

# Why Move to HALDB

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**TAKE BACK CONTROL**

**IBM INFORMATION ON DEMAND 2006**  
**October 15 – 20, 2006**  
**Anaheim Convention Center**  
**Anaheim, California**

# Why Move to HALDB?

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- **High Availability through partition independence**
- **Large Databases (Removing Size Restrictions)**
- **Designing Application for HALDB**



# Why Move to HALDB?

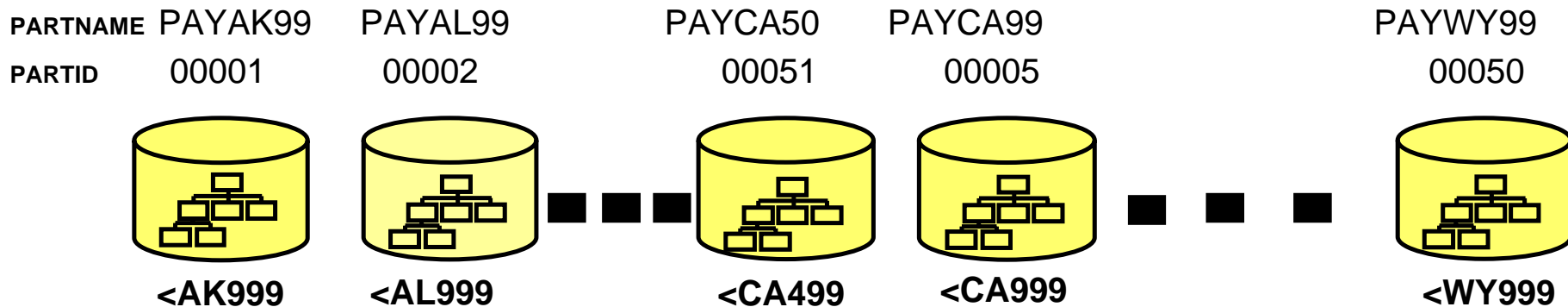
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- **High Availability through partition independence**



# High Availability Partition Independence

- **Partition Independence**
  - ▶ **DBRC authorizes on a partition level**
  - ▶ **Data partitioned into state/country**
  - ▶ **Data partitioned into time zones**



# High Availability Partition Independence

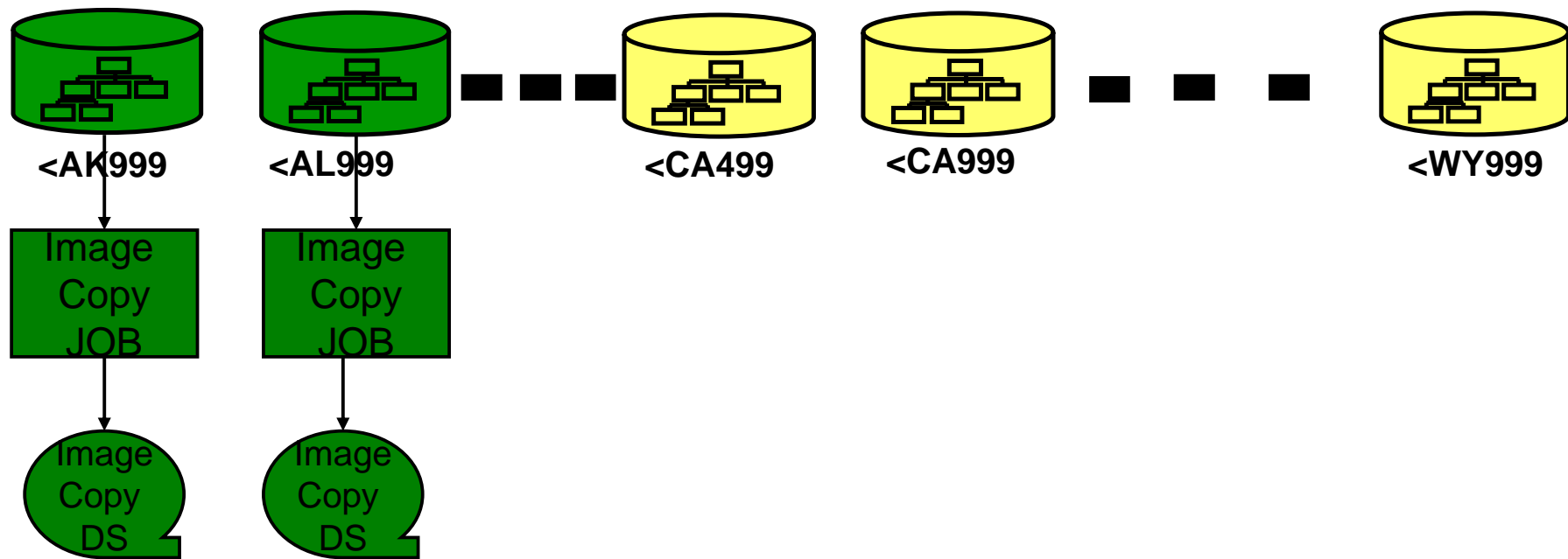
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## Backup Processing



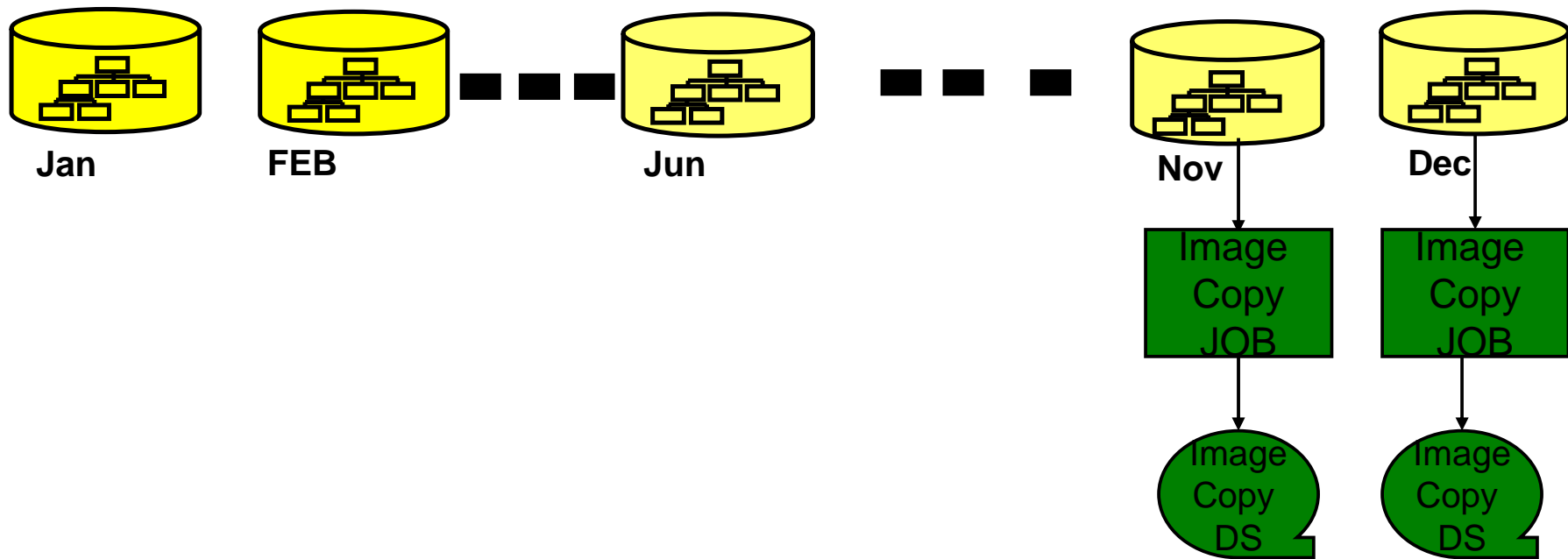
# Backup Processing

- **Partition backups are independent**
  - ▶ **Concurrent processing**
  - ▶ **Shorter elapse times**



