

# Datapoint 2200

## **PROGRAMMERS' MANUAL**

A product of COMPUTER TERMINAL CORPORATION  
9725 Datapoint Drive  
San Antonio, Texas 78229

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**PROGRAMMERS' MANUAL**

**August, 1971**

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SECTION 1

DATAPoint 2200

REFERENCE MANUAL

# PART 1

## GENERAL FEATURES

### 1.1 INTRODUCTION

The Computer Terminal Corporation Datapoint 2200 is an integrated data system consisting of an alphanumeric keyboard for data entry, a cathode ray screen for data display, two digital cassette recorders for bulk data storage, a general purpose digital computer for control, and a communications capability for interface with external devices and communications facilities.

Through programming of the control computer the Datapoint 2200 may be used for an infinite variety of data processing applications.

The achievement of a small computer with integrated keyboard, display, storage and communications at such low cost now makes possible computer sophistication for applications not previously practical - particularly in the computer terminal/data entry/communications area.

This manual describes the specific hardware details of the Datapoint 2200. For information regarding specific applications the Datapoint 2200 Programmers' Guide and specific application manuals should be referred to.

### 1.2 SYSTEM ELEMENTS

There are four basic system elements in the basic Datapoint 2200 plus the capability of interface to a number of external peripheral devices.

This manual covers the basic elements (c.r.t., keyboard, processor, cassette tape decks) and one external device (communications adaptor).

### 1.3 CRT DISPLAY

The Datapoint 2200 CRT Display provides the following features:

- a. 7" x 2-1/2" viewing area;
- b. 960 characters;
- c. 80 character by 12 line format;
- d. 4/32" x 3/32" character size;
- e. Entire 94 character ASCII set;
- f. 60 frame per second refresh rate;
- g. 5 x 7 matrix character generation;
- h. 5 x 7 solid, blinking cursor, alternates with character, nondestructive;
- i. P31 green phosphor;
- j. Single control line erasure, frame erasure, and page roll-up; and
- k. Direct control of all c.r.t. functions by the 2200 processor, providing tab, editing, form control, etc.

### 1.4 KEYBOARD

The integral keyboard provides a basic 41 key alphanumeric key group, an 11 key numeric group and five system control keys.

The keyboard provides a unique multi-key roll-over characteristic providing maximum ease of typing. Transfer of characters from the keyboard is under control of the 2200 processor. An audible click providing an acoustical feedback to the typist is available under processor control.

A programmable audio "beep" is also provided when it is desired to gain a typist's attention.

### 1.5 PROCESSOR

The integral processor provides all control functions and includes:

- a. 28 different instruction types;
- b. 7 addressable registers;
- c. 7 deep pushdown stack;
- d. 8 bit memory word length;
- e. Up to 8192 word memory;
- f. Complete parallel I/O system;
- g. Automatic power-up restart.

## 1.6 CASSETTE TAPE DECKS

Two read-write tape decks are provided for program and data storage. The deck accepts Norelco-type cassettes and provides:

- a. 47 characters per inch density;
- b. Dual-capstan forward-reverse operation;
- c. Processor controlled data transfer, direction control, and high-speed rewind.

## 1.7 COMMUNICATIONS ADAPTOR

The communications adaptor is a unique feature of the Datapoint 2200 system. There are four versions of the adaptor:

- a. EIA RS-232 interface for use with external data sets or peripherals;
- b. High-level keying interface for connection to telegraph-type communications channels or equipment;
- c. 103-type data set interface for direct connection to common carrier lines, and including automatic dialing and answering;
- d. 202-type data set interface with 150 bit/sec supervisory channel operation for direct connection to common carrier lines, and including automatic dialing and answering.

The adaptor permits program selection of the desired bit rate, character length, and character set providing the most versatile communications capability yet provided for a remote terminal.

## 1.8 GENERAL SPECIFICATIONS

The Datapoint 2200 has the following general characteristics:

- a. 105-135 v.a.c., 60 cycle, 180 watts, power input;
- b. 47 pounds weight;
- c. 9-5/8" high, 18-1/2" wide, by 19-5/8" deep outside dimensions;
- d. 0° to 50° C (32° to 122° F), 10 to 90 percent relative humidity operation environment.

## 1.9 OPTIONAL PERIPHERALS

A number of optional peripherals are available (in addition to the communications adaptor) for use with the Datapoint 2200 including a:

- a. 132 column, 30 c.p.s. impact page printer; and a
- b. IBM compatible magnetic tape deck.

For further information on these devices reference should be made to their respective reference manuals.

## PART 2

### BASIC PROCESSOR

#### 2.1 PROCESSOR ORGANIZATION

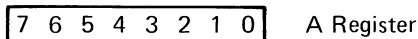
The processor contained in the Datapoint 2200 is comprised of the Arithmetic/Logic Unit, 7 program accessible registers, 2K to 8K words of read/write memory, an instruction decoder and a seven-level hardware pushdown stack used in sub-routine type operations.

#### 2.2 ARITHMETIC/LOGIC UNIT

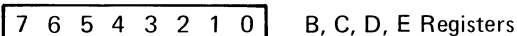
The Arithmetic/Logic Unit is capable of processing both binary integers and logical operands. All arithmetic and logical operations may take place between the A-register and any of the 7 program accessible registers (or between the A-register and memory). The A-register always contains the result of an arithmetic or logic operation, with the other register (or memory cell) being unaffected. Arithmetic and logic operations affect the Sign, Carry, Zero and Parity Flip-flops.

#### 2.3 PROCESSOR REGISTERS

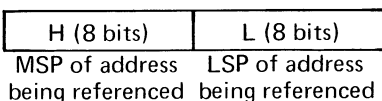
A - The Accumulator register is used to hold the result of all arithmetic and logical instructions. All data transfers into or out of the computer take place through this register.



B, C, D, E - These are general purpose registers which may be used in conjunction with the Accumulator in arithmetic and logical operations. Each register may be loaded from or stored into memory or another register. When used in conjunction with the A and H, L registers, the B, C, D and E registers may function as indexes.

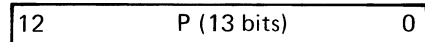


H, L - The H and L Registers are utilized to contain respectively the most significant portion (MSP) and least significant portion (LSP) of the address of a memory location being referenced. All memory reference instructions utilize these registers with the exception of CALL and JUMP commands. However, the H and L Registers may be used as general purpose registers when not being used as above.



P - The program register or "Location Counter" contains the address of the next instruction to be executed. This

register is stored in the pushdown stack upon the execution of a "CALL" instruction and is loaded with the effective address upon execution of a "JUMP", "CALL" or "RETURN" instruction. The P register is 13 bits in length and is capable of addressing up to 8K of memory.



I - The I register is the register which holds the "operation code" of the instruction currently being executed. The contents of I are gated through a decoding network to determine what operation, internal or external, is to be performed.

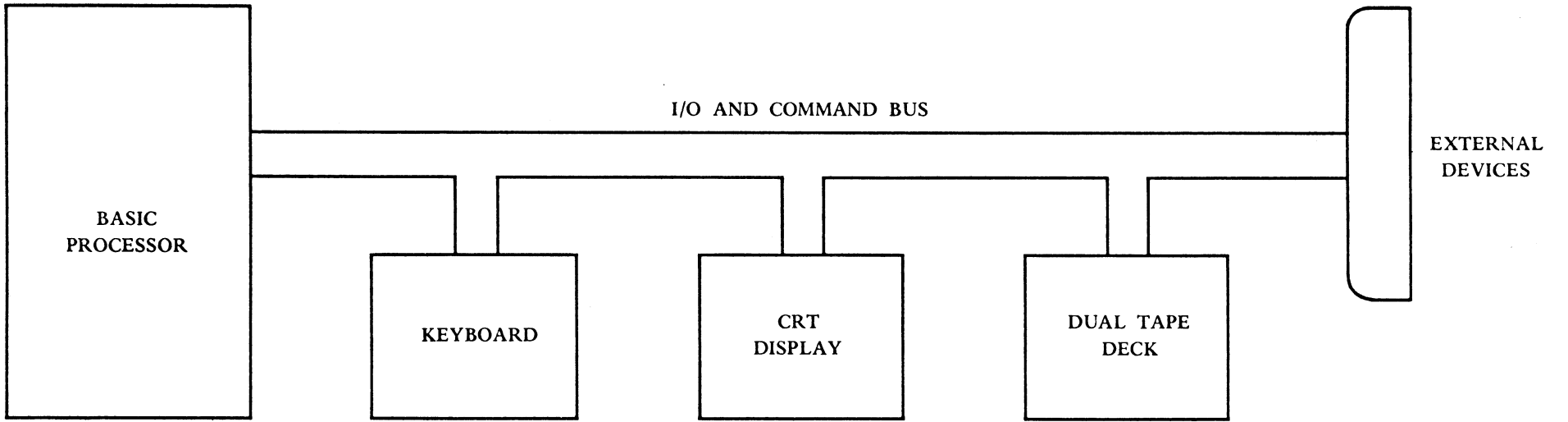
#### 2.4 MEMORY

The basic Datapoint 2200 is supplied with 2048 eight-bit words of memory. Additional modules of 2048 words each may be incorporated with the total memory capacity of the processor being 8192 words. Each 2K memory is made up of 32 individual MOS shift registers with each one having a capacity of 512 bits or 64 eight-bit words. These registers are clocked at a rate of 1.2 MHz. Data is read out in bit serial fashion with one word taking 8 microseconds. During this period of time, two clock pulse times are available for the processor to perform any necessary gating or testing functions.

The Datapoint 2200 memory might be likened to a drum type memory in some respects. It takes approximately 1/2 millisecond for the memory to completely circulate. Thus, if the current instruction referenced a memory location for data access, there would be a 1/2 millisecond delay before that instruction could be completed. However, unlike a drum memory the MOS memory may be stopped during instruction execution so that each succeeding instruction may be read from memory without delay (in 8 usec.).

Physically, instructions require a variable number of cycles for completion. In the first cycle, the instruction is fetched from memory and decoded. If the instruction involves no memory reference, it is then executed within 8 microseconds for a total completion time of 16 microseconds.

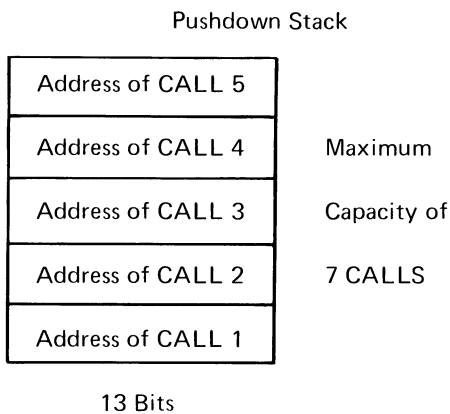
"Immediate" type instructions are the same as instructions requiring no memory reference and require a 16 usec interval for the operand fetch and execute cycle. Jump and Call type instructions require a variable amount of time for execution, depending on the difference between the old and new locations.



**FIGURE 2-1**  
DATAPOINT 2200  
BLOCK DIAGRAM

## 2.5 PUSHDOWN STACK

A unique feature of a machine this size is the incorporation into the processor's structure of a pushdown stack which is useful in any type of application which requires program subroutines. The stack automatically stores the contents of the P register upon execution of a "CALL" instruction and automatically restores P upon execution of a "RETURN". The stack is a group of "last-in/first-out" registers and has a capacity of 7 CALLS. Note that "CALLS" may be "nested", that is more than one CALL may be made before the execution of a RETURN. The execution of a "RETURN" will cause processor control to be given to the next instruction following the last executed CALL.



## 2.6 CONTROL FLIP-FLOPS

Also contained in the basic processor are four control flip-flops which reflect the state of the arithmetic logic unit and which may be tested through the execution of a conditional jump, call or return instruction. The flip-flop mnemonics with their associated functions are as follows:

**C<sub>f</sub>**-Carry Flip-flop. Set when an arithmetic operation results in either a carry (add) or borrow (subtract).\* The Carry Flip-flop also reflects the state of the most significant bit in the accumulator after completion of a shift right instruction. Likewise, it reflects the state of the accumulator least significant bit after completion of a shift left instruction.

**Z<sub>f</sub>**-Zero Flip-flop. Set when the result of an arithmetic or logical operation is equal to zero.\*

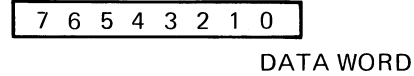
**S<sub>f</sub>**-Sign Flip-flop. This flip-flop reflects the state of bit 7 in the accumulator after an arithmetic type operation.\*

**P<sub>f</sub>**-Parity Flip-flop. Indicates the parity or "number of one bits" contained in the accumulator. If this flip-flop is set (true), the A register contains an odd number of one bits; if it is reset (false), the A register contains an even number of one bits.\*

\*In the event of a compare instruction the contents of the accumulator are not changed; however, the control flip-flops reflect the equivalent of a subtract instruction.

## 2.7 DATA FORMAT

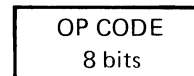
Data is represented in the Datapoint 2200 in the form of 8-bit binary integers.



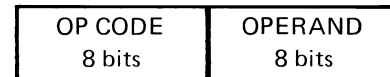
## 2.8 INSTRUCTION FORMATS (GENERAL)

Instruction formats, dependent upon the operation to be performed, may be eight, sixteen or twenty-four bits in length.

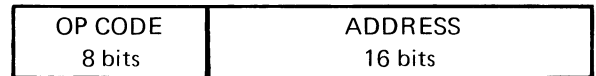
Type-1- register to register, memory reference, arithmetic, logical, shift instructions

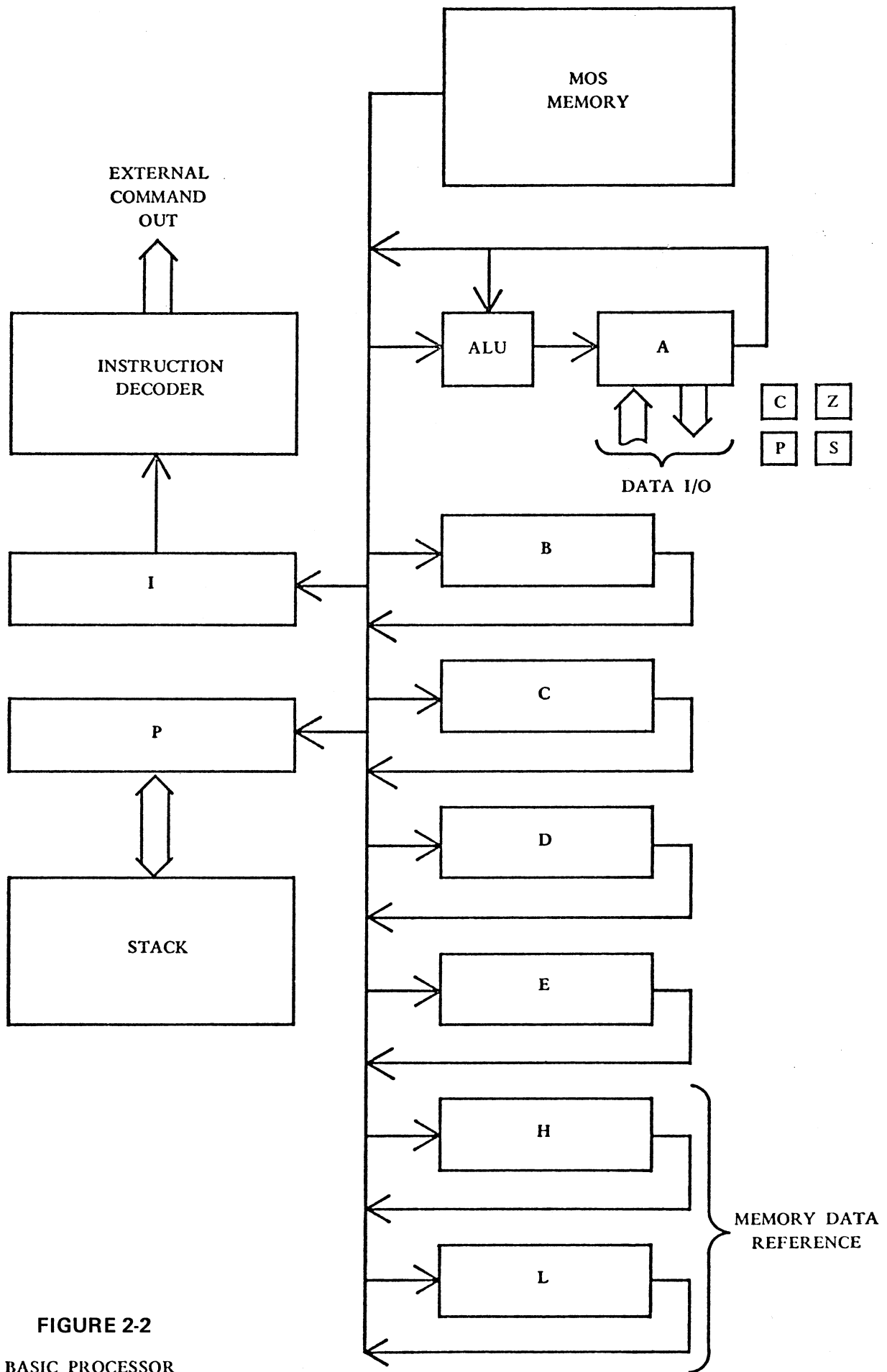


Type-2- immediate mode instructions



Type-3- JUMP & CALL instructions





**FIGURE 2-2**  
**BASIC PROCESSOR**





**LOAD IMMEDIATE:  $Lr_d$**   
 OP CODE: 0d6                      TIMING: 16 usec.  
 OPERATION:  $(P+1) \rightarrow r_d$ ,  $P+2 \rightarrow P$   
 DESCRIPTION: Transfers the contents of the memory location immediately following the instruction, to the register specified by bits 3-5(d) of the instruction.

INSTRUCTION FORMAT:

P			P+1						
7	6	5	4	3	2	1	0	7	0
0		d			6			OPERAND	

d: is the destination designator  
 d=7: is not allowed

NOTE

1. The contents of P+1 are unchanged.
2. None of the Flag Flip-flops are affected.
3. Refer to Table 3-1 for destination codes.

**LOAD:  $Lr_dM, Lr_dr_s, LMr_s$**   
 OP CODE: 3ds                      TIMING: 16 usec. for register to register transfers, 520 usec. for memory reference.

OPERATION:  $(M) \rightarrow r_d$   $s=7, d \leq 6$  ( $Lr_dM$ )  
 $(r_s) \rightarrow r_d$   $s \leq 6, d \leq 6$  ( $Lr_dr_s$ )  
 $(r_s) \rightarrow M$   $s \leq 6, d=7$  ( $LMr_s$ )

DESCRIPTION: Transfers the operand from the source specified by bits 0-2 of the instruction word to the destination specified by bits 3-5 of the instruction word.

INSTRUCTION FORMAT:

P							
7	6	5	4	3	2	1	0
3		d			s		

d: designates the destination of data.  
 s: designates the source. If either s or d=7 a memory reference is indicated and the contents of registers H&L specify the address of the memory location.

NOTE

1. The data source is unaffected.
2. s & d both = 7 results in a Halt instruction.
3. None of the Flag Flip-flops are affected by execution of this instruction.
4. s=d results in a NOP, except as stated in Note 2.

**ADD IMMEDIATE: AD**

OP CODE: 004 TIMING: 16 usec.

OPERATION: (A) + (P+1) → A, P+2 → P

DESCRIPTION: Adds to the contents of the A register the contents of the memory location immediately following the instruction, and retains the sum in the A register. Sets the C<sub>f</sub> Flip-flop if ADD overflow occurs, otherwise resets C<sub>f</sub>.

## INSTRUCTION FORMAT:

P				P+1				
7	6	5	4	3	2	1	0	
0	0	4	OPERAND				7	0

## NOTE

1. The Sign, Zero and Parity Flip-flops will indicate the status of the A register at completion.
2. The contents of P+1 are unchanged.
3. The Carry Flip-flop is cleared at the beginning of this instruction.

**ADD: AD<sub>r<sub>s</sub></sub> ADM**

OP CODE: 20s TIMING: 16 usec. if RR, 520 usec. if MR

OPERATION: (A) + (r<sub>s</sub>) → A or (A) + (M) → A

DESCRIPTION: This instruction is identical to ADD IMMEDIATE with the exception of operand source.

## INSTRUCTION FORMAT:

P							
7	6	5	4	3	2	1	0
2	0	s					

s: specifies the operand source. Refer to Table 3-1 for source codes.

**ADD WITH CARRY****IMMEDIATE: AC**

OP CODE: 014 TIMING: 16 usec.

OPERATION: (A) + (P+1) + (C<sub>f</sub>) → A, P+2 → P

DESCRIPTION: Adds the C<sub>f</sub> bit and the contents of the location immediately following the instruction to the contents of the A register, and retains the sum in the A register. If add overflow occurs, the C<sub>f</sub> Flip-flop is set, otherwise C<sub>f</sub> is reset.

## INSTRUCTION FORMAT:

P				P+1				
7	6	5	4	3	2	1	0	
0	1	4	OPERAND				7	0

## NOTE

1. The Sign, Zero and Parity Flip-flops will indicate the status of the A register at completion.
2. The contents of P+1 remain unchanged.

**ADD WITH CARRY: AC<sub>r<sub>s</sub></sub> ACM**

OP CODE: 21s TIMING: 16 usec. if RR, 520 usec. if MR

OPERATION: (A) + (C<sub>f</sub>) + (r<sub>s</sub>) → A or (A) + (C<sub>f</sub>) + (M) → A

DESCRIPTION: This instruction is identical to ADD WITH CARRY IMMEDIATE with the exception of operand source.

## INSTRUCTION FORMAT:

P							
7	6	5	4	3	2	1	0
2	1	s					

s: specifies the operand source. Refer to Table 3-1 for source codes.

**SUBTRACT IMMEDIATE:**

**SU**

OP CODE: 024

TIMING: 16 usec.

OPERATION:  $(A) - (P+1) \rightarrow A, P+2 \rightarrow P$

DESCRIPTION: Subtracts the contents of the memory location immediately following the instruction from the contents of the A register, and retains the difference in the A register. The  $C_f$  Flip-flop is set if underflow occurs.

INSTRUCTION FORMAT:

P			P+1				
7	6	5	4	3	2	1	0
0		2		4		OPERAND	

NOTE

1. The contents of P+1 is unchanged.
2. The Zero, Sign, and Parity Flip-flops represent the status of the A register at the completion of this instruction.

**SUBTRACT WITH BORROW IMMEDIATE:**

**SB**

OP CODE: 034

TIMING: 16 usec.

OPERATION:  $(A) - (P+1) - (C_f) \rightarrow A, P+2 \rightarrow P$

DESCRIPTION: Subtracts the contents of the memory location immediately following the instruction and the  $C_f$  bit, from the contents of the A register. Sets the  $C_f$  bit if underflow occurs, otherwise resets  $C_f$ .

INSTRUCTION FORMAT:

P			P+1				
7	6	5	4	3	2	1	0
0		3		4		OPERAND	

NOTE

1. The contents of P+1 are unchanged.
2. The Zero, Sign, and Parity Flip-flops represent the status of the A register at the completion of this instruction.

**SUBTRACT:**

**SU<sub>r<sub>s</sub></sub> SUM**

OP CODE: 22s

TIMING: 16 usec. if RR, 520 usec. if MR

OPERATION:  $(A) - (r_s) \rightarrow A$  or  $(A) - (M) \rightarrow A$

DESCRIPTION: This instruction is identical to SUBTRACT IMMEDIATE with the exception of operand source.

INSTRUCTION FORMAT:

P							
7	6	5	4	3	2	1	0
2		2		s			

- s: specifies the operand source. Refer to Table 3-1 for source codes.

**SUBTRACT WITH BORROW:**

**SBr<sub>s</sub> SBM**

OP CODE: 23s

TIMING: 16 usec. if RR, 520 usec. if MR

OPERATION:  $(A) - (r_s) - (C_f) \rightarrow A$  or  $(A) - (M) - (C_f) \rightarrow A$

DESCRIPTION: This instruction is identical to SUBTRACT WITH BORROW IMMEDIATE with the exception of operand source.

INSTRUCTION FORMAT:

P							
7	6	5	4	3	2	1	0
2		3		s			

- s: specifies the operand source. Refer to Table 3-1 for source codes.

**AND IMMEDIATE: ND**

OP CODE: 044 TIMING: 16 usec.

OPERATION:  $(P+1) \wedge (A) \rightarrow A, P+2 \rightarrow P$

DESCRIPTION: Forms the logical product of the contents of the A register with the contents of the memory location immediately following the instruction, and places the results in the A register.

**INSTRUCTION FORMAT:**

P				P+1			
7	6	5	4	3	2	1	0
0		4		4			OPERAND

**NOTE**

1. The Carry Flip-flop will be reset upon completion of the operation.
2. The Zero, Sign, and Parity Flip-flops will represent the status of the A register upon completion of the operation.

**SAMPLE OPERATION:**

(A Reg)	0	0	1	1
(P+1)	0	1	0	1
(A Reg)	0	0	0	1

**AND:**

OP CODE: 24s

**ND<sub>r<sub>s</sub></sub>, NDM**

TIMING: 16 usec. if RR, 520 usec. if MR

OPERATION:  $(A) \wedge (r_s) \rightarrow A$ , or  $(A) \wedge (M) \rightarrow A$

DESCRIPTION: This instruction is identical to AND IMMEDIATE with the exception of operand source.

**INSTRUCTION FORMAT:**

7	6	5	4	3	2	1	0
2		4				s	

s: specifies the operand source. Refer to Table 3-1 for source codes.

**OR IMMEDIATE: OR**

OP CODE: 064 TIMING: 16 usec.

OPERATION:  $(A) \vee (P+1) \rightarrow A, P+2 \rightarrow P$

DESCRIPTION: Forms the logical sum of the contents of the A register and the contents of the memory location immediately following the instruction, and places the result in the A register.

**INSTRUCTION FORMAT:**

P				P+1			
7	6	5	4	3	2	1	0
0		6		4			OPERAND

**NOTE**

1. The Carry Flip-flop will be reset at conclusion.
2. The Zero, Sign, and Parity Flip-flops will represent the status of the A register at completion of the operation.

**SAMPLE OPERATION:**

(A Reg)	0	0	1	1
(P+1)	0	1	0	1
(A Reg)	0	1	1	1

**OR:**

OP CODE: 26s

**OR<sub>r<sub>s</sub></sub> ORM**

TIMING: 16 usec. if RR, 520 usec. if MR

OPERATION:  $(A) \vee (r_s) \rightarrow A$ , or  $(A) \vee (M) \rightarrow A$

DESCRIPTION: This instruction is identical to OR IMMEDIATE with the exception of operand source.

**INSTRUCTION FORMAT:**

P							
7	6	5	4	3	2	1	0
2		6				s	

s: specifies the operand source. Refer to Table 3-1 for source codes.

**EXCLUSIVE OR  
IMMEDIATE:**

**XR**

OP CODE: 054

TIMING: 16 usec.

OPERATION: (A)  $\nabla$  (P+1)  $\rightarrow$  A, P+2  $\rightarrow$  P

DESCRIPTION: The logical difference of the contents of the A register and the contents of the memory location immediately following the instruction is formed, and the result replaces the contents of the A register.

**INSTRUCTION FORMAT:**

P			P+1				
7	6	5	4	3	2	1	0
0	5		4		OPERAND		

**NOTE**

1. The Carry Flip-flop will be reset at conclusion.
2. The Zero, Sign and Parity Flip-flops will represent the status of the A register upon completion of the operation.

**SAMPLE OPERATION:**

(A Reg)	0	0	1	1
(P+1)	0	1	0	1
(A Reg)	0	1	1	0

**EXCLUSIVE OR:**

**XR<sub>r<sub>s</sub></sub> XRM**

OP CODE: 25s

TIMING: 16 usec. if RR, 520 usec. if MR

OPERATION: (A)  $\vee$  (r<sub>s</sub>)  $\rightarrow$  A, (A)  $\vee$  (M)  $\rightarrow$  A

DESCRIPTION: This instruction is identical to EXCLUSIVE OR IMMEDIATE with the exception of operand source.

**INSTRUCTION FORMAT:**

P							
7	6	5	4	3	2	1	0
2	5		s				

s: specifies the operand source. Refer to Table 3-1 for source codes.

**COMPARE**

**IMMEDIATE:**

**CP**

OP CODE: 074

TIMING: 16 usec.

OPERATION: (A) : (P+1), P+2  $\rightarrow$  P

DESCRIPTION: Compares the contents of the A register with the contents of the memory location immediately following the instruction. The flag flip-flops assume the same state as they would for a Subtract instruction.

**INSTRUCTION FORMAT:**

P			P+1				
7	6	5	4	3	2	1	0
0	7		4		OPERAND		

**NOTE**

1. The contents of the A register are unaffected.

**COMPARE:**

**CP<sub>r<sub>s</sub></sub> CPM**

OP CODE: 27s

TIMING: 16 usec. if RR, 520 usec. if MR

OPERATION: (A) : (r<sub>s</sub>) or (A) : (M)

DESCRIPTION: This instruction is identical to COMPARE IMMEDIATE with the exception of operand source.

**INSTRUCTION FORMAT:**

P							
7	6	5	4	3	2	1	0
2	7		s				

s: specifies the operand source. Refer to Table 3-1 for source codes.

## UNCONDITIONAL

### JUMP: JMP

OP CODE: 104 TIMING: Variable\*

OPERATION: (P+1, P+2) → P

DESCRIPTION: An unconditional transfer of control. The contents of P+1 represent the least significant portion of the address, while the contents of P+2 represent the most significant portion.

#### INSTRUCTION FORMAT:

P						P+1		P+2			
7	6	5	4	3	2	1	0	7	0	7	0
OP CODE						ADDRESS					
1	0			4			LSP		MSP		

The three high order bits in the address are ignored, the remaining 13 bits specify the address to which control is to be transferred.

#### NOTE

\*Timing is variable dependent upon cyclic difference between instruction and effective address locations.

## JUMP IF CONDITION

### TRUE: JTf<sub>c</sub>

OP CODE: 1(c+4)0 TIMING: Variable if condition true, 24 usec. if condition false.

OPERATION: If (f<sub>c</sub>=TRUE), (P+1, P+2) → P. Otherwise, P+3 → P.

DESCRIPTION: Examines the designated flip-flop. If set, transfers control to the address designated by the contents of the two memory locations immediately following the instruction. If the selected flip-flop is reset, executes the next sequentially available instruction.

#### INSTRUCTION FORMAT:

P						P+1		P+2			
7	6	5	4	3	2	1	0	7	0	7	0
OP CODE						ADDRESS					
1	c+4			0			LSP		MSP		

c: designates which flip-flop condition is to be tested. Refer to Table 3-2 for list of Flip-flop codes.

#### NOTE

1. The condition of the selected Flip-flop is unchanged by this instruction.

**JUMP IF CONDITION**

**FALSE: JFf<sub>c</sub>**

OP CODE: 1c0

TIMING: Variable if condition false, 24 usec. if condition true.

OPERATION: If (f<sub>c</sub>=FALSE), (P+1, P+2) → P. Otherwise P+3 → P.

DESCRIPTION: Examines the designated flip-flop. If reset, transfers control to the address designated by the contents of the two memory locations immediately following the instruction. If the selected flip-flop is set, executes the next sequentially available instruction.

**INSTRUCTION FORMAT:**

P			P+1				P+2				
7	6	5	4	3	2	1	0	7	0	7	0
1	c		0			LSP				MSP	
OP CODE						ADDRESS					

c: designates which flip-flop (condition) is to be tested. Refer to Table 3-2 for list of flip-flop codes.

**NOTE**

1. The condition of the selected flip-flop is unchanged by this instruction.

**SUBROUTINE CALL: CALL**

OP CODE: 106

TIMING: Variable

OPERATION: P+3 → STACK, (P+1, P+2) → P

DESCRIPTION: Transfers the address of the next sequentially available instruction to the Pushdown Stack, and transfer control to the address specified by the contents of the two memory locations immediately following the Op Code.

**INSTRUCTION FORMAT:**

P			P+1				P+2					
7	6	5	4	3	2	1	0	7	0	0	7	0
1	0		6			LSP				MSP		
ADDRESS												

**NOTE**

1. The Stack is open-ended in operation. If it is overfilled, the deepest address will be lost.

## CONDITIONAL SUBROUTINE CALL

**IF CONDITION TRUE:**  $CTf_c$

OP CODE:  $1(c+4)2$

TIMING: Variable if condition true, 24 usec. if condition false.

OPERATION: If ( $f_c=TRUE$ ),  $P+3 \rightarrow STACK$ ,  $(P+1, P+2) \rightarrow P$ . Otherwise,  $P+3 \rightarrow P$ .

DESCRIPTION: Examines the designated flip-flop. If set, transfers the address of the next sequentially available instruction to the pushdown stack, and transfers control to the address of the two memory locations immediately following the Op Code. If the selected flip-flop is reset, executes the next sequentially available instruction.

INSTRUCTION FORMAT:

			P+1				P+2					
7	6	5	4	3	2	1	0	7	0	7	0	
1	c+4			2			LSP				MSP	

ADDRESS

c: designates which flip-flop (condition) is to be tested.

### NOTE

1. The condition of the selected flip-flop is unchanged by this instruction.
2. The stack is open-ended in operation. If it is overfilled, the deepest address will be lost.
3. Refer to Table 3-2 for list of flip-flop codes.



## CONDITIONAL SUBROUTINE CALL

### IF CONDITION FALSE: $CFf_c$

OP CODE: 1c2

TIMING: Variable if condition false, 24 usec. if condition true.

OPERATION: If ( $f_c=FALSE$ ),  $P+3 \rightarrow STACK$ ,  $(P+1, P+2) \rightarrow P$ .

DESCRIPTION: Examines the designated flip-flop. If reset, transfers the address of the next sequentially available instruction to the pushdown stack, and transfers control to the address of the two memory locations immediately following the Op Code. If the selected flip-flop is set, executes the next sequentially available instruction.

### INSTRUCTION FORMAT:

			P+1				P+2									
7	6	5	4	3	2	1	0	7			0	7			0	
1			c			2			LSP				MSP			

ADDRESS

c: designates which flip-flop (condition) is to be tested.

### NOTE

1. The condition of the selected flip-flop is unchanged by this instruction.
2. The stack is open-ended in operation. If it is overfilled, the deepest address will be lost.
3. Refer to Table 3-2 for list of flip-flop codes.

## SUBROUTINE

### RETURN:

OP CODE: 007

OPERATION: (STACK) → P

DESCRIPTION: Transfer control to the address specified by the most recent entry in the Pushdown Stack. Deletes the most recent entry from the Stack.

INSTRUCTION FORMAT:

P							
7	6	5	4	3	2	1	0
0		0					7

### NOTE

1. The effect of attempting more "RETURN" than the Stack is capable of handling is undefined.

## CONDITIONAL SUBROUTINE RETURN

### IF CONDITION TRUE:

**RTf<sub>c</sub>**

OP CODE: 0(c+4)3

TIMING: Variable if condition true, 16 usec. if condition false.

OPERATION: If (f<sub>c</sub>=TRUE), Stack → P. Otherwise P+1 → P

DESCRIPTION: Examines the designated flip-flop. If set, transfers control to the address specified by the most recent entry in the pushdown stack. Deletes the most recent entry in the stack. If the selected flip-flop is reset, executes the next sequentially available instruction.

INSTRUCTION FORMAT:

7	6	5	4	3	2	1	0
1		c+4					3

c: designates which flip-flop (condition) is to be tested.

### NOTE

1. The condition of the selected flip-flop is unchanged by this instruction.
2. The effect of attempting more "RETURN" than the stack is capable of handling is undefined.
3. Refer to Table 3-2 for list of flip-flop codes.

## CONDITIONAL SUBROUTINE RETURN

### IF CONDITION FALSE:

**RFf<sub>c</sub>**

OP CODE: 0c3

TIMING: Variable if condition false, 16 usec. if condition true.

OPERATION: If (f<sub>c</sub>=FALSE), Stack → P. Otherwise, P+1 → P

DESCRIPTION: Examines the designated flip-flop. If reset, transfers control to the address specified by the most recent entry in the stack. If the selected flip-flop is set, executes the next sequentially available instruction.

INSTRUCTION FORMAT:

7	6	5	4	3	2	1	0
1			c				3

c: designates which flip-flop (condition) is to be tested.

### NOTE

1. The condition of the selected flip-flop is unchanged by this instruction.
2. The effect of attempting more "RETURN" than the stack is capable of handling is undefined.
3. Refer to Table 3-2 for list of flip-flop codes.

**SHIFT RIGHT****CIRCULAR:****SRC**

OP CODE: 012

TIMING: 16 usec.

OPERATION:  $A_m \rightarrow A_{m-1}$ ,  $A_0 \rightarrow A_7$ ,  $A_0 \rightarrow C_f$ 

DESCRIPTION: Shifts the contents of the A register right in a circular fashion. Shifts the least significant bit into the most significant bit position. Upon completion of the operation, the Carry Flip-flop is equal to the most significant bit.

INSTRUCTION FORMAT:

P							
7	6	5	4	3	2	1	0
0		1				2	

NOTE

None of the flag flip-flops other than  $C_f$  is affected by this instruction.

**SHIFT LEFT****CIRCULAR:****SLC**

OP CODE: 002

TIMING: 16 usec.

OPERATION:  $A_m \rightarrow A_{m+1}$ ,  $A_7 \rightarrow A_0$ ,  $A_7 \rightarrow C_f$ 

DESCRIPTION: Shifts the contents of the A register left in a circular fashion. Shifts the most significant bit into the least significant bit position. Upon completion of the operation, the Carry Flip-flop is equal to the least significant bit.

INSTRUCTION FORMAT:

P							
7	6	5	4	3	2	1	0
0		0				2	

NOTE

None of the flag flip-flops other than  $C_f$  is affected by this instruction.

**NO OPERATION:****NOP**

OP CODE: 300

TIMING: 16 usec.

OPERATION:  $P+1 \rightarrow P$ 

DESCRIPTION: No instruction is executed.

INSTRUCTION FORMAT:

P							
7	6	5	4	3	2	1	0
3		0				0	

**HALT:****HALT**

OP CODE: 000,001,377

TIMING: Execution Stops

OPERATION:

DESCRIPTION: The computer halts. When the START button on the console is depressed, operation resumes at  $P+1$ .

INSTRUCTION FORMAT:

P							
7	6	5	4	3	2	1	0
0		0				0	
0		0				1	
3		3				7	

**INPUT:****INPUT**

OP CODE: 101

TIMING: 16 usec.

OPERATION:  $(I/O \text{ Bus}) \rightarrow A$ 

DESCRIPTION: Transfers the contents of the I/O Bus to the A register.

INSTRUCTION FORMAT:

P							
7	6	5	4	3	2	1	0
1		0				1	

**EXTERNAL COMMAND: EX (exp)**

OP CODE: 121-177 depending on the specific command being executed. TIMING: 16 usec.

OPERATION: Performs I/O control functions according to (exp)

DESCRIPTION: These instructions perform the functions necessary for control of the I/O system and external devices. Many of these functions are specifically related to operation of particular devices. The device oriented commands for the Keyboard, CRT Display, Cassette Tapes, and Communications Interface are explained in the sections covering these devices.

**INSTRUCTION FORMAT:**

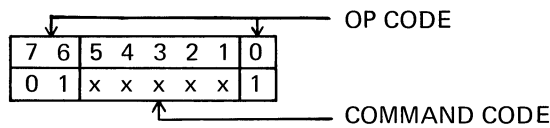


Table 3-3 is a list of External Commands used. For a detailed discussion of their use, reference should be made to Part 4 (Input/Output Operations) and to descriptions of the separate external devices.

**TABLE 3-3  
EXTERNAL COMMANDS**

**EX (exp)**

COMMAND NUMBER	(exp)	OCTAL CODE	COMMAND	DESCRIPTION	DEVICE ADDRESS
1	ADR	121	Address	Selects device specified by A-register	ALL
2	STATUS	123	Sense Status	Connects selected device status to input lines	
3	DATA	125	Sense Data	Connects selected device data to input lines	
4	WRITE	127	Write Strobe	Signals selected device that output data word is on output lines	
5	COM1	131	Command 1	Outputs a control function to selected device	
6	COM2	133	Command 2	Outputs a control function to selected device	
7	COM3	135	Command 3	Outputs a control function to selected device	
8	COM4	137	Command 4	Outputs a control function to selected device	
9	---	141	(Unassigned)	---	
10	---	143	(Unassigned)	---	---
11	---	145	(Unassigned)	---	---
12	---	147	(Unassigned)	---	---

**TABLE 3-3**  
**EXTERNAL COMMANDS**

EX (exp)

(Continued)

COMMAND NUMBER	(exp)	OCTAL CODE	COMMAND	DESCRIPTION	DEVICE ADDRESS
13	BEEP	151	Beep	Activates tone producing mechanism	341
14	CLICK	153	Click	Activates audible click producing mechanism	341
15	DECK1	155	Select Deck 1	Connects deck 1 to I/O bus	360
16	DECK2	157	Select Deck 2	Connects deck 2 to I/O bus	
17	RBK	161	Read Block	Enables read circuitry and sets tape in forward motion	
18	WBK	163	Write Block	Enables write circuitry and sets tape in forward motion	360
19	---	165	(Unassigned)	---	---
20	BSP	167	Backspace One Block	Backs up the selected tape one record	360
21	SF	171	Slew Forward	Sets selected tape deck in forward motion	
22	SB	173	Slew Backward	Sets selected tape deck in backward motion	
23	REWIND	175	Rewind	Rewinds the selected deck to beginning of tape	
24	TSTOP	177	Stop Tape	Halts motion of the selected tape deck	360

**TABLE 3-4**  
**INSTRUCTION REPERTOIRE**

OP CODE	MNEMONIC	OP CODE	MNEMONIC	OP CODE	MNEMONIC
000	HALT	050		120	JFS
001	HALT	051		121	EX ADR
002	SLC	052		122	CFS
003	RFC	053	RTZ	123	EX STATUS
004	AD	054	XR	124	
005		055		125	EX DATA
006	LA	056	LH	126	
007	RETURN	057		127	EX WRITE
010		060		130	JFP
011		061		131	EX COM1
012	SRC	062		132	CFP
013	RFZ	063	RTS	133	EX COM2
014	AC	064	OR	134	
015		065		135	EX COM3
016	LB	066	LL	136	
017		067		137	EX COM4
020		070		140	JTC
021		071		141	
022		072		142	CTC
023	RFS	073	RTP	143	
024	SU	074	CP	144	
025		075		145	
026	LC	076		146	
027		077		147	
030		100	JFC	150	JTZ
031		101	INPUT	151	EX BEEP
032		102	CFC	152	CTZ
033	RFP	103		153	EX CLICK
034	SB	104	JMP	154	
035		105		155	EX DECK1
036	LD	106	CALL	156	
037		107		157	EX DECK2
040		110	JFZ	160	JTS
041		111		161	EX RBK
042		112	CFZ	162	CTS
043	RTC	113		163	EX WBK
044	ND	114		164	
045		115		165	
046	LE	116		166	
047		117		167	EX BSP

**TABLE 3-4**  
**INSTRUCTION REPERTOIRE**  
**(Continued)**

OP CODE	MNEMONIC	OP CODE	MNEMONIC	OP CODE	MNEMONIC
170	JTP	240	NDA	310	LBA
171	EX SF	241	NDB	311	
172	CTP	242	NDC	312	LBC
173	EX SB	243	NDD	313	LBD
174		244	NDE	314	LBE
175	EX REWND	245	NDH	315	LBH
176		246	NDL	316	LBL
177	EX TSTOP	247	NDM	317	LBM
200	ADA	250	XRA	320	LCA
201	ADB	251	XRB	321	LCB
202	ADC	252	XRC	322	
203	ADD	253	XRD	323	LCD
204	ADE	254	XRE	324	LCE
205	ADH	255	XRH	325	LCH
206	ADL	256	XRL	326	LCL
207	ADM	257	XRM	327	LCM
210	ACA	260	ORA	330	LDA
211	ACB	261	ORB	331	LDB
212	ACC	262	ORC	332	LDC
213	ACD	263	ORD	333	
214	ACE	264	ORE	334	LDE
215	ACH	265	ORH	335	LDH
216	ACL	266	ORL	336	LDL
217	ACM	267	ORM	337	LDM
220	SUA	270	CPA	340	LEA
221	SUB	271	CPB	341	LEB
222	SUC	272	CPC	342	LEC
223	SUD	273	CPD	343	LED
224	SUE	274	CPE	344	
225	SUH	275	CPH	345	LEH
226	SUL	276	CPL	346	LEL
227	SUM	277	CPM	347	LEM
230	SBA	300	NOP	350	LHA
231	SBB	301	LAB	351	LHB
232	SBC	302	LAC	352	LHC
233	SBD	303	LAD	353	LHD
234	SBE	304	LAE	354	LHE
235	SBH	305	LAH	355	
236	SBL	306	LAL	356	LHL
237	SBM	307	LAM	357	LHM

**TABLE 3-4**  
**INSTRUCTION REPERTOIRE**  
(Continued)

OP CODE	MNEMONIC	OP CODE	MNEMONIC	OP CODE	MNEMONIC
360	LLA	370	LMA		
361	LLB	371	LMB		
362	LLC	372	LMC		
363	LLD	373	LMD		
364	LLE	374	LME		
365	LLH	375	LMH		
366		376	LML		
367	LLM	377	HALT		

NOTE

OP Codes shown without Mnemonics are undefined.



## PART 4

### INPUT/OUTPUT OPERATIONS

#### 4.1 GENERAL

The versatile input/output capability of the Datapoint 2200 permits it to communicate with external devices (such as the 2200 communications adaptor) through a parallel I/O system. The keyboard, c.r.t. and tape decks that are an integral part of the Model 2200 perform all operations over the same I/O system as external devices.

#### 4.2 INPUT/OUTPUT INSTRUCTIONS

Two types of instructions provide for I/O operations. One is the INPUT command (see section 3) which, upon execution, transfers whatever is on the input bus to the A-register. The second is the EXTERNAL command which is sub-divided into 24 separate command operations (8 of which are available to devices physically external to the Model 2200). Each EXTERNAL command produces a strobe pulse which may be used for control external to the processor. The actual control functions assigned to each external command are listed in Table 3-3.

#### 4.3 INPUT/OUTPUT CABLE

The parallel I/O cable carries data, input strobe, external commands, and power between the 2200 processor and external devices connected to it. A complete I/O system is structured by connecting external devices in partyline fashion as shown in Figure 2-1. The I/O cable contains 8 input data lines, 8 output data lines, 1 input strobe line, 8 (of the 24) external command lines, 1 clock line, and 7 power and ground lines.

#### 4.4 I/O DATA LINES

The data lines are broken into two groups. 8 lines are used for output and 8 lines are used for input.

The data output lines are connected (at all times) to the A-register in the processor and are used to perform three basic functions:

- a. To transfer an address to select an external device (including the keyboard, c.r.t. and tape decks);
- b. To transfer commands to an addressed device; and
- c. To transfer data to an addressed device.

