page, jumper point W5 to one of the following:

W1 for page 0 W2 for page 1 W3 for page 2 W4 for page 4

To locate each of the 1K bytes of the memory card in the selected memory page, four bank jumpers are required. The outputs of the decoder used to select one of eight 1K byte memory segments (W6-W13) must be connected to the selected 1K bytes of memory on the 2650 PC2000 (W14-W17).

Factory installed jumpers allow for immediate hook-up to a Demo System (DS1000/2000) which has 2K of memory. These jumpers have been hooked-up as follows:

W1 to W5 (page 0) W8 to W14 W9 to W15 W10 to W16 W11 to W17

PIN	NAME	PIN	NAME
1,2,A,B 4 5 6 7 8 9 10 11 17 18 20 21 22 24 33	GROUND DBUS0 DBUS1 DBUS2 DBUS3 DBUS4 DBUS5 DBUS6 DBUS7 R/W WRP OPREQ M/IO OPACK OPEX ABUS11	34 35 36 37 38 39 40 41 42 43 44 45 46 47 50,f	ABUS13 ABUS12 ABUS14 ABUS9 ABUS10 ABUS8 ABUS7 ABUS6 ABUS5 ABUS3 ABUS0 ABUS1 ABUS4 ABUS2 VCC +5V

Table 1 EDGE CONNECTOR TABLE

DESCRIPTION

The Demo System 2000 (2650 DS2000) is a hardware base for use with the 2650 CPU printed circuit board (PC1001) and allows the exercising of this card with user defined options. When the DS2000 is combined with a CPU board (PC1001) and a TTY, the user is equipped with everything he needs to exercise any of the software or hardware features of the 2650. The DS2000 has a built in power supply.

FEATURES

- User defined expansion capability from connector supplying address, data and control lines.
- RS232 and TTY interface
- Two extended and two non-extended I/O ports
- Single step capability for program debuaaina
- Display of address bus, data bus and the two non-extended I/O ports

CONNECTORS

The 2650 CPU Board (PC1001) is inserted into the J8 connector to complete the demo system. The user printed circuit board is inserted into the J7 connector. Both connectors are the same type (100 Pin Amphenol, series 225) and the numbered pins of J7 and J8 have the same signals (except pin 12). The lettered pins of J7 (pins A through \overline{g}) are not used. The sockets and connectors of the DS2000 and their associated signals are provided in this data sheet.

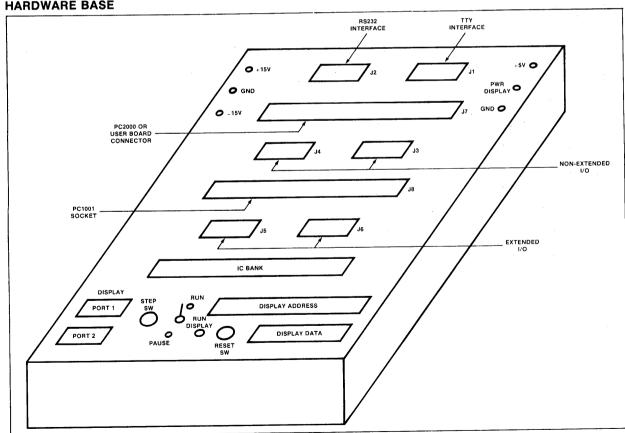
DISPLAYS

The address and data bus LED displays reeach OPREQ (beginning of an external operation). Latches store the information until another OPREQ is received. The two nonextended port displays represent data on channel C (port 2) and channel D (port 1) during the OPREQ for each I/O operation. A logic one one these displays will turn "on" the LED and a logic zero will turn them "off."

CONTROLS

The pause and step logic allows one instruction to be executed at a time by pushing the 'step' button when the Run/Pause switch is in the pause position. In this mode the Run/Wait display LED will go off. The reset switch will reset the display latches and place all zeros in the 2650 instruction address register.

HARDWARE BASE



SIGNAL NAMES FOR CONNECTORS 17 AND 18

PIN NO.	FUNCTION (J7 & J8)	PIN NO.	FUNCTION (J8 ONLY)1
1	GND	A	GND
2	GND	В	GND
3	NC2	C	NC NC
4	DBUS0	ا م	OPD0
5	DBUS1	E	OPD1
6	DBUS2	F	OPD2
7	DBUS3	Н	OPD3
8	DBUS4	J	OPD4
9	DBUS5	K	OPD5
10	DBUS6	l L	OPD6
11	DBUS7	М	OPD7
121	EIPD	N	COPD
13	D/C	P	TTY SERIAL IN +
. 14	DMA	R	TTY SERIAL IN -
15	E/NE	s	TTY SERIAL OUT +
16	INTACK	T	TTY SERIAL OUT -
17	R/W	Ù	RS232 GROUND
18	WRP	V	RS232 OUTPUT
19	RUN/WAIT	w	TTY TAPE READER OUT
20	OPREQ	x x	TTY TAPE READER OUT -
21	M/IO	Î	RS232 INPUT
22	OPACK	ż	COPC
23	CLOCK	a	OPC0
24	OPEX	b	OPC1
25	RESET	С	OPC2
26	INTREQ	d	OPC3
27	PAUSE	e	OPC4
28	NC	f	OPC5
29	NC	g	OPC6
30	NC	h	OPC7
31	NC	j	EIPC
32	NC	k	IPD0
33	ABUS11	m	IPD1
34	ABUS13	n	IPD2
35	ABUS12	р	IPD3
36	ABUS14	r	IPD4
37	ABUS9	s	IPD5
38	ABUS10	t	IPD6
39	ABUS8	u	IPD7
40	ABUS7	r	IPC0
41	ABUS6	w	IPC1
42	ABUS5	×	IPC2
43	ABUS3	у	IPC3
44	ABUS0	z	IPC4
45	ABUS1	а	IPC5
46	ABUS4	b	IPC6
47	ABUS2	С	IPC7
48	+12V	đ	+12V
49	-12V	e	-12V
50	+5V	g l	+5V

^{1.} J7 has no connections to these pins.

^{2.} NC = no connection.

EXTENDED INPUT/OUTPUT DIP SOCKETS

NON-EXTENDED INPUT/OUTPUT DIP SOCKETS

PIN NO.	FUNCTION J5	FUNCTION J6
1	DBUS0	ABUS0
2	DBUS1	ABUS1
3	DBUS2	ABUS2
4	DBUS3	ABUS3
5	DBUS4	ABUS4
6	DBUS5	ABUS5
7	DBUS6	ABUS6
8	DBUS7	ABUS7
9	OPACK	ABUS8
10	M/ĪŌ	ABUS9
11	OPREQ	ABUS10
12	RUN/WAIT	ABUS11
13	WRP	ABUS12
14	R/W	ABUS13
15	INTACK	ABUS14
16	E/NE	PAUSE
17	DMA	INTREQ
18	D/C	CLOCK

PIN NO.	FUNCTION J3		FUNCTION J4	
1	(Output Port C)	0	(Output Port D)	0
2	OPC	1	OPD	1
3	OPC	2	OPD	2
4	OPC	3	OPD	3
5	OPC	4	OPD	4
6	OPC	5	OPD	5
7	OPC	6	OPD	6
8	OPC	7	OPD	7
9	Clock Output Port	С	Clock Output Port	D
10	Enable Input Port	С	Enable Input Port	D
11	(Input Port C)	7 -	(Input Port D)	7
12	. IPC	6	IPC	6
13	IPC	5	IPC	5
14	IPC	4	IPC	4
15	IPC	3	IPC	3
16	IPC	2	IPC	2
17	IPC	1	IPC	1
18	IPC	0	IPC	0

DESCRIPTION

The 2650 PC3000 is a basic text generating system requiring only six integrated circuits including one 2650 microprocessor. The serial communication link between the 2650 and the users terminal is accomplished with the flag and sense lines on the microprocessor. The 2650 PC3000 is used to control the storage of characters entered from a terminal with either a current loop or voltage swing capability (±7.5V min).

Control Characters allow the text to be printed out on the terminal with the capabilitv for inserting unique characters at locations identified during text generation. When the text is printed out the entire text will be output unless a control character is detected. The microprocessor then stops the print-out and the operator enters the desired unique information. Another control character is then given to continue printing the text until all characters stored in memory are printed, or until another stop character is detected. The stop character is recorded in memory just like any other character; however, it is not printed during text print-out.

Additional control characters allow for the erasure of the previous character typed or the erasure of the entire memory.

FEATURES

- Total of six IC packages
- Operates at +5V at a max of 500mA
- Interface to either current loop or device capable of sending and receiving a minimum voltage swing of \pm 7.5 volts referenced to signal ground
- 250 character storage capability
- Card size less than 3" X 4" with four screwed-on stand-offs at corners
- 1MHz clock implemented with 74123 oneshot
- Variable baud rate between 110 and 300 baud by trimmer pot adjustment of clock
- PROM mounted in 24 pin socket
- Card edge connector supplied with each card
- Inputs provided for an external system reset

TERMINAL INTERFACE Voltage Mode Terminal Connection

The voltage mode interface is very similar to the standard RS232 interface except that the "signal" ground cannot be connected to "protective" ground. When a Cinch type 25pin connector (DB25P or DB25S) is used on an RS232 compatible terminal, the PC3000 should be connected as shown in Table 1.

Current Loop Terminal Connection

When a terminal is used that employs current loop transmission techniques the four wires from the terminal should be connected to the corresponding four pins on the PC3000 card: TTY OUT +, TTY OUT -, TTY IN +, and TTY IN -.

On card jumper point 'A' to 'B' and point 'D' to 'F.'

PC3000 COMMAND SUMMARY

KEY	FUNCTION
	Erase last character in memory and echo the erased character. Additional preceding characters can be erased by continuing to depress the delete key.
Control and E Control and B	Erase entire memory. Used to indicate beginning of inserted message. Is not printed but stored in memory. Stops print-out when read from memory. Required once from each unique information entry.
Control and C	Continues print-out of memory after entry of unique information.
Control and P	Prints out contents of ter- minal memory.
Control and R	Software reset.

Bell will ring if any of the following are true.

- 1. Entering more than 250 characters in memory
- Requesting print-out of an empty buffer.
- Attempting to delete more characters than there are in memory

DB25P (DB25S)	PC3000 EDGE	PC3000
PIN NO.	CONNECTOR PIN NO.	SIGNAL NAME
1 3	No connection 6	VS OUT +
2	J	VS IN +
7	K	VS OUT—(Signal GND)
5,6,8,20	Connect Together	

On card jumper point 'A' to 'C' and point 'D' to 'E.'

Table 1 VOLTAGE MODE TERMINAL PIN DESCRIPTION

PART DESCRIPTIONS

PART	FUNCTION
2650	8-bit TTL compatible N-
	Channel Microprocessor in-
	corporating a serial I/O Port.
1	(See 2650 Hardware Specifica-
1	tion Manual for complete
2606	description—2650 BM1000.)
2000	1024-bit static MOS, TTL compatible RAM memory organ-
1	ized as 256 words by 4
	bits/word.
82S115	4096-bit Bipolar TTL compati-
020110	ble PROM organized as 512
	words by 8 bits/word.
N7426	Quad 2-input high voltage
	NAND gate with open collec-
	tor capable of driving voltage and current loop interfaces
	(20mA maximum).
74123	Dual retriggerable monostable
	multivibrator with clear con-
	figured as a clock for 2650.
Poten-	Hellpot series 91C, 50K, ohm
tiometer	3/8" cermet trimming poten-
	tiometer.
DO E .	
PC Edge	Ampnenol—225-21021-401-
Connector	Cinch—251-10-30-160

Miscellaneous components consist of 11 (1/4) watt and two (1/2) watt resistors, and two mica, one ceramic and one tantilum capacitor.

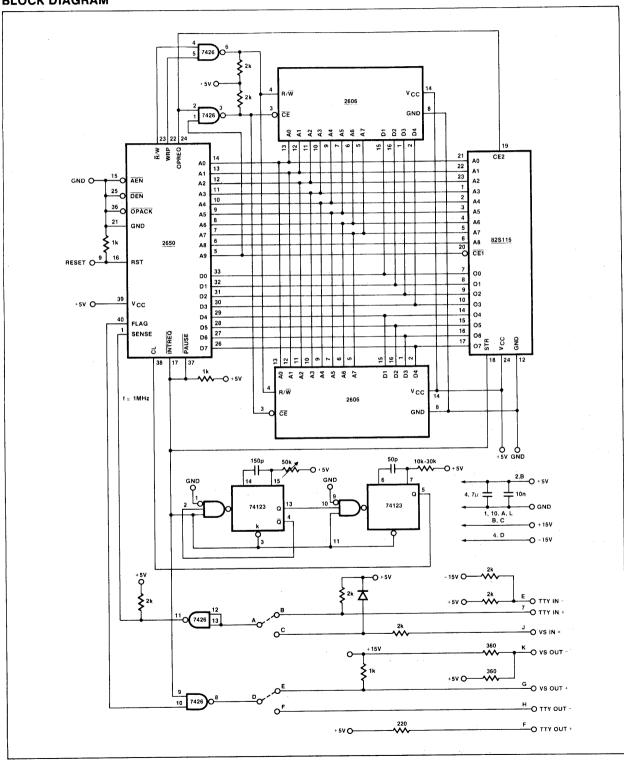
The following are required to make the board functional but are not supplied with the card:

RS232 type connector for voltage swing interface: DB25P or DB25S

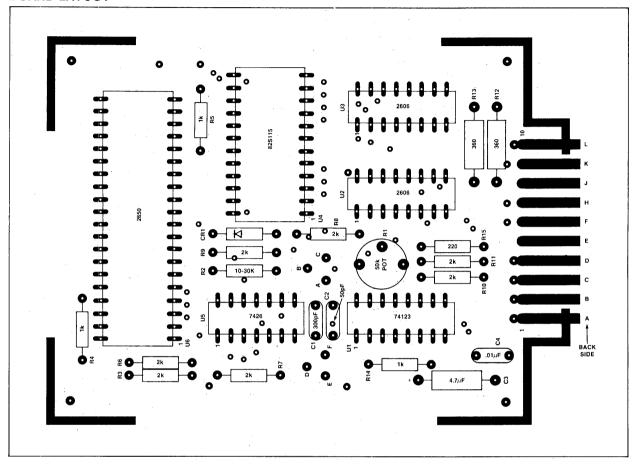
Reset switch-(normally open, connected to +5V)

Power supplies: +5V ±15V

BLOCK DIAGRAM



BOARD LAYOUT



PC3000 CONNECTOR PIN ASSIGNMENT

PIN	FUNCTION	PIN	FUNCTION
1	GND	A	GND
2	+5	В	+5
3	+15	l c	+15V
4	-15	D	-15
5	- .	. E	TTY IN -
6	VS OUT +	F	TTY OUT +
7	TTY IN +	H.	TTY OUT -
8	_	، ز	VS IN +
9	RESET	K	VS OUT - (Signal Ground)
10	GND	L	GND
VS-	Voltage Swing		

DESCRIPTION

The 2650 assembly language (PIPHASM) is a symbolic language designed specifically to facilitate the writing of programs for the Signetics 2650 microprocessor.

The AS1000 is configured to operate on 32bit or larger machines and the AS1100 is configured to operate on 16-bit machines.

The 2650 assembler is a program which accepts symbolic source code as input and produces a listing and/or an object module "Hexadecimal" format compatible to the two tape punching programs PIPHTAP (for acceptance by PIPBUG), PIPSTAP (for PROMs) and also to the simulator, PIPSIM.

The assembler is written in standard Fortran IV and is approximately 1.250 Fortran card images in length. It is modular and may be executed in an overlay mode should memory restrictions make that necessary. It operates in a two pass mode to build a symbol table, to issue helpful error messages, produce an easily readable program listing and output a computer readable object module. This version of the assembler compiles into a 12K word load module on the PDP-11/40 (16-bit words) and executes under DOS (8K) within a 28K memory.

AVAILABILITY

The 2650 assembler is available on both NCSS and GE timeshare. It is also available from Signetics on 9 track magnetic tape written in EBCDIC in 80 character unblocked records at a density of 800 bpi.

FEATURES

- Forward references
- Pseudo-Ops to aid programming
- Self-defining constants
- Symbolic machine operation codes
- Free format source code
- Syntax error checking
- Symbolic address assignment and references
- **Data creation statements**
- Storage reservation statements
- Assembly listing control statements
- Addresses can be generated as constants
- Character codes may be specified as
- Comments and remarks may be encoded for documentation

LANGUAGE REQUIREMENTS

Input Requirements

ASCII or EBCDIC

Input to the assembler consists of a sequence of characters combined to form assembly language elements. These language elements include symbols, instruc-

tion mnemonics, constants and expressions which make up the individual program statements that comprise a source program.

CHARACTERS		
Alphabetic:	A through Z	
Numeric:	0 through 9	
Special		
Characters:		blank
	(left parenthesis
)	right parenthesis
	+	add or positive value
	-	subtract or negative
		value
	*	asterisk
	,	single quote
	,	comma
	/	slash
	\$	dollar sign
	<	less than sign
	>	greater than sign

SYMBOLS

Symbols are formed from combination or characters. Symbols provide a convenient means of identifying program elements so they can be referenced by other elements.

CONSTANTS

A constant is a self-defining language element. Unlike a symbol, the value of a constant is its own "face" value and is invariant. Internal numbers are represented in 2's complement notation. There are two forms in which constants may be written: the Self-Defining Constant and the General Con-

Self-Defining Constant

The self-defining constant is a form of constant which is written directly in an instruction and defines a decimal value.

General Constant

The general constant is also written directly in an instruction, but the interpretation of its value is dictated by a code character and delimited by quotation marks. Its form can be binary, octal, decimal, hexadecimal, EBCDIC or ASCII.

EXPRESSIONS

An expression is an assembly language element that represents a value. It consists of a single term or combination of terms separated by arithmetic operators. A term may be a valid symbolic reference, a selfdefining constant or a general constant.

Fields

A statement prepared for processing by the assembler is logically divided into four fields, as indicated below. They are free form and are separated by at least one blank character. The name must begin in logical column 1.

OPERATION OPERAND COMMENTS LARFI operand(s) name opcode

1416---

Where:	
FIELD	DESCRIPTION
Label	Contains an optional label which the assembler will assign as the symbolic address of the first byte of the instruction.
Operation	Contains any of the 2650 processor mnemonic operation codes as detailed in Appendix A, or any assembler Directive. This field may include an expression which specifies a register or value as required by the instruction. All symbols used in this field must have been previously defined, i.e. no symbolic forward references are allowed.
Operand	Contains one or more operand elements such as indirect address indicator, operand expression, index register specification, auto-increment/auto-decrement indicator, constant specification, etc., depending on the requirements of the particular instruction.
Comments	Any characters following the argument field will be reproduced in the assembly listing without processing. The Comments Field must be separated from the argument field by at least one blank.

Directives

There are eleven directives which the assembler will recognize. These assembler directives, although written much like processor instructions, are simply commands to the assembler instead of to the processor. They direct the assembler to perform specific tasks during the assembly process, but have no meaning to the 2650 processor. These assembler directives are:

MNEMONIC	FUNCTION
ORG	Set location counter
EQU	Specify a symbol
	equivalence
ACON	Define address constant
DATA	Defines memory data
RES	Reserve memory storage
END	End of assembly
EJE	Eject the listing page
PRT	Printer control
SPC	Space control
TITL	Title
PCM	Punch control

DESCRIPTION

The 2650 Simulator (PIPSIM) is a Fortran IV program which allows a user to simulate the execution of his program without utilizing the 2650 processor. The simulator executes the 2650 program via host computer software by maintaining its own internal Fortran storage registers to describe the 2650 program, the microprocessor registers, the ROM/RAM memory configuration, and the input data to be read dynamically from I/O devices. Inputs to the simulator are the object module (or the 2650 program in object format) produced by the 2650 assembler and a deck of user commands. The simulator can accommodate an object module of up to 8,192 bytes.

The output consists of a listing of the user's commands and a printout of both static and dynamic information as requested by the commands. The user may request traces of the processor status, dumps of the contents of memory, and recording of program timing statistics. Multiple simulations of the same program with different parameters may be executed during one simulation run.

FEATURES

- Cycle counter for timing estimates
- Instruction fetch break points
- . Operand fetch break points
- Trace facilities
- Snapshot dumps
- Patching facility
- Statistical Information generated
- Easy-to-use command language
- Optionally selected start and end addresses
- Simulated registers may be displayed while the simulation program is executed
- Simulated registers may be altered while the program is executing
- Maintain a 8K cell to simulate a read/ write RAM
- Capability exists for configuring parts of simulator memory to look like ROM
- Incorporates a 200-byte first in, first out (FIFO) buffer to store the data read from a simulated input device
- · Establishes initial program conditions
- Monitors execution sequences

AVAILABILITY

The 2650 Simulator is available on NCSS, TYMSHARE and GE timesharing services. It is also available from Signetics on a 9 track magnetic tape written in EBCDIC in 80 character unblocked records at a density of 800bpi.

The Simulator available on tape is configured to execute 16-bit machines (2650 SM1100) and 32-bit machines (2650 SM1000).

HIGHER LEVEL LANGUAGE (PLuS)

DESCRIPTION

The Signetics higher level language is designed for use with the 2650 microprocessor. This language allows the programmer to to reduce programming effort while retaining the control and efficiency of assembly language. It is written in ANSI standard Fortran IV and will execute on most machines without alteration. Programs written in this language tend to be self-documenting and are easily altered.

