Signals & Variables (1A)

Concurrent & Sequential Signal Assignments

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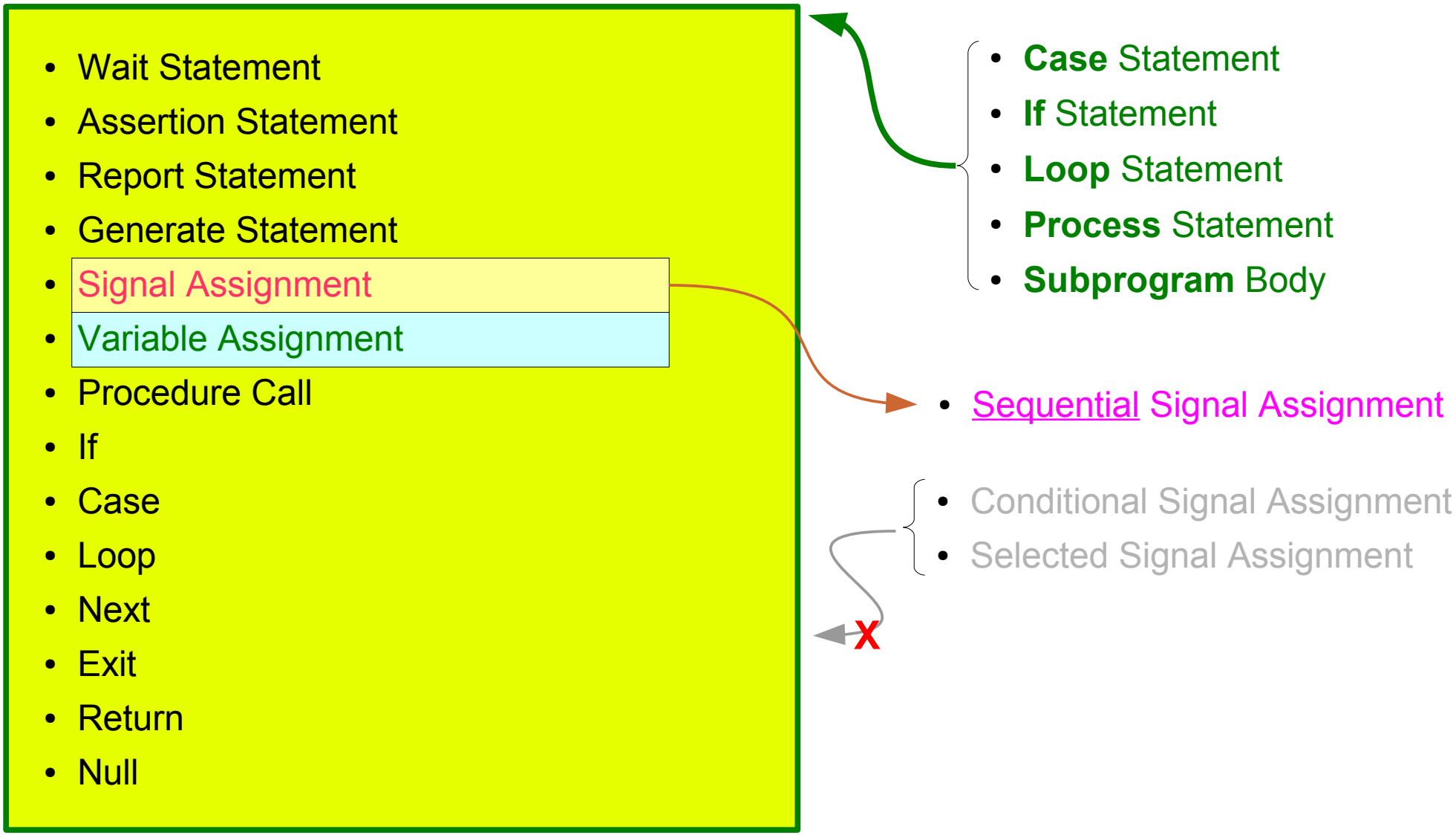
Sequential Statement



Concurrent Statement

- Block Statement
- Process Statement
- Component Statement
- Generate Statement
- **Concurrent** Signal Assignment
- Concurrent Assertion
- Concurrent Procedure Call

- **Architecture** Body
- **Block** Statement
- **Generate** Statement

- **Conditional** Signal Assignment
- **Selected** Signal Assignemnt

Concurrent Signal Assignment

- **Conditional** Signal Assignment

```
Z <=    A or B    [ after 1 ns ]   when SEL = "00"  else
        A or C    [ after 2 ns ]   when SEL = "01"  else
        A or D    [ after 2 ns ]   when SEL = "10"  else
        A or E    [ after 3 ns ]   when SEL = "11"  else
        A or F    [ after 4 ns ]   ;
```

condition

- **Selected** Signal Assignment

```
with SEL select
Z <=    A or B    [ after 1 ns ]   when "00",
        A or C    [ after 2 ns ]   when "01",
        A or D    [ after 3 ns ]   when "10",
        A or E    [ after 4 ns ]   when "11",
        A or F    [ after 5 ns ]   when others;
```

selection

Conditional Signal Assignment (1)

```
Z <=    A or B    [ after 1 ns ]  ;
```

simple concurrent statement

```
Z <=    A or B    [ after 1 ns ]  when S0 = '1' ;
```

One condition

```
Z <=    A or B    [ after 1 ns ]  when S0 = '1'  else
        C or D    [ after 2 ns ]  ;
```

One condition with 'else'

```
Z <=    A or B    [ after 1 ns ]  when S0 = '1'  else
        C or D    [ after 2 ns ]  when S1 = '1'  else
        E or F    [ after 3 ns ]  ;
```

Two conditions with 'else'

Concurrent Signal Assignment

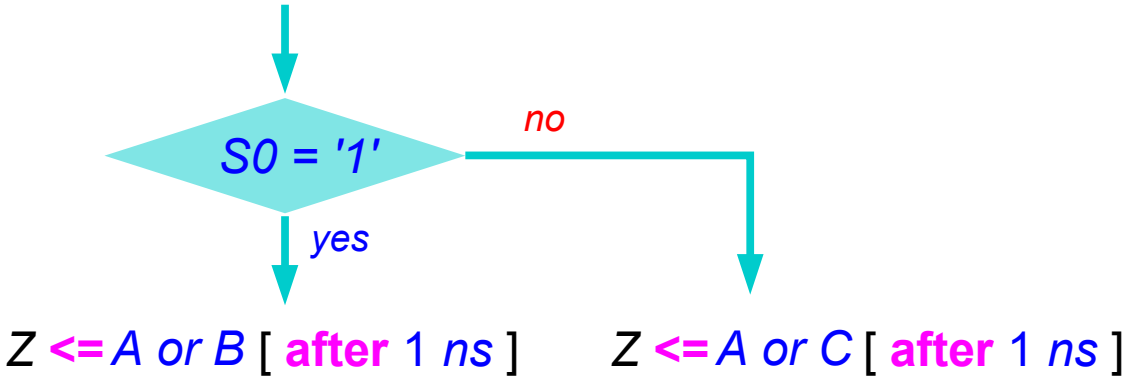
- **Conditional** Signal Assignment
- **Selected** Signal Assignment

Conditional Signal Assignment (2)

```
Z <=   A or B   [ after 1 ns ]   ;
```
simple concurrent statement

```
Z <=   A or B   [ after 1 ns ]   when S0 = '1' ;
```
One condition

```
Z <=   A or B   [ after 1 ns ]   when S0 = '1'  else
       C or D   [ after 2 ns ]   ;
```
One condition with 'else'

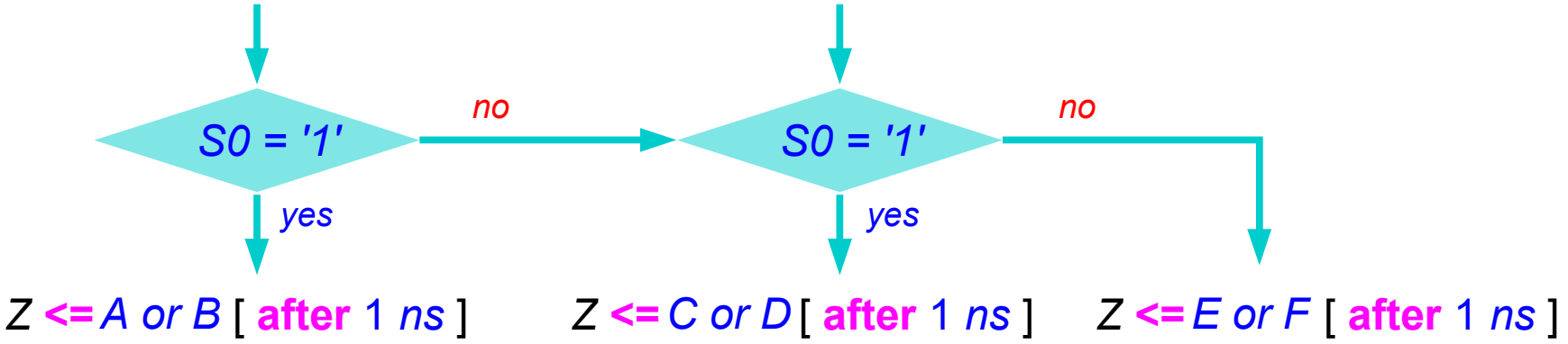


Conditional Signal Assignment (3)

```
Z <=    A or B    [ after 1 ns ]   ;
```
simple concurrent statement

```
Z <=    A or B    [ after 1 ns ]   when S0 = '1' ;
```
One condition