# Backup Processing

- **Partition backups are independent**
  - ▶ **Different timing**
  - ▶ **Only backup active partitions**



# High Availability Partition Independence

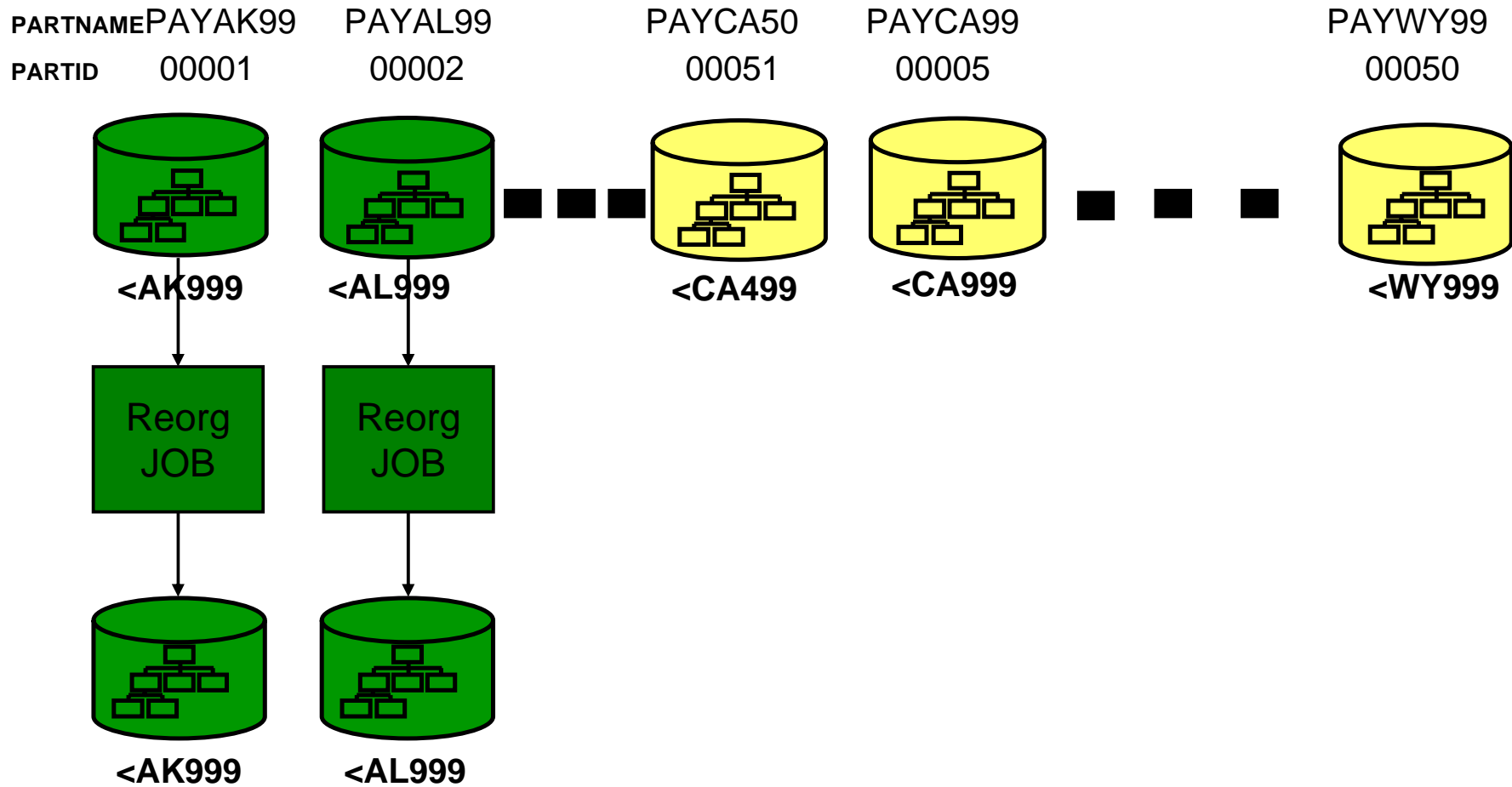
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# Reorganization Processing