AVAILABILITY

The Signetics higher level language is available on both NCSS and GE timeshare. It is also available from Signetics on magnetic tape for 32-bit machines.

FEATURES

- Written in free-form
- Block structured
- Employs procedure calls
- Byte and address data elements
- Based variables
- In line assembly language
- Macro capability
- Generates relocatable code supported by a relocating loader
- Includes PL/M as a subset
- Allows separate compilation of program modules
- Has improved control structure over PL/M
- Conditional compilation
- Compile time expression evaluation

OVERVIEW OF THE LANGUAGE

The higher level language is a sequence of "Declarations" and "Executable Statements."

The declarations allow the programmer to control allocation of storage, define simple textual substitutions (Macros), and define procedures. The language is "Block Structured": Procedures may contain further declarations which control storage allocation and define other procedures.

The procedure definition facility of the language allows modular programming: A program can be divided into sections (e.g. teletype input, conversion from binary to decimal forms, and printing output messages). Each of these sections is written as a language procedure. Such procedures are conceptually simple, easy to formulate and debug, and easily incorporated into a large program. They may form a basis for a procedure library, if a family of similiar programs is being developed. Procedures may be individually compiled.

The language handles two kinds of data, its two basic "Data Types": Byte and address. A byte variable or constant is one that can be represented as an 8-bit quantity; an address variable or constant is a 16-bit or double-byte quantity. The programmer can declare variable names to represent byte or address values. One can also declare vectors (or arrays) or type byte or address.

In general, executable statements specify the computational processes that are to take place. To achieve this, arithmetic, logical (Boolean), and comparison (relational) operations are defined for variables and constants of both types (BYTE and ADDRESS). These operators and operands are combined to form EXPRESSIONS, which resemble those of elementary algebra. Expressions are a major component of language statements.

A simple statement form is the assignment statement, which computes a result and stores it in a memory location defined by a variable name. Other statements in the language perform conditional tests and branching, loop control, and procedure invocation with parameter passing. The flow of program execution is specified by means of powerful control structures that take advantage of the block-structured nature of the language. Input and output statements read and write 8-bit values from and to input and output ports. Procedures can be defined which use these basic input and output statements to perform more complicated I/O operations.

A method of automatic text-substitution (more specifically, a "compile-time macor facility") is also provided. A programmer can declare a symbolic name to be completely equivalent to an arbitrary sequence of characters. As each occurrence of the name is encountered by the compiler, the declared character sequence is substituted, so the compiler actually processes the substituted character string instead of the symbolic name.

The compiler supports compile time expression evaluation and conditional compilation which allows selective compilation of code depending on an input parameter at compile time.

The language generates absolute and/or relocatable code. The relocatable modules may be linked by a powerful linkage editor at load time.

Additionally the language contains all machine independent features of the PL/M language as a subset, thereby enhancing portability of programs.

DESCRIPTION

The 2650 assembly language is a symbolic language designed specifically to facilitate the writing of programs for the Signetics 2650 Microprocessor.

The 2650 relocatable assembler is a program which accepts symbolic source code as input and produces a listing, crossreference table of symbols, and an object module. The object module may be in absolute format compatible with the simulator (PIPSIM), PIPBUG and the 2650 TWIN system, or it may be a relocatable format for use with PLuS and the relocating loader.

The assembler is written in ANSI standard Fortran IV and is approximately 3200 Fortran card images in length. It operates in a two-pass mode to build a symbol table, issue helpful error messages, produce an easily readable program listing and crossreference listing, and output a computer readable object module.

AVAILABILITY

The 2650 relocatable assembler is available on NCSS, GE, and TYMSHARE timesharing services. It is also available from Signetics on 9-track magnetic tape written in EBCDIC in 80 character unblocked records at a density of 800 bpi.

FEATURES

- Pseudo-ops to aid programming
- Self-defining constants
- Symbolic machine operation codes
- Free format source code
- Syntax error checking
- Symbolic address assignment and references
- Data creation statements
- Storage reservation statements
- Assembly listing control statements
- Address constants
- **ASCII character codes**
- Comments and remarks for documentation
- Six-character symbols
- Cross-reference listing of symbols
- Automatic formatting of the listing
- Complex expression evaluation
- Conditional assembly
- **Boolean operators**
- Relocatable output compatible with the PLµS compiler output

LANGUAGE REQUIREMENTS Input Requirements

Input to the assembler consists of a sequence of chartacters combined to form assembly language elements. These language elements include symbols, instruction mnemonics, constants and expressions which make up the individual program statements that comprise a source program.

CHARACTERS

A through Z Alphabetic: Numeric: 0 through 9 Special Characters: blank left parenthesis right parenthesis add or positive value subtract or negative value asterisk single quote comma slash \$ dollar sign less than sign greater than sign pound sign @ at sign question mark exclamation double quote percent sign equal semicolon not symbol colon period

SYMBOLS

Symbols are formed from combinations of characters. Symbols provide a convenient means of identifying program elements so they can be referenced by other elements.

CONSTANTS

A constant is a self-defining language element. Unlike a symbol, the value of a constant is its own "face" value and is invariant. Internal numbers are represented in 2's complement notation. There are two forms in which constants may be written: the Self-Defining Constant and the General Constant.

SELF-DEFINING CONSTANT

The self-defining constant is a form of constant which is written directly in an instruction and defines a decimal value.

GENERAL CONSTANT

The general constant is also written directly in an instruction, but the interpretation of its value is dictated by a code character and delimited by quotation marks. Its form can be binary, octal, decimal, hexadecimal, or ASCII.

EXPRESSIONS

An expression is an assembly language element that represents a value. It consists of a single term or combination of terms separated by arithmetic or boolean operators. A term may be a valid symbolic reference, a self-defining constant or a general constant.

FIELDS

A statement prepared for processing by the assembler is logically divided into four fields, as indicated below. They are free form and are separated by at least one blank character. The name must begin in logical column 1.

LABEL	OPERATION	OPERAND
name	opcode	operand(s)

Where:	
FIELD	DESCRIPTION
Label	Contains an optional label which the assembler will assign as the symbolic address of the first byte of the instruction.
Operation	Contains any of the 2650 processor mnemonic operation codes as detailed in Appendix A of the reference manual, or any assembler directive. This field may include an expression which specifies a register or value as required by the instruction. All symbols used in this field must have been previously defined, i.e., no symbolic forward references are allowed.
Operand	Contains one or more operand elements such as indirect address indicator, operand expression, index register specification, auto-increment/auto-decrement indicator, constant specification, etc., depending on the requirements of the particular instruction.
Comments	Any characters following the argument field will be reproduced in the assembly listing without processing. The Comments Field must be separated from the argument field by at least one blank.

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Directives

There are fifteen directives which the assembler will recognize. These assembler directives, although written much like processor instructions, are simply commands to the assembler instead of to the processor. They direct the assembler to perform specific tasks during the assembly process, but have no meaning to the 2650 processor. These assembler directives are:

MNEMONIC	FUNCTION
ORG	Set location counter
EQU	Specify a symbol
	equivalence
ACON	Define address constant
DATA	Defines memory data
RES	Reserve memory storage
END	End of assembly
EJE	Eject the listing page
PRT	Printer control
SPC	Space control
TITL	Title
PCH	Punch control
IF)	
ENDIF >	Conditional assembly
ELSE '	
SEG	Begin a relocatable
	segment
ENDS	End a relocatable segment
ENTRY	Define an entry point
EXTRN	Define an external
	reference

Ģ

DESCRIPTION

The TWIN (TestWare Instrument) system, shown in the block diagram of Figure 1. consists of interdependent subsystems, each contributing to the total task of implementing user microprocessor applications from initial concept to actual hardware operation. The system closely resembles a general-purpose minicomputer during the initial stages of product development, and allows source programs to be entered, edited and assembled into object programs. Object programs may be executed simply as programs, or as part of a user's product emulation. When programs have been run and debugged to the user's satisfaction, the TWIN system is capable of programming PROM devices for inclusion in the user's prototype hardware.

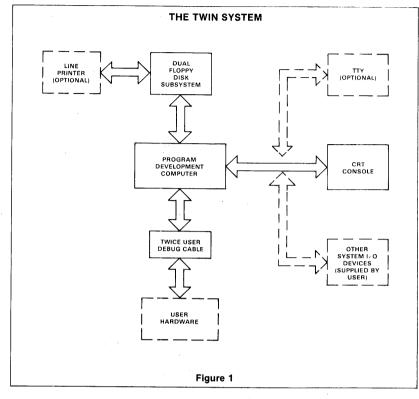
The program development computer is configured using plug-in modules, each dedicated to a specific task within the system. Each module interacts with the others by use of a common bus structure. The major buses include a data bus, an address bus, and a control bus. The program development computer acts as the controller for the entire system, and performs the necessary functions of transferring data to and from system peripherals.

To facilitate system operation, a dual-drive floppy disk subsystem is provided in all TWIN configurations. This subsystem, with integral controller, stores user program "files" and allows retrieval of these files for operations by the editor and assembler. The Signetics Disk Operating System (SDOS) software is loaded from the disk system into the development computer. SDOS provides complete control over all portions of the TWIN system.

A CRT console is the standard device for entering user programs. It is equipped with a full ASCII keyboard for data entry and a CRT display to allow the operator to view the results of his program manipulations during editing, assembly, and debug operations. The display and keyboard can be separated for user convenience. An interface is provided for an ASR-33 teletypewriter, which can be used in place of the CRT console. An optional line printer provides hard copy output.

Other system I/O devices can be added to the system to enhance its capabilities. Physical interface to the system bus is accomplished by use of the optional General-Purpose I/O card. Addition of a peripheral device requires addition of its software driver to the TWIN operating system.

TWICE (TestWare In-Circuit Emulation) denotes hardware/software elements of the



TWIN system that support the integration and check-out of the user's prototype product. A number of different TWICE operation modes support product development, integration, and production.

TWIN HARDWARE

The TWIN system is available in two configurations: the BASIC TWIN and the SUPER TWIN.

Basic TWIN Configuration

Program Development Computer consisting of:

Master CPU—16K Master Memory Slave CPU—16K Slave Memory Debug Logic Card TWICE Cable

Dual-Drive Floppy Disk Subsystem

Super TWIN Configuration

Basic TWIN configuration plus: 1702A PROM Program card 82S115 PROM Program card General-Purpose I/O Card CRT Console Line Printer

Program Development Computer

The Program Development Computer (PDC) is the principal subsystem of the TWIN. The

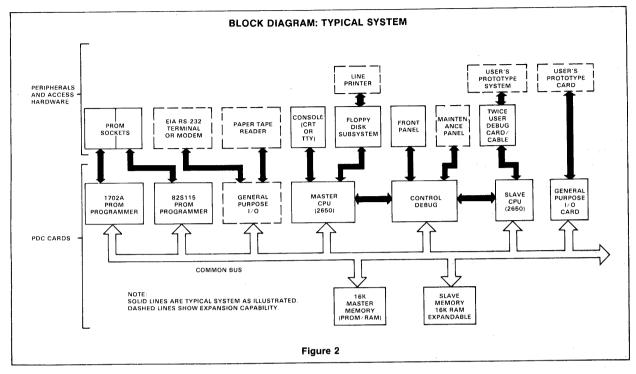
PDC employs dual-processor architecture operating on a common bus structure.

The Master/Slave architectural concept provides the following features:

- 1. Protected operating system memory and I/O.
- Complete user access to slave memory address space and I/O address space.
- 3. Operating system and Debug software are independent of user programs.
- Use of future Slave processors by adding boards and supplying additional software only.
- Multiple Slave CPUs of different types can be accommodated.
- 6. Debug and TWICE hardware.

The PDC can be divided into four distinct areas as shown in Figure 2.

- The master CPU provides access to all support peripherals (Floppy Disk, Line Printer and CRT console). By way of the common bus, the slave side of the system can make requests of these system peripherals. In addition, via the General Purpose I/O Board, the user has the option of adding system peripherals on the master side.
- The master CPU controls PROM programming by transferring data from the slave memory to the PROM sockets. Both erasable (1702A) and nonerasable (82S113) PROMs can be programmed through sockets controlled by the operating system.



3. The third functional area, known as TWICE, provides interface from the Program Development Computer (PDC) to the prototyping world. TWICE consists of a cable and driver/receiver hardware that allow in-circuit emulation of the user programs in the TWIN system. The user's microprocessor is unplugged from the prototype and replaced by the TWICE cable plugged directly into the vacant socket. The other end of the Program Development Computer cable connects to the slave CPU board. This board contains multiplexing and other logic to support the TWICE operating modes. The slave CPU thus becomes the CPU for the prototype system, allowing programs to be executed in a hardware environment nearly identical to that of the user's final product.

In addition, a General-Purpose I/O card or prototype card may be added to the system bus to provide interface with user prototype peripherals.

 The fourth subsystem is built around the control debug card that links the master CPU to the slave CPU. This card, in conjunction with the front panel and maintenance panel on the PDC, allows for Master/Slave communication under SDOS.

PROGRAM DEVELOPMENT COMPUTER SUBASSEMBLIES

Master CPU

Provides a resident microcomputer for the operating system software. Prioritizes and services all interrupts regardless of whether they were initiated from the master or slave

partitions.

Responsible for system I/O

Master Memory

This memory is reserved for the resident monitor and the overlays from disk necessary to service any valid user request. A minimum of 16K bytes of memory is required.

- Protected from slave access
- Four 4K statis RAM cards (or 16K dynamic card)
- 256 bytes of PROM—expandable to 2K bytes on one of the master memory cards.

Slave CPU

Provides an isolated microprocessor that enables the user to write and debug his program without endangering the operating system. The slave CPU also executes the user's program in his prototype via the TWICE cable.

- While the slave CPU is presently a 2650 microprocessor, future designs will support other microprocessors.
- Enables real-time execution of the user's programs in his prototype system.

Slave Memory

The primary function of the slave memory is to store programs. During PROM program-

ming, the debugged software logic in the slave memory is sent via the common bus to the PROM programming cards on the master side. During in-circuit emulation, the TWICE cable is connected through the Slave CPU to the Slave memory.

- Four 4K static RAM cards (or 16K dynamic card)
- Base address selectable in 4K increments
- · Expandable from 16K to 64K bytes.

Common Bus

These data, address and control channels provide communications between slave and master partitions as well as communications to other boards for outside world interfacing.

Debug Board

Designed to provide communications between slave and master partitions; it also supplies interface capability to the front panel.

- Two address breakpoint registers
- Two program counter registers
- Single cycle, halt, and diagnostic interrupts

PROM Programmer Boards

These cards, placed in two reserved slots of the master partition, provide access to front-panel PROM sockets on the PDC. They provide either erasable or nonerasable PROM programming of Signetics 82S115 bipolar PROMs and 1702A MOS PROMs.

- · Read, write and compare functions
- · Checks validity, power, and addresses
- Interfaces directly to bus
- Current limiting amplifier
- On-card rectifier/regulator

General-Purpose I/O Card

This card may be placed on the master side to support additional system peripherals and/or on the slave side to provide an interface to the user's hardware I/O simulator boards.

- Compatible with 20mA, 110 baud TTY interface
- Compatible with EIA/RS-232 interfaces
- Baud rates—110/150/300/600/1200, jumper selectable
- Four 8-bit input ports and four 8-bit output ports
- · Eight interrupts

Disk/Printer Controller Board

This card supports parallel I/O transfers to the floppy disk subsystem and to the optional printer. This board is contained in the floppy disk subsystem.

FUNCTIONAL HARDWARE/ SOFTWARE INTEGRATION

The TWIN system's operating functions can be grouped into three major modes as shown in Figure 3.

Software Development

Using the TWIN Assembler and Text Editor, the designer clears his programs of any syntax errors. The system I/O devices interfaced through the master side of the TWIN enable the user to develop and debug his software. The software is verified as much as possible before actually running it in the user's prototype system. All memory and I/O are contained in the TWIN. As shown in Figure 3A, there is no interconnection with the user's prototype system.

Software/Hardware Integration

The final development phases are accomplished using the TWICE capability. In these phases, the user can easily complete the normally complex function of hardware/software system debugging.

PROTOTYPE HARDWARE INTERFACE (TWICE): The TWIN memory is used by the Slave, but I/O is provided to the user's

system via the TWICE cable as shown in Figure 3B. The user's program that had been loaded into Slave Memory by SDOS is executed by the user's prototype circuits.

PROM INTEGRATION (TWICE): All program memory and prototype I/O are contained in the user's system as shown in Figure 3C. The operator can still control breakpoints, starting addresses and obtain a partial Trace. This mode is used after PROMs and ROMs have been prepared. This mode verifies the final system configuration in real-time operation.

TWIN DEVELOPMENT SOFTWARE

System software provided with the Prototype Development System includes the Signetics Disk Operating System (SDOS), Text Editor, Assembler, and Debug Package.

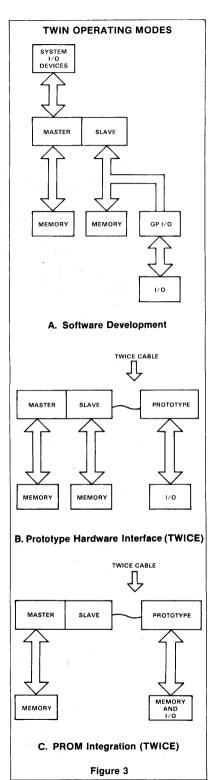
SDOS

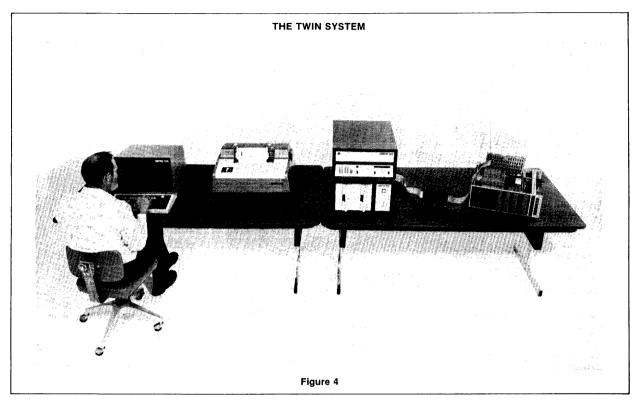
SDOS relieves the user of the necessity of understanding the detailed internal operation of the TWIN. It provides complete control over operation of all portions of the Prototype Development System. All functions relating to file handling, loading and execution are monitored and controlled by SDOS, including the in-circuit emulation and the PROM programming functions. SDOS is written in 2650 assembly language and resides in a dedicated memory consisting of a 256-byte PROM and 16K RAM running under the master CPU.

The SDOS software allows the user to create, edit, and assemble files; obtain object and listing outputs; load and execute programs; and, through the debug system, check out programs in an efficient manner.

SDOS provides a powerful command file capability that enables the user to create customized operating system commands.

SDOS controls the multi-drive floppy disk subsystem, a line printer, and the CRT console, which may be an ASR-33 TTY or an RS-232 compatible device. Software drivers are provided within SDOS for these I/O devices. The SDOS software provides a flexible input/output system that enables the user to dynamically assign any logical channel to any physical device or file within the system. Thus, system I/O devices may be dynamically assigned using SDOS commands either from the console or from within a user's program. Thus, the user may write his own driver for other peripheral devices and link them into the SDOS system by use of the optional General-Purpose I/O Card.





Text Editor

The Text Editor is a comprehensive software package that allows the user to enter and modify text files. The Text Editor is line oriented and accepts inputs from the terminal or a diskette file, performs modifications in a work space, and outputs the revised text to diskette file.

Assembler

The Prototype Development System Resident Assembler translates symbolic assembly language instructions into appropriate machine language code.

The Assembler is written in the Signetics Higher-Level Language ($PL\mu S$) and generates absolute object code. This code is in hexadecimal format and may be loaded into the system for direct execution or may be converted by an SDOS command to SMS format for PROM or ROM programming.

Debug

The debug system is a software program which provides the user with real-time program debug capabilities within both a software and hardware environment. It uses special hardware features built into the program development system to control the execution of the user's program. User pro-

PHYSICAL AND ENVIRONMENTAL SPECIFICATIONS

	VOLTAGE	VOLT AMPERES	LINE FRE- QUENCY (HERTZ)	WEIGHT (POUNDS)	DIMEN- SIONS (INCHES)
Development Computer Dual-Drive Floppy Disk CRT Console Printer		230 400 130 375	47 — 63 59.5 — 60.5 47 — 63 59.5 — 60.5	52	8 x 17 x 22 10 x 17 x 22 14 x 17 x 22 12 x 23 x 19

NOTE

Export version will operate from other line voltages and frequencies.

Table 1 SUBSYSTEM CHARACTERISTICS

grams operating under the debug system have dynamic program trace, breakpoint capabilities, and memory modification capabilities. Status reporting on the memory, the program, and the processor is also provided.

Total System Characteristics

Operating Temperature

Range: 15° — 37°C

Operating Humidity

Range: 50% — 80% NC

Basic system shipping

weight: 180 pounds

Card Cage—Mother Board Assembly

7 x 11 cards/20 slots in PDC

Common Bus Structure

External +5V ground bus (spare with rear panel terminals)

Front Panel Switches and Indicators

Key-operated primary power On/Off switch PROM power switch

Reset switch that initializes system and sets both CPUs to location zero, enables master, and pauses slave.