The data input lines are strobed into the A-register upon execution of the INPUT instruction and used to perform two basic functions.

- a. To transfer status information from an addressed external device; and
- b. To transfer data from an addressed external device.

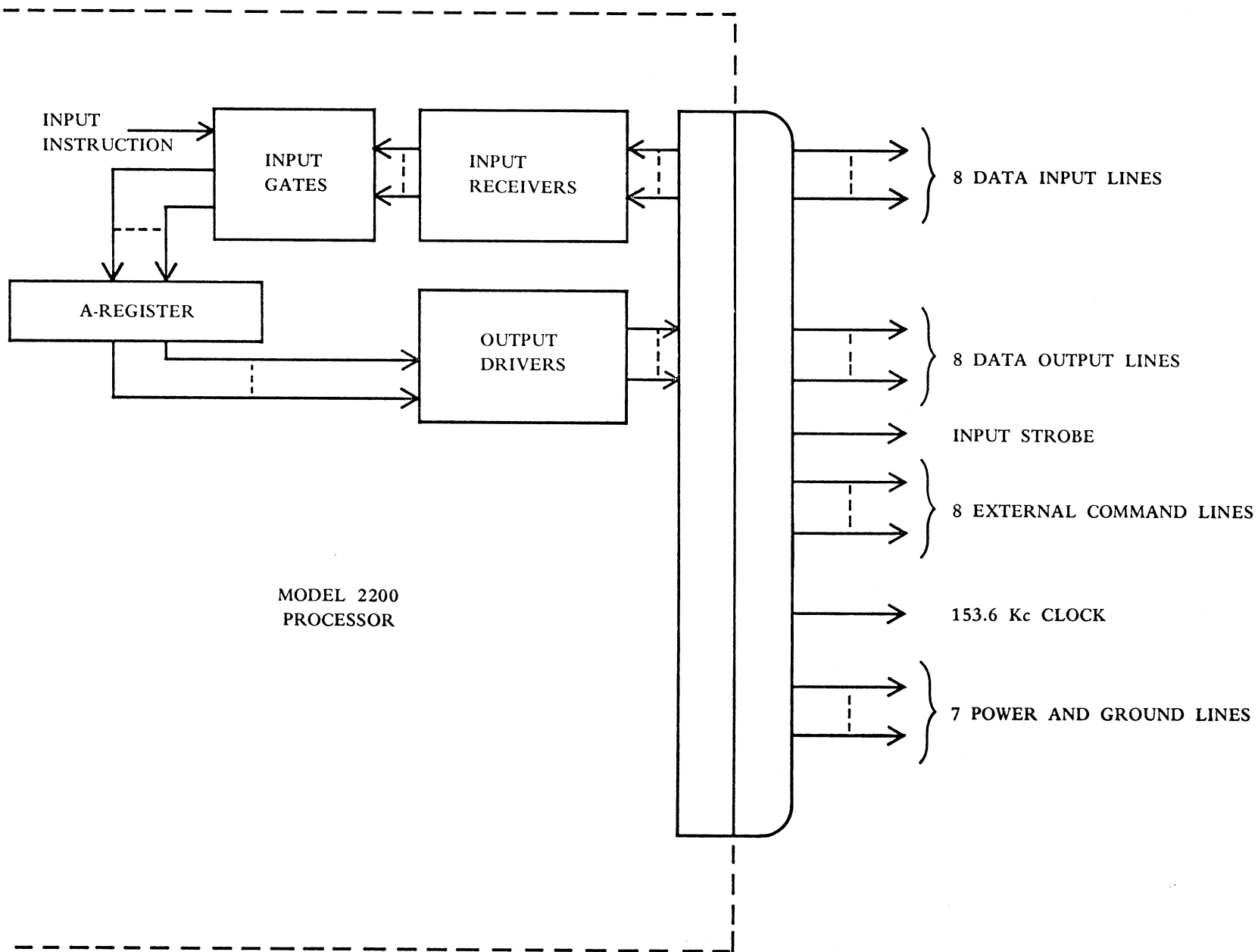
As shown in Figure 4-1, input data or status from the data input lines is processed through input receivers and gated into the A-register. Once in the A-register data can be manipulated or stored as desired. Addresses, commands, or data that is to be transferred to an external device must first be loaded into the A-register. From the A-register it is transmitted through output devices onto the data output lines. The A-register, then, is used as a buffer register between the 2200 processor and external devices for all input and output data transfers.

#### 4.5 INPUT STROBE

The INPUT STROBE carries a signal (8 usec. pulse) from the processor to the external device to indicate that whatever data is on the data input lines has been sampled and transferred into the A-register. The trailing edge of the pulse may be used by an external device to remove data from the data input line or to clear a status bit. The INPUT strobe is generated upon execution of the INPUT instruction.

#### 4.6 EXTERNAL COMMAND STROBES

The eight EXTERNAL commands used by devices physically external to the Model 2200 are given function assignments as follows:



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FIGURE 4-1

I/O SYSTEM, FUNCTIONAL DIAGRAM

**TABLE 4-1**  
**EXTERNAL COMMANDS**

**EX (exp)**

<b>COMMAND NUMBER</b>	<b>(exp)</b>	<b>OCTAL CODE</b>	<b>COMMAND</b>	<b>DESCRIPTION</b>
1	ADR	121	Address	Selects device specified by A-register
2	STATUS	123	Sense Status	Connects selected device status lines to data input bus
3	DATA	125	Sense Data	Connects selected device data lines to data input bus
4	WRITE	127	Write Strobe	Signals selected device that output data is on data output lines
5	COM1	131	Command 1	Signals selected device that a control word is on data output lines
6	COM2	133	Command 2	Signals selected device that a control word is on data output lines
7	COM3	135	Command 3	Signals selected device that a control word is on data output lines
8	COM4	137	Command 4	Signals selected device that a control word is on data output lines

Execution of an EXTERNAL instruction provides a pulse 8 microseconds long. No functions are performed within the 2200 processor during execution of an EXTERNAL instruction. The interpretation of each of the EXTERNAL instructions is as follows:

- a. Address. The address command (EX ADR) is a signal from the processor to all external devices to indicate that the information on the data output bus is to be interpreted as an external device address. Whenever an address command is executed all external devices should be disconnected from the I/O system except the device whose address appears in the A-register. (See paragraph 4.10 for discussion of address assignments).
- b. Sense Status. The sense status (EX STATUS) command is a signal from the processor to the selected external device to place status information on the data input lines. (Note: External devices should be configured such that status is connected to the data input line whenever the device is first addressed. It is only necessary to use the EX STATUS instruction when it is desired to sense status after an EX DATA instruction has been used and a new address sequence has not been executed).
- c. Sense Data. The sense data (EX DATA) command is a signal from the processor to the selected external device to place its data on the data input lines.

- d. Write Strobe. The write strobe (EX WRITE) command is a signal from the processor that data is present on the data output lines for the selected external device.

- e. Command 1 through Command 4. Command 1 through Command 4 (EX COM1, etc.) have meaning appropriate to the device selected. Reference should be made to a description of each device for specific function assignments.

#### **4.7 CLOCK LINE**

The clock line is crystal controlled 153.6 kilohertz square-wave that is available to external devices for timing purposes.

#### **4.8 I/O BUS ELECTRICAL SPECIFICATIONS**

All signals in the I/O System operate with a voltage swing of zero to +5 volts. Line drivers have a source impedance of approximately 470 ohms and line receivers have an input impedance in excess of 18,000 ohms and a decision threshold of +1.7 volts. Figure 4-2 illustrates a typical output line circuit.

All logic levels are True (logical 1) for zero (less than +1.7) volts and False (logical 0) for +5 (greater than 1.7) volts.

#### **4.9 DATA TRANSFER OPERATION**

- a. Data Output. Figure 4-3 illustrates the sequence of events that occur when data is transferred from the

2200 processor to an external device. A typical program sequence to execute a transfer is as follows:

WDATA	LA 0322	Load device address into A-register
	EX ADR	
	INPUT	Load device status into A-register
	SRC	Shift desired status bit into C flip-flop
	JFC EXIT	Exit if device not ready
	LAM	Load A-register with DATA
	EX WRITE	Write Data to device

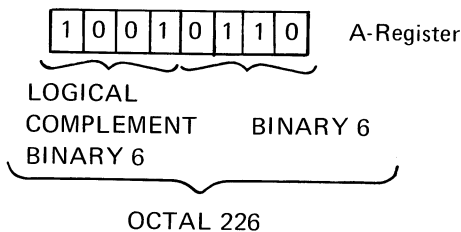
Once a device is addressed it remains addressed until another device is addressed so that succeeding commands may be transmitted to a device without re-addressing the device. Transmitting a command to a device would follow a program sequence similar to a data transfer except that EX COM<sub>n</sub> would replace EX WRITE.

- b. Data Input. Figure 4-4 illustrates a sequence of events that occur when data is transferred from an external device to the 2200 processor. A typical program sequence is as follows:

RDATA	LA 0322	Load A-register with device address
	EX ADR	
	INPUT	Load device status into A-register
	SRC	Shift status bit into C flip-flop
	JFC EXIT	Exit if device not ready
	EX DATA	Place data on input lines
	INPUT	Load A-register with data

#### 4.10 DEVICE ADDRESS NUMBERING

Address assignments in the I/O system provides for up to 16 devices external to the 2200 processor. The address word is formulated such that the low-order four bits form the binary value for the address and the high-order four bits form the logical complement of the low order bits. For example device number 6 would have an address word as follows:



This addressing system permits any device to be coded for its particular address with only a four-input gate strapped to those output lines that are set to one during the address command. In addition, all devices can be cleared by setting the A-register to all zeros and executing an EX ADR instruction.

Device addresses used in the Model 2200 are given in the following table:

**TABLE 4-2**  
**DEVICE ADDRESS ASSIGNMENTS**

DEVICE	NUMBER	BINARY	OCTAL
Cassette Tape Decks	0	11110000	360
CRT/Keyboard	1	11100001	341
Communications Adaptor	2	11010010	322
2200P Printer	3	11000011	303
2200T Tape Transport	4	10110100	264
Unassigned	5	10100101	245
"	6	10010110	226
"	7	10000111	207
"	8	01111000	170
"	9	01101001	151
"	10	01011010	132
"	11	01001011	113
"	12	00111100	074
"	13	00101101	055
"	14	00011110	036
"	15	00001111	017

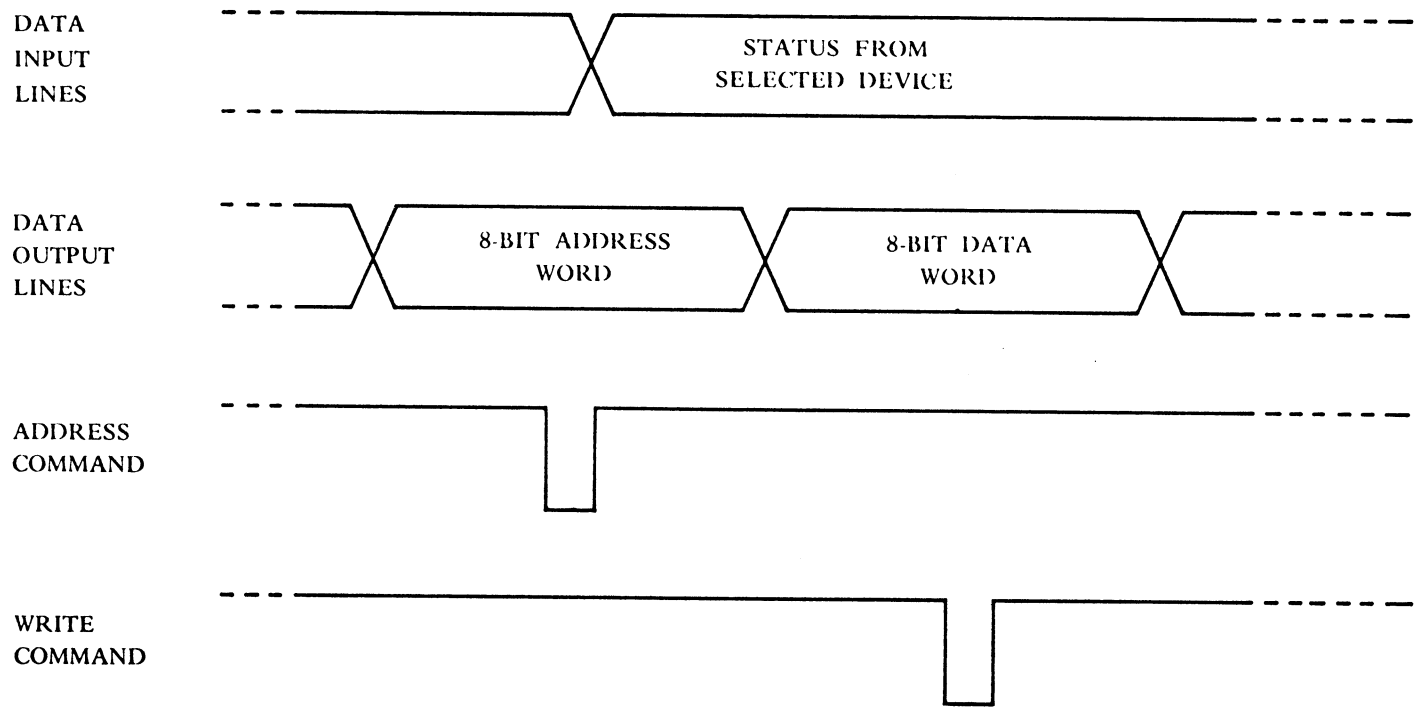
#### 4.11 I/O POWER AND GROUND LINES

The Model 2200 provides several power supply voltages for use by external devices. Table 4-3 below lists the characteristics of each power and ground line.

#### 4.12 I/O SYSTEM CONNECTOR

Connection to the I/O system is made through an Amphenol 17-20500-1 connector. The mating I/O cable should have a 50-pin Amphenol 17-10500-1 connector.

Table 5-5 lists the pin assignments.



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**FIGURE 4-3**  
TYPICAL DATA OUTPUT  
SEQUENCE

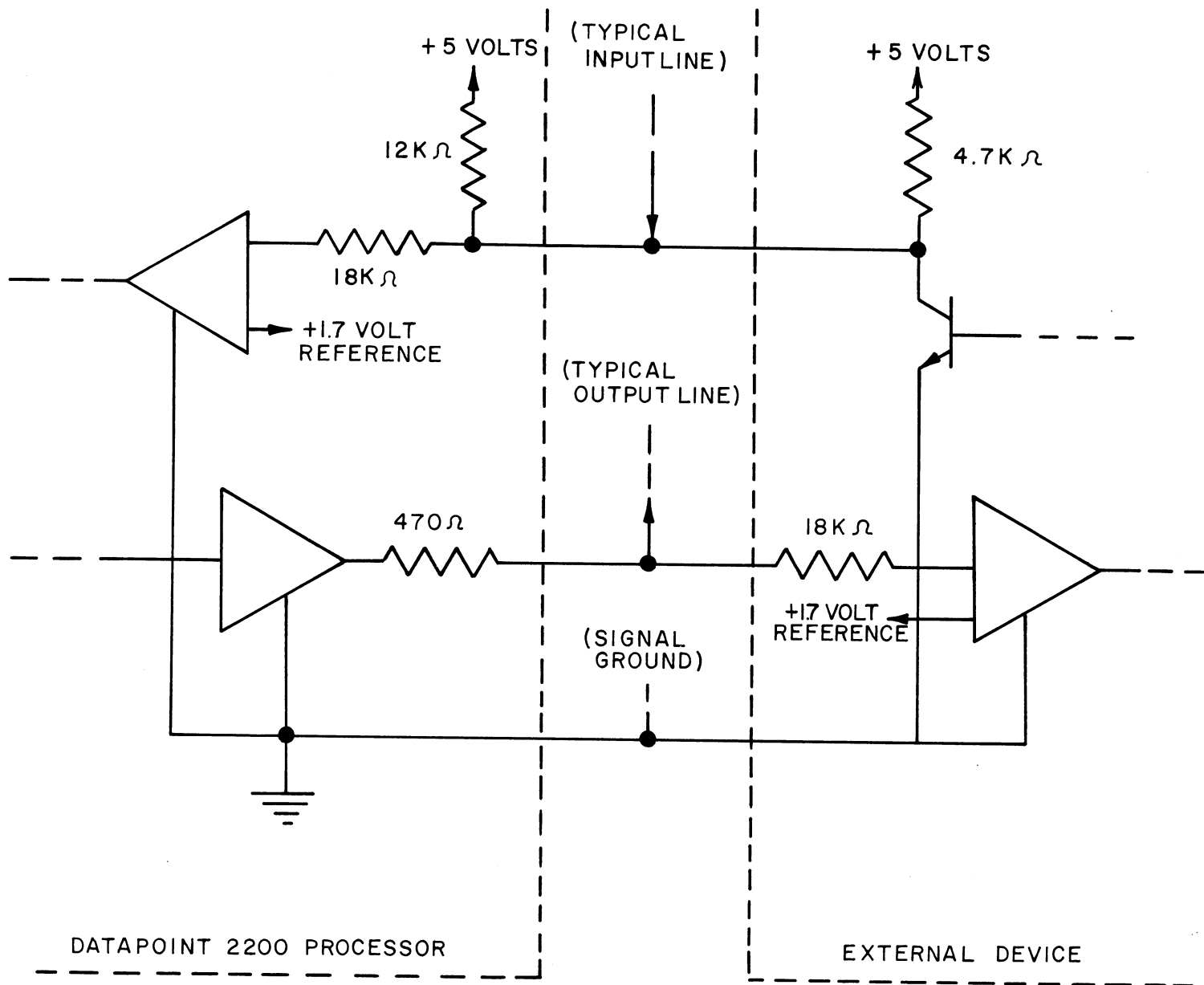


FIGURE 4-2

I/O CABLE, ELECTRICAL CHARACTERISTICS

**TABLE 4-3**  
**I/O POWER AND GROUND LINES**

VOLTAGE	MAX. CURRENT	REGULATION
-12 Volts	0.5 amps	±10%
- 5 Volts	0.1 amps	±5%
+ 5 Volts	3.4 amps	±5%
+12 Volts	0.5 amps	±10%
+24 Volts	0.1 amps	±5%
Power Ground	—	—
Signal Ground	—	—

**TABLE 4-4**  
**I/O CONNECTOR PIN ASSIGNMENTS**

ASSIGNMENT	PIN NUMBER
Data output 0	44
1	45
(A Bus Outputs) 2	46
3	29
4	30
5	31
6	32
7	33
Data Input 0	1
1	2
(A Bus Inputs) 2	3
3	4
4	5
5	6
6	7
7	18
Input Strobe (Read)	12
Address Command	15
Sense Status Command	13
Sense Data Command	14
Write Command	19
Command 1	20
Command 2	21
Command 3	22
Command 4	23
153.6 KHz Clock	39
-12v	24
-5v	27
+5v	8, 9, 10, 11
+12v	25
+24v	26
Ground (Power & Signal)	40, 41, 42, 43

DATA  
INPUT  
LINES

DATA  
OUTPUT  
LINES

ADDRESS  
COMMAND

INPUT  
STROBE

SENSE DATA  
COMMAND

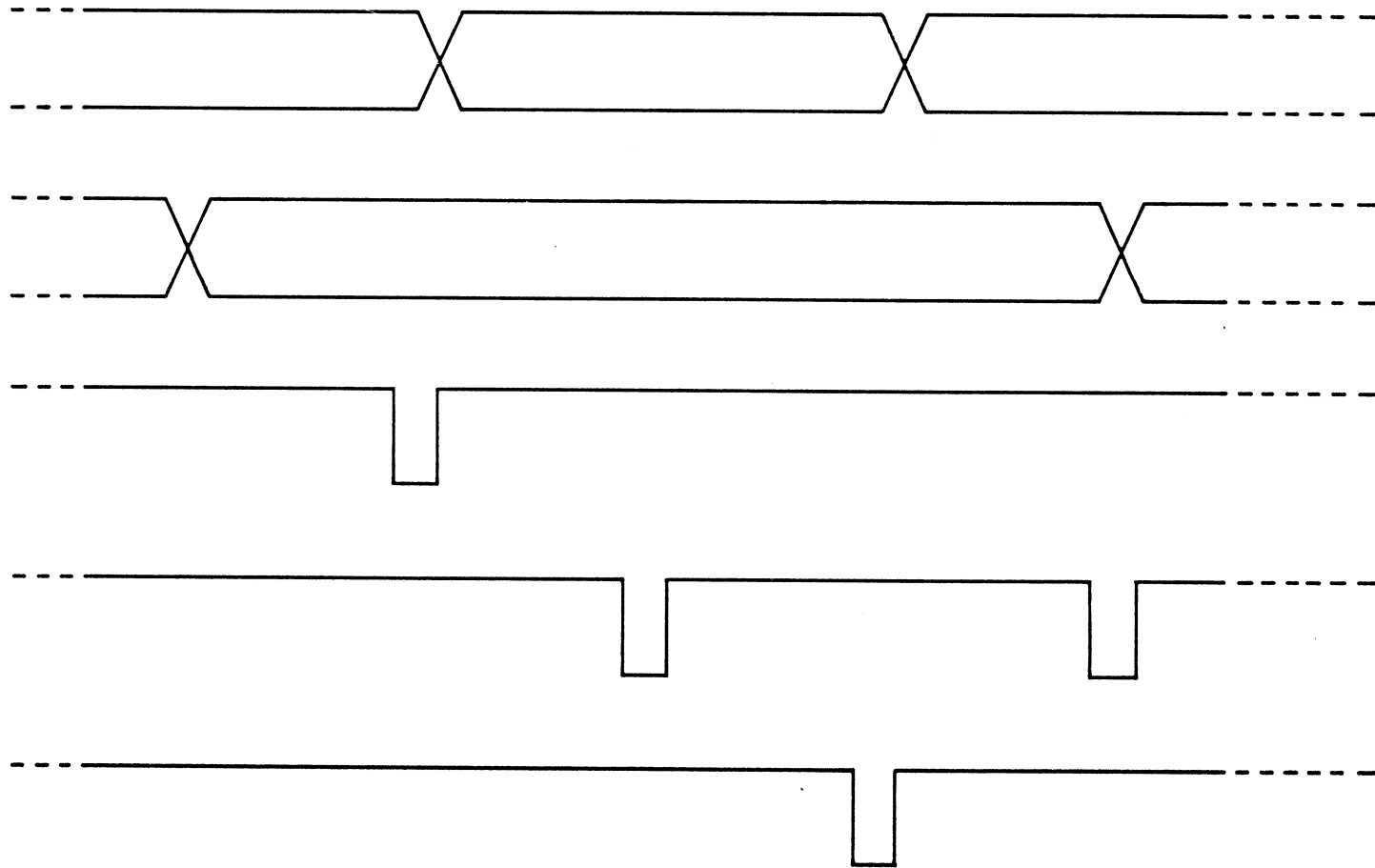


FIGURE 4-4  
TYPICAL DATA INPUT  
SEQUENCE



## PART 5

### KEYBOARD

#### 5.1 GENERAL DESCRIPTION

The keyboard on the Datapoint 2200 performs the functions of data entry and processor control. The keys are divided into three sections, each of which has its own function.

Section 1 consists of 41 standard alphabetic, numeric and special character keys found in the ASCII character set. Figure 5-1 illustrates the keyboard layout.

Section 2 consists of an 11 key matrix which is identical to a standard adding machine keyboard with the addition of a decimal point (period). The keys in this section are duplicates of certain keys found in Section 1 and are provided to facilitate entry of large amounts of numeric data.

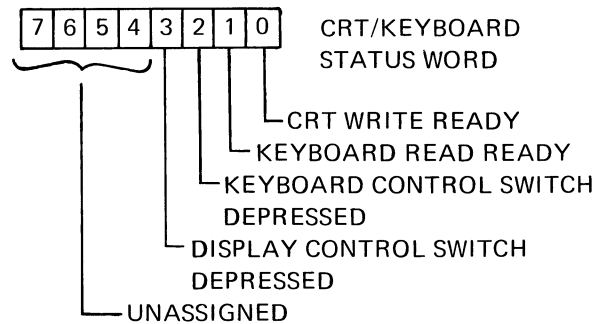
The keys in Section 3 are special function keys which exert control over the processor. Their names and associated functions are as follows:

RUN	Momentary contact switch, which when depressed, causes the processor to begin execution of the instruction located at the address in memory currently addressed by the program counter.
STOP	Momentary contact switch which, when depressed, causes instruction execution to halt at the completion of the current instruction. Care should be taken when using this switch, because any tape operation which may be in progress will be aborted.
KEYBOARD	Momentary contact switch which sets a status bit that may be tested at any time by the processor.
DISPLAY	Momentary contact switch with a function similar to that of KEYBOARD switch. Either one or both of these switches may be depressed.
RESTART	Momentary contact switch which causes the processor to halt, rewind the system or program tape mounted on Deck 1, load and execute the first record found on tape.

#### 5.2 OPERATION

The keyboard is addressed by the processor by loading the A-register with 341<sub>8</sub> and executing an EX ADR command. (The crt display also uses this address. Data transfers to the

processor are from the keyboard and transfers from the processor are to the display). Following the address sequence the c.r.t./keyboard status word can be loaded into the A-register by executing an INPUT instruction. Bit 1 of the A-register may be tested by the program to determine if a character is ready for transfer from the keyboard. Bits 2 and 3 will indicate if either the KEYBOARD or DISPLAY control switch is pressed.



The External Commands associated with the operation of the keyboard are as follows:

- a. EX BEEP. This command produces a 1500 Hertz tone for a duration of about 100 msec. The tone could be used as an error or ready signal to the keyboard operator.
- b. EX CLICK. This command produces an audible click which could be used to acknowledge receipt of a valid character when a key is depressed.
- c. EX COM1 (Command 1). Presents a control word contained in the A-register to the keyboard. Bit 5 of the control word controls the KEYBOARD switch light and bit 6 controls the DISPLAY switch light as follows:

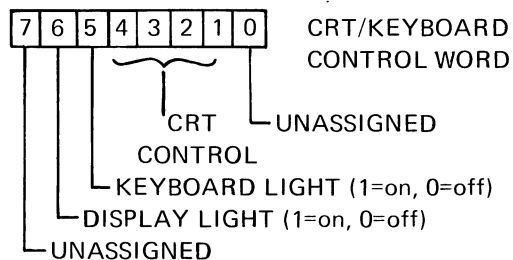


TABLE 5-1

KEYBOARD CODING (ASCII)

A-101	a -141	0-060	:	-072
B-102	b -142	1-061	;	-073
C-103	c -143	2-062	<	-074
D-104	d -144	3-063	=	-075
E-105	e -145	4-064	>	-076
F-106	f -146	5-065	?	-077
G-107	g -147	6-066	[	-133
H-110	h -150	7-067	~	-176
I -111	i -151	8-070	]	-135
J -112	j -152	9-071	^	-136
		Space-040	_	-137
K-113	k -153			
L-114	l -154	!-041	@	-100
M-115	m-155	"-042	{	-173
N-116	n -156	#-043	\	-134
O-117	o -157	\$-044	'	-140
P-120	p -160	%-045		-174
Q-121	q -161	&-046	}	-175
R-122	r -162	'-047	Enter	-015
S -123	s -163	(-050	Cancel	-030
T-124	t -164	) -051	Backspace	-010
U-125	u -165	*-052	Rubout (R.O.)	-177
V-126	v -166	+ -053		
W-127	w-167	, -054		
X-130	x -170	- -055		
Y-131	y -171	. -056		
Z-132	z -172	/ -057		

## PART 6

### CRT DISPLAY

#### 6.1 GENERAL DESCRIPTION

The display unit on the Datapoint 2200 consists of a CRT capable of displaying 12 lines of 80 characters each, a character generator, 960 cells of refresh memory (refresh rate 60 Hz), and a group of registers utilized to position the cursor. Maximum character transfer rate to the CRT is 60 characters per second.

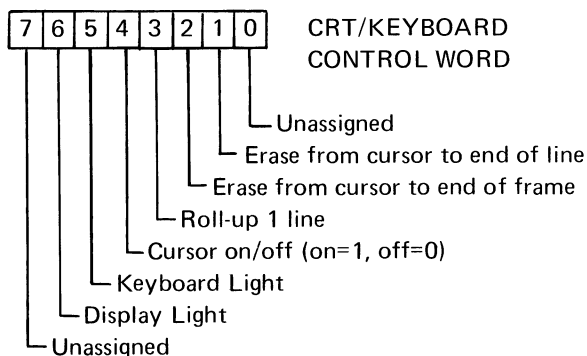
The character set utilized by the CRT display consists of the full ASCII set with both upper and lower case alphabets and all numeric and special characters.

#### 6.2 OPERATION

The CRT is addressed and status tested in the same manner as the keyboard (see paragraph 5.2). Bit 0 of the status word indicates that the CRT is ready to accept data or commands. Characters are transferred to the screen by loading the A-register with the character to be displayed and executing an EX WRITE. The character will be displayed at the current cursor location.

Control of the CRT is accomplished through the use of the three external commands - Command 1, Command 2, and Command 3. The functions performed by these commands are as follows:

- a. EX COM1 (Command 1) Transfers a control word contained in the A-register to the CRT. The applicable bit assignments and their functions are as follows:



The erase functions permit selective erasures on the screen by limiting erasures to those character positions following the current cursor position to the end of the line (or page).

The roll-up function causes all displayed characters (not the cursor) to move up one line. The top line on the screen is lost.

The cursor image may be turned on or off through the control word. The cursor position is the same in either case. The cursor image is automatically turned off whenever the processor is in the HALT state.

- b. EX COM2 (Command 2) Positions the cursor to the horizontal character slot designated by the contents of the A-register. Character position 0-79<sub>10</sub>(0-117<sub>g</sub>) are valid.
- c. EX COM3 (Command 3) Positions the cursor to the line designated by the contents of the A-register. Line number 0-11<sub>10</sub> (0-13<sub>g</sub>) are valid.

In order to write a new character, the cursor must occupy that character's position on the screen. After the character has been written, the cursor should then be moved to the next horizontal (or vertical) position desired. The CRT Write Ready status bit must be true before positioning the cursor or displaying a character.

Both the CRT and keyboard utilize the standard ASCII character set. (See Table 5-1). Any invalid character code will appear as a blank space on the CRT screen.

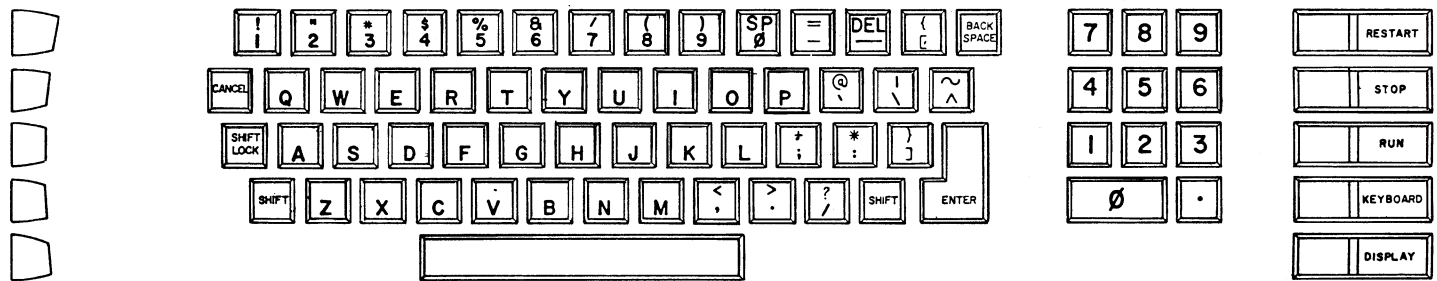


FIGURE 6-1  
KEYBOARD LAYOUT

## PART 7

### CASSETTE TAPES

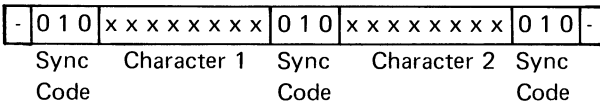
#### 7.1 GENERAL DESCRIPTION

The Datapoint 2200 contains two cassette tape recording devices for storage of programs and data. Since the hardware RESTART (section 5.1) uses the rear deck (number one), programs will typically be on it while data areas will be the front deck (number two). However, once the machine is initially loaded, either deck may be used for both purposes.

Data on the Tape is organized by record (of any length). Records are written and read at 350 eight-bit characters per second with a tape speed of approximately 7.5 inches per second. See Table 7.1 for a list of the physical specifications.

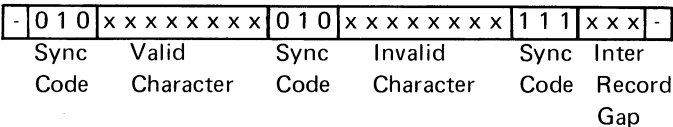
#### 7.2 OPERATION

Data is recorded or read in bit serial fashion on one track. Each eight bit character is framed by three sync bits on either side of the character:



The appearance of the correct sync code indicates that the character is valid. Any other sync code causes special action to be taken on data reads. Note that the sync codes are valid for tape motion in either direction so the tape may be read backwards although in the reverse direction the data bits will appear reverse d (bit 0 will be bit 7, 1 will be 6 etc.)

A record is a group of successive valid characters. An inter-record gap is indicated by the failure of the sync code to be zero one zero and all mark code. (ones):

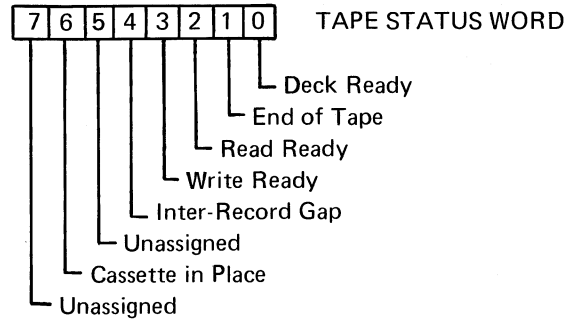


Only valid characters will be presented as data from the tape unit.

#### 7.3 STATUS

The cassette tape unit is addressed by the processor by loading the A-register with 360g and executing the EX ADR instruction. Following this sequence, the tape unit

status can be loaded into the A-register by executing an INPUT instruction. The bit assignments are as follows:



#### DECK READY

Deck ready will be set whenever the tape unit is ready to accept another command. (Only the TSTOP command should be issued if this bit is false). The tape will be stopped, a cassette in the selected deck and not wound to the clear leader at either end, and the head engaged when this bit is true. This bit should be checked after selecting a deck.

#### END OF TAPE

End of Tape indicates that the cassette has run onto leader (in either direction).

#### READ READY

Read Ready indicates that the selected deck has read another character.

#### WRITE READY

Write Ready indicates that the selected deck is ready to write another character.

#### INTER-RECORD GAP

Inter-Record Gap indicates the selected deck has come across an inter-record gap (invalid sync code).

#### CASSETTE IN PLACE

Cassette in Place indicates that a cassette is physically in place in the selected deck.

#### 7.4 CONTROL

When the cassette tape unit is addressed the following instructions will control the action of the tape:

- a. EX TSTOP causes any motion of either deck to be stopped, any read or write operations to be terminated. When everything has settled, the ready status bit will come true and operations may be resumed.
- b. EX DECK1 causes deck one (rear) to be the currently selected deck. Before commanding a deck selection, care should be taken that the currently selected deck has completed all operations.

- c. EX DECK2 causes deck two (front) to be the currently selected deck. Note the precaution in (b).
- d. EX RBK causes the currently selected deck to be set in forward motion and, after 70 msec, for the read circuitry to be enabled. The read ready status bit will come true upon appearance on the tape of the first valid character. Upon appearance of an invalid sync code, the inter-record gap status bit comes true and tape motion is automatically stopped. Note that this will happen only after at least one valid character has been found. Once the read ready status bit comes true, the character must be taken within 2.8 milliseconds or it will be overwritten with the next one. The tape read hardware double-buffers incoming characters to allow the 2.8 msec character availability.
- e. EX BSP is similar to EX RBK except that tape motion is in the reverse direction so the data bits will be reversed.
- f. EX SF is similar to EX RBK except the tape is not stopped upon appearance of an inter-record gap, and if allowed to continue will start to read the next record on the tape. In this case, the read ready status bit will come true again after the first character of the next record is read. Only an EX TSTOP will stop the motion initiated by EX SF.
- g. EX WBK causes the currently selected deck to be set in forward motion and for all status bits except the write ready to go false. A character must then be presented within 2.8 milliseconds (the first character will be accepted at once due to the buffering in the tape hardware and then there will be a pause while the tape comes up to speed), at which time the write ready will go false until the writing circuitry is ready to accept another character. An end of record is signalled to the hardware by withholding a character for a period of time longer than 2.8 milliseconds specified above. When this is done, the write ready will go false, an inter-record gap will be written, the tape motion will cease, and the deck ready status bit will come true again.
- h. EX REWIND causes the tape to be rewound to the beginning on the selected deck. Worst case rewind time is approximately 40 seconds.
- i. PUNCH TABS, on the Cassette Cartridge are used for "write protect" and "automatic restart". The punch tab on the left (as you face the terminal) inhibits the ability to write on tape, when punched. When the tab on the right is punched, it causes an automatic restart whenever a halt or power-up occurs.

**TABLE 7-1**

**TAPE UNIT PHYSICAL SPECIFICATIONS**

Density	47 characters/inch
Speed	7.5 ips
Recording Rate	350 c.p.s.
Capacity	130,000 characters (typical)
Start/Stop Time (Inter-Record Gap)	280 msec.
Start/Stop Distance (Inter-Record Gap)	2 inches
Rewind Speed	90 ips
Rewind Time (max 300 ft.)	40 sec.
Character Transfer Time	2.8 msec.

## PART 8

### COMMUNICATIONS ADAPTOR

#### 8.1 GENERAL DESCRIPTION

The 2200 Communications Adaptor is an external device, which when connected to the Datapoint 2200 Input/Output System permits asynchronous serial data interchange to other remote systems or devices.

The Communications Adaptor consists of three basic parts:

- a. The serial data transmitter and time base;
- b. The serial data receiver and time base; and
- c. The communications channel interface.

The communications channel interface may be one of four types:

- a. An EIA RS-232 type interface;
- b. An isolated high-level neutral or polar telegraph loop interface;
- c. A modem compatible with the Bell System 103 type modems;
- d. A modem compatible with the Bell System 202 type modems.

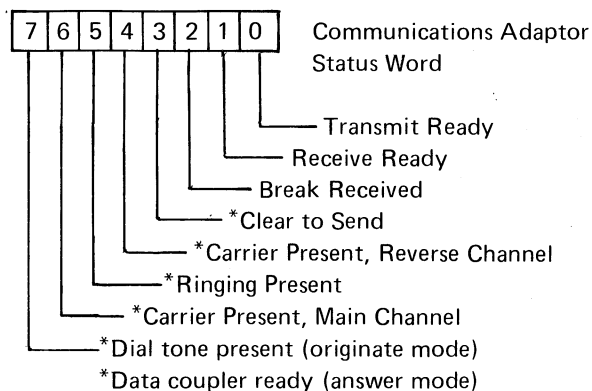
#### 8.2 OPERATION

The serial data transmitter and receiver are addressed at the same time (the address of the first used communications adaptor is 322g - see Table 4-2). Additional adaptors may be given previously unassigned addresses.

To set the bit rate desired for the transmitter time base two successive EX COM3 instructions are used to transfer two 8-bit masks from the A-register (See paragraph 8.6 for a discussion of time base mask words). For the receiver EX COM2 is used.

To set the character length for the transmitter and receiver an EX COM4 command is executed with a character length mask from the A-register (see paragraph 8.7 for a discussion of character length mask words).

The status of the communications adaptor is transmitted to the A-register with the following bit assignments:



\*Used with data set options.

#### Communications Adaptor Status Bit Description

##### Bit 0, Transmit Ready

The "true" condition of this bit indicates that the serial transmitter is ready to accept a new character for transmission. Should another write command be issued to the Communications Adaptor while this bit is "false", i.e. transmitter NOT ready, the previous character will be written over.

##### Bit 1, Receiver Ready

The Receive Ready bit, in the true state, indicates the presence of a new received character. A read command to the Communications Adaptor returns this bit to the false state. If a read command is not issued before another new character is received, the new character will replace the existing character and the status will remain true.

##### Bit 2, Break Received

The Break Received status bit simply indicates that the received data is in the "space" or "zero" condition for longer than one character time.

##### Bit 3, Clear to Send

The true state of Clear to Send status indicates that the data set (internal or external) is prepared to accept data for transmission. This bit has meaning only when an internal or external data set is in used.

#### Bit 4, Carrier Present - Reverse Channel

This status bit has significance only when operating half-duplex with either an internal or external 202 type data set (modem). The true condition indicates that the reverse (supervisory) channel carrier is being received.

#### Bit 5, Ringing Present

The true condition of Ringing Present indicates that the ringing of an incoming call has been detected. This bit has significance only when used with an internal or external (with proper options) data set.

#### Bit 6, Carrier Present - Main Channel

The true condition of this status bit indicates that the main channel carrier is being received. This status bit has meaning only when used with an internal or external data set.

#### Bit 7, (1) Dial Tone Present (Originate Mode)

##### (2) Data Coupler Ready (Answer Mode)

(1) When originating a call, the true condition of this status bit indicates that a dial tone is present and dialing may proceed; during dialing, the status will become false. Following dialing, and a 2 to 5 second delay, this bit will return to the true condition indicating connection to the telecommunication network (but does not indicate the called number has answered).

(2) When answering a call, the true condition of this status bit indicates that the data coupler is connected to the telecommunications network.

### 8.3 DATA OUTPUT

After addressing the communications adaptor transmission of each character is accomplished in the following manner:

- a. Input the status word and verify that status bit 0, Transmit Ready, is set to 1 indicating that the adaptor can accept another character.
- b. Load the A-register with the byte to be transmitted.
- c. Apply a write strobe (EX WRITE). Data present on the A-bus will be loaded into the data transmitter and data will be serially transmitted at the selected code length and bit rate.

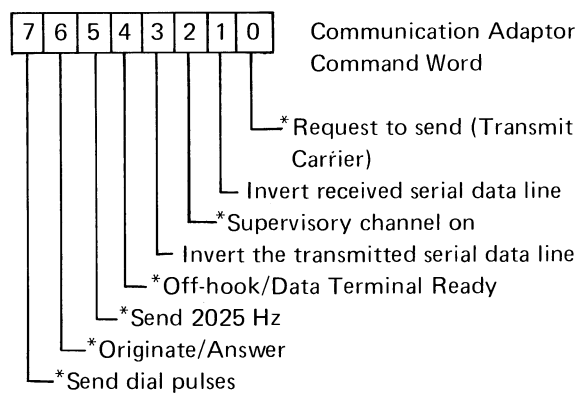
### 8.4 DATA INPUT

After addressing the communications adaptor, reception of each character is accomplished in the following manner:

- a. Input the status word and verify that status bit 1, Receiver Ready, is set to 1, indicating that a character has been received.
- b. Execute an EX DATA instruction.
- c. Execute an INPUT command, transferring the received character to the A-register.

### 8.5 COMMAND WORD

Control of the communications adaptor is accomplished through the use of a command word. The command word is transmitted to the adaptor by executing EX COM1.



\*Used with data set options

#### Communication Adaptor Command Word Description

##### Bit 0 - Request to Send

This command bit controls the transmit carrier of an internal or external data set. A "one" in this position turns on the transmit carrier and indicates to the data set that it must prepare for data transmission.

##### Bit 1 - Invert Received Serial Data Line

A "one" in this position permits data to be received normally when the received serial data line is inverted.



### Bit 2 - Supervisory Channel On

This command is used only with a 202 type modem in half-duplex operation. A "one" in this command indicates to the modem that the supervisory (or reverse) channel will be operative, transmit or receive.

### Bit 3 - Invert Transmitted Serial Data

A "one" in this command inverts the transmitted serial data.

### Bit 4 - Off-Hook

A "one" must be placed in this bit position any time a telecommunication call is to be originated or answered. This command allows connection to be made to the telecommunication network with an internal modem and a Bell System Data Access Arrangement. When using an external modem, this command provides "Data Terminal Ready" to the external modem, i.e., the system is prepared for on-line communications. This command is used only for the cases described above.

### Bit 5 - Send 2025

This command is used only with an internal 202 type modem, half-duplex operation and "answer" mode. The only use of this command is described as follows:

- 1.) following receipt of Ringing Present, Status Bit 5, the Off Hook Command, Command Bit 4, is set to a "one".
- 2.) Next, Status Bit 7, Data Coupler Ready must become "true".
- 3.) Send 2025 command must now be set to a "one" only for a period of 1/2 second to 3 seconds to inform the calling data set of our response.

### Bit 6 - Originate

This command is used only with internal data sets (modems). A "one" in this command instructs the modem that the system will originate a telecommunication call, A "zero" tells the modem the system is prepared to answer a telecommunication call.

### Bit 7 - Send Dial Pulses

This command is used only with internal data sets (modems) and is set to "one" only when dialing. Its use is described as follows:

- 1.) Off-Hook Command (Bit 4) is set to "one".
- 2.) Status Bit 7 - Dial Tone Present becomes "true".
- 3.) Bit 7, Bit 4 and Bit 3 (invert xmit), are now set to "one".
- 4.) When the last dial pulse is completely transmitted, Bit 7 and Bit 3 must be returned to "zero".

## 8.6 TIME BASE MASK WORDS

Both time base generators are programmed for their respective bit rates by the processor. Each time base is independently controlled to allow transmission and reception at different rates.

After addressing the communication interface, two eight-bit mask words are loaded into the time base registers to synthesize the selected bit rates. As each respective byte is presented, a corresponding EX COM2 instruction must be executed to load the receive time base and an EX COM3 instruction to load the transmit time base.

These two bytes are combined to form a 16 bit word which is placed in a holding register. A counter is then set to the value in the holding register. This counter is incremented at the rate of 153,600 Hz. Each time the counter overflows, i.e., goes from all one to all zeroes, a pulse is generated and the counter is reset to the value in the holding register. The time between pulses represents 1/2 clock period or 1/2 bit time. Given a bit rate (bps), the following formula can be used to determine the number N to be entered into the holding register:

$$N = 65,536 \cdot \left( \frac{76,800}{\text{bps}} \right)$$

This number N may then be converted to a 16 bit binary number and separated into the two 8-bit mask words.

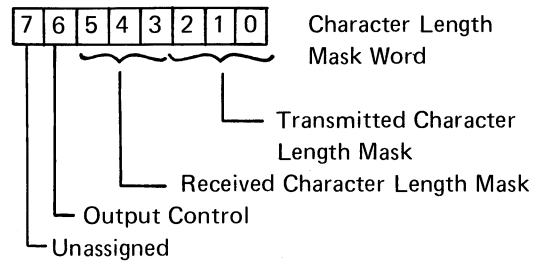
The octal codes for some of the more frequently used rates are listed below:

BIT RATE	1ST MASK WORD	2ND MASK WORD
1) 100*	375	000
2) 110	375	106
3) 220	376	243
4) 440	377	121
5) 150	376	000
6) 300	377	000
7) 600	377	200
8) 1200	377	300
9) 2400	377	340
10) 4800	377	360
11) 9600	377	370

\*(Dialing)

## 8.7 CHARACTER LENGTHS

Character lengths for the transmit and receive sections and its output control bit are determined by a character length mask word which is transmitted to the communications adaptor with an EX COM4 instruction.