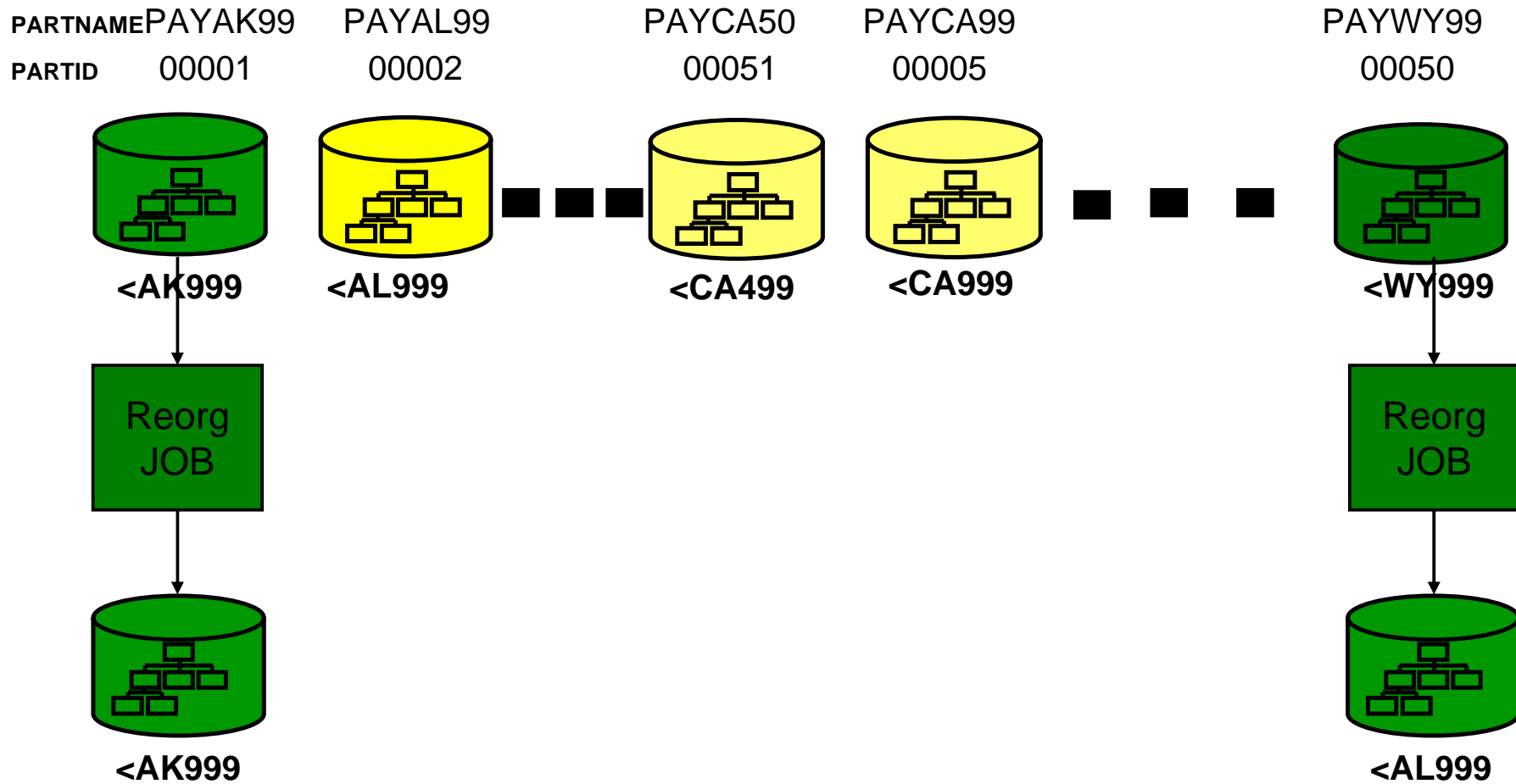




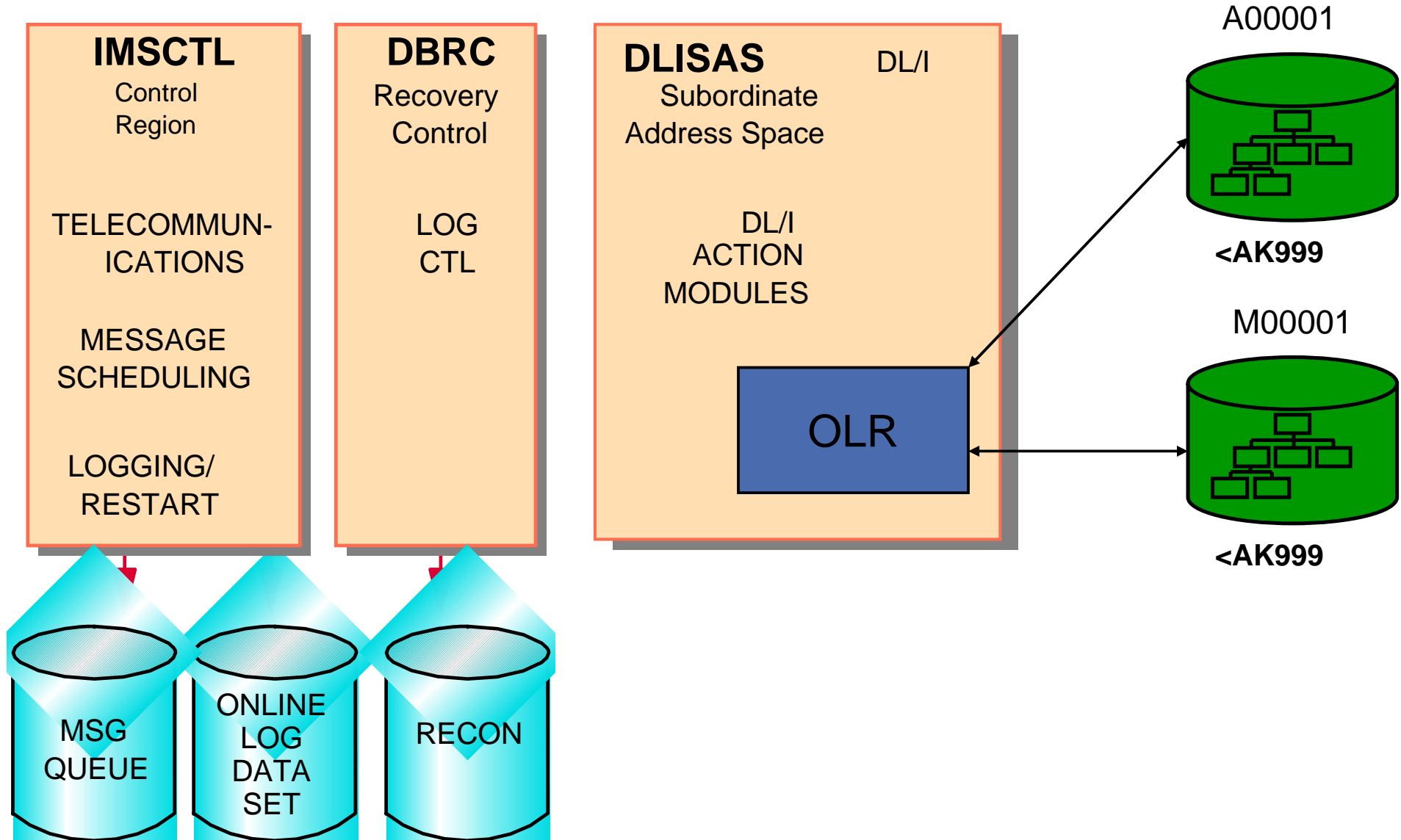
# Reorganization Processing



# Reorganization Processing



# High Availability - Online Reorganization



# High Availability Partition Independence

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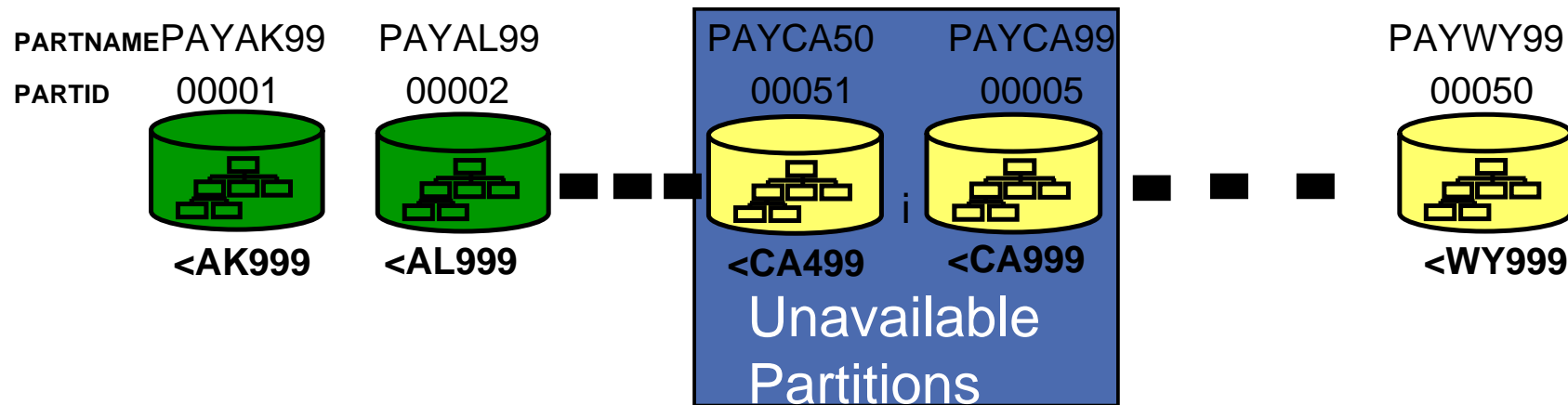
## Increased Application Processing