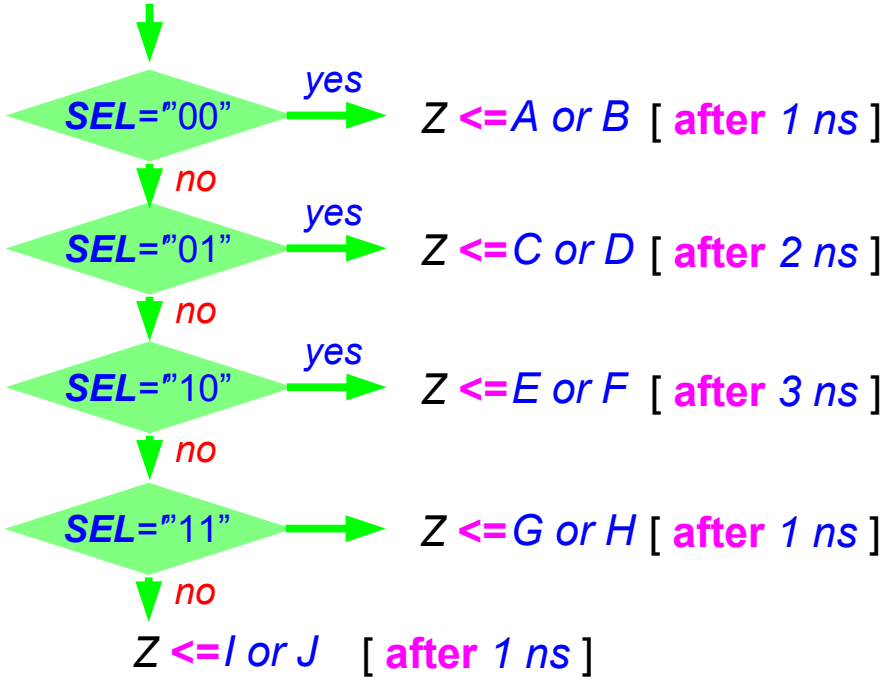
```
Z <=    A or B    [ after 1 ns ]   when S0 = '1'  else
        C or D    [ after 2 ns ]   when S1 = '1'  else
        E or F    [ after 3 ns ]   ;
```
Two conditions with 'else'



Selected Signal Assignment

```
with SEL select
Z <=    A or B   [ after 1 ns ]  when "00",
        C or D   [ after 2 ns ]  when "01",
        E or F   [ after 3 ns ]  when "10",
        G or H   [ after 4 ns ]  when "11",
        I or J   [ after 5 ns ]  when others;
```

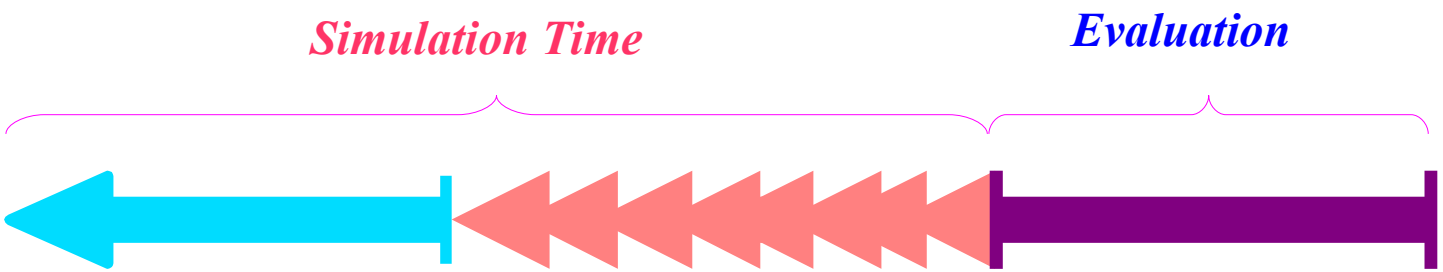
selection



Concurrent Signal Assignment

- **Conditional** Signal Assignment
- **Selected** Signal Assignment

Simulation Time (1)



Unit: ms, ns, ps, ...

$1\,ms = 1000\,ns$

$1\,ns = 1000\,ps$

Real Delay

Unitless Delta D

used for a simulator to mimic parallel activities simulator

$1\,ps \neq n \cdot \Delta$

no integer n that make n delta's equal to 1 ps.

$n \cdot \Delta = 0\,ps = 0\,ns \cdots$

Zero Delay

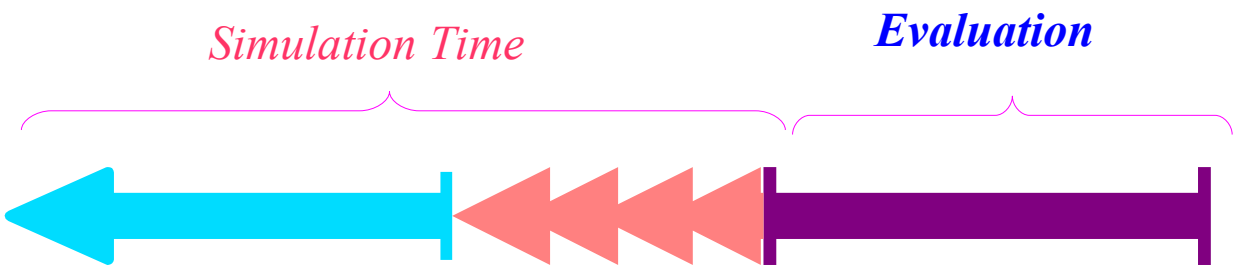
Assumed to take no time

Not related to any time

Zero Delay Assignment

```
X1 <= A or B ;
X1 <= A or B after 0 ns;
```

Simulation Time (2)



Unit: ms, ns, ps, ...

Unitless Delta

Zero Delay Assignment

```
X1  <=  A or B ;
```

X1 is updated after at least one Δ

Non-zero Delay Assignment

```
X1  <=  A or B after 2 ns;
```

X1 is updated after 2 ns

Concurrent vs Sequential (1)

architecture *arch* of entity *ent* is

• • •

begin

concurrent *signal statement*$_1$

concurrent *signal statement*$_2$

concurrent *signal statement*$_3$

outside process statement

process (A, B, C)