The functions of the mask bits are given in the following tables:

**TABLE 8-1**  
**TRANSMITTED CHARACTER LENGTH MASK BITS**

MASK BIT POSITION 210	START UNITS	INFORMATION UNITS	STOP UNITS	CODE BIT POSITIONS 76543210
000	1	8	1	87654321
001	1	8	2	87654321
010	1	7	1	7654321
011	1	6	1	654321
100	1	5	1	54321
101	—	—	—	---
110	—	—	—	---
111	—	—	—	---

When codes having 5, 6, or 7 information units are to be transmitted, the remaining high-order bits in the character byte must be coded to "1".

When a two-unit stop pulse is required for characters having 5, 6, or 7 information bits, the next larger character length is used; the remaining high-order bits (all coded 1) form the stop pulses.

When received characters contain 5, 6 or 7 information bits, the remaining low-order bits (as shown above) must be disregarded.

One additional command bit, Bit 6 (output control) of the Character Length Mask Word, is used to control the EIA RS-232 Transmitted Data and the High-Level Keyer Transmitted Data. A "one" in this command bit enables serial data to be transmitted only to the EIA RS-232 output or to

**TABLE 8-2**  
**RECEIVED CHARACTER LENGTH MASK BITS**

<b>MASK BIT POSITION</b> <b>543</b>	<b>START UNITS</b>	<b>INFORMATION UNITS</b>	<b>STOP UNITS</b>	<b>CODE BIT POSITIONS</b> <b>76543210</b>
000	1	8	1 or more	87654321
001	1	8	1 or more	87654321
010	1	7	1 or more	7654321x
011	1	6	1 or more	654321xx
100	1	5	1 or more	54321xxx
101	—	—	—	---
110	—	—	—	---
111	—	—	—	---

the High Level Keyer. A "zero" in this command bit allows serial data to be transmitted only to an internal data set (modem).

### 8.8 INTERFACE CONNECTOR

This interface is provided through an Amphenol 17-10500-1 connector. Pin assignments are as follows:

<b>LEAD</b>	<b>FUNCTION</b>	<b>INPUT/OUTPUT</b>
1	Protective Ground	----
2	Protective Ground	----
3	OH (Off Hook)	Output
4	+25v	----
5	DA (Transmission Path Request)	Output
6	R (Ring Indicator)	Input
7	CCT (Data Coupler Ready)	Input
9	DT (4 wire)	} Direct Private Line Connection
10	DT (2 wire)	
11	DR (2 wire)	
12	DR (4 wire)	
23	Clear to Send (RS-232)	Input
24	Transmitted Data (RS-232)	Output
28	Signal Ground	----
29	Signal Ground	----
32	+5v	----
33	+5v	----
40	Request to Send (RS-232)	Output
41	Received Data (RS-232)	Input
42	Data Terminal Ready (RS-232)	Output

<b>LEAD</b>	<b>FUNCTION</b>	<b>INPUT/OUTPUT</b>
44	Supervisory Transmitted Data (RS-232)	Output
45	Data Carrier Detector (RS-232)	Input
46	Supervisory Received Data (RS-232)	Input
49	Clock for 3300P	Output
50	Transmit Bit Rate Clock	Output

### 8.9 HIGH LEVEL OPTION

Interface with telegraph-type current loops is provided with the high level option. This option provides for completely isolated electronic neutral/polar output relay and a completely isolated neutral/polar input relay. Loop voltage may be as high as 400 volts across the relay and as high as 1000 with respect to ground.

Loop resistance and power is not included with the option.

For further information, refer to the Datapoint 2200 Installation Manual.

### 8.10 103-DATA SET OPTION CHARACTERISTICS

The 103-Data Set option provides for full duplex data transmission for rates up to 300 bits per second with a signalling system that is compatible with the Bell System 103 series Dataphones. Connection to the common carrier lines would normally be made through a Bell System Access Arrangement type F-58118, CBT, or 1001B. Other connections are also possible where automatic dialing or answering is not required.

The data set may be placed in either the answer mode or originator mode through the use of bit 6 of the communications adaptor command word (see paragraph 8.5). Bit 6 is set to 0 for answer mode and 1 for originator mode. The request to send command bit (bit 0) is normally set to 1 with the 103 option to maintain the transmit carrier on.

Operation of the automatic dialing and answering features discussed in paragraph 8.12 and 8.13.

Table 8-3 provides a summary of characteristics of the 103 Data Set option.

**TABLE 8-3**

**103 DATA SET OPTION CHARACTERISTICS**

Originate Mode			
Carrier Frequencies:	Transmit:	Mark: 1270 Hz	Space: 1070 Hz
	Receive:	Mark: 2225 Hz	Space: 2025 Hz
Answer Mode			
Carrier Frequencies:	Transmit:	Mark: 2225 Hz	Space: 2025 Hz
	Receive:	Mark: 1270 Hz	Space: 1070 Hz
Keying Rate:	Up to 300 bits per second		
Transmit Level:	0 to -10 dbm.		
Impedance:	600 ohms nominal		
Receive Sensitivity:	+5 to -30 dbm.		

**8.11 202-DATA SET OPTION**

The 202 Data Set option provides for either full or half duplex data transmission for rates up to 1200 bits per second (1800 bits per second on conditioned private lines). This option is compatible with Bell System 202 series Dataphones (including supervisory channel operation) and in addition provides a 150 bit per second supervisory channel when used with another Datapoint 2200 Data Set option of the same type. Connection may be directly to private lines or to common carrier lines through a Bell System Access Assignment type F-58118, CBT, or 1001B where access to the telephone switched network is desired.

Operation of the automatic dialing and answering are discussed in paragraphs 8.12 and 8.13 respectively.

Table 8-4 provides a summary of characteristics of the 202 Data Set option.

**TABLE 8-4**

**202 DATA SET OPTION CHARACTERISTICS**

Main Channel Frequencies:	Mark: 1200 Hz
	Space: 2200 Hz
	Soft Turn-Off: 880 Hz
Supervisory Channel Frequencies:	Mark: 387 Hz
	Space: 470 Hz
	Soft Turn-Off: 330 Hz
Special Command Frequency:	2025 Hz
Main Channel Keying Rate:	Up to 1200 baud (1800 baud on conditioned private lines.)
Supervisory Channel Keying Rate:	Up to 150 baud
Transmit Level:	0 to -10 dbm
Impedance:	600 ohms nominal
Receive Sensitivity:	+5 to -30 dbm

**8.12 AUTOMATIC DIALING OPERATION**

When using the Datapoint 103 or 202 data set options with the Bell System Access Arrangement type F-58118, CBT, or 1001B it is possible to automatically originate a call into the telephone switches network. The procedure for this function is as follows:

- a. Set bits 4 and 6 of the communications adaptor command word to 1 to provide an off-hook signal to the telephone network and to prepare the modem for originate operation.
- b. Test bit 7 of the communications adaptor status word for a 1 indicating dial tone present.
- c. Set the transmitter time base to 100 bits per second (see paragraph 8.6).
- d. Set the character length mask word to all zeros (ten bit length-see paragraph 8.7).
- e. Set bits 3 and 7 of the command word to 1 thus inverting the serial transmitter output and transferring this output to the dial pulse keyer.
- f. Sequentially transmit the octal byte 360 for each dial pulse required for each number (see paragraph 8.3-Data Output).
- g. Program approximately 1 second delay between each number and at the end of the last number transmitted.

- h. Re-establish the correct code length and bit rate for data transmission and set command word bits 3 and 7 to zero to restore the normal transmitter output.

### **8.13 AUTOMATIC ANSWERING OPERATION**

When using the Datapoint 103 or 202 data set options with the Bell System Access Arrangement type F-58118, CBT, or 1001B it is possible to automatically answer a call from the telephone switched network.

Ringling is detected simply by testing bit 5 of the Communications Adaptor Status Word, Response to ringing would be to set bit 4 of the Communications Adaptor Command Word to 1 to provide an off-hook signal to the telephone network.

If the 103 Option is used Command Word bit 0 is set to 1 and bit 6 is set to 0 turning on the transmit carrier and selecting the answer-mode carrier frequencies.

If the 202 Option is used bit 5 of the command word is set to 1 for 1/2 to 3 seconds to transmit a 2025 Hz tone to disable echo suppressors and to inform the calling data set of out sequence in the telephone network, after which normal data transmissions occurs.

SECTION 2

DATAPOINT 2200

OPERATING SYSTEM

## SECTION 2

### THE OPERATING SYSTEM

The operating system is a conversational mode program for the Datapoint 2200 user to have a means to catalog, load, debug and run user programs and to provide other utilities important to the use of the 2200. All other programs discussed in the "Programmers Manual" such as the Program Editor and the Assembler are programs which the user may catalog onto an operating system tape and call into use as required.

The operating system itself is a relatively long program which is generally overlayed when user programs are called in from tape (unless they are less than 2K bytes in length and properly located). However, a family of resident utility routines is loaded with the operating system that may be used by user programs to simplify frequently used functions such as reading from the keyboard, writing to the CRT screen, reading and writing tape records, etc. The detailed use of the routines and the makeup of the operating system are described in Section 5—Advanced operating system command and subroutine usage.

Section 1 will describe the command language for the operating system and does not require any particular programming skills.

#### Start-Up Procedures

When power is first applied to the Datapoint 2200, it is incapable of performing any useful function except to load a block of data from the rear tape cassette deck into the processor's memory and transferring control to it. In the operating system, this first block of data is called a **LOADER** and when control is transferred to it, it proceeds to the first check itself to see if it was loaded properly and then to load the next file on the same tape which is the rest of the operating system program. This process can be executed at any time (assuming a proper program tape is in the rear deck) by pressing the **RESTART** key on the right hand side of the 2200.

This first block of data can be used to load programs other than the operating system and is generally useful for all applications of the 2200. In order to use the operating system, a full 8K bytes of memory must be provided in the 2200, but the loader alone can be used with any size memory.

When an operating system program tape is loaded, the first thing that appears on the screen is:

```
COMPUTER TERMINAL OPERATING SYSTEM
```

### READY

At this point, any of the operating system commands discussed below may be typed into the 2200.

Each command has to be in the form of a word followed by the **ENTER** key, or a word followed by a space, or a dash and a name, or a modifier. Each of the following are valid operating system commands:

```
CATALOG  
REPLACE RST4  
REPLACE-RST4  
RUN*
```

Only the first three letters of a command are actually decoded so that the following are valid commands:

```
CAT  
REP RST4
```

If an invalid command is typed the system responds with:

```
WHAT?
```

#### Operating System Commands:

##### CATALOG

The **CATALOG** command will print out a list of programs currently available on that particular operating system tape. Up to 14 programs may be cataloged on a tape under this system and the order that they appear on the tape is the order their names appear on the screen when the catalog command is given. A typical catalog response might be:

```
CATALOG  
RST4 CODER BANDIT ANNUIT  
READY
```

Which would indicate that four programs are logged onto the tape.

##### NAME

The **NAME** command allows any program in the system to have its name changed. For example, the program named **BANDIT** in the above example can be changed to **GAME** with the command:

```
NAME-BANDIT, GAME
```

Names may have any combination of letters and digits up to six characters in length and beginning with a letter. All of the following are valid program names:

BOB  
R12345  
A  
NAME

## RUN

The RUN command causes the operating system to position the cassette tape to the program named in the run command, load the program into the 2200 memory and transfer control to it. The program being run may overlay part or all of the operating system. If it does, returning to the operating system can only be done by reloading it. This can be done by using the restart switch or by program control (See Section 5).

A run command would appear like this:

RUN-BANDIT

A program that has not been logged onto the system tape may be run by placing an assembled form of the program on the front deck and typing:

RUN\*

## IN

The IN command causes a program to be cataloged onto the operating system tape. The program must be assembled and the assembled program tape placed on the front deck. The IN command is typed giving the name to be assigned to the program as shown here:

IN RST4

The operating system tape (on the rear deck) will position itself to the end of its program library and will then copy the program from the front deck into place and add its name and position to the system catalog.

## DELETE

The DELETE command causes the program named in the command to be removed from the system library. Unless the program being deleted happens to be the last program in the library, a SCRATCH tape will be required in the front deck to copy part of the library out and back to the system tape to CLOSE-UP the space. In this case, when the command is entered, the system will write a message back:

FRONT TAPE SCRATCH?

Then the processor will stop. The stop key on the right hand side of the keyboard will be lighted. If there is a tape in the front deck that can be recorded on, press the run key on the right hand side and the system will proceed to delete the named program for you.

This will generally take a little time. When the system is through it will write READY on the screen.

## REPLACE

The REPLACE command allows a program already in the system catalog to be deleted and a new program to be put in its place in the same order on the tape. The new program does not have to be the same length as the old one. Again, this command takes time to execute due to the amount of "shuffling" of tape files to get everything in place.

## AUTO

The AUTO command allows one of the programs in the system library to be marked for an automatic RUN whenever the operating system is restarted. Once a program has been named in an auto command, it may be cleared by typing a MANUAL command. The automatic feature may be overridden during a restart by holding down the KEYBOARD key on the right hand side of the keyboard.

The automatic program calling feature is particularly valuable when the program is to be run in an unattended situation. If the knock-out tab on the back of the operating system tape is removed then whenever power is reapplied or the processor is halted for any reason (including a programming halt), an automatic restart is executed and, of course, the program named in the auto command is reloaded and given control.

If the command AUTO is typed without a name then the system will respond with:

NAME REQUIRED

If a program is already named in an auto command then the system will respond with:

AUTO SET TO (PROGRAM NAME)

## MANUAL

The MANUAL command will delete any program from the auto mode.

## OUT

The OUT command causes any named program to be copied to the front tape causing any data already on the front tape to be lost. This copy may then be cataloged onto some other



system tape or be saved for some future use. When using the OUT command the system will write FRONT TAPE SCRATCH? onto the screen and the processor will halt. If you have a usable tape in the front deck, then press the run key on the right hand side of the keyboard and the system will continue.

If the command OUT \$ is typed, then the entire operating system including the library is copied to the front deck. If the command OUT \* is typed, then the loader and the library is copied to the front deck but the operating system is deleted. This permits a program or family of programs to be used on a Datapoint 2200 with less than the full 8K of memory. (See Section 5 for details).

## **PREP**

The PREP command causes the tape in the front deck to be rewound and a NULL file to be written at the beginning of the tape, effectively ERASING the tape and making it ready for use by operating system tape routines. The operating system commands that write on the front deck execute the PREP function automatically, however, USER PROGRAMS WRITING DATA TO THE TAPE MAY REQUIRE THE TAPE TO BE "PREPped" IN ADVANCE.

## **HEX**

The HEX command allows programs generated on other machines that follow a specified hexadecimal format, to be loaded into the Datapoint 2200. Users will not normally be concerned with the HEX command. (See Section 5 for details).

## **DEBUG**

The DEBUG command transfers program control to a small sub-program within the operating system that is used as a programming aid to debug and modify programs that are loaded into the 2200 memory. The debug sub-program allows you to write the contents of memory locations to the screen, modify memory locations, load programs into memory from the library or from the front deck, transfer control to parts of a program in memory and TRAP register values upon return to the debug program. Instructions on the use of the debug sub-program are given in Section 5.

SECTION 3

DATAPOINT 2200

SOURCE CODE EDITOR

## SECTION 3

### ASSEMBLER SOURCE CODE EDITOR

The assembler source code editor program provides for the preparation and editing of source data tapes in an assembler compatible format.

The editor program is called by the Computer Terminal Operating System (CTOS), if it is cataloged therein, by typing the following command:

#### STEP 1

##### RUN EDIT

#### STEP 2

a) When the editor has been loaded, the following message will appear on the screen:

##### Compressed Source Code Editor

##### Edit (E) or Convert (C)?

Type 'C' only if you have a tape generated by EDIT (1.1). This version of the Editor generated unblocked string records (using SSFW\$). The present version generates "compressed source" records (using SNFW\$).

If conversion is required, type 'C'. The following message will appear:

##### PLACE SYMBOLIC TAPE IN FRONT DECK- WHEN READY PRESS RUN

Place the old tape in the front deck and push RUN. A converted file will be generated in the scratch area of the CTOS tape. When the 1st pass of the conversion is complete, the following message will appear:

##### PLACE SCRATCH TAPE IN FRONT DECK- WHEN READY PRESS RUN

To protect your original tape, use a new tape to record the new compressed source code. Place the new tape in the front deck and press run.

The conversion process may be repeated for several tapes.

The Editor will now accept only compressed source tapes.

b) Type 'E' to EDIT, this message will appear on screen:

##### TYPE (:NEW, :OLD, OR :DUPLICATE)?

The meaning of the possible responses are as follows:

(NOTE:) All commands to the editor must begin with colons.

**:NEW** Indicates that the tape on the front deck is to be treated as a new source data tape. Any old data on this tape will be written over by the editor program.

**:OLD** Indicates that the front deck contains a tape with assembler source data on it. The operator will be allowed to edit this tape, changing only those lines which the operator specifies.

**:DUPLICATE** Indicates to the editor that it should copy the contents of the scratch file to the source data tape. This provides copies of a single source data file. It also provides for recovery capability should a system failure occur during the editing process. See Step 4, Recovery Procedures.

#### STEP 3

a) If the response in Step 2 was :NEW the following question will appear on the screen:

##### NEW NAME?

The operator may now enter a character string of up to 40 characters. This text will become the first record of the source data tape and will appear as part of a comment line in the assembled text.

b) If the response in Step 2 was :OLD the editor will read the header record from the source data tape (front deck) and display that header on the screen in the following format:

##### OLD NAME IS XXXX XXXX...XXXX

This old name header is also written to the scratch file and is retained as the header record for the source tape.

c) If the response in Step 2b was :DUPLICATE the editor will return to Step 2b after the copying operation has been completed. See Step 4, Recovery Procedures.

After a) or b) above has been completed, a "READY" message will be written on the screen and the cursor will appear at the beginning of the bottom line. The editor is now ready to accept new text data or a command. In order to enter text, simply type the desired text. Upon pressing the enter key, the typed line will be rolled up on the screen one line and the cursor will reappear on the bottom line and accept another text line or a command. When a line of text rolls off

the top of the screen, it is written to the scratch file. Lines will be written to the scratch file in the same order as they appear on the screen, the top line being first in the file.

Commands which can be entered, and their respective functions, are listed below. Command lines are distinguished from text lines by a leading colon; therefore, it is necessary to begin any command by typing a leading colon.

The "pointed line" referred to by some command descriptions below, is the line currently being pointed to by the visible pointer at the left side of the screen (col. 0). The "point's" vertical position is controlled by the keyboard/display keys. Specifically, pressing KEYBOARD causes the pointer to move up one line, pressing DISPLAY moves it down one line. Motion in either direction is circular around the screen. (It wraps around).

Commands may be entered from the bottom line only and must be preceded by a colon (:).

The functions available and their respective descriptions are listed below:

**MANUAL SEARCH** (Not a typed command).  
MANUAL SEARCH is like a continuing find of the very next line. That is, searching continues line at a time, with the next new line going to the eleventh line and the screen rolling up. It is unlike a find in that the screen isn't cleared for each new line acquired. It is useful for manually scanning through the data to bring to the screen and, therefore, into a position to edit the data of interest. MANUAL SEARCH is achieved by holding down the KEYBOARD and DISPLAY keys simultaneously. While held down, the search will proceed until end of file at which point the keys will become inoperable.

**:FIND <TEXT>** Where <text> represents n characters of text data. The editor searches the source data tape for the first match with the desired n-<text> characters. When the desired line is found it is displayed on the bottom line of the screen. The search is circular through the data files. If no match is found the text which occupied the bottom line at the time the command was issued is restored. Leading blanks on data-lines are ignored during the search. A FIND or EOF in progress can be stopped by the manual search operation.

**:COPY or :COP** The pointed line is copied to the bottom line and is simultaneously deleted from its previous location. The cursor will occupy the pointed line and will accept new text at this position. Striking the <enter> key returns the cursor to the bottom line and rolls the screen up one line.

**:DELETE or :DEL** The pointed line is deleted from the screen and the scratch area. The cursor will occupy the pointed line and will accept new text at this position. Striking the <enter> key returns the cursor to the bottom line. The screen is not rolled up.

**:INSERT or :INS** The pointed line and all lines above it are rolled up one line. The cursor will now occupy the blank line created and is ready to accept the new text. Striking the <enter> key returns the cursor to the bottom line.

**:EOF** The editor will search the source data tape for an end-of-file. Upon finding it, the last 11 text lines are on the screen and the cursor occupies the bottom line ready to accept new text or commands. A FIND or EOF in progress can be stopped by the manual search operation.

**:END** Causes the screen and the remaining source data to be copied to the scratch file. The scratch file is then copied to the source data tape.

**:END/DEL** The same as :END except that all data on the source tape which follows the current screen data is deleted.

**:SCRATCH or :SCR** Cause all lines between the top of the screen and the pointer, inclusive, to be deleted from the screen and the scratch area. The cursor will occupy the pointed line and will accept new text at this position. Striking the <enter> key returns the cursor to the bottom line. The screen is not rolled up.

#### STEP 4

Recovery procedures using the :DUPLICATE command can be implemented should the edit program be aborted during its execution without benefit of having completed all necessary copying and end of file writing.

Causes of difficulties which could require such action are:

1. Power failure during execution of the program
2. Turning off the power execution of the program
3. Striking the restart switch during execution
4. Encountering unrecoverable tape errors during execution
5. Removal of the tape cassette during execution

The EDITOR edits from one tape [the "Source Tape," containing old data] through the screen to the "scratch tape." The identity of the "source" and "scratch" tapes or files are reversed each time the current "source" tape reaches a file marker indicating end of file. Therefore, if an even number of passes have been completed, the updated file is on the front deck. If an odd number of passes have been made, the CTOS scratch file, file 40g, is the updated file. New programs are written on the front deck.

Understanding the activity of the data with respect to the tapes will allow the operator to determine the course of action should difficulty occur.

Should the operator encounter difficulty and be rather vague as to which action to take, it is recommended that the operator choose the most valuable data tape by the following:

When unrecoverable difficulty arises, remove the data tape and replace it with a new tape which will be written over by the scratch area of the operating system tape.

Run the edit program and type :DUPLICATE in response to the original questions.

The scratch area will be deposited on the new tape.

Then the operator can manually search through the two tapes as separate data tapes and make a decision as to the most valuable to keep.

Special situations:

It is possible that the scratch area has the most valuable data on it but it is missing any file termination record. This can occur when the program was interrupted during data entry and the previous scratch-to-data tape copy (if any) was so far back as to render the scratch area the only desirable data. Should this occur, the copy will proceed from scratch to source until it runs out of data, in which case the operating system will encounter garbage on the tape. It will then write end-of-file marks on the data tape at that point. This will give a clean data tape suitable for beginning again. A note of caution: it is usually advisable, when attempting to recover valuable data, to use the two data tape approach and visually compare the two to make a value judgment as to which has the most desirable data.

SECTION 4

DATAPOINT 2200

ASSEMBLER

## SECTION 4

### THE ASSEMBLER

The 2200 assembly system consists of the ASSEMBLER, EDITOR and the OPERATING SYSTEM.

The ASSEMBLER generates a block of absolute object code which can be loaded by the operating system loader and cataloged by the operating system. It generates the object code from the symbolic source code which was generated by the editor.

The ASSEMBLER makes two passes over the source code.

The first pass generates a symbol table from the labels in the source code and checks for certain error conditions, primarily syntax and form. The symbol table is maintained in memory.

The second pass generates the program listing and the object code on the tape. It also produces further diagnostics of a more subtle nature.

Basically, the ASSEMBLER is a program that assigns numerical values to symbols and outputs these values upon input of the associated symbols. Symbols in certain fields have pre-assigned values such as the opcode mnemonics. The value assigned to an instruction mnemonic is the binary bit configuration recognized by the 2200 processor for that instruction.

For example, the following instruction mnemonics have the following octal values:

MNEMONIC	VALUE
ADB	201
RETURN	007
SUB	221

Symbols in fields other than the opcode field may be defined by the user. Pre-defined and user-defined symbols are kept separately by the ASSEMBLER so that the user may define symbols that are the same as the pre-defined symbols without encountering any difficulties.

Along with relating symbols with numbers, another major function of the ASSEMBLER is to enable one to reference a symbol that is defined later in the program. This is called FORWARD REFERENCING, and may be handled in a variety of ways. One of the simplest is to look at the source code twice. The first time determines the definitions of all the symbols and the second time uses the symbols to produce the object code. Each "look" at the source code is called a "PASS". Therefore, we end up with a two pass assembly process.

### Statements

A 2200 assembly code statement consists of a label field, an instruction field, an expression field and a comment field. An example:

```
   1   2   3   4  
LABEL1 JTC  START  THIS IS THE COMMENT FIELD
```

Field 1 is the label field

Field 2 is the instruction field

Field 3 is the expression or operand field

Field 4 is the comment field

The 2200 editor provides automatic formatting so that the fields always are justified to begin in a certain column with tabbing to that field automatic. However, the ASSEMBLER only requires the following:

A non-space in the first column means that the first field is a label, except for a leading period which designates the entire line as a comment line.

A space in the first field means a null label and the first field is an instruction.

Scanning proceeds from left to right with one or more spaces serving as field delimiters.

Terminating fields by other than a space or a line termination will result in E-flags during the assembly.

THE LABEL FIELD may consist of up to 6 characters. An excess of 6 characters will be truncated. The first character may be any alphabetic character or a \$ sign. The other characters may be any alphanumeric character or a \$ sign. For example:

LEGAL	ILLEGAL
LABEL1	1LABEL (starts with numeric)
LABEL2	LABEL* (non-alphanumeric character)
LABEL\$	LABEL. (non-alphanumeric character)
L1B2L3	

THE INSTRUCTION FIELD may be any of the instruction mnemonics listed in the Datapoint 2200 Reference Manual, compound instruction (described later) or assembler directives.

The Instruction Field may be from two to four characters. However, only the first three are scanned and consequently the user may abbreviate. For example:

LEGAL	ILLEGAL
CALL	CALL2 (instructions have no more than four characters or numeric characters in the field)
JTZ	
SET	
TP	

THE EXPRESSION or OPERAND FIELD consists of any number of strings, numbers or symbols with operators between them. If a space or line end terminates a number or a symbol, the expression is assumed to be ended. Numbers are assumed to be decimal (base 10) unless they have one or more leading zeros, in which case they are taken to be octal. That is, 123 is 123 decimal, whereas 0123 or 00123 (the octal number 123) is really 83 decimal.

String quantities are delimited (preceded and followed) by apostrophies. In expressions, only the last character of a string is used if more than one appears. If a string were to be added to a number, only the last character of the string would be added. The character value is the ASCII binary number with the parity bit always a zero. A null string is legal ( ' ' ) and results in a zero value. The forcing character, #, is used in strings to indicate that the next character should be taken as ASCII no matter what it is. This is useful for getting the characters ( ' ) and ( # ) themselves into the string. For example:

' # ' # ' # ' is the character string ' # '

There are three operators allowed in the expressions:

1. + This means addition
2. - This means subtraction
3. >8 This means shift right by 8. Use this to get the MSP of an expression.

Expressions are evaluated from left to right and all operations are assumed to have the same priority.

The operand or expression is a symbolic expression which is evaluated at assembly time and the value is used in whatever manner is required by the opcode.

THE COMMENT FIELD begins immediately after the first delimiter space after the operand. The comment field may have any character including punctuation within it. It is terminated by the end of the line which was written by the editor. Comments may take over the entire line, in which case that line must begin with a period.

ASSEMBLER DIRECTIVES are available for setting label and location counter values to other than the normal sequential location assignments and for defining constants. There are seven:

1. EQU EQUALS. Sets the value of the label on the statement to the value of the operand expression.
2. SET SET. Changes the value of the location counter to the value of the operand expression.
3. SK SKIP. Increments the value of the location counter by the value of the operand expression.
4. TP TABULATE PAGE. Increments the value of the location counter until it is a multiple of 256. This is useful for minimizing execution time and for blocking out data areas addressable by single precision.
5. DC DEFINE CONSTANT. Generates eight bit object words from one or more expressions or strings following the opcode. If the expression is terminated by a space, the DC directive returns control to the main assembly process loop which obtains another instruction. If it is terminated by a comma, another expression or string is looked for. Another special exception is made for string items found in the DC directive. All the characters of a string item are significant and as many words as necessary are generated to accommodate all the characters of the given string. Again, a comma is looked for after the closing apostrophe in a string item to see if more expressions follow. This special string item is in effect only if the expression opened with an apostrophe. String items in expressions still have only one character of significance. For example:

DC 1,2+3,2+'A','ABC'

generates the following octal values:

1,5,103,101,102,103

6. DA DEFINE ADDRESS. Generates a two byte constant which is the address, LSP first, of the expression.
7. RP REPEAT. Will cause the following line to be processed, the number of times, indicated by the operand value. For example:

RP 5  
LDA 0123

would produce the same code as:



LDA 0123  
LDA 0123  
LDA 0123  
LDA 0123  
LDA 0123

### NOTE

Repeated statements which have a label on them will result in multiple definition of that label and all that entails, including the "D" error flags.

FORWARD referencing in the expression field in assembler directives only is not permitted.

8. END      END. Indicates that there is no more input data to be processed and that the ASSEMBLER should complete generating the output. The operand field has special significance in the END statement. The value of the expression in the operand of the END statement is the starting value of the execution of the program. That is the starting address. This is, of course, optional. When no operand is specified, the results are indeterminate. It should only be left vacant when the program is to be loaded without direct transfer of execution to the program such as an overlay.

Compound instructions are instructions which directly result in the assembly producing a sequence of source code. In this case, the 2200 ASSEMBLER has two: The HL instruction and the DE instruction:

1. HL      The HL compound instruction generates the LABEL      load H-REGISTER and load L-REGISTER instruction necessary to place the address of the label LABEL in the H-REGISTER and L-REGISTER properly so that the load to and from memory will operate to that address. In doing the HL, it loads the most significant byte of the value of LABEL into the H-REGISTER and the least into L.
2. DE      The same as with HL except loads into the LABEL      D and E registers.

THE ERROR FLAGS produced by the 2200 ASSEMBLER are as follows:

The error flags can occur during either pass of the ASSEMBLER in response to bad statements.

They are:

1. D      The D flag means DIFFERENT DEFINITION. It is flagged if the label has been redefined to

a different value during the assembly. In that case, it has the second value.

2. I      The I flag means INSTRUCTION MNEMONIC UNKNOWN. The instruction was not an accepted instruction in which case a zero is inserted for this instruction.
3. E      The E flag means that an error has occurred in an expression or some unrecognizable character appeared in the wrong place. In this case a zero is substituted for the expression or in whatever was unrecognizable.
4. U      The U flag means UNDEFINED LABEL. It is used whenever a label is referenced and is not defined. This can occur in pass 1 when an assembly directive is operating on an expression containing a forward reference.

EXTERNAL COMMANDS & REFERENCES can be taken care of in two ways:

1. Directly produce the numeric value in the expression field corresponding to the reference external address (such as an operating system subroutine resident in memory) or the external command operand such as EX 1 instead of writing EX ADR.
2. Equating labels to these referenced locations using the EQU assembler directive and then referencing the labels. This can be done for external references to operating system subroutines by duplicating the operating system subroutine entry point label in your program and equating it to that address. i.e. instead of:

```
CALL 017000
```

to get the operating system keyboard string input routine, a more meaningful listing can be obtained if, at the beginning of the program, this was entered:

```
KEYIN$ EQU 017000
```

and then all references to this routine can be this way:

```
CALL KEYIN$
```

The same is true of the external commands used in the 2200. Rather than say:

```
EX 1
```

it is more meaningful to say:

```
EX ADR
```

Since it is an external command address that is desired.

A list of the external commands and the operands which the ASSEMBLER incorporates into the proper EX coding are below.

The ASSEMBLER treats external command labels differently to produce the octal command byte. For the commands, the operands are as follows:

ADR	1
STATUS	2
DATA	3
WRITE	4
COM1	5
COM2	6
COM3	7
COM4	8
UNUSED	9
UNUSED	10
UNUSED	11
UNUSED	12
BEEP	13
CLICK	14
DECK 1	15
DECK 2	16
RBK	17
WBK	18
UNUSED	19
BSP	20
SF	21
SB	22
REWIND	23
TSTOP	24

It is recommended that for those external commands used, the EQU to the table number is done at the start of the source program and then the external command references are done to the label.

## Operating The Assembler

The ASSEMBLER must have a symbolic source tape generated by the 2200 editor.

Place this tape in the front tape deck.

Run the ASSEMBLER.

It will ask for printer speed. For the Datapoint 3300P, state 300. For a model 33 or 35 Teletype, state 110. For a model 37 Teletype, state 150. For no printer or no listing desired, state 0.

The source deck will rewind and begin to read in.

At the end of the first pass the ASSEMBLER asks if the second pass should proceed. It only requires a YES or NO.

This is a convenience, since many times many errors will be uncovered by the ASSEMBLER already after the first pass and the user will desire to correct those errors before proceeding to the second pass and the listing.

If the second pass begins, the tape will rewind and begin accepting data again from the source tape, printing the listing and writing the object file on the scratch area of the rear tape.

When the tape has reached the end of the source the second time, assembly is complete and it only needs to copy the object code block on the rear tape to the area on the front tape just after the source code. This results in the rear tape being backspaced to the beginning of the block of code and then copying proceeding forward reading a block from the rear deck and writing it on the front deck.

At the end of ASSEMBLY, the operating system will be reloaded and come up running.

The front tape can be loaded into the machine to test using the operating system command RUN\*, inputted into operating system catalog or loaded using the Debug program by using the F command.

SECTION 5

DATAPOINT 2200

ADVANCED OPERATING SYSTEM COMMAND and SUBROUTINE USAGE

## SECTION 5

### 1. INTRODUCTION

The primary function of CTOS is to provide the user with an easily accessible data environment which will greatly facilitate program generation. This function is fulfilled through the use of a file handling system which is available both directly from the keyboard in the form of system commands and through program calls to file handling input/output subroutines. Note that the keyboard facility deals mainly with the system (rear) tape (using the data (front) tape mainly for input/output and scratch space) but that the program routines are generalized to allow use of either tape.

#### 1.1 KEYBOARD FACILITIES

The keyboard accessible facility allows the user to fetch and execute object files, which may be either system packages, such as the editor and assembler, or files the user has generated with either the assembler or other code generating programs. This facility also allows the user to create new files, alter or delete old ones, or perform certain utility functions. The system tracks the files on the system tape in a symbolic catalog which may be manipulated by the operator at the keyboard or used in program linking.

#### 1.2 PROGRAM FACILITIES

The program routines perform basic operations such as reading and writing records with all parity checking and generation handled for the user. Other operations such as positioning to the beginning or end of a file, backspacing over records, or rewinding the tape are also provided. Parameterization is handled in a generalized way to make subroutine usage easy and consistent.

#### 1.3 PHYSICAL LAYOUT

The memory layout of the operating system is shown below. The OS FILE HANDLER is the program accessible facility mentioned above while the OS COMMAND HANDLER is the keyboard accessible facility. Note that only 017400 and up need be in memory if only the symbol linker (which calls in an overlay by name so that its physical file number may be changed without having to rewrite the program calling in the overlay) is to be used, only 016200 and up need be in memory if only the debugging tool is to be used, and only 014000 and up need to be in memory if the keyboard facilities are not to be used (of course, 0-0777 is always reserved by the system). Also note that the user may load a program designed to fit into a 2K machine without overlaying any part of the full operating system.

### CTOS MEMORY USAGE MAP:

SYMBOLIC LINKER	017777
CATALOG	017600
KEYBOARD DISPLAY	017400
DEBUG	017000
	016200
	016200
OS FILE HANDLER	014000
OS BOOTBLOCK COPY	013000
OS COMMAND HANDLER	05000
2K UNUSED	01000
LOADER	0

### 2. THE LOADER

The loader is the heart of CTOS. It enables other programs to load files from the tapes into memory without the tape having to be at the beginning of the desired file and provides extensive error protection. It is the routine used by the bootstrap mechanism (indeed, it is part of the bootblock) to load the initial program and is also the routine used in overlay and linkup operations both by CTOS and utility packages.

#### 2.1 BOOTSTRAP ACTIONS

When a restart occurs, the rear deck is rewound and the first block on the tape (called the bootblock) is loaded into memory starting at location zero. The first 512 bytes of memory (0 to 0777 octal) have been reserved for a permanently resident program which is loaded from the bootblock. The first 40 bytes of this block constitute a program which runs a parity check on the rest of the block that should have been loaded. The processor is halted (note auto-restart implications if the auto-restart tab on the cassette

is punched out) if this routine finds a fault in the check. Otherwise, zeros are stored in the memory locations used in the parity check routine. This will cause a halt if an early data drop-out from the tape machine occurs during the next bootstrap load (typically only one or two bytes get loaded in this failure mode). After the low memory has been cleared, a routine calls the loader, which has been loaded in the bootstrapping operation, asking for file zero to be loaded from the rear deck. If file zero cannot be loaded for some reason, the program halts the processor without a whimper (no bells or whistles in any of the bootstrap operation), otherwise, it jumps to the starting address supplied with file zero. Note that if the auto-restart tab is punched out of the rear cassette, any failure along the road of bootstrapping will cause the whole process to be tried again.

## 2.2 FILE ORGANIZATION

Once file zero has been loaded from the system tape, the bootstrap program (locations 0 through 074) is never used again until the next restart operation which will overlay it. The loader, however, will be used many times. The physical layout of information on the system tape is as follows:

BOOTBLOCK/FILE0/FILE1/. . ./FILE15/FILE32/FILE127

File 0 is the one executed by the bootstrap and is typically followed by a sequential (required to be sequential by the loader) set of minimally increasing (file numbers go up by only one at a time) files up to 15 (a CTOS catalog size limitation, although the loader will load a file with any positive number), followed by a file 32, which is a system scratch file, followed by a file 127 (largest positive eight bit number), which is a dummy to mark the logical end of the tape.

## 2.3 FILE LAYOUT AND RECORD FORMAT

Each file is a group of records starting with a very special four byte record. Every record used by CTOS starts with two special bytes to indicate that it is one of three types: file marker, numeric data, or symbolic data. The file marker, which is the special four-byte record at the beginning of a file, contains two additional bytes that denote the file number. The use of two bytes for both the record type and file number provides redundancy for error control, since the second byte is simply the one's complement of the first. The record types are denoted by 0201 for file marker, 0303 for numeric data, and 0347 for symbolic data. The following table summarizes all of the various data formats used by the system. XP and CP denote the two longitudinal parity checks and will be described later. FN denotes the file number and -FN its one's complement.

FILE MARKER RECORD: 0201 / 0176 / FN / -FN  
NUMERIC DATA RECORD: 0303 / 074 / XP / CP /  
DATA  
SYMBOLIC DATA RECORD: 0347 / 030 / XP / CP /  
DATA (with VRC)  
FILE: FILE MARKER / DATA RECORD/ DATA  
RECORD / . . .  
SYSTEM TAPE: BOOTBLOCK / FILE 0 / FILE 1 /  
. . . / FILE 15 / FILE 32 / FILE 127  
DATA TAPE: FILE 0 / FILE 1 / . . . / FILE 127

## 2.4 LOADER ACTION

When the loader is told to load a given file, it begins searching the tape (the loader can load files from either deck, depending upon which entry point is used) forward until it finds a file marker record. Note that all records passed over must have a valid type number pair or an error recovery procedure will be initiated which will try up to three times to read the record correctly and then make an error exit if failure occurred all three times. Upon finding a file marker, the loader determines, from the number in that marker, whether the tape is positioned to the correct place (the number is equal to that requested), is not positioned far enough forward (the number is greater than that requested). If the tape is positioned to the correct place, the loader proceeds to load all of the numeric records it finds, obtaining the memory address of where it is to put the data from the beginning of each record, (symbolic records are ignored) until it runs across another file marker. At this point it stops the tape (which was in slew forward mode) and backs up over the file marker so a succeeding call on the loader would cause a file marker to be found immediately. If there were no numeric records in the file, an error return is made. If the tape is not positioned far enough forward, the loader searches forward for the next file marker. If the tape is positioned too far forward, the loader enters a reverse search mode. If, in this mode, the loader finds a file marker that indicates that the tape is now positioned to the correct place, tape motion is reversed and the file is loaded as in the forward search case. If it finds a file marker which indicates that the tape is not positioned backward far enough, the loader continues searching in the reverse mode for the next file. If, however, a file marker is found that indicates that the tape has been positioned too far backward, the loader decides that the file is not on the tape and makes an error return. Error returns are also made if a record can not be read without a parity failure or type indicator discrepancy (the two characters are either not the one's complement of each other or are not one of the three special numbers) occurring in all three trials or if loading the record would overstore the loader. In all of these cases, the carry condition will be true (a satisfactory load always rendering the carry condition false) and the tape will be positioned after any offending record.

## 2.5 PARITY CHECKING

The third and fourth bytes of every data record contain longitudinal parity checks. These bytes are set up by the record generation program such that the following exclusive OR sums will yield zeros: the first byte with all the data characters (data characters start with the fifth byte of the record and proceed to the end) and the second byte with the same characters except the sum is shifted right circularly one place after each exclusive OR. In the case of symbolic records, the additional condition of the vertical parity of each character being odd must also be met. One thing not mentioned in the discussion of the loader was that the first four data characters (fifth through eighth bytes in the record) are not really data but are the MSP and LSP of the starting memory address followed by the one's complement of the MSP and LSP of the starting memory address of where the data is to be loaded.

## 3. THE CATALOG, SYMBOLIC LOADER, BASIC I/O, AND DEBUG

As mentioned above, the operating system maintains a catalog of names which correspond to the files on the system tape. This catalog may be used in manipulating the files from the keyboard or in symbolically calling in overlays using the symbolic loader from a user program.

### 3.1 CATALOG CHARACTERISTICS

Each name in the catalog must start with a letter and may additionally contain from one to five alpha-numeric characters. There is room in the catalog for up to fourteen names so there is a limit of fourteen cataloged files on one system tape. The symbolic loader contains routines which will look up a given name up in the catalog and load the corresponding file. This same lookup routine is used by the command handler and is labeled LOOKUP.

### 3.2 UTILITY ROUTINES IN THE SYMBOLIC LOADER

Other utility routines in the symbolic loader area are a block transfer, labeled BLKTFR, and a routine, labeled INCSWP, which increments the H and L register pair and then swaps it with the D and E register pair. The block transfer will move the number of characters specified by the entry value in the C register from a memory address starting with the entry values in the H and L registers to a memory address starting with the entry values in the D and E registers.

### 3.3 LOADING ROUTINES

To use the symbolic loading routine, one loads into the D and E registers the address of the six characters of the desired name (trailing blanks must be included) and calls

MLOAD\$. If the zero condition is false upon return, then the given name was not in the catalog. If the zero condition is true but the carry condition is false upon return the loader could not either find or correctly load the file requested. Note that one must be certain to place the call to MLOAD\$ in a place that will not be overlaid since execution will resume following the CALL instruction after the file has been loaded.

## 3.4 OTHER SYMBOLIC LOADER FACILITIES

Another facility in the symbolic loader area will load and execute a file whose number (not name) is in the B register upon call. Calling MAUTO\$ will load the file from the system tape and calling MAUT2\$ will load the file from the data tape. If the loader could not either find or load the file, the operating system is automatically reloaded.

## 3.5 KEYBOARD AND DISPLAY ROUTINES

The operating system contains facilities to ease the burden of communicating with the operator. Two routines exist. The first accepts the characters from the keyboard, displays them on the screen, and stores them into a memory buffer. The second writes a string of characters from a memory buffer onto the screen.

### 3.5.1 KEYBOARD INPUT

The keyboard input routine, labeled KEYIN\$, accepts a specified maximum number of characters, given by the entry value of the C register, from the keyboard and puts them into memory starting at the entry value of the H and L registers and onto the screen at a starting horizontal cursor position of the entry value of the D register and vertical cursor position (which cannot be changed during the course of one input) of the entry value of the E register. Note that if the cursor collides with the right edge of the screen during entry, characters other than backspace, cancel, and ENTER will not be accepted, although they will print over each other in the last display position. The ENTER character (015) terminates input and is stored in the memory buffer to specify the end of data but is not written to the screen. Hitting the backspace key will delete the last character entered and move the cursor appropriately while hitting the delete key will delete all characters entered and also move the cursor appropriately. These two keys also back up the buffer memory pointer appropriately. Note that if one has typed a character at either the screen limit or at the maximum character count limit, hitting a backspace will cause the previous character to be erased and leave the last character still on the screen, although it will either not appear on the memory buffer or be after the 015.

## 3.5.2 DISPLAY OUTPUT

The display routine, labeled DSPLY\$, will display the string of characters stored in memory starting at the address which is the entry value of the H and L registers and terminating with a character whose numerical value is either a 3 (ETX) or 015 (ENTER). The cursor starts at the entry values in the D (horizontal) and E (vertical) registers (a cursor position that is off the screen will not be sent to the CRT) and stops after the last character printed if the terminating character was a 3 or at the beginning of the following line if the terminating character was an 015. Note that, as in KEYIN\$, the cursor stops at the right edge of the screen and the characters overwrite each other if more are available after collision. Also note that if display was occurring on the bottom line and the terminating character is an 015, then the whole screen is rolled up to force the existence of a following line and the information that was at the top of the screen is lost. After return from the display routine, the H and L registers will point to the location after the terminating character and the D and E registers will reflect the current cursor position. The cursor will be off while the display routine is writing, but it is turned back on upon exit even if it was off upon entry. Other special control characters can cause cursor positioning, line/frame erasure, and screen roll-up:

- 011 - a new horizontal position (0 to 79) follows
- 013 - a new vertical position (0 to 11) follows
- 021 - erase to the end of the frame
- 022 - erase to the end of the line
- 023 - roll the screen up one line

## 3.6 THE DEBUGGING TOOL

The debugging program allows the user to observe and modify any location in memory, to load files from either the system or data tapes, and to start execution at any place in memory. This allows him to load and debug programs with surprising ease. The major debugging technique is to insert RETURN instructions in critical places in memory so one routine at a time may be checked using the CALL command. All but two (user specifiable) of the registers may be saved upon return from the program being tested, allowing the user to determine if the proper actions are taking place by observing critical register and memory values. The registers A, B, C, D, E (subject to the H and L commands in Section 3.6.3) are stored in locations 16770, 16771, 16772, 16773 and 16774 respectively upon a return to Debug from a program which was called from Debug.

