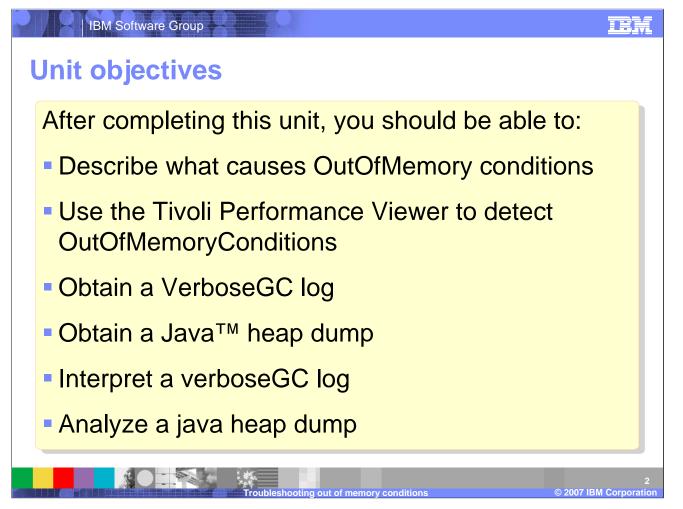


This unit covers the problem determination techniques associated with out of memory problems in WebSphere Application Server.



After completing this unit, you will be able to:

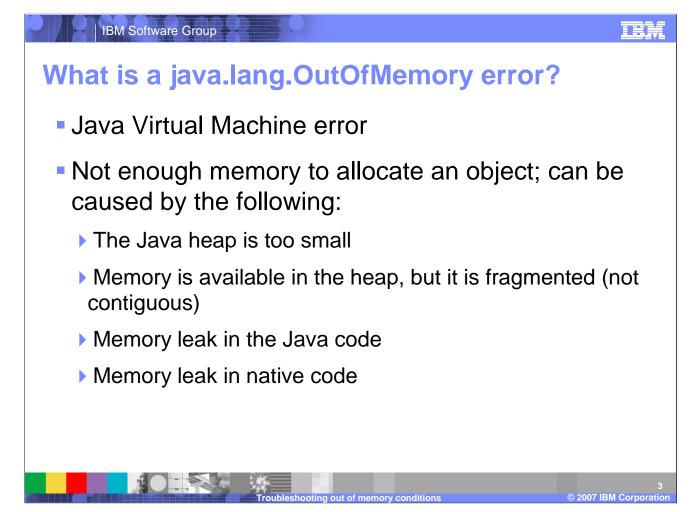
Define an out-of-memory condition

Use Tivoli Performance Viewer to detect out-of-memory conditions

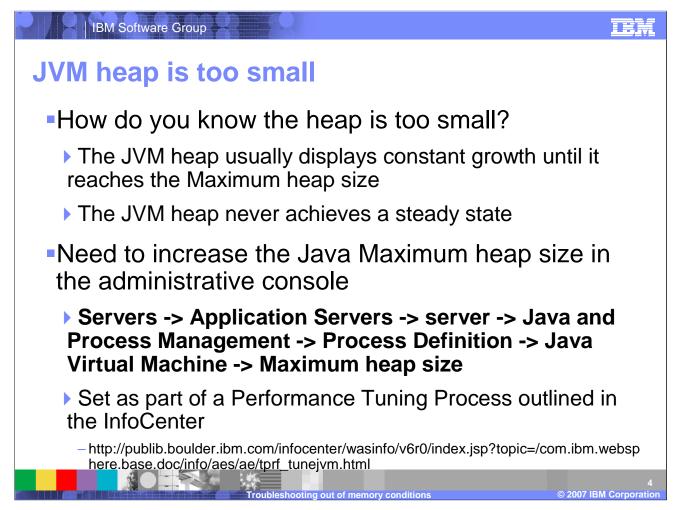
Obtain and interpret a verboseGC log

Obtain and interpret a heap dump

Discuss tools for analyzing out-of-memory problems



When the Java virtual machine (JVM) tries to allocate an object and it fails, it runs garbage collection to free up heap space from objects that are no longer being used. If the object cannot be allocated after garbage collection, the JVM throws a java.lang.OutOfMemoryError. There are four conditions that cause the OutOfMemoryError to be thrown. The first is that the JVM heap's maximum size is too small. Second, there may be enough unallocated space in the JVM heap, but it is not contiguous, which is called fragmentation. Third, the Java code could be continuously creating objects and never relinquishes them, thus causing a memory leak. The final condition is a leak in native code. In this case the JVM heap has plenty of contiguous space available, but the Operating System fails to provide memory for the object allocation.



Applications require a minimum amount of memory to reach a stable state. The stable state occurs when the heap is no longer consistently growing. To reach a stable state, the application has to run through its commonly used code paths and instantiate all of the frequently referenced objects. The load used to test for the stable state is very important. Usually more memory is required to reach a stable state under higher loads. If the JVM is configured with a maximum heap that is too small and never allows the JVM to reach a stable state, then allocation failures will occur and the JVM will throw OutOfMemoryError. Applications should go through rigorous performance tuning before being deployed into a production environment. One goal of performance tuning is to determine the application's stable state and allocate the appropriate amount of memory to the application server.

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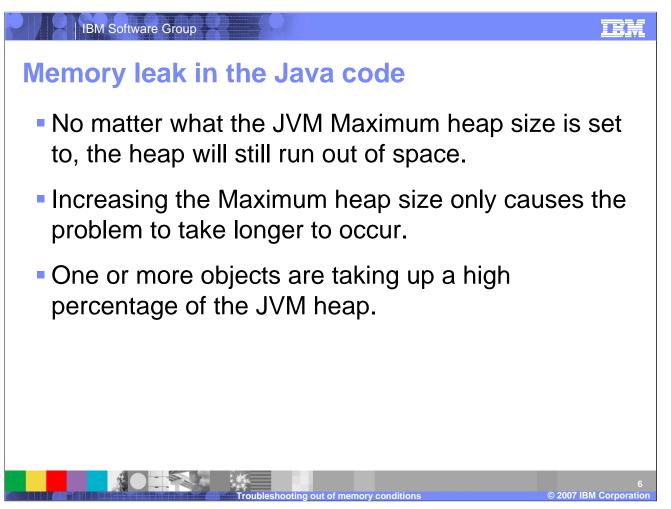
Memory fragmentation

- The difference between the Maximum heap size and the current heap size is only slightly greater than the size of the object to be allocated.
 - Some fragmentation will always occur.
 - Should be treated as if the heap was too small.
- The object being allocated is excessively large.
 - The JVM is attempting to allocate an object that takes up a significant portion of the heap by itself.
 - The application developer should attempt to reduce the size of the object being allocated, and if that's not possible, increase the JVM heap.