Diagnostic interrupt switch

Power indicator

Run indicator

Hun indicator

Master indicator Slave indicator

PROM power indicator

SDOS COMMAND SUMMARY

System Utilities

COMMAND	FUNCTION
FORMAT D	Formats the diskette on drive D.
VERIFY D	Verifies the diskette on drive D.
RENAME D	Re-identifies the diskette on drive D.
DUP	Duplicates a diskette.
LDIR (D)	Lists the directory of the diskette.
DELETE	Deletes a diskette file.
COPY	Copies data from a speci- fied input device or file to a specified output device or
DOINT (L)	file.
PHINT (L)	Prints from a disk file to a specified output device. If (L) is included, the lines are numbered.

System Control

COMMAND	FUNCTION
SUSPEND	Suspends execution of an
	active program.
CONT	Continues execution of a
	suspended program.
ABORT	Aborts an active SDOS or
	user program.
ESCAPE*	Interrupts the execution of
	a program.
DOUBLE	Suspends execution of all
ESCAPE*	programs.
SPACE*	Stops or starts the console
	output.
STATUS	Gives the status of the pro-
	gram being executed by the
	Slave CPU.
ASSIGN	Connects a logical channel
	to a specified device.
CLOSE	Closes the specified chan-
	nels.
	*Keys on console

Object Program Utilities

COMMAND	FUNCTION
RHEX	Loads an absolute hexa- decimal object file into slave memory.
WHEX	Outputs an absolute hexadecimal data file from slave memory.
WSMS	Outputs a 512-byte block of memory in SMS format.
CSMS	Reads an SMS file and compares it with the contents of memory.
MODULE	Creates an absolute binary program load module from the Slave CPUs memory.

Debug Commands

COMMAN	D FUNCTION
DEBUG	Enables debug mode for
	user programs.
BKPT	Selects a program break-
	point.
	Clears breakpoint.
SET	Modifies Slave CPU register(s).
DOTAT	• •
DOTAL	Displays the Slave CPU
	registers and the selected
TD 4 0.5	breakpoints.
	Sets program trace control.
	Alters slave memory.
	Dumps slave memory.
EXAM	Examines and optionally
	alters slave memory.
SLAVE	Designates the active Slave
	CPU and sets its mode of
	operation.
RESET	Resets the Slave CPU.
LOAD	Loads a program load mod-
	ule into the Slave CPUs
	memory.
XEQ	Loads and executes a pro-
	gram load module.
GO	Executes a previously
	loaded load module or
	transfers control to an ab-
	solute memory location in
	slave memory.

Command File

COMMAN	ID FUNCTION
TYPE	Controls command line typeout during execution of a command file.
	Aborts or continues a command file if an error is encountered. Implements comments within a command file.

Editor

COMMAND	FUI	NCTION
EDIT Loads Text E		executes the

Assembler

COMMAI	ND	FUI		
_	Loads 2650 as		executes oler.	the

PROM Utilities

COMMAND	FUNCTION
RPROM	Reads the contents of PROM into slave memory.
CPROM	Compares the contents of slave memory with the contents of PROM
WPROM	Burns a PROM with the contents of slave memory.
System O _l	otions

Sy

CC	MMAND	FUNCTION
	SYSTEM	Designates the system drive.
	DEVICE	Informs SDOS of device status.
1	SEARCH	Specifies file search option.

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TEXT EDITOR COMMAND SUMMARY

COMMAN	D FUNCTION		COMMAND	FUNCTION
AGAIN	Repeats the last "repeatable" command entered.		LIST N	Lists n lines or lines p through q on the line print-
BEGIN	Sets the current line pointer			er.
	to the first line of the work			Creates a macro.
	space.		MACRO m	Executes the previously de-
BRIEF	Reverses the initial mode of			fined macro m.
	operation.		N	Types the current line prin-
COPY N	Copies n lines, or lines p			ter position on the console.
	through q, from the input		PUT(K) N	Puts n lines starting at the
	file FILENAME ₁ to the out-			current line or puts lines p
	put file FILENAME ₂ .			through q from the work-
DOWN n	Moves the current line			space to the primary output
	pointer down n lines.			file.
END	Sets the current line pointer	1	QUIT	Closes the primary input
	to the end of text (last line			and primary output files
	plus one).			and quits the editor.
FILE	Transfers all workspace		REPLACE	Replaces the current line
	text and remaining primary			with a character string.
	input file text to the primary		SDOS	Suspends the text editor
	output file and quits the			and returns control to
E11.15	editor.		0007.7175	SDOS.
FIND	Moves the current line		SOBSTITUTE	Substitutes text on the cur-
	pointer down to the line		T4D.0	rent line.
	containing the first occur-		TABC	Defines the character C as
OFT N	rence of a character string.		TARC	the tab character.
GETN	Gets n lines or lines p			Sets tab positions to speci- fied columns.
	through q into the work-			Types n lines or lines p
INDLIT	space. Establishes the "input"			through q on the console.
INFOI	mode.			Moves the current line
INSEDT	Inserts a given string above		01 11	pointer up n lines.
INOLITI	the current line.		m<>	Repeats the command
KILLN	Deletes n lines starting with			string enclosed by the
NILL IN	the current line, or deletes			brackets m times.
	lines p through q.			Types the current I/O sta-
	mico p unough q.	1		tus on the console.

CHAPTER 3 Standard Support Circuits

Signetics 277

INTRODUCTION

In addition to the dedicated support circuits available for the various Signetics microprocessors, Signetics offers a complete line of standard circuits to complete the design of a microcomputer system.

A complete line of Schottky-clamped TTL, read/write memory arrays is offered. All feature open collector or tri-state output options for optimization of word expansion in bused organizations. Memory expansion is further enhanced by full on-chip address decoding, chip enable function and PNP input transistors which reduce input loading requirements. All devices offer high performance read access and write cycle times making these devices ideally suited in high speed memory applications such as "cache," buffers, scratch pads, writable control store, main store, etc.

Signetics offers the industry's broadest line of bipolar high performance ROMs, PROMs

and FPLAs. The PROMs and FPLAs are field programmable, which means that custom patterns are immediately available by following the provided fusing procedures. Signetics PROMs are supplied with all outputs at logical 0. Outputs are programmed to a logic 1 at any specified address by fusing a Ni-Cr link matrix. All bipolar ROMs, PROMs and FPLAs are fully TTL compatible, and include on-chip decoding and chip enable functions for ease of memory expansion. Tri-state and open collector output functions are available, and low input currents reduce input buffer requirements. Most Signetics PROMs and FPLAs also have pin and performance compatible ROMs and PLAs. offering the user the ultimate in flexibility and cost reduction.

Signetics n-channel MOS products include a complete family of 1k static RAMs and 8k static ROMs. These feature TTL compatible inputs and outputs, and require only a single +5V power supply. A variety of 4k dynamic RAMs is also available for system configurations requiring large amounts of read/write memory.

The 8T series of interface devices includes display drivers, bus drivers, input/output ports, level converters, and special purpose devices. A complete line of standard and low power Schottky (LS) 74 series devices is available in addition to the 8200 series of MSI devices. Many devices from the Signetics analog product line are also suitable for use in microcomputer systems. These include voltage regulators, operational amplifiers, comparators and timers.

This chapter provides product line summaries and data for a representative selection of Signetics standard support circuits.

BIPOLAR MEMORY SELECTION GUIDE

DEVICE	DEVICE ORGANI- OUTF		OUTPUT Logic ²	ACCESS TIME [ns] ⁴	TEMPERATURE RANGE ³	PACKAGE	NO. OF PINS	MAX. ICC [mA] ⁴
CAMS 10155	8X2	0E	_	13	С	F,N	18	140
SAMS 82S12 82S112	8X4 8X4	OC TS	T T .	40 40	C C	F,N F,N	24 24	160 160
RAMS 82S25 3101A 54/74S89 54/74S189	16X4 16X4 16X4 16X4	0C 0C 0C TS	В В Т В	50 35 50 35	M,C M,C M,C M,C	F,N F,N F,N	16 16 16 16	105 105 105 110
82S21 82S16 82S116 82S17 82S117 54/74S200 54/74S201 54/74S301	32X2 256X1 256X1 256X1 256X1 256X1 256X1 256X1	0C TS TS 0C 0C TS TS	T T T T B B B	50 50 40 50 40 50 50 50 50	C M,C C M,C C M,C M,C M,C	F,N F,N F,N F,N F,N F,N	16 16 16 16 16 16 16	130 115 115 115 115 130 130
82S09 82S10 82S110 82S11	64X9 1024X1 1024X1 1024X1	0C 0C 0C TS	Т В В В	45 45 35 45	M,C M,C C M,C	I,N F,N F,N F,N	28 16 16 16	190 170 170 170
82S111 93415A 93425A	1024X1 1024X1 1024X1	TS OC TS	B B B	35 45 45	C M.C M.C	F.N F.N F.N	16 16 16	170 170 170
82S208* 82S210* 82S400* 82S401*	256X8 256X9 4096X1 4096X1	TS TS OC TS	B B B B	60 60 70 70	C C C	F F,N I	22 24 18 18	185 185 155 155
ROMS 82S226 82S229	256X4 256X4	OC TS	<u>-</u>	50 50	M,C M,C	F,N F,N	16 16	120 120
82S214 82S230	256X8 512X4	TS OC		60 50	M,C M,C	F,N	24 16	175 140 140
82S231 82S215 82S240 82S241	512X4 512X8 512X8 512X8	TS TS OC TS	_ _ _ _	50 60 60 60	M.C M.C M.C M.C	F,N F,N F,N F,N	16 24 24 24 24	175 175 175
8228 82\$280 82\$281	1024X4 1024X8 1024X8	TTL OC TS	_ _ _	50 70 70	C M,C M,C	F,N F,N	16 24 24	170 140 140
82S290 82S291	2048X8 2048X8	OC TS		80 80	M,C M,C	F,N F,N	24 24	170 170

^{*}To be announced

NOTES

3. Temperature range:
C = Commercial (0 ° C to +75 ° C)
M = Military (-55 ° C to +125 ° C)
All ECL 10,000 series (-30 ° C to +85 ° C)

4. Commercial (0°C to +75°C)

^{1.} Output circuit:

OE = Open emitter

OC = Open collector

TS = Tri-state

2. Output logic:

T = Transparent—Input data appears on output during Write

B = Blanked—output is blanked during Write

BIPOLAR MEMORY SELECTION GUIDE (Cont'd)

DEVICE	ORGANI- ZATION	OUTPUT CIRCUIT	OUTPUT Logic ²	ACCESS TIME (ns)	TEMPERATURE RANGE ³	PACKAGE	NO. OF PINS	MAX.	EQUIVA- Lent rom
PROMS 82S23 82S123 10139	32X8 32X8 32X8	OC TS OE	_ _ _	50. 50 15.	M,C M,C C	F,N F,N F,N	16 16 16	77 .77 .145	
82S27 82S126 82S129 10149	256X4 256X4 256X4 256X4	0C 0C TS 0E	_ _ 	40 50 50 20	C M,C M,C C	F F,N F,N F	16 16 16 16	140 120 120 150	82S226 82S229 —
82S114 82S130 82S131	256X8 512X4 512X4	TS OC TS	_ 	60 50 50	M,C M,C M,C	F,N F,N F,N	24 16 16	175 140 140	82\$214 82\$230 82\$231
82S115 82S140 82S141	512X8 512X8 512X8	TS OC TS	_ _ _	60 60 60	M,C M,C M,C	F,N F,N F,N	24 24 24	175 175 175	82S215 82S240 82S241
82S136 82S137	1024X4 1024X4	OC TS		60 60	M,C M,C	F,N F,N	18 18	140 140	
82S180 82S181 82S2708	1024X8 1024X8 1024X8	0C TS TS	<u> </u>	70 70 70	M,C M,C M,C	F,N F,N F,N	24 24 24	175 175 175	82\$280 82\$281 —
82S 184 82S 185	2048X4 2048X4	OC TS	_	100 100	M,C M,C	·	18 18	120 120	· — —
82S190 82S191	2048X8 2048X8	OC TS	<u> </u>	70 70	M,C M,C	-	24 24	175 175	82S290 82S291
FPLAS 82S100 82S101	16X48X8 16X48X8	TS 0C	<u></u>	50 50	M,C M,C	1,N 1,N	28 28	170 170	<u>-</u> -
PLAS 82S200 82S201	16X48X8 16X48X8	TS 0C	_	.50 50	M,C M,C	I,N I,N	28 28	170 170	
FPGAS 82S102 82S103	16X9 16X9	OC TS	_	30 30	M,C M,C	1,N 1,N	28 28	170 170	_

^{*}To be announced

NOTES

^{1.} Output circuit:

OE = Open emitter
OC = Open collector
TS = Tri-state

^{2.} Output logic:

T = Transparent—input data appears on output during Write
B = Blanked—output is blanked during Write

^{3.} Temperature range: C = Commercial (0° C to +75° C) M = Military (-55° C to +125° C)

All ECL 10,000 series (-30° C to +85° C)

^{4.} Commercial (0°C to +75°C)

MOS MEMORY SELECTION GUIDE

OS WIEWO	HT SELECT	ION GOID	<u> </u>				г	
DEVICE	ORGANIZATION	OUTPUT Circuit'	ACCESS/CYCLE TIME (ns)	TEMPERATURE RANGE?	PACKAGE	NO. OF PINS	CLOCK/CE/TTL COMPATABILITY	POWER Supplies (V)
RAMS								
Static								
2501	256X1	TTL	1000/1000	С	I, N	16	Yes	+5, -9
25L01	256X1	TTL	1000/1000	Č	I, N	16	Yes	±5, -12
2101	256X4	TS	1000/1000	Č	F. N	22	Yes	+5, Gnd
2101-1	256X4	TS	500/500	Č	F, N	22	Yes	+5, Gnd
2101-2	256X4	TS	650/650	č	F. N	22	Yes	+5, Gnd
2111	256X4	TS	1000/1000	Č	I. N	18	Yes	+5, Gnd
2111-1	256X4	TS	500/500	Č	I, N	18	Yes	+5, Gnd
2111-2	256X4	TS	650/650	Č	I, N	18	Yes	+5, Gnd
2112	256X4	TS	1000/1000	Č	F. N	16	Yes	+5, Gnd
2112-1	256X4	TS	500/500	l č	F, N	16	Yes	+5, Gnd
2112-1	256X4 256X4	TS	650/650	C	F, N	16	Yes	+5, Gnd
2606	256X4	TS	750/750	C	F, I, N	16	Yes	+5, Gnd
		TS	500/500	Č	F, I, N	16	Yes	+5, Gnd
2606-1	256X4				F, I, N	16	Yes	+5; Gnd
2102	1024X1	TS	1000/1000	C		16	Yes	+5, Gnd
2102-1	1024X1	TS	500/500	, ,	F, I, N	16	Yes	+5, Gnd
2102-2	1024X1	TS	650/650	C	F, I, N		Yes	+5, Gnd
2102A	1024X1	TS	350/350		F, I, N	16	Yes	+5, Gnd
2102AL	1024X1	TS	350/350		F, I, N	16		+5, Gnd
2102A-2	1024X1	TS	250/250	CCC	F, I, N	16	Yes	+5, Gnd
2102AL-2	1024X1	TS	250/250		F, I, N	16	Yes	+5, Gnd
2102A-4	1024X1	TS	450/450		F, I, N	16	Yes	+5, Gnd
2102AL-4	1024X1	TS	450/450	000000000000000000000000000000000000000	F, I, N	16	Yes	+5, Gnd
2102A-6	1024X1	TS	650/650	C	F, I, N	16	Yes	
21F02	1024X1	TS	350/350	C	F, I, N	16	Yes	+5, Gnd
210F02-2	1024X1	TS	250/250	C	F, I, N	16	Yes	+5, Gnd
21F02-4	1024X1	TS	450/450	C	F, I, N	16	Yes	+5, Gnd
21L02	1024X1	TS	1000/1000	C	F, I, N	16	Yes	+5, Gnd
21L02-1	1024X1	TS	500/500	C	F, I, N	16	Yes	+5, Gnd
21L02-2	1024X1	TS	650/650	C	F, I, N	16	Yes	+5, Gnd
21L02-3	1024X1	TS	400/400	C	F, I, N	16	Yes	+5, Gnd
2115*	1024X1	OD	45/45	C	F, I, N	16	Yes	+5, Gnd
2115L*	1024X1	OD	45/45	C	F, I, N	16	Yes	+5, Gnd
2125*	1024X1	TS	45/45	C	F, I, N	16	Yes	+5, Gnd
2125L*	1024X1	TS	45/45	C	F, I, N	16	Yes	+5, Gnd
2614*	1024X4	TS	200/200	C	F, I, N	18		+5, Gnd
2613*	4096X1	TS	200/200	С	F, I, N	18		+5, Gnd
Dynamic								
1103	1024X1	OD	300/480	С	I, N	18	No	+20, +16, Gno
2660	4096X1	TS	250/375	Č	F, I, N	16	Yes	+12, ±5, Gno
2660-1	4096X1	TS	300/425	Č	F, I, N	16	Yes	+12, ±5, Gnd
2660-2	4096X1	TS	350/500	l č	F, I, N	16	Yes	+12, ±5, Gno
2660-3	4096X1	TS	140/375	C	F, I, N	16	Yes	+12, ±5, Gnd
2680	4096X1	TS	200/400	Č	F, I, N	22	No	+12, ±5, Gnd
2680-1	4096X1	TS	270/470	Č	F. I. N	22	No	+12, ±5, Gnd
2680-1	4096X1	TS	350/800	C	F, I, N	22	No	+12, ±5, Gnd
2600-2 2627*	4096X1 4096X1	15	150/320	C	F F	16		+12, ±5, Gnd
2627-1*	4096X1 4096X1	_	200/200	C	F	16	_	+12, ±5, Gnd
2627-1* 2627-2*	4096X1 4096X1	_	250/250	C	F	16	_	+12, ±5, Gno
2627-2 2690*	16,384X1	_	150/375	Č		16	_	+12, ±5, Gnd
7050	1 10,50471	_	1 100/0/0	, 0	1	1 10	i .	1 12, -0, dillo

*To be announced

NOTES

NOTES

1. Output circuit:
TS = Tri-state
OD = Open drain
BD = Bare drain
PD = Pull down
PP = Push-pull
2. Temperature range:
C = Commercial (0 °C to +75 °C)
M = Military (-55 °C to +125 °C)

MOS MEMORY SELECTION GUIDE (Cont'd)

DEVICE	ORGANIZATION	OUTPUT CIRCUIT'	ACCESS/CYCLE TIME (ns)	TEMPERATURE RANGE ²	PACKAGE	NO. OF PINS	CLOCK/CE/TTL COMPATABILITY	POWER Supplies (V)
ROMS					· · · · · · · · · · · · · · · · · · ·			
Static								
2530	512X8	TS	700/700	l c	I, N	24	Yes	±5, -12
2609	128X9X7	TS	500/500	C	F, I, N	24	Yes	+5, Gnd
2607	1024X8	TS	500/500	l c l	F, I, N	24	Yes	+5, Gnd
2608	1024X8	TS	550/550	C	F, I, N	24	Yes	+5, Gnd
2608-1	1024X8	TS	450/450	C	F, I, N	24	Yes	+5, Gnd
2580	2048X4	TS	950/950	l c l	i, N	24	Yes	+5, -12
2600	2048X8	TS	500/500	l c l	F, I, N	24	Yes	+5, Gnd
2600-1	2048X8	TS	300/300	C .	F, I, N	24	Yes	+5, Gnd
2616	2048X8	TS	450	C	F, I, N	24	Yes	+5, Gnd
2616-1	2048X8	TS	350	C	F, I, N.	24	Yes	+5, Gnd
2617	2048X8	TS	450	l c l	F, I, N	- 24	Yes	+5, Gnd
2617-1	2048X8	TS	350	C	F, I, N	24	Yes	+5, Gnd
2632*	4096X8	_	500/500	C . C	i, N	24	_	+5, Gnd
2633*	4096X8		450/450	C	1, N	24	_	+5, Gnd
CHARACTER	GENERATORS					***************************************		
2513	64X8X5	TS	600/600	· c	I, N	24	Yes	±5, -12
2516	64X6X8	TS	600/600	Č	I, N	24	Yes	±5, -12
2526	64X9X9	TS	700/700	· · · C	i, N	24	Yes	+5, -12
UV EPROMS		:						
1702A	256X8	TS	1000/1000	C .	ı	24	Yes	+5, -9
2704	512X8	TS	450/450	C C	i	24	Yes	+12, ± 5, Gnd
2708	1024X8	TS	450/450	Č l	i	24	Yes	+12, ±5, Gnd

