

Is It Virtual, Real or Auxiliary?

zSTSU 2004

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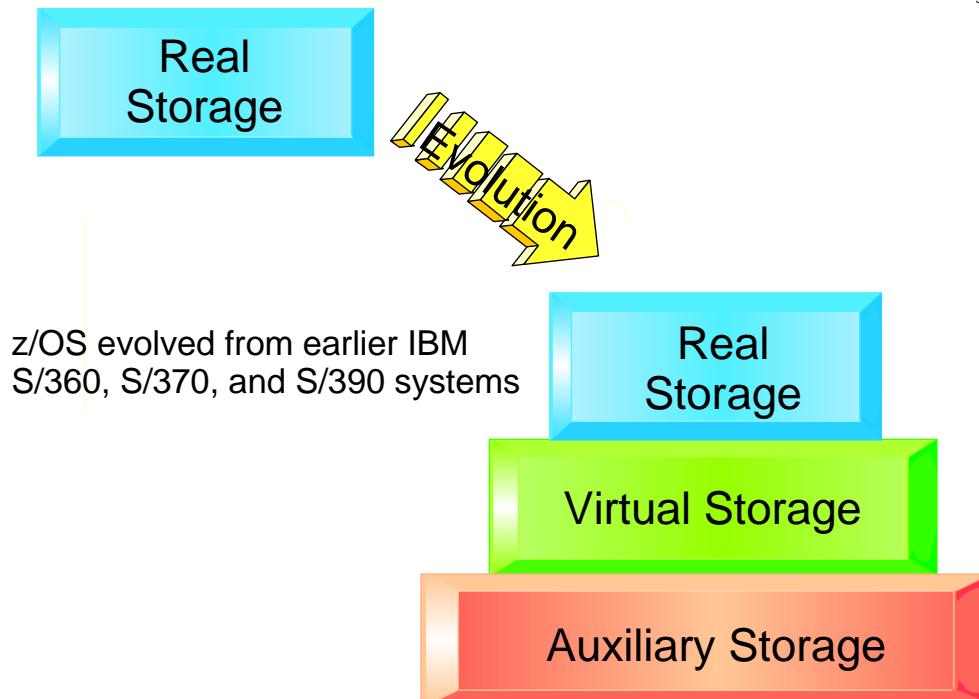
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Agenda



- ▶ **Storage evolution**
- ▶ **Storage overview**
 - Virtual storage
 - Real storage
 - Auxiliary storage
- ▶ **Paging**
 - Demand paging
 - Page stealing
 - Relationship of virtual, real, and auxiliary storage
 - Reports for paging activity
- ▶ **Summary**



In the Past



- ▶ **S/360 MVT - only real storage**
- ▶ **Real storage contained**
 - Operating system code and control blocks
 - Application programs and data
- ▶ **Disadvantages**
 - Storage occupied by inactive routines
 - Limited number of applications could run concurrently
 - Storage fragmentation, prevented larger applications from running

Real Storage is size of storage on processor

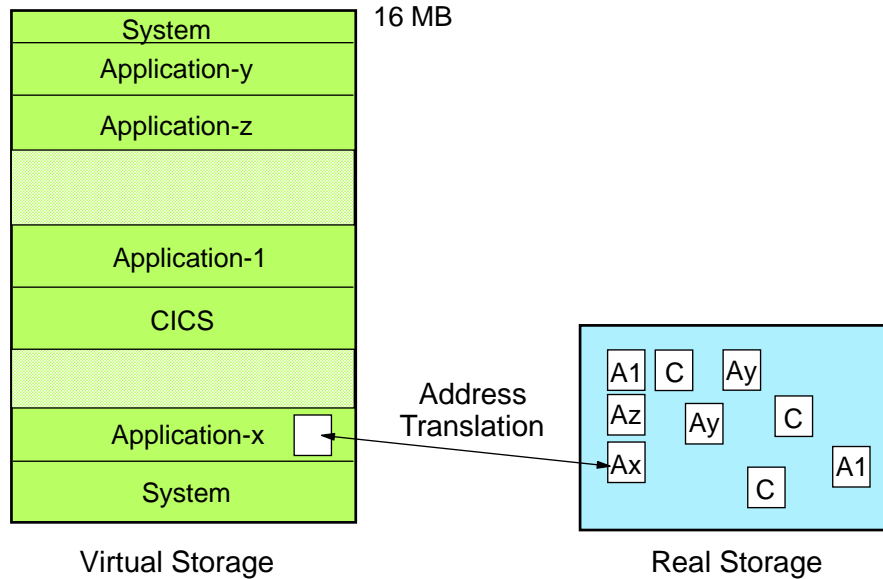
System
Application-y
Application-z
Application-x
System