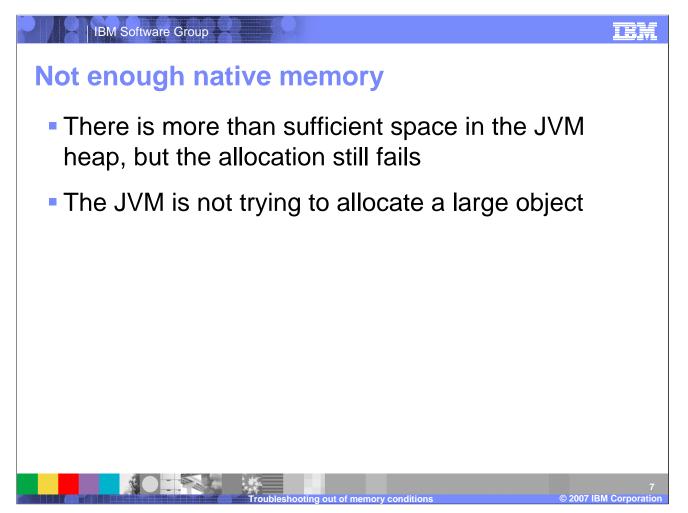
When memory is allocated for an object, the object gets a range of contigous memory from the heap. The next object will get another range of contiguous memory addresses from the heap. At the same time memory is being allocated, it is being deallocated and returned to the heap. Because memory is constantly in the process of being allocated and deallocated, it is impossible for all available memory in the heap to be in one contiguous block. Instead, chunks of allocated memory are fragmented throughout the heap. To be able to allocate an object, memory must be in a contiguous block. It is possible that there is enough available memory in the heap to fulfill an allocation request, but there is not a single block of contiguous memory available so the request fails. There are two specific cases of fragmentation that are commonly seen. The first occurs when the difference between the current heap size and the maximum heap size is only slightly greater than the size of the object to be allocated. Since some fragmentation will always occur within any memory management system, this case should be treated as if the heap size was too small. The second case is when the JVM continuously attempts to allocate excessively large objects. To illustrate the problem of large object allocations, imagine the java heap as a stack of seven blocks. Each block represents one unit of memory. Now imagine that you have an allocation request for three units of memory units. Now imagine that the only contiguous three blocks available are the center three, so you allocate them. Soon after, all other blocks become deallocated, thus there are four available units of memory units. Now imagine that the only contiguous units of memory units of memory units. Now imagine that the only contiguous three blocks available are the center three, so you allocate them. Soon after, all other blocks become deallocated, thus there are four available units of memory. Unfortunately, since the three units in the middle of the heap are allocated, that leaves only t

IER

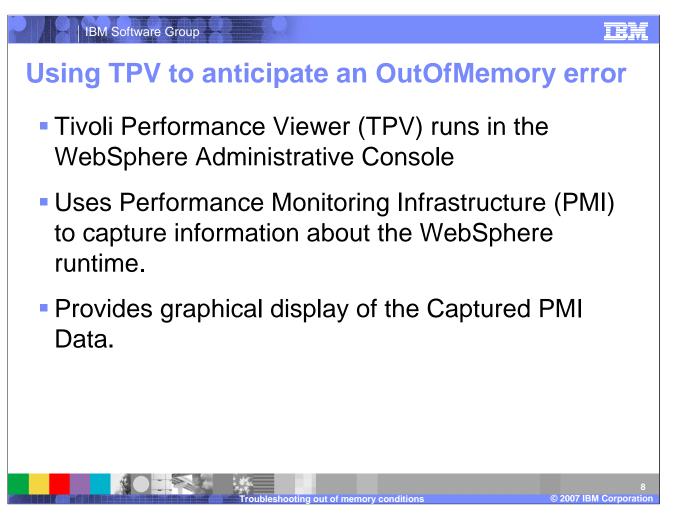
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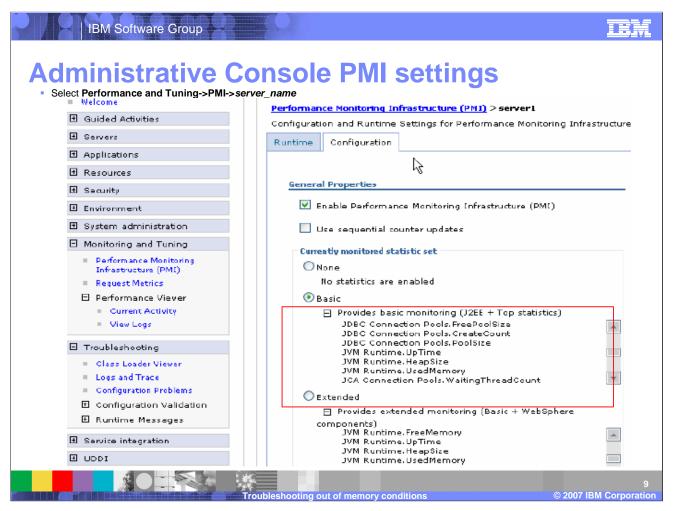
Memory is not explicitly allocated and deallocated in Java, like in C and C++,but it is still possible to create a memory leak. One example would be to save an object into some type of collection. If the collection is a class object, and the class always stays loaded, the object will never be removed from the collection. If objects are continuously added to the collection, the collection could grow until it consumes a significant portion of the java heap. The misuse of object caches is a common cause of memory leaks seen in applications running in WebSphere. One example is the configuration of a large PreparedStatement cache. The PreparedStatement cache size is the total number of PreparedStatements that WebSphere will maintain in the cache. PreparedStatement objects are the parent objects of ResultSet objects, where data from a PreparedStatement is stored. If the ResultSet objects associated with the PreparedStatements in cache are fairly large, and the PreparedStatement cache size is large, then the PreparedStatement cache can grow very quickly and consume significant amounts of available memory. Another commonly seen misuse of cache in WebSphere is the HTTP max Sessions parameter in WebSphere. HTTP Sessions can be written to store very large objects. If the HTTP max sessions parameter is set too high, then it will again consume a large portion of the JVM heap.



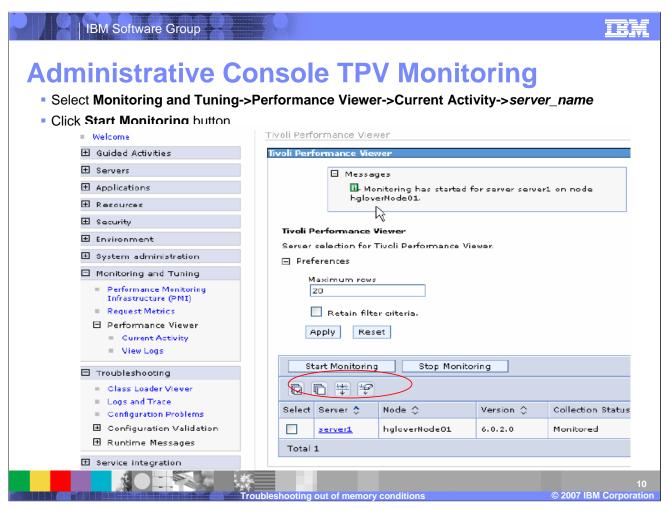
When the Java Virtual Machine is unable to find enough contiguous memory in the heap for object allocation, and it has already called the garbage collection routine, it will then attempt to grow the JVM heap. In order to grow the JVM heap the JVM must request, and be granted, system memory from the operating system Memory Management Unit (MMU). If the MMU is unable to provide memory to the JVM for heap expansion, the JVM will throw OutOfMemoryError. Every Java thread running within a JVM also has an associated thread stack, which requires system memory. If there is not enough memory in the system for the allocation of a Java thread, the system will again throw OutOfMemoryError. Finally, through the Java Native Interface (JNI) the application can access system libraries, which require native system memory to be loaded. If the application fails to remove unused references to JNI objects it will cause a native memory leak. If the system is low on native memory to begin with and is unable to provide memory for the object allocation, the JVM will also throw OutOfMemoryError.