begin

Sequential *signal statement*$_1$

Sequential *signal statement*$_2$

Sequential *signal statement*$_3$

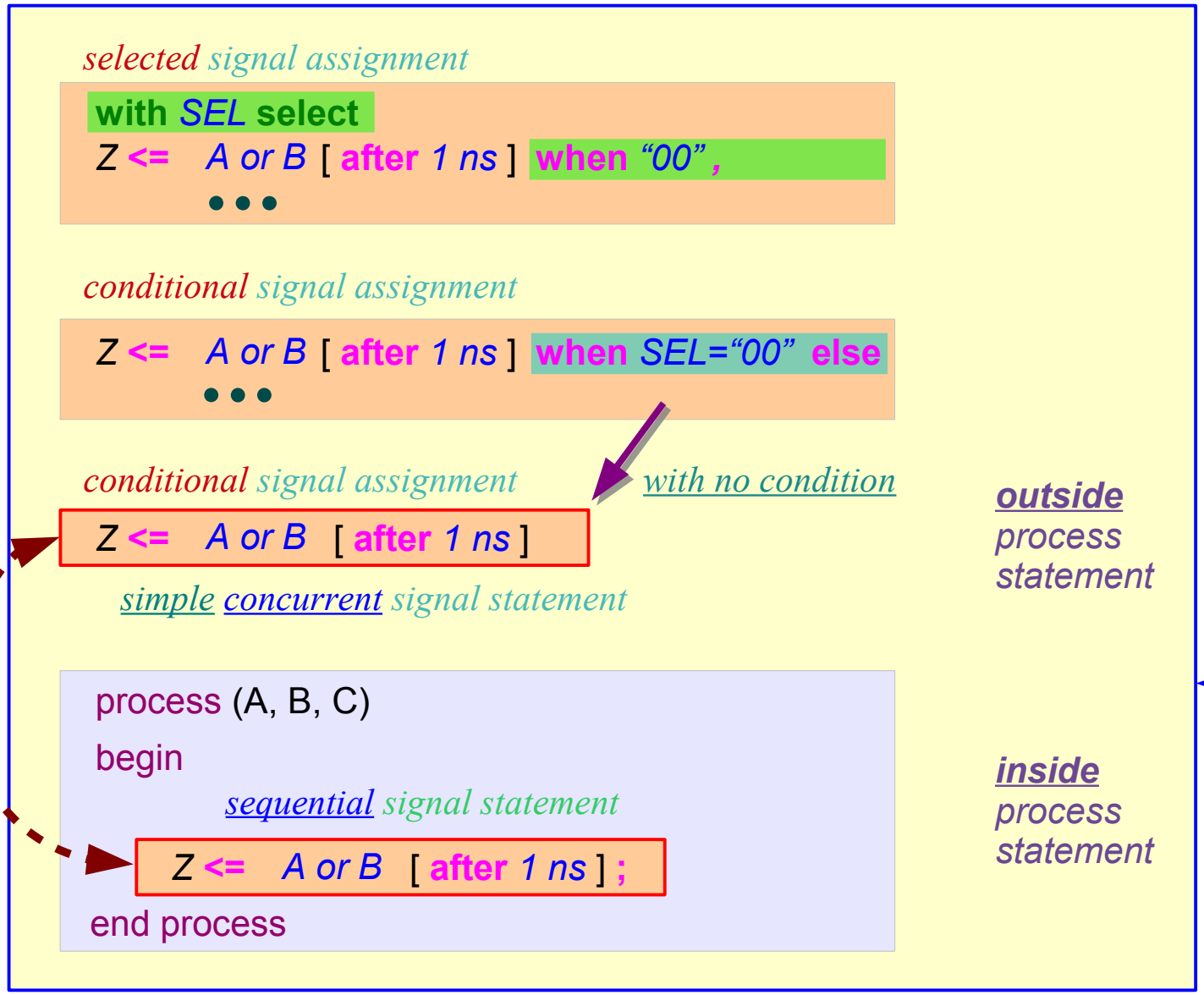
inside *process statement*

end process

end

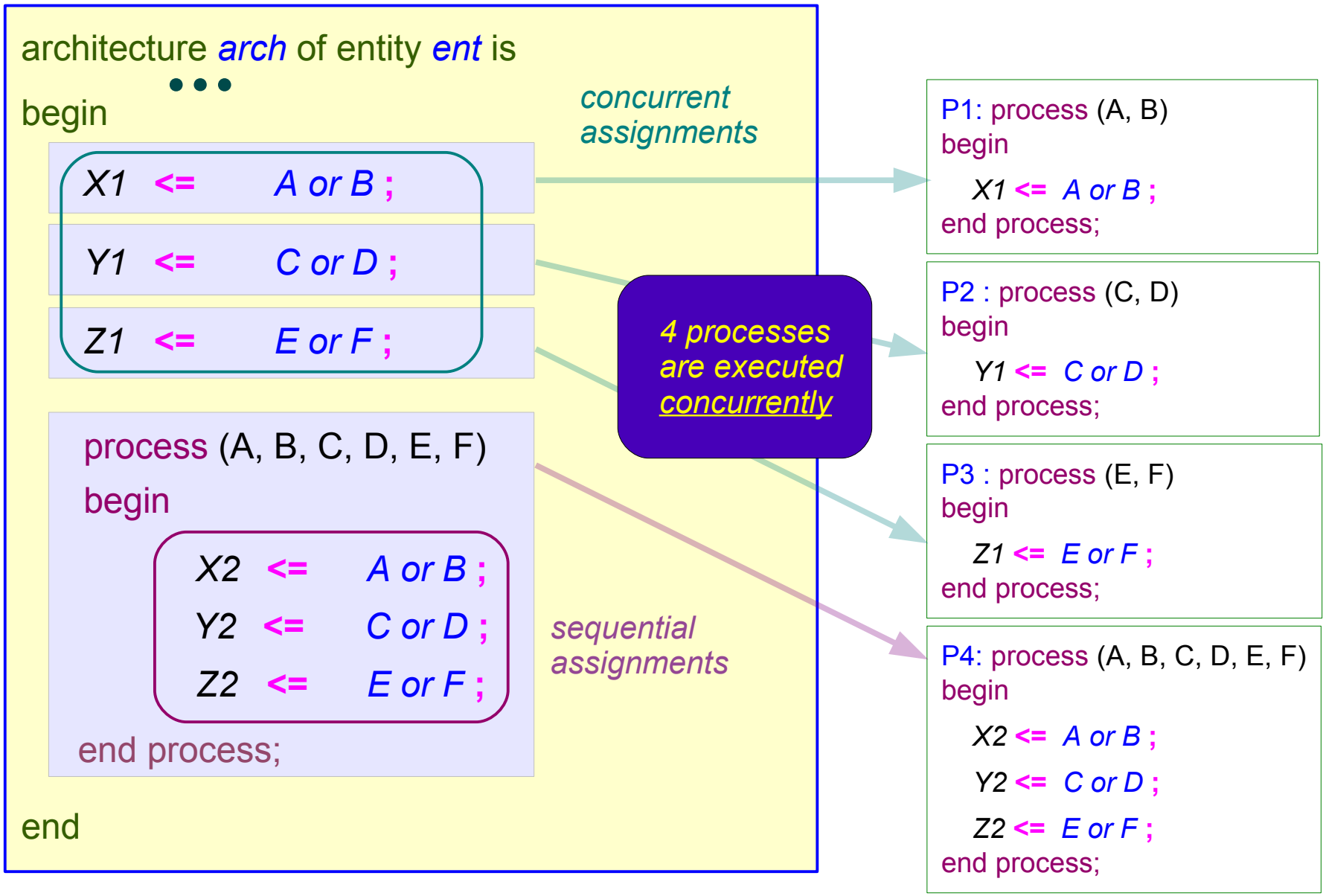
- Architecture Body
- Block Statement
- Generate Statement

Concurrent vs Sequential (2)



- Architecture Body
- Block Statement
- Generate Statement

Concurrent vs Sequential (3)



```
X1  <= A or B     after 1 ns ;

Y1  <= C or D     after 1 ns ;

Z1  <= E or F     after 1 ns ;
```

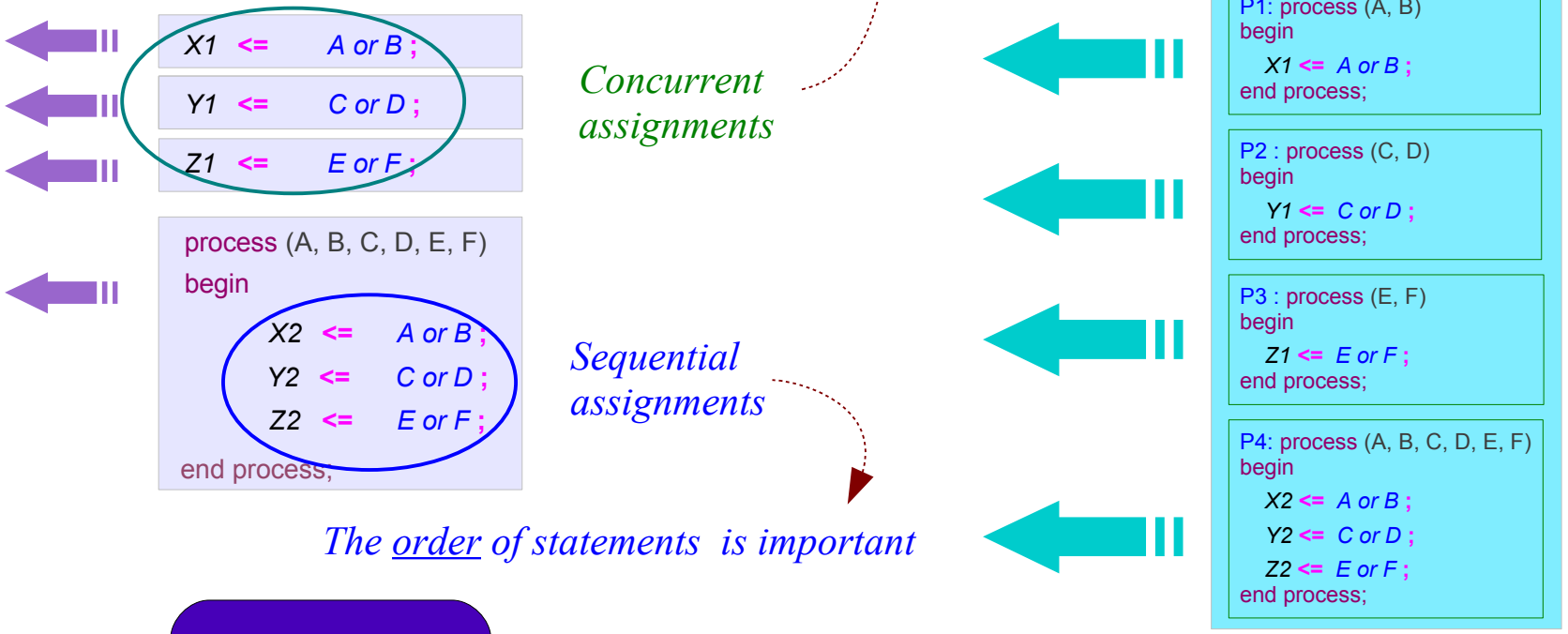
```
process (A, B, C, D, E, F)
begin
        X2  <= A or B     after 1 ns ;
        Y2  <= C or D     after 1 ns ;
        Z2  <= E or F     after 1 ns ;
end process;
```

Concurrent vs Sequential (4)

Simulation of parallel activities

Non-deterministic Execution Order

Don't know which process executes first among P1 ~ P4.