### 3.6.1 INPUT SYNTAX AND ERROR ACTION

The debugging program is entered from the command handler as explained in a later section or by processor execution control being passed to the location labeled DEBUG\$. At

this time the bottom line of the display will be erased and the current location and its contents will be displayed there. The program is now ready to accept input in the format <number><command>. The number is assumed to be octal and the absence of any digits between zero and seven implies a value of zero for the number. Only sixteen bits of significance are kept for the input value. If more are entered, the first digits entered are lost. Some commands use only the lower order eight bits. The number is terminated by the first character that is not between zero and seven and this character is taken to be the command. Note that leading spaces are not permitted. This line is read in using the KEYIN\$ routine previously discussed, thus enabling the use of the backspace and cancel keys but requiring the ENTER key to be struck to obtain a response. In one case the ENTER character is the command and in some others the number is disregarded. If the command is not recognized, the program simply ignores it and the old current address and its contents are displayed again. After every command, control is returned to the entry point of the debugging program which will display the now current address and its contents.

### 3.6.2 THE CURRENT ADDRESS

Two memory locations in the debug contain an address (initialized to zero upon loading) which points to a memory location which is the current center of interest. Available commands allow one to change the contents of this memory location and move the pointer as well as perform other functions.

### 3.6.3 COMMAND MEANINGS

The following is a list of each command character with its effect and the number (in parenthesis) of bits of the given number used:

- ENTER - set the current address to the number given (16)
- I - increment the current address by one (0)
- D - decrement the current address by one (0)
- M - change the current address contents to the number given (8)
- .- - do the M followed by the I command (8)
- L - upon return from a C command, cause the L register to be stored into the register whose number is given (3)
- H - same as the L command but for the H register (3)
- G - load from the system tape the file whose number is given (8)
- F - load file one from the data tape (0)
- O - return to the operating system command handler (be sure it is there) (0)
- C - execute a CALL instruction to the location whose number is given (16)

## 4. KEYBOARD FACILITIES (OS COMMAND HANDLER)

The operating system contains a program which will interpret user commands given at the keyboard and perform the tasks indicated. These tasks mainly involve copying new files from the data tape onto the system tape, copying files from the system tape onto the data tape, deleting and updating files on the system tape, and executing programs kept in these files, as well as several other functions.

### 4.1 SYNTAX RULES AND ENTRY ERROR ACTIONS

The command input format is purposely made quite strict to reduce the chance of causing unwanted action which could be catastrophic to the user's data. The command must start with the first character entered (leading spaces are illegal) and any alphabetic after the third character is ignored (thus DEBACLE will be interpreted as the DEB command just as well as DEBUG). The first non-alphabetic character must be either an ENTER, a space, or a dash (minus sign). Some commands will not allow the ENTER but typing a non-alphabetic other than these three will always net you an error message of WHAT?. This will also appear if a command that has legal syntax but is not one of those defined is entered. If the command is to be parameterized, the first name must follow the dash or space immediately and must be terminated with an ENTER if that is the only parameter. The name must start with an alphabetic but may contain any number of alpha-numeric characters even though all after the six will be ignored. If the command has two parameters, the first must be terminated by an ENTER. If a name is terminated by characters other than those specified, the error message BAD NAME will be displayed. If a name is not supplied but the command requires one, the error message NAME REQUIRED will be displayed. If the name given is required to be in the file catalog but is not, the error message NO SUCH NAME will be displayed. If the inverse is true, the error message NAME IN USE will be displayed.

### 4.2 OPERATING SYSTEM COMMAND INSTRUCTIONS

The following paragraphs describe the usage and effect of each command in the system. Each paragraph is titled by what must be entered to use the corresponding routine. Note that, for clarity, more than just the necessary three characters have been shown.

#### 4.2.1 CATALOG

The CATALOG command lists the names of all files that are currently on the system tape. They are listed across the screen in the physical order in which the files appear on the tape. Any parameters supplied are ignored.

#### 4.2.2 NAME (old), (new)

The NAME command will change the name of the file specified by the first name given to the second name given. This command requires that the first and not the second name be in the catalog. The catalog file on the system tape (a one record file (number one) that immediately follows the operating system file) will be overwritten with the new catalog. Note that this leaves the system tape positioned before the file marker of any existing first cataloged file. This operation is performed by all commands that change the catalog.

#### 4.2.3 RUN (name)

The RUN command uses the loader to load the file specified by the name given and then transfers processor control to the starting address indicated to the loader by the file information. Note that it is the responsibility of the loaded program to return to or reload the operating system if this is desired. There is a special case to the RUN command that breaks the general syntax rules. If the name consists of exactly one asterisk terminated by an ENTER (RUN-\*), the loader will be directed to load physical file 1 from the data tape. This provision is made to allow the user to run a program he has generated without having to load it onto the system tape. This, along with the F command in the debugging tool, eliminates a lot of tape movement when debugging programs.

#### 4.2.4 IN (name)

Note that exactly the characters shown must be typed to execute this command since the space which must be the third command character will also terminate the command. This command will position the system tape after the last cataloged file and the data tape to the beginning of physical file 1. (The data tape convention is that physical file 0 will be the first piece of information on the tape, containing the users symbolic data for a given program, and that physical file 1 will be the second piece of information on the tape and will contain the users object data for a given program, and all tape after this is to be considered a scratch area which is properly terminated by physical file 127 to indicate the logical end of the tape.) The command then copies all records in the file from the data tape onto the system tape creating a file on the system tape (a file marker being written before the data was copied) which has the next available physical file number. Following this new file, file markers 32 and 127 are written on the system tape to indicate the new start of system scratch and logical end of tape. If the system tape contained no cataloged files before this command was issued, the file entered will be physical file 2 and immediately follow the catalog file. After the new file has been written, the new name is entered into the catalog and the catalog file is updated as in the NAME command. Note that if the catalog was full when the command was entered, the error message LIBRARY FULL will be displayed and no



other action will occur. The name supplied must not already be in the catalog.

#### 4.2.5 OUT (name)

The OUT command first executes the PREPARE command to provide itself with a null data tape which can be handled by the file handling routines. It then positions the system tape to the beginning of the given file (the name must have been in the catalog) and the data tape to the beginning of physical file one and copies all the records in the file on the system tape onto the data tape. It then places a file marker 127 on the data tape and quits. Note that the catalog file is not updated for this command. This command is provided to allow moving a file from one system tape to another through the associated use of the IN command.

There are two special cases to the OUT command that break the general syntax rules. If the name consists of exactly one dollar sign terminated by an ENTER (OUT-\$) then an exact copy is made of the system tape up to file marker 32 at which time the copy is terminated by file markers 32 and 127 (which causes any scratch data on the old system tape to be removed). If the name is exactly one asterisk terminated by an ENTER (OUT-\*), the action is similar to the previous case except physical files 0 and 1 (namely, the operating system) are deleted and the file numbers of all following data files (not file 32 or 127) are lowered by two. Note that if this tape is now bootstrap loaded, the first program loaded will be what was the first file cataloged in the operating system. This is most useful in preparing bootstrap tapes that will be used in machines with less than 8K of memory.

#### 4.2.6 DELETE (name)

The DELETE command takes two different courses of action depending on whether or not the file deleted is the last one cataloged. If it is, the system tape is moved to the end of the next to the last cataloged file and file markers 32 and 127 are written, thus logically destroying the last file. The name is then deleted from the catalog and the catalog file is updated. If the file is not the last one cataloged, the PREPARE command is called to obtain a fresh data tape, as in the OUT command, and the system tape is positioned to the end of the named file. The rest of the system tape (up to the file 32 marker) is then copied onto the data tape and the data tape is terminated with a file marker 127. Note that the data tape file numbers start out at one and increase by one for each succeeding file copied onto the data tape. These numbers are not used since all the copy back part needs to know is file delimitation since it is getting its file number information from catalog positions. The copying onto the data tape is followed by the system tape being positioned to the end of the file before the one named and the data tape being positioned to the beginning of file one. A file marker having a value one greater than the previous

marker is then written on the system tape and then the data tape is copied back onto the system tape with every file marker encountered on the data tape causing a file marker of value one greater than the previous marker to be written on the system tape. This process terminates when a file marker 127 is encountered on the data tape which causes file markers 32 and 127 to be written on the system tape. The given name is deleted from the catalog, all following entries are dropped down one place to correspond to the similar shift in file numbers that took place, and the catalog file is updated.

#### 4.2.7 REPLACE (name)

The REPLACE command is quite similar to the DELETE command except that instead of preparing the data tape with the PREPARE command, it positions it to the end of file 1 and then writes a file marker 2. Now, copying all the files after the named one onto the data tape in a fashion similar to the DELETE command and copying the data tape back onto the system tape in exactly the same fashion as in the DELETE command will replace the named file by file 1 on the data tape, with any necessary physical expansion or contraction taking place. Even though the catalog is not changed in this operation, it is updated anyway since this is an easy way to position the system tape to a place before file marker 127. Without this, a succeeding call on the loader would run into trouble since the system tape would be left positioned after file marker 127 and the loader always starts by searching a tape forward which in this case would be off the logical end of the tape. The loader starts with a forward search because the very first time it is used, the tape is positioned just after the boot-block and a backward search for a file marker would cause trouble. The operating system routine which searches for files can start with a reverse search to avoid the problem since the tape will never be resting before file zero.

#### 4.2.8 AUTO or AUTO (name)

There is a word in the catalog which contains the physical file number of a file which should be loaded and executed immediately upon loading and execution of the operating system. This enables a user program to be run after restart without interaction with the operation system being required. If this word is a zero or the keyboard switch is being depressed upon initial execution of the operating system, the normal entry is made into the operating system and the start up message and response request are displayed.

If the AUTO command is given with no name and the auto pointer is zero then the error message NAME REQUIRED will be displayed. Otherwise the name of the file being pointed to will be displayed in the message AUTO SET TO (name). If the auto command is given with a name (which must be in the catalog) and the auto pointer is a zero, the

pointer will be changed to the corresponding file number and the catalog (which contains the pointer) will be updated. If the auto pointer is non-zero, the name is ignored and the AUTO SET TO (name) will be displayed as in the no-name case.

#### 4.2.9 MANUAL

The MANUAL command will zero the auto pointer and update the catalog if the auto pointer was non-zero. Otherwise, the message AUTO NOT SET will be displayed.

#### 4.2.10 PREPARE

The PREPARE command first asks the operator if the data tape contains anything of value and then halts. (Note that the auto-restart tab should not be broken out of the operating system tape because it will prevent use of the OUT, DELETE, or PREPARE commands since halting the processor will cause an auto-restart.) After the operator hits the RUN button as a response, it is assumed that the data tape is of no value as it is rewound and file markers 0, 1, and 127 are written on it. This is needed since the operating system routines require file markers for which they can search in using the data tape.

#### 4.2.11 HEX (name)

The HEX command is similar to the IN command except that the data tape is formatted in symbolic records with no parity checking. This is useful in loading onto the system tape data produced by sources other than the 2200. There are four types of records accepted. The type is determined by the second character (the first must always be an 012 (LF)): asterisk means ignore the record; pound sign denotes the logical end of the tape; plus sign means the following four hexadecimal characters are a new starting address (these must be terminated by an 023 (XOFF)); and a hexadecimal character denotes a data record. All other cases are assumed to be data read errors. A data record must always contain an even number of only hexadecimal (0 through 9 and A through F) characters terminated by either an 023 or a plus sign. The characters are paired up to form successive bytes of eight bit data. If the terminating character is an 023 then the block of data bytes is written out in loader format and the starting address is incremented by the number of data bytes. If the terminating character is a plus sign then the data remains in the buffer and the following record will be appended to it. This allows blocks of larger than 36 bytes (128 is the upper limit) to be written when the device which writes the tape is limited to lines of 72 characters. Note that there is no buffer overflow protection and it is the responsibility of the program generating the symbolic data to keep the total number of continued bytes to 128 or less (128 hexadecimal character pairs). Also note that if a continuation line is followed by a new address line, the data will remain in the buffer but the starting address will change. This combination will cause incorrect results

since even if the buffer did not overflow will also overwrite critical pointers which will cause the operating system to produce an error message (because it will be called with incorrect parameters when the critical pointers are overwritten) and be reloaded. If a read error is detected, the data tape is backspaced one record and read again. This will go on until the data appears correctly or the keyboard switch is depressed. Depression of the keyboard switch causes the same action as reading from the data tape a record starting with a pound sign.

#### 4.2.12 DEBUG

The DEBUG command causes the debugging tool described earlier to be entered.

### 4.3 SYMBOLIC OPERATING SYSTEM AND EXTENDED COMMAND INSTRUCTIONS

The overlay program SOSX is available to extend the operating system command set. The following paragraphs describe the usage and effect of each new command. Each paragraph is titled by what must be entered to use the corresponding routine. Not that, for clarity, more than just the necessary three characters have been shown.

#### 4.3.1 CHOP (name)

The CHOP command deletes the named file and all subsequent files.

#### 4.3.2 INSERT (new, (old)

The INSERT command proceeds like a REPLACE command except it includes the old named file as one of the files written after file 1 on the front deck. When the front deck is copied back onto the CTOS tape a new object file has been inserted.

#### 4.3.3 APPEND (name)

The APPEND command appends the object file from deck 2 onto the end of the named file on the CTOS tape. Like the DELETE command, it has two possible courses of action, depending on whether or not the file being appended is the last cataloged file. If it is, the tape is positioned to the end of the cataloged file and a new object file is copied from the front deck. New file 32 and 127 markers are written. If the named file is not the last cataloged file, the operation proceeds like REPLACE except that the CTOS tape is positioned to the end of the named file before the copy is performed.

#### 4.3.4 LGO (name [, name, name . . . ] )

The LGO command makes a tape with a loader and the named file(s) in the sequence named in the command. The files will have sequential file markers starting with 0. There is a limitation

of 23 characters on the command length, thus to name many files in the LGO command it may be necessary to temporarily rename the files with one character labels. LGO \* is not permitted. OUT \* has the desired effect of generating a load and go tape of all cataloged files.

### 4.3.5 SYMBOLIC (name)

The SYMBOLIC command adds a compressed source file (file #0) to the CTOS tape (in a fashion similar to IN). The name in the internal catalog will have an 'S' in the seventh (not displayed) position to identify the file as symbolic.

### 4.3.6 SREPLACE (name)

The SREPLACE, symbolic replace, command is performed exactly as the REPLACE except the compressed source file (file #0) is used instead of the object file. File 1 may be overwritten.

### 4.3.7 SINSERT (new), (old)

The SINSERT, symbolic insert, command is performed exactly as the INSERT except the compressed source file is inserted instead of the object file. File 1 will be overwritten.

### 4.3.8 ATTACH (name [, name, name . . . ] )

The ATTACH command positions the front deck to the end of file 0 and (without file markers) copies specified file(s) from the CTOS tape to the front deck. When all specified files are copied, the question 'END (LABEL OR :)?' will appear. A six character label may be entered. If ':' is typed no end statement will be added. The ATTACH \* form of the command will attach, in cataloged sequence, all symbolic files to file 0 on the front deck.

## 5. PROGRAM FACILITIES (OS FILE HANDLER)

The operating system contains a set of routines which will perform all of the various input/output functions needed to maintain the files of data on the tapes. These routines are packed in the upper 2K of memory and are made available to the user if he wishes to handle his mass storage problems in conformance with the conventions of the operating system. All routines are uniformly parameterized and are accessed through an entry point table (a group of JUMP instructions to the actual routine locations) so any updates to the operating system will not have any effect upon the user's code.

### 5.1 ROUTINE PARAMETERIZATION

Routine parameterization consists of a memory location in the D (MSB) and E (LSB) registers of the first byte of a group of four bytes (called a packet) which parameterizes the call

more explicitly. This method reduces the number of memory locations required to perform a routine call since, in a typical program, one needs only a few different packets but will have many different calls. The parameterization of some routines is not as extensive as that of others, but the same packet can generally be used for the different calls when they are affecting the same file.

#### 5.1.1 LOGICAL FILE NUMBERS

The first byte in the packet is the logical file number and must be between zero and seven or an internal error H will occur upon calling any routine using this packet. This error condition usually occurs when the user has either failed to load the D and E registers at all or has loaded them with an erroneous value before calling the routine. The second and third bytes in the packet contain the LSB and MSB (respectively) of the first location in memory to be used as a data buffer. Actually, the two bytes previous to this location will be used by some of the routines as discussed later. This data buffer may be located anywhere in memory. The fourth byte in the packet specifies the length of the data buffer when numeric data is being handled. Note that using only one byte for the length implies that numeric records may not contain more than 256 data bytes. Actually, the maximum number of data bytes specified may not be greater than 254 for reasons that are made clear in the numeric routine instructions. The four bytes of the packet may be located anywhere in memory.

#### 5.1.2 PHYSICAL DEVICE AND FILE NUMBERS

The logical file number specified in the packet is converted by each routine, via an internal transformation table, into physical file and device numbers. The physical device number specifies whether the operation is to be performed on deck on (rear) or deck two (front) and the physical file number specifies which file is to be treated on the given deck. Actually, not all routines use all of this information since, for instance, when one is reading records from a file he assumes that he is using the file to which the tape was last positioned. The internal transformation table is initialized at load time to the following values:

LOGICAL FILE	PHYSICAL FILE	PHYSICAL DEVICE	GENERAL USE
0	0	0	Unassigned
1	0	1	General deck one
2	0	2	General deck two
3	1	1	CTOS catalog
4	0	2	Symbolic data
5	1	2	Object data
6	0	0	Unassigned
7	32	1	System scratch

It is shown that logical files 1 and 2 are specified for use of any physical file, even though 0 is shown in the table. This can be done by use of a routine that will change the physical file number of a given logical file. A routine also exists to allow the physical device number to be changed, thus allowing the user to set up logical files in any physical configuration needed. Note, however, that one must have logical files 1 through 5 and 7 in the state shown (except for the physical device numbers of logical files 1 and 2) if one returns control to the operating system command handler, since the loaded values are assumed by this program. Logical files 0 and 6 may be used freely but must be set before the first call utilizing them. The following is an example of a packet usage as it would be expressed in the assembler: (Note all calls to CTOS tape routines must, as in the following example, be preceded by a DE to the first byte of the packet. Note also that the packet consists of 4 bytes: Logical file number, LSP of buffer, MSP of buffer and length of buffer)

	LA	2	Set up logical file six to be used as physical file 3 on the
	DE	PACKET	front deck
	CALL	CPDN\$	
	LA	3	
	DE	PACKET	
	CALL	CPFNS\$	
	DE	PACKET	Position to the beginning of the file
	CALL	PBOFS\$	
LOOP	DE	PACKET	Read a record of symbolic from it into BUFFER
	CALL	SSFR\$	
	JTC	DONE	Quit if to the next file marker
	JFZ	TERR	Exit if type error
	.		
	.		Action taken for each record
	.		
	JMP	LOOP	
DONE	.		
	.		Action taken when file completely in
	.		
TERR	.		
	.		Type error action
	.		
PACKET	DC	6	Logical file 6
	DA	BUFFER	Buffer address
	DC	0	Length (not used)
	DC	0,0	Room for parity check generation
BUFFER	SKIP	128	Buffer area

## 5.2 ROUTINE USAGE INSTRUCTIONS

To use a routine, one sets up whatever is required for proper parameterization and then calls the desired location in the entry point table. The locations are labeled with the first word in the following paragraph titles followed by a \$. For example, to call the serial numeric file read, one would say CALL SNFR\$. The routine will either perform the requested task or take one of two error exit paths. The first path is

taken in the case of fatal errors, for which it is decided that the only recourse is to reload the operating system. This is called an internal error and the message INTERNAL ERROR (letter) is written on the bottom line of the display before the system is reloaded. The various letters which may appear are the following:

- A - Illegal device specification
- B - Illegal record format
- D - Unrecoverable parity error
- G - Unfindable file
- H - Illegal logical file specification

The other path is non-fatal and simply returns with certain condition flags in states other than normal to indicate that something unusual happened. Since every routine uses a common subroutine (labeled GETPKT) to get the parameters from the specified packet, common internal errors can occur. If the logical file number is not between zero and seven, an

internal error H occurs. If the physical device number is not either a 1 or a 2, an internal error A occurs. Other than for these error actions, the following paragraphs described the effects of and the exact parameterization needed for each routine.

### 5.2.1 SNFR - SERIAL NUMERIC FILE READ

This routine reads the next record from the specified device. If the record is of type symbolic, the zero and carry conditions

are set false and return occurs with no parity checking or data storage being performed. If the record is a file marker, the carry condition is set true and the tape is backed up to where it was before the routine was called. Again, return occurs with no data storage being performed. If the type is numeric, the two parity bytes followed by the data are read into the buffer. If the parity checking fails or the record type is bad, three efforts are made at reading the record by backing up to its beginning and starting over. If recovery is not made in one of these efforts, an internal error D occurs. If the record is read successfully, return occurs with the zero condition true, the carry condition false, and the H and L registers containing the memory location of the byte following the last one loaded from the tape. To calculate the length of the buffer area used, one must subtract the buffer starting address from returned values in the H and L registers. Remember that the first two characters in the buffer are not data characters but are the two longitudinal check sums. To obtain the number of data characters loaded, one must subtract the buffer starting address plus two from the returned values in the H and L registers. The parity checks are stored because the SBFW routine uses them instead of regenerating them from the data, thus shortening the time required to copy numeric records from one deck to the other.

### **5.2.2 SSFR - SERIAL SYMBOLIC FILE READ**

This routine reads the next record from the specified device. If the record is of type numeric or file marker, the action taken will be the same as when SNFR reads a symbolic or file marker record. Action similar to that taken by SNFR is also taken if parity or type faults occur. If the record is read satisfactorily, only data characters will be in the buffer starting at the address specified. An 015 will mark the end of the data string and all vertical parity bits will be zero. The same normal exit conditions as in SNFR will occur.

### **5.2.3 SBFW - SERIAL BLOCK FILE WRITE**

This routine writes a record of type numeric on the specified device. The total number of bytes, including the parity initialization sums as the first two, must be in the fourth byte of the packet. Note that inclusion of the parity initialization sums implies that the total number of actual data bytes cannot exceed 254. This routine assumes that the first two bytes in the buffer are the correct parity initialization sums since it does not generate them from the data. There are no error exits from this routine which implies that writing off the end of the tape will not be caught and that read-after-write checking is not performed.

### **5.2.4 SNFW - SERIAL NUMERIC FILE WRITE**

This routine performs in a fashion similar to the SBFW routine except the two parity bytes are not included

in the data buffer and the length specifies the number of actual data characters. The routine generates the two longitudinal parity sum initialization values and inserts them in the two locations preceding the buffer. It then writes on the specified device a record of type numeric containing the two parity bytes generated, followed by the number of data bytes specified. Note that the length is adjusted to accommodate the two parity bytes so, as in the SBFW routine, only 254 actual data bytes may be written. If one specifies a length of 255 or 0, the only bytes (besides the record type) written on the tape will be respectively the first or both parity initialization sums. No error exits are made from this routine.

### **5.2.5 SSFW - SERIAL SYMBOLIC FILE WRITE**

This routine performs in a fashion similar to the SNFW routine except that an 015 character in the data string rather than a specified value is used to determine the buffer length, vertical (in addition to longitudinal) parity generation is performed, and a record of type symbolic rather than numeric is written. The terminating 015 character is not included in the set of characters written to the tape, but remember that it will appear again if the SSFR routine is used to read the record.

### **5.2.6 PEOF - POSITION TO END OF FILE**

This routine searches forward on the specified device until it finds a file marker. It then backspaces the tape until it is between the next to the last and the last record in the file. It then forward spaces the tape one record which puts it at the end of the file, having arrived there via forward tape motion. This forward arrival is important to observe when one plans to append one record after another and still maintain physical interrecord gap integrity. Note that every record passed over by the PEOF routine must have a valid record type or it will be read again in action similar to parity failure action in the SNFR routine.

### **5.2.7 PBOF - POSITION TO BEGINNING OF FILE**

This routine searches for a file marker in a fashion similar to the loader except it starts by searching backwards. The file number searched for is specified by the physical file number supplied by the generalized parameterization. Note that since this routine starts by searching backwards, it will not decide that the requested file is not on the tape until it has found in the search forward mode a file marker that specifies a file number greater than the one desired, if indeed the file is not on the tape. Also note that if the leader is found in the search backward mode, the tape is positioned forward past the first record and the backward search is continued. If the first record is not a file marker (operating system convention requires it to be) or is a file marker whose value is greater than the one desired, the first record on the tape will be passed over back and forth until external intervention is imposed.

Otherwise, all search rules and error exit conditions of the loader routine apply here. If, upon return, the carry condition is true, then the file was not found. Otherwise, the tape will be positioned at the interrecord gap following the file marker, having approached that point with forward tape motion for the reasons expressed in the PEOF routine instructions.

### **5.2.8 BSP - BACKSPACE**

This routine simply backspaces the tape one record using the hardware backspace function. No checking is made to see if the record was of proper type or if the tape ran onto the leader.

### **5.2.9 CPDN - CHANGE PHYSICAL DEVICE NUMBER**

This routine stores the entry value of the A register (note the break from generalized parameterization) into the physical device number entry for the specified logical file in the internal transformation table.

### **5.2.10 CPFN - CHANGE PHYSICAL FILE NUMBER**

This routine stores the entry value of the A register (note the break from generalized parameterization) into the physical file number entry for the specified logical file in the internal transformation table.

### **5.2.11 TRW - TAPE REWIND**

This routine performs a hardware high speed rewind of only the front deck. If the rear deck (physical device 1) is specified, an internal error A will occur. Upon exit from this routine, the tape will be positioned to the clear leader.

### **5.2.12 TFNR - TAPE FILE NUMBER READ**

This routine acts in a fashion similar to PEOF until it finds the file marker. At this point, it simply reads the value of that marker and leaves the tape positioned after the marker record. The value read is returned in the C register. Error exits similar to the PEOF routine can occur.

### **5.2.13 TFWN - TAPE FILE NUMBER WRITE**

This routine will write on the specified deck the special four byte file marker record containing the physical file number specified. No error exits will occur.

SECTION 6

DATAPoint 2200

TRACE

## SECTION 6

### DATAPOINT 2200 TRACE PROGRAM

#### Introduction

TRACE is an interactive octal debugging aid for the Data-point 2200. It operates under the Computer Terminal Operating System (CTOS) and occupies memory space between 5600g and 13777g. The normally resident operating system subroutines are not overlaid and are callable by the program being traced.

TRACE accepts commands from the keyboard and displays its results in the rightmost eight columns of the CRT display. It allows a user to trace the execution of a program, to examine and change the contents of the registers and memory.

#### Entering Commands

TRACE commands consist of up to two octal operands followed by a single letter operation. If there are two operands, a comma shall separate the two. An operand may be a 13-bit address or an 8-bit byte value, either expressed in octal. If the operand is an address, it may be given in two parts, separated by a blank. The first part consisting of the five most significant bits, and the second part consisting of the remaining eight bits. An address may also be given as a single octal number, but it is displayed in two parts as described. Leading zeros need not to be entered.

Examples:

TYPED VALUE	DISPLAYED VALUE	
	FIRST PART	SECOND PART
101A	—	101A
707J	01	307 J
0331, 1477W	{ — 03	331 , 077 W

The command being entered is displayed on lines 10 and 11 of the CRT display, as shown in Figure 1. Line 10 shows the first operand. Line 11 shows the second operand and operation. If there is only one operand, line 11 will be blank. If there are no operands, line 11 shows only the operation and line 10 will be blank. NOTE: In all ensuing examples, the display format is used to exemplify the referenced operation.

If an illegal character is typed, the beep signal will sound and the character will be ignore. The CANCEL key will cause the command being entered to be discarded and another command can be entered.

The ENTER key will cause the command just entered to be executed. The command may be CANCELED at any time before the ENTER key is depressed.

#### Command To Modify Registers Or Memory

##### A, B, C, D, E, H, L

The operations A, B, C, D, E, H or L take a single byte value operand. The register specified by the operation is set to the operand value.

Examples:     1 7 3 B sets B to 173  
                   0 0 1 A sets A to 001  
                   3 7 5 L sets L to 375

##### Operation F

The operation F takes a single address operand. The Zero, Sign, and Parity flip-flops are set as if the lower 8 bits of the address were the result in A of some arithmetic instruction. The Carry flip-flop is set to the rightmost bit of the first part of the address.

Examples:             0 0 0 F sets Zero  
   resets Sign, Carry, and Parity

0 1 0 0 0 F sets Zero and Carry  
   resets Sign, and Parity

2 0 0 F sets Parity and Sign  
   resets Zero and Carry

##### Operation O

The operation O takes a single address operand. It opens the specified location for possible modification. The contents of the location are shown on line 12 of the CRT display. A byte value can be entered followed by ; , <, or > followed then by ENTER. The location is set to that value. If the terminating character is ; , TRACE will accept another command. If it is <, the previous location is now opened. If it is >, the next location is opened. If the CANCEL key is used, the currently open location remains open and any modification for it is discarded. The modifying byte value is shown as it is entered, following the contents of the open location on the CRT display.

Examples:

01 115 O opens 01 115  
                   < now opens 01 114  
                   > reopens 01 115  
                   57; sets 01 115 to 057



```

00 017 O opens 00 017
5> sets 0 017 to 005,
opens 0 020
20; sets 0 020 to 20

```

## Command To Displayed Memory

### Operation M

The operation M takes two operands, both addresses. They are the lower and upper bounds respectively of the region of memory to be displayed. Sixteen bytes are displayed across the entire width of the CRT display. The address is given on one line followed by the memory contents on a second line. The display continues being built and rolled up unless the DISPLAY key is depressed. The display then stops until the DISPLAY key is depressed again. The KEYBOARD key terminates the memory display.

Example: 000, displays the first 256 bytes of  
00 377 M memory

## Transfer Of Control Commands

### Operation K

The operation K takes one address operand. It causes a Call instruction to be performed to the address given as the operand. Return is to TRACE.

Example: 02 000 K calls routine at 2 000

### Operation J

The operation J optionally takes one address operand. If the operand is absent, the content of the P register is used. It causes a Jump instruction to be performed to the address given as the operand or in P.

Examples: 03 101 J jumps to 3 101  
(octal 1501)  
J jumps to address in P

## Mode Setting Commands

### Operation X

The operation X optionally takes two address operands. They are the lower and upper bounds respectively of a region in memory. Any Call instructions into this memory region are actually executed rather than being simulated. Since TRACE loses control at this point, it is imperative that routines called in this region return. All registers are properly set when the Call is performed. The contents of H and L are lost on the Return. If the operation X is given without operands, any region in effect is removed.

Example: 0 1 000 , sets the special CALL region to  
to 1 000 to 3 377  
03 377 X (octal 400 to 1777)

### Operation W

The operation W optionally takes two address operands. They are the lower and upper bounds respectively of a region in memory. At the completion of any instruction in this region, the registers contents are shown in the first nine (9) lines of the CRT display as shown in Figure 1. The contents of the Carry, Zero, Sign, and Parity flip-flops are shown as the letters C, Z, S, and P respectively on the right-hand side of the CRT display if set and blank is reset. If the operation W is given without operands, any region in effect is removed. If the regions for the X operation and W operation overlap, the X operation takes precedence.

Example: 12 000 , display the register contents after  
12 377 W every instruction in 12 000 to  
12 377 (octal 5000 to 5377)

### Operation S

The operation S optionally takes one address operand. Before TRACE executes the instruction at the given address, the register contents are shown as in a W operation region and TRACE stops to accept commands. A J operation with no operands will restart the program. If the S command is given without an operand, any stop address in effect is removed. If the stop address falls within an X operation region, the X operation takes precedence.

Example: 7 011 S the program will stop before the in-  
struction at 7 011 is executed

To stop a program when TRACE is in control, depress the KEYBOARD and DISPLAY Keys at the same time. This has the same effect as an S operation for the current program address.

If a Halt instruction is executed by TRACE, the result is the same, as if an S operation was set for the Halt otherwise ignored.

## Starting Trace

TRACE is loaded like any other program from the operating system library (RUN TRACE). Once started, it will request the name of the program desired to TRACE. The name of the program must be typed, followed by ENTER. CANCEL will cause TRACE to ask again. The named program will be loaded using the symbolic linker and loader in the operating system, and TRACE will show the register contents as if the program had been stopped at its entry point. P and I will be the only registers with non-zero contents.

To TRACE a program already in memory, simply depress the ENTER key without entering a program name, then jump to the entry point of the traced program, using the J command.

To TRACE a program located on the front deck type \* and depress the ENTER key. The object file (file #1) from the front tape will be loaded using the operating system loader, and TRACE will show the register contents as if the program had been stopped at its entry point. P and I will be the only registers with non-zero contents.

### Operational Summary

A	Set A to operand
B	Set B to operand
C	Set C to operand
D	Set D to operand
E	Set E to operand
H	Set H to operand
L	Set L to operand
F	Set flip-flops from operands
O	Open location
J	Jump
K	Call
M	Display memory
X	Set special CALL region
W	Set register display region
S	Set stop address

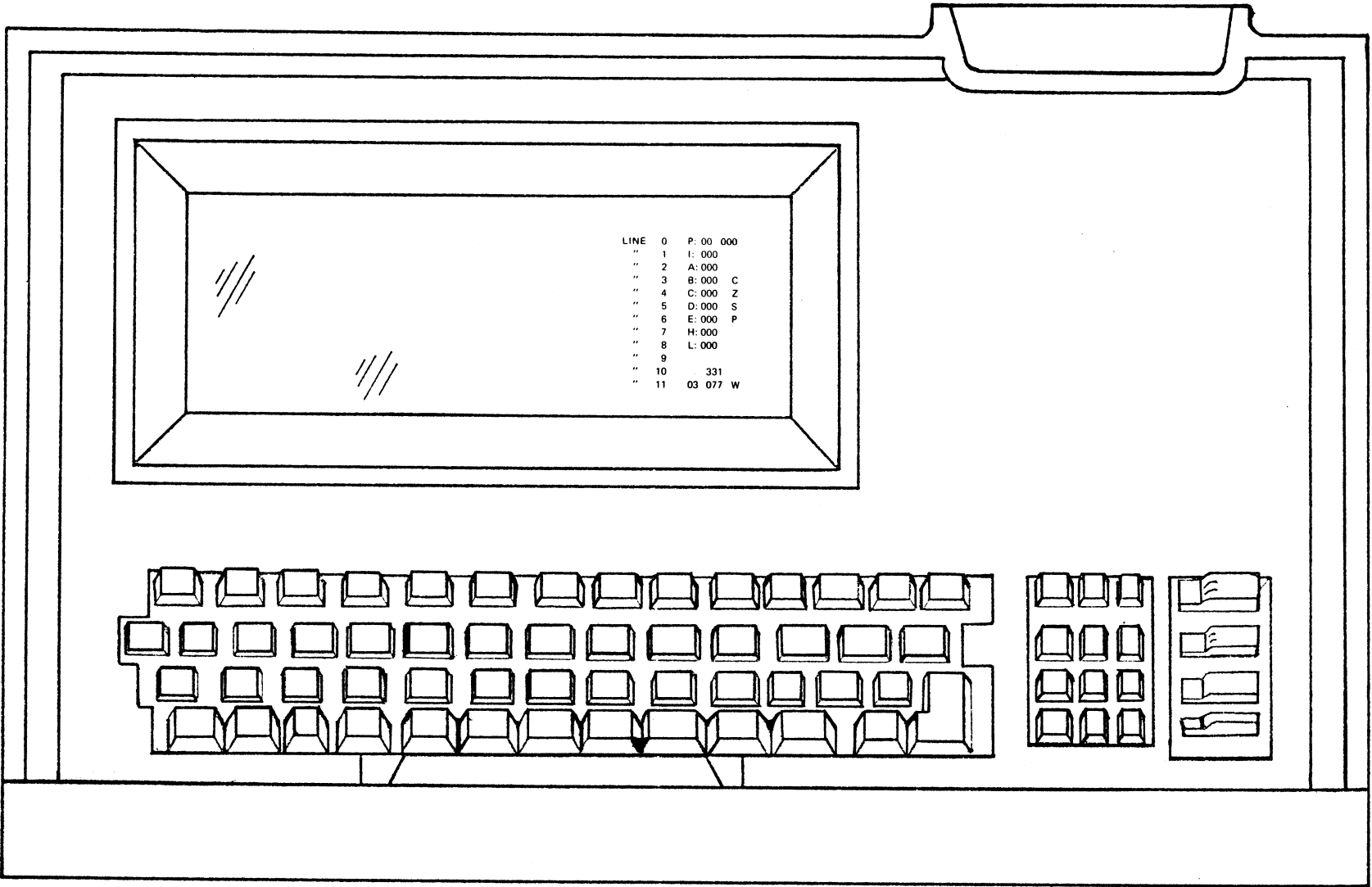


Figure 1. Operation W

SECTION 7

DATAPoint 2200

ARITHMETIC SUBROUTINES

## SECTION 7

### 7.1.1 INTRODUCTION

STATH is a subroutine package specifically designed to provide formatted keyboard input, screen display, checksum and arithmetic operations on numeric strings. Each function of STATH is obtained by calling the entry point associated with that function.

Following is a list of the functions available through STATH. The labels given to their entry points and the sections incorporating their usage parameters:

Entry Point	Function
ADD\$	Addition
COM\$	Compare Magnitude
DIV\$	Division
DSP\$	Display on screen
KEY\$	Keyboard formatted Input
MOD10\$	MOD 10 checksum calculation
MOD11\$	MOD 11 checksum calculation
MOV\$	Move string
MUL\$	Multiply
SUB\$	Subtract

#### 7.1.1.1 INTRODUCTION TO STRINGS - NUMERIC AND OTHERWISE

The purpose of a 'string' is to carry around a 'package' of text. A string is an individual block of text and just like a string it has a definite beginning and end. The composition of the string is an uninterrupted sequence of ASCII characters. That is, between the beginning and end of the string only ASCII characters are allowed. The ASCII character may be any of the 95 plus space (blank) characters listed in Section 1 of the Programmers' Manual (2200 Reference Manual).

The string is bounded at the beginning and end in different ways. The end is determined by the first occurrence of the ASCII 'ETX' which is equal to (003g) in the sequence of characters called the string. The 003 tells STATH that the string is ended. The CTOS will also accept a carriage return character (015g) in place of the 003 but STATH only accepts the 003.

The following are valid strings. The contents of the parentheses are intended to be the byte value of the ASCII character for single character values or the octal value of the octal triple such as 003.

(N) (O) (W) ( ) (I) (S) ( ) (T) (H) (E) ( ) (T) (I) (M)  
(E) (003)

(0) (1) (2) (3) (4) (5) (6) (7) (8) (9) (0) (003)

Which are in octal:

116,117,127,040,111,123,040,124,110,105,040,124,111,  
115,105,003

and

060,061,062,063,064,065,066,067,070,071,003

Although a string has an inherent end built into itself, the 003, there is no beginning. At least no beginning which itself is part of the string of characters in memory. The beginning is combined with the pointer to the string itself. That is, a string is referred to by calling out a location in memory. That location is the first character of the string. In the above samples, for 'now is the time' to be referred to beginning with the word 'now' the location of the letter 'N' would be specified. It is clear that specifying only the 'N' yields a complete description which is:

'Begin with the character in the location specified and continue until a 003 is reached.'

Beginning with 'N' and continuing to the 003 gives: 'Now is the time'. If the location of the letter 'W' in now were specified, the string string resulting would be 'w is the time'.

Therefore, to specify a string to a routine (like STATH) which is going to use the string, the user must only transfer the address of the first character of the string or the character in the string the user wants to begin the string (it may not be the first) to the routine. Also, if the user created the string, he must be assured that there is a terminating 003 byte immediately following the last character of the string in memory.

STATH differentiates between two categories of strings:

1) Numeric strings

and

2) Non-numeric strings

Where numeric strings are only regular strings with the character set restricted the characters 012345689 with an optional single period representing the decimal point and/or a single hyphen leading the string representing a minus sign.

A non-numeric string is any string which is not numeric by the above definition.

A numeric string (omitting temporarily the 003) can look like:

00000034567788888777.9999999999991

or

-123.45

or

34.5000000000

There is a size limit as to the number of characters a string may have in STATH. This is not true of ordinary text strings in CTOS where a string may, for some strange purpose, have thousands of characters in it. STATH is a mathematic package and the numeric strings represent numbers. The largest number of digits, therefore, is limited in STATH and that limit is 126.

### 7.1.1.2 INTRODUCTION TO THE FUNCTIONS OF STATH

STATH functions fall in the following four categories. The categories are listed with their appropriate functions below.

#### Arithmetic Analysis Manipulative Input/Output

Addition	Compare	Move	Keyboard formatted input
Subtraction	MOD10 Check		Display on screen
Division	MOD11 Check		
Multiplication			

The arithmetic functions are the normal functions with which everyone is familiar.

The analysis functions permit decisions to be made on the content of a number. MOD10 and MOD11 verify the check-sum Modulo 10 or 11 as is used in many business applications. Compare will permit comparing two numbers to determine equality or relative magnitude.

The move function is necessary as a preparation for using the multiplication and division functions in STATH, applicable for general use in the user's program to move numeric strings from one location to another and to format and round them in the process.

The input/output functions provide the user with simple techniques for bringing numbers into memory from the keyboard and displaying string numbers in memory onto the screen.

## 7.1.2 STATH FUNCTIONS AND ARGUMENTS

Each routine takes one or two arguments. An argument consists of a CTOS-compatible string. The argument strings are bounded at the end by an ASCII ETX (=003), and the beginning boundary is determined by the address contained in the register-pair associated with that argument. The maximum size for any STATH string is 126 characters. This means arguments and results are limited each to 126 digits.

Except for the routine DSP\$, all strings must be 'numbers' which means a sequence only of ASCII numeric digits (0123456789) with an optional decimal point. Optional leading minus sign and optional leading blanks (an octal 040). The number must be right justified in the argument string. All strings except for DSP\$ set the condition flags as follows:

Flag	Indication
Zero	The result was zero
Sign	The result was negative
Carry	An overflow occurred
Parity	One or both arguments were improperly formatted

If parity is not set at the end of an operation, HL and DE contain the addresses of the location in memory past their respective ETX's. In the case of KEY\$ and DSP\$, D contains the column and E contains the row of the position immediately beyond the display area used. MUL\$ and DIV\$ leave D and E with junk in them. MOD10 and MOD11 leave H and L containing the address of the check digit position.

### 7.1.2.1 EXAMPLES ON THE USE OF STATH

Following is a 488-byte program which is a useful desk calculator using STATH. It is included as an example of a program calling STATH functions.

'DCLAC', the desk calculator, inputs a numeric string and provides addition, subtraction, multiplication or division of that inputted string against an accumulator. 'DCALC' always inputs the string from the keyboard into a string labeled 'input'. The accumulator is in a string labeled 'accum'.

The four arithmetic operations performed in the program are routines labeled as 'ADDOP', 'SUBOP', 'MULOP', and 'DIVOP'. The routines are very short but demonstrate the use of STATH.

	SET	01000
BOOT\$	EQU	064
MOV\$	EQU	010000
ADD\$	EQU	010003
SUB\$	EQU	010006
MUL\$	EQU	06000
DIV\$	EQU	06003
KEY\$	EQU	010014
DSP\$	EQU	010017
KEYIN\$	EQU	017000
DSPLY\$	EQU	017151
MLOAD\$	EQU	017620
BEEP	EQU	13
HEADING	DC	021,811,20,2 2 0 0
	DC	013,5,011,31,'Total'
	DC	013,7,011,28,'Keyboard'
	DC	013,2,011,28,'0 To 9'
	DC	'Decimal Places?'
DECPL	DC	'0',3
OVFMSG	DC	'Overflow',3
BLANK	DC	' ',3
CLEAR	DC	022,3
INPUT	DC	'0000000000',3
ACCUM	DC	'00000000000',3
DIVID	DC	'0000000000000000000000',3
NAME1	DC	'Stath'
OPCODE	DC	' ',815
DCALC	DE	NAME1
	CALL	MLOAD\$
	JFZ	BOOT\$
DCALCH	DE	0
	HL	HEADING
	CALL	DSPLY\$
	LD	51
	LE	2
	HL	DECPL
	CALL	KEY\$
	LL	INPUT
	CALL	FILLIN
	LL	ACCUM
	CALL	FILLIN
	LL	DIVID
	CALL	FILLIN
	LL	DECPL
	LAM	
	SU	'0'
	LBA	
	LC	'.'
	LA	INPUT+10
	SUB	
	LLA	
	LMC	
	LA	ACCUM+10
	SUB	
	LLA	
	LMC	
	LAB	
	SLC	

	LBR	
	LA	DIVID+20
	SUB	
	LLA	
	LMC	
	LD	28
	LE	2
	HL	CLEAR
	CALL	DSPLY\$
	DE	ACCUM
	HL	ACCUM
	CALL	SUB\$
DCALCL	LD	38
	LE	5
	HL	ACCUM
	CALL	DSP\$
	LD	50
	LE	7
	HL	BLANK+6
	CALL	DSP\$
	LE	38
	LE	7
	HL	INPUT
	CALL	KEY\$
	LC	1
	LE	50
	LE	7
	HL	OPCODE
	CALL	KEYIN\$
	HL	OPCODE
	LAM	
	CP	015
	JTZ	ADDOP
	CP	'A'
	JTZ	ADDOP
	CP	'S'
	JTZ	SUBOP
	CP	'M'
	JTZ	MULOP
	CP	'D'
	JTZ	DIVOP
	CP	'E'
	JTZ	MOVOP
	CP	'R'
	JTZ	DCALCH
	EX	BEEP
	JMP	DCALCL
ADDOP	DE	INPUT
	HL	ACCUM
	CALL	ADD\$
OVFTST	JFC	NOOVF
	LD	36
	LE	3
	HL	OVFMSG
	CALL	DSP\$
	EX	BEEP
	JMP	DCALCL
NOOVF	LE	36

```

LE      3
HL      BLANK
CALL    DSP$
SUBOP   JMP    DCALCL
        DE    INPUT
        HL    ACCUM
        CALL  SUB$
MULOP   JMP    OVFTST
        DE    ACCUM
        HL    ACCUM
        CALL  MOV$
        DE    INPUT
        HL    ACCUM
        CALL  MUL$
        JMP   OVFTST
MOVOP   DE    INPUT
        HL    ACCUM
        CALL  MOV$
        JMP   OVFTST
DIVOP   DE    ACCUM
        HL    DIVID
        CALL  MOV$
        DE    INPUT
        HL    ACCUM
        CALL  DIV$
        JMP   OVFTST
FILLIN LAM
        CP    3
        RTZ
        LA   '0'
        LMA
        LAL
        AD   1
        LLA
        JMP  FILLIN
        END  DCALC

```

Observe the addition routine, 'ADDOP'. To add together the inputted string 'input' to the accumulator 'accum' the user only writes the following code as found at 'ADDOP'.

```

ADDOP   DE    INPUT
        HL    ACCUM
        CALL  ADD$

```

Executing this code will cause string 'input' to be added to the string 'accum' with the result in the string 'accum'. The accumulator, it must be realized, is simply a string which the writer of 'DCALC' is using as his result string and he preferred to call it an accumulator.