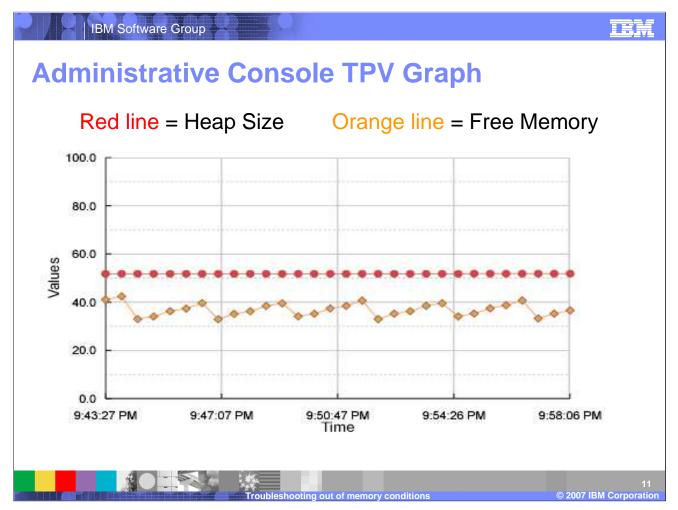
The Tivoli Performance Viewer is embedded within the WebSphere Administrative Client. It uses the PMI infrastructure to capture information about the WebSphere runtime, such as the JVM heap size. We can change the PMI settings in the WebSphere console to capture the information we want to see, set the TPV to begin logging the PMI request information, and view the graphical representation of the data collected. In this case, we want to monitor the size of the JVM heap.



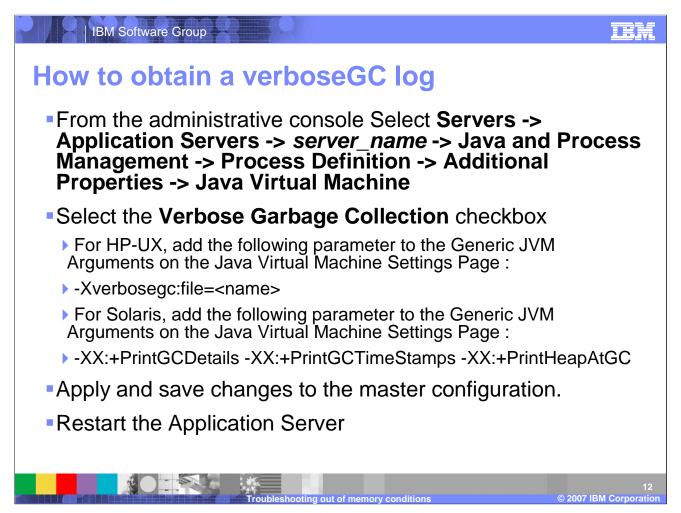
Under Monitoring and tuning, select the `Performance Monitoring Infrastructure' (PMI) link. Under the `Configuration' tab, select the checkbox to enable PMI. Then select 'Basic' from the `Currently monitored statistic set.' From the picture we see this includes `JVM Runtime.HeapSize' statistics.



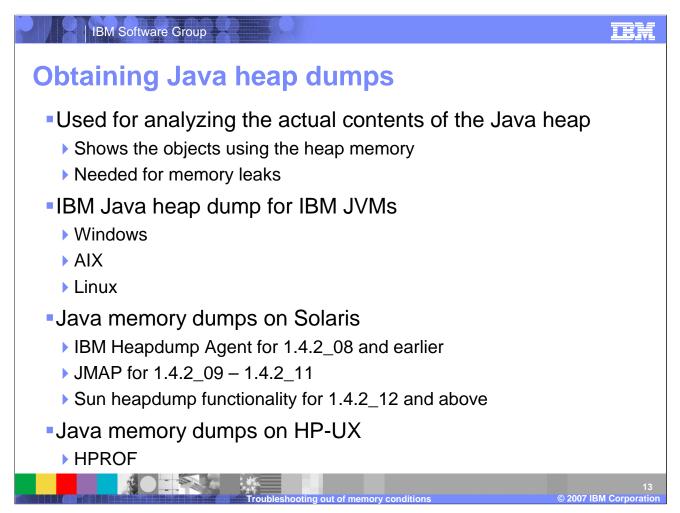
To start monitoring go to `Performance Viewer' and select the `Current Activity' link. Select the Application Server you wish to monitor, and hit the `Start Monitoring' button. Then select the link to your Application Server, in this case `server1', and hit the `Start Logging' button.



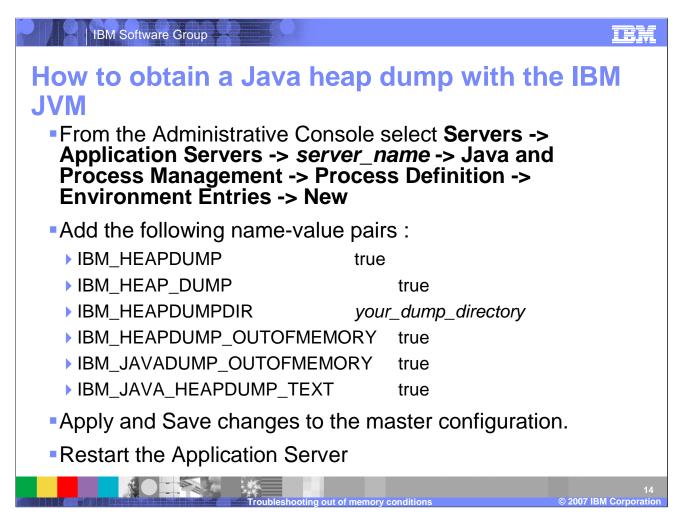
You may need to leave that page and come back in, but the TPV will provide a Graph that inlcudes the JVM Heap Size of the WebSphere runtime over time. In this case the Heap Size is the red line, and the Free Memory is represented by the orange line. We can monitor these values to determine when an OutOfMemory condition is about to occur. In this case, the heap usage is behaving normally. The heap size grows until grabage collection and then returns to the previous relative minimum. If the heap size does not return to the same size as before, the sawtooth patten is continuously increasing, then there is probably a memory leak in the application server.



The verboseGC (verbose garbage collection) output is the diagnostic data used to identify what type of OutOfMemoryError condition is occurring on the system. The verboseGC output shows the size of the object that is being allocated, how much memory is available on the Java Heap, and the JVM's response to the allocation request. In many cases the VerboseGC output will give us a good idea how to resolve the problem, and even if it doesn't it will help us to narrow the scope of the problem. There are some small differences in the VerboseGC output between the Sun,HP-UX, and IBM JVMs. The IBM JVM (Windows®, AIX®, and Linux®) produces highly detailed output by default, and places the information in the native_stderr.log file. The HP-UX JVM also provides detailed ouput, but you have to tell it where to write the output. The Sun JVM on Solaris will write the verboseGC output to the native_stdout.log file, but it does not provide detailed output by default. You have to provide the paramters to tell it to print out the detailed information about the GC event, the timestamp of the GC event, and the size of the JVM heap at the time of the GC event.