The order of statements is important

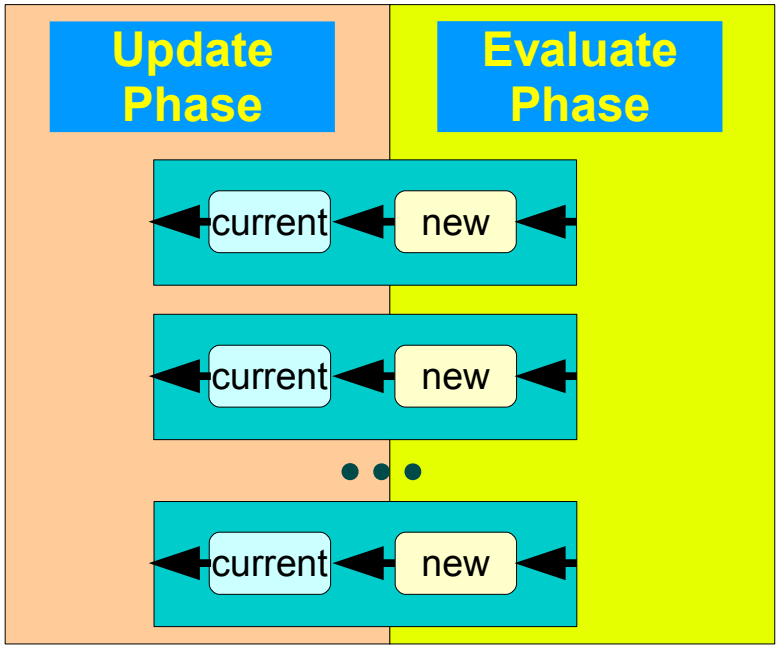
4 processes are executed concurrently

Evaluate – Update (1)

*When X or Y is changed, the assignments are **evaluated** using the current values, not the new values of X or Y*

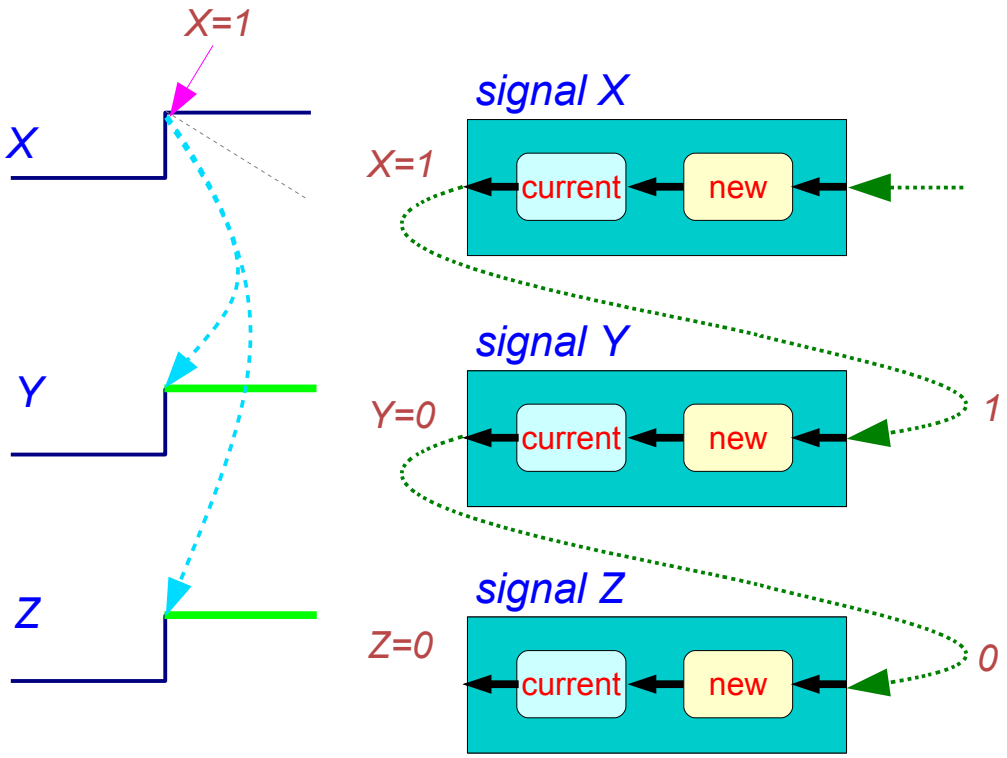
Non-Blocking Assignments

```
process (X, Y)
begin
        Y   <=   X ;
        Z   <=   Y ;
end process;
```



process (X, Y)

Only when X or Y is changed (updated), two assignments are evaluated and updated.

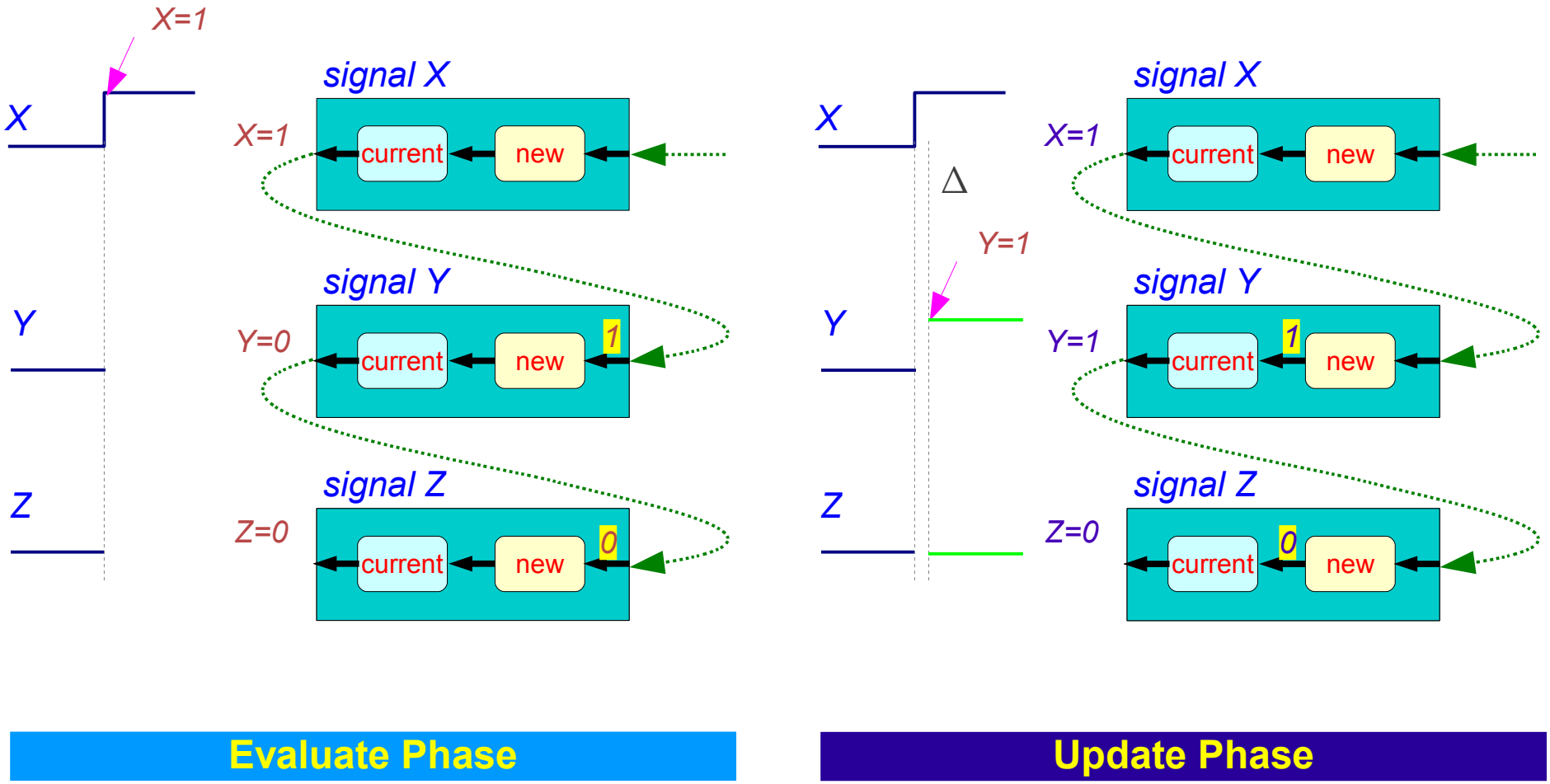


Evaluate – Update (2)

process (X, Y)