Introduction of Virtual Storage



- ▶ **S/370 SVS (Single Virtual Storage)**
- ▶ **Virtual storage is a concept created by the architecture and implemented through system page tables and the dynamic address translation (DAT) hardware feature**
- ▶ **Allowed more efficient use of real storage**
 - Inactive portions of programs are no longer in real storage
 - Unused portions of the region are not in real storage
- ▶ **Single virtual storage region of 16 MB**
 - Virtual storage larger than real storage
 - Divided into pages
- ▶ **Disadvantage**
 - Limited number of applications could run concurrently

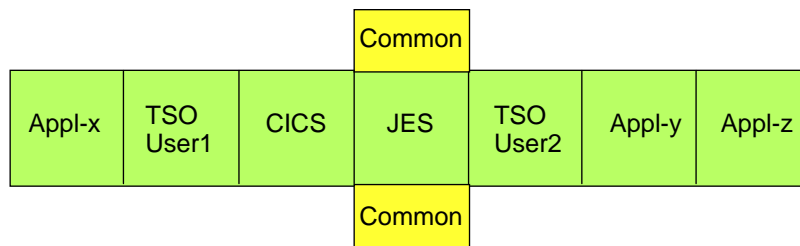
Single Virtual Storage

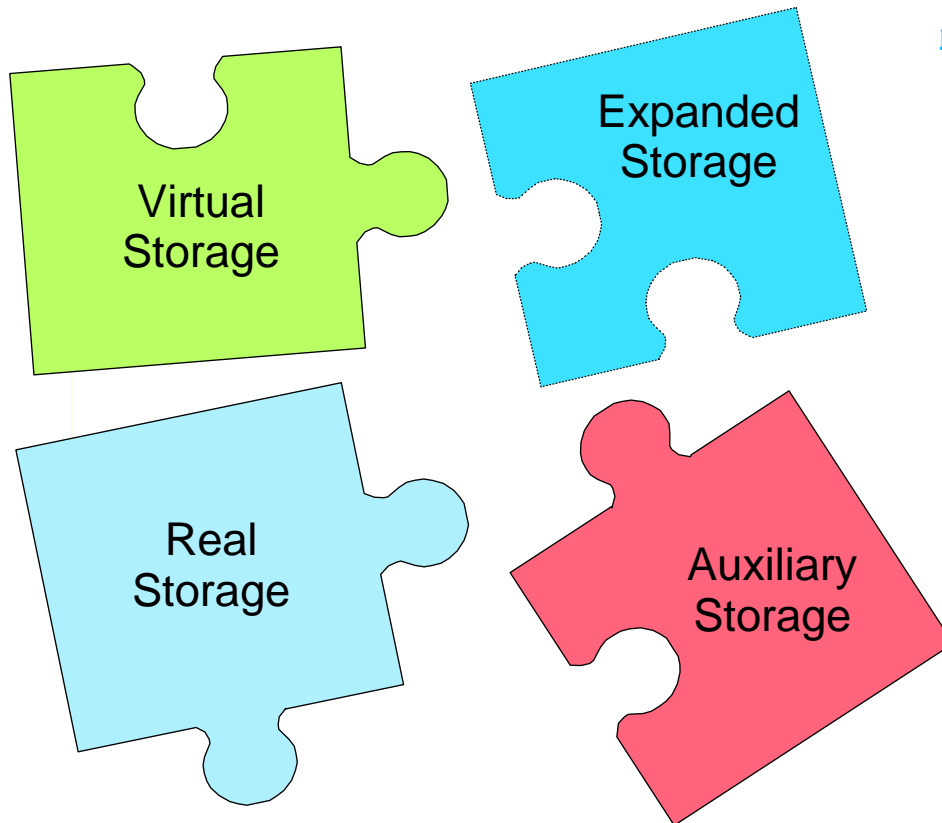


Multiple Virtual Address Spaces



- ▶ **S/370 MVS (Multiple Virtual Storage)**
- ▶ **Each user has its own virtual 16MB address space**
 - Common storage shared by all address spaces
 - Private area is unique for each address space
 - Isolated each application's virtual storage from other applications
- ▶ **Increased number of concurrent applications**