# Increased Application Processing

## ➤ INIT status call

- Allow programs to handle unavailable partitions
  - Get “BA” status code
  - Avoid the ‘3303’ abend
  - GN after a “BA” moves to next available partition
  - “NA” – Master database not available



# Increased Application Processing

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```
DCL 1 STATUS_CALL
,5 LL1      BIT(08)      INIT('00000000'B)
,5 LL2      BIT(08)      INIT('00000000'B)
,5 LL3      BIT(08)      INIT('00000000'B)
,5 LL4      BIT(08)      INIT('00010001'B)
,5 ZZ1      FIXED BIN(15) INIT(0)
,5 TEXT     CHAR(17)     INIT('STATUS GROUPA  ')
;
CALL PLITDLI($3,INIT,IOPCP_PTR,STATUS_CALL);
```

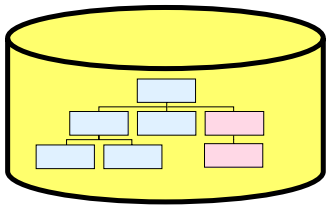


# Increased Application Processing

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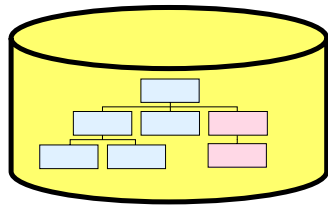
- **Processing partitions independently**
  - ▶ **DFSHALDB DD statement**

SYR2001



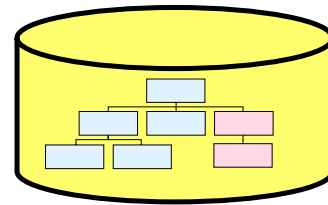
**High Key**  
**20011231**

SYR2002



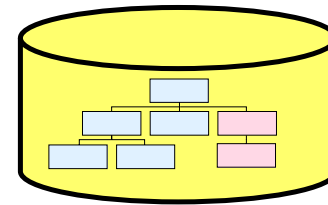
**High Key**  
**20021231**

SYR2003



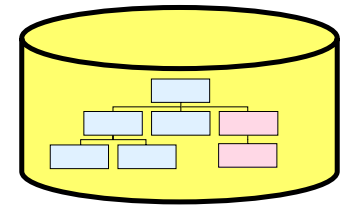
**High Key**  
**20031231**

SYR2004