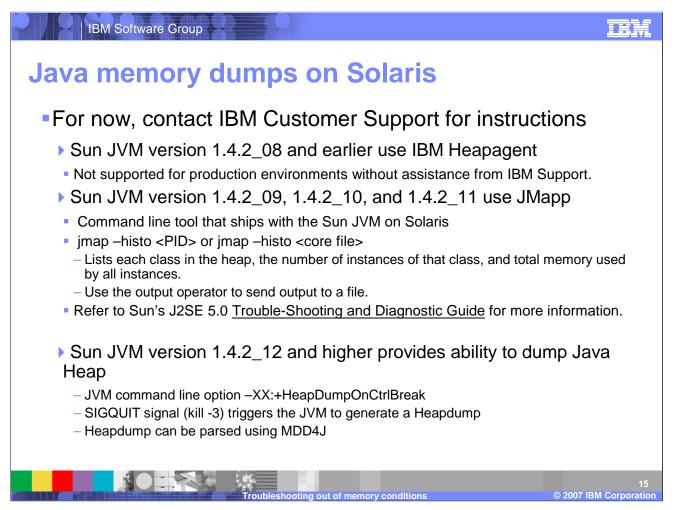


A memory dump of the JVM heap is the primary diagnostic data for identifying memory leak culprits. With the use of special tools, you can analyze the heap dumps to tell us what objects are taking up significant chunks of heap memory. We can also tell who is responsible for instantiating those objects. For IBM JVM's, which run on Windows, AIX, and Linux platforms, you obtain IBM Java Heap Dumps. The JVM is easily configured within the Administrative console to generate an IBM Heap Dump. The advantage of the IBM Heapdump is that there are tools available to parse and process the heapdump file. On Solaris, the tool depends on the version of the Sun JVM. Depending on the version of the JVM you will use the IBM Heapdump Agent, JMAP, or Sun's own heapdump functionality. For HP-UX, the suggested tool for obtaining a dump of the java heap is the Heapdump agent.

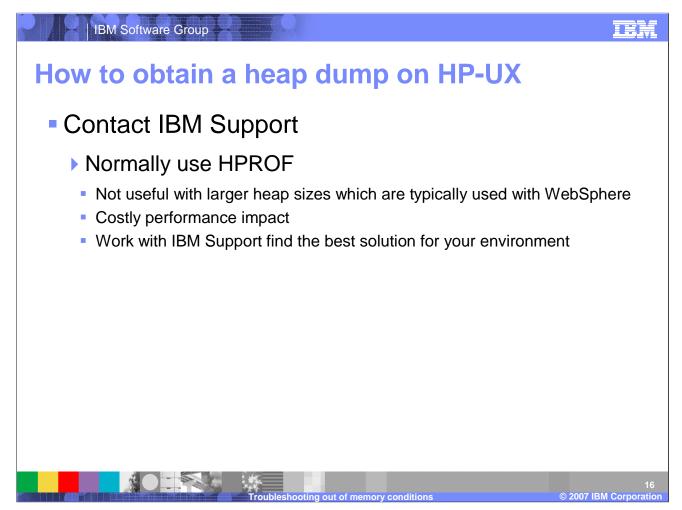


For the IBM JVM, command line parameters must be passed to the Java executable on startup in order to obtain an IBM Java Heapdump. The command line parameters can be set within the WebSphere Administrative console. After setting the parameters, the Application Server must be restarted in order for the JVM to be initialized with the heapdump settings.

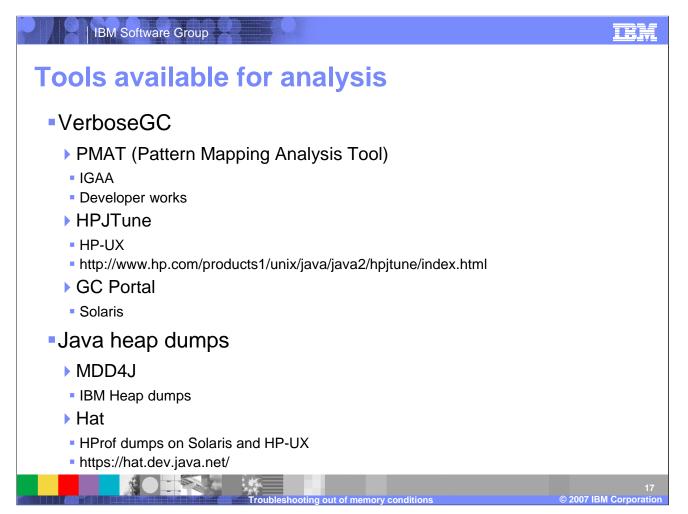
Once setup, the IBM JVM will automatically dump a Java heapdump when an OutOfMemoryError occurs or can be forced by sending a kill -3 to the application server process. There are several options for configuring the Heap Dump. You can specify the location, format, and if one should automatically be generated when an out of memory occurs.



The Sun 1.4.2 JVM has recently undergone a series of changes to it's serviceability code. Because of that, depending on the the version of the service release of the 1.4.2 JVM, there are three different options for obtaining a dump of the JVM heap. For Sun JVM version 1.4.2_08 and earlier, IBM provides the IBM Heapdump agent which can be used to obtain an IBM Java heap dump. The advantage of the IBM Java heapdump is there are tools available to parse and analyze the file. The Heapdump Agent is not supported for production environments without assistance from IBM, so users should contact IBM Support in order to obtain a heapdump for Sun JVM 1.4.2_08 and earlier. Changes in the JVM code for Sun JVM version 1.4.2_09 will cause the JVM to crash if you attempt to attach the IBM Heapdump Agent. Because of this, the IBM Java Development teams suggests that you use JMapp to obtain a dump of the Java heap instead. JMapp is a command line tool that is packaged along with the Sun JVM on Solaris. JMapp can be run on a running process or a core file to obtain a dump of the Java heap. Beginning with Sun JVM 1.4.2_12, Sun provides it's own heapdump mechanism. To configure the Sun Java heapdump, set the JVM command line option, restart the WebSphere Application Server, and then issue a kill -3 against the WebSphere Application Server JVM process to generate a heapdump. The advantage of the Sun Java heapdump is that you will be able to parse and analyze it using MDD4J.



Obtaining a dump of the Java heap on HP-UX is more difficult than with the IBM or Sun JVMs. HPROF is the best tool available, but it often hangs with Java heaps greater than 500 MB. Since most WebSphere Application Server JVM heap sizes are greater than 500 MB, HPROF is rarely useful. It also incurs a heavy performance cost, which negatively affects the behavior of the application. One advantage for HPROF is that there is a tool, HAT, available to parse and analyze HPROF dumps. Users that need to obtain a dump of the Java heap for WebSphere running on HP-UX should contact IBM Support.



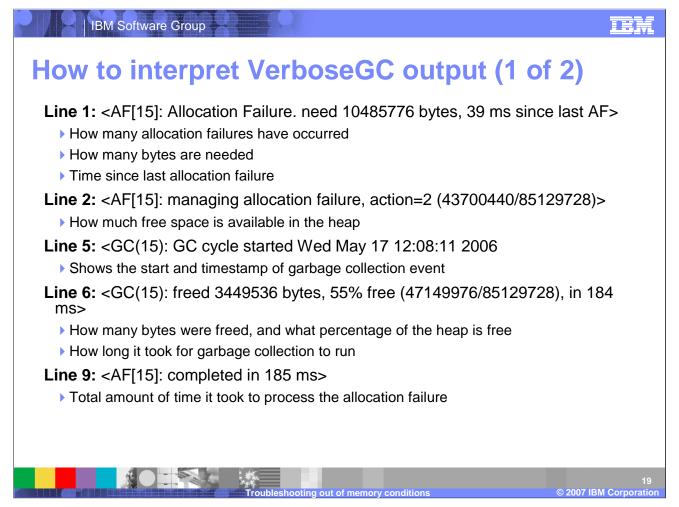
There are several tools available for analyzing diagnostic data for OutOfMemoryError conditions. These tools provide expert analysis, and in the case of MDD4J and HAT, they make analyzing the objects in a Java heap possible.

PMAT is a common tool used to analyze Verbose:gc output for IBM JVMs. It is packaged within ISA and can be launched from within IGAA. PMAT will let you know if you are dealing with insufficient heap space or large object allocations.