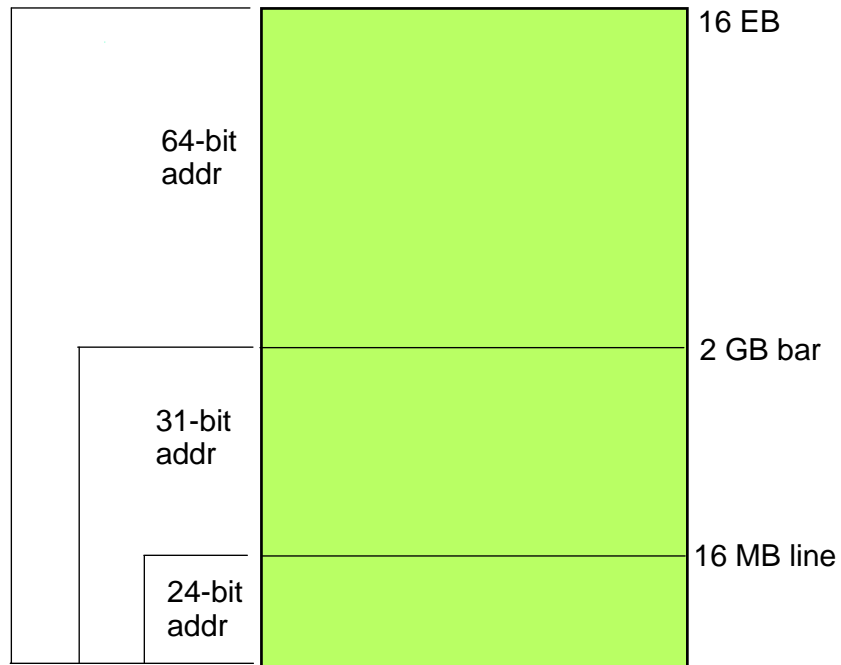


Virtual Storage

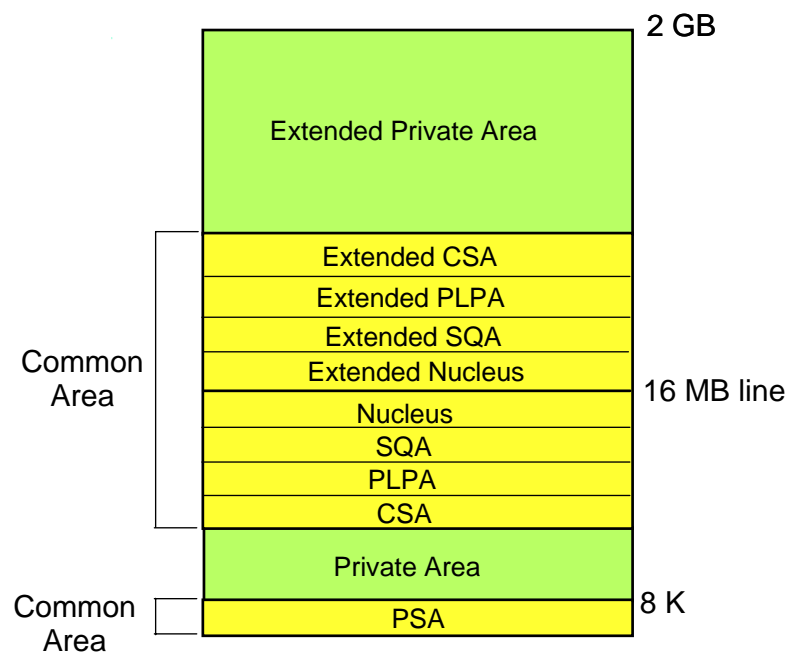
- ▶ Logical view of address space storage
- ▶ Divided into 4K increments called pages
- ▶ Common areas are shared by all address spaces
- ▶ Private areas are unique to each address space
- ▶ Programs are placed in contiguous virtual storage
- ▶ Addressing mode determines size of virtual address space
 - AMODE 24 = 16 bit address = 16 megabytes
 - AMODE 31 = 31 bit address = 2 gigabytes
 - AMODE 64 = 64 bit address = 16 exabytes
- ▶ Storage above the 2 GB bar is used for data and is obtained in 1 MB blocks
- ▶ Virtual storage is managed by the virtual storage manager (VSM)