**High Key**  
**20041231**

SYR2005



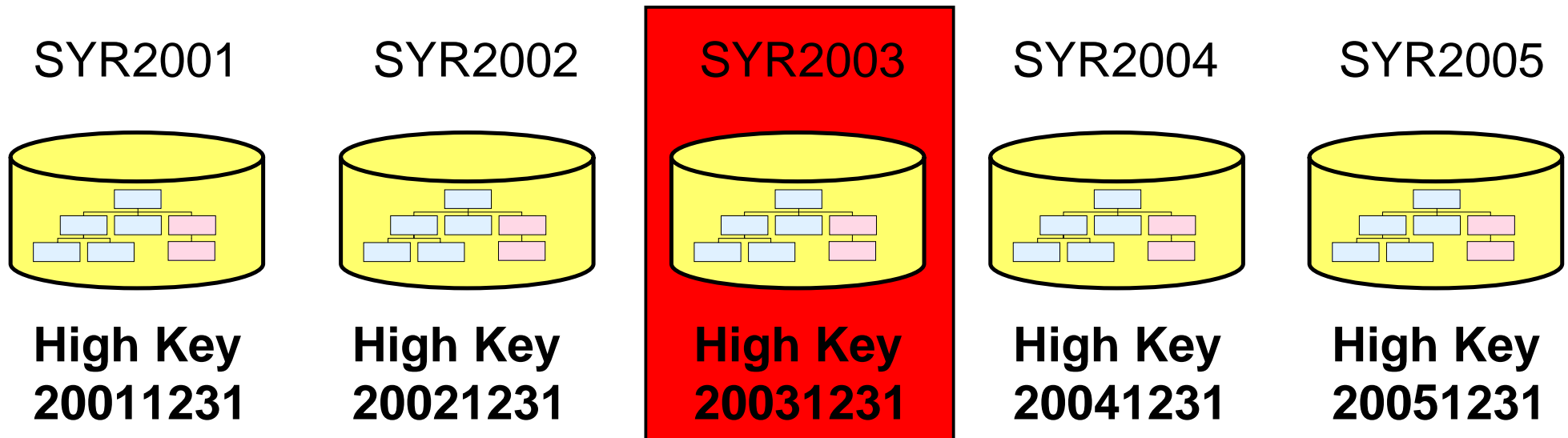
**High Key**  
**20051231**



# Increased Application Processing

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- **Processing partitions independently**
  - ▶ **DFSHALDB DD statement**
  - ▶ **Single partition processing**

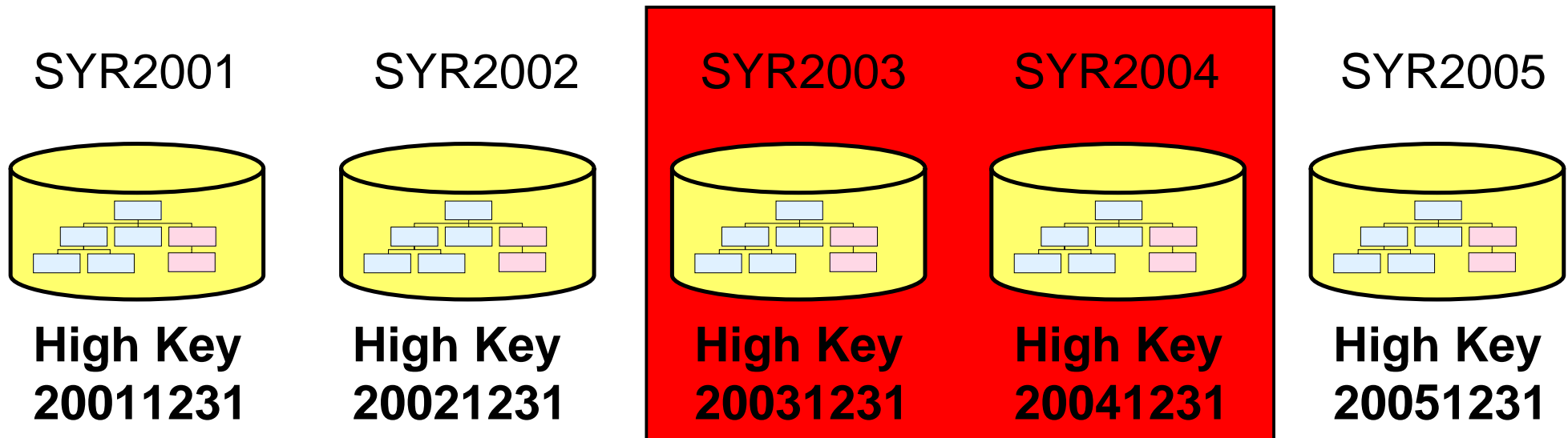




# Increased Application Processing

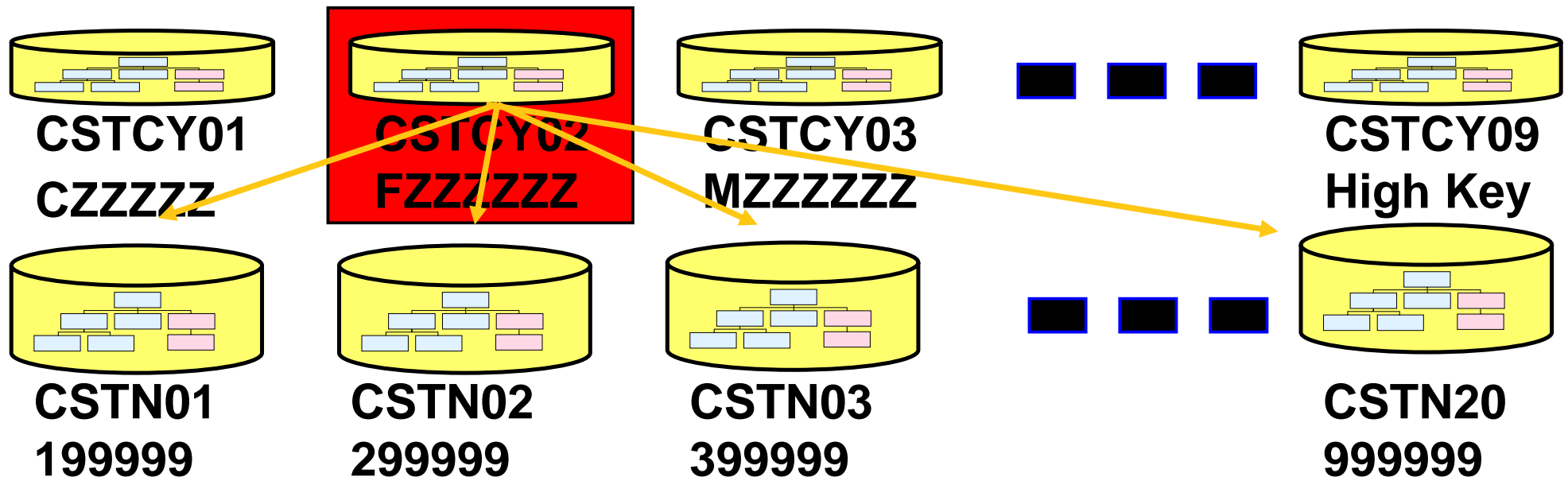
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- **Processing partitions independently**
  - ▶ **DFSHALDB DD statement**
  - ▶ **Single partition processing\**
  - ▶ **Partition name,xx - sequential partitions**
  - ▶ **IMS V9 PK04880**



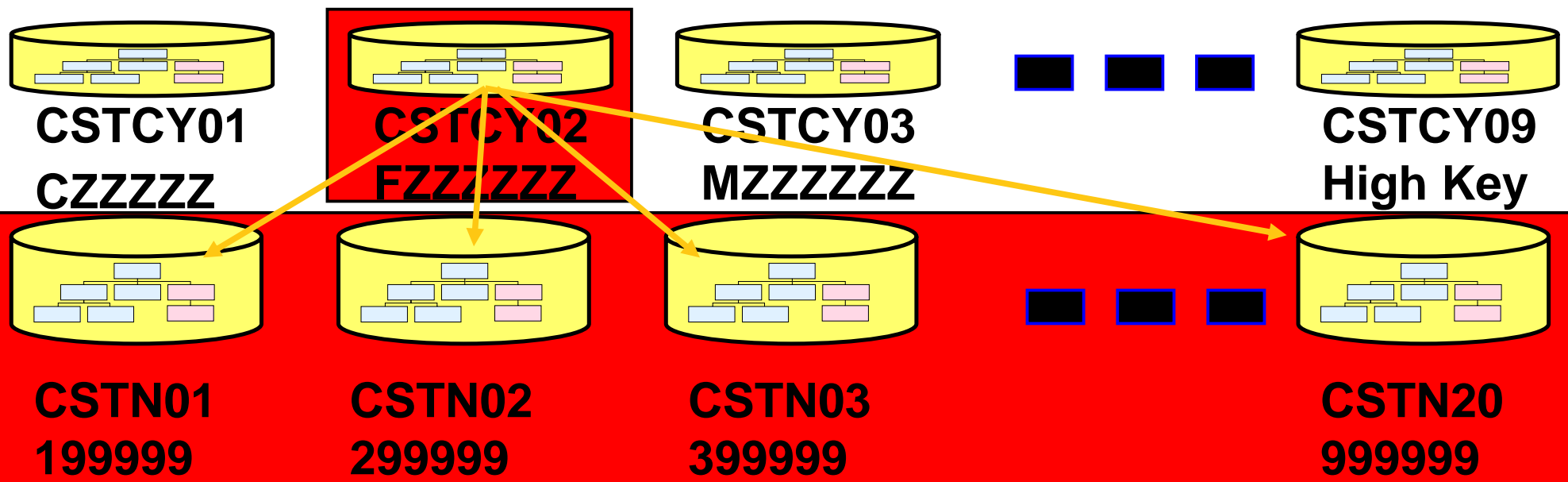
# Increased Application Processing

- ▶ **Processing Partitions Independently**
  - ▶ **Processing is limited to the partition of the secondary index**



# Increased Application Processing

- ▶ **Processing Partitions Independently**
  - ▶ **Source HALDB segments are not limited by the DFSHALDB parm**



# Why Move to HALDB?

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## Large Databases (Removing Size Restrictions)