STANDARD ROM CODE

DEVICE	CODE NO.	DESCRIPTION
STATIC ROM		
2530	CM3530	Code Converter, ASCH to EBCDIC and EBCDIC to ASCII
2608	CN0000	10X7 Upper and Lower Case ASCII Character Generator
2609	CN6571	128 ASCII Characters in 7X9 Matrix Count Down
	CN6571A	128 ASCII Characters in 7X9 Matrix Count Up
	CN6575	128 ASCII Characters in 7X9 Matrix Count Up with Special Characters
2580	CMXXXX	Random code pattern for evaluation purposes
CHARACTER GI	NERATOR	
2513	CM2140	New ASCII Character Generator, Upper Case, 7X5, Horizontal Scan
	CM2170	ASCII Character Generator, Upper Case with Yen Sign, 7X5, Horizontal Scan
	CM3021	ASCII Character Generator, Lower Case, 7X5, Horizontal Scan
	CM3030	Old ASCII Character Generator, Upper Case, 7X5, Horizontal Scan
05.40	CM4800	Katakana Character Generator, 7X5, Horizontal Scan
2516	CM2150	ASCII Character Generator, Upper Case, 5X7, Vertical Scan
	CM3001/3010	ASCII Character Generator, Upper Case, 10X7, Vertical Scan (2 chips)
	CM3041	ASCII Character Generator, Lower Case, 10X7, Vertical Scan
2506	CM3970/3980	ASCII Character Generator, Upper Case, 12X8, Vertical Scan (2 chips)
2526	CM3400	ASCII Character Generator with EBCDIC and BAUDOT code translations, Upper Case, 7X9, Vertical Scan
	CM3940 CM6760	ASCII Character Generator, Upper Case, 7X9, Horizontal Scan
	CIVIO/OU	Katakana Character Generator, 7X9, Horizontal Scan

^{*}To be announced NOTES

NOTES

1. Output circuit:
TS = Tri-state
OD = Open drain
BD = Bare drain
PD = Pull down
PP = Push-pull
2. Temperature range:
C = Commercial (0 ° C to +75 ° C)
M = Military (-55 ° C to +125 ° C)

MOS MEMORY SELECTION GUIDE (Cont'd)

DEVICE	ORGANIZATION	OUTPUT CIRCUIT'	ON CHIP RECIRCULATE	TEMPERATURE RANGE ²	PACKAGE	NO. OF PINS	NO. OF CLOCKS	TYPICAL SPEED (MHz)	CLOCK/CE/TTL Compatability	POWER Supplies (V)
SHIFT REG	SISTERS									
Static				ļ						
2518	32X6	BD	Yes	С	N	16	1	3.0	Yes	+5, -12
2519	40X6	BD	Yes	С	N	16	1	3.0	Yes	+5, -12
2509	50X2	TS	Yes	C	N. K	14/10	1	3.0	Yes	+5, -5, -12
2532	80X4	PP	Yes	C C C	N I	16	1	3.0	Yes	+5, -12
2510	100X2	TS	Yes	C	N. K	14/10	1	3.0	Yes	+5, -5, -12
2521	128X2	PP	Yes	C	N	8	1	3.0	Yes	+5, -12
2522	132X2	PP	Yes	C C C C C	N	8	1	3.0	Yes	+5, -12
2511	200X2	TS	Yes	С	N, K	14/10	1	3.0	Yes	+5, -5, -12
2527	240X2	PP	Yes	C	N	8	1	2.5	Yes	+5, -12
2528	250X2	PP	Yes	C	N	8	1	2.5	Yes	+5, -12
2529	256X2	PP	Yes) C	N	8	1	3.0	Yes	+5, -12
2533	1024X1	PP	Jumper	C	N	8	1	2.0	Yes	+5, –12
Dynamic										
2506	100X2	BD	No	C	T. N	8	2	4.0	No	+5, -5
2507	100X2	7.5KPD	No	C	T. N	8	2	4.0	No	+5, -5
2517	100X2	20KPD	No	C	T, N	8	2	4.0	No	+5, -5
2505	512X1	BD	Yes	C	K	10	2	3.0	No	+5, -5
2524	512X1	BD	Yes	C	N	8	2	5.0	No	+5, -5 +5, -5
2502	256X4	BD	No	C C C C	_ N	16	2 2	10.0	No No	+5, -5 +5, -5
2503	512X2	BD	No	C	TA. N	8	2	10.0	No No	+5, -5 +5, -5
2504	1024X1	BD	No	C	TA, N	8	2	10.0	No No	+5, -5
2512	1024X1	BD	Yes	C	K	10	2 2	5.0 3.0	No	+5, -5 +5, -5
2525	1024X1	BD	Yes	C .	N	8	2	3.0	INU	10, -0

*To be announced

NOTES

NOTES

1. Output circuit
TS = Tri-state
OD = Open drain
BD = Bare drain
PD = Pull down
PP = Push-pull
2. Temperature range:
C = Commercial (0 ° C to +75 ° C)
M = Military (-55 ° C to +125 ° C)

7400 SERIES

		7400	74H	74LS	748
DEVICE	DESCRIPTION	PLASTIC Cerdip	PLASTIC CERDIP	PLASTIC Cerdip	PLASTIC CERDIP
74178	4-Bit Parallel Access Shift Reg (8270)	ΝF			N F
74179	4-Bit Parallel Access Shift Register (8271)	NF			NF
74180 74181	8-Bit Odd/Even Parity Checker	NF			
74181 74182	4-Bit Arithmetic Logic Unit	NF		NF	NF
74102	Look-Ahead Carry Generator	NF			N F
74190	Synchronous Up/Down Counter (BCD) Synchronous Up/Down Counter (Binary)	N F N F	# ·	NF	
74192	Synchronous Decade Up/Down Counter Synchronous Decade Up/Down Counter	N F	$H^{\prime\prime}$ T_{12}	N F N F	
74193	Synchronous 4-Bit Binary Up/Down Counter	N F N F		NF	
74194	4-Bif Bidirectional Universal Shift Reg.	NF		N F	N F
74195	4-Bit Parallel-Access Shift Register	NF	1 111 227 (8)	N F	N F
74196	Presettable Decade Counter/Latch (8290)	NF		NF	N F
74197	Presettable Binary Counter/Latch (8291)	N F		NF	NF
74198	8-Bit Shift Register	NF			
74199	8-Bit Shift Register	NF			
74221	Dual Monostable Multivibrator	N.F	400	N F	#4
74251	Data Selector/Mux with 3-State Outputs	4.4	4.	NF	NF
74253	Dual 4-Line to 1-Line Data Selector/Mux			N F N F	NF
74257	Quad 2-Line to 1-Line Data Selector/Mux			NF	NF
74258	Quad 2-Line to 1-Line Data Selector/Mux	150		NF	N F
74260	Dual 5-Input NOR Gate			N F	ΝF
74261	2X4 Parallel Binary Multiplier			NF	
74266 7 4273	Quad Exclusive NOR Gate			NF	
74273	Octal D Flip-Flop			NF	
74279	Quad S-R Latch	NF			real/dus
74283	9-Bit Odd/Even Parity Generator/Checker 4-Bit Adder	#####	-		NF
74290	Decade Counter			N F	
74293	4-Bit Binary Counter			N F	
74295A	4-Bit Right-Shift Left-Shift Register			N F	it Ti
74298	Quad 2-Input Mux with Storage	NF	W. 7.	N F N F	i l ta nstani
74365	Hex Buffer w/Common Enable (3-State)	NF		NF	
74366	Hex Buffer w/Common Enable (3-State)	NF		NF	
74367	Hex Buffer, 4-Bit and 2-Bit (3-State)	N F		NF	
74368	Hex Inverter, 4-Bit and 2-Bit (3-State)	NF		N F	
74375	Quad Latch	3 1987 118	4/4/	NF	W W
74377	Octal D Flip-Flop	$\mathbb{Z}[\mathbb{Z}]$		NF	
74386	Quad Exclusive-OR Gate	147,44	41 IE	N F	
74395	4-Bit Cascadeable Shift Register (3-State)		$-\downarrow$	NF	
74670	4X4 Register File (Tri-state)	=		NF	
8.50m3.00.00 EUS-968.00 第2 3 3		图 3 开始 6 克克尔克克	953 8 M X V	923862683	Na at Properties

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LOGIC-8200/9300/9600 SERIES

		JM38510		AN LIFIED*	JA PROCI	N ESSED	RE	MIL L/883 TEMP
DEVICE	DESCRIPTION	SLASH	Dip	Flat Pack	Dip	Flat Pack	Dip	Flat Pack
8200	Dual 5-Bit Buffer Register	_	-		_		1	Q
8201	Dual 5-Bit Buffer Register with D Inputs						!	Q Q
8202	10-Bit Buffer Register	_	_		_		!	Q
8203	10-Bit Buffer Register with D Inputs		-		_	w		w
8230	8-Input Digital Multiplexer	/01402	esa in winds yet		F -7.7542390	VV	F F	W
8231	8-Input Digital Multiplexer		+170	+ //) (1) F	W
8232	8-Input Digital Multiplexer				Sign Tall		F	W
8233	2-Input 4-Bit Digital Multiplexer		F Pini		- T		112	W.
8234	2-Input 4-Bit Digital Multiplexer	Listensia			$m_{\tilde{t}}$		2	W
8235	2-Input 4-Bit Digital Multiplexer		1 T			and There	F	W
8241	Quad Exclusive-OR Gate	_	_		_		F	W
8242	Quad Exclusive-NOR Gate	_			_			
8243	8-Bit Position Scaler		_		_	 W	F	Q W
8250	Binary-to-Octal Decoder	/15204	2	2 :	F	W		W
8251	BCD-to-Decimal Decoder	/15205	2	2	i		F	W
8252	BCD-to-Decimal Decoder	/15206	2	2	F.//	W,		
8260	Arithmetic Logic Unit							O
8261	Fast Carry Extender							W
8262	9-Bit Parity Generator and Checker		-	<u>-</u>	T.		F	W
8263	3-Input 4-Bit Digital Multiplexer	Large set	14		7 70			Q
8264	3-Input 4-Bit Digital Multiplexer		_		-			Q
8266	2-Input 4-Bit Digital Multiplexer	_	-	-	_		F	W
8267	2-Input 4-Bit Digital Multiplexer	_	_		_		F	W
8268	Gated Full Adder	_	-		_	_	F	Q
8269	4-Bit Comparator	-	l –				F	W
8270	4-Bit Shift Register			# ## ·	π		F	W
8271	4-Bit Shift Register				#1	#	F	W
8273	10-Bit Serial-In, Parallel-Out Shift Register			₩.	1 H		a E	W
8274	10-Bit Parallel-In, Serial-Out Shift Register		-	+		- 10 - 1	F	W
8275	Quad Bistable Latch					+ +	F	W
8276	8-Bit Serial Shift Register	-	-	-	_		F	
8277	Dual 8-Bit Shift Register	_	-		-	_	[F	
8280	Presettable Decade Counter	1 -	-	_	_	_	F	W
8281	Presettable Binary Counter	_		_	_	·	F	W
8284	Binary Up/Down Counter	_	-		_		F .	W
8285	Decade Up/Down Counter		187-83		-	- 11	F	W
8288	Divide-by-Twelve Counter		1147	ANH-E	11-1	\mathbb{N}/\mathbb{Z}	F	W
8290	Presettable High Speed Decade Counter		1 4	4 - 1				W
8291	Presettable High Speed Binary Counter			T 1 4 1	10.4	-	F	, in
8292	Presettable Low Power Decade Counter		- 1		/ /-	in the second	F	W
8293	Presettable Low Power Binary Counter	_	I -		_	_	F	W
9300	4-Bit Shift Register	/15901	*	*	F	W	F	W
9301	BCD to Decimal Decoder	/15206	2	2	F	W	F	W
9308	Dual 4-Bit Latch w/Clear				-		1	C
9309	Dual 4-Input Multiplexer	/01404		1	F	W	F	V
9310	4-Bit Decade Counter			4.4	1. 44	— —	F	V
9312	8-Input Digital Multiplexer	/01402			F.	W	Fil	W
9316	4-Bit Binary Counter						F	W
4-2-2-2-3	Data Selector-Multiplexer	15552003			1 4		F	W .
9322	5-Bit Comparator	/15002			F	WF	W	- 11
9324	8-Bit Addressable Latch	/16001	10 ka 259 490:	veltnagg- 195 —	p (50 km (54 m) 0 —	k) t = ty2#7kd; ————————————————————————————————————	F	٧
9334	Dual Monostable Multivibrator	710001	·	•	F	W	F	W
9602	Lual Monostable Muttivibrator	1	1		1 '		1	

NOTE

Per QFL 38510-28 dated 1 Apr. 1977 1 = Level 1 Qualification 2 = Level 2 Qualification

8T00 SERIES INTERFACE

The 8T00 Series provide a line of integrated circuits for applications such as line driving and receiving, level shifting and tri-state bus interfacing. Both Gold Doped and Schottky technologies are used to produce these devices.

		Г	
DEVICE	DESCRIPTION	PLASTIC	CERAMIC
8T04	7-Segment Decoder Display Driver (Active-Low Outputs)	N.	F
8T05	7-Segment Decoder Display Driver (Active-Hi Outputs)	N	F
8T06	7-Segment Decoder Display Driver (Active-Low Outputs)	N	F
8T09	Quad Bus Driver with Tri-State Outputs	N	F
8T10	Quad D-Type Bus Latch (Tri-State)	N	F
8T13	Dual Line Driver	N	F
8T14	Triple Line Receiver/Schmitt Trigger	N	F
8T15	Dual Communications EIA/MIL Line Driver	N	F
8T16	Dual Communications EIA/MIL Line Receiver	N	
8T18	Dual 2-Input NAND (High Voltage to TTL Interface)	N	F
8T20	Bidirectional Monostable Multivibrator (Diff. Input)	N	F
8T22	Retriggerable Monostable Multivibrator (74122/9601)	N	F
8T23	Dual Line Driver for IBM 360/370 Interface (75123)	N	- F
8T24	Triple Line Receiver for IBM 360/370 Interface (75124)	N	F
8T25	Dual MOS Sense Amplifier with Latch (Tri-State Outputs)	V	_
8T26A	Quad Bus Driver/Receiver (Tri-State Outputs)	N	i f
8T28	Quad Non-Inverting Bus Driver/Receiver (Tri-State Outputs)	N	F
8T30	Dual TTL/DTL to MOS Transceiver/Port Controller	N	F
8T31	8-Bit Bidirectional I/O Port	- N	F
8T32	Programmable 8-Bit 1/0 Port (3-State)	N	F
8T33	Programmable 8-Bit I/O Port (Open Collector)	- N.	F
8T34	Quad Bus Transceiver (DM8834)	N	F
8T35	Asynchronous Programmable 8-Bit I/O Port (Open Collector)	N	F
8T36	Asynchronous Programmable 8-Bit I/O Port (3-State)	N	F
8T37	Hex Bus Receiver with Hysteresis—Schmitt Trigger (DM8837)	N.	F
8T38	Quad Bus Transceiver (Open Collector) (DM8838)	N	F
8T39	Bus Extender/Repeater	N	F
8T40	Pipeline I/O Port	N N	F
8T58	Working Storage Bus Extender	. N.	F
8T80	Quad 2-Input NAND Gate (High Voltage)	N	F
8T90	Hex Inverter (High Voltage)	. N	F
8T95	High Speed Hex Buffers/Inverters (74365/DM8095)	N	F
8T96 8T97	High Speed Hex Buffers/Inverters (74366/DM8096)	N	F
8197 8T98	High Speed Hex Buffers/Inverters (74367/DM8097)	N	F
8198 8T363	High Speed Hex Buffers/Inverters (74368/DM8098)	N	F
8T380	Dual Zero Crossing Detector	N	F
01300	Quad Bus Receiver with Hysteresis	. N	HESH

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ANALOG (LINEAR)

CONSUMER/COMMUNICATION CIRCUITS		
NE/SE540	Power Driver	
NE541	Power Driver	
NE542	Dual Preamp	
NE543	Servo Amplifier	
NE544/644	Servo Amplifier	
NE544/644 NE546 NE570/571	AM Radio Analog Compandor	
N/S5596	Balanced Modulator/Demodulator	
μA758	Stereo Decoder	
ULN2111 ULN2208 ULN2209	FM Detector/Limiter FM Gain Block FM Gain Block	
ULN2211 TBA120S/U	TV Sound Channel TV Sound IF	
TBA1440	TV Video IF	
TBA327/395/396	TV PAL Chroma Set	
TCA440	AM Receiver	
TDA2541	TV Video IF	
CA3089	FM IF System	
PA239	Dual Preamp	
LM381, 381A	Dual Preamp	
LM382	Dual Preamp	
LM387	Dual Preamp	
MC1496/1596	Balanced Modulator/Demodulator	

Also see D-MOS FETs

OP AMPS		
NE /CES31	High Slaw Pata On Amn	
NE/SE531	High Slew Rate Op Amp	
*NE/SE/SA532	Dual Op Amp Dual Op Amp	
NE/SE532A *NE/SE535	• •	
	High Slew Rate Op Amp FET Input Op Amp	
NE/SU536 NE/SE538	High Slew Rate Op Amp	
NE/SE5534	Low Noise Op Amp	
*MC1456/1556	High Performance Op Amp	
*MC1458/1558/SA1458	Dual Op Amp	
*μA709/709C/SA709C	Op Amp	
μΑ740C	FET Input Op Amp	
*μΑ741/741C/SA741C	General Purpose Op Amp	
*μΑ747/747C/SA747C	Dual Op Amp	
*μA748/748C/SA748C	General Purpose Op Amp	
LF155/255/355	BIFET Input Op Amp	
LF156/256/356	BIFET Input Op Amp	
LF157/257/357	BIFET Input Op Amp	
*LM101/201	High Performance Op Amp	
*LM101A/201A/301A	High Performance Op Amp	
*LM107/207/307	General Purpose Op Amp	
*LM108/208/308	Precision Op Amp	
*LM108A/208A/308A	Precision Op Amp	
*LM124/224/324	Quad Op Amp	
LM124A/224A/324A	Quad Op Amp	
*LM158/258/358	Dual Op Amp	
SA1458	Dual Op Amp	
SA534	Quad Op Amp	
*SA709	Op Amp	
*SA741	General Purpose Op Amp	
*SA747	Dual Op Amp	

PERIPHERAL INTERFACE		
NE/SE501	Video Amp	
NE/SE592	Video Amp	
*μA733/733C	Video Amp	
55450B/51B/52B/	·	
53B/54B	Dual Peripheral Drivers	
75450B/51B/52B/		
53B/54B	Dual Peripheral Drivers	
DS3611/12/13/14	Dual Peripheral Drivers—8∪√ Output	
UDN5711/12/13/14	Dual Peripheral Drivers—80V Output	
MC1488	Quad Line Driver	
MC1489/1489A	Quad Line Receiver	
75\$107/108	Dual Line Receiver	
*DM7820/8820	Dual Differential Line Receiver	
*DM7820A/8820A	Dual Differential Line Receiver	
*DM7830/8830	Dual Differential Line Driver	

DISPLAY INTERFACE	
DM8880 DM8880-1 NE584 NE585 NE582	Display Decoder Driver Display Decoder Driver—100V Outputs Gas Discharge Segment Driver Gas Discharge Digit Driver LED Digit Driver

MEMORY INTERFACE		
3207A	MOS Clock Driver	
3207A-1	MOS Clock Driver	
7520	Dual Core Memory Sense Amp	
7522	Dual Core Memory Sense Amp	
7524	Dual Core Memory Sense Amp	
7528	Dual Core Memory Source Amp	
75232	Dual Core Memory Source Amp	
75234	Dual Core Memory Source Amp	
75S207/208	Dual MOS Memory Sense Amp	
75324	Core Memory Driver	
55/75325	Switched Paid Quad Core Driver	

MOSFET-ANALOG/DIGITAL SWITCHES (D-MOS)		
*SD210/211	Switch Driver	
*SD212/213	Switch	
*SD214/215	Switch	
*SD5000/5001/5002	Quad Switch Array IC	
*SD5100/5101	Quad Multiplexer IC	
*SD5200	Quad Switch Driver IC	
*SD5301	8x2 Crosspoint Switch	
SD5350	8 Channel Multiplexer w/Shift Register Control	
DMP/DMS4025	Power FET	

D/A/D CONVERTERS		
MC1408-7	8-Bit D/A Converter, 1 LSB Accuracy	
MC1408-8 MC1508-8	8-Bit D/A Converter, ½ LSB Accuracy 8-Bit D/A Converter, ½ LSB Accuracy	
NE5007	8-Bit D/A Converter, 1 LSB Accuracy 8-Bit D/A Converter, ½ LSB Accuracy	
NE/SE5008 NE/SE5009	8-Bit D/A Converter, ½ LSB Accuracy	

^{*}Available with Military processing (see Military section)

ANALOG (LINEAR) (Cont'd)

MOSFET-RF (D-MOS)				
SD200/201 SD202/203 SD300 SD303 SD304 SD305 SD306 SD6000 DMP/DMS4025	Single Gate UHF Single Gate UHF Dual Gate UHF Dual Gate UHF Dual Gate VHF Dual Gate VHF Dual Gate VHF Mixer Dual Gate VHF Amp VHF Mixer/Amp IC Power FET			

TIMERS				
NE/SE/SA558 NE/SE/SA559 *NE/SE/SA555 *NE/SE/SA556	Quad Timer Quad Timer Timer Dual Timer			

PHASE LOCKED LOOPS				
NE560 NE561 NE562 NE/SE564 NE/SE565 NE/SE566 *NE/SE567	Phase Locked Loop Phase Locked Loop Phase Locked Loop Phase Locked Loop Phase Locked Loop Function Generator Tone Decoder PLL			

HYBRIDS				
*LH2101AF *LH2108F *LH2108AF *LH2111AF	Dual Op Amp Dual Precision Op Amp Dual Precision Op Amp Dual Comparator			

TRANSISTOR ARRAYS					
ULN2001/2/3/4 N5501/2/3/4	Darlington Transistor Array High Voltage Darlington Transistor Array				

DIFFERENTIAL AMPLIFIERS				
*NE/SE510	Dual Differential Amp			
*NE/SE511	Dual Differential Amp			
*NE/SE515	Differential Amp			
*µA733	Video Amp			

VOLTAGE REGULATORS					
5V Voltage Regulator					
Precision Voltage Regulator					
Three-Terminal Positive					
Voltage Regulators					
Three-Terminal Positive					
Voltage Regulators					
High Input Breakdown (1 Amp)					
Positive Voltage Regulators					
Three-Terminal Positive					
Voltage Regulators					
High Input Breakdown (500mA)					
Positive Voltage Regulators					
Three Terminal Negative					
Voltage Regulator					
Three Terminal Negative					
Voltage Regulator (500mA)					
1 Amp Adjustable Positive					
Voltage Regulator					
500mA Adjustable Positive					
, , ,					
,					
Voltage Regulator					
1 5 5					
Dual Tracking Regulator					
	5V Voltage Regulator Precision Voltage Regulator Precision Voltage Regulator Three-Terminal Positive Voltage Regulators Three-Terminal Positive Voltage Regulators High Input Breakdown (1 Amp) Positive Voltage Regulators Three-Terminal Positive Voltage Regulators High Input Breakdown (500mA) Positive Voltage Regulators Three Terminal Negative Voltage Regulator Three Terminal Negative Voltage Regulator (500mA) 1 Amp Adjustable Positive Voltage Regulator				

COMPARATORS					
*NE521/522 *NE/SE526 *NE/SE527 *NE/SE529 *\(\mu\)A710/710C *\(\mu\)A711/711C *LM111/211/311 *LM119/219/319 *LM139/239/339 *LM139/239/339 LM2903 MC3302 SA539	High Speed Dual Differential Comparator Analog Voltage Comparator High Speed Voltage Comparator High Speed Voltage Comparator Voltage Comparator Dual Voltage Comparator Voltage Comparator Dual Voltage Comparator Quad Voltage Comparator Quad Voltage Comparator Dual Voltage Comparator Dual Voltage Comparator Dual Voltage Comparator Dual Voltage Comparator Quad Voltage Comparator Quad Voltage Comparator Quad Voltage Comparator				

^{*}Available with Military processing (see Military section).