Note that after each operation there is a jump to 'OVFTST' or as in 'ADDOP', the code is immediately after and executed right after 'ADDOP'. Observe that the first instruction

```
OVFTST JFC NOOVF
```

of the overflow test is the actual test: If the carry isn't set then there was no overflow resulting from the operation. If the carry was set, in 'DCALC' the message 'overflow' is printed on the screen as is seen from the code following the 'JFC NOOVF'.

Subtraction behaves the same as addition except for the CALL to SUB\$.

Multiplication and division are slightly different from addition and subtraction but operate similar to each other. Observe the following code as taken from 'DCALC'.

```

MULOP   DE    ACCUM
        HL    ACCUM
        CALL  MOV$
        DE    INPUT
        HL    ACCUM
        CALL  MUL$
        JMP   OVFTST

```

This demonstrates the requirement, as stated in 7.1.5, that, in MUL\$ and DIV\$, the argument #2 must be the result of the previous move. The reason for this is that multiplication and division really require three 'registers' or strings: The two strings being multiplied and the result. The 'MOV\$' move operation makes a copy of whatever is being moved, during the move, in an internal STATH 'register' string. Therefore, note that the first three instructions in 'MULOP' cause the accumulator to be 'MOV\$' moved to itself. Frequently the user can save time by utilizing this fact in making the last move before calling 'MUL\$' a move of a string involving argument #2. (Again, argument #2 is the argument associated with the H and L registers).

Also note that 'MULOP' tests overflow using the same routine that is used for the other three arithmetic routines 'OVFTST' as described above.

### 7.1.3 LOADING STATH

STATH may be loaded in memory in either of two ways:

- 1) Incorporating the source code of STATH into the problem source code.
- 2) Catalog STATH as an object file and call it in through the operating system.

The second is preferred and simpler, as is done in 'DCALC'. Once cataloged, the following calls STATH into memory:

```

NAME1   DC    'STATH'
        DE    NAME1
        CALL  BOOT$

```



### 7.1.3 ADDITION

Entry Point Name	ADD\$
Entry Point Address	10003 Octal
Argument #1 Address	D-E Registers
Argument #2 Address	H-L Registers
Result Location	Argument #2
Arithmetic Function	(Argument #2) = (Argument #2) + (Argument #1)

#### Action:

Adds two numeric string numbers, rounds, and installs leading blanks and trailing zeros when needed in the result.

#### Typical calling sequence:

```
ADD$   EQU    010003
      .
      DE     ARG1
      HL     ARG2
      CALL   ADD$
```

#### Arguments:

Arguments must be each numeric strings of less than 126 characters in length. Argument 1 is addressed by the D and E Registers. Argument 2 is addressed by the H and L Registers and will contain the result.

#### Result:

The contents of argument 1 (D and E) will remain unchanged.

The contents of argument 2 (H and L) will contain the sum of arguments 2 and 1 and will have leading blanks and trailing zeros when needed.

#### Changes:

The contents of argument 2 are changed to contain the result.

#### Errors Recognized:

Improper argument format (parity bit set)

Overflow occurrence (carry bit set)

#### Comparison Flags:

Result was zero (zero bit set)

Result was negative (sign bit set)

### 7.1.4 SUBTRACTION

Entry Point Name	SUB\$
Entry Point Address	10006 Octal
Argument #1 Address	D-E Registers
Argument #2 Address	H-L Registers
Result Location	Argument #2
Arithmetic Function	(Argument #2) = (Argument #2) - (Argument #1)

#### Action:

Subtracts one numeric string number from another, rounds and installs leading blanks and trailing zeros when needed in the result.

#### Typical calling sequence:

```
SUB$   EQU    010006
      .
      DE     ARG1
      HL     ARG2
      CALL   SUB$
```

#### Arguments:

Arguments must be each numeric strings of less than 126 characters in length. Argument 1 is addressed by the D and E Registers. Argument 2 is addressed by the H and L Registers and will contain the result.

#### Result:

The contents of argument 1 (D and E) will remain unchanged.

The contents of argument 2 (H and L) will contain the difference of arguments 2 and 1 and will have leading blanks and trailing zeros when needed.

#### Changes:

The contents of argument 2 are changed to contain the result.

#### Errors Recognized:

Improper argument format (parity bit set)

Overflow occurrence (carry bit is set)

#### Comparison Flags:

Result was zero (zero bit is set)

Result was negative (sign bit is set)

## 7.1.5 MULTIPLICATION

Entry Point Name	MUL\$
Entry Point Address	6000 Octal
Argument #1 Address	D-E Registers
Argument #2 Address	H-L Registers
Result Location	Argument #2
Arithmetic Function	(Argument #2) = (Argument #2) X (Argument #1)
Argument Restrictions	Argument #2 must be result of last MOV\$ call

### Action:

Multiplies two numeric string numbers, rounds and installs leading blanks and trailing zeros when needed in the result.

### Typical calling sequence:

```
MUL$   EQU   06000
        DE   ARG1
        HL   ARG2
        CALL MOV$

        DE   ARG1,
        HL   ARG2
        CALL MUL$
```

### Arguments:

Arguments must be each numeric strings of less than 126 characters in length. Argument 1 is addressed by the D and E Registers. Argument 2 is addressed by the H and L Registers and will contain the result. Argument 2 must have been involved in the previous move operation.

### Result:

The contents of argument 1 (D and E) will remain unchanged.

The contents of argument 2 (H and L) will contain the product of arguments 2 and 1 and will have leading blanks and trailing zeros when needed.

### Changes:

The contents of argument 2 are changed to contain the result.

### Errors Recognized:

Improper argument format (parity bit set)  
Overflow occurrence (carry bit set)

### Comparison Flags:

Result was zero (zero bit set)  
Result was negative (sign bit set)

## 7.1.6 DIVISION

Entry Point Name	DIV\$
Entry Point Address	6003 Octal
Argument #1 Address	D-E Registers
Argument #2 Address	H-L Registers
Result Location	Argument #2
Arithmetic Function	(Argument #2) = (Argument #2) / (Argument #1)
Argument Restrictions	Argument #2 must be result of last MOV\$ call

### Action:

Divides one numeric string number into another, rounds and installs leading blanks and trailing zeros when needed in the result.

### Typical calling sequence:

```
MOV$   EQU   010000
      DE    ARG1
      HL    ARG2
      CALL  MOV$
      .
DIV$   EQU   06003
      DE    ARG1
      HL    ARG2
      CALL  DIV$
```

### Arguments:

Arguments must be each numeric strings of less than 126 characters in length. Argument 1 is addressed by the D and E Registers. Argument 2 is addressed by the H and L Registers and will contain the result. Argument 2 must have been involved in the previous move operation.

### Result:

The contents of argument 1 (D and E) will remain unchanged.

The contents of argument 2 (H and L) will contain the result of the division of argument 1 into argument 2 and will have leading blanks and trailing zeros when needed.

The number of decimal places in the result is equal to the number of decimal places in the dividend less the number of decimal places in the divisor. This number may not be negative and if it is, the number of decimal places is extended to make the difference zero.

The size of the result equals the size of the extended dividend less the size of the divisor.

Note that the string '10.0' divided by the string '3.0' is the string '3'. It is rounded to ZERO decimal places.

### Changes:

The contents of argument 2 are changed to contain the result.

### Errors Recognized:

- Improper argument format (parity bit set)
- Overflow occurrence (carry bit set)

### Comparison Flags:

- Result was zero (zero bit set)
- Result was negative (sign bit set)

## 7.1.7 COMPARE

Entry Point Name	COM\$
Entry Point Address	10011 Octal
Argument #1 Address	D-E Registers
Argument #2 Address	H-L Registers
Result Location	Arguments unchanged. Only sets condition code
Arithmetic Function	(cond-code) = (cond [ (Argument #2) - (Argument #1) ])

### Action:

Compares two numeric string numbers as to magnitude. No change to arguments results. Changes are only made to the condition flags.

### Typical calling sequence:

```
COM$    EQU    010011
        DE     ARG1
        HL     ARG2
        CALL   COM$
```

### Arguments:

Arguments must be each numeric strings of less than 126 characters in length. Argument 1 is addressed by the D and E Registers. Argument 2 is addressed by the H and L Registers and will contain the result.

### Result:

The contents of both arguments will remain unchanged. Only the condition code will change and will obtain the exact same condition as if a call to SUB\$ were done. Therefore, the resultant condition flags will behave as if the result were to be rounded.

### Changes:

The contents of both arguments remain unchanged. Only the condition flags are changed.

### Errors Recognized:

Improper argument format (parity bit set)  
Overflow occurrence (carry bit set)

### Comparison Flags:

Result was zero (zero bit set)  
Result was negative (sign bit set)

## 7.1.8 MOVE

Entry Point Name	MOV\$
Entry Point Address	10000 Octal
Argument #1 Address	D-E Registers
Argument #2 Address	H-L Registers
Result Location	Argument #2
Arithmetic Function	(Argument #2) = (Argument #1)

### Action:

Replaces the numeric string number in argument 2 with that of argument 1, rounds and installs leading blanks and trailing zeros when needed in the result.

### Typical calling sequence:

```
MOV$    EQU    010000
        DE     ARG1
        HL     ARG2
        CALL   MOV$
```

### Arguments:

Arguments must be each numeric strings of less than 126 characters in length. Argument 1 is addressed by the D and E Registers. Argument 2 is addressed by the H and L Registers and will contain the result.

### Result:

The contents of argument 1 (D and E) will remain unchanged.  
The contents of argument 2 (H and L) will contain the number of argument 1 rounded and reformatted if necessary.

### Changes:

The contents of argument 2 are changed to contain the result.

### Errors Recognized:

Improper argument format (parity bit set)  
Overflow occurrence (carry bit set)  
[Note that overflow can occur in a MOV\$ if a move from a larger to smaller field is attempted]

### Comparison Flags:

Result was zero (zero bit set)  
Result was negative (sign bit set)

### 7.1.9 MOD10 CHECKSUM CALCULATION

Entry Point Name	MOD10\$
Entry Point Address	6006 Octal
Argument #1 Address	H-L Registers
Result Location	A-Register (no reformatting of argument)
Arithmetic Function	(A Reg) = Check-MOD-10 (Argument #1)

**Action:**

Checks validity of Modulo 10 checksum of a numeric string number.

**Typical calling sequence:**

```
MOD10$ EQU 06006
        HL ARG1
        CALL MOD10$
```

**Arguments:**

The argument must be a numeric string of less than 126 characters in length. Argument 1 is addressed by the H and L Registers.

**Result:**

- The contents of the argument remains unchanged.
- The carry bit is set if the check digit is 10.
- The zero bit is set if the check digit is not 10.
- The check digit is in the A-Register upon return.

**Changes:**

The contents of the argument remain unchanged.

**Errors Recognized:**

Improper argument format (parity bit set)

**Comparison Flags:**

- Check digit was 10 (carry bit set)
- Check digit was not 10 (zero bit set)

### 7.1.10 MOD11 CHECKSUM CALCULATION

Entry Point Name	MOD11\$
Entry Point Address	6011 Octal
Argument #1 Address	H-L Registers
Result Location	A-Register (no reformatting of argument)
Arithmetic Function	(A Reg) = Check-MOD-10 (Argument #1)

**Action:**

Verifies the Modulo 11 checksum of the numeric string number.

**Typical calling sequence:**

```
MOD11$ EQU 06011
        HL ARG1
        CALL MOD11$
```

**Arguments:**

The argument must be a numeric string of less than 126 characters in length. The argument is addressed by the H and L Registers.

**Result:**

- The contents of the argument remains unchanged.
- The carry bit is set if the check digit is 11.
- The zero bit is set if the check digit is not 11.
- The A-Register contains the check digit.

**Changes:**

The contents of the argument remain unchanged.

**Errors Recognized:**

Improper argument format (parity bit set)

**Comparison Flags:**

- Check digit was 11 (carry bit set)
- Check digit was not 11 (zero bit set)

### 7.1.11 KEYBOARD FORMATTED INPUT

Entry Point Name	KEY\$
Entry Point Address	10014 Octal
Argument #1 Address	H-L Registers
Extra Parameters	(D Reg) = Column. (E Reg) = Row for cursor
Input Function	(Argument #1) = (Keyed in number)
Input Restrictions	Screen format and, therefore, keyed in number has same format as originally in Argu- ment #1

#### Action:

Provides formatted input from the keyboard into a numeric string. The format is maintained on the screen and only a number fitting the format can be entered. The inputted numeric string is placed in argument 1.

#### Typical calling sequence:

```
KEY$    EQU    010014
        LD     COLUMN
        LE     ROW
        HL     ARG1
        CALL   KEY$
```

#### Arguments:

The argument must be a formatted numeric string. The D and E Registers must contain the column and row of the cursor position of the first character to be typed in.

#### Result:

The contents of argument 1 are replaced by the inputted number. Striking the enter key with no input will cause the argument to be replaced with a zero. The H and L Registers are pointing immediately after the ETX.

#### Changes:

The contents of the argument are replaced with the inputted string

#### Errors Recognized:

Improper argument format (parity bit set)

#### Comparison Flags:

Result was zero (zero bit set)  
Result was negative (sign bit set)

### 7.1.12 DISPLAY STRING

Entry Point Name	DSP\$
Entry Point Address	10017 Octal
Argument #1 Address	H-L Registers
Extra Parameters	(D Reg) = Column. (E Reg) = Row for cursor
Input Functions	(Display starting at D,E) = (Argument #1)
Input Restrictions	None. May even be non-numeric string

#### Action:

Displays a string onto the screen. String may be non-numeric.

#### Typical calling sequence:

```
DSP$    EQU    010017
        LD     COLUMN
        LE     ROW
        HL     ARG1
        CALL   DSP$
```

#### Arguments:

The argument may be a numeric or non-numeric string as long as it terminates with an ETX. The D and E Registers contain the column and row of the location of the first character of the string.

#### Result:

The string in argument 1 is displayed on the screen starting at the cursor location beginning with the column and row specified by the D and E Registers. The H and L Registers point the location immediately after the ETX in argument 1.

#### Changes:

The contents of the argument remain unchanged.

#### Errors Recognized:

None

#### Comparison Flags:

None

## 7.2.1 INTRODUCTION

FPAK is a subroutine package which gives the Datapoint 2200 the capability of performing numerical operations with numbers in the range of  $-10^{-38}$  to  $10^{37}$ . This is accomplished by representing all numbers in a form called "floating point." Floating point notation is a shorthand method of number representation and is very similar to the familiar "scientific notation" used in technical work.

FPAK also provides conversion of floating point numbers to and from 16 bit binary integers, particularly attractive for analyzing binary data gathered by the 2200 from instrumentation systems.

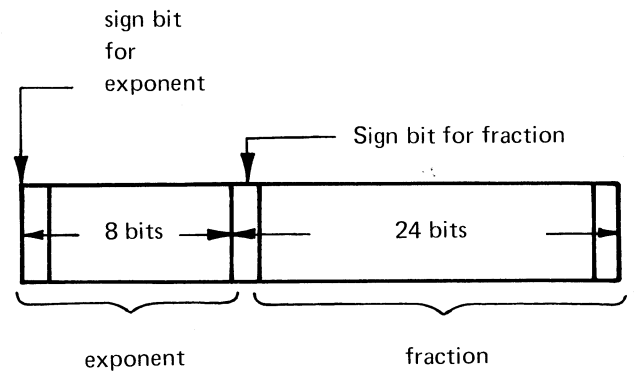
Supplied with FPAK is FCON which supplies the user with simple conversion of ASCII numeric string representation (suitable for displaying or printing) to and from the internal floating point representation.

The CTC 2200 floating point software consists of two main sections: FCON, the conversion section and FPAK, the arithmetic section. The conversion section converts a floating point number to an ASCII string, and visa versa. ASCII is the character code used by the 2200 keyboard and CRT display, and this section of the floating point software allows the user to interface with the computer. The user can enter numbers in a form familiar to him and read the results in a similar form. The ASCII string which the user enters through the keyboard (or from tape or some other means) is converted into the internal floating point form outlined in Section 7.2.1.1. When all arithmetic operations are completed, the user can request that the result be converted back into ASCII for display on the CRT (or for output to tape or for some other use).

### 7.2.1.1 INTRODUCTION TO FLOATING POINT REPRESENTATION

A number (N) in floating point form consists of two parts within the computer's memory--the "exponent" (e) and the "fraction" (f)--such that:  $N = f \cdot 2^e$  (where the \* denotes the multiplication operation). The exponent occupies one byte (word) in the 2200 and is an 8 bit signed integer. Thus, exponents on the 2200 can have a range of 127 to -128.

The fraction (sometimes called the "mantissa") on the 2200 occupies three bytes and is a 24 bit signed quantity. Like a decimal fraction (such as .5 or .0001), the fraction in a floating point number has a "decimal point," although "decimal point" is not what it is called. Its proper name, in a binary fraction, is a "binary point." In the notation used on the 2200, the binary point is located immediately to the right of the sign bit (high order bit) of the fraction. Thus, a picture of a 2200 floating point number would look like this:



The exponent and the fraction are separate parts of the number, and one can be positive while the other is negative. On the 2200, negative numbers are represented in their 2's complement form. Since the floating point representation requires more than one byte (word) on the 2200, a convention is used to address a floating point number in memory. The address of a floating point number is the byte (word) address of the exponent byte of the number. The software will use the addressed byte, and the three bytes immediately following, in whatever operation is being performed. Later in the documentation, reference is made to addressing floating point numbers. In such cases, this means that the MSP of the address of the exponent should be in the H or D register and the LSP of the address should be in the L or E register, depending upon whether the HL or DE pair is being used.

## 7.2.2 FCON - FLOATING POINT/STRING CONVERSION

### 7.2.2.1 INTRODUCTION TO FCON - FLOATING POINT/STRING CONVERSION

LOCATION	FUNCTION
FISC (0441)	Floating Internal to String Conversion entry point
FSIC (04444)	Floating String to Internal Conversion entry point
FSCE (04460)	Floating Set Conversion Error Branch entry point
OPER (013403)	Location of Floating point number to be converted to or from ASCII

```

FISC      EQU      04441
          HL        String
          CALL     FISC
          LA        015
          LMA
    
```

Note the LA 015 and LMA will install a 015 as the terminating character to the resultant string which is the ASCII representation of the floating point number in OPER.

### 7.2.2.2 FISC - FLOATING INTERNAL TO STRING CONVERSION

The Floating Internal to String Conversion routine has been designed so that the user need not specify the type of number he is going to supply. That is, as long as the ASCII characters being converted represent a valid, decimal number, the conversion routine can decide what type of number it is (i.e., integer, fraction) and perform the proper conversion without any further instructions. This type of input is referred to as "free form" input.

The result of all Floating Point Arithmetic routines end up at location 013403, labeled 'OPER'. The conversion routine, FISC, converts floating point numbers at OPER into a string beginning at the location specified by the H and L Registers upon execution of the CALL to FISC.

For example, should OPER (and the subsequent 3 bytes) contain the floating point number represented by 123,450,000,000,000,000 the string resulting from a call to FISC would look like this: 1.2345E20 where the ASCII number 1, an octal 61, would appear in the location specified by H and L and the period (an octal 056) in H and L plus 1 etc. A note of caution, FISC does not put a terminating 003 or 015 after the string. To be compatible with the CTOS string routines, the string must be terminated with either 003 or 015. However, FISC, upon return from being called, leaves the H and L registers pointing to the location immediately after the last character in the string. This enables the user to immediately store the terminating character of his choice (003 or 015) in that location upon a return. The following call to FISC will illustrate:



Name: Floating Internal to String Conversion (FISC)

### 7.2.2.3 FSCI - FLOATING STRING TO INTERNAL CONVERSION

Action:

Converts a floating point number to its ASCII character representation.

The Floating String to Internal Conversion routine has been designed to convert floating point numbers into the proper ASCII representation. If the floating point numeric string is a small integer, it will be converted to an integer, with no decimal point in the representation. If the numeric string is a large integer, or a noninteger, it will be converted into scientific notation, or more precisely what is known as the FORTRAN E format, such as 1.3456E17.

Calling Sequence:

FISC	EQU	04441
	HL	String
	CALL	FISC
	LA	015
	LMA	

FSIC converts to internal floating point representation an ASCII numeric string with optional leading minus sign, optional decimal point, and optional trailing FORTRAN E, type exponent, i.e. -1.2345E20. The H and L registers must point to the first character of the string. The result goes into the FPAK 'register' called OPER starting at 013403, ready to be used by FPAK. FSCE, Floating Set Conversion Error Branch, should be set first to cover format problems in the string being converted. A simple call to FSCE with the D and E registers specifying the location of your error recovery routine will set the error branch.

Arguments:

OPER contains the number to be converted to ASCII. The H and L registers contain the address of the location, in memory, where the first (leftmost) ASCII character should be placed.

Result:

The floating point number in OPER is converted to its ASCII representation, and the ASCII characters comprising this representation are placed in memory, beginning at the address specified by the contents of the H and L registers and continuing in sequential memory locations. H and L end up pointing to the next location after the last string character enabling the user to store the string termination character of his choice up on the return from FISC.

Changes:

The contents of OPER are destroyed; the previous contents of the output string are destroyed. At the end of the execution of this routine, H and L contain the address of the memory location immediately after the last ASCII character in the converted number.

Errors Recognized:

None.

Comments:

Numbers are represented to six significant (decimal) digits and are rounded where appropriate. The format of the output is "free," with small integer in FORTRAN I format, floating point numbers with decimal exponents between -6 and 6 in FORTRAN F format, and other numbers in FORTRAN E format.

Name: Floating String to Internal Conversion (FSIC)

Action:

Converts an ASCII string, which represents a decimal number, into that number's floating point form.

Calling Sequence:

FSIC	EQU	04444	
FSCE	EQU	04460	
	DE	ERROU	Location of error routine
	CALL	FSCE	
	HL	String	
	CALL	FSIC	

Arguments:

The H and L registers contain the address of the first byte (character) of the ASCII string which represents the number to be converted.

Result:

The character string, if it represents a valid number, is converted to a floating point number, and that value is left in OPER. The result in OPER is normalized and rounded.

Changes:

The original contents of OPER are destroyed; the ASCII string is left unchanged, and upon successful conversion, the H and L registers contain an address of the character which caused termination of the number (i.e., was a character not allowed in the ASCII representation of a number).

Errors Recognized:

Invalid character found while converting from ASCII to floating point.

Comments:

The ASCII string may be in free form, that is, in FORTRAN I, F, or E format. All of those formats will be properly converted by this routine. Conversion stops when an invalid character (something other than a digit, "E", +, -, or .) is encountered after a valid number has been found. An invalid character encountered before a valid number has been found will generate an error. Some of the above characters can be considered invalid if used incorrectly (i.e., a "." in an exponent, such as 1.333E1.5, is an error) and will generate an error condition.

#### 7.2.2.4 FSCE - FLOATING SET CONVERSION ERROR BRANCH

Name: Floating Set Conversion Error Branch (FSCE)

Action:

Specifies the location of the user's routine to be branched to in the event an invalid character is encountered while converting an ASCII representation of a number to the floating point representation of that number.

Calling Sequence:

	DE	ERROU	Location of error routine
	CALL	FSCE	

Arguments:

The D and E registers contain the address of the error routine.

Result:

The location in the floating point software which specifies the location of the error routine is set to the address provided by the user in the D and E registers.

Changes:

The previous error routine address is destroyed.

Errors Recognized:

None.

Comments:

In the event the user does not specify an error routine location, the floating point software will execute a return (RET instruction) if the error condition arises, and the arithmetic routine called by the user will continue to completion. At the conclusion of that routine, the contents of the A register will be non-zero, and the result in OPER will, in general, be erroneous. The user's error routine may end with a return if the user wishes to continue execution immediately after the call to the routine which generated the error.

## 7.2.3 FPAK - FLOATING POINT ARITHMETIC PACKAGE

### 7.2.3.1 INTRODUCTION TO FPAK - FLOATING POINT ARITHMETIC PACKAGE

LOCATION	FUNCTION
FCMP (04422)	Floating Point Compare
FADD (04400)	Floating Point Addition
FSUB (04403)	Floating Point Subtraction
FMUL (04406)	Floating Point Multiplication
FDIV (04411)	Floating Point Division
FLOD (04414)	Floating Point Load [memory to 'OPER']
FSTO (04417)	Floating Point Store ['OPER' to memory]
FNEG (04425)	Floating Point Negate [Two's complement]
FABS (04430)	Floating Point Absolute Value
FSTL (04463)	Floating Point Set Tolerance [For Equal Flag]
FFIX (04433)	Floating Point Fix [to 16 bit integer]
FFLT (04436)	Floating Point Float Conversion from 16 bit integer
FSOV (04447)	Floating Point Set Overflow Error Branch
FSUN (04452)	Floating Point Set Underflow Error Branch
FSDV (04455)	Floating Point Set Divide Check Error Branch

The second section of the floating point software is the arithmetic part. This section contains the routines for performing the common arithmetic operations of add, subtract, multiply, divide, compare, negate, and absolute value, and two routines for converting between integer and floating point formats (an integer, in the floating point software, is a 16 bit (2-byte) signed quantity which is addressed by specifying the address of the high order byte).

Within the floating point software package is a 4-byte area called OPER. OPER is to the floating point software what the A register is to the 2200 processor. Floating point operations are performed on numbers in OPER, or on pairs of numbers, one of which is in OPER and the other in memory. The software supplies two routines, FLOD and FSTO which provide the user with the capability of copying numbers from memory to OPER and from OPER to memory.

With two exceptions, all of the routines in the arithmetic part of the floating point software, which take floating point numbers as their arguments, expect their operands to be "normalized." Normalization is nothing more than an agreed upon standard for writing a floating point num-

ber. A number is considered normalized if the sign bit of the fraction and the bit immediately to the right of the sign bit (the high order bit of the fraction) are unequal. Thus, a positive fraction (sign bit 0) has a 1 as its high order bit, and a negative fraction (sign bit 1) has a 0 as its high order bit. This convention makes sure that the maximum precision possible is maintained in all floating point operations.

As a rule, all routines expect their floating point operands to be normalized. The significant exceptions to this rule are the add and subtract routines, FADD and FSUB. If the user is adding or subtracting two numbers, the numbers should be normalized for a result with the greatest accuracy possible. However, if the user has a floating point number which is not normalized, he can convert the number to its normalized form by adding or subtracting a "normal" 0 to or from the unnormalized number. A normal 0 has a fraction equal to 0 and an exponent of -128 (200 octal). Except in this case, it is not recommended that the user perform operations on unnormalized numbers.

### 7.2.3.2 ERROR CONDITIONS

There are several error conditions that can arise during the course of executing routines in the floating point software package. These errors are:

- exponent overflow
- exponent underflow
- divisor of 0 (in FDIV)

For these errors, a flag (see below) is set to 1 when the error is detected. For all of these errors, an "error branch" is provided. When the error condition arises, the appropriate flag (or A register) is set, and a jump is made to a location in the floating point software package. This location contains a jump to the address of either a user-specified error routine or a return instruction (the default case if the user does not supply an error routine). There is a separate location for each error condition, and there are three routines -- FSOV, FSUN, and FSDV -- which are used to set or change the address of the error routines.

The error conditions and their respective flags are:

Exponent Underflow	UNFLO	Location 013400
Exponent Overflow	OVFLO	Location 013401
Divide by 0	DVDCK	Location 013402

If an error condition arises, the flag is set to 1 and a branch is made to the error routine address. If no error condition arises, the flag is set to 0, and a normal return from the routine occurs.

### 7.2.3.2.1 FSOV - FLOATING SET OVERFLOW ERROR BRANCH

Name: Floating Set Overflow Error Branch (FSOV)

Action:

Specifies the location of the user's routine to be branched to in the event an operation causes exponent overflow (the value of the binary exponent in the result is greater than 127).

Calling Sequence:

Execute CALL instruction location 04447.  
See 7.2.3.2.

Arguments:

The D and E registers contain the address of the error routine.

Result:

The location in the floating point software which specifies the location of the error routine is set to the address provided by the user in the D and E registers.

Changes:

The previous error routine address is destroyed.

Errors Recognized:

None.

Comments:

In the event the user does not specify an error routine location, the floating point software will execute a return (RET instruction) if the error condition arises, and the arithmetic routine called by the user will continue to completion. At the conclusion of that routine, the appropriate error flag will be set to 1, and the result in OPER will, in general, be erroneous. The user's error routine should not end with a return since that would cause processing to continue in the floating point software with incorrect values in the machine registers.

### 7.2.3.2.2 FSUN - FLOATING SET UNDERFLOW ERROR BRANCH

Name: Floating Set Underflow Error Branch (FSUN)

Action:

Specifies the location of the user's routine to be branched to in the event an operation causes exponent underflow (the value of the binary exponent in the result is less than -128).

Calling Sequence:

Execute CALL instruction to location 04452.  
See 7.2.3.4.

Arguments:

The D and E registers contain the address of the error routine.

Result:

The location in the floating point software which specifies the location of the error routine is set to the address provided by the user in the D and E registers.

Changes:

The previous error routine address is destroyed.

Comments:

In the event the user does not specify an error routine location, the floating point software will execute a return (RET instruction) if the error condition arises, and the arithmetic routine called by the user will continue to completion. At the conclusion of that routine, the appropriate error flag will be set to 1, and the result in OPER will, in general, be erroneous. The user's error routine should not end with a return since that would cause processing to continue in the floating point software with incorrect values in the machine registers.

### 7.2.3.2.3 FSDV - FLOATING SET DIVIDE CHECK ERROR BRANCH

Name: Floating Set Divide Check Error Branch (FSDV)

Action:

Specifies the location of the user's routine to be branched to in the event the divisor in a floating divide operation is 0.

Calling Sequence:

Execute CALL instruction to location 04455.  
See 7.2.3.7.

Arguments:

The D and E registers contain the address of the error routine.

Result:

The location in the floating point software which specifies the location of the error routine is set to the address provided by the user in the D and E registers.

Changes:

The previous error routine address is destroyed.

Errors Recognized:

None.

Comments:

In the event the user does not specify an error routine location, the floating point software will execute a return (RET instruction) if the error condition arises, and the arithmetic routine called by the user will continue to completion. At the conclusion of that routine, the appropriate error flag will be set to 1, and the result in OPER will, in general, be erroneous. The user's error routine may end with a return if the user wishes to continue execution immediately after the call to the routine which generated the error.

### 7.2.3.3 FLOATING COMPARE

Name: Floating Compare (FCMP)

Action:

Compares, algebraically, two floating point numbers.

Calling Sequence:

FSTL	EQU	04463	Only necessary to EQU
FCMP	EQU	04422	once per program
	LA	TLRNC	Where TLRNC is the
	CALL	FSTL	comparison tolerance
			only necessary once per
			program if tolerance
			doesn't change
	HL	NUMBER	Number will be com-
	CALL	FCMP	pared with OPER

Arguments:

OPER contains one of the floating point numbers being compared, and the contents of the H and L registers address the other floating point number being compared.

Result:

Floating Compare sets the Sign and Zero flip-flops as if a subtraction of the floating point number addressed by the contents of the H and L registers from the floating point number in OPER had taken place. However, if the absolute value of the difference is less than or equal to the tolerance specified (see the description of the routine FSTL for an explanation of how the tolerance is specified), then the Sign and Zero flip-flops are set as if both floating point numbers were found to be equal.

Changes:

Neither operand is altered by the Floating Compare operation.

Errors Recognized:

None.

Comments:

Since representations of decimal fractions in a binary machine are approximate, the Floating Compare operation allows for an "approximate" compare by allowing the user to specify how close two numbers may be before they are considered equal.

### 7.2.3.4 FLOATING ADD

Name: Floating Add (FADD)

Action:

Adds two floating point numbers, rounds and normalizes the result.

Calling Sequence:

FSOV	EQU	04447	Only necessary to EQU once per program
FSUN	EQU	04452	
FADD	EQU	04400	
DE	OVERR		Only necessary to set these once per program or until it is desired to change.
CALL	FSOV		
DE	UNERR		Where OVERR and UNERR are addresses of user and recovery routines.
CALL	FSUN		
HL	NUMBER		
CALL	FADD		Number will be added to OPER

Arguments:

OPER contains one of the floating point numbers, and the contents of the H and L registers address the other floating point number.

Result:

The contents of OPER and the floating point number addressed by the contents of the H and L registers are added together with the result left in OPER.

Changes:

The contents of OPER are altered; the floating point number addressed by the contents of the H and L registers is unchanged.

Errors Recognized:

Exponent overflow, exponent underflow.

Comments:

Maximum precision is maintained by having both operands normalized; however, an unnormalized number may be converted to its normalized form by using this routine to add a "normal" 0 to the unnormalized number.

### 7.2.3.5 FLOATING SUBTRACT

Name: Floating Subtract (FSUB)

Action:

Subtracts two floating point numbers, rounds and normalizes the result.

Calling Sequence:

FSUB is identical to FADD except the program must now contain a FSUB EQU 04403 and the last statement in calling sequence is:

CALL	FSUB	Number will be subtracted from OPER
------	------	-------------------------------------

Arguments:

OPER contains the minuend and the contents of the H and L registers address the subtrahend.

Result:

The floating point number addressed by the contents of the H and L registers is subtracted from the floating point number in OPER, and the result is left in OPER.

Changes:

The contents of OPER are altered; the floating point number addressed by the contents of the H and L registers is unchanged.

Errors Recognized:

Exponent overflow, exponent underflow.

Comments:

Maximum precision is maintained by having both operands normalized; however, an unnormalized number may be converted to its normalized form by using this routine to subtract a "normal" 0 from the unnormalized number.

### 7.2.3.6 FLOATING MULTIPLY

Name: Floating Multiply (FMUL)

Action:

Multiplies two floating point numbers, rounds and normalizes the result.

Calling Sequence:

FMUL is identical to FADD except the program must now contain a FMUL EQU 04406 and the last statement in the calling sequence is:

```
CALL    FMUL    Number will multi-
                    ply OPER
```

Arguments:

OPER contains the multiplicand, and the H and L registers contain the address of the multiplier.

Result:

The floating point of OPER and the floating point number addressed by the contents of the H and L registers are multiplied together with the result left in OPER.

Changes:

The contents of OPER are altered; the floating point number addressed by the contents of the H and L registers is unchanged.

Errors Recognized:

Exponent overflow, exponent underflow.

Comments:

This routine expects both operands to be normalized. If one or both of the operands is not normalized, erroneous results may occur.

### 7.2.3.7 FLOATING DIVIDE

Name: Floating Divide (FDIV)

Action:

Forms the quotient of two floating point numbers, rounds and normalizes the result.

Calling Sequence:

FSOV	EQU	04447	Only necessary to
FSUN	EQU	04452	EQU these
FSDV	EQU	04455	once per
FDIV	EQU	04411	program
DE	OVERR		Only necessary to set
CALL	FSOV		these once per program
DE	UNERR		or when it is desired to
CALL	FSUN		change recover routine.
DE	CKERR		Where OVERR, UNERR,
CALL	FSDV		and CKERR are address-
			es of user error recovery
			routines.
HC	NUMBER		
CALL	FDIV		Number divides OPER

Arguments:

OPER contains the dividend, and the H and L register contain the address of the divisor.

Result:

The floating point number in OPER is divided by the floating point number addressed by the contents of the H and L registers with the result left in OPER.

Changes:

The contents of OPER are altered; the floating point number addressed by the contents of the H and L registers is unchanged.

Errors Recognized:

Exponent overflow, exponent underflow, divisor equal to 0.

Comments:

This routine expects both operands to be normalized. If one or both of the operands is not normalized, erroneous results may occur.

### 7.2.3.8 FLOATING LOAD

Name: Floating Load (FLOD)

Action:

Copies a floating point number from its location in memory to OPER.

Calling Sequence:

FLOD	EQU	04414	Only necessary to EQU this once per program.
	HC CALL	NUMBER FLOD	Number is loaded into OPER

Arguments:

The H and L registers contain the address of the floating point number that is to be copied into OPER.

Result:

The floating point number addressed by the H and L registers is copied into OPER.

Changes:

The original contents of OPER are destroyed. The floating point number addressed by the contents of the H and L registers is unchanged.

Errors Recognized:

None.

Comments:

None.

### 7.2.3.9 FLOATING STORE

Name: Floating Store (FSTO)

Action:

Copies a floating point number from OPER to memory.

Calling Sequence:

FSTO	EQU	04417	Only necessary to EQU this once per program.
	HL CALL	NUMBER FSTO	Number is loaded from OPER

Arguments:

The H and L registers contain the address of the location, in memory, to which the floating point number is to be copied.

Result:

The floating point number is copied into the location addressed by the contents of the H and L registers.

Changes:

The original contents of memory (4 bytes) addressed by the H and L registers are destroyed. The contents of OPER are unchanged.

Errors Recognized:

None.

Comments:

None.



### 7.2.3.10 FLOATING NEGATE

Name: Floating Negate (FNEG)

Action:

Forms the two's complement of the floating point number in OPER.

Calling Sequence:

FNEG	EQU	04425	Only necessary to EQU this once per program.
	CALL	FNEG	OPER is negated

Arguments:

OPER contains the floating point number to be negated.

Result:

The number in OPER is converted to two's complement form and then this result is normalized. The final result is left in OPER.

Changes:

The original contents of OPER are destroyed.

Errors Recognized:

Exponent overflow, exponent underflow.

Comments:

None.

### 7.2.3.11 FLOATING ABSOLUTE VALUE

Name: Floating Absolute Value (FABS)

Action:

Forms the absolute value of a floating point number.

Calling Sequence:

FABS	EQU	04430	Only necessary to EQU this once per program
	CALL	FABS	OPER becomes the absolute value of OPER

Arguments:

OPER contains the floating point number whose absolute value is to be computed.

Result:

If the contents of OPER are greater than or equal to zero, then they are left unchanged. Otherwise, the contents of OPER are negated (see the description of FNEG). In the latter case, the original contents of OPER are destroyed.

Changes:

Contents of OPER are destroyed if they are less than zero; otherwise, the contents of OPER are unchanged.

Errors Recognized:

Exponent overflow, exponent underflow.

Comments:

None.

### 7.2.3.12 FLOATING FIX

Name: Floating Fix (FFIX)

Action:

Converts a floating point number into a 16 bit integer.

Calling Sequence:

FFIX	EQU	04433	Only necessary to EQU this once per program
	HL	NUMBER	Number and number+1 will contain the 16 bit integer made from OPER
	CALL	FFIX	

Arguments:

OPER contains the floating point number to be fixed (converted to an integer), and the H and L registers contain the address, in memory, of the high order byte (upper eight bits of the integer) where the integer is to be placed.

Result:

The floating point number is converted to a 16 bit integer. If the number has a fractional part, that part is lost. The 16 bit integer is stored in memory beginning at the byte addressed by the contents of the H and L registers.

Changes:

The original contents of the 16 bits addressed by the contents of the H and L registers are destroyed. The contents of OPER are unchanged.

Errors Recognized:

None.

Comments:

If the number in OPER is such that it cannot be represented in 16 bits, only the low order 16 bits are stored in memory. Any higher order bits are lost.

### 7.2.3.13 FLOAT

Name: Float (FFLT)

Action:

Converts a 16 bit integer into a normalized floating point number.

Calling Sequence:

FFLT	EQU	04436	Only necessary to EQU this once per program
	HL	NUMBER	The 16 bit integer in number and number+1 will be converted to floating point in OPER
	CALL	FFIX	

Arguments:

The H and L registers contain the address, in memory, of the high order byte (high order eight bits) of the 16 bit integer that is to be converted.

Results:

The 16 bit integer is converted from its integer form to the floating point form, and the result is normalized and left in OPER.

Changes:

The original contents of OPER are destroyed. The 16 bit integer addressed by the H and L registers is unchanged.

Errors Recognized:

None.

Comments:

None.

#### 7.2.3.14 FLOATING SET COMPARE TOLERANCE

Name: Floating Set Compare Tolerance (FSTL)

Action:

Specifies a range in which the difference of two floating point numbers must lie for the two numbers to be considered equal.

Calling Sequence:

See 7.2.3.3.

Arguments:

The A register contains the tolerance as a positive eight bit integer (the high order bit of the A register must be 0).

Result:

The location in the floating point software which specifies the floating point compare tolerance is set to reflect the value provided by the user in the A register.

Changes:

The previous value of the tolerance is destroyed.

Errors Recognized:

None.

Comments:

When the floating point package is initialized, the tolerance is set as if the user had called FSTL with a 2 in the A register. If the value in the A register is less than .0 when FSTL is called, erroneous results may occur when using the floating compare routine, FCMP.

SECTION 8

DATAPOINT 2200

COMMUNICATIONS SUBROUTINES

## SECTION 8

### 1. INTRODUCTION

Interfacing the Datapoint 2200 with a wide range of communication facilities is a simple task. All that is needed is the 2210 ACA Communications Adaptor with the required data set or keyer option and the necessary software subroutines to drive it. The software subroutines may or may not have been written for a particular application. However, it seems likely that most users will choose to develop their own to fit their particular needs. This chapter is devoted to aiding the user in fulfilling this goal.

Understanding communications subroutines is useful for many reasons. It enables use of communication disciplines not previously used to fill a specialized need. It enables a user to develop routines that are most efficient for his particular application, it permits a user to modify previously written routines for special purposes and provides greater insight into how the communications system functions.

There is nothing difficult about the communications routines. They are just another part of the user's applications program. The routines are given special treatment here because they are used so frequently and because the terminology and hardware used for communication is foreign to many users.

In addition to the material covered in this chapter the user should be familiar with material covered in other publications on the subject of data communications. Two references that are highly recommended before embarking on any communications oriented 2200 applications are:

Bell System Data Communications Technical Reference Manual\*  
Martin, James; Teleprocessing Network Organization;  
Prentice - Hall, 1970

### 2. TYPES OF SUBROUTINES

As in most modern computers, the input/output devices used with the Datapoint 2200 are much slower than the 2200 processor. In order for an input/output (I/O) routine to be efficient it must be possible for the processor to perform other tasks (including other I/O operations) while a given I/O routine is active. One approach is to use an interrupt system in the processor to stop one routine and give control to another when an I/O operation is needed. The Datapoint 2200 does not have an interrupt system but in its place it has a very powerful subroutine calling mechanism

which permits many separate I/O routines to be "scanned" during normal execution of a program so that several I/O or other subroutines can be active at the same time.

This leads to the two possible types of communications subroutines: "in-line" and "interleaved". In-line subroutines are those routines which are written in such a way that whenever they are called they "capture" the processor until their function is complete and hence do not permit any other subroutine to be active at the same time. In many situations in-line subroutines are all that is required (such as during an automatic dialing operations when the 2200 has no other functions to perform). Interleaved subroutines are written in such a way that they return to the calling routine at regular intervals while they are active - to be called again to complete their work. Return points in communications subroutines frequently occur following status checks of external devices so that the communications subroutine does not sit in a "tight-loop" waiting for some external operation to be completed.

All of the I/O routines in the CTOS (Operating System) are in-line and would not be used during interleaved operations.

### 3. INPUT/OUTPUT OPERATIONS

In order to write any type of I/O subroutine for the Datapoint 2200 it is necessary to have a working knowledge of the input/output section of the processor. All 2200 I/O devices (including the CRT, keyboard and tape cassette decks as well as the Communications Adaptor) operate alike and have the same general I/O structure.

The basic physical details of the I/O structure are given in Part 4 of the Datapoint 2200 Reference Manual. We will deal with this system here from a programmer's point of view.

#### 3.1 Data Buses

Data flow to and from the processor takes place over a set of I/O data lines connected to the A-register in the processor. Output data is transmitted from the A-register by eight wires which at all times reflect the contents of the A-register. Whenever the content of the A-register is to be transmitted to an I/O device, one of the external command instructions is executed, which causes one of the External Command Strobes to pulse a signal to the I/O device, informing it that the data on the output bus is for it, and should be read.

---

\*Obtained through Engineering Director—Data Communications, American Telephone and Telegraph Co.,  
195 Broadway, N.Y., N.Y. 10007

Input data is transmitted to the A-register in the processor by eight wires which form a bus connected to all I/O devices. Each I/O device is so arranged that only the one currently addressed will have access to this bus. Normally, when an I/O device is first addressed, a status word is placed on this bus. The status word (or whatever is placed on the bus) is loaded into the A-register whenever an INPUT instruction is executed.

ADR	EQU	1
STATUS	EQU	2
DATA	EQU	3
WRITE	EQU	4
COM1	EQU	5
COM2	EQU	6
COM3	EQU	7
COM4	EQU	8

### 3.2 External Command Strokes

The Datapoint 2200 processor has 24 External Command Strokes in its I/O structure, only eight of which are brought to devices outside of the 2200 proper (e.g. the Communications Adaptor) and need be considered here.

These eight command lines are physically identical, and their functions are pre-assigned in the table below for the sake of consistency between I/O devices.

(In all examples following in this chapter it is assumed that all External Command labels have been defined.)

When an External Command is executed, physically all that occurs in the processor is a pulse (or strobe) on the indicated command line. All other action occurs in one of the I/O devices.

a. EX ADR is the only command strobe acted upon by all I/O devices at the same time. All other command strobes affect only the I/O device that is currently addressed.

## EXTERNAL COMMAND

COMMAND NUMBER	(exp)	OCTAL CODE	COMMAND	DESCRIPTION
1	ADR	121	Address	Selects device specified by A-register
2	STATUS	123	Sense Status	Connects selected device data lines to data input bus
3	DATA	125	Sense Data	Connects selected device data lines to data input bus
4	WRITE	127	Write Strobe	Signals selected device that output data is on data output lines
5	COM1	131	Command 1	Signals selected device that a control word is on data output lines
6	COM2	133	Command 2	Signals selected device that a control word is on data output lines
7	COM3	135	Command 3	Signals selected device that a control word is on data output lines
8	COM4	137	Command 4	Signals selected device that a control word is on data output lines

When external commands are to be used in a program the names or labels for the commands used should be defined to the assembler at the beginning of the source code listing as in the following example:

b. EX STATUS causes the selected device to place its status word on the input bus (it may already be on the bus in which case the EX STATUS does nothing).

c. EX DATA causes the selected device to place its data word on the input bus. This data will remain there until an EX STATUS or an EX ADR is executed.

d. EX WRITE- The write strobe command is a signal from the processor that data is present on the data output lines for the selected external device.

e. EX COM1 thru EX COM4 are used generally to load command words into I/O device command word registers. Depending on the device, however, they may be used for any purpose.

Device addressing in the Datapoint 2200 follows an unusual convention which the programmer should be aware. Up to 16 devices may be addressed, and the first four (low order) bits of the address word indicate which address is selected (zero through fifteen). The second four (high order) bits of the address word must contain the binary complement of the first four bits. Some of the sixteen possible addresses are reserved for specific devices. The remaining ones may be assigned as needed for a particular application.