# Large Databases

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- **Removes the 4GB VSAM and 8GB OSAM limits**
- **Removes the need for data set groups if they were introduced for space reasons**
  - **4GB X 10 data sets = 40GBs**
  - **4GB X 1001 data sets = 4004GBs (without dataset groups)**



# Large Databases

- Partitions allows the application to retain more data
- Historical data retained instead of archiving to a history database

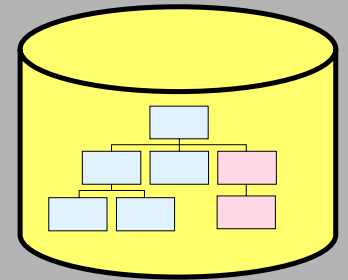
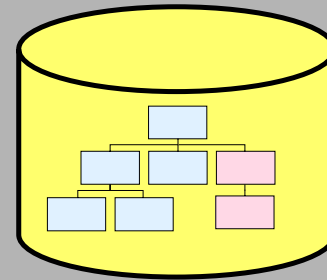
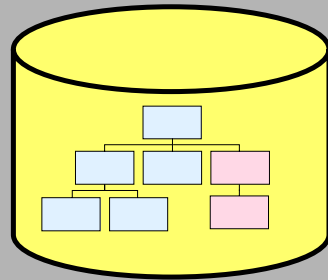
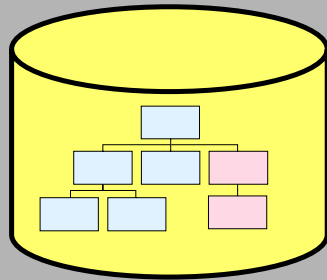
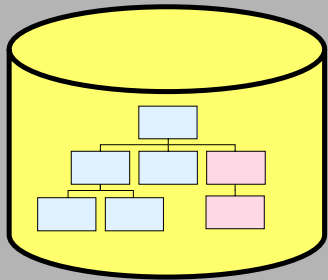
SYR2001

SYR2002

SYR2003

SYR2004

SYR2005



High Key  
20011231

High Key  
20021231

High Key  
20031231

High Key  
20041231

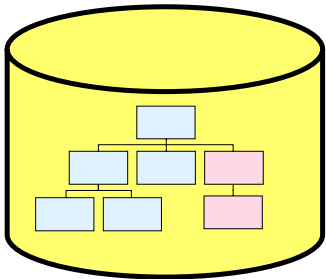
High Key  
20051231

# Large Databases

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- PHIDAM with ever increasing root key values?
- Invoice #s, Acct Payable #s,

INV001



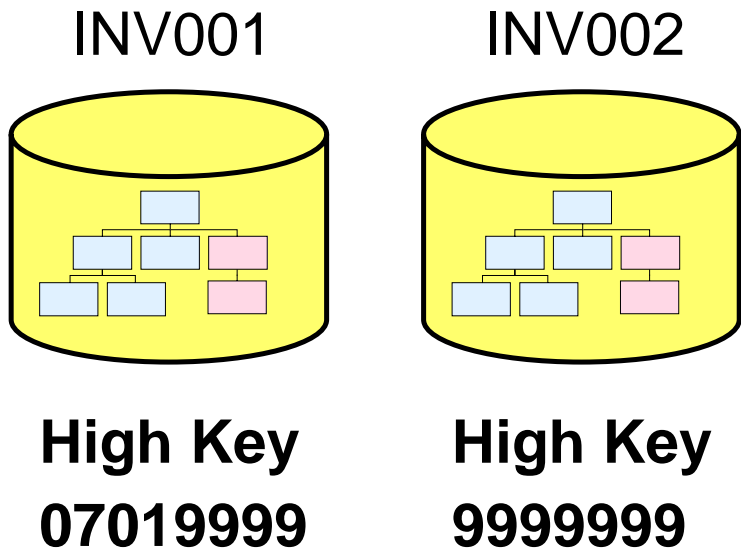
**High Key**  
**9999999**



# Large Databases

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- PHIDAM with ever increasing root key values?