Virtual is Unreal ...

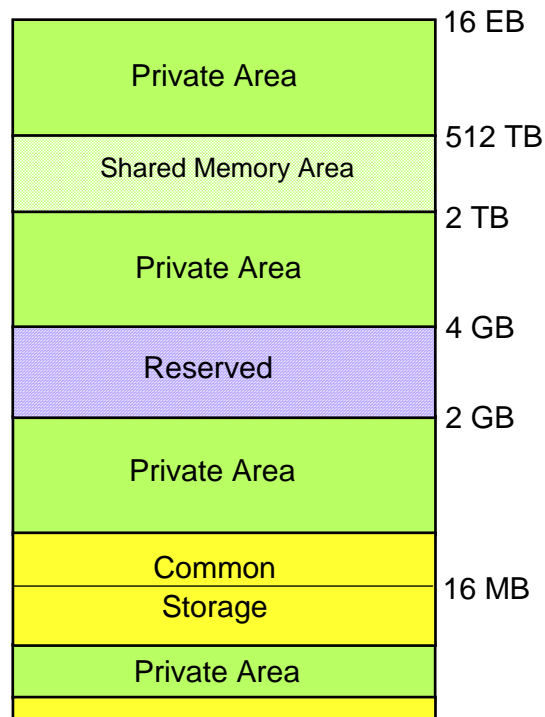
Virtual Storage Addressing



2 GB Address Space



64-bit Address Space



Common Storage

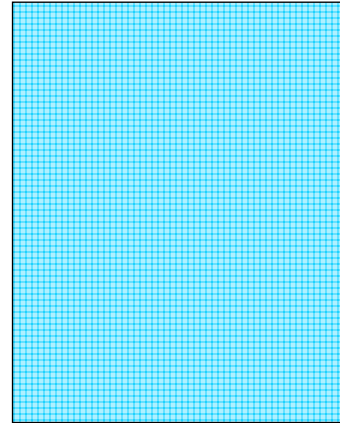


- ▶ **PSA - prefixed storage area**
Contains old PSWs, new PSWs, and register save areas
- ▶ **CSA / Extended CSA - common service area**
Contains data and programs shared across address spaces
Size specified with CSA parameter in IEASYSxx
- ▶ **PLPA / Extended PLPA - pageable link pack area**
Contains read-only, reentrant programs
Programs read into storage at IPL
- ▶ **SQA / Extended SQA- system queue area**
Contains system tables and queues
Size may be specified on SQA parameter in IEASYSxx
If no SQA is available, system will extend SQA into CSA
- ▶ **Nucleus - fixed and non-swappable**
Contains nucleus load module and extensions
Created during IPL

Real Storage



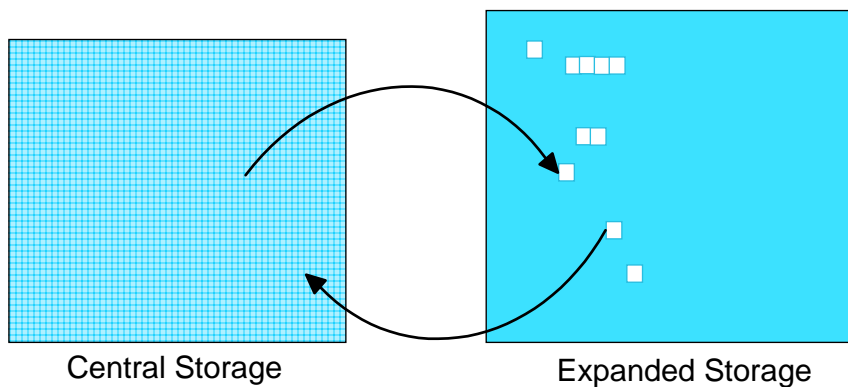
- ▶ Storage with directly addressable byte locations
- ▶ Each byte location is assigned a real storage address
- ▶ Also called central storage or main storage
- ▶ Size depends upon amount of processor storage purchased and amount assigned to logical partition
- ▶ Storage is divided into 4K blocks called page frames
- ▶ Programs must reside in real storage to be executed
- ▶ Real storage is managed by the real storage manager (RSM)