MIUTARY

INTRODUCTION

Manufacturing integrated circuits to the stringent requirements of the Defense and Aerospace industries has been a dedication of Signetics since its founding in 1961. Today, Signetics remains unchanged in charter and continues to produce integrated circuits for this market segment with state-of-the-art process technologies which result in unequalled overall reliability.

Signetics' dedicated approach to the mili-

tary market is evidenced by the fact that it is the only major semiconductor manufacturer with a separate Military Products Division—with its own production facilities, engineering staff and marketing organization. Rather than dealing with each individual product division, as they must at many companies, Signetics' military customers can deal with one group of personnel familiar with the company's broad product capability as well as with the special requirements of the military market.

Assessing the individual products of any high-reliability supplier is enhanced by a complete understanding of the company's technological processes and the degree of its commitment to quality control and reliability.

PROCESSING LEVELS AND REQUIREMENTS

MILITARY PRODUCTS/ PROCESS LEVELS

The Signetics Mil 38510/883 Program is organized to provide a broad selection of processing options, structured around the most commonly requested customer flows. The program is designed to provide our customers:

- Standard processing flows to help minimize the need for custom specs
- Cost savings realized by using standard processing flows in lieu of custom flows.
- Better delivery lead times by minimizing spec negotiation time, plus allows customers to buy product off-the-shelf or in various stages of production rather than waiting for devices started specifically to custom specs.

The following explains the different processing options available to you. Special device marking clearly distinguishes the type of screening performed. Refer to Tables 2, 3 and 4.

JAN QUALIFIED (JB)

JAN Qualified product is designed to give you the optimum in quality and reliability. The JAN processing level is offered as the result of the government's product standardization programs, and is monitored by the Defense Electronic Supply Center (DESC), through the use of industry-wide procedures and specifications.

JAN Qualified products are manufactured, processed and tested in a government certified facility to Mil-M 38510, and appropriate device slash sheet specifications. Design documentation, lot sampling plans, electrical test data and qualification data for each specific part type has been approved by the Defense Electronic Supply Center (DESC) and products appear on the DESC Qualified Products List (OPL-38510).

Group B testing, per Mil-Std-883 Method 5005, is performed on each six weeks of production on each slash sheet for each package type. Group C, per Mil-Std-883 Method 5005, is performed every ninety days for each microcircuit group. Group D testing, per Mil-Std-883 Method 5005, is performed every six months for each package type.

In addition to the common specs used throughout the industry for processing and testing, JAN Qualified products also possess a requirement for a standard marking used throughout the IC industry.

JAN PROCESSING (JBX)

This option is extremely useful when the reliability and screening of a JAN device is required, however, Signetics is not listed on the QPL for the product needed. Processing is performed to Mil-Std-883 Method 5004, and product is 100% electrically tested to the appropriate JAN slash sheet.

Group B, C and D data for JAN processed and the other military processing levels

JAN CASE OUTLINE	SIGNETICS MILITARY PACKAGE TYPES					
AND				•		
LEAD FINISH	8-Pin	10-Pin	14-Pin	16-Pin	24-Pin	
СВ			F		_	
EB	_	_	_	F	l —	
JB	-	_		_	1/F*	
DB	_		W	_	<u> </u>	
FB	-		_	W		
ZC	_	_	-	-	Q	
GC	T	_	_	-	-	
IC .	-	K	_	ļ. —		

*The gold plated versions of these packages will be available for a limited time. All products listed in Military section are also available in Die form.

Table 1 MILITARY PACKAGE AVAILABILITY

	JB	JB JBX		RB	s.	
	JAN Qualified	JAN Processed	JAN Rei	/883	Mil Temp	
54/54H	X	×	Х	X	X	
54LS	×	l x	X	X	Х	
548	X	l x	Х	Х	X	
82/8T	X	X	X	Х	X	
93XX	X	l x	Х	Х	Х	
96XX	_		X	Х	X	
Linear	Planned	l x	Х	Х	Х	
Bipolar Memory	Planned		X	X	Х	
Microprocessor	l	_	Х	X	х	

Table 2 MILITARY SUMMARY

which follow, consist of Group B, C and D testing performed per Mil-Std-883 Method 5005, in accordance with the Signetics Military Data Program. Offshore assembly is allowed.

JAN REL (RBX)

Processing to this option is ideal when no JAN slash sheets are released on devices required. Product is processed to Mil-Std-883 Method 5004, and is 100% electrically tested to industry data sheets.

/883B (RB)

This is a lower priced version of the JAN Rel option described above. Processing is identical with the only exceptions being the dc electrical testing over the temperature range and ac electrical testing at room temperature are performed as a part of Group A instead of 100%.

MIL TEMP/883C (S/RC)

If you need a Military temp. range device, but do not require all the high reliability screening performed in the other processing options, our Mil-Temp. product is ideal. Mil-Temp. parts are the standard full Mil-Temperature range product guaranteed to a 1% AQL to the Signetics data sheet parameters.

MILITARY GENERIC DATA

Signetics has a new program for those customers who require quality conformance data on their products. This program

allows our customers to obtain reliability information without the necessity of running Groups B, C and D inspections for their particular purchase order. It provides for the customer something that has not been readily available before in the semiconductor industry in that all Military Generic Data is controlled and audited by both Government Inspection in the case of JAN data and Signetics Quality Assurance.

Signetics Military Generic Data is compiled by the Military Products Division based on data from 1) JAN quality conformance lots, and 2) Data generated by quality conformance lots run for other reliability programs. Refer to Table 4.

A Military Generic family is defined as consisting of die function and package type families.

Military Generic Data

- Allows our customers to qualify Signetics products based on existing quality conformance data performed at Signetics.
- Allows our customers to reduce costs and improve deliveries.
- Provides assurance that all Signetics die function families and packages meet
- Mil-M-38510 and customer reliability requirements.

Provides an attributes summary to the customer backed by lot identity and traceability.

PROCESS LEVEL AND MARKING	PRE-CAP VISUAL	BURN IN	FUNCTIONAL TEST	DC/AC @25°C	DC/AC @TEMP	QPL	OFF SHORE
JB JM38510/XXXXX	883B	Yes	100%	100%	100%	Yes	No
JBX M38510/XXXXX	883B	Yes	100%	100%	100%	No	Yes
RBX M38510/P/N	883B	Yes	100%	100%	100%	No	Yes
RB SXXXX/883B	883B	Yes	100%	100% dc Sample ac	Sample dc only	No	Yes
RC/S SXXXX/883C	883B	No	100%	100% dc Sample ac	Sample dc only	No	Yes

Table 3 MILITARY PRODUCTS PROCESSING MATRIX

QUALIFIED SUB-GROUPS	QUALIFIES	OPTION 1	OPTION 2
A*	Electrical Test		
В	Package—Same package construction and lead finish.	Data selected from devices manufactured within 6 weeks of the manufacturing period on the same production line through final seal.	Data selected from devices manufactured within 24 weeks of manufacturing period.
С	Die/Process—Devices representing the same process families.	Data selected from representa- tive devices from the same mi- crocircuit group and sealed within 12 weeks of the man- ufacturing period.	Data selected from the representative devices from the same microcircuit group and sealed within 48 weeks of the manufacturing period.
D	Package—Same package construction and lead finish.	Data selected from the devices representing the same package construction and lead finish manufactured within the 24 weeks of manufacturing period. If specific data not available, Option 2 will be supplied.	Data selected from the devices representing the same package construction and lead finish manufactured within the 52 weeks of manufacturing period.

NOTE*

Group A is performed on each lot or sublot of Signetics devices.

Table 4 DEFINITION AND QUALIFYING MANUFACTURING PERIODS FOR GENERIC DATA

					PROCE	SSING LE	VELS	
DESCRIPTION OF REQUIREMENTS AND SCREENS	MIL-M-38510 AND MIL-STD- 883 REQUIREMENTS, METH- ODS AND TEST CONDITIONS	REQUIRE- MENT	CLASS A	JAN Qualified (JB)	JAN Proc- essed (JBX)	JAN Rel (RBX)	/883B (RB)	/883C (RC)
General Mil-M-38510 1. Pre-Certification A.Product Assur- ance Program Plan B. Manufacturer's Certification	The manufacturer shall establish and implement a Products Assurance Program Plan and provide for a manufacturer survey by the qualifying activity, Para. 3.4.1.1	_	X	X	N/A	N/A	N/A	N/A
2. Certification	Received after manufacturer has completed a successful survey, Para. 3.4.1.2		Х	X	N/A	N/A	N/A	N/A
3. Device Qualification	Device qualification shall consist of subjecting the desired device to groups A, B, C & D of method 5005 to tightened LTPD, Para. 3.4.1.2	_	x	X	N/A	N/A	N/A	N/A
4. Traceability	Traceability maintained back to a production lot Para. 3.4.6	_	X	x	Х	х	×	X
5. Country of Origin	Devices must be manufactured, assembled, and tested within the U.S. or its territories, Para. 3.2.1	. —	х	X	N/A	N/A	N/A	N/A
Screening Per Method 5004 of Mil-Std-883			,					
6. Internal Visual (Pre-	2010, Cond. A or B	100%	XA	XB	ХВ	XB	ХВ	ХВ
7. Stabilization Bake	1008, Cond. C Min; (24 Hrs @ 150°C)	100%	x	x	X	×	х	x
8. Temperature Cycling* *For Class B and C devices thermal shock may be substituted, 1011, Cond. A; (15 cycles, 0° to +100°C)	1010, Cond. C; (10 cycles, -65°C to +150°C)	100%	X	X	X	X	х	X
9. Constant Accelera-	2001, Cond. E; (30 _{kg} in YI Plane)	100%	X :	, X	, X	х	х	Х
10. Visual Inspection	There is no test method for this screen; it is intended only for the removal of "Catastrophic Failures" defined as "Missing Leads, Broken Packages or Lids Off."	100%	, X	X	x	×	X	X
11. Seal (Hermeticity) A. Fine B. Gross	1014 Cond. A or B (5.0 X 10-8CC/Sec) Cond. C2 Min.	100% 100%	X X	x x	X	X X	X X	X X
12. Interim Electricals (Pre Burn-In)	Per applicable device specification	100% Optional	100% Read &	Slash Sheet	Slash Sheet	Data Sheet	Data Sheet	N/A
13. Burn-In	1015, Cond. as specified (160 hrs. Min. at 125°C)	100%	Record 100%	x	x	×	×	N/A
14. Final Electricals	Per applicable Device Specification	100%	240 hrs. 100% Read &	Slash Sheet	Slash Sheet	Data Sheet	Data Sheet	Data Sheet
A. Static Tests @ 25°C B. Static Tests	Sub Group 1	1	Record X	×	x	x	х	×
@ +125°C C. Static Tests	Sub Group 2		х	x	х	x	N/A	N/A
@ -55°C	Sub Group 3		х	x	Х	х	N/A	N/A

Table 5 REQUIREMENTS AND SCREENING FLOWS FOR STANDARD CLASS B PRODUCTS

			.,		PROC	ESSING LE	VELS	
DESCRIPTION OF REQUIREMENTS AND SCREENS	MIL-M-38510 AND MIL-STD- 883 REQUIREMENTS, METH- ODS AND TEST CONDITIONS	REQUIRE- MENT	CLASS A	JAN Qualified (JB)	JAN Proc- essed (JBX)	JAN Rel (RBX)	/883B (RB)	/883C (RC)
D. Dynamic Test @ 25° C	Sub Group 4 (for Linear Product Mainly)		х	х	х	Х	х	х
E. Functional Test @ 25°C F. Switching	Sub Group 7		×	х	х	×	х	х
Test @ 25°C	Sub Group 9		Х	×	X	х	N/A	N/A
15. Percent Defective allowable (PDA) 16. Marking	A PDA of 10% is a normal requirement applied against the static tests @ 25°C (A-1). This is controlled by the slash sheets for JB & JBX products. For RBX & RB 10% is standard. Fungus Inhibiting Paint	10% 100%	5% As Reg'd	X JM38510/	X M38510/	. X M38510/	X SXXXX/	N/A SXXXX/
, and the second	rangao minaning rami			XXXX Slash Sheet #	XXXX Slash Sheet #	SXXXX Sig. Basic #	883B Sig. Basic #	883C Sig. Basic #
17. X-Ray	2012		100%	N/A	N/A	N/A	N/A	N/A
18. External Visual	2009	100%	x	х	×	×	X	Х
Quality Conformance Inspection per Method 5005 of Mil-Std-883						-		
19. Group A	Electrical Tests-Final Electri- cals (#14 above) repeated on a sample basis. (Sub Groups 1 thru 12 as specified.)	Each Lot	х	х	Х	Х	х	Х
20. Group B	Package functional and constructional related test I.E. package dimensions, resistance to solvents, internal visual & mechanical, bond strength & solderability.	Every 6 week per microcircuit group	х	х	Gener	ic Data Ava	ilable	
21. Group C	Die related tests I.E. 1,000 hr. op- erating life, temperature cycling, & constant acceleration.	Every 3 months per package type	х	x	Gener	ic Data Ava	ilable	
22. Group D	Package related tests I.E. physical dimensions, lead fatigue, thermal shock, temperature cycle, moisture resistance, mechanical shock, vibration variable frequency constant acceleration, & salt atmosphere.	Every 6 months per package type	×	x	Genei	ric Data Ava	ilable	

Table 5 REQUIREMENTS AND SCREENING FLOWS FOR STANDARD CLASS B PRODUCTS (Cont'd)

LOGIC

LOGIC—54LS SERIES

			JA	AL	JA PRO ESS	oc-	M REL	/883
DEVICE	DESCRIPTION	JM38510 SLASH SHEET	DIP	FLATPACK	DIP	FLATPACK	DIP	FLATPACK
54LS00 54LS01	Quad 2-Input NAND Gate Quad 2-Input NAND Gate	/30001	2	2	F F	w w	F F	W W
54LS02 54LS03	with o/c Quad 2-Input NOR Gate Quad 2-Input NAND Gate with o/c	/30301 /30002	2	2	F F	w w	F	W W
54LS04 54LS05 54LS08 54LS09 54LS10 54LS11	Hex Inverter Hex Inverter with o/c Quad 2-Input AND Gate Quad 2-Input AND Gate with o/c Triple 3-Input NAND Gate Triple 3-Input NAND Gate	/30003 /30004 /31004 /30005 /31001	1 1 2 - 1 2	1 12 - 12	F F F F F F	>>>	т. Т. т. т. т. т.	W W W W
54LS12 54LS13	Triple 3-Input NAND Gate with o/c Dual NAND Schmitt Trigger	/30006	1	1	F	W	F ·	W
54LS14 54LS15	Hex Schmitt Trigger Triple 3-Input AND Gate with o/c	/31302 /31002	2	2	F	W	F F	W
54LS20 54LS21 54LS22 54LS26 54LS27 54LS28	Dual 4-Input NAND Gate Dual 4-Input AND Gate Dual 4-Input NAND Gate with o/c Quad 2-Input NAND Gate with o/c Triple 3-Input NOR Gate Quad 2-Input NOR Buffer	/30007 /31003 /30008 /32101 /30302 /30204	1 2 1 * 2 .	1 2 1	T T T T T T T T T T T T T T T T T T T	W W W W	T T T T T T T	W W W W W
54LS30 54LS32 54LS33	8-Input NAND Gate Quad 2-Input OR Gate Quad 2-Input NOR Buffer with o/c	/30009 /30501 —	2	2 2 —	F F	W W —	FF	W W W
54LS37 54LS38	Quad 2-Input NAND Buffer Quad 2-Input NAND Buffer with o/c	/30202 /30203	2	2	F	W W	F	W W
54LS40 54LS42 54LS51 54LS54 54LS55 54LS73	Dual 4-Input NAND Buffer BCD-to-Decimal Decoder Dual 2-Wide 2-Input A01 Gate 4-Wide 2-Input A01 Gate 2-Wide 4-Input A01 Gate Dual J-K Master-Slave Flip-Flop	/30201 /30703 /03401 /30402 — /30101	2 2 -	2 2 2 —	F.F.	₩ ₩ ₩ —	T T T T T T	₩ ₩ ₩ ₩ ₩
54LS74	Dual D-Type Edge-Triggered Flip-Flop	/30102	•	*	F	W	F	W
54LS75 54LS76	Quad Bistable Latch Dual J-K Master-Slave Flip-Flop	/30110	*	*	F	W	F	W
54LS78 54LS83A 54LS85 54LS86 54LS90 54LS92 54LS93	Quad Bistable Latch 4-Bit Binary Full Adder 4-Bit Magnitude Comparator Quad 2-Input Exclusive-OR Gate Decade Counter Divide-by-Twelve Counter 4-Bit Binary Counter	/31201 /31101 /30502 /31501 /31510 /31502		_		W W W W W		W W W W W
54LS95 54LS96	4-Bit Left-Right Shift Register 5-Bit Shift Register	/30603		*	F	w	F	w

NOTE

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Per QPL 38510-28 dated
1 April 1977.
1 = Level 1 Qualification
2 = Level 2 Qualification
• = In Process

LOGIC-54LS SERIES (Cont'd)

			JA	AL	PR	AN OC- SED		IIL _/883
DEVICE	DESCRIPTION	JM38510 SLASH SHEET	dia	FLATPACK	DIP	FLATPACK	DIP	FLATPACK
54LS107	Dual J-K Master-Slave Flip-Flop	/30108	*	*	F	w	F	w
54LS109	Dual J-K Positive Edge- Triggered Flip-Flop	/30109	*	*	F	W	F	W
54LS112	Dual J-K Negative Edge- Triggered Flip-Flop	/30103	*	*	F	W	F	W
54LS113	Dual J-K Negative Edge- Triggered Flip-Flop	/30104	*	*	F	W	F	W
54LS114	Dual J-K Negative Edge- Triggered Flip-Flop	/30105	*	*	F	W	·F	W
54LS122	Retriggerable Monostable Multivibrator	/31403	F 149					
54LS125	Quad Bus Buffer Gate w/Tri-State Outputs	/32301					F	W
54LS126	Quad Bus Buffer Gate w/Tri-State Outputs	/32302	* : 40		*		F	W
54LS132 54LS136	Quad Schmitt Trigger Quad Exclusive-OR with o/c	/31303 —	*	***	F	W	F	W
54LS138	3-to-8 Line Decoder/Demux	/30701	*	*	*	*	F	W
54LS139	Dual 2-to-4 Line Decoder/ Demux	/30702	*	*	*	*	F	W
54LS145	BCD to Decimal Decoder/Dye		_	-	_		F	W
54LS151	8-Line to 1-Line Mux	/30901			F	W	F	W
54LS153	Dual 4-Line to 1-Line Mux	/30902			<u>ا</u> ا	VV		
54LS154	4-Line to 16-Line Decoder/ Demux		_	: بست آرائ				Q
54LS155	Dual 2-Line to 4-Line Decoder/Demux						F	W
54LS157	Quad 2-Input Data Selector (non-inv.)	/30903	9		F	W	F	W
54LS158	Quad 2-Input Data Selector (inv.)	/30904			F	W	F.	W
54LS160	Synchronous 4-Bit Decade Counter	/31503					F	W
54LS161	Synchronous 4-Bit Binary Counter	/31504	*	*	F	W	F	W
54LS162	Synchronous 4-Bit Decade Counter	/31511	*	*	*	*	F	W
54LS163	Synchronous 4-Bit Binary Counter	/31512	*	*	*	*	F	W
54LS164	8-Bit Parallel-Out Serial Shift Register	/30605	*	*	F	W	F	W
54LS170	4X4 Register File		—	_	-		F	W
54LS173	Quad D-Type Flip-Flop		120	18.12		1 4 5 T	F	W
E 41 C 4 3 4	(Tri-State) (8T10)	100100		eafeko) Lawaa				1
54LS174	Hex D-Type Flip-Flop with Clear	/30106			F	W	F	W
54LS175	Quad D-Type Edge-Triggered	/30107			F	W	F	W
54LS181	4-Bit Arithmetic Logic Unit	/30801	2				1	Q
54LS190	Synchronous Up/Down Counter (BCD)	/31513			Manager Comments	W	F	w
54LS191	Synchronous Up/Down Counter (Binary)	/31509	*	*	F	W	F	W