### DEVICE ADDRESS ASSIGNMENTS

DEVICE	NUMBER	BINARY	OCTAL
Cassette Tape Decks	0	11110000	360
CRT/Keyboard	1	11100001	341
Communications Adaptor	2	11010010	322
2200P Printer	3	11000011	303
2200T Tape Transport	4	10110100	264
Unassigned	5	10100101	245
"	6	10010110	226
"	7	10000111	207
"	8	01111000	170
"	9	01101001	151
"	10	01011010	132
"	11	01001011	113
"	12	00111100	074
"	13	00101101	055
"	14	00011110	036
"	15	00001111	017

By way of example, to address (or select) the Communications Adaptor (and de-address all other devices) the following instructions are all that is required:

```
LA      0322
EX      ADR
```

### 3.3 The Input Command

In order to load the A-register with whatever is on the input bus an INPUT instruction is executed. In addition to loading the A-register with a new value, it transmits a strobe to the selected external device to inform it that the input bus has been read. Generally, if the status word is on the input bus, the input strobe is of no interest to the I/O device. However, if the data word from the I/O device is on the input bus, then the input strobe informs the I/O device that it has been read by the processor and the device then clears the read ready status bit.

### 3.4 Command Words

Through the use of the EX COM1 through EX COM4 strobes, it is possible to load command words in an I/O device, which causes the device to carry out specific instructions, or to assume some specific configuration. An excellent example of a command word structure is shown in paragraph 8.5 of the Datapoint 2200 Reference Manual. Each bit of the command word affects some aspect of the Communications Adaptor configuration and the entire operating mode of the adaptor is determined by the content of the Command Word Register at any given time.

When an EX COMn is executed, all of the bits in the affected command word register are loaded from the A-register so care must be taken that all eight bits are accounted for whenever a change is made in a command word register. Generally, when a word is loaded into a command word register, it remains there until another one replaces it. In some devices, (not the Communications Adaptor) a bit set to one will return to zero automatically when some function is carried out.

To give an example; suppose it was desired to instruct the Communications Adaptor to go "off-hook" and to "send 2025 Hz". The device would be addressed:

```
LA      0322
EX      ADR
```

and then a command word loaded

```

LA      060
EX      COM1

```

Where 060 is the octal value of the command word.

### 3.5 The Status Words

The status word provides a means of communicating to the processor the state of an I/O device at any given time. The status word is placed on the input bus whenever an I/O device is addressed, and remains there until the device is de-addressed or an EX DATA is executed. If the status of the device changes while it is selected, the value on the input bus changes with it, and may be read into the A-register without re-addressing the device. Paragraph 8.2 of the 2200 Reference Manual provides a detailed example of the status word structure used in the Communications Adaptor.

If it were desired to jump to a subroutine if the "Ring-ing Present" bit of this status word were to come true it could be coded as follows:

```

LA      0322      ADDRESS DEVICE
EX      ADR
INPUT   INPUT STATUS WORD
ND      040      MASK OFF ALL OTHER BITS
JFZ     SUBR     JUMP TO SUBROUTINE IF A CONTAINS A ONE

```

### 3.6 Character Buffers

An I/O device generally has one or more registers or buffers used to hold characters (also called "data") which are being transmitted or received by the device. Slow devices such as the keyboard usually have only one character buffer since the processor has plenty of time to read a character from the buffer before another is loaded.

Faster devices such as the Communications Adaptor have a double character buffer for transmitting or receiving data so that the processor may be reading (or writing) from one buffer while a data set (or some other external equipment) is writing (or reading) to the other buffer. This means that the processor always has at least one full character time in which to service the Communications Adaptor between data transfers.

Some even faster devices (such as the 2200T IBM Compatible Tape Deck) buffer an entire string of characters (up to 1024 in this example).

## 4. SIMPLE COMMUNICATIONS ADAPTOR ROUTINES

In writing any routine for the Communications Adaptor some simple rules must be followed. Reference should be made to Section 8 of the Datapoint 2200 Programmer's Manual in order to understand the following discussion.

Whenever data is to be transmitted or received through the Communications Adaptor, the device must first be configured for the mode of operation to be used. This is generally done with a prep subroutine which sets the Communications Adaptor Command Word (EX COM1). The transmit and receive time base registers (EX COM3 and EX COM2) and the Character Length Mask Word (EX COM4).

### 4.1 External Printers

Suppose it is desired to drive an external printer such as the Datapoint 3300P from the Communications Adaptor (The Datapoint 2200P connects directly to the I/O bus and does not use the Communications Adaptor). The 3300P is an EIA RS-232 interface serial printer, operates at 300 baud (bits/second)

uses an 8-information-bit code, and works best with two stop units.

Referring to paragraph 8.5 of the Reference Manual we see that the Command Word can be all zeros. (No data set is involved and neither transmit or received data is inverted).

Referring to paragraph 8.6 we see that to transmit 300 baud the transmit time base must be loaded with 377 followed by 000. The receive time base need not be set since we are only transmitting to a printer.

Referring to paragraph 8.7 we see that the transmitted character length mask must be 001 (binary) and the receive character length can be 000 (binary) since we are not going to receive anything. Bit 6 must be 1 since we are using the EIA-RS-232 output. The binary value of this word then is 01000001 (binary) or 101 (octal).

The following subroutine will therefore configure the Communications Adaptor for the 3300P printer:



```

PREP1  LA      0322
        EX      ADR      ADDRESS DEVICE
        LA      0
        EX      COM1     SET COMMAND WORD
        LA      0377
        EX      COM3     SET TRANSMIT TIME BASE
        LA      0
        EX      COM3
        LA      0101
        EX      COM4     SET CHARACTER LENGTH
        RET

```

This routine need only be executed once at the start of the use of the printer.

Once the Communications Adaptor is configured a subroutine must be called to transmit data to the printer. An in-line subroutine could look like this:

```

PRINT1  HL      MSG      LOAD H AND L WITH BUFFER ADDRESS OF MESSAGE TO BE TRANSMITTED
LOOP    LA      0322
        EX      ADR
        INPUT
        ND      1        MASK FOR TRANSMIT READY
        JTZ     LOOP     LOOP BACK IF NOT READY
        LAM
        EX      WRITE    TRANSMIT TO COMMUNICATIONS ADAPTOR
        CP      015     COMPARE WITH END OF MESSAGE CHARACTER
        RTZ
        CALL    INCHL*   INCREMENT H AND L
        JMP     LOOP     LOOP BACK IF NOT END OF MESSAGE

```

The above example assumes that a message has been stored in a buffer area in memory and is transmitted to the printer to the exclusion of all other activity.

A more general routine might be to transmit a single character to the printer and the return to the calling program for other activity while the printer is printing. An example of this might be as follows:

```

PRINT   LA      0322
        EX      ADR      ADDRESS DEVICE
        INPUT
        ND      1        MASK FOR TRANSMIT READY
        RTZ
        LAB
        EX      WRITE
        OR      1        MAKE SURE Z-FLAG IS SET TO 0
        RET      RETURN TO CALLING PROGRAM

```

---

\*See end of chapter for frequently used utility routines

Before calling this subroutine, B is loaded with the character to be printed and when the subroutine returns the Z-flag can be tested to see if the printer accepted the character.

#### 4.2 Non-Automatic Data Sets

Data sets that are not automatically controlled from the software such as acoustic couplers or external data sets using private line connections are generally the easiest to program and will be used as our first examples of programming for data sets.

For an example, let us program a Datapoint 2200 to interface with an acoustic coupler which will be used to call a time-sharing service and operate full-duplex at 110 baud. (This program will make the 2200 look like a typical KSR teletype machine). The main program might be written like this:

START1	CALL	PREP2	CONFIGURE COMM ADAPTOR
	LA	012	(LINE FEED)
SCAN1	CALL	DISPLY	CLEAR BOTTOM LINE OF CRT
	CALL	READ1	INPUTS CHAR FROM COMM
			ADAPTOR IF ONE READY
	JTZ	SCAN2	GO TO KYBD CHECK
HDX	CALL	DISPLY	WRITE CHAR IN A-REG TO CRT
SCAN2	CALL	KEYIN	INPUTS CHAR FROM KYBD
			IF ONE READY
	JTZ	SCAN1	CHECK COMM ADAPTOR
	CALL	WRITE1	OUTPUT KYBD CHAR TO
			COMM ADAPTOR
	JMP	SCAN1	

This is all there needs to be to the main program. When starting, a prep subroutine is called to configure the Communications Adaptor. A scanning loop is then entered which looks for characters from the Communications Adaptor or the keyboard and transmits them to their respective destinations.

If it were desired to operate the program in a half-duplex mode where the characters are displayed directly on the CRT rather than full-duplex where the characters are transmitted back from the remote computer then the last instruction in the main program should be JMP HDX rather than JMP SCAN1.

In this particular mode of operation the Command Word would have bits 0 and 4 set to one and all others set to zero (Paragraph 8.5, 2200 Reference Manual). The time base mask words would be 375 and 106 for both transmit and receive, and the Character Length Mask word would be 111 (octal). (Bit 6 is set to one since the acoustic coupler is an external data set and uses the EIA-RS-232 interface.)

The PREP2 subroutines would therefore be coded as follows:

```

PREP2  LA      0322
       EX      ADR      ADDRESS DEVICE
       LA      021      OUTPUT COMMAND WORD
       EX      COM1     '
       LA      0375     SET TRANSMIT AND RECEIVE
       EX      COM2     TIME BASES TO 110 BAUD
       EX      COM3     '
       LA      0106     '
       EX      COM2     '
       EX      COM3     '
       LA      0111     SET CHAR LENGTH MASK

       EX      COM4     TO 11-UNIT CODE
       RET

```

To input characters from the Communications Adaptor a subroutine READ1 is written. It will test the Communications Adaptor to see if a character is ready, and if so, read it. If no character is found the Z-flag is returned set to 1 and if a character is read it is returned set to zero. The code is as follows:

```

READ1  LA      0322
       EX      ADR      ADDRESS DEVICE
       IN
       ND      2        CHECK READ READY AND
                       RETURN IF NOT READY
       RTZ
       EX      DATA    PUT DATA ON INPUT BUS
       IN
       ORA
       RET          SET Z-FLAG IF CHAR = 0

```

To output characters to the Communications Adaptor a subroutine WRITE 1 is written. It will accept a character in the A-register, transmit it to the Communications Adaptor, and return to the main program when the task is finished with the character remaining in the A-register. It is coded as follows:

```

WRITE1  LBA     SAVE A IN B
        LA      0322
        EX      ADR      ADDRESS DEVICE
RETRY   IN
        ND      1        TEST FOR TRANSMIT READY
        JTZ     RETRY    AND RETRY IF BUSY
        LAB
        EX      WRITE    SEND CHAR OUT
        RET

```

The subroutines DISPLY and KEYIN are shown at the end of this chapter for information purposes. Since they do not involve the Communications Adaptor they will not be discussed here.

#### 4.3 The High Level Keyer

When the high level keyer is used it operates in every respect like an external data set except that the Command Word is set to all zeros. Bit 6 of the Character Length Mask is set to one.

NUMBER DC '9\*5125551234',015

would cause 9 to be dialed, then a pause, then 512-555-1234 to be dialed then control transferred to the calling program.

All other characters in the buffer area are ignored.

### 5. AUTOMATIC DATA SET OPERATION

One of the major features of the Datapoint 2200 is its ability to operate with the telephone network, providing completely automatic call origination and answering.

#### 5.1 Automatic DDD Network Call Origination.

Automatic Call origination requires the Communications Adaptor to be provided with either a 103 or 202 internal data set option. These data sets interface with the telephone network through a Bell System Direct Access Arrangement (DAA). (See the Datapoint 2200 Installation Manual for specific details).

To automatically originate a call the following events must occur:

- a. The DAA must have been on-hook long enough to assure complete termination of any previous call.
- b. The Communications Adaptor must be configured for an automatic dialing mode.
- c. The DAA must be set "Off-hook" and the dial tone present bit tested for ready (one).
- d. The desired number transmitted.
- e. The Communications Adaptor configured for the type of data set used and the connection confirmed (answered by another data set). (If the call is not confirmed within a reasonable time, usually about 30 seconds, a retry is probably indicated about 3 to 5 times).
- f. Normal data transmission occurs.
- g. The DAA is set to "on-hook" as soon as the connection is no longer desired.

The following code (page 8-9) provides an example of a complete automatic call origination sequence up to the point of reconfiguring the Communications Adaptor for the particular data set used (Step e. above). The number to be dialed is assumed to have been previously stored in an ASCII character sequence in a buffer area in memory beginning at NUMBER. An ASCII '\*' (052) between digits results in an extra delay between dial pulses when such might be required to obtain an outside line in a private exchange or for some other reason. The end of the number is indicated by an ASCII return (015). For example:

DIAL	HL	PHNUMB	BUFFER POINTER
	LA	0322	
	EX	ADR	ADDRESS DEVICE
	SUA		
	EX	COM1	SET DAA ON-HOOK
	DE	10000	
	CALL	DELAY	DELAY 5 SECONDS
	LA	0330	CONFIGURE FOR DIALING; OFF-HOOK,
	EX	COM1	INVERT DATA, SEND DIAL PULSES.
	LA	0375	SET 100 BAUD (10 CPS) DIAL RATE
	EX	COM3	
	LA	0	
	EX	COM3	
	EX	COM4	
DTONE	IN		WAIT FOR DIAL TONE
	ND	0200	'
	JTZ	DTONE	'
	JMP	LDIG	GET FIRST DIGIT
NEXDIG	CALL	INCHL	INCREMENT H AND L
LDIG	LAM		
	CP	'**'	IF THE A-REG CONTAINS '**' THEN
	JFZ	CMPR	CALL 5 SECOND DELAY
	DE	10000	
	CALL	DELAY	'
	JMP	NEXDIG	
CMPR	CP	015	
	RTZ		RETURN IF END OF NUMBER
	CP	'9'+1	TEST FOR VALID DIGIT
	JFS	ERR1	'
	CP	'0'	'
	JTS	ERR1	'
	JFZ	MASK	
	LA	10	CHANGE ZERO TO TEN
MASK	ND	017	MASK-OFF HIGH ORDER BITS
	LBA		SAVE A IN B
PLOOP	IN		WAIT FOR TRANSMIT READY
	ND	1	
	JTZ	PLOOP	'
	LA	0360	
	EX	WRITE	SEND DIAL PULSE
	LAB		DECREMENT PULSE COUNTER
	SU	1	'
	LBA		
	DE	2000	
	JFZ	PLOOP	
	CALL	DELAY*	DELAY ONE SECOND
	JMP	NEXDIG	

Upon returning from the DIAL subroutine the Communications Adaptor should be reconfigured for the type of data set used and the status bit tested for main channel carrier present. If it is not received within 30 seconds the call should be terminated and retried. The following code shows how this could be done for a 103 type data set operating at 150 baud.

---

\*See end of chapter for frequently used utility routines



c. If ringing is detected configure the Communications Adaptor for the type of data set used and set Command Word bit 5 for off-hook (1).

d. Depending on the type of data set, test for a received carrier (main channel in a type 103, main or supervisory channel depending on initial direction of communications in a 202 type). If no carrier is received after 30 seconds return to step a. above). If normal carrier is received then continue with normal communications.

## 6. FREQUENTLY USED SUBROUTINES

### 6.1 INCHL

This subroutine is used to increment the value stored in the H and L register as a double precision (16-bit) number.

```
INCHL  LAL
        AD      1
        LLA
LA      LAH
        AC      0
        LHA
        RET
```

### 6.2 DELAY

This subroutine provides a means for a time delay up to 30 seconds. Before calling the routine a double precision number is loaded into the D and E registers using the DE macro. This number is decremented at a rate of 2000 counts per second until D and E are zero and then the subroutine returns to the calling program.

```
DELAY  LAE
        SU      1
        LEA
        LAD
        SB      0
        LDA
        JFZ     DELAY
        ADE
        RTZ
        JMP     DELAY
```

### 6.3 DISPLY

This routine accepts a character in the A-register and displays it on the CRT screen at the current cursor position and then increments the cursor to the next position. Characters are always entered on the bottom line of the screen and the screen is rolled up one line whenever an ASCII line-feed is received (012). The character displayed is in the A-register when the routine returns.

DISPLY	LBA		SAVE A IN B
	LA	0341	ADDRESS DEVICE
	EX	ADR	
	LAB		LOAD A FROM B AND
	ND	0177	MASK PARITY BIT
	CP	015	TEST FOR CR
	JTZ	CRDET	
	CP	012	TEST FOR LF
	JTZ	LFDET	
	CP	040	TEST FOR VALID
	RTS		ASCII CHARACTER
	CP	0177	(RUBOUT)
	RTZ		
	EX	WRITE	
	HL	CURPOS	INCREMENT CURSOR POS
	LAM		
	AD	1	
	CP	80	
	JFS	OFDET	
	LMA		SAVE CURSOR POS
	LCA		SAVE A IN C
WCOMP	IN		TEST FOR WRITE DONE
	ND	1	
	JTZ	WCOMP	
	HL	CMDWRD	GET COMMAND WORD
	LAM		
	EX	COM1	
	LA	020	
	LMA		RESTORE COMMAND WORD
	LAC		
	EX	COM2	WRITE NEW CURSOR POS
	LA	11	MAINTAIN CURSOR ON
	EX	COM3	BOTTOM LINE
	LAB		RESTORE CHAR TO A
	RET		
LFDET	HL	CMDWRD	
	LA	030	SET NEW
	LMA		COMMAND WORD
	HL	CURPOS	
	LCM		LOAD CURSOR POS
	JMP	WCOMP	
LFDET	HL	CMDWRD	
	LA	030	SET NEW
	LMA	030	COMMAND WORD
	LC	0	
	HL	CURPOS	SET NEW CURSOR
	LMC		POSITION AND STORE
	JMP	WCOMP	
CURPOS	DC	0	
CMDWRD	DC	020	
CRDET	EQU	LFDET	



#### 6.4 KEYIN

This subroutine is used to scan the keyboard and if a character is present return it to the calling program in the A-register. If the keyboard switch is held down during a keyboard entry, bit six of the data word is set to 0 allowing upper case ASCII characters to be converted to ASCII control characters (e.g., upper case J is converted to ASCII line-feed). The subroutine exits with the Z-flag set to one if no character is input and set to zero if a character is present.

KEYIN	LA	0341	ADDRESS DEVICE
	EX	ADR	
	IN		INPUT STATUS
	LBA		SAVE STATUS IN B
	ND	2	
	RTZ		RETURN IF READ NOT READY
	LAB		RESTORE STATUS
	ND	4	MASK FOR KYBD SENSE SW
	EX	DATA	
	IN		READ DATA FROM KYBD
	JFZ	CCONT	JUMP IF KYBD SW SET
	ORA		RESET Z-FLAG
	RET		
CCONT	ND	077	MASK BIT 6
	RET		

**SECTION 9**

**DATAPoint 2200**

**OPERATING SYSTEM LISTING**

. PARITY CHECK THE BOOTSTRAPED DATA

00000	066	050	056	000	CKLOAD:	HL	PSTART	
00004	036	000				LD	\$\$	INITIALIZE XOR CHECK
00006	046	000				LE	\$\$	INITIALIZE CIRCLE CHECK
00010	307				CKLOOP:	LAM		GET A BYTE
00011	320					LCA		SAVE IT
00012	253					XRD		ACCUMULATE THE XOR PARITY
00013	330					LDA		
00014	302					LAC		
00015	254					XRE		ACCUMULATE THE CIRCLE PARITY
00016	012					SRC		
00017	340					LEA		
00020	306					LAL		INCREMENT HL
00021	004	001			HALT:	AD	1	
00023	360					LLA		
00024	305					LAH		
00025	014	000				AC	0	
00027	350					LHA		
00030	074	002				CP	PEND>8	STOP WHEN PAST END
00032	110	010	000			JFZ	CKLOOP	
00035	306					LAL		
00036	074	000				CP	PEND	
00040	110	010	000			JFZ	CKLOOP	
00043	303					LAD		CHECK THE PARITY ACCUMULATIONS
00044	264					ORE		
00045	110	022	000			JFZ	HALT+1	
00050	066	054	056	000	PSTART:	HL	SCLOOP	CLEAR LOW CORE TO HALT SHORT LOADS
00054	306				SCLOOP:	LAL		DECREMENT MEMORY POINTER
00055	024	001				SU	1	
00057	360					LLA		
00060	373					LMD		CLEAR THE LOCATION
00061	110	054	000			JFZ	SCLOOP	GO UNTIL LOCATION ZERO CLEAR

. BOOTSTRAP LOADS THE ZEROth FILE

00064	016	000			BOOT\$#	LB	0	LOAD FILE ZERO
00066	106	100	000			CALL	LOAD\$	
00071	100	075	000			JFC	RUN\$	EXECUTE IF LOAD WAS OKAY
00074	377					HALT		
00075	104	064	000		RUN\$#	JMP	BOOT\$	OVERSTORED WITH STARTING ADDRESS

. 2200 BINARY IMAGE FILE LOADER  
 . UPON ENTRY THE B REGISTER SHOULD CONTAIN  
 . THE DESIRED FILE NUMBER (POSITIVE)  
 . FILE LABEL RECORD FORMAT: 0201/0176/N/-N  
 . DATA RECORD FORMAT: 0303/074/XP/CP/H/L/-H/-L/DATA...  
 . THE 0303/074 INDICATES NUMERIC TYPE DATA  
 . H AND L DEFINE THE STARTING ADDRESS  
 . XP IS THE XOR PARITY AND CP IS THE CIRCULAR PARITY  
 . FOR THE CHARACTERS FOLLOWING THE CP

00100	006	360			LOAD\$#	LA	0360	ADDRESS THE CASSETTE MECHANISM
-------	-----	-----	--	--	---------	----	------	--------------------------------

00102	121			EX	ADR		
00103	106	322	001	CALL	STOP	STOP ANY TAPE MOTION	
00106	155			EX	DECK1	SELECT THE SYSTEM DECK	
00107	104	121	000	JMP	LOAD		
00112	006	360		LOAD2\$#	LA	0360	ADDRESS THE CASSETTE MECHANISM
00114	121			EX	ADR		
00115	106	322	001	CALL	STOP	STOP ANY TAPE MOTION	
00120	157			EX	DECK2	SELECT THE DATA DECK	
00121	106	323	001	LOAD:	CALL	DWAIT	WAIT FOR DECK SELECTION
00124	301				LAB		THE REQUESTED FILE NUMBER MUST BE
00125	260				ORA		POSITIVE
00126	160	270	001		JTS	ARGH	
00131	066	077	056	000	HL	RUN\$+2	INITIALIZE THE STARTING LOCATION MSB
00135	250				XRA		FOR 'NOTHING LOADED' FLAG
00136	370				LMA		
00137	104	012	001		JMP	FSTART	

## . SEARCH FOR THE DESIRED FILE

00142	106	360	001	FWAIT:	CALL	GETCH	WAIT FOR END OF RECORD
00145	100	142	000		JFC	FWAIT	
00150	106	350	001	FNEXT:	CALL	RTINIT	INITIALIZE THE RE-TRY COUNT
00153	026	006		FREAD:	LC	6	WAIT FOR DATA OR LEADER
00155	106	325	001		CALL	TWAIT	
00160	044	002			ND	2	QUIT IF LEADER
00162	110	270	001		JFZ	ARGH	
00165	106	360	001		CALL	GETCH	GET THE RECORD TYPE
00170	330				LDA		SAVE IT
00171	106	360	001		CALL	GETCH	GET THE RECORD TYPE COMPLEMENTED
00174	054	377			XR	0377	UN-COMPLEMENT IT
00176	273				CPD		THE TWO MUST MATCH
00177	110	244	000		JFZ	FSTOP	
00202	074	303			CP	0303	IGNORE NUMERIC RECORDS
00204	150	142	000		JTZ	FWAIT	
00207	074	347			CP	0347	IGNORE SYMBOLIC RECORDS
00211	150	142	000		JTZ	FWAIT	
00214	074	201			CP	0201	ELSE IT MUST BE AN EOF RECORD
00216	110	244	000		JFZ	FSTOP	
00221	106	360	001		CALL	GETCH	GET THE FILE NUMBER
00224	330				LDA		SAVE IT
00225	106	360	001		CALL	GETCH	GET THE FILE NUMBER COMPLEMENTED
00230	054	377			XR	0377	UN-COMPLEMENT IT
00232	273				CPD		MAKE SURE THE TWO MATCH
00233	110	244	000		JFZ	FSTOP	
00236	106	360	001		CALL	GETCH	MAKE SURE THIS IS THE END OF THE RECORD
00241	140	262	000		JTC	WCHWAY	
00244	106	322	001	FSTOP:	CALL	STOP	STOP THE TAPE
00247	167				EX	BSP	BACK UP OVER THE RECORD
00250	106	334	001		CALL	DECRTO	DECREMENT THE RE-TRY COUNT
00253	160	270	001		JTS	ARGH	QUIT IF TOO MANY RE-TRIES
00256	171				EX	SF	RE-INITIATE FORWARD MOTION
00257	104	153	000		JMP	FREAD	
00262	303			WCHWAY:	LAD		SEE IF WE ARE THERE YET
00263	271				CPB		

00264	160	150	000	JTS	FNEXT	KEEP GOING IF NOT FAR ENOUGH
00267	150	037	001	JTZ	NXTREC	START LOADING IF THERE
00272	106	322	001	CALL	STOP	ELSE STOP THE TAPE
00275	173			EX	SB	AND START SEARCHING BACKWARD
00276	106	350	001	BWAIT:	CALL RTINIT	INITIATE THE RE-TRY COUNT
00301	026	006		LC	6	WAIT FOR DATA OR LEADER
00303	106	325	001	CALL	THAIT	
00306	044	002		ND	2	QUIT IF LEADER
00310	110	270	001	JFZ	ARGH	
00313	365			BREAD:	LLH	PUSH THE CHAR ONTO THE STACK
00314	354			LHE		
00315	343			LED		
00316	330			LDA		
00317	106	360	001	CALL	GETCH	GET THE NEXT RECORD CHARACTER
00322	100	313	000	JFC	BREAD	
00325	304			LAE		GET THE RECORD TYPE COMPLEMENTED
00326	054	377		XR	0377	UN-COMPLEMENT IT
00330	273			CPD		IT MUST MATCH THE TYPE
00331	110	021	001	JFZ	BSTOP	
00334	074	303		CP	0303	IGNORE NUMERIC RECORDS
00336	150	276	000	JTZ	BWAIT	
00341	074	347		CP	0347	IGNORE SYMBOLIC RECORDS
00343	150	276	000	JTZ	BWAIT	
00346	074	201		CP	0201	ELSE IT MUST BE AN EOF RECORD
00350	110	021	001	JFZ	BSTOP	
00353	306			LAL		GET THE FILE NUMBER COMPLEMENTED
00354	054	377		XR	0377	UN-COMPLEMENT IT
00356	225			SUH		MAKE SURE IT MATCHES THE FILE NUMBER
00357	110	021	001	JFZ	BSTOP	
00362	340			LEA		FLIP OVER THE FILE NUMBER
00363	026	010		LC	8	
00365	305			FLIP:	LAH	
00366	012			SRC		
00367	350			LHA		
00370	304			LAE		
00371	210			ACA		
00372	340			LEA		
00373	302			LAC		
00374	024	001		SU	1	
00376	320			LCA		
00377	110	365	000	JFZ	FLIP	
00402	304			LAE		COMPARE IT TO THE DESIRED FILE NUMBER
00403	271			OPB		
00404	160	270	001	JTS	ARGH	IT AINT THERE
00407	110	276	000	JFZ	BWAIT	WE HAVEN'T GONE BACK FAR ENOUGH
00412	106	322	001	FSTART:	CALL STOP	ELSE STOP THE TAPE
00415	171			EX	SF	AND START GOING FORWARD AGAIN
00416	104	150	000	JMP	FNEXT	
00421	106	322	001	BSTOP:	CALL STOP	TRY THAT RECORD IN REVERSE AGAIN
00424	161			EX	RBK	
00425	106	334	001	CALL	DECRTO	DECREMENT THE RE-TRY COUNT
00430	160	270	001	JTS	ARGH	QUIT IF TOO MANY RE-TRIES
00433	173			EX	SB	RE-INITIATE BACKWARD MOTION
00434	104	313	000	JMP	BREAD	

## . READ IN A DATA RECORD HEADER

00437	106	350	001	NXTREC:	CALL	RTINIT	INITIALIZE THE RE-TRY COUNT
00442	026	020		NXTWAIT:	LC	020	WAIT FOR IRG
00444	106	325	001		CALL	TWAIT	
00447	106	360	001	NEXTRY:	CALL	GETCH	GET THE RECORD TYPE
00452	140	047	001		JTC	NEXTRY	WAIT FOR DATA
00455	330				LDA		SAVE THE RECORD TYPE
00456	106	360	001		CALL	GETCH	GET THE RECORD TYPE COMPLEMENTED
00461	054	377			XR	0377	UN-COMPLEMENT IT
00463	273				CPD		THE TWO MUST MATCH
00464	110	304	001		JFZ	AGAIN	
00467	074	347			CP	0347	IGNORE SYBOLIC RECORDS
00471	150	042	001		JTZ	NXTWAIT	
00474	074	303			CP	0303	LOAD NUMERIC RECORDS
00476	150	130	001		JTZ	NXTONE	
00501	074	201			CP	0201	QUIT ON EOF MARKER
00503	110	304	001		JFZ	AGAIN	
00506	106	322	001		CALL	STOP	STOP THE TAPE
00511	167				EX	BSP	BACK UP TO THE END OF THE FILE
00512	106	323	001		CALL	DWAIT	
00515	066	077	056	000	HL	RUN\$+2	MAKE SURE SOMETHING WAS LOADED
00521	307				LAM		
00522	260				ORA		
00523	150	270	001		JTZ	ARGH	ERROR EXIT IF NOT
00526	250				XRA		ELSE SET THE ZERO CONDITION
00527	007				RET		AND QUIT
00530	106	360	001	NXTONE:	CALL	GETCH	GET THE PARITY INITIALIZATION VALUES
00533	350				LHA		IN H (XP) AND L (CP)
00534	106	360	001		CALL	GETCH	
00537	360				LLA		
00540	106	360	001		CALL	GETCH	GET THE STARTING ADDRESS IN DE
00543	330				LDA		
00544	106	360	001		CALL	GETCH	
00547	340				LEA		
00550	106	360	001		CALL	GETCH	GET IT AGAIN FOR A CHECK
00553	054	377			XR	0377	IT IS COMPLEMENTED THIS TIME
00555	273				CPD		
00556	110	304	001		JFZ	AGAIN	
00561	106	360	001		CALL	GETCH	
00564	140	304	001		JTC	AGAIN	CATCH THE RECORD BEING OVER ALREADY
00567	054	377			XR	0377	UN-COMPLEMENT
00571	274				CPE		
00572	110	304	001		JFZ	AGAIN	
00575	306				LAL		SAVE THE PARITY ACCUMULATORS
00576	325				LOH		
00577	066	076	056	000	HL	RUN\$+1	STORE THE STARTING ADDRESS IN RUN\$ JUMP
00603	374				LME		
00604	066	077			LL	RUN\$+2	
00606	373				LMD		
00607	353				LHD		SET STORAGE POINTER TO STARTING ADDRESS
00610	364				LLE		
00611	332				LDC		RESTORE THE PARITY ACCUMULATORS

00612	255			XRH		ACCUMULATE IN THE STARTING ADDRESS
00613	012			SRC		
00614	256			XRL		
00615	012			SRC		
00616	255			XRH		
00617	012			SRC		
00620	256			XRL		
00621	012			SRC		
00622	340			LEA		
00623	305			LAH		

LOAD A RECORD ACCUMULATING PARITY

00624	106	360	001	NXTBYT:	CALL	GETCH	GET A BYTE OF DATA
00627	140	277	001		JTC	EOR	CATCH END OF RECORD
00632	320				LCA		ELSE SAVE IT
00633	253				XRD		ACCUMULATE THE PARITIES
00634	330				LDA		
00635	302				LAC		
00636	254				XRE		
00637	012				SRC		
00640	340				LEA		
00641	306				LAL		PREVENT LOADING INTO THE LOADER
00642	024	000			SU	PEND	
00644	305				LAH		
00645	034	002			SB	PEND>8	
00647	160	270	001		JTS	ARGH	
00652	372				LMC		STORE THE DATA IF ADDRESS OKAY
00653	306				LAL		INCREMENT THE MEMORY ADDRESS
00654	004	001			AD	1	
00656	360				LLA		
00657	305				LAH		
00660	014	000			AC	0	
00662	044	037			ND	037	DO MEMORY WRAP-AROUND
00664	350				LHA		
00665	104	224	001		JMP	NXTBYT	GET THE NEXT DATA BYTE
00670	106	322	001	ARGH:	CALL	STOP	STOP THE TAPE
00673	064	001			OR	1	INDICATE ABORTIVE EXIT WITH CARRY TOGGLE
00675	012				SRC		
00676	007				RET		
00677	303			EOR:	LAD		CHECK PARITY ACCUMULATIONS
00700	264				ORE		
00701	150	037	001		JTZ	NXTREC	
00704	106	322	001	AGAIN:	CALL	STOP	TRY THAT RECORD AGAIN
00707	167				EX	BSP	
00710	106	334	001		CALL	DECRTC	DECREMENT THE RE-TRY COUNT
00713	160	270	001		JTS	ARGH	QUIT IF TOO MANY RE-TRIES
00716	171				EX	SF	RE-INITIATE FORWARD MOTION
00717	164	047	001		JMP	NEXTRY	AND TRY THE RECORD AGAIN

UTILITY ROUTINES

00722	177			STOP:	EX	TSTOP	STOP THE TAPE
00723	026	001		DWAIT:	LC	1	WAIT FOR DECK READY

00725	123			TWAIT:	EX	STATUS	
00726	101			WAITL:	IN		
00727	242				NDC		
00730	150	326	001		JTZ	WAITL	WAIT FOR SPECIFIED STATUS
00733	007				RET		
00734	106	323	001	DEORTC:	CALL	DWAIT	WAIT FOR I/O OPERATION
00737	066	377	056 001		HL	RTC	DECREMENT THE RE-TRY COUNT
00743	307				LAM		
00744	024	001			SU	1	
00746	370				LMA		
00747	007				RET		
00750	066	377	056 001	RTINIT:	HL	RTC	INITIATE THE RE-TRY COUNT
00754	006	003			LA	3	TO TRY FOUR TIMES
00756	370				LMA		
00757	007				RET		
00760	123			GETCH:	EX	STATUS	GET A CHARACTER
00761	101				IN		
00762	044	024			ND	024	WAIT FOR DATA OR IRG
00764	150	360	001		JTZ	GETCH	
00767	002				SLC		
00770	002				SLC		
00771	002				SLC		
00772	002				SLC		
00773	043				RTC		END OF RECORD
00774	125				EX	DATA	ELSE GET THE CHARACTER
00775	101				IN		
00776	007				RET		

## . SYSTEM STORAGE

00777	000			RTC:	DC	0	RE-TRY COUNT
01000				PEND:	EGU	\$	END OF LOADER LOCATION

DONE



## OPERATING SYSTEM COMMAND DECODER

05000			SET	05000		
05000	016	001	LB	1	LOAD THE TAPE DIRECTORY	
05002	106	100	CALL	LOAD\$		
05005	100	017	JFC	GOODL	IT LOADED OKAY	
05010	066	001	HL	BDCMSG	ELSE PRINT CAT UN-LOADABLE MSG	
05014	104	046	JMP	NOCAT		
05017	006	341	GOODL:	LA	0341	KEYBOARD SWITCH OVERRIDES AUTO-LOAD
05021	121		EX	ADR		
05022	101		IN			
05023	044	004	ND	4		
05025	110	042	JFZ	OS\$		
05030	066	171	HL	ALPFN	RUN ANY AUTO-LOAD PROGRAM	
05034	307		LAM			
05035	260		ORA			
05036	310		LBA			
05037	110	201	JFZ	MAUTO\$		
05042	066	303	OS\$#:	HL	OSMSG	PRINT THE START-UP MESSAGE
05046	106	151	NOCAT:	CALL	DSPLY\$	
05051	066	367	NXTCMD:	HL	RDYMSG	PRINT 'READY'
05055	106	151	CALL	CALL	DSPLY\$	
05060	066	151	HL	CMDBUF	INPUT THE COMMAND	
05064	046	013	DE	11	POSITION THE CURSOR FOR ENTRY	
05070	026	024	LC	20	ONLY ACCEPT 20 CHARACTER	
05072	106	000	CALL	CALL	KEYIN\$	
05075	066	362	HL	CRLF		
05101	106	151	CALL	CALL	DSPLY\$	DO CRLF AFTER COMMAND ENTRY
05104	250		XRA		KEEP THE CURSOR OFF	
05105	131		EX	COM1		
05106	066	150	HL	INPTR	INITIALIZE THE SCANNER POINTER	
05112	006	151	LA	CMDBUF		
05114	370		LMA			
05115	106	316	CALL	GETSYM	GET THE COMMAND SYMBOL	
05120	066	200	HL	SYMBOL+6	CHECK THE TERMINATING CHARACTER	
05124	307		LAM			
05125	074	015	CP	015	IT MUST BE AN ENTER	
05127	150	144	JTZ	FNDCMD		
05132	074	055	CP	'-'	A DASH	
05134	150	144	JTZ	FNDCMD		
05137	074	040	CP	' '	OR A SPACE	
05141	110	264	JFZ	BADCMD		
05144	066	175	FNDCMD:	HL	SYMBOL+3	USE ONLY THE FIRST THREE CHARACTERS
05150	016	040	LB			
05152	026	003	LC	3		
05154	106	040	CALL	BLKSET		
05157	046	000	DE	CMDLST	LOOK IT UP IN THE COMMAND LIST	
05163	106	264	CALL	LOOKUP		
05166	306		LAL			
05167	044	370	ND	0370	POINT THE MEMORY POINTER TO THE	
05171	004	006	AD	6	BRANCH ADDRESS	
05173	360		LLA			
05174	347		LEM			

05175 106 353 036  
 05200 337  
 05201 066 214 056 012  
 05205 374  
 05206 066 215 056 012  
 05212 373  
 05213 106 264 012  
 05216 104 051 012

CALL INCHL  
 LDM  
 HL CBI+1  
 LME  
 HL CBI+2  
 LMD  
 CBI: CALL BADCMD  
 JMP NXTCMD

PUT THE ADDRESS IN THE JUMP INSTRUCTION

. ERROR MESSAGES

05221 066 073 056 013  
 05225 104 270 012  
 05230 066 111 056 013  
 05234 104 270 012  
 05237 066 031 056 013  
 05243 104 270 012  
 05246 066 042 056 013  
 05252 104 270 012  
 05255 066 057 056 013  
 05261 104 270 012  
 05264 066 217 056 013  
 05270 036 000  
 05272 046 013  
 05274 106 151 036  
 05277 151  
 05300 104 051 012

NAMREQ: HL NROMSG  
 JMP BADSPL  
 NONAME: HL NONMSG  
 JMP BADSPL  
 BADNAM: HL BDNMSG  
 JMP BADSPL  
 CATFUL: HL CFLMSG  
 JMP BADSPL  
 DUPNAM: HL DUPMSG  
 JMP BADSPL  
 BADCMD: HL BCMMSG  
 BADSPL: LD 0  
 LE 11  
 CALL DSPLY\$  
 EX BEEP  
 JMP NXTCMD

05303 011 000 013 000  
 05314 103 117 115 120  
 05362 011 000 013 013  
 05367 011 000 013 013  
 05401 011 000 013 013  
 05431 102 101 104 040  
 05442 114 111 102 122  
 05457 116 101 115 105  
 05473 116 101 115 105  
 05511 116 117 040 123  
 05526 101 125 124 117  
 05543 101 125 124 117  
 05557 040 040 040 040  
 05566 011 000 013 013  
 05617 127 110 101 124

OSMSG: DC 011,0,013,0,021,011,23,013,11  
 DC 'COMPUTER TERMINAL OPERATING SYSTEM',023,023,023,015  
 CRLF: DC 011,0,013,11,015  
 RDYMSG: DC 011,0,013,11,'READY',015  
 BDCMSG: DC 011,0,013,11,022,'CATALOG UNLOADABLE',015  
 BDNMSG: DC 'BAD NAME',015  
 CFLMSG: DC 'LIBRARY FULL',015  
 DUPMSG: DC 'NAME IN USE',015  
 NROMSG: DC 'NAME REQUIRED',015  
 NONMSG: DC 'NO SUCH NAME',015  
 NOAMSG: DC 'AUTO NOT SET',015  
 AUTMSG: DC 'AUTO SET TO '  
 AUTENT: DC ',015  
 OBTMSG: DC 011,0,013,11,022,'FRONT TAPE SCRATCH?',015  
 BCMMSG: DC 'WHAT?',015

05625 001  
 05626 002 016  
 05630 000  
 05631 002  
 05632 002 016  
 05634 000  
 05635 003  
 05636 004 037  
 05640 166  
 05641 005

D1PKT: DC 1 DECK ONE IS LOGICAL FILE ONE  
 DA TFRBUF  
 DC 0  
 D2PKT: DC 2 DECK TWO IS LOGICAL FILE TWO  
 DA TFRBUF  
 DC 0  
 CATPAK: DC 3 CATALOG IS LOGICAL FILE THREE  
 DA CATH  
 DC ALPFN-CATH+1  
 OBJPKT: DC 5 OBJECT FILE IS LOGICAL FILE FIVE

05642 002 016  
05644 000

DA TFRBUF  
DC 0

. CALCULATE A PHYSICAL FILE NUMBER FROM CATALOG ADDRESS

05645 024 010  
05647 012  
05650 012  
05651 012  
05652 004 002  
05654 007

NOALC: SU CAT  
SRC  
SRC  
SRC  
AD 2  
RET

. SCAN OFF A NAME AND LOOK IT UP

05655 106 316 013  
05660 074 015  
05662 110 237 012  
05665 066 172  
05667 307  
05670 074 040  
05672 150 221 012  
05675 046 010 036 037  
05701 106 264 037  
05704 306  
05705 044 007  
05707 150 230 012  
05712 306  
05713 044 370  
05715 007

GETNAM# CALL GETSYM GET THE NAME  
GETNAM: CP 015  
JFZ BADNAM TERMINATING CHARACTER MUST BE AN 015  
LL SYMBOL GET THE FIRST CHARACTER  
LAM  
CP  
JTZ NAMREQ THERE MUST BE A NAME  
GETNAX: DE CAT LOOK IT UP IN THE CATALOG  
CALL LOOKUP  
LAL  
ND 7  
JTZ NONAME IT ISN'T THERE  
LAL SET TABLE POINTER TO BEGINNING OF ENTRY  
ND 0370  
RET

. OPERATING SYSTEM LEXICAL SCANNING SUBROUTINES

05716 016 040  
05720 026 007  
05722 066 172 056 037  
05726 106 040 014  
05731 026 172  
05733 106 010 014  
05736 074 101  
05740 160 002 014  
05743 074 133  
05745 120 002 014  
05750 056 037  
05752 362  
05753 370  
05754 302  
05755 074 200  
05757 014 000  
05761 320  
05762 106 010 014  
05765 074 060  
05767 160 002 014  
05772 074 072  
05774 160 350 013

GETSYM: LB ' ' BLANK THE SYMBOL STORAGE  
LC 7  
HL SYMBOL  
CALL BLKSET  
LC SYMBOL INITIALIZE THE SYMBOL STORAGE POINTER  
CALL GETCH GET THE FIRST CHARACTER  
GETLTR: CP 'A' OR BETWEEN A AND Z  
JTS GETERM  
CP 'Z'+1  
JFS GETERM  
GETNBR: LH SYMBOL>8 STORE THE CHARACTER  
LLC  
LMA  
LAC BUMP THE STORAGE INDEX  
CP SYMBOL+6 UNLESS IT IS AT THE END OF THE STORAGE  
AC 0  
LOA  
CALL GETCH GET THE NEXT CHARACTER  
CP '0' CHECK IT'S RANGE BETWEEN 0 AND 9  
JTS GETERM  
CP '9'+1  
JTS GETNBR

05777 104 336 013  
 06002 066 200 056 037  
 06006 370  
 06007 007

JMP GETLTR  
 GETERM: HL SYMBOL+6 STORE THE TERMINATING CHARACTER  
 LMA  
 RET

. GET THE NEXT CHARACTER

06010 066 150 056 015  
 06014 307  
 06015 310  
 06016 004 001  
 06020 370  
 06021 361  
 06022 307  
 06023 074 015  
 06025 013  
 06026 066 150  
 06030 307  
 06031 024 001  
 06033 370  
 06034 250  
 06035 006 015  
 06037 007

GETCH: HL INPTR GET THE INPUT POINTER  
 LAM  
 LBA SAVE IT  
 AD 1 BUMP IT TO THE NEXT CHARACTER  
 LMA  
 LLB GET THE CHARACTER POINTED TO  
 LAM  
 CP 015 EXIT IF NOT OR  
 RFZ  
 LL INPTR ELSE DECREMENT THE CHARACTER POINTER  
 LAM  
 SU 1  
 LMA  
 XRA AND EXIT WITH ZERO CONDITION TRUE  
 LA 015 AND WITH A 015  
 RET

. SET A BLOCK OF CORE TO THE B REGISTER CONTENTS  
 . STARTING ADDRESS IN HL; NUMBER OF POSITIONS IN C

06040 371  
 06041 106 353 036  
 06044 302  
 06045 024 001  
 06047 320  
 06050 110 040 014  
 06053 007

BLKSET: LMB  
 CALL INCHL  
 LAC  
 SU 1  
 LCA  
 JFZ BLKSET  
 RET

. STORAGE

06400  
 06400 103 101 124 040  
 06406 000 017  
 06410 116 101 115 040  
 06416 127 017  
 06420 122 125 116 040  
 06426 267 022  
 06430 111 116 040 040  
 06436 263 017  
 06440 117 125 124 040  
 06446 036 020  
 06450 104 105 114 040  
 06456 147 021  
 06460 122 105 120 040  
 06466 341 020  
 06470 101 125 124 040  
 06476 344 022

TP  
 CMDLST: DC 'CAT COMMAND LIST  
 DA CATCMD  
 DC 'NAM  
 DA NAMCMD  
 DC 'RUN  
 DA RUNCMD  
 DC 'IN  
 DA INCMD  
 DC 'OUT  
 DA OUTCMD  
 DC 'DEL  
 DA DELCMD  
 DC 'REP  
 DA REPCMD  
 DC 'AUT  
 DA AUTCMD

06500 115 101 116 040  
 06506 031 023  
 06510 120 122 105 040  
 06516 073 023  
 06520 110 105 130 040  
 06526 250 023  
 06530 104 105 102 040  
 06536 200 034  
 06540 040 040 040 040  
 06546 264 012  
 06550 000  
 06551  
 06577 000  
 06600 000  
 06601 040 040 040 040  
 06612 000  
 06613 000  
 06614 000  
 07000  
 07000 000 000  
 07002

DC 'MAN  
 DA MANCMD  
 DC 'PRE  
 DA PRECMD  
 DC 'HEX  
 DA HEXCMD  
 DC 'DEB  
 DA DEBUG\$  
 DC  
 DA BADCMD  
 INPTR: DC 0  
 CMDBUF: SKIP 22  
 CATPTR: DC 0  
 OSCPTR: DC 0  
 CATSPS: DC  
 ENTSAV: DC 0  
 PFNSL: DC 0  
 PFNCTR: DC 0  
 TP  
 DC 0.0  
 TFRBUF: SKIP 254