Expanded Storage



- ▶ Storage purchased on a processor may be divided into central storage and expanded storage
- ▶ Expanded storage is an extension to central storage, which is addressed in 4K blocks
- ▶ Expanded storage is not available in z/Architecture mode
- ▶ The contents of a 4K storage frame (a page) may be moved between central and expanded storage



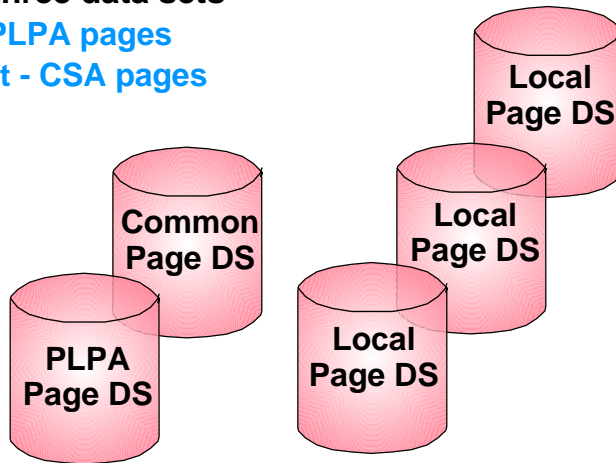
Auxiliary Storage



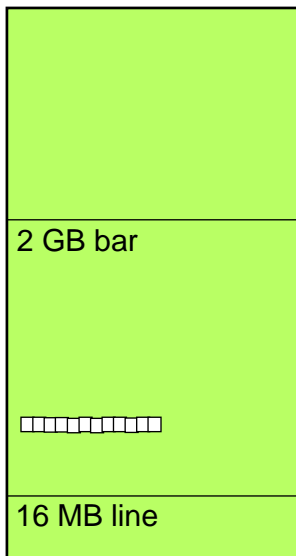
- ▶ VSAM data sets residing on DASD called page data sets
- ▶ Each data set is divided into 4K slots
- ▶ Contain 4K pages for common storage and address spaces
- ▶ Specified on the PAGE parameter in IEASYSxx
- ▶ Must specify at least three data sets

PLPA page data set - PLPA pages
Common page data set - CSA pages
Local page data set

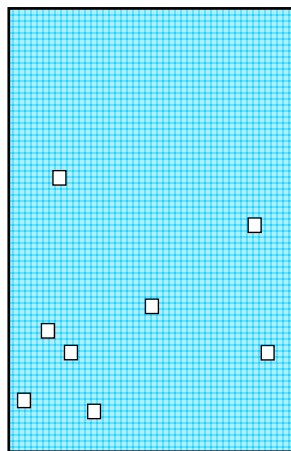
- ▶ May have multiple local page data sets
- ▶ Auxiliary storage is managed by the auxiliary storage manager (ASM)



