

UNIVERSITY OF ILLINOIS
DIGITAL COMPUTER

LIBRARY ROUTINE H 5 - 85

TITLE	Minimization of a Function of n Variables (DOI or SADOI)	
TYPE	Closed	
NUMBER OF WORDS	89	
TEMPORARY STORAGE	$n + 1$ words at S3 $n + 1$ words at S4 $n + 1$ words at S5	
ACCURACY	Depends upon the condition of the function	
DURATION	A "minor" cycle takes $T = (.5 + 1.4n + t)$ ms where n is the number of variables and $t =$ time to compute $f(x_1, \dots, x_n)$. A "major" cycle takes $(1.8 + 4.5n + 2nt + NT)$ where N is the number of "minor cycles" in a major cycle.	
PRESET PARAMETER	S3 - SN	
	3 - 00 F 00 aF	where $a, a+1, \dots, a+n$ are the addresses of $f(x_1, \dots, x_n), x_1, x_2, \dots, x_n$ on entry and exit;
	4 - 00 F 00 bF	$b, b+1, \dots, b+x$ are the addresses of $\delta, \delta_{x_1}, \dots, \delta_{x_n}$;
	5 - 00 F 00 cF	$c, c+1, \dots, c+n$ are working spaces;
	6 - 00 F 00 α J	α is the factor by which the mesh size is decreased;
	7 - 00 F 00 ϵ J	ϵ is an end constant such that only mesh sizes <u>larger</u> than ϵ will be used;
	8 - 00 F 00 sF	s is the address of a <u>closed</u> subroutine which takes x_1 from $c+1$ (1S5), x_2 from $c+2$ (2S5), \dots, x_n from $c+n$ (nS5) and places $f(x_1, x_2, \dots, x_n)$ in R_1 ;
	9 - 00 F 00 (a+n)F	n is the number of variables
	K - 00 F 00 (b+n)F	
	S - 00 F 00 (c+n)F	
	N - 00 F 00 rF or	The preset parameter in N allows the programmer some choice in scale; if
	00 F 40 F	

it is 00 F 00 rF

δ_{x_1} will be computed as:

$$\delta_{x_1} = \frac{2^{-r} f_{x_1}}{\sum_1 |2^{-r} f_{x_1}|}$$

whereas if it is 00 F 40 F we set $r = 0$ in the above expression. In the above:

$$f_{x_1} = f(\dots x_1 - \delta \dots) - f(\dots x_1 + \delta \dots)$$

In either case if the denominator exceeds scale it is replaced by $1 - 2^{-39}$.

ENTRY

$x_1^0, x_2^0, \dots, x_n^0$ the initial values of x_1, x_2, \dots, x_n are to be loaded in 1S3, 2S3, ..., nS3. δ_0 , the initial value of δ is to be in R_1 , and the main program should contain:

p	50 rF
	50 pF
p+1	26 qF

where q is the address of this routine and r is the address to the left hand side of which control will be transferred before every decrease in mesh size. At this address the programmer may place a routine to assess or print intermediate results. This routine should return control to the left hand side of $(q+71)$. This can be done automatically since when control is transferred to r (L.H.) the appropriate return address will be in the right hand address position of R_2 and can be utilized by S5, 42 into a 26 at the end of the programmer's interlude routine. The best current values of $f(x_1, \dots, x_n), x_1, \dots, x_n$ are always in S3, 1S3, ..., nS3. If it is not desired to leave this code, set $r = q + 71$. The routine is finally left with the best values of $f(x_1, x_2, \dots, x_n), x_1, x_2, \dots, x_n$ in S3, 1S3, 2S3, ..., nS3 respectively.

DESCRIPTION OF METHOD

- 1 - f_{x_1} are computed
- 2 - δ_{x_1} are computed
- 3 - $f(x_1 + \delta_{x_1}, \dots, x_n + \delta_{x_n})$ is computed and tested against previous $f(x_1, \dots, x_n)$. If the function has decreased the old values of function and argument are replaced by the new. The process is repeated with the same δ_{x_1} . This step comprises a "minor" cycle.
- 4 - When the above process ceases to improve the function, we return to step 1 and compute new δ_{x_1} . Steps 1-4 comprise a "major" cycle.
- 5 - When the first "minor" cycle of a given "major" cycle fails to improve the function, indicating that no further improvement is to be expected using the current mesh size, control is transferred to the interlude as explained above.
- 6 - Upon returning to the code, δ , the mesh size will be replaced by $\alpha\delta$ where α is entered in S6 during read-in. The steps 1-5 are then repeated.
- 7 - Steps 1-6 will be repeated $(N+1)$ times where N is the largest integer for which $\alpha^N \delta_0 > \epsilon$. ϵ is to be entered in S7 during read-in.

NOTE

A function, poorly conditioned in having a very small gradient with respect to some argument, can deceive this routine in that this coordinate of the minimum will be very poorly found. If the programmer suspects that such a condition exists, this routine can be used for roughing and Routine H 6 - 86, a brute force approach to the minimum by varying one argument at a time, may be used for finishing. All parameters (except SK and SN), entry and contents of storage location of this code and Routine H 6 are identical, so that they may be used at the same time in the machine for the above purpose.