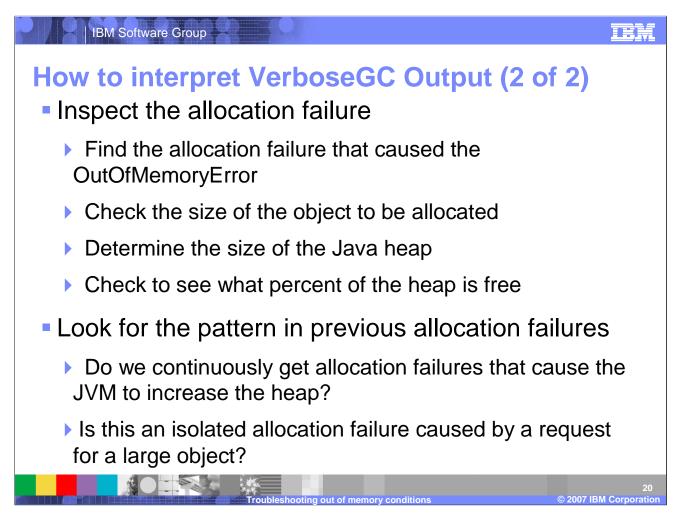
HP-UX and Solaris have similar tools. HPJTune will analyze GC analysis formated by the JVM running on HP-UX. Likewise, the Solaris GC Portal will analyze GC output from a Sun JVM. The GC Portal is a Web Application that must be installed and run on a Web Application Server. A Java heap dump is impossible to interpret without an analysis tool. MDD4J is the preferred tool for analyzing IBM heapdumps from IBM JVMs or the IBM Heapdump Agent, as well as Solaris 1.4.2_12 and above heapdumps. MDD4J is installed under ISA and can automatically be launched from within IGAA. The Heap Analysis Tool, or HAT, is a tool designed to parse Solaris and HP-UX Hprof dumps into a human readable object tree. See the HAT documentation for more information.

```
EM
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Sample VerboseGC output: allocation failure
 <AF[15]: Allocation Failure. need 10485776 bytes, 39 ms</pre>
 since last AF>
 <AF[15]: managing allocation failure, action=2</pre>
 (43700440/85129728)>
   <GC(15): freeing class
 sun.reflect.GeneratedMethodAccessor8(105962D8)>
   <GC(15): unloaded and freed 1 class>
   <GC(15): GC cycle started Wed May 17 12:08:11 2006
   <GC(15): freed 3449536 bytes, 55% free
 (47149976/85129728), in 184 ms>
   <GC(15): mark: 171 ms, sweep: 13 ms, compact: 0 ms>
   <GC(15): refs: soft 0 (age >= 6), weak 0, final 0, phantom
 0>
 <AF[15]: completed in 185 ms>
                                                         © 2007 IBM Corporation
                            ooting out of memory conditions
```

Here is a sample verboseGC output generated by the PlantsByWebSphere application. This sample shows a single allocation failure, and the subsequent garbage collection.



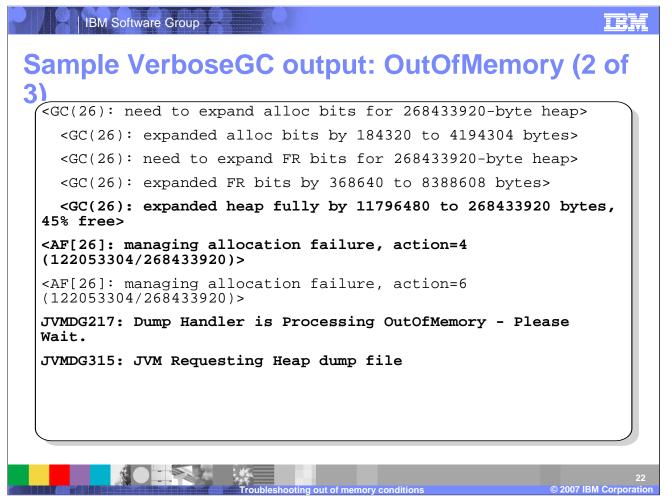
We should start analyzing the verboseGC output by first identifying the key fields in an allocation failure. The first line indicates that an allocation failure occured. The number in brackets is the current count of allocation failures. This line also shows us how many bytes are needed in order to satisfy the memory requirement for the object allocation request, and the time in milliseconds since the last allocation failure. On line 2, it provides more detailed information about the current state of the java heap. The information in parenthesis indicates that total heap size is over 85,000,000 bytes, and that there are over 43,000,000 bytes available. In this case, there is not enough contiguous space available, so garbage collection is started. Lines 5 and 6 show details of the garbage collection. On line 5, garbage collection starts at 12:08:11 on May 17, 2006. On line 6, GC has completed after freeing over 3,000,000 bytes in 184 ms. At this time, 55 % of the heap is now free. Finally, on line 9, it shows that it took a total of 185 ms to process allocation failure 15.



Once we locate the OutOfMemoryError in the log, we want to first examine the Allocation Failure (AF) that caused the event to occur. We need to check the size of the object to be allocated, and the current size of the Java heap. We can also observe what percentage of the heap is currently free. This allows us to determine whether Garbage Collection should have been able to free enough memory for the object allocation, and if not, why. For instance, it would be easy to determine that the maximum heap size was not large enough for all of the requests, or if the object to be allocated is too big relative to the JVM heap. Once we come to a conclusion on the final AF, we should begin looking at AF events leading up to the OutOfMemoryError. Does the JVM heap keep increasing until we finally run into the Max Heap Size? Or does the AF that leads to the OutOfMemoryError seem to be an isolated event? By asking these questions we can determine which of the 4 causes of OutOfMemoryErrors we are experiencing.

iem IBM Software Group Sample VerboseGC output: OutOfMemory (1 of 3) <AF[26]: Allocation Failure. need 167772176 bytes, 55031 ms</pre> since last AF> <AF[26]: managing allocation failure, action=2</pre> (49959168/256637440)> <GC(26): GC cycle started Thu Jun 15 13:50:39 2006</pre> <GC(26): freed 60297656 bytes, 42% free (110256824/256637440), in 1442 ms> <GC(26): mark: 704 ms, sweep: 16 ms, compact: 722 ms> <GC(26): refs: soft 200 (age >= 4), weak 0, final 146, phantom 10> <GC(26): moved 314304 objects, 23247400 bytes, reason=1,</pre> used 48 more bytes> <AF[26]: managing allocation failure, action=3</pre> (110256824/256637440)> <GC(26): need to expand mark bits for 268433920-byte heap> <GC(26): expanded mark bits by 184320 to 4194304 bytes> shooting out of memory conditions © 2007 IBM Corporatio

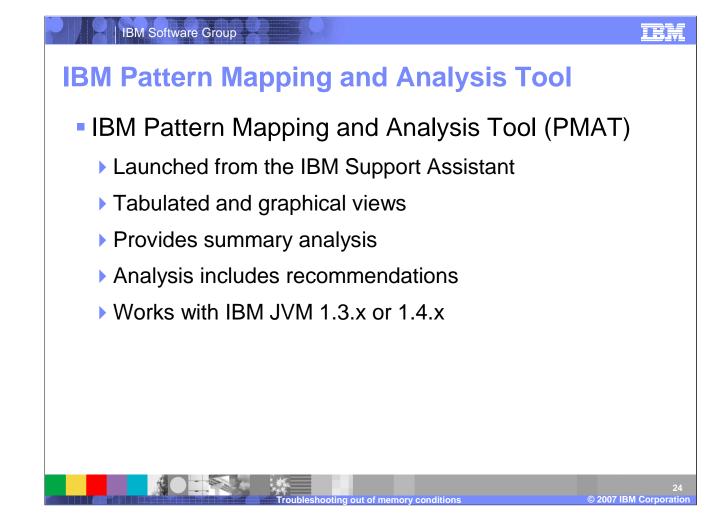
Here is another example that shows the processing of an OutOfMemory condition. The verbose garbage collection output comes from an IBM 1.4.2 JVM. The output will vary for other JVM's, but they will generally produce the same type of information. The first remark in bold in the verbose garbage collection output indicates that there was an allocation failure trying to allocate an object that's requesting 167.7 M of memory. The second line in bold shows the state of the Java heap at that time. The heap size is 256 M, of which only about 50 M of memory is free. The third line in bold shows that the GC cycle freed about 60 M giving a total of 110 M, or 42 % of the 256 M Java heap, available for allocation. This is still insufficient for an allocation request of 167.7, so in the 4 th line in bold the output shows that the garbage collector is starting to manage the allocation failure. The first thing that the garbage collector does is request that the JVM memory management unit expand the heap. Here we can see in the fifth line in bold that the JVM is attempting to expand the heap to 268 M, which in this case happens to be the maximum value.