*Event on X – **X** changed into new value '1'*

Induces a new event on Y



Evaluate Phase

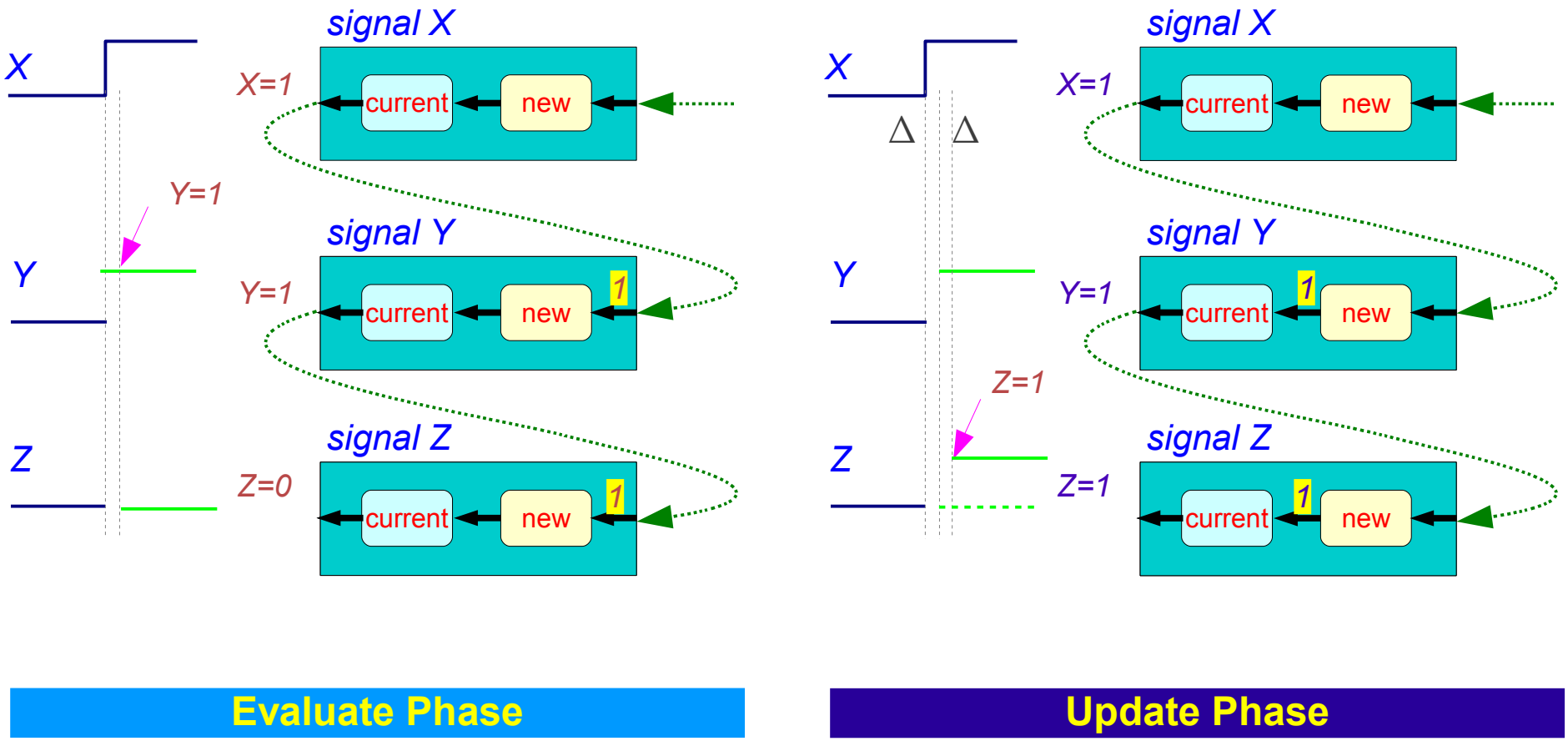
Update Phase

Evaluate – Update (3)

process (X, Y)

*Event on Y – **Y** changed into new value '1'*

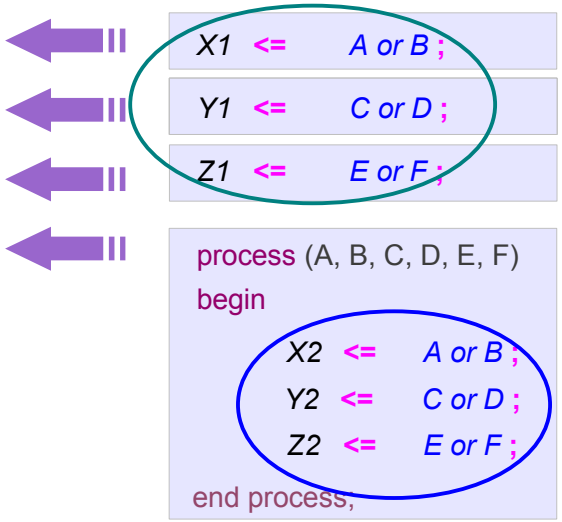
Induces a new event on Z



Zero vs Non-zero Delay Assignments (1)

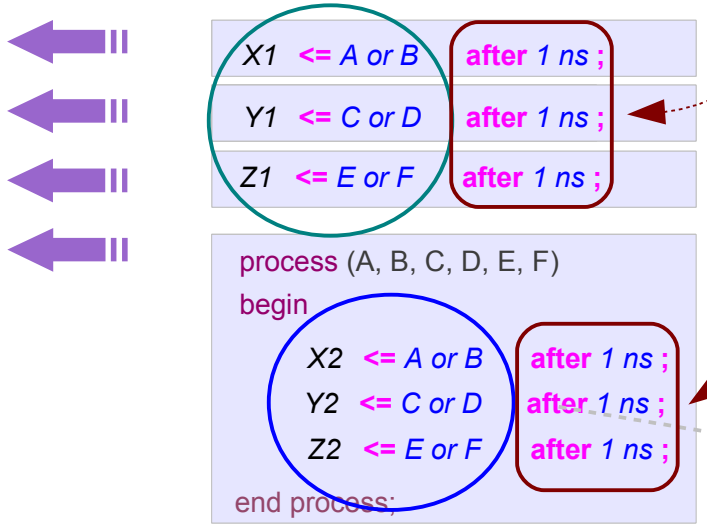
When A, B, C, D, E, or F is changed, the assignments are evaluated using the current values, not the new values of A, B, C, D, E, F

zero delay assignments



Updated values are observable after at least one delta time

nonzero delay assignments



Updated values are observable after 1 ns

Zero vs Non-zero Delay Assignments (2)

When A, B, C, D, E, or F is changed, the assignments are evaluated using the current values, not the new values of A, B, C, D, E, F

zero delay assignments

```
X1 <= A or B ;
Y1 <= C or D ;
Z1 <= E or F ;

process (A, B, C, D, E, F)
begin
    X2 <= A or B ;
    Y2 <= C or D ;
    Z2 <= E or F ;
end process;
```

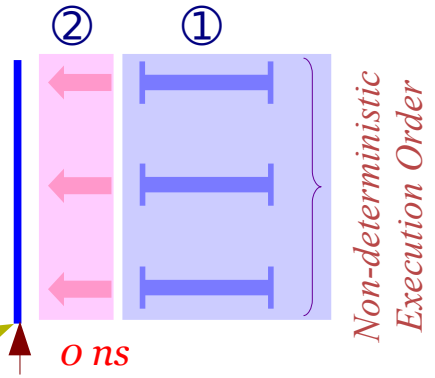
* *The order of statements is important*
  1. *scheduled after some delta time*
  2. *non-blocking assignment*

only lumped view ②①

lumped view ④

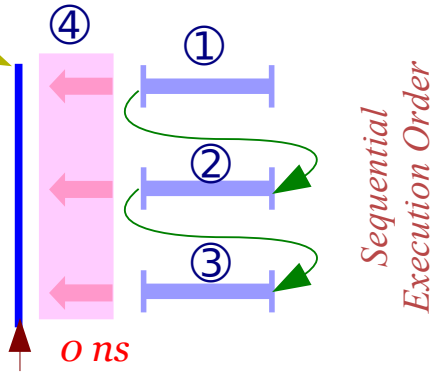
Updated values are observable after at least one delta time

Concurrent assignments



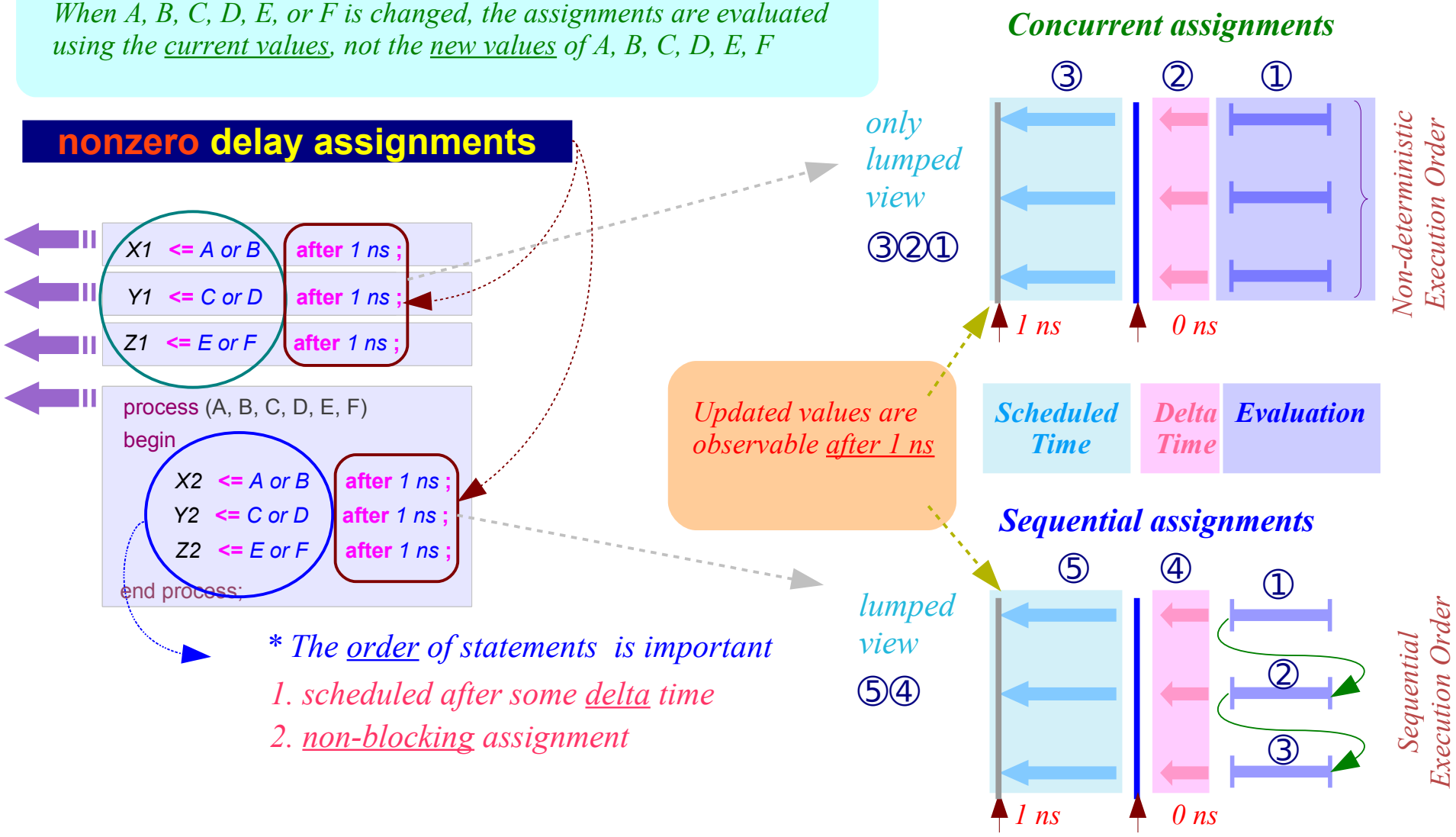
Scheduled Time | ***Delta Time*** | ***Evaluation***