INPUT SCANNER INDEX  
 LIBRARY CATALOG POINTER  
 CATALOG SCREEN POINTER  
 CATALOG NAME PRINT STRING  
 CATALOG ENTRY ADDRESS STORAGE  
 PHYSICAL FILE NUMBER SELECTED  
 PHYSICAL FILE NUMBER COUNTER  
 PARITY STORAGE FOR I/O ROUTINES  
 I/O TRANSFER BUFFER

. LIST THE CATALOG

07400 066 177 056 015  
 07404 006 010  
 07406 370  
 07407 066 200  
 07411 250  
 07412 370  
 07413 066 177 056 015  
 07417 367  
 07420 056 037  
 07422 307  
 07423 074 040  
 07425 150 117 017  
 07430 074 052  
 07432 150 117 017  
 07435 046 201 036 015  
 07441 026 006  
 07443 106 345 037  
 07446 066 200 056 015  
 07452 337  
 07453 303  
 07454 074 111  
 07456 160 070 017  
 07461 066 362 056 012  
 07465 106 151 036  
 07470 046 013  
 07472 066 201 056 015  
 07476 106 151 036  
 07501 066 200 056 015  
 07505 373  
 07506 066 177

CATCMD: HL CATPTR  
 LA CAT  
 LMA  
 LL OSCPTR  
 XRA  
 LMA  
 CATLOP: HL CATPTR  
 LLM  
 LH CAT>8  
 LAM  
 CP  
 JTZ CATEND  
 CP '\*'  
 JTZ CATEND  
 DE CATSPS  
 LC 6  
 CALL BLKTRF  
 HL OSCPTR  
 LDM  
 LAD  
 CP 73  
 JTS CATMOR  
 HL CRLF  
 CALL DSPLY\$  
 CATMOR: LE 11  
 HL CATSPS  
 CALL DSPLY\$  
 HL OSCPTR  
 LMD  
 LL CATPTR

INITIALIZE THE CATALOG POINTER  
 INITIALIZE THE SCREEN POSITION  
 GET THE ADDRESS OF THE NEXT CAT ENTRY  
 GET THE FIRST CHARACTER  
 LISTING IS FINISHED IF IT IS A SPACE  
 OR AN ASTERISK  
 TRANSFER NAME INTO PRINT STRING  
 GET THE CURSOR POSITION  
 SEE IF WE NEED TO GO TO A NEW LINE  
 PUT OUT CR LF IF SO  
 ALWAYS PRINT ON LINE 11  
 PRINT THE NAME  
 UPDATE THE CURSOR POSITION  
 UPDATE THE CATALOG ENTRY POSITION

07510 307  
 07511 004 010  
 07513 370  
 07514 104 013 017  
 07517 066 362 056 012  
 07523 106 151 036  
 07526 007

LAM  
 AD 8  
 LMA  
 JMP CATLOP  
 CATEND: HL CRLF  
 CALL DSPLY\$  
 RET

DO NEXT ENTRY  
 MAKE ROOM FOR NEXT COMMAND

. CHANGE THE FILE NAME

07527 106 316 013  
 07532 074 054  
 07534 110 237 012  
 07537 066 172  
 07541 307  
 07542 074 040  
 07544 150 221 012  
 07547 046 010 036 037  
 07553 106 264 037  
 07556 306  
 07557 044 007  
 07561 150 230 012  
 07564 335  
 07565 306  
 07566 044 370  
 07570 066 177 056 015  
 07574 370  
 07575 066 200  
 07577 373  
 07600 106 316 013  
 07603 074 015  
 07605 110 237 012  
 07610 066 172 056 037  
 07614 307  
 07615 074 040  
 07617 150 221 012  
 07622 046 010 036 037  
 07626 106 264 037  
 07631 306  
 07632 044 007  
 07634 110 255 012  
 07637 066 177 056 015  
 07643 347  
 07644 066 200  
 07646 337  
 07647 066 172 056 037  
 07653 026 006  
 07655 106 345 037  
 07660 104 054 023

NAMCMD: CALL GETSYM  
 CP  
 JFZ BADNAM  
 LL SYMBOL  
 LAM  
 CP  
 JTZ NAMREQ  
 DE CAT  
 CALL LOOKUP  
 LAL  
 ND 7  
 JTZ NONAME  
 LDH  
 LAL  
 ND 0370  
 HL CATPTR  
 LMA  
 LL CSOPTR  
 LMD  
 CALL GETSYM  
 CP 015  
 JFZ BADNAM  
 HL SYMBOL  
 LAM  
 CP  
 JTZ NAMREQ  
 DE CAT  
 CALL LOOKUP  
 LAL  
 ND 7  
 JFZ DUPNAM  
 HL CATPTR  
 LEM  
 LL CSOPTR  
 LDM  
 HL SYMBOL  
 LC 6  
 CALL BLKTRF  
 JMP UPDAT

GET THE OLD NAME

IT MUST BE TERMINATED BY A COMMA

THERE MUST BE A NAME  
 LOOK IT UP

IT MUST BE IN CATALOG  
 SAVE THE CATALOG POINTER

GET THE NEW NAME

THE NEW NAME MUST BE TERMINATED BY 015

THERE MUST BE A NEW NAME  
 IT MUST NOT ALREADY BE IN THE CATALOG

RESTORE THE CATALOG POINTER

TRANSFER THE SYMBOL INTO THE CATALOG

UPDATE THE CATALOG FILE

. BRING A NEW OBJECT FILE INTO THE SYSGEM

07663 106 300 017  
 07666 046 241 036 013

INCMD: CALL INGET  
 DE OBJPKT

DO THE PART COMMON WITH HEXCMD  
 GET TO THE BEGINNING OF THE INPUT FILE

07672	106	022	030	CALL	PBOF\$	
07675	104	076	021	JMP	REPFIL	
.						
07700	106	316	013	INGET: CALL	GETSYM	GET THE NAME SYMBOL
07703	074	015		CP	015	
07705	110	237	012	JFZ	BADNAM	TERMINATING CHARACTER MUST BE 015
07710	066	172		LL	SYMBOL	GET THE FIRST CHARACTER.
07712	307			LAM		
07713	074	040		CP		
07715	150	221	012	JTZ	NAMREQ	THERE MUST BE A NAME
07720	046	010	036 037	DE	CAT	LOOK UP THE NAME IN THE CATALOG
07724	106	264	037	INEXT: CALL	LOOKUP	
07727	074	052		CP	'*'	
07731	150	246	012	JTZ	CATFUL	CATALOG FULL IF FIRST CHARACTER IS *
07734	306			LAL		
07735	044	007		ND	7	
07737	110	255	012	JFZ	DUPNAM	ENTRY MUST NOT BE IN THE TABLE
07742	335			LDH		PUT THE NEW NAME IN CATALOG
07743	306			LAL		BUMP MEMORY POINTER TO START OF ENTRY
07744	044	370		ND	0370	
07746	340			LEA		
07747	066	212	056 015	HL	ENTSAV	SAVE THE CATALOG ADDRESS
07753	370			LMA		
07754	066	172	056 037	HL	SYMBOL	
07760	026	006		LC	6	
07762	106	345	037	CALL	BLKTR	
07765	066	212	056 015	HL	ENTSAV	CALCULATE THE SELECTED FILE NUMBER - 1
07771	307			LAM		
07772	106	245	013	CALL	NCALC	
07775	370			LMA		SAVE THE SELECTED FILE NUMBER
07776	024	001		SU	1	
10000	046	225	036 013	DE	D1PKT	POSITION DECK ONE TO THAT FILE
10004	106	033	030	CALL	CPFN\$	
10007	046	225	036 013	DE	D1PKT	
10013	106	022	030	CALL	PBOF\$	
10016	046	225	036 013	DE	D1PKT	GET TO THE END OF THAT FILE
10022	106	017	030	CALL	PEOF\$	SO READY TO APPEND THE NEW ONE
10025	066	212	056 015	HL	ENTSAV	AFTER THE NEW FILE MARKER RECORD
10031	307			LAM		
10032	106	174	023	CALL	D1FNH	
10035	007			RET		
.						
. OUTPUT AN ELEMENT						
10036	106	316	013	OUTCMD: CALL	GETSYM	GET THE ELEMENT NAME
10041	074	052		CP	'*'	CHECK THE TERMINATING CHAR
10043	150	166	020	JTZ	OUTALL	COPY WHOLE SYSTEM TAPE IF *
10046	074	044		CP	'\$'	
10050	150	166	020	JTZ	OUTALL	COPY ALL BUT OS AND CAT IF \$
10053	106	260	013	CALL	GETNAM	ELSE DO THE REST OF GETNAM
10056	106	245	013	CALL	NCALC	CALCULATE THE PHYSICAL FILE NUMBER
10061	046	225	036 013	DE	D1PKT	POSITION SYSTEM TAPE TO THAT FILE
10065	106	033	030	CALL	CPFN\$	

10070	046	225	036	013	DE	D1PKT	
10074	106	022	030		CALL	PBOF\$	
10077	106	073	023		CALL	PRECMD	PREP THE DATA TAPE
10102	046	241	036	013	DE	OBJPKT	POSITION TO THE OUTPUT FILE
10106	106	022	030		CALL	PBOF\$	
10111	046	225	036	013	OUTTFR: DE	D1PKT	PUT OUT THE FILE
10115	106	000	030		CALL	SNFR\$	READ A RECORD FROM THE SYSTEM TAPE
10120	140	145	020		JTC	OUTEND	CATCH END OF FILE
10123	306				LAL		CALCULATE THE LENGTH
10124	024	002			SU	TFRBUF	
10126	066	244	056	013	HL	OBJPKT+3	PUT IT IN THE OUTPUT FILE LENGTH
10132	370				LMA		
10133	046	241	036	013	DE	OBJPKT	WRITE OUT THE RECORD
10137	106	006	030		CALL	SBFN\$	
10142	104	111	020		JMP	OUTTFR	DO THE NEXT RECORD
10145	046	231	036	013	OUTEND: DE	D2PKT	PUT FILE MARKER 127 ON OUTPUT FILE
10151	006	177			LA	127	
10153	106	033	030		CALL	CPFN\$	
10156	046	231	036	013	DE	D2PKT	
10162	106	044	030		CALL	TFNWS	
10165	007				RET		
10166	066	172			OUTALL: LL	SYMBOL	THERE MUST NOT HAVE BEEN A NAME
10170	307				LAM		
10171	074	040			CP		
10173	110	237	012		JFZ	BADNAM	
10176	066	166	056	013	HL	CBTMSG	MAKE SURE THE FRONT TAPE IS SCRATCH
10202	106	151	036		CALL	DSPLY\$	
10205	151				EX	BEEP	
10206	377				HALT		
10207	006	360			LA	0360	ADDRESS DECK 2
10211	121				EX	ADR	
10212	106	146	024		CALL	DWAIT	
10215	157				EX	DECK2	
10216	106	146	024		CALL	DWAIT	
10221	175				EX	REWIND	REWIND THE TAPE
10222	106	146	024		CALL	DWAIT	
10225	066	000	056	026	HL	BOOTS	WRITE THE BOOT BLOCK
10231	046	000	036	030	DE	BOOTE	
10235	106	213	023		CALL	WBLOK	
10240	106	146	024		CALL	DWAIT	
10243	066	200	056	037	HL	SYMBOL+6	SEE IF THIS IS A FULL COPY
10247	307				LAM		OR JUST FILES 2 TO THE END
10250	024	044			SU	'\$'	
10252	150	257	020		JTZ	OUTSYS	START COPYING FROM FILE ZERO
10255	006	002			LA	2	START COPYING FROM FILE TWO
10257	046	225	036	013	OUTSYS: DE	D1PKT	
10263	106	033	030		CALL	CPFN\$	
10266	046	225	036	013	DE	D1PKT	
10272	106	022	030		CALL	PBOF\$	
10275	066	214	056	015	HL	PFNCTR	COPY THE TAPE USING FIRST HALF OF UPDATE
10301	006	377			LA	-1	
10303	370				LMA		SET UP TO START WRITING FILE MARKERS AT Z
10304	106	346	021		CALL	UPDAT0	
10307	006	177			LA	127	TERMINATE THE DATA TAPE



10311 046 231 036 013  
 10315 106 033 030  
 10320 046 231 036 013  
 10324 106 044 030  
 10327 106 146 024  
 10332 175  
 10333 106 146 024  
 10336 104 054 023

DE D2PKT WITH FILE MARKER 127  
 CALL CPFN\$  
 DE D2PKT  
 CALL TFN\$  
 CALL DWAIT  
 EX REWIND REWIND DECK 2  
 CALL DWAIT  
 JMP UPGAT

REPLACE THE NAMED FILE

10341 106 255 013  
 10344 066 212 056 015  
 10350 004 010  
 10352 370  
 10353 024 010  
 10355 106 245 013  
 10360 066 213 056 015  
 10364 370  
 10365 046 241 036 013  
 10371 106 022 030  
 10374 066 212 056 015  
 10400 367  
 10401 056 037  
 10403 307  
 10404 074 040  
 10406 150 053 021  
 10411 074 052  
 10413 150 053 021  
 10416 046 241 036 013  
 10422 106 017 030  
 10425 066 213 056 015  
 10431 307  
 10432 046 231 036 013  
 10436 106 033 030  
 10441 046 231 036 013  
 10445 106 044 030  
 10450 104 066 022  
 10453 066 213 056 015  
 10457 307  
 10460 046 225 036 013  
 10464 106 033 030  
 10467 046 225 036 013  
 10473 106 022 030  
 10476 046 241 036 013  
 10502 106 000 030  
 10505 140 132 021  
 10510 306  
 10511 024 002  
 10513 066 230 056 013  
 10517 370  
 10520 046 225 036 013  
 10524 106 006 030  
 10527 104 076 021

REPCMD: CALL GETNAM GET THE FILE NAME  
 HL ENTSAV SAVE THE CATALOG ENTRY ADDRESS  
 AD 8  
 LMA  
 SU 8  
 CALL NCALC CALCULATE THE PHYSICAL FILE NUMBER  
 HL PFNSEL SAVE IT  
 LMA  
 DE OBJPKT POSITION TO THE INPUT FILE  
 CALL PBOF\$  
 HL ENTSAV SEE IF THIS IS THE LAST ENTRY IN THE DATA  
 LLM  
 LH CAT>8  
 LAM  
 CP  
 JTZ REPUP DO SPECIAL UPDATE IF IT IS  
 CP '\*'  
 JTZ REPUP  
 DE OBJPKT POSITION TO THE END OF THE INPUT FILE  
 CALL PEOF\$  
 HL PFNSEL PUT OUT A FILE MARKER AFTER IT  
 LAM  
 DE D2PKT  
 CALL CPFN\$  
 DE D2PKT  
 CALL TFN\$  
 JMP UPDATE AND THEN DO THE NORMAL UPDATE  
 HL PFNSEL GET SELECTED FILE NUMBER  
 LAM  
 DE D1PKT POSITION SYSTEM TAPE TO THAT FILE  
 CALL CPFN\$  
 DE D1PKT  
 CALL PBOF\$  
 REPFIL: DE OBJPKT READ AN INPUT RECORD  
 CALL SNFR\$  
 JTC REPEND CATCH END OF FILE  
 LAL CALCULATE THE LENGTH  
 SU TFRBUF  
 HL D1PKT+3  
 LMA  
 DE D1PKT  
 CALL SBFW\$ WRITE THE RECORD  
 JMP REPFIL DO THE NEXT RECORD

10532 006 040  
 10534 106 174 023  
 10537 006 177  
 10541 106 174 023  
 10544 104 054 023

REPEND: LA 32  
 CALL D1FNW  
 LA 127  
 CALL D1FNW  
 JMP UPDAT

FOLLOW THE FILE BY FILE MARKERS  
 32 AND 127

UPDATE THE CATALOG FILE

DELETED A NAMED FILE

10547 106 255 013  
 10552 340  
 10553 066 212 056 015  
 10557 370  
 10560 106 245 013  
 10563 066 213 056 015  
 10567 370  
 10570 066 171 056 037  
 10574 227  
 10575 110 204 021  
 10600 370  
 10601 104 213 021  
 10604 120 213 021  
 10607 307  
 10610 024 001  
 10612 370  
 10613 304  
 10614 004 010  
 10616 056 037  
 10620 360  
 10621 307  
 10622 074 040  
 10624 150 275 021  
 10627 074 052  
 10631 150 275 021  
 10634 026 010  
 10636 106 345 037  
 10641 307  
 10642 074 040  
 10644 150 254 021  
 10647 074 052  
 10651 110 234 021  
 10654 364  
 10655 006 040  
 10657 370  
 10660 106 073 023  
 10663 046 241 036 013  
 10667 106 022 030  
 10672 104 066 022  
 10675 066 213 056 015  
 10701 307  
 10702 024 001  
 10704 046 225 036 013  
 10710 106 033 030  
 10713 046 225 036 013  
 10717 106 022 030

DELCMD: CALL GETNAM  
 LEA  
 HL ENTSAV  
 LMA  
 CALL NCALC  
 HL PFNSEL  
 LMA  
 HL ALPFN  
 SUM  
 JFZ DELDEC  
 LMA  
 JMP DELAUT  
 DELDEC: JFS DELAUT  
 LAM  
 SU 1  
 LMA  
 DELAUT: LAE  
 AD 8  
 LH CAT>8  
 LLA  
 LAM  
 OP  
 JTZ DELAST  
 OP '\*'  
 JTZ DELAST  
 DELMOV: LC 8  
 CALL BLKTRF  
 LAM  
 OP  
 JTZ DELEND  
 OP '\*'  
 JFZ DELMOV  
 DELEND: LLE  
 LA  
 LMA  
 CALL PRECMD  
 DE OBJPKT  
 CALL PBOF\$  
 JMP UPDATE  
 DELAST: HL PFNSEL  
 LAM  
 SU 1  
 DE D1PKT  
 CALL OPFN\$  
 DE D1PKT  
 CALL PBOF\$

GET THE NAMED FILE  
 SAVE IT  
 SAVE THE CATALOG ENTRY ADDRESS  
 CALCULATE THE PHYSICAL FILE NUMBER  
 SAVE IT  
 KILL AUTO PTR IF IT IS POINTING  
 TO THE FILE TO BE DELETED  
 DELETED FILE AFTER AUTO-POINTED FILE  
 ELSE BUMP DOWN THE AUTO POINTER  
 TO CORRESPOND TO CATALOG SHIFT  
 SEE IF AN ENTRY FOLLOWS  
 TAKE SPECIAL ACTION IF NOT  
 SHIFT DOWN THE CATALOG  
 DONE WHEN NO NEXT ENTRY  
 OR AT CATALOG STOP ENTRY  
 CLEAR THE LAST ENTRY VACATED  
 BY THE MOVE  
 PREP THE DATA TAPE  
 POSITION FRONT DECK TO OBJECT FILE  
 ANN DO THE NORMAL UPDATE  
 SCROG THE LAST FILE  
 POSITION THE SYSTEM TAPE TO THE  
 SELECTED FILE MINUS ONE

10722	046	225	036	013	DE	D1PKT	POSITION TO THE END OF THE FILE
10726	106	017	030		CALL	PEOF\$	
10731	066	212	056	015	HL	ENTSAV	DELETE THE ENTRY FROM THE CATALOG
10735	367				LLM		
10736	056	037			LH	CAT>8	
10740	006	040			LA		
10742	370				LMA		
10743	104	132	021		JMP	REPEND	TERMINATE TAPE AND UPDATE CATALOG
. UPDATE THE SYSTEM TAPE							
10746	066	214	056	015	UPDAT0:	HL	PFNCTR
10752	307				LAM		WRITE THE CURRENT PFN ON DECK TWO
10753	004	001			AD	1	INCREMENT THE CURRENT PFN
10755	370				LMA		
10756	046	231	036	013	DE	D2PKT	
10762	106	033	030		CALL	OPFN\$	
10765	046	231	036	013	DE	D2PKT	WRITE IT ON DECK 2
10771	106	044	030		CALL	TFNW\$	
10774	046	225	036	013	UPDAT1:	DE	D1PKT
11000	106	000	030		CALL	SNFR\$	READ A RECORD FROM DECK 1
11003	140	030	022		JTC	UPDAT2	CATCH EOF
11006	306				LAL		CALCULATE ITS LENGTH
11007	024	002			SU	TFRBUF	
11011	066	234	056	013	HL	D2PKT+3	AND PUT IT IN THE WRITE PACKET
11015	370				LMA		
11016	046	231	036	013	DE	D2PKT	
11022	106	006	030		CALL	SBFN\$	WRITE THE RECORD INCLUDING PARITIES
11025	104	374	021		JMP	UPDAT1	DO THE NEXT RECORD
11030	046	225	036	013	UPDAT2:	DE	D1PKT
11034	106	041	030		CALL	TFNR\$	READ FILE NUMBER FROM DECK 1
11037	302				LAC		
11040	074	040			OP	32	
11042	160	346	021		JTS	UPDAT0	MORE TO GO IF LESS THAN 32
11045	006	040			LA	32	ELSE PUT FILE MARKER 32 ON DECK 2
11047	046	231	036	013	DE	D2PKT	
11053	106	033	030		CALL	OPFN\$	
11056	046	231	036	013	DE	D2PKT	
11062	106	044	030		CALL	TFNW\$	
11065	007				RET		
. UPDATE:							
11066	066	213	056	015	UPDATE:	HL	PFNSEL
11072	307				LAM		GET THE SELECTED PHYSICAL FILE NUMBER
11073	066	214	056	015	HL	PFNCTR	INITIALIZE THE PFN COUNTER
11077	370				LMA		
11100	004	001			AD	1	
11102	046	225	036	013	DE	D1PKT	POSITION TO THE FILE AFTER THE ONE' SELECT
11106	106	033	030		CALL	OPFN\$	
11111	046	225	036	013	DE	D1PKT	
11115	106	022	030		CALL	PBOF\$	
11120	106	374	021		CALL	UPDAT1	COPY SYSTEM TAPE TO DATA TAPE
11123	046	241	036	013	DE	OBJPKT	POSITION DATA TAPE TO THE OBJECT FILE
11117	106	022	030		CALL	PBOF\$	

11132	066	213	056	015	HL	PFNSEL	RE-INITIALIZE THE FILE COUNTER	
11136	307				LAM			
11137	066	214	056	015	HL	PFNCTR		
11143	370				LMA			
11144	046	225	036	013	DE	D1PKT	POSITION DECK 1 TO SELECTED FILE	
11150	106	033	030		CALL	CPFN\$		
11153	046	225	036	013	DE	D1PKT		
11157	106	022	030		CALL	PBOF\$		
11162	104	174	022		JMP	UPDAT4		
11165	046	225	036	013	UPDAT3:	DE	D1PKT	WRITE A FILE NUMBER ON DECK 1
11171	106	044	030		CALL	TFNW\$		
11174	046	231	036	013	UPDAT4:	DE	D2PKT	READ A RECORD FROM DECK 2
11200	106	000	030		CALL	SNFR\$		
11203	140	230	022		JTC	UPDAT6	CATCH EOF	
11206	306				LAL		CALCULATE IT'S LENGTH	
11207	024	002			SU	TFRBUF		
11211	066	230	056	013	HL	D1PKT+3	PUT IT IN THE WRITE PACKET	
11215	370				LMA			
11216	046	225	036	013	DE	D1PKT	WRITE THE FILE	
11222	106	006	030		CALL	SBFW\$	INCLUDING THE PARITY CHARACTERS	
11225	104	174	022		JMP	UPDAT4	DO THE NEXT RECORD	
11230	066	214	056	015	UPDAT6:	HL	PFNCTR	INCREMENT THE CURRENT PFN COUNTER
11234	307				LAM			
11235	004	001			AD	1		
11237	370				LMA			
11240	046	225	036	013	DE	D1PKT	CHANGE THE PACKET NUMBER	
11244	106	033	030		CALL	CPFN\$		
11247	046	231	036	013	DE	D2PKT	READ THE NEXT FILE NUMBER FROM DECK 2	
11253	106	041	030		CALL	TFNR\$		
11256	302				LAC			
11257	074	040			CP	32		
11261	160	165	022		JTS	UPDAT3	DO THE NEXT FILE IF IT IS LESS THAN 32	
11264	104	132	021		JMP	REPEND	ELSE TERMINATE TAPE AND UPDATE CATALOG	
. LOAD AND EXECUTE A FILE								
11267	106	316	013		RUNCMD:	CALL	GETSYM	GET THE FILE NAME
11272	074	052			CP	'*'		LOAD OBJECT FILE IF *
11274	150	311	022		JTZ	RUNOBJ		
11277	106	260	013		CALL	GETNAM	ELSE LOOK UP NAME	
11302	106	245	013		CALL	NCALC	CALCULATE THE PFN	
11305	310				LBA		RUN IT	
11306	104	201	037		JMP	MAUTO\$		
11311	066	172			RUNOBJ:	LL	SYMBOL	MAKE SURE THERE
11313	307				LAM			WAS NO NAME BESIDES *
11314	074	040			CP	' '		
11316	110	237	012		JFZ	BADNAM		
11321	046	241	036	013	DE	OBJPKT	POSITION THE FILE FOR THE LOADER	
11325	106	022	030		CALL	PBOF\$		
11330	046	231	036	013	DE	D2PKT		
11334	106	025	030		CALL	BSP\$		
11337	016	001			LB	1	RUN THE OBJECT FILE	
11341	104	212	037		JMP	MAUT2\$	ON THE FRONT DECK	

## . SET THE AUTO-LOAD POINTER

```

11344 066 171 056 037 AUTCMD: HL ALPFN GET THE POINTER
11350 307 LAM
11351 260 ORA
11352 110 377 022 JFZ AUTDUP ERROR IF ALREADY SET
11355 106 255 013 CALL GETNAM ELSE GET THE NAME
11360 024 010 SU CAT CALCULATE THE FILE NUMBER
11362 012 SRC
11363 012 SRC
11364 012 SRC
11365 004 002 AD 2
11367 066 171 056 037 HL ALPFN AND SET THE POINTER
11373 370 LMA
11374 104 054 023 JMP UPCAT AND UPDATE THE CATALOG FILE
11377 024 002 AUTDUP: SU 2 CALCULATE TABLE ADDRESS
11401 002 SLC
11402 002 SLC
11403 002 SLC
11404 004 010 AD CAT
11406 360 LLA
11407 056 037 LH CAT>8
11411 046 157 036 013 DE AUTENT
11415 026 006 LC 6
11417 106 345 037 CALL BLKTFR PUT TABLE ENTRY IN STRING
11422 066 143 056 013 HL AUTMSG
11426 104 270 012 JMP BADSPL AND PRINT IT

```

## . RESET THE AUTO-LOAD POINTER

```

11431 066 171 056 037 MANCMD: HL ALPFN
11435 307 LAM
11436 260 ORA
11437 066 126 056 013 HL NOAMSG
11443 150 270 012 JTZ BADSPL AUTO IS NOT SET
11446 066 171 056 037 HL ALPFN
11452 250 XRA
11453 370 LMA

```

## . UPDATE THE CATALOG FILE

```

11454 046 235 036 013 UPCAT: DE CATPAK
11460 106 022 030 CALL PBOF$
11463 046 235 036 013 DE CATPAK
11467 106 011 030 CALL SNFW$
11472 007 RET

```

## . PREPARE A BLANK DATA TAPE

```

11473 066 166 056 013 PRECMD: HL CBTMSG WAIT FOR BLANK TAPE
11477 106 151 036 CALL DSPLY$
11502 151 EX BEEP
11503 377 HALT
11504 046 231 036 013 DE D2PKT REWIND THE DATA TAPE

```

11510	106	036	030	CALL	TRW\$	
11513	046	231	036	DE	D2PKT	WRITE A FILE NUMBER 0 ON IT
11517	006	000		LA	0	
11521	106	033	030	CALL	CPFN\$	
11524	046	231	036	DE	D2PKT	
11530	106	044	030	CALL	TFNW\$	
11533	046	231	036	DE	D2PKT	WRITE A FILE NUMBER 1 ON IT
11537	006	001		LA	1	
11541	106	033	030	CALL	CPFN\$	
11544	046	231	036	DE	D2PKT	
11550	106	044	030	CALL	TFNW\$	
11553	006	177		LA	127	
11555	046	231	036	DE	D2PKT	WRITE A FILE NUMBER 127 ON IT
11561	106	033	030	CALL	CPFN\$	
11564	046	231	036	DE	D2PKT	
11570	106	044	030	CALL	TFNW\$	
11573	007			RET		

. WRITE A FILE MARKER ON DECK 1

11574	046	225	036	013	D1FNW:	DE	D1PKT
11600	106	033	030		CALL	CPFN\$	
11603	046	225	036	013	DE	D1PKT	
11607	106	044	030		CALL	TFNW\$	
11612	007				RET		

. WRITE A BLOCK TO TAPE

11613	163			WBLOK:	EX	WBK	FIRE UP THE WRITE
11614	317			WNEXT:	LBM		GET THE DATA CHARACTER
11615	123			WWAIT:	EX	STATUS	WAIT FOR WRITE READY
11616	101				IN		
11617	044	010			ND	010	
11621	150	215	023		JTZ	WWAIT	
11624	301				LAB		WRITE THE DATA CHARACTER
11625	127				EX	WRITE	
11626	306				LAL		BUMP THE MEMORY POINTER
11627	004	001			AD	1	
11631	360				LLA		
11632	305				LAH		
11633	014	000			AC	0	
11635	350				LHA		
11636	273				OPD		SEE IF AT END OF BLOCK YET
11637	110	214	023		JFZ	WNEXT	NO CHANCE
11642	306				LAL		
11643	274				CPE		TRY LSB
11644	110	214	023		JFZ	WNEXT	
11647	007				RET		ELSE WE ARE DONE

. PUT A TSB TAPE INTO THE LIBRARY

11650	106	300	017	HEXCMD:	CALL	INGET	DO THE PART THAT IS LIKE INCMD
11653	046	231	036	DE	D2PKT		
11657	106	036	030	CALL	TRW\$		

11662	106	157	024	HEXASR:	CALL	HEXRBK	SEARCH FOR THE FIRST STARTING ADDRESS
11665	066	007			LL	HEXBUF+1	
11667	307				LAM		
11670	074	053			CP	'+'	THE FIRST CHARACTER MUST BE A +
11672	110	262	023		JFZ	HEXASR	
11675	066	010		HEXGAD:	LL	HEXBUF+2	GET THE STARTING ADDRESS
11677	106	256	024		CALL	HEXCON	
11702	140	123	024		JTC	HEXERR	IT MUST BE FOUR GOOD HEX CHARACTERS
11705	321				LCB		SAVE MSB
11706	106	256	024		CALL	HEXCON	
11711	140	123	024		JTC	HEXERR	
11714	066	004			LL	HEXADR	SAVE THE ADDRESS
11716	372				LMC		
11717	066	005			LL	HEXADR+1	
11721	371				LMB		
11722	106	157	024	HEXREC:	CALL	HEXRBK	LOAD A RECORD
11725	066	006			LL	HEXBUF	GET THE FIRST CHARACTER
11727	307				LAM		
11730	074	012			CP	012	IT MUST BE A LINE FEED
11732	110	123	024		JFZ	HEXERR	
11735	066	007			LL	HEXBUF+1	GET THE SECOND CHARACTER
11737	307				LAM		
11740	074	052			CP	'*'	IGNORE RECORD IF *
11742	150	322	023		JTZ	HEXREC	
11745	074	053			CP	'+'	GET ADDRESS IF +
11747	150	275	023		JTZ	HEXGAD	
11752	074	043		E	CP	'###'	END OF FILE IF #
11754	150	132	021		JTZ	REPEND	
11757	066	114			LL	HEXWBP	CONVERT THE HEX IN HEXBUF TO BINARY IN HEXWBF
11761	347				LEM		
11762	066	007			LL	HEXBUF+1	
11764	106	256	024	HEXCL:	CALL	HEXCON	QUIT IF NON-HEX CHARACTER
11767	140	006	024		JTC	HEXEC	SWAP E AND L
11772	306				LAL		
11773	364				LLE		
11774	340				LEA		
11775	371				LMB		STORE BINARY NUMBER
11776	306				LAL		INCREMENT AND SWAP L AND E
11777	004	001			AD	1	
12001	364				LLE		
12002	340				LEA		
12003	104	364	023		JMP	HEXCL	DO NEXT HEX PAIR
12006	307			HEXEC:	LAM		TERMINATING CHAR MUST BE 023
12007	074	023			CP	023	
12011	150	027	024		JTZ	HEXWRT	
12014	074	053			CP	'+'	UNLESS THIS BLOCK IS TO BE CONTINUED
12016	110	123	024		JFZ	HEXERR	
12021	066	114			LL	HEXWBP	IN WHICH CASE, JUST UPDATE THE WRITE BUFFER POINTER
12023	374				LME		
12024	104	322	023		JMP	HEXREC	
12027	066	114		HEXWRT:	LL	HEXWBP	ELSE RESET THE WRITE BUFFER PTR
12031	036	123			LD	HEXWBF+4	

12033	373	LMD		
12034	066 004	LL	HEXADR	PUT THE STARTING ADDRESS IN BUFFER
12036	307	LAM		
12037	066 117	LL	HEXWBF	
12041	370	LMA		
12042	054 377	XR	0377	
12044	066 121	LL	HEXWBF+2	
12046	370	LMA		
12047	066 005	LL	HEXADR+1	
12051	307	LAM		
12052	066 120	LL	HEXWBF+1	
12054	370	LMA		
12055	054 377	XR	0377	
12057	066 122	LL	HEXWBF+3	
12061	370	LMA		
12062	304	LAE		CALCULATE THE CORE BLOCK LENGTH
12063	024 123	SU	HEXWBF+4	
12065	340	LEA		
12066	066 005	LL	HEXADR+1	UPDATE THE CORE ADDRESS
12070	307	LAM		
12071	204	ADE		
12072	370	LMA		
12073	066 004	LL	HEXADR	
12075	307	LAM		
12076	014 000	AC	0	
12100	370	LMA		
12101	304	LAE		CALCULATE THE WRITE BLOCK LENGTH
12102	004 004	AD	4	COMPENSATE FOR HL GIVEN TWICE
12104	066 003 056 025	HL	HEXPKT+3	PUT THE LENGTH IN THE PACKET
12110	370	LMA		
12111	046 000 036 025	DE	HEXPKT	WRITE THE BUFFER
12115	106 011 030	CALL	SNFW\$	
12120	104 322 023	JMP	HEXREC	AND DO THE NEXT RECORD
12123	106 146 024	HEXERR: CALL	DWAIT	TRY THAT RECORD AGAIN
12126	167	EX	BSP	
12127	106 146 024	CALL	DWAIT	
12132	006 341	LA	0341	UNLESS KEYBOARD SWITCH DEPRESSED
12134	121	EX	ADR	
12135	101	IN		
12136	044 004	ND	4	
12140	150 322 023	JTZ	HEXREC	
12143	104 132 021	JMP	REPEND	
12146	026 001	DWAIT: LC	1	DECK WAIT LOOP
12150	123	TWAIT: EX	STATUS	
12151	101	IN		
12152	242	NDC		
12153	150 150 024	JTZ	TWAIT	
12156	007	RET		
12157	006 360	HEXRBK: LA	0360	MAKE SURE THE CASSETTE IS ADDRESSED
12161	121	EX	ADR	
12162	106 146 024	CALL	DWAIT	READ A BLOCK



12165	157	EX	DECK2	FROM DECK 2
12166	106 146 024	CALL	DWAIT	
12171	066 006 056 025	HL	HEXBUF	INTO HEXBUF
12175	161	EX	RBK	
12176	026 024	HEXRNX: LC	024	WAIT FOR IRG OR DATA
12200	106 150 024	CALL	TWAIT	
12203	044 020	ND	020	
12205	110 146 024	JFZ	DWAIT	QUIT IF IRG
12210	125	EX	DATA	ELSE PUT DATA INTO BUFFER
12211	101	IN		
12212	044 177	ND	0177	STRIP THE PARITY
12214	370	LMA		
12215	306	LAL		BUMP THE MEMORY POINTER
12216	004 001	AD	1	
12220	360	LLA		
12221	104 176 024	JMP	HEXRNX	
12224	024 060	HEXGET: SU	'0'	CONVERT HEX TO 4-BIT BINARY
12226	160 252 024	JTS	HEXCEN	
12231	074 012	CP	10	
12233	160 250 024	JTS	HEXLOW	
12236	024 007	SU	7	
12240	160 252 024	JTS	HEXCEN	
12243	074 020	CP	16	
12245	120 252 024	JFS	HEXCEN	
12250	260	HEXLOW: ORA		CLEAR THE CARRY TIGGLE
12251	007	RET		
12252	064 001	HEXCEN: OR	1	SET THE CARRY TIGGLE
12254	012	SRC		
12255	007	RET		
12256	307	HEXCEN: LAM		GET THE FIRST CHARACTER
12257	106 224 024	CALL	HEXGET	CONVERT IT TO BINARY
12262	043	RTC		QUIT IF NOT HEX
12263	012	SRC		PUT IT IN LEFT HALF OF BYTE
12264	012	SRC		
12265	012	SRC		
12266	012	SRC		
12267	310	LBA		SAVE IT
12270	306	LAL		BUMP THE MEMORY POINTER
12271	004 001	AD	1	
12273	360	LLA		
12274	307	LAM		GET THE SECOND CHARACTER
12275	106 224 024	CALL	HEXGET	CONVERT IT TO BINARY
12300	043	RTC		QUIT IF NOT HEX
12301	261	ORB		MERGE THE TWO HALVES
12302	310	LBA		LEAVE RESULT IN B REGISTER
12303	306	LAL		BUMP THE MEMORY POINTER AGAIN
12304	004 001	AD	1	
12306	360	LLA		
12307	007	RET		
12400		TP		

12400	001	HEXPKT: DC	1	OUTPUT FILE IS LOGICAL FILE ONE
12401	117 025	DA	HEXWBF	WRITE FROM WRITE BUFFER
12403	000	DC	0	
12404	000 000	HEXADR: DA	0	CURRENT CORE ADDRESS
12406		HEXBUF: SKIP	70	
12514	123	HEXWBP: DC	HEXWBF+4	WRITE BUFFER POINTER
12515	000 000	DC	0.0	ROOM FOR PARITY CHECKS
12517	000 000 000 000	HEXWBF: DC	0.0.0.0	ROOM FOR H AND L
12523		SKIP	128	ROOM FOR THE DATA
12723		HEXWBE: EQU	\$	
13000		SET	013000	ROOM FOR THE BOOT BLOCK
13000		BOOTS: SKIP	01000	
14000		BOOTE: EQU	\$	

DONE

14000

SET 014000

## . OPERATING SYSTEM ROUTINE ENTRY POINT TABLE

14000	104	052	030	SNFR\$#	JMP	SNFRX
14003	104	230	030	SSFR\$#	JMP	SSFRX
14006	104	376	030	SBFW\$#	JMP	SBFWX
14011	104	005	031	SNFW\$#	JMP	SNFWX
14014	104	072	031	SSFWS\$#	JMP	SSFWSX
14017	104	366	031	PEOF\$#	JMP	PEOFX
14022	104	375	031	PBOF\$#	JMP	PBOFX
14025	104	004	032	BSP\$#	JMP	BSPX
14030	104	016	032	CPDN\$#	JMP	CPDNX
14033	104	030	032	CPFN\$#	JMP	CPFNX
14036	104	075	034	TRW\$#	JMP	TRWX
14041	104	114	034	TFNR\$#	JMP	TFNRX
14044	104	127	034	TFNW\$#	JMP	TFNWX
14047	104	104	032	ERR\$#	JMP	ERRX

## . SERIAL NUMERIC FILE READ

14052	106	355	031	SNFRX:	CALL	RTCI	INITIALIZE THE RE-TRY COUNT
14055	106	152	032	SNFRS:	CALL	GETPKT	GET THE PACKET PARAMETERS
14060	106	027	033		CALL	RBK\$	START READING THE RECORD
14063	106	363	032		CALL	READ\$	GET THE RECORD TYPE
14066	330				LDA		SAVE IT
44067	106	363	032		CALL	READ\$	GET THE RECORD TYPE COMPLEMENTED
14072	054	377			XR	0377	UN-COMPLEMENT IT
14074	273				OPD		MAKE SURE THEY MATCH
14075	110	217	030		JFZ	SNFRR	TRY AGAIN IF THEY DON'T
14100	074	201			OP	0201	SEE IF IT IS A FILE MARKER
14102	150	062	032		JTZ	FEACT	QUIT IF IT IS
14105	074	347			OP	0347	SEE IF IT IS A SYMBOLIC RECORD
14107	150	074	032		JTZ	TEACT	TYPE ERROR IF IT IS
14112	074	303			OP	0303	MAKE SURE IT IS A NUMERIC RECORD
14114	110	217	030		JFZ	SNFRR	
14117	106	363	032		CALL	READ\$	GET THE PARITY CHECKS
14122	330				LDA		
14123	370				LMA		STORE PARITY IN FIRST BYTE OF BUFFER
14124	306				LAL		
14125	004	001			AD	1	
14127	360				LLA		
14130	305				LAH		
14131	014	000			AC	0	
14133	350				LHA		
14134	106	363	032		CALL	READ\$	
14137	140	217	030		JTC	SNFRR	TRY AGAIN IF RECORD OVER ALREADY
14142	340				LEA		
14143	370				LMA		STORE PARITY IN SECOND BYTE OF BUFFER
14144	306				LAL		
14145	004	001			AD	1	
14147	360				LLA		

14150	305			LAH			
14151	014	000		AC	0		
14153	350			LHA			
14154	106	363	032	SNFRL: CALL	READ\$	READ THE REST OF THE RECORD	
14157	140	205	030	JTC	SNFRE	QUIT IT AT END OF RECORD	
14162	370			LMA		STORE THE BYTE OF DATA	
14163	320			LCA		SAVE IT	
14164	253			XRD		ACCUMULATE THE PARITIES	
14165	330			LDA			
14166	302			LAC			
14167	254			XRE			
14170	012			SRC			
14171	340			LEA			
14172	306			LAL		BUMP THE MEMORY POINTER	
14173	004	001		AD	1		
14175	360			LLA			
14176	305			LAH			
14177	014	000		AC	0		
14201	350			LHA			
14202	104	154	030	JMP	SNFRL	DO THE NEXT BYTE	
14205	303			SNFRE: LAD		CHECK THE PARITY TOTALS	
14206	264			ORE			
14207	110	217	030	JFZ	SNFRR	TRY AGAIN IF THEY ARENT BOTH ZERO	
14212	106	016	033	CALL	WAIT\$	ELSE WAIT FOR THE OPERATION TO BE COMPLET	
14215	250			XRA		CLEAR THE CARRY TIGGLE	
14216	007			RET		AND RETURN	
14217	106	324	031	SNFRR: CALL	DEORTC	BACK UP AND TRY AGAIN	
14222	120	055	030	JFS	SNFRS	UNLESS RTC IS NEGATIVE	
14225	104	102	032	JMP	PEACT	IN WHICH CASE. PARITY ERROR EXIT	

## . SERIAL SYMBOLIC FILE READ

14230	106	355	031	SSFRX: CALL	RTCI	INITIALIZE THE RE-TRY COUNT	
14233	106	152	032	SSFRS: CALL	GETPKT	GET PACKET PARAMETERS	
14236	106	027	033	CALL	RBK\$	START THE READ	
14241	106	363	032	CALL	READ\$	GET THE RECORD TYPE	
14244	330			LDA		SAVE IT	
14245	106	363	032	CALL	READ\$	GET THE RECORD TYPE COMPLIMENTED	
14250	054	377		XR	0377	UN-COMPLEMENT IT	
14252	273			CPD		THEY MUST MATCH	
14253	110	365	030	JFZ	SSFRR		
14256	074	201		OP	0201	QUIT IF IT IS AN EOF RECORD	
14260	150	062	032	JTZ	FEACT		
14263	074	303		OP	0303	TYPE ERROR IF IT IS A NUMERIC RECORD	
14265	150	074	032	JTZ	TEACT		
14270	074	347		OP	0347	MAKE SURE IT IS A SYMBOLIC RECORD	
14272	110	365	030	JFZ	SSFRR		
14275	106	363	032	CALL	READ\$	INITIALIZE THE PARITY ACCUMULATORS	
14300	330			LDA			
14301	106	363	032	CALL	READ\$		
14304	340			LEA			
14305	140	365	030	JTC	SSFRR	TRY AGAIN IF THE RECORD IS OVER ALREADY	
14310	106	363	032	SSFRL: CALL	READ\$	READ THE REST OF THE RECORD	
14313	140	350	030	JTC	SSFRE	QUIT IF THE RECORD IS ENDED	

14316	260			ORA		CHECK THE VERTICAL PARITY	
14317	130	365	030	JFP	SSFRR	TRY AGAIN IF IT IS FALSE	
14322	320			LCA		SAVE THE BYTE	
14323	044	177		ND	0177	STRIP THE VERTICAL PARITY	
14325	370			LMA		STORE THE BYTE	
14326	302			LAC		ACCUMULATE THE PARITIES	
14327	253			XRD			
14330	330			LDA			
14331	302			LAC			
14332	254			XRE			
14333	012			SRC			
14334	340			LEA			
14335	306			LAL		BUMP THE MEMORY POINTER	
14336	004	001		AD	1		
14340	360			LLA			
14341	305			LAH			
14342	014	000		AC	0		
14344	350			LHA			
14345	104	310	030	JMP	SSFRL	DO THE NEXT CHARACTER	
14350	006	015		SSFRE:	LA	015	TERMINATE STRING WITH AN 015
14352	370			LMA			
14353	303			LAD		CHECK THE PARITY SUMS	
14354	264			ORE			
14355	110	365	030	JFZ	SSFRR	TRY AGAIN IF BOTH ARENT ZERO	
14360	106	016	033	CALL	WAIT\$	ELSE WAIT FOR THE OPERATION TO COMPLETE	
14363	250			XRA		CLEAR THE CARRY TOGGLE	
14364	007			RET		AND RETURN	
14365	106	324	031	SSFRR:	CALL	DECRTO	BACK UP AND TRY AGAIN
14370	120	233	030	JFS	SSFRS	UNLESS RTC IS NEGATIVE	
14373	104	102	032	JMP	PEACT	IN WHICH CASE. PARITY ERROR EXIT	
. SERIAL BLOCK FILE WRITE							
14376	106	152	032	SBFWX:	CALL	GETPKT	
14401	342			LEC			PUT THE LENGTH IN THE E REGISTER
14402	104	026	031	JMP	SBFWE		
. SERIAL NUMERIC FILE WRITE							
14405	106	152	032	SNFWX:	CALL	GETPKT	GET THE PACKET PARAMETERS
14410	106	277	032	CALL	SAVHL		SAVE THE BUFFER STARTING ADDRESS
14413	036	000		LD	0		INITIALIZE THE PARITY ACCUMULATORS
14415	046	000		LE	0		
14417	307			SNFWPG:	LAM		GENERATE THE PARITY TOTALS
14420	106	213	031	CALL	PARGEN		
14423	110	017	031	JFZ	SNFWPG		
14426	106	034	033	SBFWE:	CALL	WBK\$	START UP THE WRITE
14431	036	303		LD	0303		WRITE OUT RECORD TYPE NUMERIC
14433	106	002	033	CALL	WRITE\$		
14436	036	074		LD	074		WRITE OUT ITS COMPLEMENT
14440	106	002	033	CALL	WRITE\$		
14443	337			SNFWL:	LDM		WRITE OUT THE REST OF THE RECORD
14444	106	002	033	CALL	WRITE\$		
14447	306			LAL			BUMP THE MEMORY POINTER