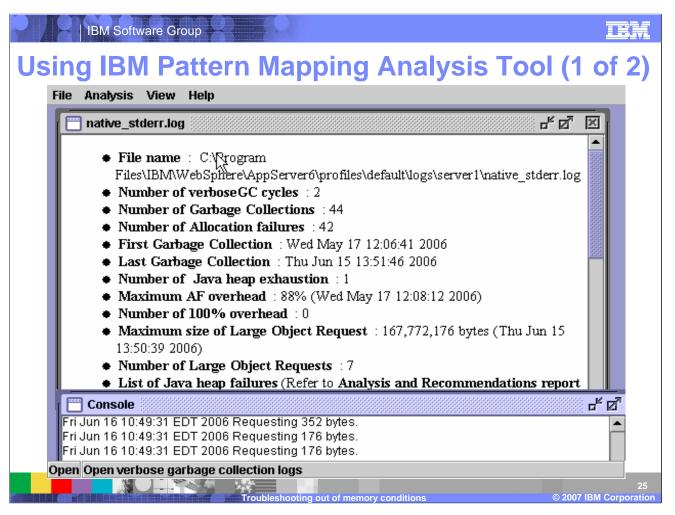
In this first bold line we see that the JVM managed to expand the heap by about 12 meg to 268 M, and 45 % of it is free. On the next line, though, we see that the total free space is still only about 122 M, which is not enough to satisfy the 167 M request. The allocation failure routine then begins OutOfMemory processing. On the last line we see the JVM getting ready to produce a heap dump file.

```
tem
     IBM Software Group
Sample VerboseGC output: OutOfMemory (3 of
JVMDG318: Heap dump file written to C:\Program
Files\IBM\WebSphere\AppServer6\profiles\default\
heapdump.20060615.135039.1448.phd
JVMDG303: JVM Requesting Java core file
JVMDG304: Java core file written to C:\Program
Files\IBM\WebSphere\AppServer6\profiles\default\
javacore.20060615.135054.1448.txt
JVMDG274: Dump Handler has Processed
OutOfMemory.
<AF[26]: Insufficient space in Javaheap to
satisfy allocation request>
<AF[26]: completed in 18805 ms>
        ---
                                            © 2007 IBM Corporatio
```

The verboseGC output shows us where the heap dump file is written to disk. The output then shows the JVM dumping a Java core file, and also identifies the location on disk.



PMAT is a tool written by IBM Support that parses VerboseGC output from IBM 1.3.X or 1.4.X JVM's. PMAT will display all GC events in text, tabulated, and graphical views. One nice feature of PMAT is that it will also interpret the verboseGC information and provide recommendations. PMAT is available as a tool pluggin within IBM Support Assistant.



The PMAT tools contains two windows, one contains the the summary report and the other displays console events.

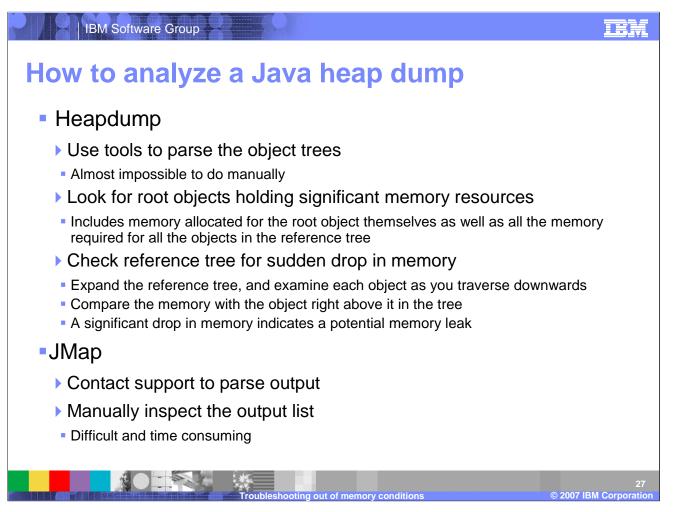
IBM Software Group

Using IBM Pattern Mapping Analysis Tool (2 of 2)

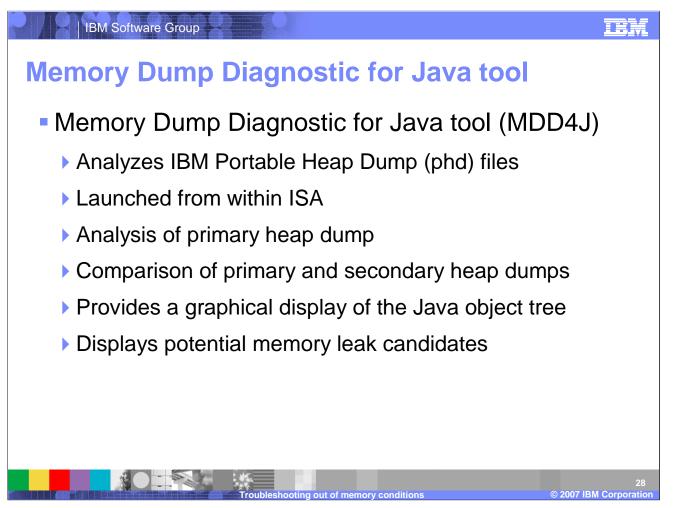
Analysis and Recommendations report				
Garbage collection start / finish	Analysis	Recommendations		
Wed May 17 12:06:41 2006 Wed May 17 12:08:12 2006	No Java heap exhaustion found	There seems to be a steady increase in Java heap usage. (ratio(%): 162.68823 with percentage error(%): 3.9455755)		
Thu Jun 15 13:42:50 2006 Thu Jun 15 13:51:46 2006	Java heap shortage. 167,772,176 bytes requested while 110,256,824 bytes available Thu Jun 15 13:50:39 2006	Increase maximum Java heap size using -Xmx option. If it does not work, review Java heap dump		
ri Jun 16 10:49:3 ri Jun 16 10:49:3	2006 1 EDT 2006 Requesting 352 byte 1 EDT 2006 Requesting 176 byte 1 EDT 2006 Requesting 176 byte	S.		

If you scroll to the bottom of the summary page, the tool provides recommendations. In this example the tool first suggests increasing the java heap, and if that fails it instructs the user to review a Java heap dump.