Sequential assignments



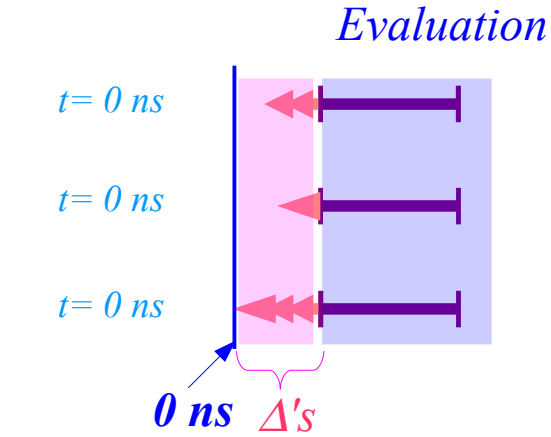
Zero vs Non-zero Delay Assignments (3)

When A, B, C, D, E, or F is changed, the assignments are evaluated using the current values, not the new values of A, B, C, D, E, F

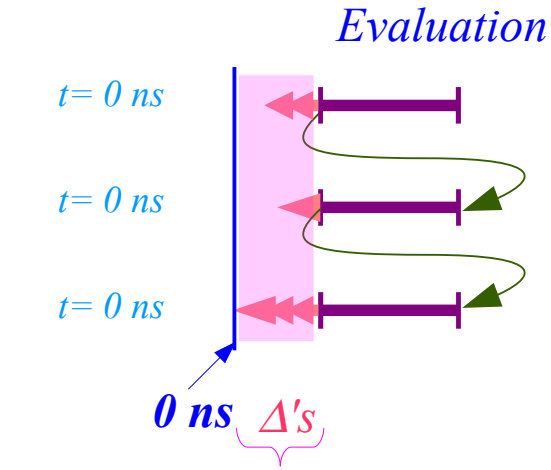


Zero Delay Assignment

```
architecture arch of entity ent is
   • • •
begin
   X1 <= A or B ;
   Y1 <= C or D ;
   Z1 <= E or F ;

   process (A, B, C, D, E, F)
   begin
      X2 <= A or B ;
      Y2 <= C or D ;
      Z2 <= E or F ;
   end process;

end
```



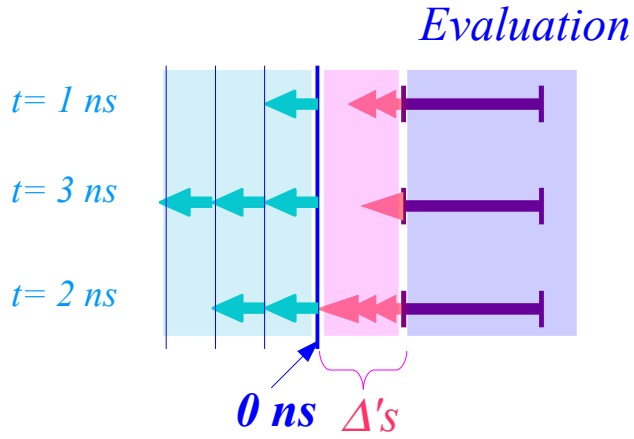
The exact no of delta is determined by the simulator and the context



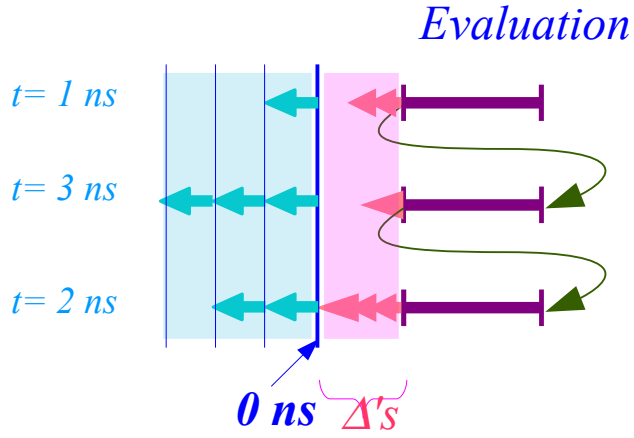
Not yet updated

Non-Zero Delay Assignment

```
architecture arch of entity ent is
• • •
begin
    X1 <= A or B  after 1 ns;
    Y1 <= C or D  after 3 ns;
    Z1 <= E or F  after 2 ns;

    process (A, B, C, D, E, F)
    begin
        X2 <= A or B  after 1 ns;
        Y2 <= C or D  after 3 ns;
        Z2 <= E or F  after 2 ns;
    end process;

end
```



The exact no of delta is determined by the simulator and the context



Not yet updated

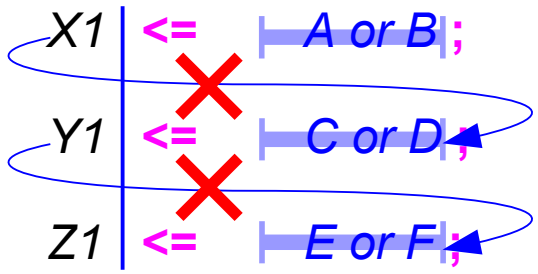
Non-blocking Assignment (1)

```
architecture arch of entity ent is
  • • •
begin
  X1 <= A or B;
  Y1 <= C or D;
  Z1 <= E or F;

  process (A, B, C, D, E, F)
  begin
    X2 <= A or B;
    Y2 <= C or D;
    Z2 <= E or F;
  end process;
end
```

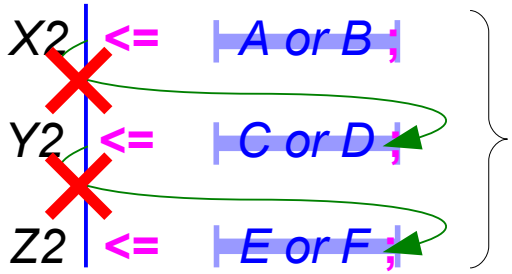
Concurrent assignment

Sequential assignment

When A, B, C, D, E, or F is changed, the assignments are evaluated using the current values, not the new values of A, B, C, D, E, F



non-blocking assignment

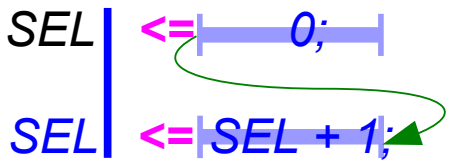


Non-blocking Assignment (2)

```
process (A, I0, I1)
begin
   SEL <=  0;
   if (A='1') then SEL <= SEL + 1; end if;
   case SEL is
        when 0
             Q <= I0;
        when 1
             Q <= I1;
   end case;
end process;
```

Scheduled on the next delta time

➡ *SEL value will not be **updated***
until the next delta time



Non-blocking Assignment

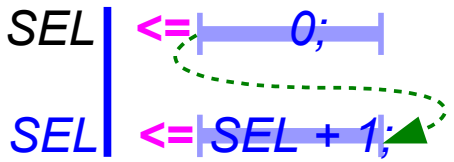
Without waiting *the next delta time,*
it can continue to process the next
sequential statement
(processed with the wrong value of SEL)

Non-blocking Assignment (3)

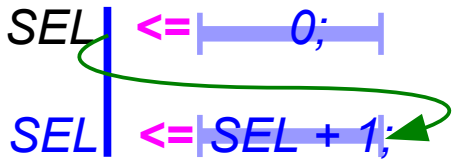
```
process
begin
    SEL      <=   A or B;
    wait for 0 ns;
    if (A='1') then SEL <= SEL + 1; end if;
    wait for 0 ns;
    case SEL is
        when 0
            Q <= I0;
        when 1
            Q <= I1;
    end case;
    wait on A, I0, I1;

end process;
```

Wait for one delta time

Non-blocking
*: next statement **before** update*



wait for *0 ns;*

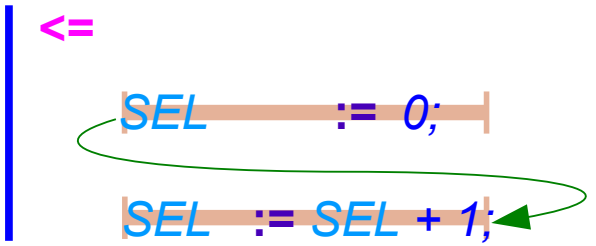


Blocking
*: next statement **after** update*

Non-blocking Assignment (4)

```
process (A, I0, I1)
  variable SEL : integer range 0 to 1;
begin
    SEL       :=    A or B;
    if (A='1') then SEL := SEL + 1; end if;
    case SEL is
          when 0
                Q <= I0;
          when 1
                Q <= I1;
    end case;

end process;
```

Variable SEL changes its value immediately.