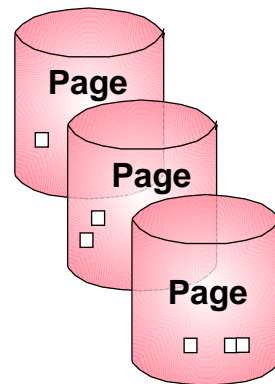
z/OS Storage



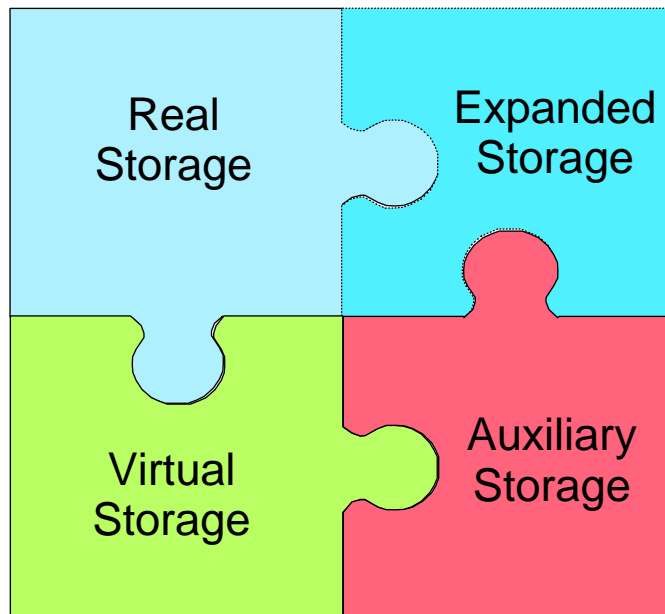
Virtual Storage Pages



Real Storage Page Frames

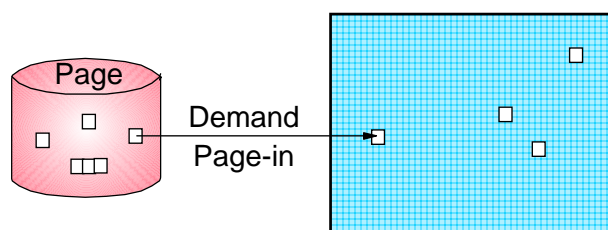


Auxiliary Storage Page Slots



Demand Paging

- ▶ The virtual storage for all the address spaces on the z/OS system is greater than the size of central storage
- ▶ Programs and their data must be in central storage to run
- ▶ When a program is running and references a page which is not in storage a page fault occurs
- ▶ To resolve the page fault, the page demanded by the application is brought into central storage from auxiliary storage - this is called demand paging
- ▶ This results in a page-in, as it is bringing a page into a real storage frame



Least Recently Used Pages



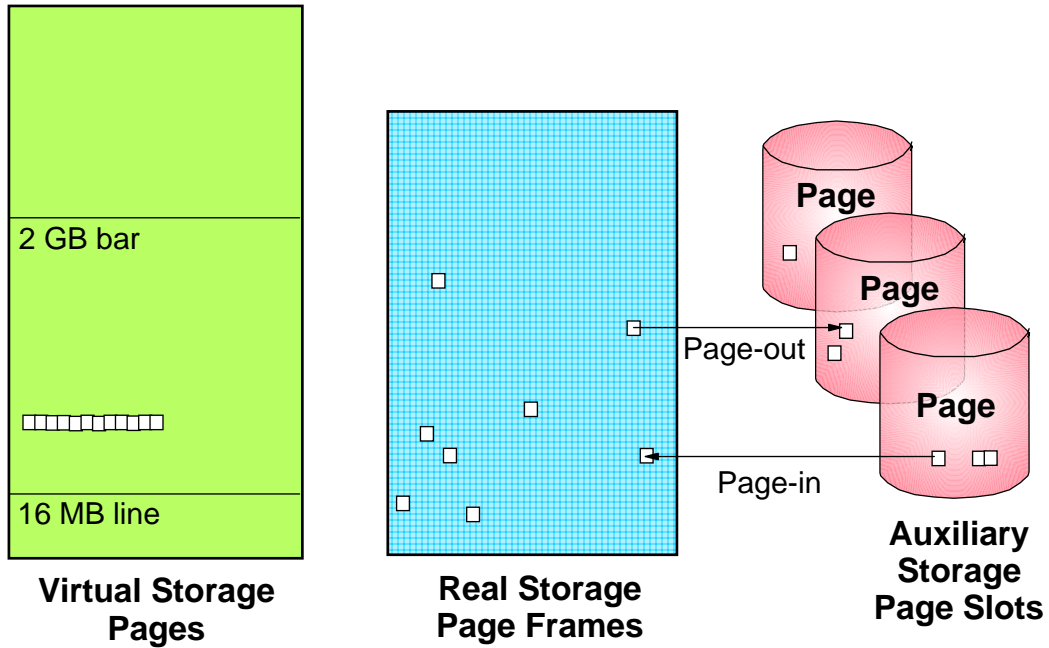
- ▶ **z/OS keeps information about each page in real storage**
 - Change bit - whether the page has been changed
 - Reference bit - whether the page has been referenced
 - Unreferenced Interval Count (UIC) - measure of how long a page has not been referenced
- ▶ **The change bit is turned on when a change is made to a page**
- ▶ **The reference bit is turned on when a program references a page**
- ▶ **Periodically the storage frames are scanned**
 - If the reference bit is on, then the UIC is set to 0
 - If the reference bit is not on, then the UIC is incremented
 - UIC is a measure of how frequently frames are referenced, thus it is an indicator of the "stress" on real storage