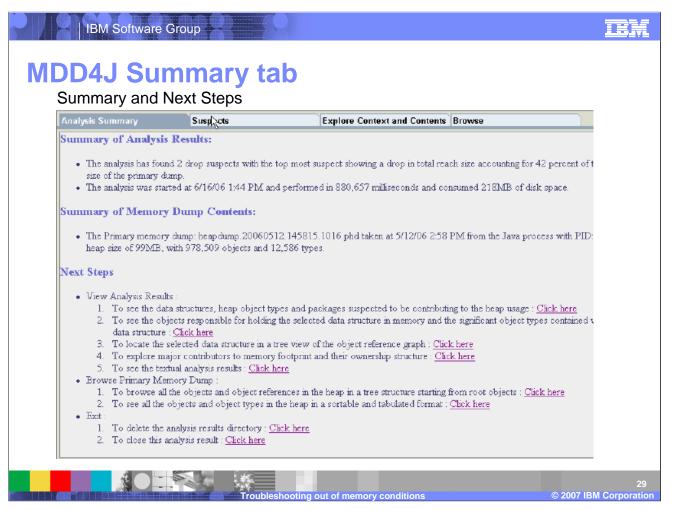
IEM



The most important diagnostic information for debugging memory leaks is a dump of the JVM heap. With proper tools, we can inspect the objects allocated in the JVM heap, traverse the object tree, and look for potential memory leak candidates. We do this by searching for root objects holding a significant portion of the memory in the heap, and traverse the object tree until we find a a location where the memory drops significantly from parent to child. When we find such a drip in memory, the parent becomes a memory leak suspect. There are two types of output that are usefull to IBM support. The first is the IBM formatted heap dump file (php). We can use the MDD4J tool to parse and display the information within the heap dumps. The other type of output is Sun's JMAP output. Parsing JMAP output is difficult and time consuming, and requires IBM development assistance.



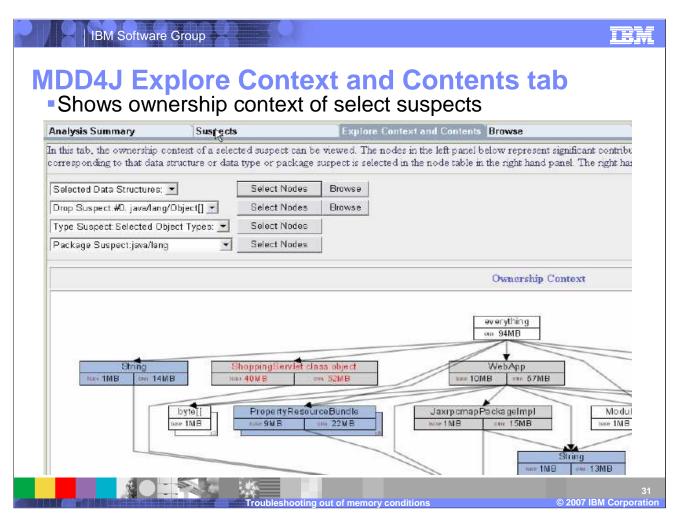
The MDD4J tool provides good analysis and interpretation of phd files. There are two forms of analysis that can be performed. The first is the analysis of a single heap dump from a single OutOfMemoryError failure. The secod is a comparison of a primary (or failure) and secondary (or baseline) heap dumps. This is useful when we are unable to determine the leaking object from the analysis of a single heapdump from the OutOfMemory failure. The baseline heap dump is provided from a JVM running in a healthy state, before the leak has consumed signifcant memory resources. We then compare a failure heap dump from the baseline to determine what the change in allocated objects is. We are focused on the analysis of primary heap dump files. The MDD4J tool will parse the heapdump and provide the results to the user. The first information of interest is the list of memory leak suspects. The tech preview version of MDD4J lists a single suspect, and the subsequent release will list 5 potential suspects. The tool also displays the data graphicaly and within a traversable object tree.



The Analysis Summary page gives a quick summary of information extracted from the heap dump. This information includes the size of the memory dumped to file, and the number of objects contained in memory. It's important to verify that the heap size dumped to file is correct, otherwise we are looking at a truncated heap dump which will not contain useful information.

The usual suspects	Suspects		splore Context and		1
ata structures with large	e drops in reach si	ize			
- Fallenting to be a factor of the					11
				zes from the parent object to its ire reachable from that object di	
e object references in the heap	. These parent objects	with large dro	ops can be indicative	e of array based container object	ts holding or
				of the encapsulating data structu	re.
# Object type of susj	pected container	Reach size	e of the container ob 42MB	ject Drop in reach size 42MB	
<u>0</u> java/lang/Object[] <u>1</u> com/ibm/ws/management/e			42MB 84MB	42MB	
· · · · · · · · · · · · · · · · · · ·	, serving o allo anotholer sp		0 - EALD		
			0-11120		
bject Types that contrib	oute most to heap	size			
bject Types that contrib # Suspected Object Type	nute most to heap	<mark>size</mark> Bytes			
bject Types that contrib # Suspected Object Type <u>0</u> java/lang/String	Number of instances 253,672	size Bytes 7,102,816			
bject Types that contrib # Suspected Object Type <u>0</u> java/lang/String <u>1</u> char[]	Number of instances 253,672 247,718	size Bytes 7,102,816 25,614,366			
bject Types that contrib # Suspected Object Type 0 java/lang/String 1 char[] 2 java/util/HashMapSEntry	Number of instances 253,672 247,718 137,789	size Bytes 7,102,816 25,614,366 3,858,092			
bject Types that contribution # Suspected Object Type 0 java/lang/String 1 char[] 2 java/lutl/HashMapSEntry 3 java/utl/HashMapSEntry[]	Number of instances 253,672 247,718 137,789 26,073	size Bytes 7,102,816 25,614,366 3,858,092 2,641,348			
bject Types that contribution # Suspected Object Type 0 java/lang/String 1 char[] 2 java/utl/HashMapSEntry	Number of instances 253,672 247,718 137,789	size Bytes 7,102,816 25,614,366 3,858,092			
bject Types that contribution # Suspected Object Type 0 java/lang/String 1 char[] 2 java/lang/String 3 java/lang/String 3 java/lang/String 4 java/lang/String	Number of instances 253,672 247,718 137,789 26,073 19,818	size Bytes 7,102,816 25,614,366 3,858,092 2,641,348 951,264			
bject Types that contributed # Suspected Object Type 0 java/lang/String 1 char[] 2 java/utl/HashMapSEntry 3 java/utl/HashMapSEntry[] 4 java/utl/HashMapSEntry[] 4 java/utl/HashMapSEntry[] 4 java/utl/HashMap	Number of instances 253,672 247,718 137,789 26,073 19,818 most to heap size	size Bytes 7,102,816 25,614,366 3,858,092 2,641,348 951,264			
bject Types that contribute # Suspected Object Type 0 java/ang/String 1 char[] 2 java/utl/HashMapSEntry 3 java/utl/HashMapSEntry[] 4 java/utl/HashMapSEntry[] 4 java/utl/HashMapSEntry[] 4 java/utl/HashMap * ackages that contribute # Suspected Package	Number of instances 253,672 247,718 137,789 26,073 19,818 most to heap size Number of instances	size Bytes 7,102,816 25,614,366 3,858,092 2,641,348 951,264			
Dbject Types that contributed # Suspected Object Type 0 java/lang/String 1 char[] 2 java/utl/HashMapSEntry 3 java/utl/HashMapSEntry[] 4 java/utl/HashMap	Number of instances 253,672 247,718 137,789 26,073 19,818 most to heap size	size Bytes 7,102,816 25,614,366 3,858,092 2,641,348 951,264			

The Suspects tab is where MDD4J shows us who it thinks is leaking the memory. The result provided by the current technical preview version of MDD4J only shows one potential suspect. In most cases the first suspect identified will not be the leaking object. In the next version, the top 5 suspects will be displayed, giving us much more reliable information.