NOTE
Per QPL 38510-28 dated 1 April 1977.
1 = Level 1 Qualification
2 = Level 2 Qualification
* = In Process

LOGIC-54LS SERIES (Cont'd)

			j . J/	AN	PR	OC-	M	IL
			QL	JAL		SED	REL	/883
DEVICE	DESCRIPTION	JM38510 SLASH SHEET	DIP	FLATPACK	DIP	FLATPACK	DIP	FLATPACK
54LS192	Synchronous Decade Up/Down Counter	/31507	*	*	F	W	F	W
54LS193	Synchronous 4-Bit Binary Up/Down Counter	/31508	*	*	F	W	F	W
54LS194	4-Bit Bidirectional Universal Shift Register	/30601	*	*	*	*	*	*
54LS195	4-Bit Parallel-Access Shift Register	/30602	*	*	*	*	*	*
54LS196	Presettable Decade Counter/ Latch (8290)	/31601	*	* .	*	*	*	*
54LS197	Presettable Binary Counter/ Latch (8291)	/31602		*	*	*	*	
54LS221	Dual Monostable Multivibrator	/31402	*		*			
54LS251	Data Selector/Mux with 3-State Outputs	/30905			*	*		
54LS253	Dual 4-Line to 1-Line Data Selector/Mux	/30908			F	W	F	W
54LS257	Quad 2-Line to 1-Line Data Selector/Mux	/30906			F	W	F	W
54LS258	Quad 2-Line to 1-Line Data Selector/Mux	/30907	*	*	*	*	F	W
54LS260	Dual 5-Input NOR Gate	_	-	_	_	_	F	W
54LS261	2X4 Parallel Binary Multiplier	_	—		-	_	F	W
54LS266	Quad Exclusive-NOR Gate	/30303	2	2	F	W	F	W
54LS279	Quad S-R Latch					enematiche ve	F	W
54LS280	9-Bit Odd/Even Parity Generator/Checker		, -		*	_		
54LS283	4-Bit Adder	/31202					F	W
54LS290	Decade Counter	/32003			F	W	F	W
54LS293	4-Bit Binary Counter	/32004			F	W	F	W
54LS295A	4-Bit Right-Shift Left-Shift	/30606			4/4		F	W
	Register	March 1	100		11			
54LS298	Quad 2-Input Mux with Storage		l –		-	_	F	W
54LS365	Hex Buffer w/Common Enable (3-State)	/32201	*	•			F	W
54LS366	Hex Buffer w/Common Enable (3-State)	/32202	*				F	W
54LS367	Hex Buffer, 4-Bit and 2-Bit (3-State)	/32203					F	W
54LS368	Hex Buffer, 4-Bit and 2-Bit (3-State)	/32204	*	*		· ·	F	W
54LS375	Quad Latch			777	177		F	W
54LS386 54LS395	Exclusive-OR Gate 4-Bit Cascadeable Shift Register (3-State)	 /30607			7		F	W W
54LS445 54LS670	BCD to Decimal Decoder/Dye 4X4 Register File (Tri-State)						F	W W

NOTE

Per QPL 38510-28 dated 1 April 1977.

1 = Level 1 Qualification

2 = Level 2 Qualification

* = In Process

LOGIC—54S SERIES

			į.	AN JAL	PR	AN OC- SED	REL	
DEVICE	DESCRIPTION	JM38510 SLASH SHEET	OIP	FLATPACK	OID	FLATPACK	DIP	FLATPACK
54S00	Quad 2-Input NAND Gate	/07001	1	1	F	w	F	W
54S02	Quad 2-Input NOR Gate	/07301	2	2	F	W	F	W
54S03	Quad 2-Input NAND Gate with o/c	/07002	2	2	F	W	F	W
54S04	Hex Inverter	/07003	1	1	F	W	F	W
54S05	Hex Inverter with o/c	/07004	1	1	F	W	F	W
54S08	Quad 2-Input AND Gate	/08003	3.		F	W	F	W
54809	Quad 2-Input AND Gate with o/c	/08004		-			F	W
54S10	Triple 3-Input NAND Gate	/07005	2	2	F	W	F	W
54S11	Triple 3-Input NAND Gate	/08001	2	2	F	W	F	W
54S15	Triple 3-Input AND Gate with o/c	/08002	2	2	"F	W	F	W
54S20	Dual 4-Input NAND Gate	/07006	2	2	F	W	F	W
54\$22	Dual 4-Input NAND Gate with o/c	/07007	1	1	F	W	F	W
54S30	8-Input NAND Gate	/07008	_			_		_
54S32	Quad 2-Input OR Gate	_	_		_	_	F	W
54\$40	Dual 4-Input NAND Buffer	/07201	2	2	F	W	F	W
54851	Dual 2-Wide 2-Input	/07401	2	- 2	F.	w.	F	W
54864	A01 Gate 4-2-3-2 Input A01 Gate	/07402	2	2	F	W	F	w
54S65	4-2-3-2 Input A01 Gate	/07403	2	2	F	W	F	W
54874	Dual D-Type Edge-Triggered Flip-Flop	/07101	2	2	F	W	F	W
54S85	4-Bit Magnitude Comparator	/08201			F	No.	F	
54S86	Quad 2-Input Exclusive-OR Gate	/07501	2	2	F	w	F.	W
548112	Dual J-K Negative Edge- Triggered Flip-Flop	/07102	_	_	-	_	F	W
54S113	Dual J-K Negative Edge- Triggered Flip-Flop	/07103	_	_	_	_	F	W
54\$114	Dual J-K Negative Edge- Triggered Flip-Flop	/07104	_	_	_	_	F	W
54S133	13-Input NAND Gate	/07009	2	2	F	W	F	W
548134	12-Input NAND Gate w/Tri- State Outputs	/07010	2	2	F	W	F	W
54S135	Quad Exclusive-OR/NOR Gate	/07502		2			27	
54S138	3-to-8 Line Decoder/Demux	/07701				74 <u></u> -		
548139	Dual 2-to-4 Line Decoder/	/07702	-		# - -2.		F	W
54S140	Dual 4-Input NAND Line Driver	/08101	2	2		W	F	W
54S151	8-Line to 1-Line Mux		1	_		. 19.50		W
54S153	Dual 4-Line to 1-Line Mux	/07901 /07902	2	2	F F	W	F	W
54S157	Quad 2-Input Data Selector (non.inv.)	/07903	2	2	F	W	F	W
54S158	Quad 2-Input Data Selector	/07904	*	*	F	W	F	W
54S174	Hex D-Type Flip-Flop with Clear	/07106	_		_	_	*	*
548175	Quad D-Type Edge-Triggered Flip-Flop	/07105	1 <u>-</u>				19	
54\$181	4-Bit Arithmetic Logic Unit	/07801	*	· <u> </u>	4	السا	D.	7.0
545182	Look-Ahead Carry Generator	/07802		البلق	المقاوة	أشياه	3.7	
548194	4-Bit Bidirectional Universal Shift Register	/07601						
54\$195	4-Bit Parallel-Access Shift Register	/07602	_				-	

NOTE

Per QPL 38510-28 dated 1 April 1977.
1 = Level 1 Qualification
2 = Level 2 Qualification
• = In Process

LOGIC-54S SERIES (Cont'd)

				QUAL E		AN OC- SED		
DEVICE	DESCRIPTION	JM38510 SLASH SHEET	DIP	FLATPACK	DIP	FLATPACK	DIP	FLATPACK
54S251	Data Selector/Mux with 3-State	/08905	-	_	-	_	I	_
54S253	Outputs Dual 4-Line to 1-Line Data Selector/Mux	_	-	_	-	_	F	W
54S257	Quad 2-Line to 1-Line Data Selector/Mux	/07906	-	_	-		-	_
54S258	Quad 2-Lint to 1-Line Data Selector/Mux	/07907	_	_	_	-	-	_
54S260	Dual 5-Input NOR Gate	_	l _	_	_	_	F	W
54S280	9-Bit Odd/Even Parity Generator/Checker	/07703					-	
54\$350	4/6 Bit Shifter-Tri-State		-		-		F	

NOTE

Per QPL 38510-28 dated 1 Apr. 1977

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LOGIC-8200/9300/9600 SERIES

		JM38510		AN .ified*	JA PROCE		RE	MIL L/883 TEMP
DEVICE	DESCRIPTION	SLASH SHEET	Dip	Flat Pack	Dip	Flat Pack	Dip	Flat Pack
8200	Dual 5-Bit Buffer Register						1	Q
8201	Dual 5-Bit Buffer Register with D Inputs			_	_	_	ı	Q
8202	10-Bit Buffer Register	_	_	_	_	_	1	Q
8203	10-Bit Buffer Register with D Inputs		_		_	_	I	Q
8230	8-Input Digital Multiplexer	/01402	*	*	F	W	F	W
8231	8-Input Digital Multiplexer		1137 118		H_{i}		77 HE	W
8232	8-Input Digital Multiplexer		197				///⊑//	W
8233	2-Input 4-Bit Digital Multiplexer			HIT III		H. Tali	Ē	W
8234	2-Input 4-Bit Digital Multiplexer				A_{ij}	1/2	F	W
8235	2-Input 4-Bit Digital Multiplexer Quad Exclusive-OR Gate				1477		2	Ŵ
8241 8242	Quad Exclusive-OH Gate Quad Exclusive-NOR Gate		_		_	_	F	W
8243	8-Bit Position Scaler	_	_		_		i i	Q
8250	Binary-to-Octal Decoder	/15204	2	2	F	w	F	w
8251	BCD-to-Decimal Decoder	/15205	2	2	F	W	F	W
8252	BCD-to-Decimal Decoder	/15206	2	2	667 F 766	W	F	W
8260	Arithmetic Logic Unit			4				Q
8261	Fast Carry Extender		4.1		1.4		F	W
8262	9-Bit Parity Generator and Checker			######################################		<u>-</u>	∐ F ≥ 1	W
8263	3-Input 4-Bit Digital Multiplexer		0.4	#	11 11 11			\mathbf{q}
8264	3-Input 4-Bit Digital Multiplexer	- 1. WUSAFEND ** I - 4					1	Q
8266	2-Input 4-Bit Digital Multiplexer				_	_	F	W
8267	2-Input 4-Bit Digital Multiplexer		-	_		_	F	W
8268	Gated Full Adder		_			_	F	Q
8269	4-Bit Comparator					— Meg a. gro	F	W
8270	4-Bit Shift Register			1			Ē	W
8271	4-Bit Shiff Register	$\Pi \Pi \Pi \Pi \Pi$	The second		$H_{i}(\overline{D})$		F	W
8273 8274	10-Bit Serial-In, Parallel-Out Shift Register 10-Bit Parallel-In, Serial-Out Shift Register	$\bar{\mathcal{I}}$	WE :					W
8275	Quad Bistable Latch			ΠU	_ _		F	W
8276	8-Bit Serial Shift Register					######################################	F	, (# () () () () () () () () () () () () ()
8277	Dual 8-Bit Shift Register	_				_	F	_
8280	Presettable Decade Counter		_	_	_	_	F	W
8281	Presettable Binary Counter	_			l —		F	W
8284	Binary Up/Down Counter		_	_		_	F	W
8285	Decade Up/Down Counter				0.0 1. 01	Diffic		W
8288	Divide-by-Twelve Counter			HH H	- il		F	W
8290	Presettable High Speed Decade Counter		1-11		-	<u> </u>	#F	W
8291	Presettable High Speed Binary Counter	1 - 1 - 1 - 1 - 1 - 1 - 1 - 1 - 1 - 1 -			77. T		F	W
8292	Presettable Low Power Decade Counter		pi , t jih	HHI	-	#####################################	F ::	W.
8293	Presettable Low Power Binary Counter		-		_	141	F	W
9300	4-Bit Shift Register	/15901		_	F	W	F	W
9301	BCD to Decimal Decoder	/15206	2	2	F	W	F	W
9308	Dual 4-Bit Latch w/Clear		<u> </u>	_	_			Q
9309	Dual 4-Input Multiplexer	/01404	 	 	F	W	F	W
9310	4-Bit Decade Counter			170		w		W
9312 9316	8-Input Digital Multiplexer 4-Bit Binary Counter	/01402			F	vv	F	W
9310	Data Selector-Multiplexer				\mathbb{Z}	## 1 5#		W
9322	5-Bit Comparator	/15002	7		F	WF	w	Д. <u>Ж</u> .
9334	8-Bit Addressable Latch	/16001	-073637333		**************************************	*********	F	**************************************

NOTE

Per QPL 38510-28 dated 1 Apr. 1977

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^{* =} In Process

LOGIC-8T INTERFACE SERIES

		JAN M38510		EL/883 TEMP
DEVICE	DESCRIPTION	SHEET	Dip	Flat Pack
8T04	7-Segment Decoder Display Driver (Active-Low Outputs)			l w
8T05	7-Segment Decoder Display Driver (Active-Low Outputs) 7-Segment Decoder Display Driver (Active-Hi Outputs)	_	[w
8T06	7-Segment Decoder Display Driver (Active-In Outputs) 7-Segment Decoder Display Driver (Active-Low Outputs)		[w
8T09	Quad Bus Driver with Tri-State Outputs			l w
8T10	Quad D-Type Bus Latch (Tri-State)		E	l w
8T13	Dual Line Driver			W
8T14	Triple Line Receiver/Schmitt Trigger		e l	W
8T18	Dual 2-Input NAND (High Voltage to TTL Interface)	$\mu_{ij}(i)$	e la	W
8T20	Bidirectional Monostable Multivibrator (Diff. Input)	经租赁		
8T22	Retriggerable Monostable Multivibrator (54122/9601)		F I	w
8T26A	Quad Bus Driver/Receiver (Tri-State Outputs)		F	l w
8T28	Quad Non-Inverting Bus Driver/Receiver (Tri-State Outputs)	_	F	w
8T31	8-Bit Bidirectional I/O Port		*	*
8T32	Programmable 8-Bit, I/O Port (3-State)		1	*
8T33	Programmable 8-Bit, I/O Port (Open Collector)	_	ı	*
8T35	Asynchronous Programmable 8-Bit I/O Port (Open Collector)	en en u nione		(24)
8T37	Hex Bus Réceiver with Hysteresis—Schmitt Trigger (DM8837)		∦F	W
8T38	Quad Bus Transceiver (Open Collector) (DM8838)		F	W
8T80	Quad 2-Input NAND Gate (High Voltage)		F/	W
8T90	Hex Inverter (High Voltage)		(f) (F)	W
8T95	High Speed Hex Buffers/Inverters (74365/DM8095)		F	W
8T96	High Speed Hex Buffers/Inverters (74366/DM8096)	_	F	W
8T97	High Speed Hex Buffers/Inverters (74367/DM8097)		F	W
8T98	High Speed Hex Buffers/Inverters (74368/DM8098)		F	W

^{* =} Qualification planned

LOGIC—5400 SERIES

				AN AL*	PRO ESS		ı	IL /883
DEVICE	DESCRIPTION	JM38510 SLASH SHEET	DIP	FLATPACK	DIP	FLATPACK	DIP	FLATPACK
5400 5401	Quad 2-Input NAND Gate Quad 2-Input NAND Gate with o/c	/00104 /00107	1	1	F	W	F	W
5402 5403	Quad 2-Input NOR Gate Quad 2-Input NAND Gate with o/c	/00101 /00109	1	1	F	w *	F	W —
5404	Hex Inverter	/00105	1	1	F	W	F	W
5405 5406	Hex Inverter with o/c Hex Inverter w/Buffer/Driver with o/c	/00108 /00801	1	1	F	W	F	W
5407	Hex Buffer/Driver with o/c	/00803	المكرا	1	F	W	F	W
5408 5409	Quad 2-Input AND Gate Quad 2-Input AND Gate	/01601 /01602	1	1	F	W W	F	W W
5410	with o/c Triple 3-Input NAND Gate	/00103	1	1	F	W	F	W
5411 5412	Triple 3-Input NAND Gate Triple 3-Input NAND Gate	/00103 — /00106	_ _	<u>-</u>	_	_ _	F	W W
5440	wih o/c	45464					_	
5413	Dual NAND Schmitt Trigger	/15101	*	*	F	W	F	W
5414 5416	Hex Schmitt Trigger Hex Inverter Buffer/Driver	/15102 /00802		r <u>.</u> rear	F	W	F	W
	with o/c	700602		15			48	
5417	Hex Buffer/Driver with o/c	/00804	3	الشيار المسائر	F	W	F	W
5420	Dual 4-Input NAND Gate	/00102	4		F	W	F	W
5421	Dual 4-Input AND Gate						F	W
5426	Quad 2-Input NAND Gate	/00805	11		F	وكنيت	F	
	with o/c				_			TH.
5427	Triple 3-Input NOR Gate	/00404	ļ -	•	F	W	F	W
5428 5430	Quad 2-Input NOR Buffer 8-Input NAND Gate	/16201 /00101	1	1	F	W	F	W
5432	Quad 2-Input OR Gate	/16101		*	[_	~~	F	W
5433	Quad 2-Input NOR Buffer with o/c	_		_	-	_	F	W
5437 5438	Quad 2-Input NAND Buffer Quad 2-Input NAND Buffer with o/c	/00302 /00303	1	1	F	W	F	W
5439	Quad 2-Input NAND Buffer		42		1	1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1	F	W
5440	Dual 4-Input NAND Buffer	/00301	1	1	F	W	F	W
5442	BCD-to-Decimal Decoder	/01001	, 1	1 1.	F	W	(F)	W
5443 5444	Excess 3-to-Decimal Decoder Excess 3-Gray-to-Decimal Decoder	/01002 /01003	1	1	F	W	F	W
5445	BCD-to-Decimal Decoder/Driver with o/c	/01004	-	_	F	W	F	W
5446A	BCD-to-7 Segment Decoder/ Driver	/01006	_	_	F	W	F	W
5447A	BCD-to-7 Segment Decoder/ Driver	/01007	-		F	W	F	W
5448	BCD-to-7 Segment Decoder/ Driver	/01008			F	W	F	W
5450	Driver Expandable Dual 2-Wide 2- Input A01	/00501	1	1	F	W	F	W
5451 5453	Dual 2-Wide 2-Input A01 Gate 4-Wide 2-Input A01 Gate	/00502 /00503		1	F	W W	F	W W
5454 5455	(Expandable) 4-Wide 2-Input A01 Gate 2-Wide 4-Input A01 Gate	/00504 /04005	1	1	F	W	F	w