- Invoice #s, Acct Payable #s,

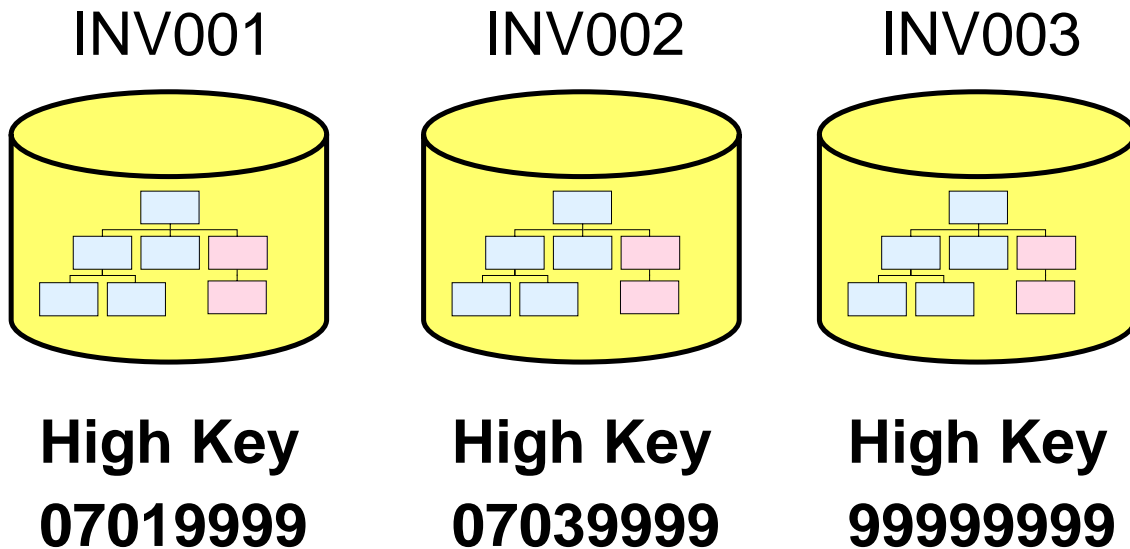




# Large Databases

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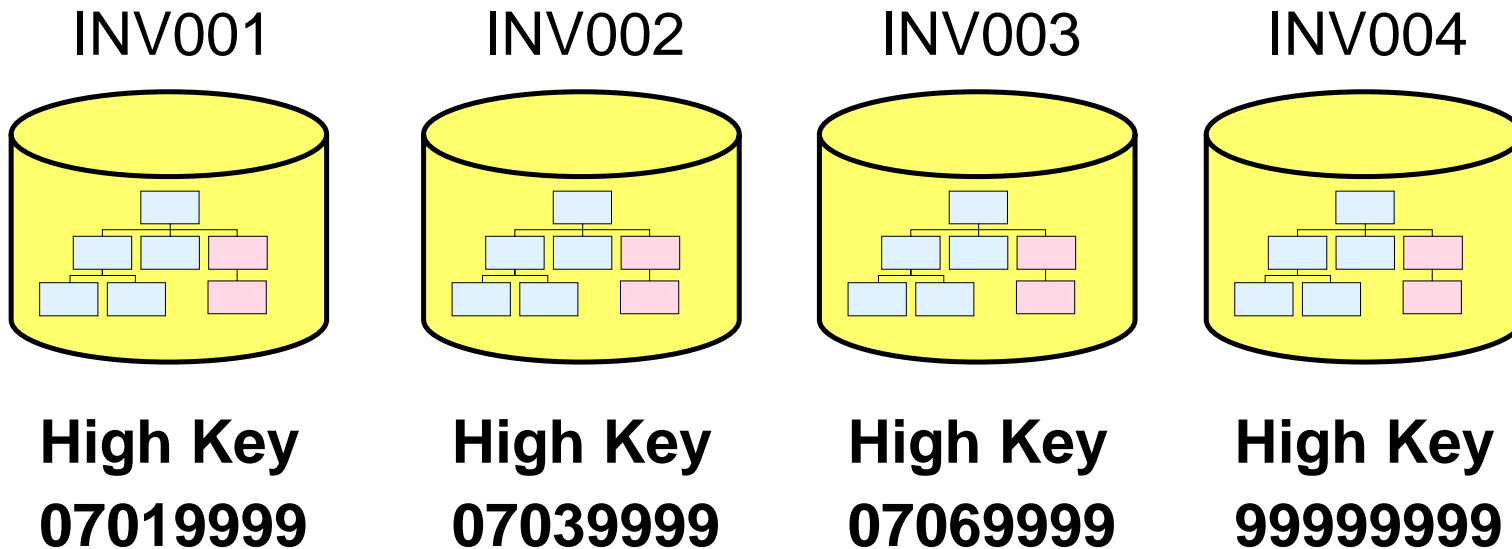
- PHIDAM with ever increasing root key values?
- Invoice #s, Acct Payable #s,



# Large Databases

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- PHIDAM with ever increasing root key values?
- Invoice #s, Acct Payable #s,



# Database Designs for HALDB

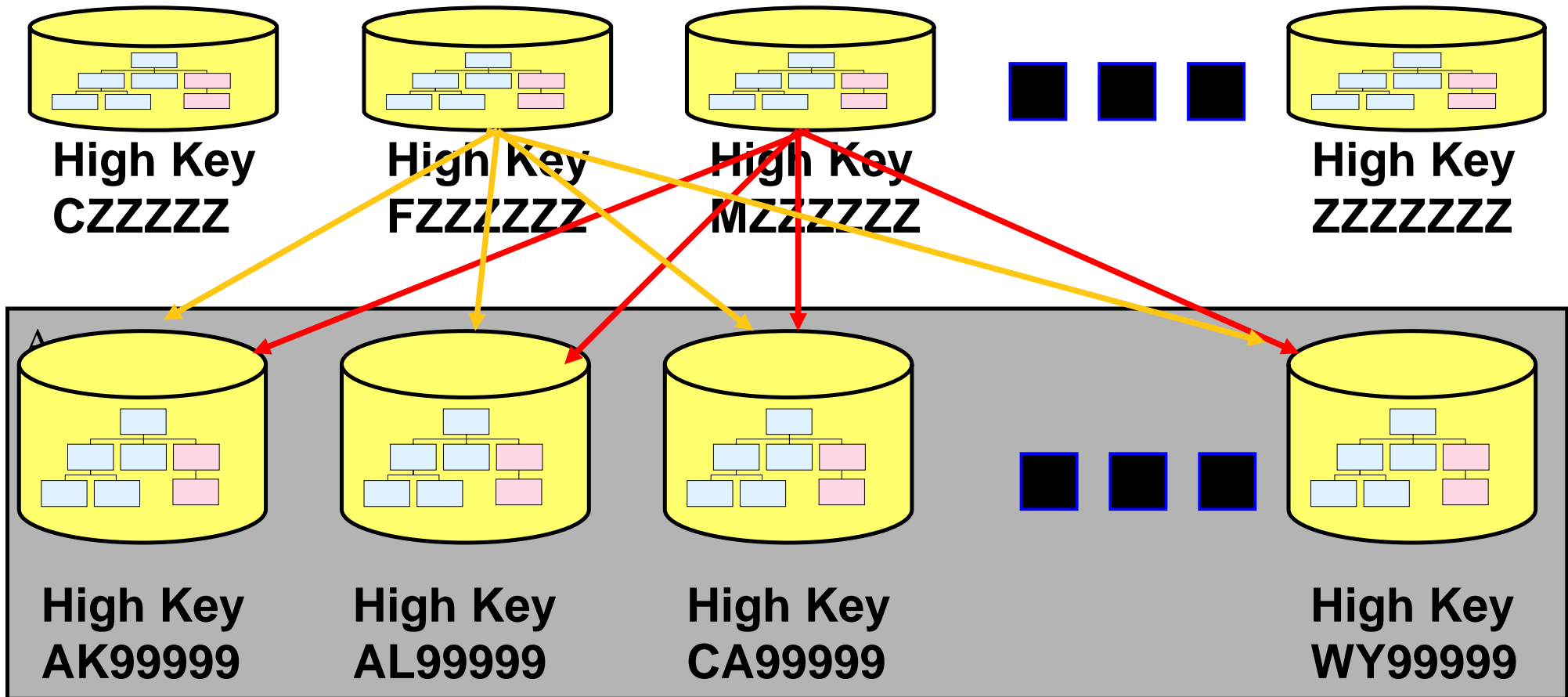
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## Designing databases for HALDB



# Secondary Index independence

- Processing via a secondary index



# Secondary Index independence

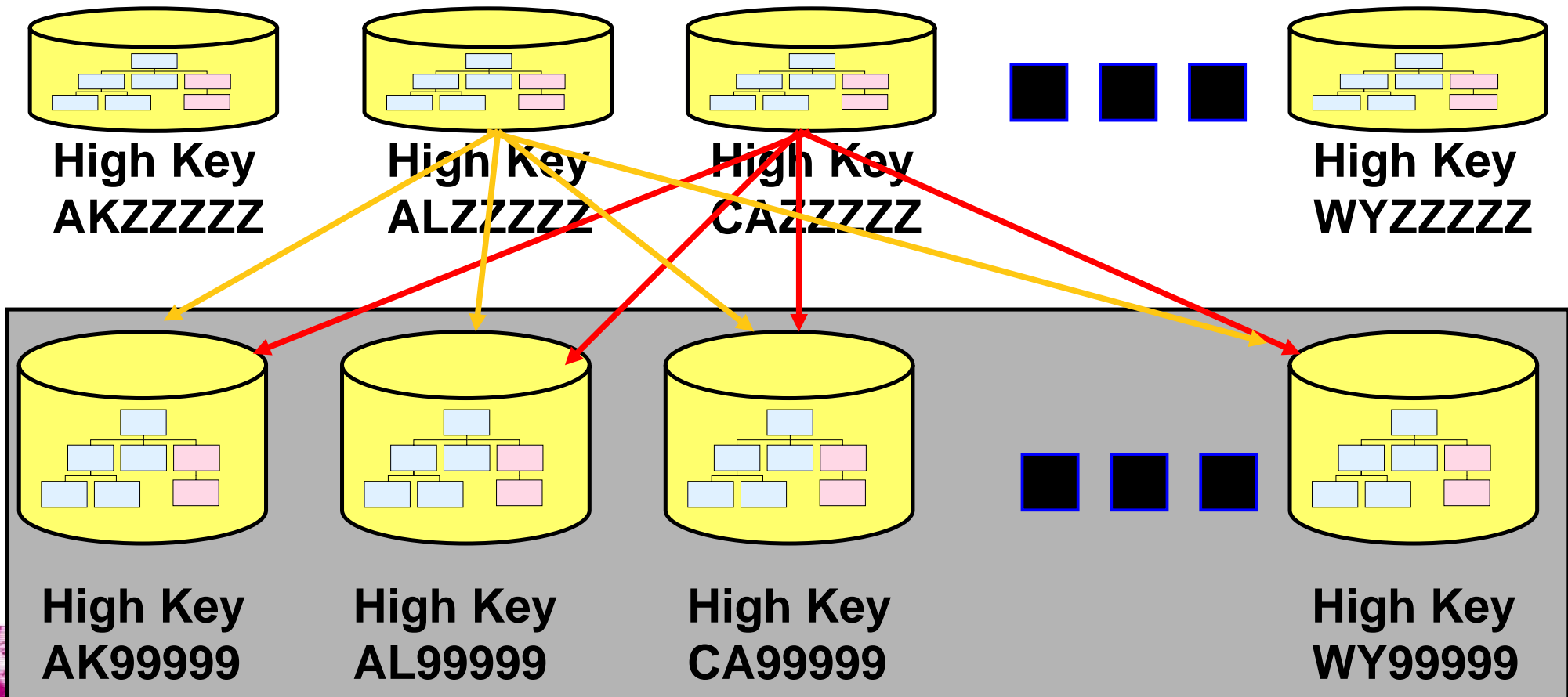
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- LCHILD NAME=(CITYINDX,Z1STAY),PTR=INDX
- XDFLD NAME=STATCITY,SRCH=(CITYNAME),SUBSEQ=/SX1
  
- LCHILD NAME=(CITYINDX,Z1STAY),PTR=INDX
- XDFLD NAME=STATCITY,SRCH=(STATEID,CITYNAME),  
SUBSEQ=/SX1



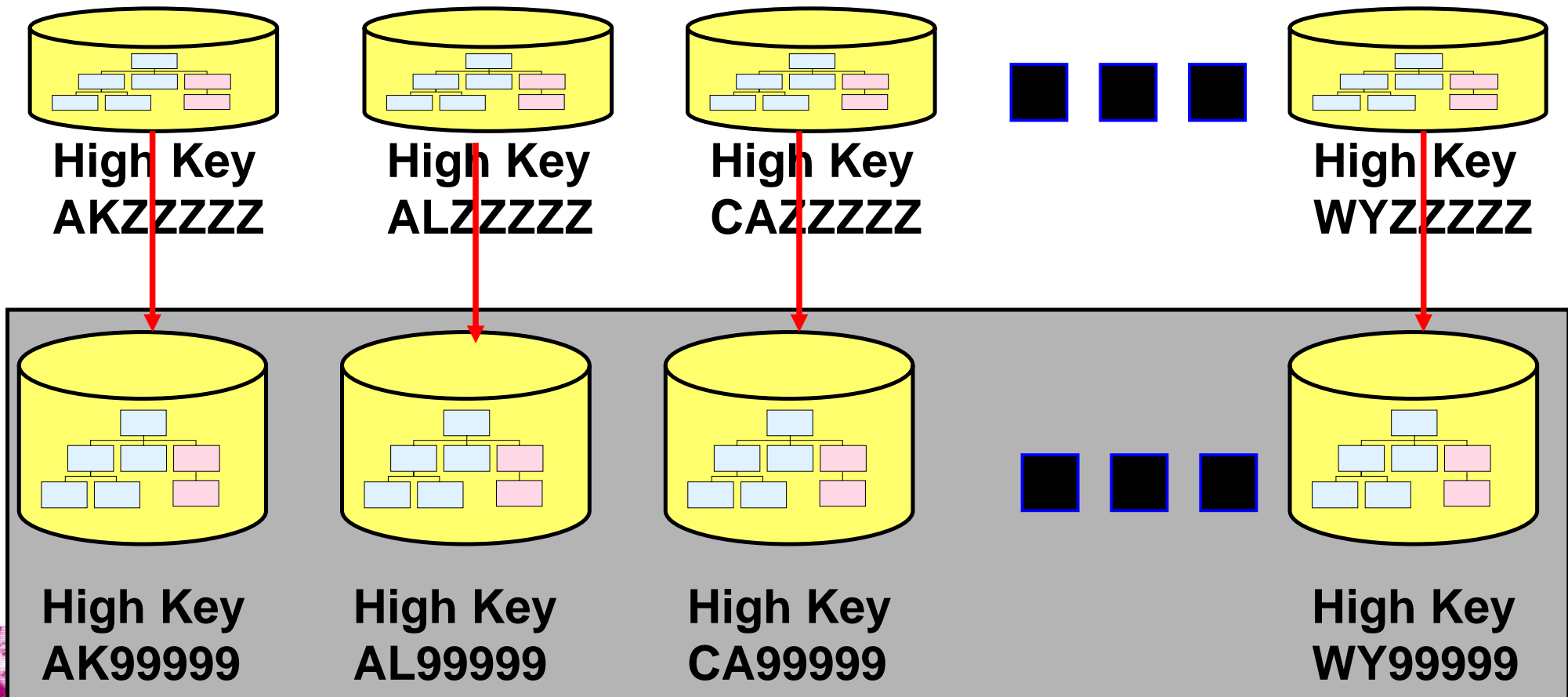
# Secondary Index independence

- Partition independence with secondary index processing



# Secondary Index independence

- Partition independence with secondary index processing