Page Stealing



- ▶ **To ensure a real storage frame is available when needed for a page, RSM maintains a queue of available frames**
- ▶ **To replenish the available frame queue, RSM will periodically steal frames containing pages**
 - A least recently used (LRU) algorithm is used to identify the pages to be stolen
 - The UIC is used to identify the least recently used pages
- ▶ **If a stolen page has the change bit on, the page will be written to auxiliary storage - this is called a page-out**
- ▶ **If a stolen page does not have the change bit on, it is not written to auxiliary storage as there is a current copy of it already on auxiliary storage**

Paging - Big Picture



Paging Considerations



- ▶ There is a "cost" for obtaining a page from DASD
- ▶ On a page fault the application waits for the page to be brought into storage
- ▶ If no expanded storage, paging will be to auxiliary (DASD) storage
- ▶ System paging information is available on the Resource Measurement Facility (RMF) Paging Activity Report
 - Page-in rates
 - Page-out rates
 - Page fault rate

Paging Activity - Page Rates



PAGING ACTIVITY

z/OS V1R4 SYSTEM ID SYSA START 12/12/2003-12.54.00 INTERVAL 001.25.59
 RPT VERSION V1R2 RMF END 12/12/2003-14.20.00 CYCLE 1.000 SECONDS
 OPT = IEAOPT00 MODE = ESAME CENTRAL STORAGE PAGING RATES - IN PAGES PER SECOND

CATEGORY	PAGE IN				PAGE OUT				
	SWAP	NON SWAP, BLOCK	NON SWAP, BLOCK	TOTAL RATE	% OF TOTL SUM	SWAP	NON SWAP	TOTAL RATE	% OF TOTL SUM
PAGEABLE SYSTEM AREAS (NON VIO)									
LPA		0.00	0.00	0.00	0				
CSA		0.00	0.00	0.00	0	0.00		0.00	0
SUM		0.00	0.00	0.00	0	0.00		0.00	0
ADDRESS SPACES									
HIPERSPACE		0.00		0.00	0		0.00	0.00	0
VIO		0.00		0.00	0		0.00	0.00	0
NON VIO	0.00	0.97	0.02	0.99	100	0.00	0.80	0.80	100
SUM	0.00	0.97	0.02	0.99	100	0.00	0.80	0.80	100
TOTAL SYSTEM									
HIPERSPACE		0.00		0.00	0		0.00	0.00	0
VIO		0.00		0.00	0		0.00	0.00	0
NON VIO	0.00	0.97	0.02	0.99	100	0.00	0.80	0.80	100
SUM	0.00	0.97	0.02	0.99	100	0.00	0.80	0.80	100
SHARED			0.00	0.00			0.00	0.00	
PAGE MOVEMENT WITHIN CENTRAL STORAGE				6.64		PAGE MOVEMENT TIME %			0.0
AVERAGE NUMBER OF PAGES PER BLOCK				9.9					
BLOCKS PER SECOND				0.10					
PAGE-IN EVENTS (PAGE FAULT RATE)				0.12					