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LOGIC-5400 SERIES (Cont'd)

			l l	AN	PR	AN OC-	I	IIL
			Ql	JAL*	ES	SED.	REL	./88 ¥
		JM38510		۸C		AC		0.40
DEVICE	DESCRIPTION	SLASH SHEET	OIP	FLATPAC	alo	FLATPAC	PIP	EI ATDA
5460 5470	Dual 4-Input Expander		_	_	T	_	F	W
5470	J-K Flip-Flop J-K Master-Slave Flip-Flop	/00206 /00201	1	1 1	F	W	F	W
5473	Dual J-K Master-Slave Flip-Flop	/00207	1	1	F	W	F	W
5474	Dual D-Type Edge-Triggered Flip-Flop	/00205	1	1	F	W	F	w
5475	Quad Bistable Latch	/01501	1		F	W	F	W
5476 5477	Dual J-K Master-Slave Flip-Flop Quad Bistable Latch	/00204 /01502	1	1	F	W W	F	W
5480	Gated Full Adder	701002			I_{\perp}	44	F	W
5483	4-Bit Binary Full Adder	/00602	1	47	F	W	F	W
5485	4-Bit Magnitude Comparator	/15001	1	1	F	W	F	W
5486	Quad 2-Input Exclusive-OR Gate	/00701	1	1	F	W	F	W
5490 5491	Decade Counter	/01307	*	*	F	W	F	W
5491	8-Bit Shift Register Divide-by-Twelve Counter	/01301	1	1	F	w	F	W
5493	4-Bit Binary Counter	/01301	1	1 1/197	F	W.	F	W
5494	4-Bit Shift Register (PISO)		1974				F	W
5495	4-Bit Left-Right Shift Register	/00901	1	#	F	<u>.44</u>	F	W
5496	5-Bit Shift Register	/00902	1	1	F	W	F	W
54100 54107	4-Bit Bistable Latch (Dual)	10 3 17 111	47	+		4	F	W
54107	Dual J-K Master-Slave Flip-Flop	/00203	1		F		F	
54109	Dual J-K Positive Edge-		_		_	_	F	w
	Triggered Flip-Flop							••
54116	Dual 4-Bit Latch with Clear	/01503	2		. 1		1	
54121	Monostable Multivibrator	/01201	1	1	F	W	F	W
54122	Retriggerable Monostable Multivibrator	/01202	_	_	—	_	—	_
54123	Retriggerable Monostable	/01203	1	1	F	W	F	W
	Multivibrator	701200						
54125	Quad Bus Buffer Gate	/15301	2	2	F	W	F	W
	w/Tri-State Outputs							
54126	Quad Bus Buffer Gate	/15302	2	2	F	W	F	W
54128	w/Tri-State Outputs Quad 2-Input NOR Buffer							
54132	Quad Schmitt Trigger	/15103			F	8	F F	W
54145	BCD-to-Decimal Decoder/Driver	/01005	47.28L	6321XI	F	W	<i>⊪</i> 633.	W
	with o/c					•		•
54147	10-Line to 4-Line Priority	/15601	*	*	F	W	F	W
54148	Encoder	/15000			_		_	14/
UT 170	8-Line to 3-Line Priority Encoder	/15602	•	^	F	W	F	W
54150	16-Line to 1-Line Mux	/01401	2		1	_	1	
54151	8-Line to 1-Line Mux	/01406	2	2	F	w	F	w
54152	8-Line to 1-Line Mux	Muthiall	-	+	#		F	W
54153	Dual 4-Line to 1-Line Mux	/01403	2	2	F	W	F	W
54154	4-Line to 16-Line Decoder/ Demux	/15201		W.	H_{i}			Q
54155	Dual 2-Line to 4-Line	/15202	2	2	F	W	F	W
	Decoder/Demux							
54156	Dual 2-Line to 4-Line	/15203	2	2	F	W	F	W
54157	Decoder/Demux Quad 2-Input Data Selector	/01405	1	1	F	w	F	W
54450	(non-inv.)						_	
54158	Quad 2-Input Data Selector (inv.)	_	_	_	-	-	F	W
- 4400		/01202	4	4	F	w	F	W
54160	Synchronous 4-Bit Decade	/01303	1	1	_	VV I		

NOTE

Per QPL 38510-28 dated 1 Apr. 1977

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LOGIC-5400 SERIES (Cont'd)

				N IAL	PR	AN OC- SED		IIL ./883
DEVICE	DESCRIPTION	JM38510 SLASH SHEET	DIP	FLATPACK	DIP	FLATPACK	OIP	FLATPACK
54161	Synchronous 4-Bit Binary Counter	/01306	1	1	F	W	F	W
54162	Synchronous 4-Bit Decade Counter	/01305	1	1	F	W	F	W
54163	Synchronous 4-Bit Binary Counter	/01304	1	1	F	W	F	W
54164	8-Bit Parallel-Out Serial Shift Register	/00903	1	-	F	-	F	_
54165	Parallel-Load 8-Bit Shift Register	/00904	*	_	F —	W	F	W
54166	8-Bit Shift Register	l marine	_		2 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1	2 <u></u>	F	
54170	4X4 Register File	/01801			-	9 30		085
54174	Hex D-Type Flip-Flop with Clear	/01701	1	1	F	W	F	W
54175	Quad D-Type Edge-Triggered Flip-Flop	/01702	1	1	F	W,	F	- W
54180	8-Bit Odd/Even Parity Checker	/01901		2	F	W	F	W
54181	4-Bit Arithmetic Logic Unit	/01101	1		1		1	
54182	Look-Ahead Carry Generator	/01102	1	1	F	W	F	W
54190	Synchronous Up/Down Counter (BCD)	_	-	_	-	_	*	*
54191	Synchronous Up/Down Counter (Binary)		_		-		*	*
54192	Synchronous Decade Up/Down Counter	/01308	*	*	F	W	F	W
54193	Synchronous 4-Bit Binary Up/Down Counter	/01309	*		F	W	F	W
54194	4-Bit Bidirectional Universal Shift Register	/00905	*	***	F	W	F	W
54195	4-Bit Parallel-Access Shift Register	/00906	*	*	F	W	F	W
54198	8-Bit Shift Register						程	Q
54199	8-Bit Shift Register			-	 _	- 15 ₂ 2 	1-	
54221	Dual Monostable Multivibrator	_	_		_		F	W
54279	Quad S-R Latch	_	l —	_	 _	_	F	W
54298	Quad 2-Input Mux with Storage	-	_	_	 _		F	W
54365	Hex Buffer w/Common Enable (3-State)	/16301	*	*	*	*	*	*
54366	Hex Buffer w/Common Enable (3-State)	/16302	*	*	*	*	F	W
54367	Hex Buffer, 4-Bit and 2-Bit (3-State)	/16303	*		*		F	W
54368	Hex Buffer, 4-Bit and 2-Bit (3-State)	/16304					F	W

NOTE

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LOGIC—54H SERIES

			1	N.	PR	AN OC-		IIL (000
			QU	JAL ¥	ES	SED ¥	KEL	./883 ⊻
DEVICE	DESCRIPTION	JM38510 SLASH SHEET	DIP	FLATPACI	DIP	FLATPACK	PIO	FLATPACK
54H00	Quad 2-Input NAND Gate	/02304	1	1	F	W	F	W
54H01	Quad 2-Input NAND Gate with o/c	/02306	1	1	F	W	F	W
54H04	Hex Inverter	/02305	1	1	F	W	F	W
54H05	Hex Inverter with o/c			_	_	_	F	W
54H08	Quad 2-Input AND Gate	/15501	2		F		F	W
54H10 54H11	Triple 3-Input NAND Gate Triple 3-Input NAND Gate	/02303	1 2	1	E	W	F	W
54H20	Dual 4-Input NAND Gate	/15502 /02302	1	7	F F	w	F	W
54H21	Dual 4-Input AND Gate	/15503	2	16/168	F	y y	F	W
54H22	Dual 4-Input NAND Gate	/02307	1		F	w	3F3	W
9	with o/c		e de la constante de la consta		300			
54H30	8-Input NAND Gate	/02301	1	1	F	W	F	W
54H40	Dual 4-Input NAND Buffer	/02401	1	1	F	W	F	W
54H50	Expandable Dual 2-Wide 2-Input A01	/04001	1	1	F	W	F	W
54H51	Dual 2-Wide 2-Input A01 Gate	/04002	1	1	F	W	F	W
54H52	Expandable 4-Wide 2-2-2-3		_		-	_	F	W
om sanaanse v	Input AND-OR Gate	A Paris and a Secretary of the						
54H53	4-Wide 2-Input A01 Gate	/04003	1	1	F	W	F	W
54H54	(Expandable) 4-Wide 2-Input A01 Gate	/04004	1	4				
54H55	2-Wide 2-Input A01 Gate	/04004	1	1	F	W	F	W
54H60	Dual 4-Input Expander	704000	100	Ŀ	1120	VV	F	W
54H61	Triple 3-Input Expander	<u> 1</u>					100	W
54H62	3-2-2-3 Input AND-OR		_	_			F	W
	Expander							•••
54H71	J-K Master-Slave Flip-Flop	_		_	—	_	F	W
	with AND-OR Inputs							
54H72	J-K Master-Slave Flip-Flop	/02201	1	1	F	W	F	W
54H73	Dual J-K Master-Slave	/02202	1	1	F	W	F	W
54H74	Flip-Flop Dual D-Type Edge-Triggered	/02202	4		F	w	F	w
341174	Flip-Flop	/02203	1	1	۲	VV	-	VV
54H76	Dual J-K Master-Slave	/02204	11	1	F	W	F	w
121	Flip-Flop					90 (de 190 (de 190 (de	Ш	
54H101	J-K Negative Edge-Triggered	/02205	1	1	F	W	F	W
	Flip-Flop	Filliation						
54H102	J-K Negative Edge-Triggered Flip-Flop					10	F	W
54H103	Dual J-K Negative Edge-	/02206	1.		F	W	F	W
	Triggered Flip-Flop						100	
54H106	Dual J-K Negative Edge-		12.	4.	$\frac{1}{2}$	# 1	F	W
5411400	Triggered Flip-Flop	1 日本東京 1 日本マンスト	Policy Property Policy			43		
54H108	Dual J-K Negative Edge-	-	_	-	_	-	F	-
	Triggered Flip-Flop							

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BIPOLAR MEMORY

BIPOLAR MEMORY

DEVICE	ORGANIZATION	PAC	KAGE*	OUTPUT	NUMBER OF PINS
PROMs					
82S23	32X8	F	R	oc	10
82S115	512X8	1 1	R	TS	16 24
82\$123	32X8	F	R	TS	16
82S126	256X4	F	R	oc	16
82S129	256X4	F	R	TS	16
82S130	512X4	III EU	TO BEST	OC	16
82S131	512X4	l F	p	78	16
82S136	1024X4	Fil	R	oc	18
82S137	1024X4	E,E	R	TS	18
825140	512X8	1	A	OC.	24
82S141	512X8	ı	R	TS	24
82S180	1024X8	1	R	- oc	24
82S181	1024X8	1	R	TS	24
82S184	2048X4		R	OC	18
82S185	2048X4		R	TS	18
82S2708	1024X8	1.1	R	TS	24
FPLAs					
82\$100	16X48X8		Я	TS	
825101	16X48X8		B	00 00	28
	10/40/0				28
PLAs					
82\$200	16X48X8	1	R	TS	28
82S201	16X48X8		R'	oc ha	28
RAMs					
3101A	16X4	F	R	oc	16
54\$89	16X4	F	R	oc	16
54S189	16X4	F	R	TS	16
54S200	256X1	F	R	TS	16
54S201	256X1	F F	R	TS	16
54\$301	256X1	F	R	HI OC	16
82809	64X9	1	A A	TS	28
82S10	1024X1	F.I	R	OC	16
82S11	1024X1	F;1:::1	R	TS	1.111116
82S16	256X1	F	R	TS	16
82817	256X1	F	R	oc	16
82S25	16X4	F	R	OC	16
ROMs					
828215	512X8	1	R	TS	24
828223	32X8	F	R	oc	16
82S224	32X8	F	R	TS	16
828226	256X4	F	R	00	16
82S229	256X4	p.	R	TS	16
82S230	512X4	F	R	oc	16
82S231	512X4	F	R	TSL	16
82S280	1024X8		R	oc	24
82S281	1024X8	I	R	TS	24
NOTE					

*NOTE

R = BeO Flat Pack F = Cerdip I = Ceramic DIP

BIPOLAR MEMORY CROSS REFERENCE

AMD	SIGNETICS
2700/27LS00	82S16
2701/27LS01	82S17
27S08/27LS08	82S23
27S09/27LS09	82S123
27S10	82S126
27S11	82S129
3101	82S25
3101A/27S02	3101A
93415	82S10
93415A	82S10
93425	82S11
93425A	82S11

FAIRCHILD	SIGNETICS
93403	82S25
93406	82S226
93411	82S17
93415	82S10
93415A	82S10
93417	82S126
93419	82809
93421	82S16
93425A	82S11
93427	82S129
93431	82S130
93436	82S130
93438	82S140
93441	82S131
93442	82S240
93446	82S131
93448	82S141
93452	82S 136
93453	82S137
93454	82S280
93457	82S126
93464	82S281
93467	82S129

HARRIS	SIGNETICS		
0064	82S25		
HM7602-2	82S23		
HM7603-2	82S123		
HM7608-2	82S2708		
HM7610-2	82S 126		
HM7611-2	82S129		
HM7620-2	82S 130		
HM7621-2	82S131		
HM7640-2	82S140		
HM7641-2	82S141		
HM7642-2	82S136		
HM7643-2	82S137		
HM7647-2	82S115		
HM7680-2	82S180		
HM7681-2	82S181		
HM7684-2	82S184		
HM7685-2	82S185		

INTEL	SIGNETICS
2708	82S2708
2716	82S2716
3101	82S25
3101A	3101A
3106/A	82S16
3107/A	82S17
3301A	82S226
3302	825230
3322	82S231
3601	82S126
3602	82S130
3604	82S140
3605	82S 136
3608	82S180
3622	82S131
3624	82S141
3625	82S137
3628	82S181

INTERSIL	SIGNETICS
5501	82S25
5508	82S10
5508A	82S10
5518	82S11
5518A	82S11
5523	74S201
5523A	82S16
5533	74S301
5533A	82S17
5600	82S23
5603A	82S126
5604	82S130
5605	82S140
5610	82S123
5623A	82S129
5624	82S131
5625	82S141
56506	82S136
56526	82S137

MOTOROLA	SIGNETICS
4004A	82S226
4064	82S25
4256	82S16
5005	82S 126

MMI	SIGNETICS
5200	82S226
5201	82S229
5205	82S230
5206	82S231
5300-1	82S126
5301-1	82S129
5305-1	82S130
5306-1	82S131
5330	82S23
5331	82S123

MMI	SIGNETICS
5340-1	82S140
5341	82S141
5348	82S146
5349	82S147
5352	82S136
5353	82S137
5380	82S180
5381	82S181
5385	82S2708
5530	82S17
5531	82S16
5560	82S25/3101A

NATIONAL	SIGNETICS
54187	82S226
54S188	82S23
54S287	82S129
54S288	82S123
54S387	82S126
54S570	82S130
54S571	82S131
54S572	82S136
54S573	82S137
8582	82S17
86L99	82S25
87S295	82S140
87S296	82S141

T.I.	SIGNETICS
54187	82S226
54\$188	82S23
54\$189	74S189
54S200	74S200
54S201	74S201
54S209	82S11
54S270	82S230
54S287	82S129
54S288	82S123
54S289	3101A
54S301	74S301
54S309	82S10
54S370	82S231
54S387	82S126
54S472	82S147
54S473	82S146
54S474	82S141
54S475	82S140
54S476	82S137
54S477	82S136

NOTE

Parts are pin for pin functional replacements except where noted.

BIPOLAR MICROPROCESSOR

BIPOLAR MICROPROCESSORS

PRODUCT	DESCRIPTION	AVA	AVAILABILITY	
		Dip	Flat Pack	
3001	Microprogram Control Unit	ı	R	
3002	Central Processing Element (2-bit slice)	1 1	R	
8X300	Interpreter/Microcontroller	1	*	
2901-1	Central Processing Element (4-bit slice)	*	*	

^{*}Under development

MICROPROCESSOR SUPPORT CIRCUITS

		AVAILABILITY	
PRODUCT	DESCRIPTION	Dip	Flatpack
LOGIC			
54123	Retriggerable Monostable Multivibrator	F	w
54180	8-Bit Odd/Even Parity Checker	F	w
54298	Quad 2-Input Mux with Storage	F	w
54S182	Look-Ahead Carry Generator	*	*
54S194	4-Bit Bidirectional Shift Register	*	*
54S195	4-Bit Parallel Access Shift Register	*	*
54LS365	High Speed Hex Tri-State Buffer	F	*
54LS366	High Speed Hex Tri-State Buffer	F	*
54LS367	High Speed Hex Tri-State Buffer	F	*
54LS368	High Speed Hex Tri-State Buffer	F	*
8262	9-Bit Parity Generator Checker	F	w
8281	Presettable Binary Counter	F	w
8291	Presettable High Speed Binary Counter	F	w
9602	Dual Monostable Multivibrator	F	w
INTERFACE		i i	•••
8T09	Quad Bus Driver with Tri-State Output	F	w
8T10	Quad D-Type Bus Latch (Tri-State Outputs)	F	w
8T13	Dual Line Driver	F	w
8T14	Triple Line Receiver/Schmitt Trigger	F	w
8T26A	Quad Bus Driver/Receiver (Tri-State)	F	w
8T28	Quad Bus Non-Inverting Driver/Receiver (Tri-State)	F	w
8T32	Programmable 8-Bit I/O Port (3-State)	i	*
8T33	Programmable 8-Bit I/O Port (Open Collector)	i	*
8T35	Asynchronous Programmable 8-Bit I/O Port (Open Collector)	i	*
8T95	High Speed Hex Buffer (Tri-State)	F	*
8T96	High Speed Hex Inverter (Tri-State)	F	*
8T97	High Speed Hex Buffer (Tri-State)	F	*
8T98	High Speed Hex Inverter (Tri-State)	F	*

^{*}Under development

ANALOG

LINEAR INDUSTRY CROSS REFERENCE

FAIRCHILD	SIGNETICS
μΑ111	LM111
μΑ139	LM139
μ A710	μA710
μΑ711	μA711
μ Α733	μ Α733
μ AF155/156	LF155/156
μ A101	LM101
μΑ101Α	LM101A
μΑ107	LM107
μΑ108	LM108
μ A108A	LM108A
MC1556	MC1556
μ Α1558	MC1558
μ Α709	μΑ709
μ A709A	μ Α709Α
μ Α741	μ Α741
μ Α747	μ Α747
μ Α748	μ Α748
MC1555	SE555
μ Α556	SE556
μ Α109	LM109
μ A79XX	79XX (7)
μ Α723	μΑ723

MOTOROLA	SIGNETICS
MLM111	LM111
MC1710	μ A710
MC1711	μ Α711
MC1733	μ Α733
LF155/56	LF155/156
MLM101	LM101
MLM101A	LM101A
MLM107	LM107
MC1558	MC1558
MC1709	μ A709
MC1709P2	μΑ709Α
MC1741	μ Α741
MC1747	μ Α747
MC1748	μΑ748
MC3556	SE556
MLM109	LM109
MC78XX	78XX (7)
MC79XX	79XX (7)
MC1723	μΑ723
MC1508	MC1508-8

ICL	
NATIONAL	SIGNETICS
LM161	SE527
LH2111	LH2111
LM111	LM111
LM119	LM119
LM139	LM139
LM193	LM193/193A
LM710	μA710
LM711	μΑ711
LM733	μ Α733
LF155/56	LM155/156
LH2101A	LH2101A
LH2108A	LH2108A
LM101A	LM101
LM101	LM101A
LM107	LM107
LM108	LM108
LM108A	LM108A
LM124	LM124
LM158	LM158
LM1558	MC1558
LM1581	SE532
LM709	μΑ709
LM741	μΑ741
LM747	μ Α747
LM748	μ Α748
LM567	SE567
DM7820	DM7820
DM7830	DM7830
LM555	SE555
LM109	LM109
LM723	μΑ723

PMI	SIGNETICS
SSS1508	MC1508-8
DAC-08	SE5008

RAYTHEON	SIGNETICS
LM111	LM111
LM139	LM139
RM710	μ A710
RM711	μ A711
RM733	μ Α733
LF155/56	LF155/156
LM101	LM101
LM101A	LM101A
LM107	LM107
LM108	LM108
LM108A	LM108A
LM124	LM124
RM1556	MC1556
RM1558	MC1558
RM709	μ A709
RM741	μ Α741
RM747	μ Α747
RM555	SE555
LM109	LM109
RM723	μ Α723

T.I.	SIGNETICS
LM111	LM111
SN52710	μA710
SN52711	μA711
SN52733	μ Α733
LF155/56	LF155/156
SN52101A	LM101A
SN52107	LM107
SN52709	μΑ709
SN52741	μ Α741
SN52748	μ Α748
SN55182	DM7820
SN55183	DM7830
SN52555	SE55
SE556	SE556
LM109	LM109
μ A79XX	μ A79XX (7)
SN52723	μ Α723

LINEAR PRODUCTS

DEVICE	DESCRIPTION	PACKAGE*	
	COMPARATORS		
SE521	Dual Comparator	F	
SE526	Analog Voltage Comparator	F	K
SE527	Analog Voltage Comparator	F	K
SE529	Analog Voltage Comparator	F	K
LH2111	Dual Comparator	F	
LM111	Comparator	F	Т
LM119	Dual Comparator	F	K
LM139	Quad Comparator	F	
LM193/193A	Dual Comparator		Т
μA710	Differential Voltage Comparator	F	Т
μA711	Comparator	F	K
	DIFFERENTIAL AMPLIFIERS		
SE510	Dual Differential Amplifier	F	
SE511	Dual Differential Amplifier	F	
SE515	Differential Amplifier	F	К
μA733	Video Amplifier	Ė	ĸ
μΑτοσ	'	'	1
	OPERATIONAL AMPLIFIERS	l _	
LF155/156	FET Op Amp	Т	
LH2101A	Dual Op Amp	F	
LH2108A	Dual Op Amp	F.	
LM101	High Perf. Op Amp	F	Т
LM101A	High Perf. Op Amp	F	Т
LM107	General Purpose Op Amp	F	F
LM108	Precision Op Amp	F	Т
LM108A	Precision Op Amp	F	Т
LM124	Quad Op Amp	F	
LM158	Dual Op Amp		Т
MC1556	Op Amp	F	Т
MC1558	Dual Op Amp	F	Т
SE532	Dual Op Amp	-	Т
SE535	Hi Slew Rate Op Amp	T	
SE538	Hi Slew Rate Op Amp	Ι Τ	
μΑ709	Op Amp	F	T
μA709A	Op Amp	F	T
μΑ741	General Purpose Op Amp	F	T
μA747	Dual Op Amp	F	K
μA748	General Purpose Op Amp	F	Т

DEVICE	DESCRIPTION	PACK	AGE
SE567	PHASE LOCKED LOOPS Tone Decoder P11	F	Т
DM7820 DM7830	LINE RECEIVERS Dual Differential Line Receiver Dual Differential Line Receiver	F	
SE555 SE556 SE558/9	TIMERS Timer Dual Timer Quad Timer	F F	Т
LM109 SE5554 78XX (7) 79XX (7) 79MXX (7) µA723	VOLTAGE REGULATORS 5 Volt Regulator Dual Track Reg Positive Reg Negative Reg Med Power Reg Precision Voltage Regulator	DA F DA DA DB F	L
DS1611-1614	DRIVERS Peripheral Drivers		т
MC1508-8 SE5008 SE5009	D/A 8-Bit D/A 8-Bit D/A 8-Bit D/A	F F F	

*NOTE
F = Cerdip
K/T/L = Metal can
DA/DB = TO-3 can
Flat pack available—special request

CHIP PROGRAM

signetics

CHIP PROGRAM

Signetics recognizes the industry need for integrated circuit die requirements. All standard military and commercial products can be supplied in chip form.