14450	004	001	AD	1	
14452	360		LLA		
14453	305		LAH		
14454	014	000	AC	0	
14456	350		LHA		
14457	304		LAE		DECREMENT THE BUFFER LENGTH COUNT
14460	024	001	SU	1	
14462	340		LEA		
14463	110	043	JFZ	SNFWL	
14466	106	016	CALL	WAIT\$	WAIT FOR THE OPERATION TO BE COMPLETE
14471	007		RET		

## . SERIAL SYMBOLIC FILE WRITE

14472	106	152	032	SSFWX:	CALL	GETPKT	GET THE PACKET PARAMETERS
14475	106	277	032		CALL	SAVHL	SAVE THE START OF BUFFER ADDRESS
14500	036	000			LD	0	INITIALIZE THE PARITY ACCUMULATORS
14502	046	000			LE	0	
14504	307			SSFWPG:	LAM		GENERATE THE PARITY TOTALS
14505	074	015			CP	015	CHECK FOR END OF BUFFER
14507	150	131	031		JTZ	SSFWPS	
14512	260				ORA		GENERATE THE VERTICAL PARITY BIT
14513	170	120	031		JTP	SSFWPT	
14516	054	200			XR	0200	
14520	370			SSFWPT:	LMA		WRITE OUT CORRECTLY PARITIED CHAR
14521	026	002			LC	2	FAKE OUT PARGEN LENGTH COUNTER
14523	106	213	031		CALL	PARGEN	
14526	104	104	031		JMP	SSFWPG	
14531	106	232	031	SSFWPS:	CALL	PARSTO	
14534	106	034	033		CALL	WBK\$	START UP THE WRITE
14537	036	347			LD	0347	PUT OUT RECORD TYPE SYMBOLIC
14541	106	002	033		CALL	WRITE\$	
14544	036	030			LD	030	PUT OUT THE TYPE COMPLEMENTED
14546	106	002	033		CALL	WRITE\$	
14551	046	002			LE	2	DONT CHECK FOR 015 IN 1ST TWO PARITY BYTE
14553	304			SSFWL:	LAE		
14554	024	001			SU	1	DECREMENT FUDGE COUNTER
14556	340				LEA		
14557	307				LAM		GET CHARACTER FROM BUFFER
14560	120	174	031		JFS	SSFWH	EREG NOT NEG SO DONT CHECK FOR 015
14563	074	015			CP	015	CHECK FOR END OF STRING
14565	110	174	031		JFZ	SSFWH	NOT END OF STRING SO WRITE IT OUT
14570	106	016	033		CALL	WAIT\$	ITS A 015 SO END OF STRING
14573	007				RET		SO RETURN.
14574	330			SSFWH:	LDA		WRITE THE BUFFERED CHARACTER
14575	106	002	033		CALL	WRITE\$	
14600	306				LAL		BUMP THE MEMORY POINTER
14601	004	001			AD	1	
14603	360				LLA		
14604	305				LAH		
14605	014	000			AC	0	
14607	350				LHA		
14610	104	153	031		JMP	SSFWL	DO THE NEXT CHARACTER

14613	310			PARGEN:	LBA		SAVE THE BYTE
14614	253				XRD		
14615	330				LDA		
14616	301				LAB		
14617	254				XRE		
14620	012				SRC		
14621	340				LEA		
14622	106	353	036		CALL	INCHL	
14625	302				LAC		DECREMENT THE BUFFER LENGTH COUNT
14626	024	001			SU	1	
14630	320				LCA		
14631	013				RFZ		DO NEXT BYTE IF NOT ZERO
14632	306			PARSTO:	LAL		CALCULATE NUMBER OF SHIFT MOD 8
14633	066	333	056 032		HL	HLSAV+1	
14637	227				SUM		
14640	044	007			ND	7	
14642	320				LCA		
14643	302			PSLOOP:	LAC		SHIFT CIRCULATING PARITY BACK THAT MANY
14644	024	001			SU	1	
14646	320				LCA		
14647	160	260	031		JTS	PSTORE	
14652	304				LAE		
14653	002				SLC		
14654	340				LEA		
14655	104	243	031		JMP	PSLOOP	
14660	106	316	032	PSTORE:	CALL	RESHL	STORE THE CIRC. PARITY
14663	106	364	036		CALL	DECHL	
14666	374				LME		
14667	106	364	036		CALL	DECHL	STORE THE XOR PARITY
14672	373				LMD		
14673	066	254	056 032		HL	PKTADR	GET THE PACKET PARAMETERS AGAIN
14677	347				LEM		
14700	066	255	056 032		HL	PKTADR+1	
14704	337				LDM		
14705	106	152	032		CALL	GETPKT	
14710	302				LAC		INIT THE BUFFER LENGTH
14711	004	002			AD	2	COMPENSATE FOR THE TWO PARITY ACCUMS
14713	340				LEA		PUT LENGTH IN E-REGISTER
14714	106	364	036		CALL	DECHL	BACK UP BUFFER POINTER TO PARITY ACCUMS
14717	106	364	036		CALL	DECHL	
14722	250				XRA		RETURN WITH ZERO CONDITION TRUE
14723	007				RET		
							BACK UP AND DECREMENT THE RE-TRY COUNT
14724	106	041	033	DEORTC:	CALL	BKSP\$	BACK UP ONE RECORD
14727	106	016	033	DCORTC:	CALL	WAIT\$	WAIT FOR IT
14732	066	365	056 031		HL	RTC	DECREMENT THE RE-TRY COUNT
14736	307				LAM		
14737	024	001			SU	1	
14741	370				LMA		
14742	066	254	056 032		HL	PKTADR	RESTORE THE PACKET ADDRESS
14746	347				LEM		
14747	066	255	056 032		HL	PKTADR+1	

14753 337  
14754 007

LDM  
RET

. INITIATE THE RE-TRY COUNT

14755 066 365 056 031  
14761 006 003  
14763 370  
14764 007

RTCI: HL RTC SET THE RE-TRY COUNT TO THREE  
LA 3  
LMA  
RET

14765 000

RTC: DC 0 RE-TRY COUNT STORAGE

. POSITION TO THE END OF THE FILE

14766 106 152 032  
14771 106 310 033  
14774 007

PEOFX: CALL GETPKT  
CALL PEF\$  
RET

. POSITION TO THE BEGINNING OF THE FILE (AFTER FILE NUMBER RE

14775 106 152 032  
15000 106 056 033  
15003 007

PBOFX: CALL GETPKT  
CALL PBF\$  
RET

. BACKSPACE ONE RECORD

15004 106 152 032  
15007 106 041 033  
15012 106 016 033  
15015 007

BSPX: CALL GETPKT  
CALL BKSP\$  
CALL WAIT\$  
RET

. CHANGE PHYSICAL DEVICE NUMBER

15016 320  
15017 353  
15020 364  
15021 046 257 036 032  
15025 104 037 032

OPDNX: LCA SAVE THE PFN  
LHD  
LLE  
DE LFT INIT THE LFT INDEX  
JMP CLFT THE REST IF LIKE OPFN\$

. CHANGE PHYSICAL FILE NUMBER

15030 320  
15031 353  
15032 364  
15033 046 260 036 032  
15037 307  
15040 260  
15041 160 247 032  
15044 074 010  
15046 120 247 032  
15051 002  
15052 204  
15053 360  
15054 303

OPFNX: LCA SAVE THE PFN  
LHD  
LLE  
DE LFT+1 INIT THE LFT INDEX  
CLFT: LAM GET THE LOGICAL FILE NUMBER  
ORA  
JTS GDFNER CHECK IT'S RANGE  
CP 8  
JFS GDFNER  
SLC INDEX INTO THE LFT  
ADE  
LLA  
LAD



15055 014 000  
 15057 350  
 15060 372  
 15061 007

AC 0  
 LHA  
 LMC CHANGE THE PFN  
 RET

. END ACTION RETURN POINTS

15062 106 041 033  
 15065 106 016 033  
 15070 064 001  
 15072 012  
 15073 007  
 15074 106 016 033  
 15077 064 001  
 15101 007  
 15102 006 004

FEACT: CALL BKSP\$ BACK UP TO THE END OF FILE  
 CALL WAIT\$ WAIT FOR IT  
 OR 1 SET THE CARRY TOGGLE  
 SRC  
 RET  
 TEACT: CALL WAIT\$ WAIT FOR RECORD TO FINISH  
 OR 1 TYPE ERROR RETURNS NON-ZERO  
 RET  
 PEACT: LA 4 INTERNAL ERROR D IF PARITY ERROR

. INTERNAL ERROR HANDLER

15104 066 150 056 032  
 15110 004 100  
 15112 370  
 15113 066 125 056 032  
 15117 106 151 036  
 15122 104 064 000  
 15125 011 000 013 013  
 15150 040 015

ERRX: HL ERRS  
 AD 'A'-1  
 LMA  
 HL ERRMSG  
 CALL DSPLY\$  
 JMP BOOT\$  
 ERRMSG: DC 011,0,013,11,'INTERNAL ERROR  
 ERRS: DC ' ',015

- . GET THE DEVICE NUMBER IN THE B REGISTER
- . THE PHYSICAL FILE NUMBER IF 'PFN'
- . THE LENGTH IN THE C REGISTER
- . THE BUFFER STARTING ADDRESS IN HL

15152 066 254 056 032  
 15156 374  
 15157 066 255 056 032  
 15163 373  
 15164 353  
 15165 364  
 15166 307  
 15167 260  
 15170 160 247 032  
 15173 074 010  
 15175 120 247 032  
 15200 002  
 15201 004 257  
 15203 360  
 15204 006 032  
 15206 014 000  
 15210 350  
 15211 317  
 15212 106 353 036  
 15215 327  
 15216 066 256 056 032

GETPKT: HL PKTADR SAVE THE PACKET ADDRESS  
 LME  
 HL PKTADR+1  
 LMD  
 LHD GET THE LOGICAL FILE NUMBER  
 LLE  
 LAM  
 ORA CATCH LOGICAL FILE NUMBER OUT OF RANGE  
 JTS GDFNER  
 CP 8  
 JFS GDFNER  
 SLC INDEX INTO THE LOGICAL FILE TABLE  
 AD LFT  
 LLA  
 LA LFT>8  
 AC 0  
 LHA  
 LBM GET THE DEVICE NUMBER IN THE B REGISTER  
 CALL INCHL  
 LCM GET THE PHYSICAL FILE NR IN THE C REG  
 HL PFN SAVE IT IN CORE

```

15222 372          LMC
15223 353          LHD          GET THE BUFFER STARTING ADDRESS
15224 364          LLE
15225 106 353 036  CALL INCHL
15230 347          LEM
15231 106 353 036  CALL INCHL
15234 337          LDM
15235 106 353 036  CALL INCHL
15240 327          LCM          GET THE LENGTH
15241 353          LHD          PUT THE BSA IN HL
15242 364          LLE
15243 106 334 032  CALL ADR$          SELECT THE PROPER PHYSICAL DEVICE
15246 007          RET
15247 006 010      GDFNER: LA 8          LOGICAL FILE NUMBER OUT OF RANGE NETS
15251 104 047 030 JMP ERR$          YOU AN INTERNAL ERROR NUMBER EIGHT

15254 000 000      PKTADR: DA 0          CURRENT PACKET ADDRESS STORAGE
15256 000          PFN:   DC 0          CURRENT PHYSICAL FILE NUMBER STORAGE

. OPERATING SYSTEM LOGICAL FILE TABLE
.
15257 000 000      LFT:   DC 0.0          LF0 IS A NULL DEVICE
15261 001 000      DC 1.0          LF1 FOR DECK 1
15263 002 000      DC 2.0          LF2 FOR DECK 2
15265 001 001      DC 1.1          LF3 IS CTOS CATALOG
15267 002 000      DC 2.0          LF4 IS CTOS DATA SOURCE FILE
15271 002 001      DC 2.1          LF5 IS CTOS DATA OBJECT FILE
15273 000 000      DC 0.0
15275 001 040      DC 1.32         LF7 IS ASM OBJECT SCRATCH FILE

. UTILITY ROUTINES
.
15277 305          SAVHL: LAH
15300 316          LBL
15301 066 332 056 032 HL HLSAV
15305 370          LMA
15306 066 333 056 032 HL HLSAV+1
15312 371          LMB
15313 361          LLB
15314 350          LHA
15315 007          RET

15316 066 332 056 032 RESHL: HL HLSAV
15322 307          LAM
15323 066 333 056 032 HL HLSAV+1
15327 367          LLM
15330 350          LHA
15331 007          RET

15332 000 000      HLSAV: DA 0

. CASSETTE MECHANISM DRIVER
.
15334 006 360      ADR$:  LA 0360          ADDRESS THE CASSETTE MECHANISM

```

```

15336 121          EX  ADR
15337 301          LAB
15340 074 001      CP      1
15342 150 357 032 JTZ  DEK1AD
15345 074 002      CP      2
15347 150 361 032 JTZ  DEK2AD
15352 006 001      LA      1
15354 104 047 030 JMP  ERR$
15357 155          DEK1AD: EX  DECK1
15360 007          RET
15361 157          DEK2AD: EX  DECK2
15362 007          RET

```

. READ A CHARACTER INTO THE A REGISTER

```

15363          READ$:
15363 123          DEKRED: EX  STATUS
15364 101          IN
15365 044 024      ND      024
15367 150 363 032 JTZ  DEKRED
15372 002          SLC
15373 002          SLC
15374 002          SLC
15375 002          SLC
15376 043          RTC
15377 125          EX  DATA
15400 101          IN
15401 007          RET

```

. WRITE A CHARACTER FROM THE D REGISTER

```

15402          WRITE$:
15402 123          DEKWRT: EX  STATUS
15403 101          IN
15404 044 011      ND      011
15406 150 002 033 JTZ  DEKWRT
15411 012          SRC
15412 043          RTC
15413 303          LAD
15414 127          EX  WRITE
15415 007          RET

```

```

15416          WAIT$:
15416 026 001      DEKWAT: LC  1
15420 123          WAIT:  EX  STATUS
15421 101          IN
15422 242          NDC
15423 150 026 033 JTZ  WAIT
15426 007          RET

```

. FIRE UP BLOCK READ

```

15427          RBK$:
15427 106 016 033 DEKRBK: CALL DEKWAT

```

```

15432 161          EX   RBK          FIRE UP THE READ BLOCK
15433 007          RET

. FIRE UP BLOCK WRITE
.
15434          WBK$:
15434 106 016 033 DEKW BK: CALL  DEKWAT      WAIT FOR THE DECK TO BE READY
15437 163          EX   WBK          FIRE UP BLOCK WRITE
15440 007          RET

. BACKSPACE ONE RECORD
.
15441          BKSP$:
15441 106 016 033 DEKBSP: CALL  DEKWAT
15444 167          EX   BSP
15445 007          RET

. REWIND THE TAPE
.
15446          REWIND$:
15446 106 016 033 DEKREW: CALL  DEKWAT
15451 175          EX   REWIND
15452 106 016 033 CALL  DEKWAT
15455 007          RET

. POSITION TO THE BEGINNING OF THE FILE
.
15456          PBF$:
15456 106 016 033 DEKPBF: CALL  DEKWAT      WAIT FOR ANY PREVIOUS OPERATIONS
15461 066 256 056 032 HL   PFN          GET THE DESIRED FILE NUMBER
15465 317          LBM
15466 173          EX   SB          START SEARCHING BACKWARDS
15467 104 131 033 JMP   BWAIT

.
15472 106 333 033 BBACK: CALL  DEKSTP      STOP THE TAPE
15475 106 355 033 CALL  DEKFNS      SEARCH FOR A FILE MARKER
15500 104 111 033 JMP   FSKIP
15503 106 355 031 FNEXT: CALL  RTCI      INITIALIZE THE RE-TRY COUNT
15506 106 375 033 CALL  DEKFNN      SEARCH FOR NEXT FILE MARKER
15511 303          FSKIP: LAD          SEE IF WE ARE THERE YET
15512 271          CPB
15513 160 103 033 JTS   FNEXT      STILL FURTHER TO GO
15516 150 274 033 JTZ   DEKTHE      WE ARE THERE
15521 106 333 033 CALL  DEKSTP      ELSE STOP THE TAPE
15524 006 007          LA   7          ERROR EXIT SEVEN
15526 104 047 030 JMP   ERR$
15531 106 355 031 BWAIT: CALL  RTCI      INITIALIZE THE RE-TRY COUNT
15534 026 006          LC   6          WAIT FOR READ READY OR LEADER
15536 106 020 033 CALL  WAIT
15541 044 002          ND   2
15543 110 256 033 JFZ   BSTOP      CATCH LEADER
15546 365          BREAD: LLH          PUSH THE CHARACTER ONTO THE STACK
15547 354          LHE
15550 343          LED

```

15551	330	LDA		
15552	106 363 032	CALL	DEKRED	GET THE NEXT RECORD CHARACTER
15555	100 146 033	JFC	BREAD	
15560	304	LAE		GET THE SECOND RECORD CHARACTER
15561	054 377	XR	0377	UN-COMPLEMENT IT
15563	273	CPD		SEE IF IT MATCHES THE FIRST
15564	110 256 033	JFZ	BSTOP	
15567	074 303	CP	0303	IGNORE NUMERIC RECORDS
15571	150 131 033	JTZ	BWAIT	
15574	074 347	CP	0347	IGNORE SYMBOLIC RECORDS
15576	150 131 033	JTZ	BWAIT	
15601	074 201	CP	0201	ELSE IT MUST BE A FILE MARKER
15603	110 256 033	JFZ	BSTOP	
15606	306	LAL		GET THE FILE NUMBER COMPLEMENTED
15607	054 377	XR	0377	
15611	275	CPH		IT MUST MATCH THE FILE NUMBER
15612	110 256 033	JFZ	BSTOP	
15615	046 000	LE	0	FLIP OVER THE FILE NUMBER
15617	036 010	LD	8	
15621	305	FLIP:	LAH	
15622	012		SRC	
15623	350		LHA	
15624	304		LAE	
15625	210		ACA	
15626	340		LEA	
15627	303		LAD	
15630	024 001		SU	1
15632	330		LDA	
15633	110 221 033		JFZ	FLIP
15636	304		LAE	
15637	271		CPB	COMPARE IT TO THE DESIRED FILE NUMBER
15640	160 072 033		JTS	BBACK
15643	110 131 033		JFZ	BWAIT
15646	106 333 033		CALL	DEKSTP
15651	161		EX	RBK
15652	106 016 033		CALL	DEKWAT
15655	007		RET	WAIT FOR IT
15656	106 333 033	BSTOP:	CALL	DEKSTP
15661	161		EX	RBK
15662	106 341 033		CALL	DCKRTC
15665	160 070 034		JTS	DEKBAD
15670	173		EX	SB
15671	104 146 033		JMP	BREAD
15674	106 333 033	DEKTHE:	CALL	DEKSTP
15677	167		EX	BSP
15700	106 016 033		CALL	DEKWAT
15703	161		EX	RBK
15704	106 016 033		CALL	DEKWAT
15707	007		RET	AND QUIT
				AND QUIT
				POSITION TO THE END OF THE FILE
15710		PEF\$:		
15710	106 355 033	DEKPEF:	CALL	DEKFNS
				SEARCH FOR THE NEXT FILE MARKER

15713	106	333	033	CALL	DEKSTP	STOP THE TAPE
15716	167			EX	BSP	POSITION IT TO AFTER THE LAST RECORD
15717	106	016	033	CALL	DEKWAT	IN A FORWARD DIRECTION
15722	167			EX	BSP	
15723	106	016	033	CALL	DEKWAT	
15726	161			EX	RBK	
15727	106	016	033	CALL	DEKWAT	
15732	007			RET		

. STOP THE TAPE AND RE-SELECT THE PROPER DECK

15733	177			DEKSTP:	EX	TSTOP	STOP THE TAPE
15734	106	016	033	CALL	DEKWAT	WAIT FOR IT TO STOP	
15737	007			RET			

. BACK UP THE TAPE AND DECREMENT THE RE-TRY COUNT

15740	167			DEKRTC:	EX	BSP
15741	106	016	033	DCKRTC:	CALL	DEKWAT
15744	066	365	056	HL	RTC	
15750	307			LAM		
15751	024	001		SU	1	
15753	370			LMA		
15754	007			RET		

. SEARCH FORWARD FOR A FILE MARKER

15755	106	355	031	DEKFNS:	CALL	RTCI	INITIATE THE RE-TRY COUNT
15760	106	016	033	DEKFNA:	CALL	DEKWAT	WAIT FOR THE DECK TO BE READY
15763	171			EX	SF	START FORWARD MOTION	
15764	104	375	033	JMP	DEKFNN	INSPECT THE NEXT RECORD	
15767	106	363	032	DEKFNW:	CALL	DEKRED	WAIT FOR BLOCK TO BE OVER
15772	100	367	033	JFC	DEKFNW		
15775	026	004		DEKFNN:	LC	4	WAIT FOR DATA
15777	106	020	033	CALL	WAIT		
16002	106	363	032	CALL	DEKRED	GET THE RECORD TYPE	
16005	330			LDA		SAVE THE CHARACTER	
16006	106	363	032	CALL	DEKRED	GET THE RECORD TYPE COMPLEMENTED	
16011	054	377		XR	0377	UN-COMPLEMENT IT	
16013	273			OPD		THEY MUST MATCH	
16014	110	057	034	JFZ	DEKFNE		
16017	074	303		CP	0303	IGNORE NON-FILE MARKERS	
16021	150	367	033	JTZ	DEKFNW		
16024	074	347		CP	0347		
16026	150	367	033	JTZ	DEKFNW		
16031	074	201		CP	0201	ELSE IT MUST BE A FILE MARKER	
16033	110	057	034	JFZ	DEKFNE		
16036	106	363	032	CALL	DEKRED	GET THE FILE NUMBER	
16041	330			LDA		SAVE IT	
16042	106	363	032	CALL	DEKRED	GET THE FILE NUMBER COMPLEMENTED	
16045	054	377		XR	0377	UN-COMPLEMENT IT	
16047	273			OPD		THEY MUST MATCH	
16050	110	057	034	JFZ	DEKFNE		
16053	106	363	032	CALL	DEKRED	THIS MUST BE THE END OF THE RECORD	

16056 043  
 16057 106 333 033  
 16062 106 340 033  
 16065 120 360 033  
 16070 006 002  
 16072 104 047 030

RTC  
 DEKFNE: CALL DEKSTP  
 CALL DEKRTC  
 JFS DEKFNA  
 DEKBAD: LA 2  
 JMP ERR\$

STOP THE TAPE  
 BACK UP AND COUNT TRY  
 TRY AGAIN IF NOT ALREADY TOO MANY  
 ELSE UNLOADABLE RECORD

## SPECIAL TAPE ROUTINES

16075 106 152 032  
 16100 301  
 16101 074 002  
 16103 006 001  
 16105 110 047 030  
 16110 106 046 033  
 16113 007

TRWX: CALL GETPKT  
 LAB  
 CP 2  
 LA 1  
 JFZ ERR\$  
 CALL REWIND\$  
 RET

REWIND THE TAPE

ONLY REWIND THE FRONT DECK

16114 106 152 032  
 16117 106 355 033  
 16122 106 333 033  
 16125 323  
 16126 007

TFNRX: CALL GETPKT  
 TRWFNR: CALL DEKFNS  
 CALL DEKSTP  
 LCD  
 RET

READ A FILE NUMBER  
 FIND A FILE MARKER  
 STOP THE TAPE AFTER IT  
 PUT THE FILE NUMBER IN THE C REGISTER

16127 106 152 032  
 16132 106 016 033  
 16135 163  
 16136 036 201  
 16140 106 002 033  
 16143 036 176  
 16145 106 002 033  
 16150 066 256 056 032  
 16154 337  
 16155 106 002 033  
 16160 303  
 16161 054 377  
 16163 330  
 16164 106 002 033  
 16167 106 016 033  
 16172 007

TFNRX: CALL GETPKT  
 CALL DEKWAT  
 EX WBK  
 LD 0201  
 CALL DEKWRT  
 LD 0176  
 CALL DEKWRT  
 HL PFN  
 LDM  
 CALL DEKWRT  
 LAD  
 XR 0377  
 LDA  
 CALL DEKWRT  
 CALL DEKWAT  
 RET

WRITE A FILE NUMBER

FIRE UP A WRITE  
 WRITE OUT THE FILE MARKER

WRITE OUT ITS COMPLEMENT

WRITE OUT THE FILE NUMBER

WRITE OUT ITS COMPLEMENT

TERMINATE THE WRITE OPERATION

DONE

16200

SET 016200

. ITSEY BITSEY DEBUG

16200	066	273	056	035	DEBUG\$*	HL	CURADR
16204	347					LEM	
16205	066	274				LL	CURADR+1
16207	337					LDM	
16210	026	005				LC	5
16212	066	306				LL	DSPADR+4
16214	106	226	035			CALL	CONBIN
16217	066	273				LL	CURADR
16221	347					LEM	
16222	066	274				LL	CURADR+1
16224	357					LHM	
16225	364					LLE	
16226	347					LEM	
16227	036	000				LD	0
16231	026	003				LC	3
16233	066	313	056	035		HL	DSPDAT+2
16237	106	226	035			CALL	CONBIN
16242	066	275				LL	DISP
16244	106	151	036			CALL	DSPLY\$
16247	066	315				LL	INBUF
16251	026	010				LC	8
16253	106	000	036			CALL	KEYIN\$
16256	066	315				LL	INBUF
16260	106	155	035			CALL	CONOCT
16263	074	015				CP	015
16265	152	355	034			CTZ	NEWADR
16270	074	111				CP	'I'
16272	152	364	034			CTZ	INCADR
16275	074	104				CP	'D'
16277	152	002	035			CTZ	DECADR
16302	074	115				CP	'M'
16304	152	020	035			CTZ	MODIFY
16307	074	056				CP	' '
16311	152	031	035			CTZ	ENTER
16314	074	114				CP	'L'
16316	152	121	035			CTZ	LSAVE
16321	074	110				CP	'H'
16323	152	144	035			CTZ	HSAVE
16326	074	107				CP	'G'
16330	152	111	035			CTZ	GET
16333	074	106				CP	'F'
16335	152	101	035			CTZ	FRONT
16340	074	117				CP	'O'
16342	150	042	012			JTZ	OS\$
16345	074	103				CP	'C'
16347	150	042	035			JTZ	GOTO
16352	104	200	034			JMP	DEBUG\$
16355	066	273			NEWADR:	LL	CURADR
16357	374					LME	



16360	066	274		LL	CURADR+1
16362	373			LMD	
16363	007			RET	
16364	066	273	INCADR:	LL	CURADR
16366	307			LAM	
16367	004	001		AD	1
16371	370			LMA	
16372	066	274		LL	CURADR+1
16374	307			LAM	
16375	014	000		AC	0
16377	370			LMA	
16400	250			XRA	
16401	007			RET	
16402	066	273	DECADR:	LL	CURADR
16404	307			LAM	
16405	024	001		SU	1
16407	370			LMA	
16410	066	274		LL	CURADR+1
16412	307			LAM	
16413	034	000		SB	0
16415	370			LMA	
16416	250			XRA	
16417	007			RET	
16420	066	273	MODIFY:	LL	CURADR
16422	337			LDM	
16423	066	274		LL	CURADR+1
16425	357			LHM	
16426	363			LLD	
16427	374			LME	
16430	007			RET	
16431	106	020 035	ENTER:	CALL	MODIFY
16434	056	035		LH	CURADR>8
16436	106	364 034		CALL	INCADR
16441	007			RET	
16442	066	051	GOTO:	LL	BRANCH+1
16444	374			LME	
16445	066	052		LL	BRANCH+2
16447	373			LMD	
16450	106	000 000	BRANCH:	CALL	0
16453	300		LSAVI:	LAA	
16454	300		HSAVI:	LAA	
16455	066	370 056 035		HL	ASAVE
16461	370			LMA	
16462	066	371		LL	BSAVE
16464	371			LMB	
16465	066	372		LL	CSAVE
16467	372			LMC	
16470	066	373		LL	DSAVE
16472	373			LMD	
16473	066	374		LL	ESAVE
16475	374			LME	
16476	104	200 034		JMP	DEBUG\$
16501	016	001	FRONT:	LB	1

16503	106	112	000		CALL	LOAD2\$
16506	104	115	035		JMP	GETLOD
16511	314			GET:	LBE	
16512	106	100	000		CALL	LOAD\$
16515	053			GETLOD:	RTZ	
16516	151				EX	BEEP
16517	250				XRA	
16520	007				RET	
16521	066	053	056	035	LSAVE:	HL LSAVI
16525	016	306			LB	0306
16527	304			HLSAVM:	LAE	
16530	044	007			ND	7
16532	002				SLC	
16533	002				SLC	
16534	002				SLC	
16535	074	070			OP	070
16537	023				RFS	
16540	261				ORB	
16541	370				LMA	
16542	250				XRA	
16543	007				RET	
16544	066	054	056	035	HSAVE:	HL HSAVI
16550	016	305			LB	0305
16552	104	127	035		JMP	HLSAVM

## : CONVERT OCTAL TO BINARY

16555	036	000		CONOCT:	LD	0
16557	343				LED	
16560	317			CONLOP:	LBM	
16561	106	353	036		CALL	INCHL
16564	301				LAB	
16565	074	070			OP	'8'
16567	023				RFS	
16570	074	060			OP	'0'
16572	063				RTS	
16573	044	007			ND	7
16575	320				LCA	
16576	303				LAD	
16577	044	037			ND	037
16601	002				SLC	
16602	002				SLC	
16603	002				SLC	
16604	330				LDA	
16605	304				LAE	
16606	002				SLC	
16607	002				SLC	
16610	002				SLC	
16611	340				LEA	
16612	044	007			ND	7
16614	263				ORD	

16615	330		LDA	
16616	304		LAE	
16617	044	370	ND	0370
16621	262		ORC	
16622	340		LEA	
16623	104	160 035	JMP	CONLOP

. CONVERT BINARY TO OCTAL (RIGHT TO LEFT)

16626	304		CONBIN: LAE	
16627	044	007	ND	7
16631	004	060	AD	'0'
16633	370		LMA	
16634	106	364 036	CALL	DECHL
16637	304		LAE	
16640	012		SRC	
16641	012		SRC	
16642	012		SRC	
16643	044	037	ND	037
16645	340		LEA	
16646	303		LAD	
16647	012		SRC	
16650	012		SRC	
16651	012		SRC	
16652	330		LDA	
16653	044	340	ND	0340
16655	264		ORE	
16656	340		LEA	
16657	303		LAD	
16660	044	037	ND	037
16662	330		LDA	
16663	302		LAC	
16664	024	001	SU	1
16666	320		LCA	
16667	110	226 035	JFZ	CONBIN
16672	007		RET	

. STORAGE

16673	000	000	CURADR: DA	0
16675	011	000 013 013	DISP: DC	011.0.013.11.021
16702	040	040 040 040	DSPADR: DC	'
16711	040	040 040 015	DSPDAT: DC	' .015
16715	040	040 040 040	INBUF: DC	'
16770			SET	016770
16770	000		ASAVE: DC	0
16771	000		BSAVE: DC	0
16772	000		CSAVE: DC	0
16773	000		DSAVE: DC	0
16774	000		ESAVE: DC	0
16775	001	002	DC	1.2

DONE

17000

SET 017000

## . KEYBOARD ENTRY ROUTINE

. ACCEPTS A STRING OF CHARACTERS FROM THE KEYBOARD AND PUTS  
 . THEM IN MEMORY STARTING WITH THE ADDRESS GIVEN IN THE H  
 . AND L REGISTERS AND AT A DISPLAY POSITION DESCRIBED BY THE  
 . D (HORZ) AND E (VERT) REGISTERS. THE MAXIMUM NUMBER OF  
 . CHARACTERS ACCEPTED IS TAKEN FROM THE C REGISTER UPON ENTRY  
 . OVERFLOW OFF THE END OF A DISPLAY LINE IS NOT PERMITTED  
 . AND IF EITHER THE MAXIMUM COUNT OR DISPLAY BOUNDARY IS  
 . EXCEEDED, SUCCESSIVE CHARACTERS WILL GO IN THE LAST  
 . POSITION OVER AND OVER. AN 015 WILL TERMINATE INPUT REQUEST  
 . THE CURSOR IS TURNED ON UPON ENTRY AND OFF UPON EXIT.

17000	006	341	KEYIN\$#	LA	0341	ADDRESS THE KEYBOARD
17002	121			EX	ADR	
17003	312			LBC		LOAD THE MAX COUNT INTO THE CURRENT COUNT
17004	006	020		LA	020	TURN ON THE CURSOR
17006	131			EX	COM1	
17007	106	326 036	KILOOP:	CALL	CHAIT	MAKE SURE THE DISPLAY IS READY
17012	123		KWLOOP:	EX	STATUS	GET A CHARACTER FROM THE KEYBOARD
17013	101			IN		
17014	044	002		ND	2	
17016	150	012 036		JTZ	KWLOOP	
17021	125			EX	DATA	
17022	101			IN		
17023	074	010		CP	010	CATCH BACKSPACE
17025	150	105 036		JTZ	KBSP	
17030	074	030		CP	030	CATCH DELETE
17032	150	113 036		JTZ	KDEL	
17035	074	100		CP	0100	REVERSE THE SHIFT KEY FUNCTION
17037	160	044 036		JTS	KSTORE	
17042	054	040		XR	040	
17044	370		KSTORE:	LMA		STORE THE CHARACTER
17045	074	015		CP	015	CATCH THE ENTER KEY
17047	150	102 036		JTZ	KEND	
17052	127			EX	WRITE	ELSE DISPLAY THE CHARACTER
17053	303			LAD		CATCH CURSOR AT SCREEN BOUNDARY
17054	074	117		CP	79	
17056	120	007 036		JFS	KILOOP	
17061	301			LAB		DECREMENT THE CHARACTER COUNT
17062	024	001		SU	1	
17064	160	007 036		JTS	KILOOP	ALREADY ABOVE THE MAXIMUM
17067	310			LBA		
17070	303			LAD		BUMP THE CURSOR POSITION FOR REAL
17071	004	001		AD	1	
17073	330			LDA		
17074	106	353 036		CALL	INCHL	BUMP THE MEMORY LOACTION
17077	104	007 036		JMP	KILOOP	DO THE NEXT CHARACTER
17102	250		KEND:	XRA		TURN OFF THE CURSOR
17103	131			EX	COM1	
17104	007			RET		

17105	106	124	036	KBSP:	CALL	KBSPR	BACKSPACE ONE CHARACTER
17110	104	007	036		JMP	KILOOP	
17113	106	124	036	KDEL:	CALL	KBSPR	BACKSPACE TO THE BEGINNING OF THE ENTRY
17116	110	113	036		JFZ	KDEL	
17121	104	007	036		JMP	KILOOP	
.							
17124	301			KBSPR:	LAB		INCREMENT THE CHARACTER COUNTER
17125	272				CPC		UNLESS AT THE BEGINNING OF THE ENTRY
17126	053				RTZ		
17127	004	001			AD	1	
17131	310				LBA		
17132	303				LAD		DECREMENT THE SCREEN POSITION
17133	024	001			SU	1	
17135	330				LDA		
17136	106	364	036		CALL	DECHL	DECREMENT THE MEMORY POINTER
17141	106	326	036		CALL	CHAIT	MAKE SURE THE DISPLAY IS READY
17144	006	040			LA	040	ERASE THE CHARACTER
17146	127				EX	WRITE	
17147	260				ORA		RETURN WITH ZERO CONDITION FALSE
17150	007				RET		

. CRT DISPLAY ROUTINE

. DISPLAYS A STRING OF CHARACTERS WHICH ARE IN MEMORY STARTIN  
 . WITH THE ADDRESS GIVEN IN THE H AND L REGISTERS AND AT THE  
 . POSITION DESCRIBED BY THE D (HORZ) AND E (VERT) REGISTERS.  
 . OVERFLOW OFF THE END OF A LINE IS NOT PERMITTED.  
 . SPECIAL CONTROL CHARACTERS TERMINATE THE LINE AND ALLOW  
 . MOVEMENT OF THE CURSOR, ERASURE OF THE SCREEN OR LINE,  
 . AND ROLL-UP OF THE ENTIRE SCREEN.

. ENTRY VALUES:           D - HORIZONTAL CURSOR POSITION (0 TO 79)  
 .                           E - VERTICAL CURSOR POSITION (0 TO 11)  
 .                           HL - FIRST CHARACTER LOCATION IN STRING  
 . EXIT VALUES:           DE - CURSOR POSITION AFTER LAST CHAR  
 .                           HL - MEMORY LOCATION AFTER TERM CHAR  
 . CONTROL CHARACTERS: 003 - END OF THE STRING  
 .                           011 - A NEW HORIZONTAL POSITION FOLLOWS  
 .                           013 - A NEW VERTICAL POSITION FOLLOWS  
 .                           015 - END OF LINE (DOES CR/LF)  
 .                           021 - ERASE TO THE END OF THE FRAME  
 .                           022 - ERASE TO THE END OF THE LINE  
 .                           023 - ROLL UP THE SCREEN ONE LINE

17151	006	341	DSPLY\$#	LA	0341	ADDRESS THE DISPLAY
17153	121			EX	ADR	
17154	250			XRA		TURN OFF THE CURSOR
17155	131		DOCOM:	EX	COM1	DO THE CONTROL COMMAND
17156	106	326	DLOOP:	CALL	CHAIT	MAKE SURE THE DISPLAY IS READY
17161	317			LBM		GET A CHARACTER FROM THE STRING
17162	106	353		CALL	INCHL	BUMP THE STRING POINTER
17165	301			LAB		CHECK FOR CONTROL CHARACTERS
17166	044	177		ND	0177	STRIP ANY PARITY

17170	074	003	CP	3	
17172	150	265	JTZ	ENDOS	END OF STRING
17175	074	011	CP	011	
17177	150	300	JTZ	PHORZ	POSITION HORIZONTALLY
17202	074	013	CP	013	
17204	150	274	JTZ	PVERT	POSITION VERTICALLY
17207	074	015	CP	015	
17211	150	245	JTZ	ENDOL	END OF LINE
17214	074	021	CP	021	
17216	150	307	JTZ	EEOF	ERASE TO THE END OF THE FRAME
17221	074	022	CP	022	
17223	150	314	JTZ	EEOL	ERASE TO THE END OF THE LINE
17226	074	023	CP	023	
17230	150	321	JTZ	ROLLUP	ROLL UP THE SCREEN
17233	127		EX	WRITE	PUT OUT THE CHARACTER
17234	303		LAD		BUMP THE CURSOR POSITION
17235	074	117	CP	79	UNLESS AT THE END OF THE LINE
17237	014	000	AC	0	
17241	330		LDA		
17242	104	156	JMP	DLOOP	
17245	036	000	ENDOL:	LD 0	RETURN CURSOR TO START OF NEXT LINE
17247	304		LAE		BUMP THE LINE COUNTER
17250	004	001	AD	1	
17252	340		LEA		
17253	074	014	CP	12	
17255	160	265	JTS	ENDOS	THERE IS ROOM FOR THE NEXT LINE
17260	046	013	LE	11	ELSE KEEP THE LINE COUNTER AT ELEVEN
17262	006	010	LA	010	AND ROLL THE SCREEN UP ONE LINE
17264	131		EX	COM1	
17265	106	326	ENDOS:	CALL CHAIT	MAKE SURE THE DISPLAY IS READY
17270	006	020	LA	020	TURN ON THE CURSOR
17272	131		EX	COM1	
17273	007		RET		RETURN
17274	347		PVERT:	LEM	SET THE VERTICAL POSITION
17275	104	301	JMP	NCHAR	
17300	337		PHORZ:	LDM	SET THE HORIZONTAL POSITION
17301	106	353	NCHAR:	CALL INCHL	BUMP THE STRING POINTER TO THE NEXT CHAR
17304	104	156	JMP	DLOOP	
17307	006	004	EEOF:	LA 4	
17311	104	155	JMP	DOCOM	
17314	006	002	EEOL:	LA 2	
17316	104	155	JMP	DOCOM	
17321	006	010	ROLLUP:	LA 010	
17323	104	155	JMP	DOCOM	
17326	123		CHAIT:	EX STATUS	WAIT FOR THE DISPLAY TO BE READY
17327	101		IN		
17330	012		SRC		
17331	100	326	JFC	CHAIT	
17334	303		LAD		MAKE SURE CURSOR IS IN CORRECT POSITION
17335	260		ORA		PREVENT CURSOR POSITIONS OUT OF RANGE
17336	063		RTS		

17337	074	120		CP	80	
17341	023			RFS		
17342	133			EX	COM2	
17343	304			LAE		
17344	260			ORA		
17345	063			RTS		
17346	074	014		CP	12	
17350	023			RFS		
17351	135			EX	COM3	
17352	007			RET		
17353	306		INCHL*	LAL		BUMP MEMORY POINTER UP
17354	004	001		AD	1	
17356	360			LLA		
17357	305			LAH		
17360	014	000		AC	0	
17362	350			LHA		
17363	007			RET		
17364	306		DECHL*	LAL		BUMP MEMORY POINTER DOWN
17365	024	001		SU	1	
17367	360			LLA		
17370	305			LAH		
17371	034	000		SB	0	
17373	350			LHA		
17374	007			RET		

DONE

## . LIBRARY CATALOG

17404					SET	017404		
17404	037				CATW#	DC	037	STARTING ADDRESS FOR LOADER
17405	010					DC	010	
17406	340					DC	0340	STARTING ADDRESS COMPLEMENTED
17407	367					DC	0367	
17410					CAT#	RPT	14	SPACE FOR 14 ENTRIES
17410	040	040	040	040		DC	'	
17420	040	040	040	040		DC	'	
17430	040	040	040	040		DC	'	
17440	040	040	040	040		DC	'	
17450	040	040	040	040		DC	'	
17460	040	040	040	040		DC	'	
17470	040	040	040	040		DC	'	
17500	040	040	040	040		DC	'	
17510	040	040	040	040		DC	'	
17520	040	040	040	040		DC	'	
17530	040	040	040	040		DC	'	
17540	040	040	040	040		DC	'	
17550	040	040	040	040		DC	'	
17560	040	040	040	040		DC	'	
17570	052					DC	'*'	
17571	000				ALPFN#	DC	0	AUTO-LOAD PHYSICAL FILE NUMBER
								. END OF PHYSICAL FILE 1
17572	040	040	040	040	SYMBOL#	DC		ITEM SYMBOL STORAGE
								. LOAD AND EXECUTE
17601	106	100	000		MAUTO\$#	CALL	LOAD\$	LOAD THE GIVEN FILE
17604	100	075	000		MAUTO:	JFC	RUN\$	EXECUTE IT IF GOOD LOAD
17607	104	074	000			JMP	BOOT\$	ELSE RE-LOAD THE OPERATING SYSTEM
17612	106	112	000		MAUT2\$#	CALL	LOAD2\$	LOAD DECK TWO FILE
17615	104	204	037			JMP	MAUTO	
								. SYMBOLIC FILE LOADER
17620	353				MLOAD\$#	LHD		GET PACKET ADDRESS
17621	364					LLE		
17622	046	172	036	037		DE	SYMBOL	PUT THE NAME IN THE LOOKUP ITEM
17626	026	006				LC	6	
17630	106	345	037			CALL	BLKTFR	
17633	046	010	036	037		DE	CAT	LOOK IT UP IN THE LIBRARY CATALOG
17637	106	264	037			CALL	LOOKUP	
17642	074	005				CP	5	SEE IF IT IS THE THE CATALOG
17644	013					RFZ		ZERO FLAG FALSE IF IT ISN'T
17645	306					LAL		CALCULATE THE FILE NUMBER
17646	044	370				ND	0370	
17650	024	010				SU	CAT	



17652	012		SRC		
17653	012		SRC		
17654	012		SRC		
17655	004	002	AD	2	FIRST ENTRY IS PHYSICAL FILE TWO
17657	310		LBA		
17660	106	100	CALL	LOAD\$	
17663	007		RET		

. SYMBOL LOOKUP ROUTINE

17664	353		LOOKUP*	LHD	CHECK FIRST ENTRY IN TABLE
17665	364			LLE	
17666	104	334		JMP	LSTART
17671	066	172	LOOKPU:	HL	SYMBOL
17675	327		LSLOOP:	LCM	GET THE ITEM STARTING ADDRESS
17676	106	365		CALL	INCSWP
17701	307			LAM	GET THE NEXT ITEM CHARACTER
17702	272			OPC	GET THE NEXT TABLE ADDRESS
17703	110	322		JFZ	LDIFF
17706	306			LAL	THEY DON'T MATCH
17707	044	007		ND	7
17711	074	005		OP	5
17713	053			RTZ	THE ITEM HAS BEEN FOUND IF SO
17714	106	365		CALL	INCSWP
17717	104	275		JMP	LSLOOP
17722	306		LDIFF:	LAL	GET THE NEXT ITEM ADDRESS
17723	044	370		ND	0370
17725	004	010		AD	8
17727	360			LLA	
17730	305			LAH	
17731	014	000		AC	0
17733	350			LHA	
17734	307		LSTART:	LAM	GET THE TABLE FIRST CHARACTER
17735	074	101		OP	'A'
17737	063			RTS	END OF TABLE IF IT IS NOT ALPHA
17740	335			LDH	SAVE THE TABLE ADDRESS
17741	346			LEL	
17742	104	271		JMP	LOOKPU
					AND TRY NEXT TABLE ENTRY

. BLOCK TRANSFER FROM HL TO DE C CHARACTERS

17745	317		BLKTFR#	LBM	GET A SOURCE CHARACTER
17746	106	365		CALL	INCSWP
17751	371			LMB	GET NEXT DESTINATION LOCATION
17752	106	365		CALL	INCSWP
17755	302			LAC	PUT IT IN A DESTINATION LOCATION
17756	024	001		SU	1
17760	320			LCA	GET NEXT SOURCE ADDRESS
17761	110	345		JFZ	BLKTFR
17764	007			RET	DECREMENT THE COUNT
					DO NEXT CHAR IF NOT ZERO

. INCREMENT HL AND SWAP IT WITH DE

17765	306		INCSWP*	LAL	
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17766	004	001	AD	1
17770	364		LLE	
17771	340		LEA	
17772	305		LAH	
17773	014	000	AC	0
17775	353		LHD	
17776	330		LDA	
17777	007		RET	

DONE