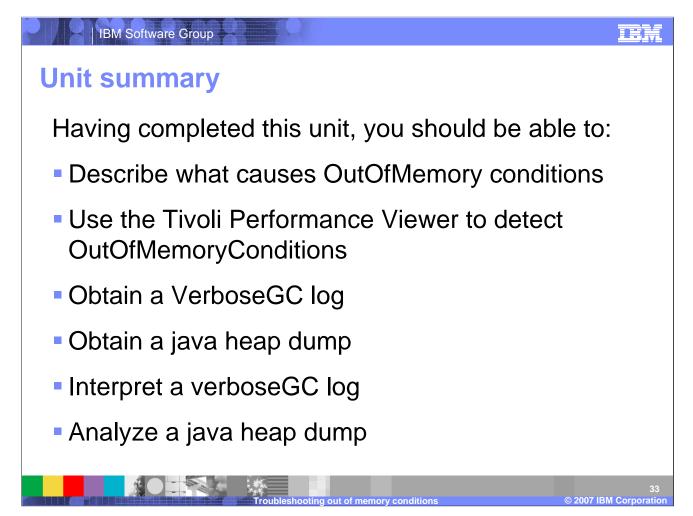


The Explore Context and Contents Tab shows the ownership context of selected suspects. Users can selet nodes listed on the drop down list, and graphically explore the ownership context to identify objects consuming significant memory.

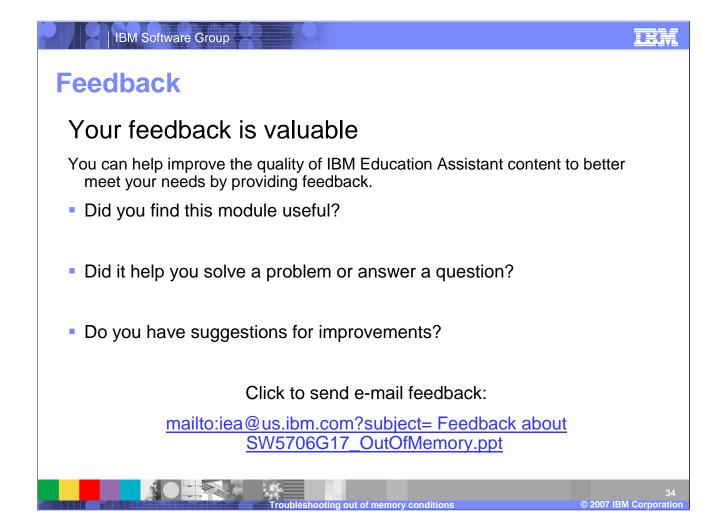
MDD4J Browse tab Traverse the object tree

Find Address Find Address Find Address Image: Solution in the intermediate in the intermediate in the intermediate in the intermediate intermedi	Analysis Summary Suspects charge Ex		xplore Context and Contents Browne			
Address 0x1058c188 Object Class object Name: com/bm/ws/classloader/ProtectionClassLoader Number of children 8 Size (bytes) 100 Total Reach Size (bytes) 92,600,764 Parent Address • Execute Parent Address • Execute 0x103a9ee0 class com/bm/ws/classloader/CompoundClassLoader 0x1041df78 object com/bm/ws/classloader/CompoundClassLoader 0x1041df78 object com/bm/ws/classloader/CompoundClassLoader 0x1041df78 object com/bm/ws/classloader/CompoundClassLoader 0x1041df78 object com/bm/ws/classloader/CompoundClassLoader 0x1041df78 object com/bm/ws/classloader/CompoundClassLoader 0x1041df78 object com/bm/ws/classloader/CompoundClassLoader	Find Address Bookmarks:		The objects and object references in the primary memory fum browsed here in a tree structure. Each node in the tree represe in the Java heap. Its children represent all the outgoing reference object sorted according to their reach sizes. Its parent is any of object from which there is an outgoing reference to this object, details of any particular object (including all its parents) select a tree.			
Address Oxfore from Object Class object Name: com/bm/ws/classloader/ProtectionClassLoader Number of children 8 Size (bytes) 100 Total Reach Size (bytes) 92,600,764 Parent Address 92,600,764 Parent Address Execute Parent Address Parent Object Mame 0s103a9ee0 class com/bm/ws/classloader/CompoundClassLoader 0s1041df78 object com/bm/ws/classloader/CompoundClassLoader 0s1041df78	The following table shows details of a selected object in the tree:					
Number of children 8 Size (bytes) 100 Total Reach Size (bytes) 92,600,764 Parent Address 92,600,764 Parent Address Execute Parent Address Parent Object Name 0x103a9se0 clast com/ibm/ws/classloader/CompoundClassLoader 0x1041df78 object com/ibm/ws/classloader/CompoundClassLoader 0x1041df78 object com/ibm/ws/classloader/CompoundClassLoader	Object Class obj	:ct	🖻 😑 0x11+77808 object iavaluti/WeaktashPep 🛱 2⇒0x11477788 java/uti/WeaktashMapβEntry[]			
Size (synes) 100 Total Reach 92,600,764 Size (bytes) 92,600,764 Actions: Execute Execute Coll325ab0 object conjour/ws/dassbader/Compound/ClassLoader Os1041d78 object object object conjour/ws/classloader/Compound/ClassLoader Os1041d78 object object Solution (Compound/ClassLoader Solutio	Number of 2		# 🗀 0x1050c100 object com/ibm/ws/classloader/ProtectionClassLoader			
Total Reach Size (bytes) 92,600,764 Actions: Execute Parent Address Parent Object Name Oc103a9ee0 class com/bm/ws/classloader/CompoundClassLoader Object com/bm/ws/classloader/CompoundClassLoader Oc1041df78 object java/uti/WeakHastMactEntry Oc1041df78 object java/uti/WeakHast	Size (bytes) 100					
Parent Address Parent Object Name 0x103a9ee0 class com/ibm/ws/cache/ServerCache 0x1041df78 object 0x1041df78 ob	924	\$00,764	💷 😂 0x1525aab8 object java/ubl/WeakHad=Mac\$Entry			
Parent Parent Object Name Image: Construction of the state of the	Actions:	Execute	🌐 🛱 📮0x1525abe0 object org/eclpse/enf/eccre/inpl/EPacka			
0x103a9ee0 class com/ibm/ws/cache/ServerCache III Guitt470750 Guitt47050 Guitt4705050 Guitt4705050 Guitt4705050 Guitt4705050 Guitt4705050 Guitt4705050		Parent Object Name	🖶 📛 0x116d8dc0 object java/utl/WeakHashMap‡Enkry			
	<u>Ox1041df78</u> object com/ibm/ws/classloader/CompoundClassLoader 0~1041df78 object		 Call 1470750 object jevejuti/WeekHashMacdEniny Call 1478660 object jevejuti/WeekHashMacdEniny Call 1478660 object jevejuti/WeekHashMacdEniny Call 1478660 object jevejuti/WeekHashMacdEniny Call 5258680 object jevejuti/WeekHashMacdEniny 			

The Browse tab allows us to traverse the object tree looking for significant drops in memory usage. On the left, we can see the details of the highlighed object in the tree on the right. The key information is the Total Reach Size, which tells us how much memory is being used by the highlighted object, as well as all of the referenced objects below it in the tree. We can identify the leaking object by traversing down the tree until we go from a parent to one of it's children, and the Total Reach Size drops significantly.



Congratulations, you have completed the unit on problem determination for out of memory errors. You should now be able to: Identify what caused an Out of memory condition, use Tivoli Performance viewer to detect OOM conditions, and obtain and analyze the pertinent diagnostic information.



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