# Self Deleting Partitions

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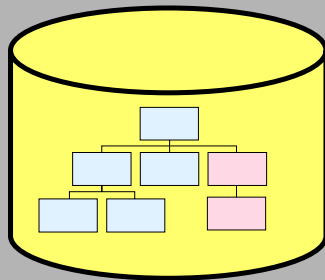
- Increasing Keys “move new data to right end partitions
- Oldest partitions become obsolete
- `DELETE.PART PART(SYR2001)`
- Limitation: 1001 partition IDs for the life of the HALDB

SYR2002

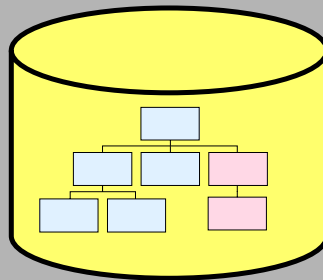
SYR2003

SYR2004

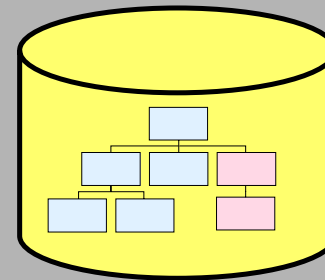
SYR2005



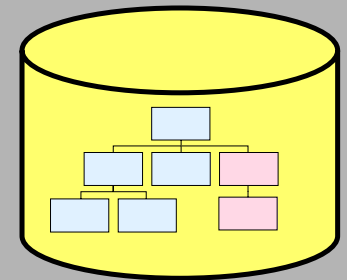
High Key  
20021231



High Key  
20031231



High Key  
20041231



High Key  
20051231

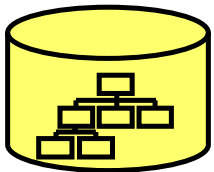


# Time Dependent Data

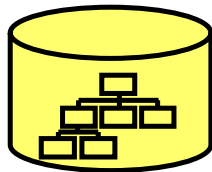
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- Activity database partitions
  - Partition selection exits works out day of week
  - Allows unload/clean up on “OFF” daysd

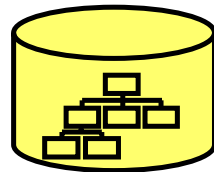
ISRT Roots



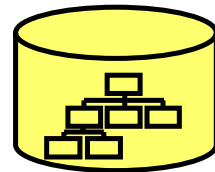
Monday



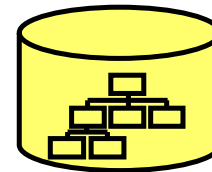
Tuesday



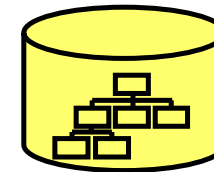
Wednesday



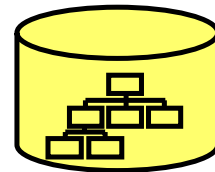
Thursday



Friday



Saturday



Sunday