General MUX model

```
process (A, I0, I1)
begin
    case A is
        when '0'
            Q <= I0;
        when '1'
            Q <= I1;
    end case;
end process;
```

Variable & Signal Assignments

When A, B, C, D, E, or F is changed, the assignments are evaluated using the current values, not the new values of A, B, C, D, E, F

Variable assignments

```
process (A, B, C, D, E, F)
   variable X2, Y2, Z2 : bit;
begin
      X2  :=  A or B  ;
      Y2  :=  C or D  ;
      Z2  :=  E or F  ;
end process;
```

*Updated values of X2, Y2, Z2 are observable **immediately***

Signal assignments

```
process (A, B, C, D, E, F)
begin
      X2  <=  A or B  ;
      Y2  <=  C or D  ;
      Z2  <=  E or F  ;
end process;
```

Updated values X2, Y2, Z2 are observable after at least on delta time.

Variable Assignment (1)

```
architecture arch of entity ent is
   • • •
begin
    X1  <=      A or B   after 1 ns;
    Y1  <=      C or D   after 3 ns;
    Z1  :=      E or F  ;

    process (A, B, C, D, E, F)
        variable Y2, Z2 : bit;
    begin
        X2  <=      A or B   after 1 ns;
        Y2  :=      C or D   after 3 ns;
        Z2  :=      E or F  ;
    end process;

end
```

*The Variable assignment is a sequential statement and cannot be used **outside a process** statement.*

*The variable is **declared** here. They are used for **local storage** in process statement and subprogram body.*

*The variable assignment has **nothing to do with time**. It executes immediately.*

*It **cannot** have an **after** clause*

Variable Assignment (2)

```
process (A, B, C, D, E, F)
    variable Z2 : bit;
begin
    X2  <=  A or B  after 1 ns;
    Y2  <=  C or D  after 3 ns;
    Z2  :=  E or F  ;
end process;
```

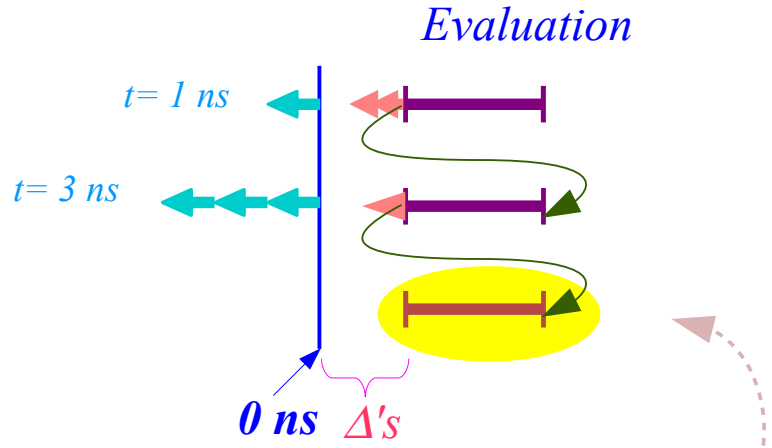
```
X2  <=  A or B  after 1 ns;
Y2  <=  C or D  after 3 ns;
        Z2  :=  E or F  ;
```

```
process (A, B, C, D, E, F)
    variable Y2 : bit;
begin
    X2  <=  A or B  after 1 ns;
    Y2  :=  C or D  ;
    Z2  <=  E or F  after 2 ns;
end process;
```

```
X2  <=  A or B  after 1 ns;
        Y2  :=  C or D  ;
Z2  <=  E or F  after 2 ns;
```

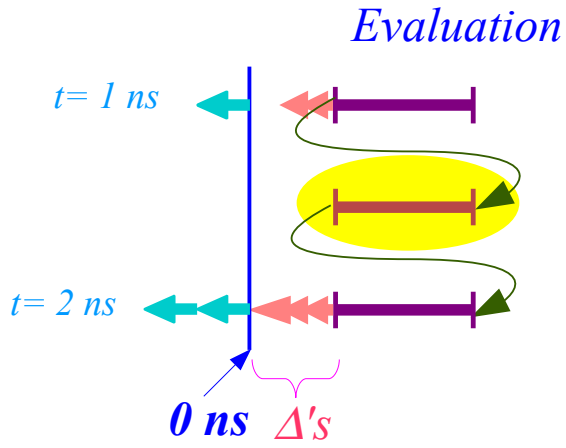
Variable Assignment (3)

```
process (A, B, C, D, E, F)
    variable Z2 : bit;
begin
    X2  <=   A or B   after 1 ns;
    Y2  <=   C or D   after 3 ns;
    Z2  :=   E or F   ;
end process;
```



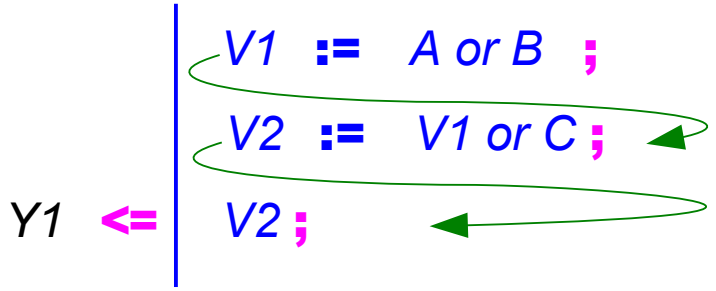
The variable assignment has nothing to do with time. It executes immediately.

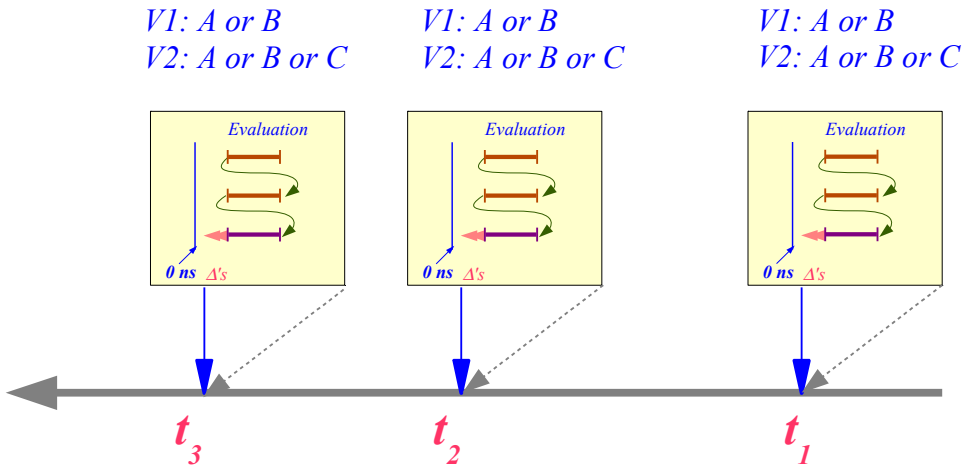
```
process (A, B, C, D, E, F)
    variable Y2 : bit;
begin
    X2  <=   A or B   after 1 ns;
    Y2  :=   C or D   ;
    Z2  <=   E or F   after 2 ns;
end process;
```



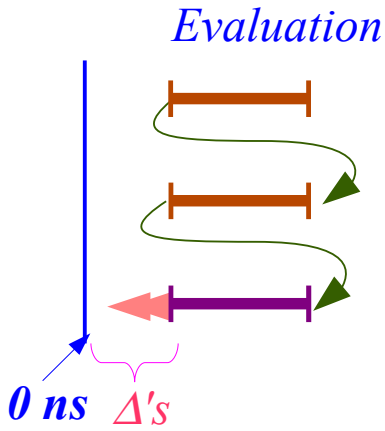
Mixed Assignments Example (1)

```
process (A, B, C)
    variable V1, V2 : bit;
begin
    V1  :=  A or B  ;
    V2  :=  V1 or C ;
    Y1  <=  V2 ;
end process;
```