Routines H 3, H 4, H 5, and H 6 constitute a set sufficient (perhaps not necessary) to minimize any function of n variables. Due to economy of storage, Routine H 3 should be used for $n = 1$ or 2 , Routine H 4 for $n = 3$ or 4 , and Routine H 5 for $n \geq 5$.

lgr

DATE	July 21, 1954	rt. 3/5/59
CODED BY	J. N. Snyder	
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LOCATION	ORDER		NOTES	PAGE 1	K 5
0	00 K(H5) 40 S4 S5 1F		Store s_0 Move in link		
1	46 60L L4 L		Plant interlude address		
2	42 73L L5 3L		Plant link Set 10L for initial entry		
3	50 S8 46 10L				
4	L5 77L 40 5L	From 12, 73	Set initial addresses in transfer orders		
5	L5 F 40 F	From 8 By 4, 7	x_1^0, \dots, x_n^0 to 1S5, ..., nS5		
6	L5 5L L4 76L		Advance addresses in transfer orders		
7	40 5L L5 80L				
8	L0 5L 36 5L		Test for $i = n$		
9	50 131 50 9L				
10	26 F 40 S3	By 3, 11	$f(x_1, \dots, x_n)$ to S3		
11	L5 9L 47 10L		Set 10L for subsequent entry		
12	26 4L 00 F				
13	L5 60L 46 63L	From 10	Set switch for failure on first try		
14	L5 77L 46 19L				
15	46 23L 40 27L		Set $i = 1$ in 19 - 27		
16	L5 22L 46 20L				

LOCATION	ORDER	NOTES	PAGE 2
17	46 24L L5 79L		
18	42 22L 40 26L		
19	L5 F L4 S4	By 14, 29 From 35	
20	40 F 50 20L	By 16, 31	
21	26 S8 10 SN		Compute f_{x_1}
22	50 1S5 40 F	By 18, 34	
23	L5 F L0 S4	By 15, 30	
24	40 F 50 24L	By 17, 32	
25	26 S8 10 SN		
26	L0 F 40 F	By 18, 33 By 18, 33	
27	L5 F 40 F	By 15, 29 By 15, 29	
28	L5 27L L4 76L		
29	40 27L 46 19L		Advance i to i+1 in 19-27
30	46 23L L5 20L		
31	L4 76L 46 20L		
32	46 24L L5 26L		
33	L5 76L 40 26L		

LOCATION	ORDER	NOTES	PAGE 3
34	42 22L L5 80L		Test for $i = n$
35	L0 27L 36 19L		
36	L5 79L 46 38L		Set $i = 1$ in 38
37	19 38F 40 S5		
38	L7 F L4 S5	By 36, 40 From 42	
39	40 S5 L5 38L		
40	L4 76L 46 38L		Compute $\sum_i f_{x_i} $
41	L5 81L L0 38L		
42	36 38L 26 85L		
43	46 44L 42 45L	From 87	Set $i = 1$ in 44-45
44	L5 F 66 S5	By 43, 47	Compute δ_{x_i}
45	7J S4 40 F	By 43, 48	
46	L5 44L L4 76L		Advance i to $i+1$ in 44-45
47	46 44L L5 45L		
48	L4 76L 42 45L		
49	L5 82L L0 45L		Test for $i = n$
50	36 44L L5 77L	From 70	

LOCATION	ORDER		NOTES	PAGE 4
51	46 54L L5 79L		Set $i = 1$ in 54-55	
52	42 54L L5 22L			
53	46 55L 50 71L			
54	L5 F	By 51, 56 From 59		
55	L4 F 40 F	By 52, 56 By 53, 58	Set $x_i + \delta_{x_i}$ in 185	
56	L5 54L L4 76L 40 54L		Advance i to $i+1$ in 54-55	
57	L5 55L L4 76L			
58	46 55L L5 83L		Test for $i = n$	
59	L0 55L 36 54L			
60	50 F 50 60L	By 1	Interlude address $f(\dots x_i + \delta_{x_i} \dots)$	
61	26 88 40 85		Form $f = f$ (before)	
62	L0 83 50 53L			
63	36 F L5 78L	By 13, 69	To interlude	
64	42 65L 46 65L		Set $i = 1$ in 65	
65	L5 F 40 F	By 64, 67 From 68 By 64, 67	Replace old values by new	
66	L5 65L L4 76L		Advance i to $i+1$ in 65	

LOCATION	ORDER		NOTES	PAGE 5
67	40 65L			
	L5 84L		Test for $i = n$	
68	L0 65L			
	36 65L			
69	L5 70L		Open failure switch and repeat	
	46 63L			
70	50 4L		Minor cycle	
	22 50L			
71	50 74L		Form $\delta_{i+1} = \alpha \delta_i$ and test against ϵ	
	7J S4			
72	40 S4			
	L0 75L			
73	36 4L			
	22 F	By 2.	Link	
74	00 F			
	00 S6		= α	
75	00 F			
	00 S7		= ϵ	
76	00 1F		Address advance	
	00 1F			
77	L5 1S3			
	40 1S5			
78	00 S5		Starting constants	
	00 S3			
79	L0 1S4			
	40 1S4			
80	L5 S9			
	40 SS			
81	L7 SK			
	L4 S5		End constants	
82	7J S4			
	40 SK			
83	40 SS			
	L5 54L			

LOCATION	ORDER	NOTES	PAGE 6
84	L5 SS 40 S9		
85	L5 S5 36 87L	From 42	Test $\sum_i f_{x_i} $
86	L5 88L 40 S5		If > 1, replace by $1-2^{-39}$
87	L5 79L 26 43L		
88	7L 4095F LL 4095F		