Central Storage Planning



- ▶ **To reduce paging, reduce demand for central storage**
 - Increase the amount of central storage
 - Decrease the amount of central storage used by applications
- ▶ **Estimate amount of central storage needed for applications**
 - Number of real storage pages required by each application also called the working set
 - Evaluate impact of paging on system and applications
- ▶ **Central storage contains fixed pages and pageable pages**
 - Fixed pages reduce the number of pageable frames available
- ▶ **Calculate amount of central storage required for fixed pages and the working sets of active applications**
- ▶ **Central storage usage information is available on the RMF Paging Activity Report**

Paging Activity - Storage Usage



PAGING ACTIVITY

z/OS V1R4 SYSTEM ID SYSA START 12/12/2003-12.54.00 INTERVAL 001.25.59
RPT VERSION V1R2 RMF END 12/12/2003-14.20.00 CYCLE 1.000 SECONDS
OPT = IEAOPT00 MODE = ESAME CENTRAL STORAGE PAGING RATES - IN PAGES PER SECOND

FRAME AND SLOT COUNTS

CENTRAL STORAGE

(585 SAMPLES)	MIN	MAX	AVG
AVAILABLE	185,685	2145457	257,978
SQA	8,400	8,648	8,524
LPA	4,489	4,489	4,489
CSA	12,106	12,624	12,346
LSQA	15,411	23,933	23,586
REGIONS+SWA	957,636	2908558	2836569
TOTAL FRAMES	3145728	3145728	3145716

FIXED FRAMES

NUCLEUS	2,227	2,227	2,227
SQA	7,172	7,418	7,293
LPA	78	78	78
CSA	619	806	622
LSQA	10,909	18,550	18,270
REGIONS+SWA	4,247	7,700	4,370
BELOW 16 MEG	86	122	110
BETWEEN 16M-2G	14,767	17,648	16,792
TOTAL FRAMES	25,257	36,399	32,861

Expanded Storage Considerations



- ▶ **Expanded storage is available in ESA/390 mode**
 - Systems with expanded storage may experience less paging activity to auxiliary storage
 - Page movement between expanded and central is much faster than between auxiliary and central storage
 - The maximum central storage allowed is 2 GB
- ▶ **Expanded storage is not available in z/Architecture**
 - Larger central storage sizes
 - Paging will be done to auxiliary storage page data sets
 - Page data set configuration should be able to efficiently handle the system paging
 - There must be sufficient space on the page data sets

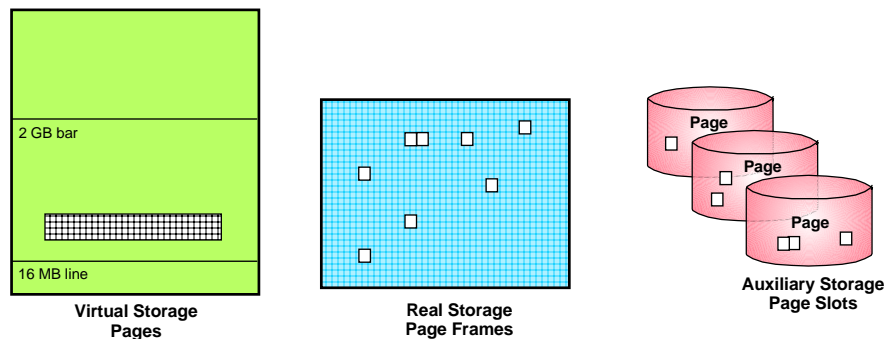
Summary



Virtual storage is implemented by the architecture and is a logical layout for address space storage. The virtual storage pages used by the address space reside in central or auxiliary storage.

Paging is the movement of between central and auxiliary storage.

Information on real storage frame usage, auxiliary storage slot usage, and paging rates is available on RMF reports.



Publications



Product Information - z/OS

[z/OS MVS Initialization and Tuning Guide](#) , SA22-7591

[z/OS MVS Extended Addressability Guide](#), SA22-7614

[z/OS RMF Report Analysis](#), SC33-7991

[z/Architecture Principles of Operation](#), SA22-7832

Websites

z/OS publications

www.ibm.com/servers/eserver/zseries/zos/bkserv/

Technical Document: "z/OS Performance Managing Processor Storage in a 64-bit Environment"

www.ibm.com/support/techdocs/atsmastr.nsf/WebIndex/WP100269