The matrix below outlines process options available for all chip products. Refer to Signetics Die Program catalog for additional information regarding: testing, lot qualification, product pin-outs and ordering information.

FLOW DESIGNATOR	VIS LEVEL	GLASS	GOLD BACKING
СВ	Н	Yes	. No
CD	X	Yes	No
CE	Н	Yes	Yes
CF	X	Yes	Yes
CG	S	No	Yes
CH	X	No	Yes
CJ	Н	No	Yes
CK	S	Yes	No
CN	s	Yes	Yes
CP	X (SEM)	Yes	No
CW	н	No	No
CX	s	No	No
CY	x	No	No

Table 6 PROCESS MATRIX

NOTES

- 1. Visual levels:
 - X = Mil std 883 2010.2 cond. A
 - H = Mil std 883 2010.2 cond. B
- S = Signetics standard criteria
- 2. Temperature options:
 - 0°C to +70°C -55°C to +125°C

signetics

INTRODUCTION

The following information applies to all packages unless otherwise specified on individual package outline drawings.

General

- Dimensions shown are metric units (millimeters), except those in parentheses which are English units (inches).
- 2. Lead spacing shall be measured within this zone.
 - a. Shoulder and lead tip dimensions are to centerline of leads.
- 3. Tolerances non-cumulative
- 4. Thermal resistance values are determined by utilizing the linear temperature dependence of the forward voltage drop across the substrate diode in a digital device to monitor the junction temperature rise during known power application across VCC and ground. The values are based upon 120 mils square die for plastic packages and a 90 mils square die in the smallest available cavity for hermetic packages. All units were solder mounted to P.C. boards, with standard stand-off, for measurement.

Plastic Only

- Lead material: Alloy 42 or equivalent, solder dipped.
- 6. Body material: Plastic
- 7. Round hole in top corner denotes lead No. 1.
- 8. Body dimensions do not include molding flash.

Hermetic Only

- 9. Lead material
 - a Alloy 52—gold plated, or solder dipped.
 - b.ASTM alloy F-15 (KOVAR) or equivalent—gold plated, tin plated, or solder dipped.
 - c. ASTM alloy F-30 (Alloy 42) or equivalent—tin plated.
 - d.ASTM alloy F-15 (KOVAR) or equivalent gold plated.
- e. ASTM alloy F-15 (KOVAR) or equivalent tin plated.
- 10. Body Material
 - a. 1010 Steel—nickel plated or tin plate over nickel.
 - b. Eyelet, ASTM alloy F-15 or equivalent—gold or tin plated.
 - c. Eyelet, ASTM alloy F-15 or equivalent—gold or tin plated, glass body.
 - d.Ceramic with glass seal at leads.
 - e. BeO ceramic with glass seal at leads.
 - f. Ceramic with ASTM alloy F-15 or equivalent.
- 11. Lid Material
 - a.1010 steel, nickel plated, or tin-plate over nickel, weld seal.
 - b.Nickel or tin plated nickel, weld seal.
 - c. Ceramic, glass seal.
 - d.ASTM alloy F-15 or equivalent, gold plated.
 - e. BeO Ceramic with glass seal.
 - f. Translucent A12O3, glass seal.

PLASTIC PACKAGES

NO. OF LEADS	PACKAGE CODE	θ _{ja} /θ _{jc} (^o C/w)	DESCRIPTION ¹	PAGE
Standard D	ual-in-Line			4
8	NE	162/65		3
14	NH	150/65	TO-116/MO-001	3
16	NJ	137/53	MO-001	3
18	NK	135/53		3
20	NL	135/53		3
22	NM	120/53		3
24	NN	116/53	MO-015	4
28	NQ	116/53	MO-015	4
40	NW ³	110/50	MO-015	4
Power Dua	l-in-Line			
14	NHA ²	95/33	Butterfly	3
16	NJA^2	95/33	Butterfly	3
18	NKA ^{2,3}	90/26	Butterfly	3
20	NLA ^{2,3}	90/26	Butterfly	3
24	NNA ²	60/23	Butterfly	4
28	NQA ²	56/21	Butterfly	4
Power				
3	S	200/70	TO-92	5
3	U	75/3	TO-220	5
3 + GND	GB ³	95/15	Single-in-Line (SIL)	5
4 + GND	GC ³	95/15	Single-in-Line (SIL)	5
12 + GND	PH/PHA3	95/15	Batwing	5

- 12. Signetics symbol, angle cut, or lead tab denotes Lead No. 1.
- 13. Recommended minimum offset before lead bend.
- 14. Maximum glass climb .010 inches.
- 15. Maximum glass climb or lid skew is .010 inches.
- 16. Typical four places.
- 17. Dimension also applies to seating plane.

HERMETIC PACKAGES

NO. OF LEADS	PACKAGE CODE	θ _{ja} / θ _{jc (} °C/W)	DESCRIPTION ¹	PAGE
Metal Head	lers			_
2	DA	TBD	TO-3 Solid Header	6
3	DB	TBD	TO-39 Solid Header, Short Can	6
4	DC	TBD	TO-72 Solid Header	6
4	DE	TBD	TO-72 Glass Filled Header	6
8	T	150/25	TO-99 Header (.200 Dia.)	7
10	K	150/25	TO-100 Header, Short Can	7
10	L	150/25	TO-100 Header, Tall Can	7
Flat Packs				
10	WF	240/50	Flat Ceramic	8
14	WH	205/50	Flat Ceramic	8
16	WJ	200/50	Flat Ceramic	8
24	WN	155/40	Flat Ceramic	8
16	RJ/RJA	133/30	Flat Ceramic, BeO	8
18	RKA ³	TBD	Flat Ceramic, BeO	_
24	RNA	TBD	Flat Ceramic, BeO	8
28	RQA	TBD	Flat Ceramic, BeO	9
40	RWA	TBD	Flat Ceramic, BeO	9
10	QF	230/55	Flat Ceramic	9
14	QH	185/45	Flat Ceramic	9
16	QJ	170/45	Flat Ceramic	9
24	QN	155/44	Flat Ceramic	9
10	QFA	230/55	Flat Ceramic Laminate	10
14	QHA	185/45	Flat Ceramic Laminate	10
16	QJA	170/45	Flat Ceramic Laminate	10
24	QNA	155/44	Flat Ceramic Laminate	10
Cerdip Far	nily			
14	FH	110/30	Dual in-Line Ceramic	11
16	FJ	100/30	Dual-in-Line Ceramic	11
18	FK	93/27	Dual-in-Line Ceramic	11
22	FM	75/27	Dual-in-Line Ceramic	11
24	FN	60/26	Dual-in-Line Ceramic	11
Laminated	Ceramic, Side	Brazed Lead		
8	IEA	100/30	Dip Laminate	12
14	IHA	95/25	Dip Laminate	12
16	IJA	90/25	Dip Laminate	12
18	IKA	88/25	Dip Laminate	12
22	IMA	80/25	Dip Laminate	12
24	INC/INE	65/25	Dip Laminate	12
28	IQA	60/25	Dip Laminate	13
40	IWA	55/25	Dip Laminate	13
50	IZA	TBD	Dip Laminate	13

NOTES

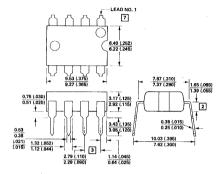
Dual-in-Line packages unless otherwise described

3. Package not yet available, scheduled for 1977 release

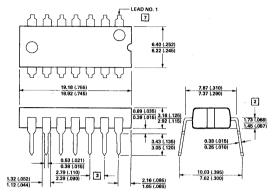
Package outline is the same as corresponding standard Dual-in-Line package with identical number of leads

PLASTIC: Standard and Power Dual-In-Line

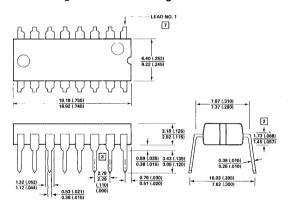
NE Package



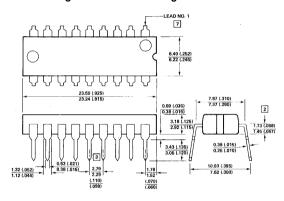
NH Package and NHA Package



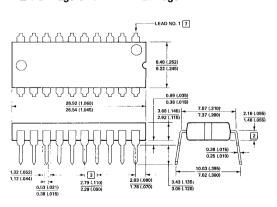
NJ Package and NJA Package



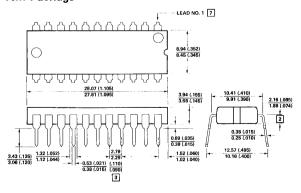
NK Package and NKA Package



NL Package and NLA Package

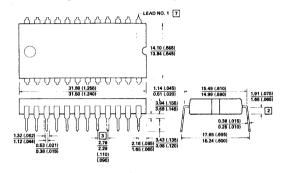


NM Package

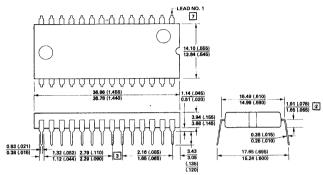


PLASTIC: Standard and Power Dual-In-Line (cont'd.)

NN Package and NNA Package



NQ Package and NQA Package

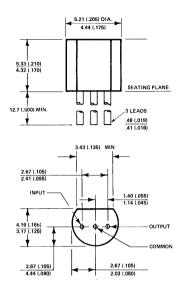


NW Package

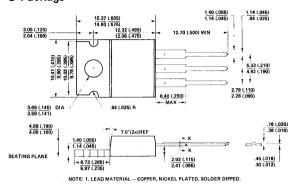
Package not yet available Scheduled for 1977 release

PLASTIC: Power (Not Dual-In-Line)

S Package



U Package



GB Package

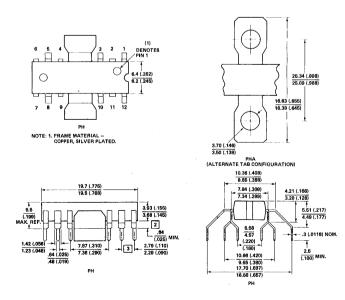
Package not vet available

GC Package

Scheduled for 1977 release

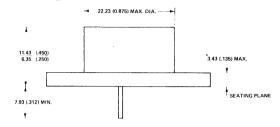
Package not yet available Scheduled for 1977 release

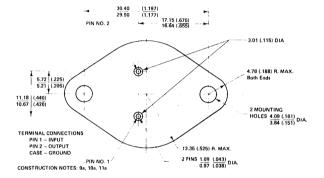
PH/PHA Package



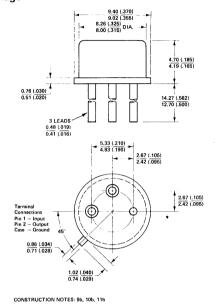
HERMETIC: Metal Headers

DA Package

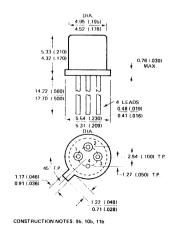




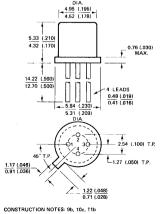
DB Package



DC Package



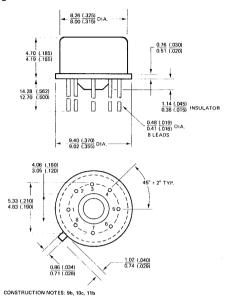
DE Package



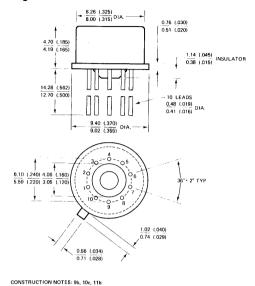
CONSTRUCT

HERMETIC: Metal Headers (cont'd.)

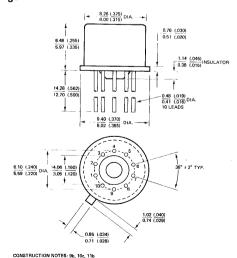
T Package



K Package

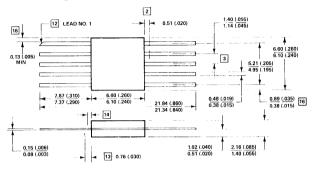


L Package



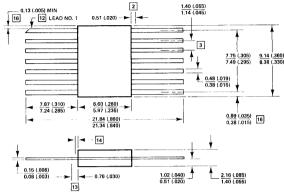
HERMETIC: Flat Packs

WF Package



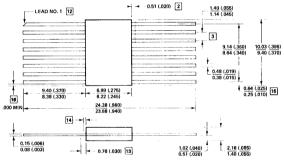
CONSTRUCTION NOTES: 9c, 10d, 11c

WH Package



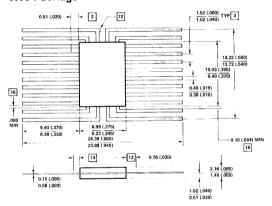
CONSTRUCTION NOTES: 9c, 10d, 11c

WJ Package



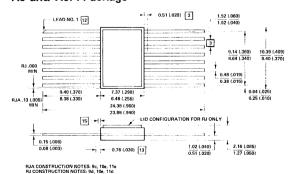
CONSTRUCTION NOTES: 9c, 10d, 11e

WN Package

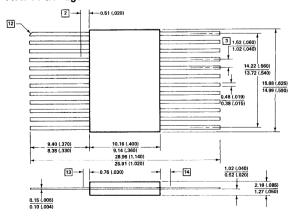


CONSTRUCTION NOTES: 9c, 10d, 11c

RJ and RJA Package

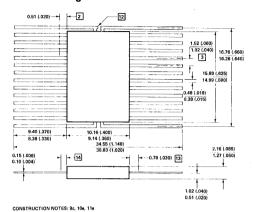


RNA Package

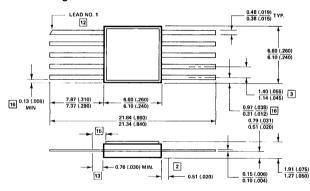


HERMETIC: Flat Packs (cont'd.)

RQA Package

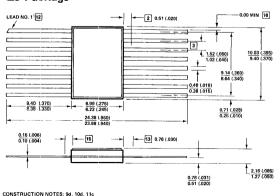


QF Package

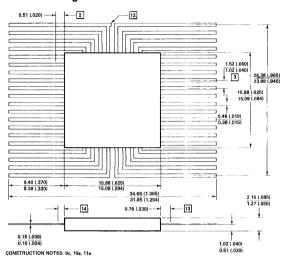


CONSTRUCTION NOTES: 9d, 10d, 11c

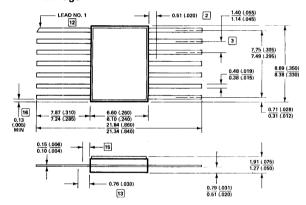
QJ Package



RWA Package

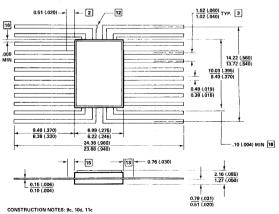


QH Package



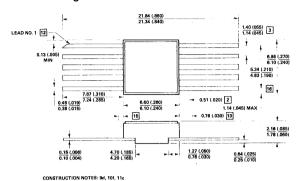
CONSTRUCTION NOTES: 9d, 10d, 11c

QN Package

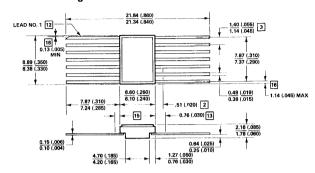


HERMETIC: Flat Packs (cont'd.)

QFA Package

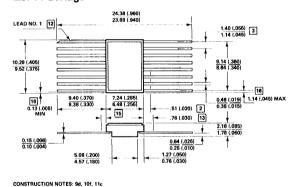


QHA Package

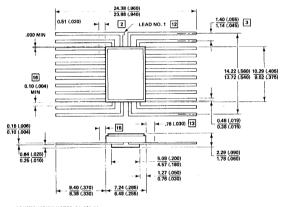


CONSTRUCTION NOTES: 9d, 10f, 11c

QJA Package



QNA Package



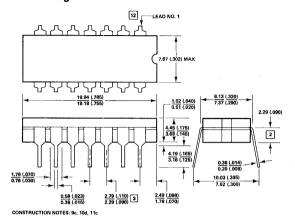
CONSTRUCTION NOTES: 9d, 10f, 11c

RKA Package

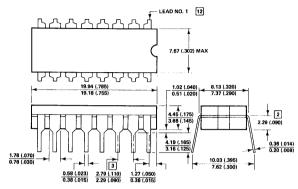
Package not yet available Scheduled for 1977 release

HERMETIC: Cerdip

FH Package

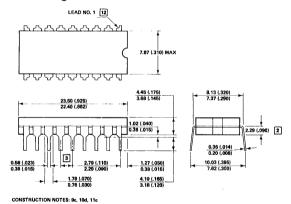


FJ Package

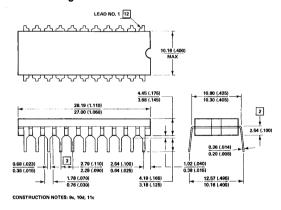


CONSTRUCTION NOTES: 9c, 10d, 11c

FK Package

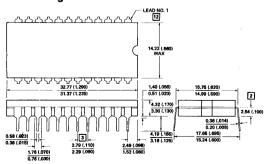


FM Package



FN Package

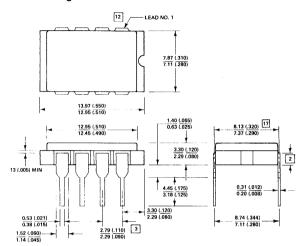
CONSTRUCTION NOTES: 90, 10d, 11c



signetics

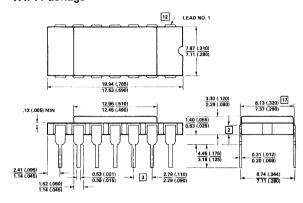
HERMETIC: Laminated Ceramic, Side Brazed Lead

IEA Package



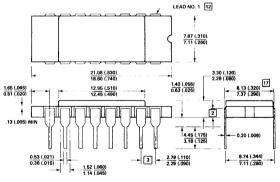
CONSTRUCTION NOTES: 9e, 10f, 11c

IHA Package



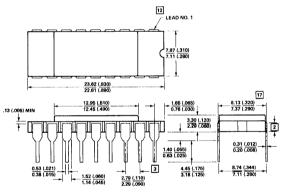
CONSTRUCTION NOTES: 9e. 10f, 11c

IJA Package



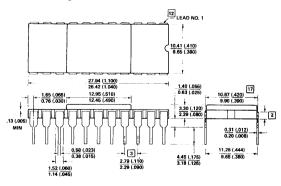
CONSTRUCTION NOTES: 9e, 10f, 11c

IKA Package



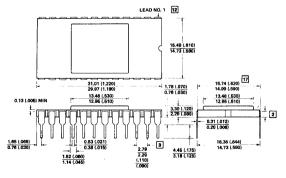
CONSTRUCTION NOTES: 9a, 10f, 11c

IMA Package



CONSTRUCTION NOTES: 96, 10f, 11c

INC Package and IND Package

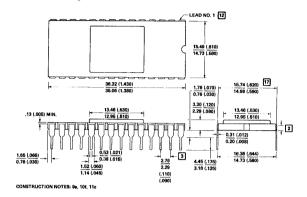


CONSTRUCTION NOTES: 9e, 10f, 11c, 11f (IND)

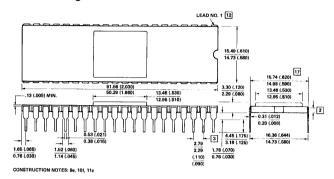
PACKACIES

HERMETIC: Laminated Ceramic, Side Brazed Lead (cont'd.)

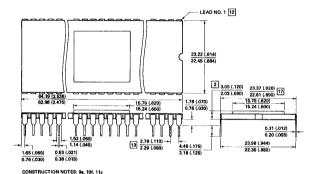
IQA Package



IWA Package



IZA Package



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