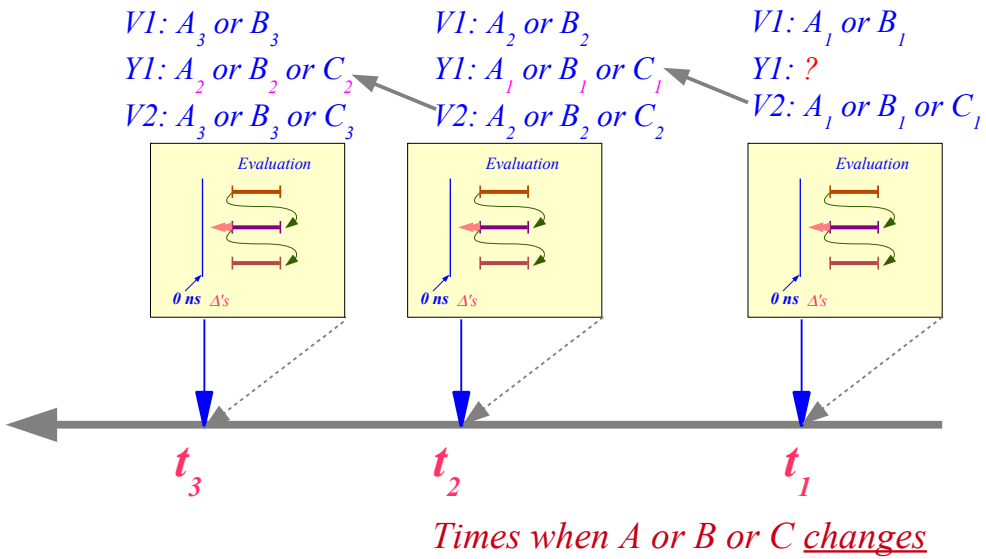
Times when A or B or C changes

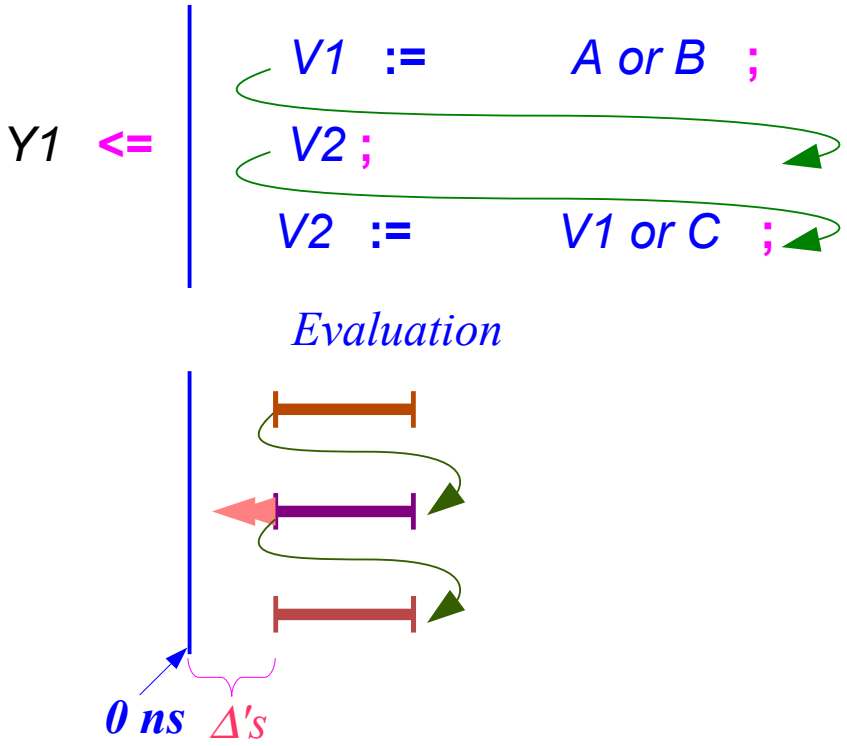


Y1 <= A or B or C ;

Mixed Assignments Example (2)

```
process (A, B, C)
    variable V1, V2 : bit;
begin
    V1  :=  A or B  ;
    Y1  <=  V2 ;
    V2  :=  V1 or C  ;
end process;
```

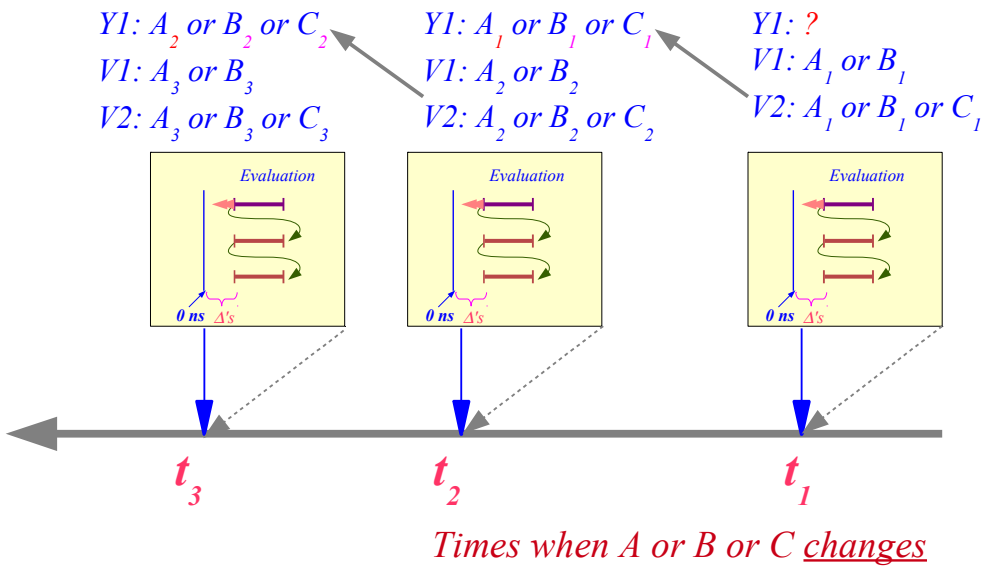


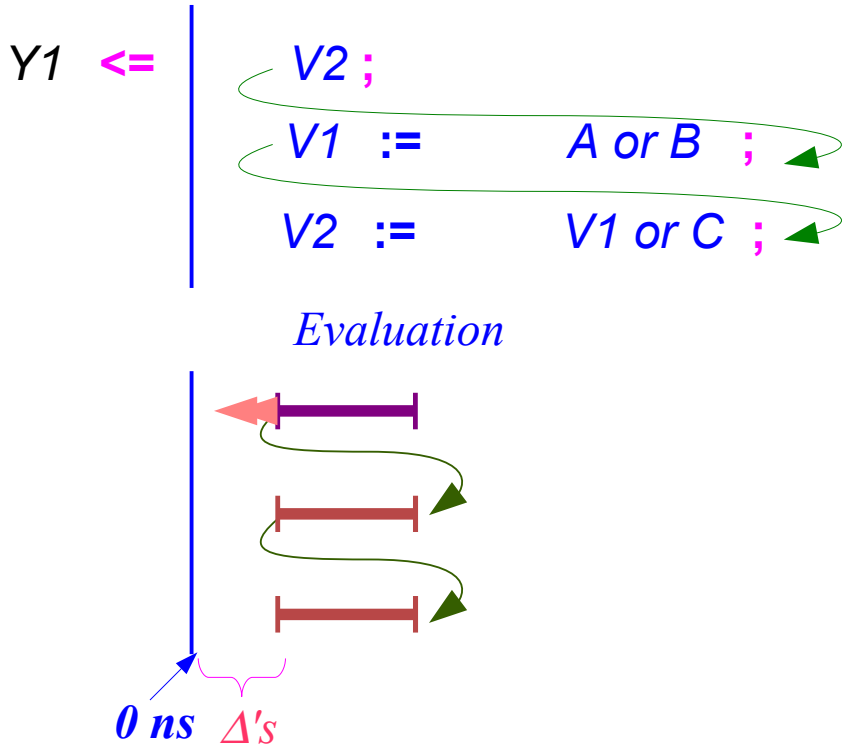


Y1 <= A or B or C ;

Mixed Assignments Example (3)

```
process (A, B, C)
    variable V1, V2 : bit;
begin
    Y1  <=  V2 ;
    V1  :=  A or B ;
    V2  :=  V1 or C ;
end process;
```





Y1 <= A or B or C ;

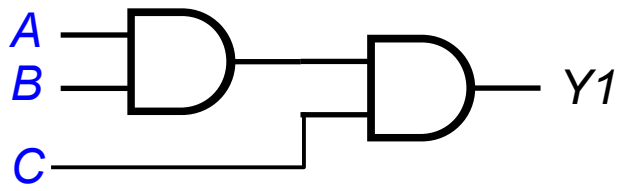
Mixed Assignments Example (4)

```
process (A, B, C)
    variable V1, V2 : bit;
begin
    V1  :=  A or B  ;
    V2  :=  V1 or C ;
    Y1  <=  V2 ;
end process;
```

```
process (A, B, C)
    variable V1, V2 : bit;
begin
    Y1  <=  V2 ;
    V1  :=  A or B  ;
    V2  :=  V1 or C ;
end process;
```

```
process (A, B, C)
    variable V1, V2 : bit;
begin
    V1  :=  A or B  ;
    Y1  <=  V2 ;
    V2  :=  V1 or C ;
end process;
```

Same Synthesis Result



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